

[week8] File System API



- OS provides the file system abstraction
- How do application processes access the file system?

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Can we directly write() to a dir file?



- ... using system calls (e.g. write)
- This is a permission question, not a privilege question
- The answer is no
- Explanation:
 - Any process independent of file system shouldn't be allowed to directly write in a directory the way they can write in files
 - Since such processes don't know how exactly the internals of directory are laid out. They should only see files, create files, and delete files.

System Calls to UNIX File Systems



• 19 system calls into 6 categories:

Return file desp.	Assign inodes		Process input/ output	Change file system	Modify view of file system
open close creat pipe dup	creat link unlink	chown chmod stat fstat		mount umount	chdir chroot

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Roadmap

- Functionality (API)
 - Basic functionality
 - Disk layout
 - File operations (open, read, write, close)
 - Directories
- Performance
 - Disk allocation
 - Buffer cache
 - File System interface
 - Disk scheduling
- Reliability
 - FS level
 - Disk level: RAID

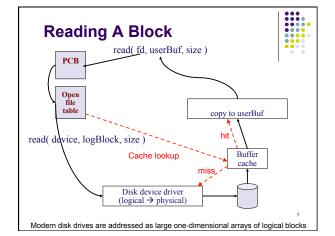
[week1] Roadmap for ECE469 **lectures**



Introduction to OS components Individual components

- Process management
- Memory managementFile management
- Secondary-storage management
- (Device management)
- (shell)
- (Networking)

Hardware support for OS interspersed (before relevant topics)



File System Interface



- How do application programs typically access file data?
 - Explicit read/write operations (conventional)

Read / Write Interface

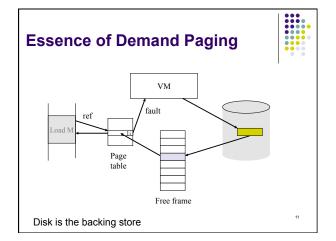


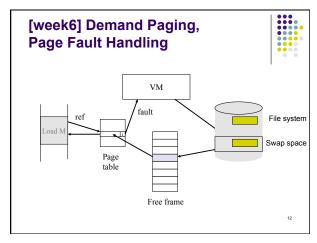
- File data is explicitly copied between disk file and process memory
- Programs cannot directly access file data

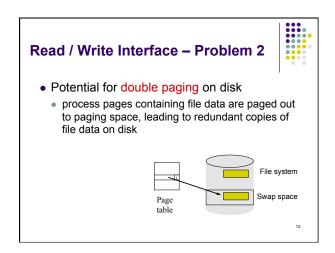
FileDescriptor fhandle; int offset, length; char *buffer = (char *) malloc(1000000);

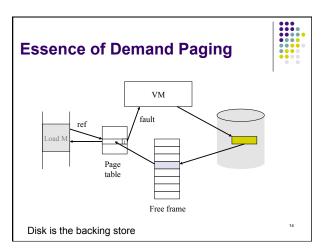
fhandle = open("pathname");
read(fhandle, offset, buffer, length);
{read file data in buffer to do important computation};
Write(fhandle, offset, buffer, length)
close(fhandle);

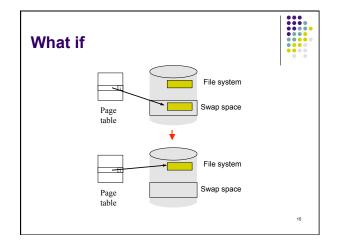
Read / Write Interface - Problem 1 • Potential for double copies in mem read(fd, userBuf, size) Open file table copy to userBuf hith Cache lookup Disk device driver (logical → physical)











Memory-mapped Files



- File is "mapped" into application's address space
 - by initializing virtual memory so that the file (directly) serves as backing store for a region of the application's address space

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Memory-mapped Files (cont')



• Elegant integration of file system and virtual memory

```
FileDescriptor fhandle;
int offset, length;
                                       This is like after
char *address;
                                       address = malloc()
fhandle = open("pathname");
mmap(fhandle, offset, address, length);
{read/write file data by accessing memory range
```

[address, address + length]}; munmap(address, length);

close(fhandle);

Implementation of mmap()



Memory-mapped Files



- File is "mapped" into application's address space
 - by initializing virtual memory so that the file (directly) serves as backing store for a region of the application's address space
- File data is demand paged upon access to the mapped file
 - No double paging on disk
- Memory-mapped files do not go through buffer cache
 - No double copy in mem
 - Program accesses file data directly

Effects and Semantics of Memory-mapped Files



- Processes that map the same file share physical memory that caches file data
- Writes may not be immediately writen to the file on disk
 - Update periodically
 - Closing the file results in writing all to disk and removing the VM mapping

Fun with memory-mapped files



- Inter-process communication
 - Virtual addresses of diff processes mapped to the same file
- File copying as Memory copying
 - Map files to virtual addresses
 - Do memory copying

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Reading



• Chapters 11-12

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