Data Compression

- basics
- run-length coding
- Huffman compression
- LZW compression

Data compression

Compression reduces the size of a file:

- To save space when storing it.
- To save time when transmitting it.
- Most files have lots of redundancy.

Who needs compression?

- Moore's law: # transistors on a chip doubles every 18-24 months.
- Parkinson's law: data expands to fill space available.
- Text, images, sound, video, …

> " *All of the books in the world contain no more information than is broadcast as video in a single large American city in a single year. Not all bits have equal value.* " *— Carl Sagan*

Basic concepts ancient (1950s), best technology recently developed.

Applications

Generic file compression.

- Files: GZIP, BZIP, BOA.
- Archivers: PKZIP.
- File systems: NTFS.

Multimedia.

- Images: GIF, JPEG.
- Sound: MP3.
- Video: MPEG, DivX™, HDTV.

Communication.

- ITU-T T4 Group 3 Fax.
- V.42bis modem.

Databases. Google.

Lossless compression and expansion

uses fewer bits (you hope)

Message. Binary data B we want to compress.

Compress. Generates a "compressed" representation $C(B)$.

Expand. Reconstructs original bitstream B.

Basic model for data compression

Compression ratio. Bits in $C(B)$ / bits in B.

Ex. 50-75% or better compression ratio for natural language.

Food for thought

Data compression has been omnipresent since antiquity:

- Number systems.
- Natural languages.
- Mathematical notation.

has played a central role in communications technology,

- Braille.
- Morse code.
- Telephone system.

and is part of modern life.

- MP3.
- MPEG.

Q. What role will it play in the future?

Data Compression

- basics
- run-length coding
- Huffman compression
- LZW compression

Data representation: genomic code

Genome. String over the alphabet { A, C, T, G }.

Goal. Encode an N-character genome: **ATAGATGCATAG**...

Standard ASCII encoding.

- 8 bits per char.
- 8 N bits.

char	hex	binary
A	41	01000001
C	43	01000011
T	54	01010100
G	47	01000111

Two-bit encoding.

- 2 bits per char.
- 2 N bits.

char	binary
A	00
C	01
T	10
G	11

Amazing but true. Initial genomic databases in 1990s did not use such a code!

Fixed-length code. k-bit code supports alphabet of size 2^k.

Reading and writing binary data

Binary standard input and standard output. Libraries to read and write bits from standard input and to standard output.

`public class BinaryStdIn`

`boolean`	`readBoolean()`	*read 1 bit of data and return as a* `boolean` *value*
`char`	`readChar()`	*read 8 bits of data and return as a* `char` *value*
`char`	`readChar(int r)`	*read* `r` *bits of data and return as a* `char` *value*
[similar methods for `byte` *(8 bits);* `short` *(16 bits);* `int` *(32 bits);* `long` *and* `double` *(64 bits)]*		
`boolean`	`isEmpty()`	*is the bitstream empty?*
`void`	`close()`	*close the bitstream*

`public class BinaryStdOut`

`void`	`write(boolean b)`	*write the specified bit*
`void`	`write(char c)`	*write the specified 8-bit* `char`
`void`	`write(char c, int r)`	*write the* `r` *least significant bits of the specified* `char`
[similar methods for `byte` *(8 bits);* `short` *(16 bits);* `int` *(32 bits);* `long` *and* `double` *(64 bits)]*		
`void`	`close()`	*close the bitstream*

Writing binary data

Date representation. Different ways to represent 12/31/1999.

A character stream (StdOut)

```
StdOut.print(month + "/" + day + "/" + year);
```

`00110001|00110010|00101111|00110111|00110001|00101111|00110001|00111001|00111001|00111001`

1 2 / 3 1 / 1 9 9 9

80 bits

Three ints (BinaryStdOut)

```
BinaryStdOut.write(month);
BinaryStdOut.write(day);
BinaryStdOut.write(year);
```

`00000000|00000000|00000000|00001100|00000000|00000000|00000000|00011111|00000000|00000000|00000111|11001111`

12 31 1999

96 bits

Two chars and a short (BinaryStdOut)

```
BinaryStdOut.write((char) month);
BinaryStdOut.write((char) day);
BinaryStdOut.write((short) year);
```

`00001100|00011111|00000111|11001111`

12 31 1999

32 bits

A 4-bit field, a 5-bit field, and a 12-bit field (BinaryStdOut)

```
BinaryStdOut.write(month, 4);
BinaryStdOut.write(day, 5);
BinaryStdOut.write(year, 12);
```

`11001111|10111110|01111000`

12 31 1999

21 bits (+ 3 bits for byte alignment at close)

Four ways to put a date onto standard output

Binary dumps

Q. How to examine the contents of a bitstream?

Standard character stream

```
% more abra.txt
ABRACADABRA!
```

Bitstream represented as 0 and 1 characters

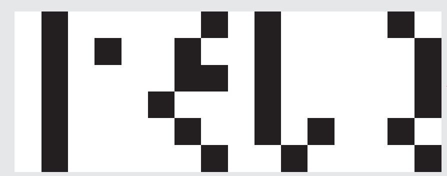
```
% java BinaryDump 16 < abra.txt
0100000101000010
0101001001000001
0100001101000001
0100010001000001
0100001001010010
0100000100100001
96 bits
```

Bitstream represented with hex digits

```
%  java HexDump 4 < abra.txt
41 42 52 41
43 41 44 41
42 52 41 21
12 bytes
```

Bitstream represented as pixels in a `Picture`

```
% java PictureDump 16 6 < abra.txt
```



16-by-6 pixel window, magnified

```
96 bits
```

	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
0	NUL	SOH	STX	ETX	EOT	ENQ	ACK	BEL	BS	HT	LF	VT	FF	CR	SO	SI
1	DLE	DC1	DC2	DC3	DC4	NAK	SYN	ETB	CAN	EM	SUB	ESC	FS	GS	RS	US
2	SP	!	"	#	$	%	&	'	(	)	*	+	,	-	.	/
3	0	1	2	3	4	5	6	7	8	9	:	;	<	=	>	?
4	@	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O
5	P	Q	R	S	T	U	V	W	X	Y	Z	[	\	]	^	_
6	`	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o
7	p	q	r	s	t	u	v	w	x	y	z	{	|	}	~	DEL

Hexadecimal to ASCII conversion table

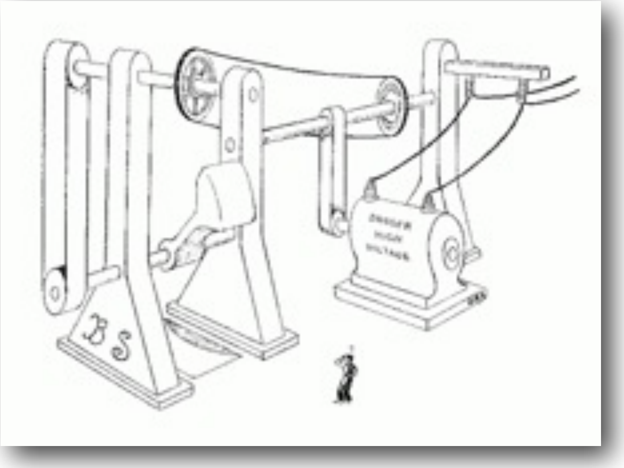
Universal data compression

US Patent 5,533,051 on "Methods for Data Compression", which is capable of compressing all files.

Slashdot reports of the Zero Space Tuner™ and BinaryAccelerator™.

> *" ZeoSync has announced a breakthrough in data compression that allows for 100:1 lossless compression of random data. If this is true, our bandwidth problems just got a lot smaller.... "*

Physical analog. Perpetual motion machines.



Gravity engine by Bob Schadewald

Universal data compression

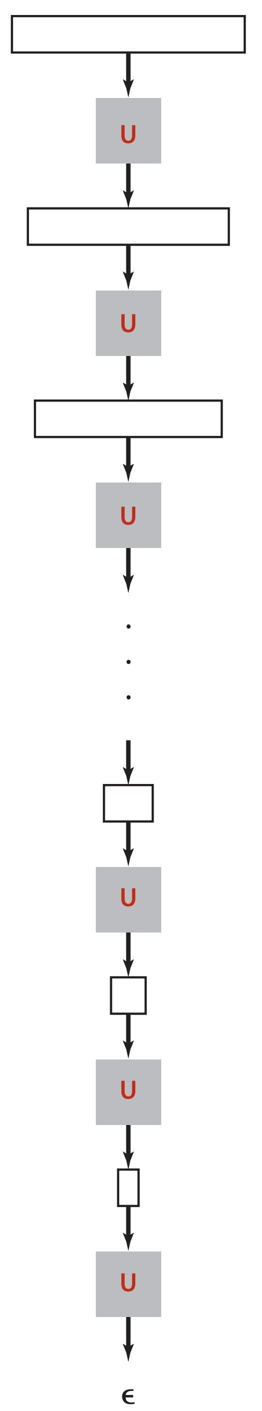
Proposition. No algorithm can compress every bitstring.

Pf 1. [by contradiction]

- Suppose you have a universal data compression algorithm U that can compress every bitstream.
- Given bitstring B_0, compress it to get smaller bitstring B_1.
- Compress B_1 to get a smaller bitstring B_2.
- Continue until reaching bitstring of size 0.
- Implication: all bitstrings can be compressed with 0 bits!

Pf 2. [by counting]

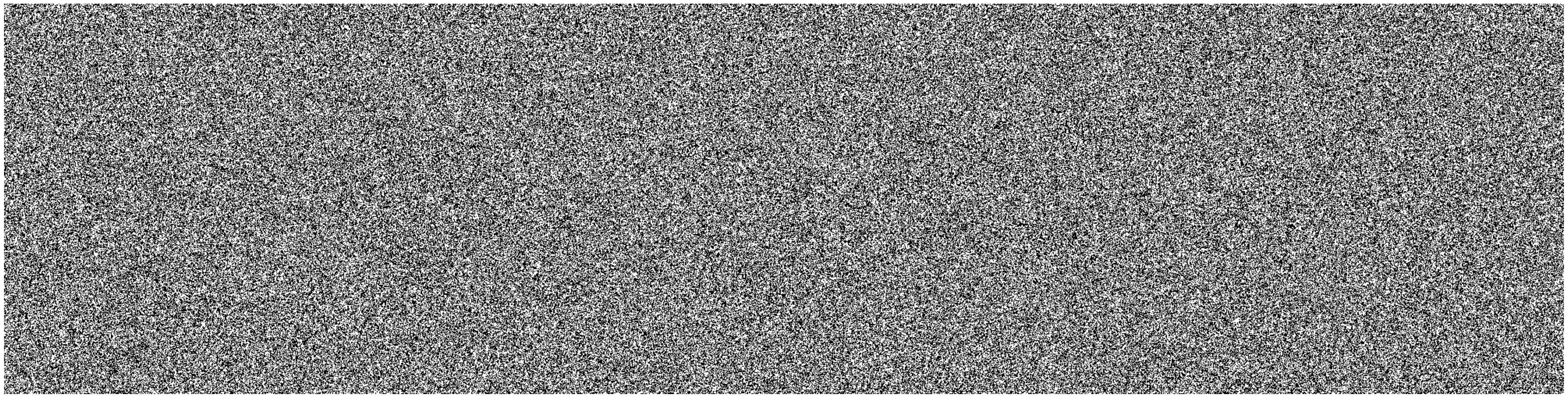
- Suppose your algorithm that can compress all 1000-bit strings.
- 2^{1000} possible bitstrings with 1000 bits.
- Only $1 + 2 + 4 + \ldots + 2^{998} + 2^{999}$ can be encoded with ≤ 999 bits.
- Similarly, only 1 in 2^{499} bitstrings can be encoded with ≤ 500 bits!



Universal data compression?

Undecidability

```
% java RandomBits | java PictureDump 2000 500
```



```
1000000 bits
```

A difficult file to compress: one million (pseudo-) random bits

```
public class RandomBits
{
   public static void main(String[] args)
   {
      int x = 11111;
      for (int i = 0; i < 1000000; i++)
      {
         x = x * 314159 + 218281;
         BinaryStdOut.write(x > 0);
      }
      BinaryStdOut.close();
   }
}
```

Rdenudcany in Enlgsih Inagugae

Q. How much redundancy is in the English language?

> *" ... randomising letters in the middle of words [has] little or no effect on the ability of skilled readers to understand the text. This is easy to denmtrasote. In a pubiltacion of New Scnieitst you could ramdinose all the letetrs, keipeng the first two and last two the same, and reibadailty would hadrly be aftcfeed. My ansaylis did not come to much beucase the thoery at the time was for shape and senqeuce retigcionon. Saberi's work sugsegts we may have some pofrweul palrlael prsooscers at work. The resaon for this is suerly that idnetiyfing coentnt by paarllel prseocsing speeds up regnicoiton. We only need the first and last two letetrs to spot chganes in meniang. "* — *Graham Rawlinson*

A. Quite a bit.

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Run-length encoding

Simple type of redundancy in a bitstream. Long runs of repeated bits.

`0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 1 1 1 1 1 1 0 0 0 0 0 0 0 1 1 1 1 1 1 1 1 1 1 1`

Representation. Use 4-bit counts to represent alternating runs of 0s and 1s: 15 0s, then 7 1s, then 7 0s, then 11 1s.

$$\underbrace{1\,1\,1\,1}_{15}\;\underbrace{0\,1\,1\,1}_{7}\;\underbrace{0\,1\,1\,1}_{7}\;\underbrace{1\,0\,1\,1}_{11} \longleftarrow \text{16 bits (instead of 40)}$$

Q. How many bits to store the counts?

A. We'll use 8.

Q. What to do when run length exceeds max count?

A. If longer than 255, intersperse runs of length 0.

Applications. JPEG, ITU-T T4 Group 3 Fax, ...

Run-length encoding: Java implementation

```
public class RunLength
{
   private final static int R = 256;

   public static void compress()
   {  /* see textbook */   }

   public static void expand()
   {
      boolean b = false;
      while (!BinaryStdIn.isEmpty())
      {
         char run = BinaryStdIn.readChar();      // read 8-bit count from standard input
         for (int i = 0; i < run; i++)
            BinaryStdOut.write(b);               // write 1 bit to standard output
         b = !b;
      }
      BinaryStdOut.close();
   }

}
```

read 8-bit count from standard input

write 1 bit to standard output

An application: compress a bitmap

Typical black-and-white-scanned image.

- 300 pixels/inch.
- 8.5-by-11 inches.
- 300 × 8.5 × 300 × 11 = 8.415 million bits.

Observation. Bits are mostly white.

Typical amount of text on a page.

40 lines × 75 chars per line = 3,000 chars.

```
% java BinaryDump 32 < q32x48.bin
00000000000000000000000000000000    32
00000000000000000000000000000000    32
00000000000000011111110000000000    15  7 10
00000000000011111111111111100000    12 15  5
00000000001111000011111111100000    10  4  4  9  5
00000000111100000000011111100000     8  4  9  6  5
00000001110000000000001111100000     7  3 12  5  5
00000011110000000000001111100000     6  4 12  5  5
00000111100000000000001111100000     5  4 13  5  5
00001111000000000000001111100000     4  4 14  5  5
00001111000000000000001111100000     4  4 14  5  5
00011110000000000000001111100000     3  4 15  5  5
00111110000000000000001111100000     2  5 15  5  5
00111110000000000000001111100000     2  5 15  5  5
00111110000000000000001111100000     2  5 15  5  5
00111110000000000000001111100000     2  5 15  5  5
00111110000000000000001111100000     2  5 15  5  5
00111110000000000000001111100000     2  5 15  5  5
00111110000000000000001111100000     2  5 15  5  5
00111110000000000000001111100000     2  5 15  5  5
00111110000000000000001111100000     2  5 15  5  5
00111111000000000000001111100000     2  6 14  5  5
00111111000000000000001111100000     2  6 14  5  5
00011111100000000000001111100000     3  6 13  5  5
00011111100000000000001111100000     3  6 13  5  5
00001111110000000000001111100000     4  6 12  5  5
00001111111000000000001111100000     4  7 11  5  5
00000111111100000000001111100000     5  7 10  5  5
00000011111111000000011111100000     6  8  7  6  5
00000001111111111111111111100000     7 20  5
00000000011111111111001111100000     9 11  2  5  5
00000000000000000000001111100000    22  5  5
00000000000000000000001111100000    22  5  5
00000000000000000000001111100000    22  5  5
00000000000000000000001111100000    22  5  5
00000000000000000000001111100000    22  5  5
00000000000000000000001111100000    22  5  5
00000000000000000000001111100000    22  5  5
00000000000000000000001111100000    22  5  5
00000000000000000000001111100000    22  5  5
00000000000000000000001111100000    22  5  5
00000000000000000000001111100000    22  5  5
00000000000000000000001111100000    22  5  5
00000000000000000000011111110000    21  7  4
00000000000000000011111111111100    18 12  2
00000000000000000111111111111110    17 14  1
00000000000000000000000000000000    32
00000000000000000000000000000000    32
1536 bits
```

7 1s

17 0s

A typical bitmap, with run lengths for each row

Data Compression

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- Huffman compression
- LZW compression

Variable-length codes

Use different number of bits to encode different chars.

Ex. Morse code: • • • – – – • • •

Issue. Ambiguity.

SOS ?

IAMIE ?

EEWNI ?

V7 ?

In practice. Use a medium gap to separate codewords.

Letters		Numbers	
A	•—	1	•————
B	—•••	2	••———
C	—•—•	3	•••——
D	—••	4	••••—
E	•	5	•••••
F	••—•	6	—••••
G	——•	7	——•••
H	••••	8	———••
I	••	9	————•
J	•———	0	—————
K	—•—		
L	•—••		
M	——		
N	—•		
O	———		
P	•——•		
Q	——•—		
R	•—•		
S	•••		
T	—		
U	••—		
V	•••—		
W	•——		
X	—••—		
Y	—•——		
Z	——••		

codeword for S is a prefix of codeword for V

Variable-length codes

Q. How do we avoid ambiguity?

A. Ensure that no codeword is a prefix of another.

Ex 1. Fixed-length code.

Ex 2. Append special stop char to each codeword.

Ex 3. General prefix-free code.

Codeword table

key	*value*
!	101
A	0
B	1111
C	110
D	100
R	1110

Compressed bitstring

011111110011001000111111100101 ← *30 bits*

A B RA CA DA B RA !

Codeword table

key	*value*
!	101
A	11
B	00
C	010
D	100
R	011

Compressed bitstring

11000111101011100110001111101 ← *29 bits*

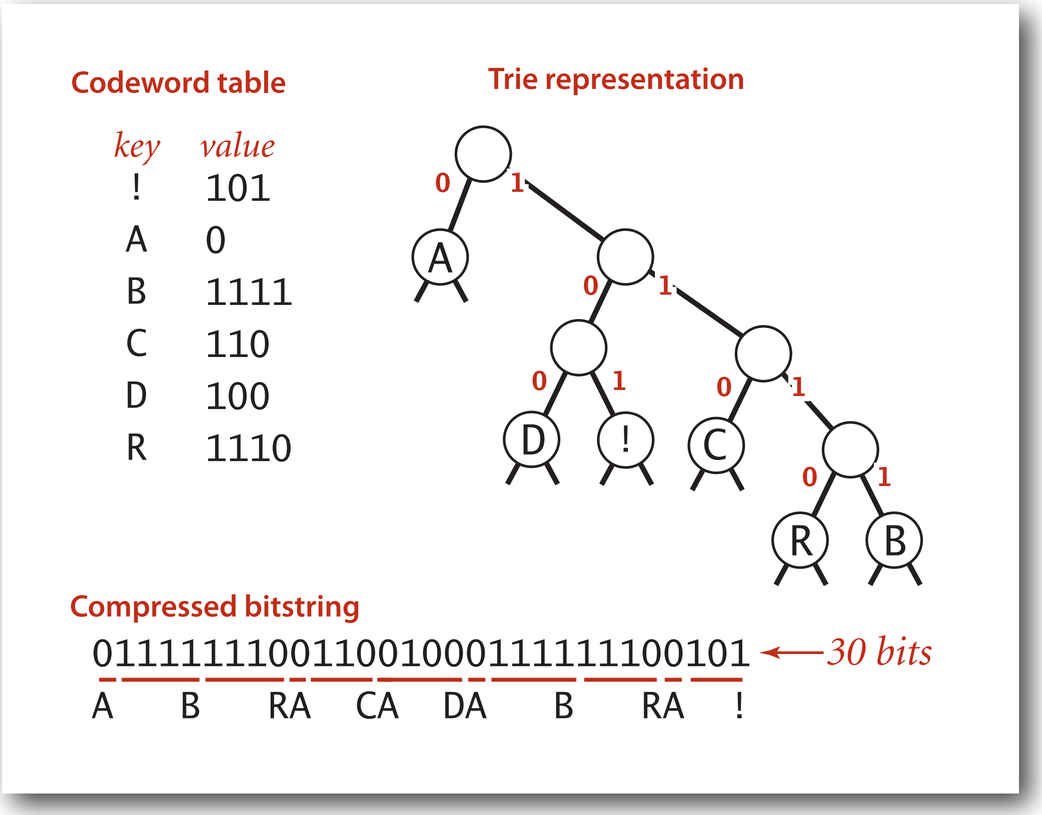
A B R A C A D A B R A !

Prefix-free codes: trie representation

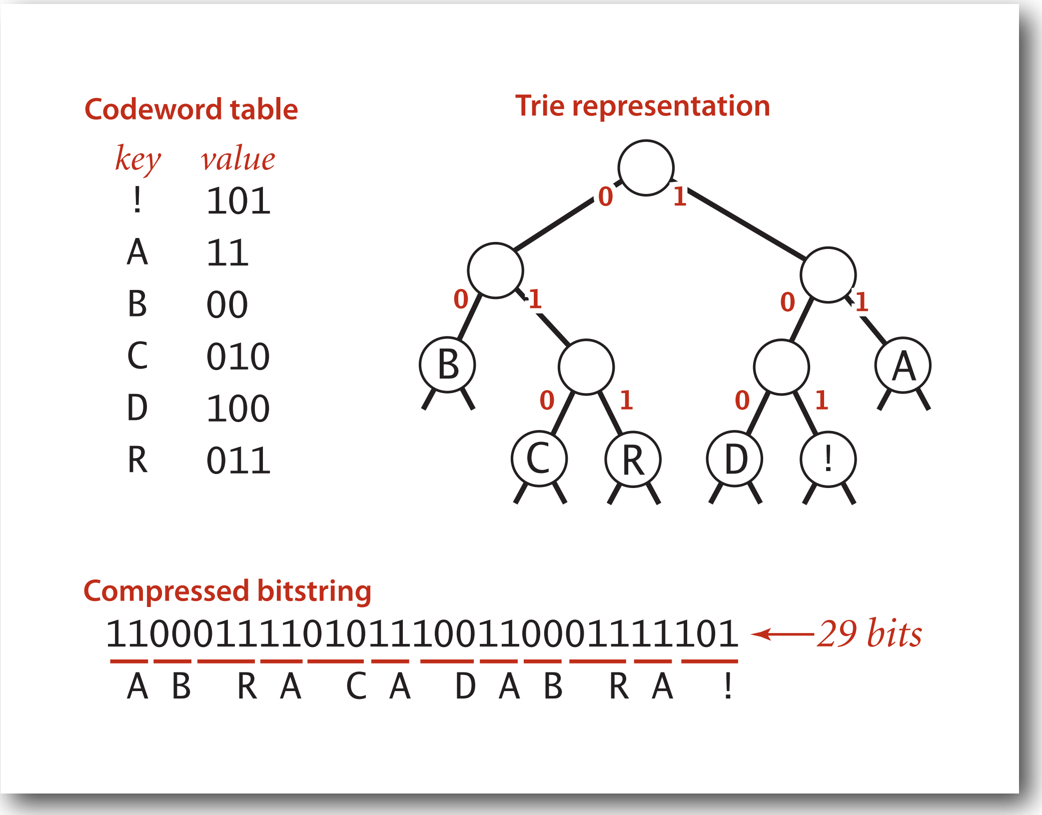
Q. How to represent the prefix-free code?

A. A binary trie!

- Chars in leaves.
- Codeword is path from root to leaf.



key	*value*
!	101
A	0
B	1111
C	110
D	100
R	1110



key	*value*
!	101
A	11
B	00
C	010
D	100
R	011

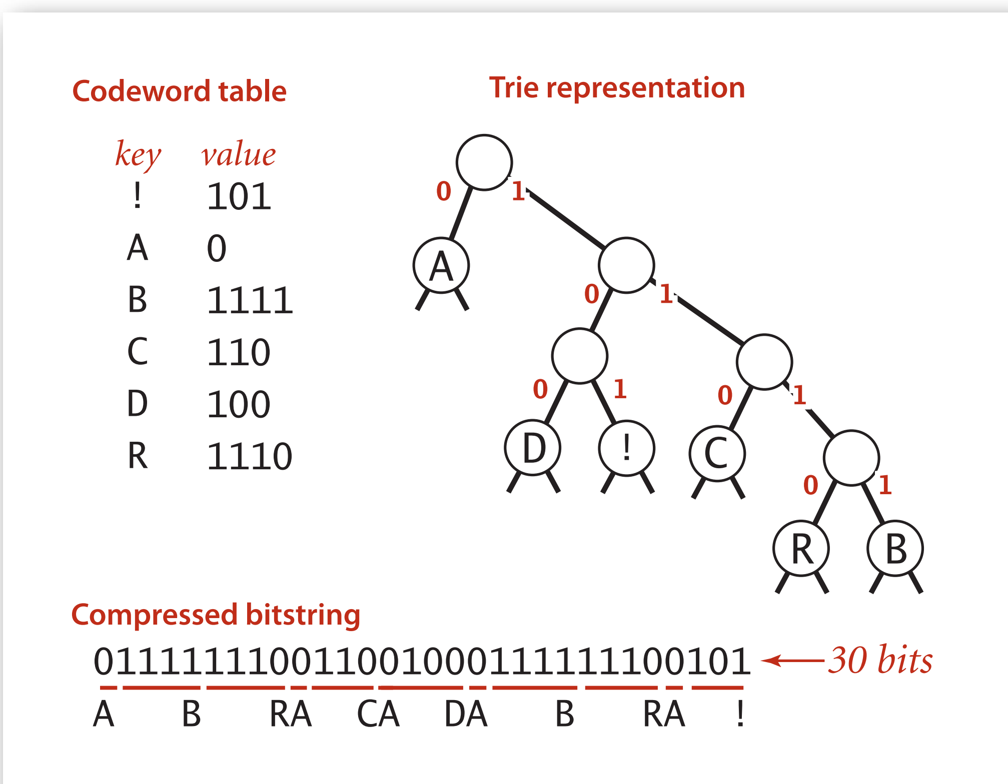
Prefix-free codes: compression and expansion

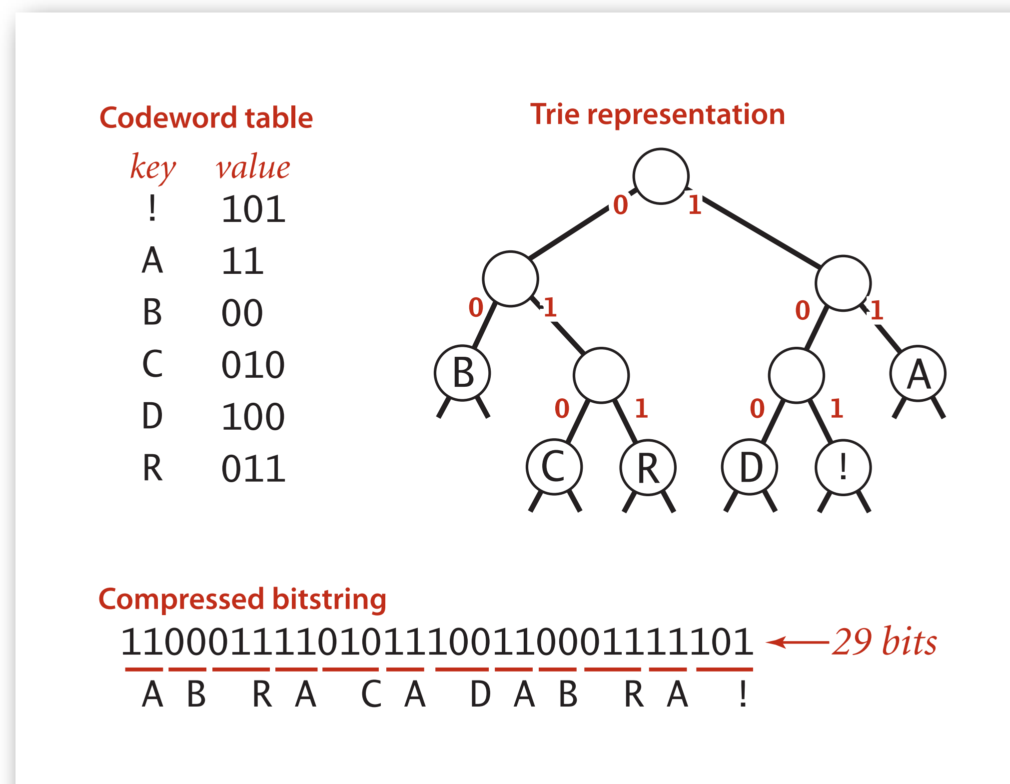
Compression.

- Method 1: start at leaf; follow path up to the root; print bits in reverse.
- Method 2: create ST of key-value pairs.

Expansion.

- Start at root.
- Go left if bit is 0; go right if 1.
- If leaf node, print char and return to root.





Huffman trie node data type

```
private static class Node implements Comparable<Node>
{
   private char ch;   // Unused for internal nodes.
   private int freq;  // Unused for expand.
   private final Node left, right;

   public Node(char ch, int freq, Node left, Node right)
   {
      this.ch    = ch;
      this.freq  = freq;
      this.left  = left;
      this.right = right;
   }

   public boolean isLeaf()
   {  return left == null && right == null; }

   public int compareTo(Node that)
   {  return this.freq - that.freq;  }
}
```

initializing constructor

is Node a leaf?

compare Nodes by frequency
(stay tuned)

Prefix-free codes: expansion

```
public void expand()
{
   Node root = readTrie();
   int N = BinaryStdIn.readInt();

   for (int i = 0; i < N; i++)
   {
      Node x = root;
      while (!x.isLeaf())
      {
         if (!BinaryStdIn.readBoolean())
            x = x.left;
         else
            x = x.right;
      }
      BinaryStdOut.write(x.ch);
   }
   BinaryStdOut.close();
}
```

read in encoding trie

read in number of chars

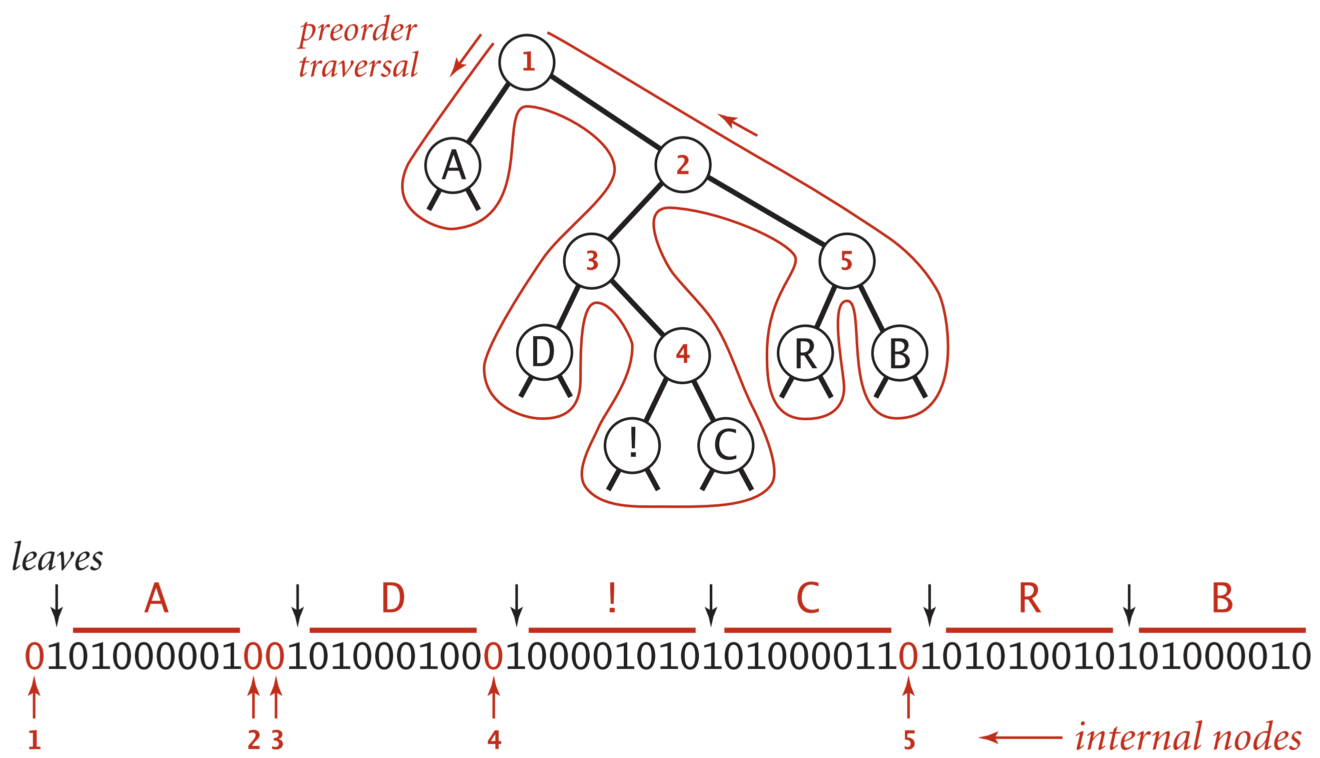
expand codeword for i^{th} char

Running time. Linear in input size (constant amount of work per bit read).

Prefix-free codes: how to transmit

Q. How to write the trie?

A. Write preorder traversal of trie; mark leaf and internal nodes with a bit.



Using preorder traversal to encode a trie as a bitstream

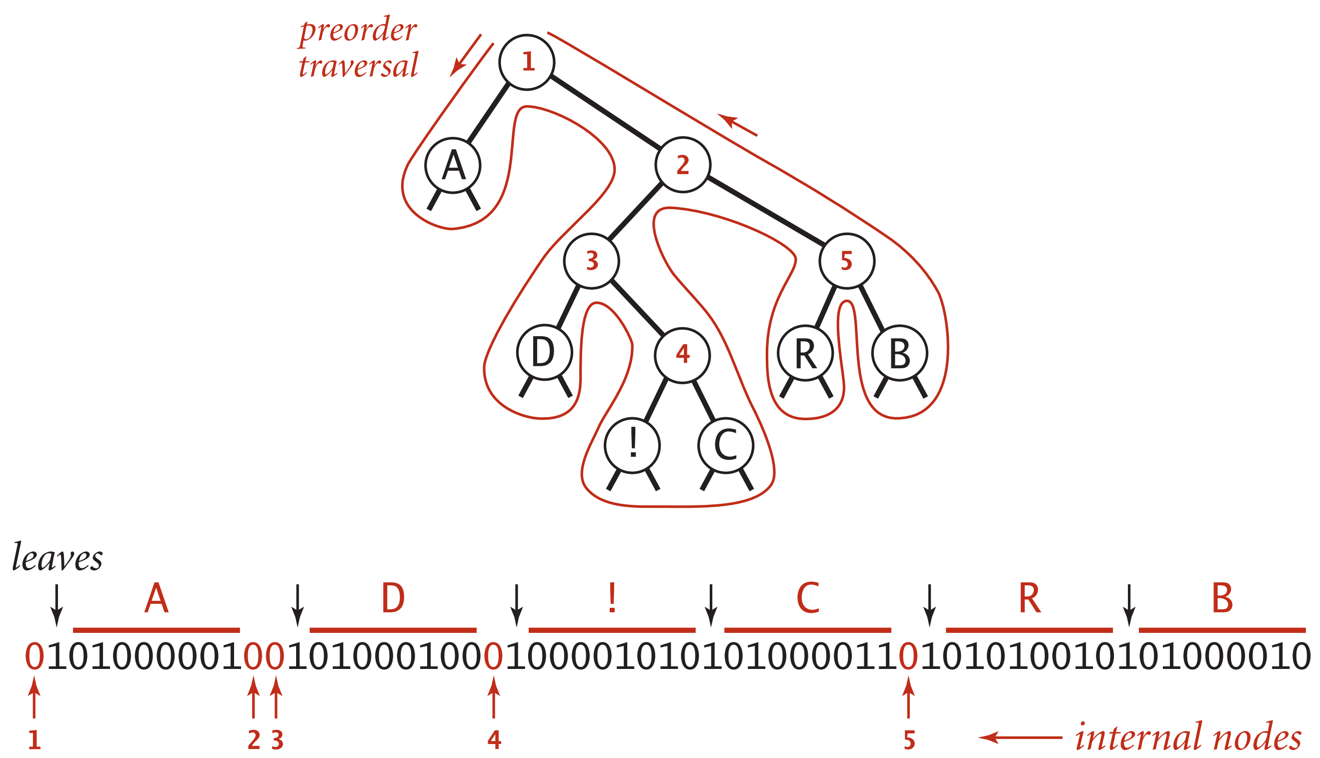
```
private static void writeTrie(Node x)
{
   if (x.isLeaf())
   {
      BinaryStdOut.write(true);
      BinaryStdOut.write(x.ch);
      return;
   }
   BinaryStdOut.write(false);
   writeTrie(x.left);
   writeTrie(x.right);
}
```

Note. If message is long, overhead of transmitting trie is small.

Prefix-free codes: how to transmit

Q. How to read in the trie?

A. Reconstruct from preorder traversal of trie.



Using preorder traversal to encode a trie as a bitstream

```
private static Node readTrie()
{
   if (BinaryStdIn.readBoolean())
   {
      char c = BinaryStdIn.readChar();
      return new Node(c, 0, null, null);
   }
   Node x = readTrie();
   Node y = readTrie();
   return new Node('\0', 0, x, y);
}
```

not used

Shannon-Fano codes

Q. How to find best prefix-free code?

Shannon-Fano algorithm:

- Partition symbols S into two subsets S_0 and S_1 of (roughly) equal frequency.
- Codewords for symbols in S_0 start with 0; for symbols in S_1 start with 1.
- Recur in S_0 and S_1.

Problem 1. How to divide up symbols?

Problem 2. Not optimal!

Huffman codes

Q. How to find best prefix-free code?

David Huffman

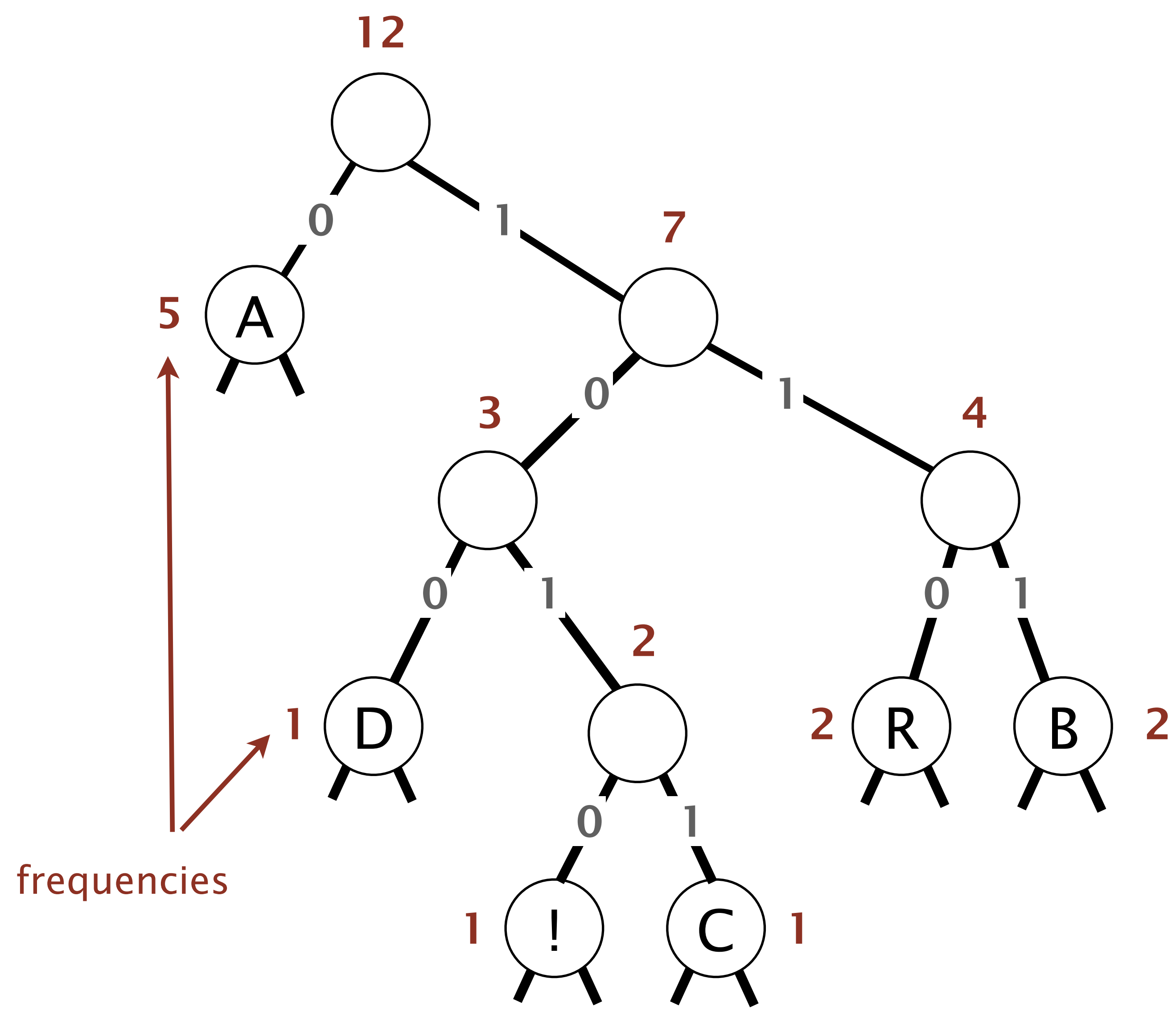
Huffman algorithm:

- Count frequency **`freq[i]`** for each char **`i`** in input.
- Start with one node corresponding to each char **`i`** (with weight **`freq[i]`**).
- Repeat until single trie formed:
  - select two tries with min weight **`freq[i]`** and **`freq[j]`**
  - merge into single trie with weight **`freq[i] + freq[j]`**

Applications. JPEG, MP3, MPEG, PKZIP, GZIP, PDF, …

Constructing a Huffman encoding trie

char	freq	encoding
A	5	0
B	2	1 1 1
C	1	1 0 1 1
D	1	1 0 0
R	2	1 1 0
!	1	1 0 1 0



Huffman code construction for A B R A C A D A B R A !

Constructing a Huffman encoding trie: Java implementation

```
private static Node buildTrie(int[] freq)
{
    MinPQ<Node> pq = new MinPQ<Node>();
    for (char i = 0; i < R; i++)
       if (freq[i] > 0)
          pq.insert(new Node(i, freq[i], null, null));

    while (pq.size() > 1)
    {
       Node x = pq.delMin();
       Node y = pq.delMin();
       Node parent = new Node('\0', x.freq + y.freq, x, y);
       pq.insert(parent);
    }

    return pq.delMin();
}
```

initialize PQ with singleton tries

merge two smallest tries

not used

total frequency

two subtries

Huffman encoding summary

Proposition. [Huffman 1950s] Huffman algorithm produces an optimal prefix-free code.

Pf. See textbook.

no prefix-free code uses fewer bits

Implementation.

- Pass 1: tabulate char frequencies and build trie.
- Pass 2: encode file by traversing trie or lookup table.

Running time. Using a binary heap $\Rightarrow$ $O(N + R \log R)$.

input size

alphabet size

Q. Can we do better? [stay tuned]

Data Compression

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- LZW compression

Statistical methods

Static model. Same model for all texts.

- Fast.
- Not optimal: different texts have different statistical properties.
- Ex: ASCII, Morse code.

Dynamic model. Generate model based on text.

- Preliminary pass needed to generate model.
- Must transmit the model.
- Ex: Huffman code.

Adaptive model. Progressively learn and update model as you read text.

- More accurate modeling produces better compression.
- Decoding must start from beginning.
- Ex: LZW.

Lempel-Ziv-Welch compression example

input	A	B	R	A	C	A	D	A	B	R	A	B	R	A	B	R	A
matches	A	B	R	A	C	A	D	A	B	R	A	B	R	A	B	R	A
value	41	42	52	41	43	41	44	81		83		82		88			41

LZW compression for ABRACADABRABRABRA

key	value
...	
A	41
B	42
C	43
D	44
...	

key	value
AB	81
BR	82
RA	83
AC	84
CA	85
AD	86

key	value
DA	87
ABR	88
RAB	89
BRA	8A
ABRA	8B

codeword table

Lempel-Ziv-Welch compression

LZW compression.

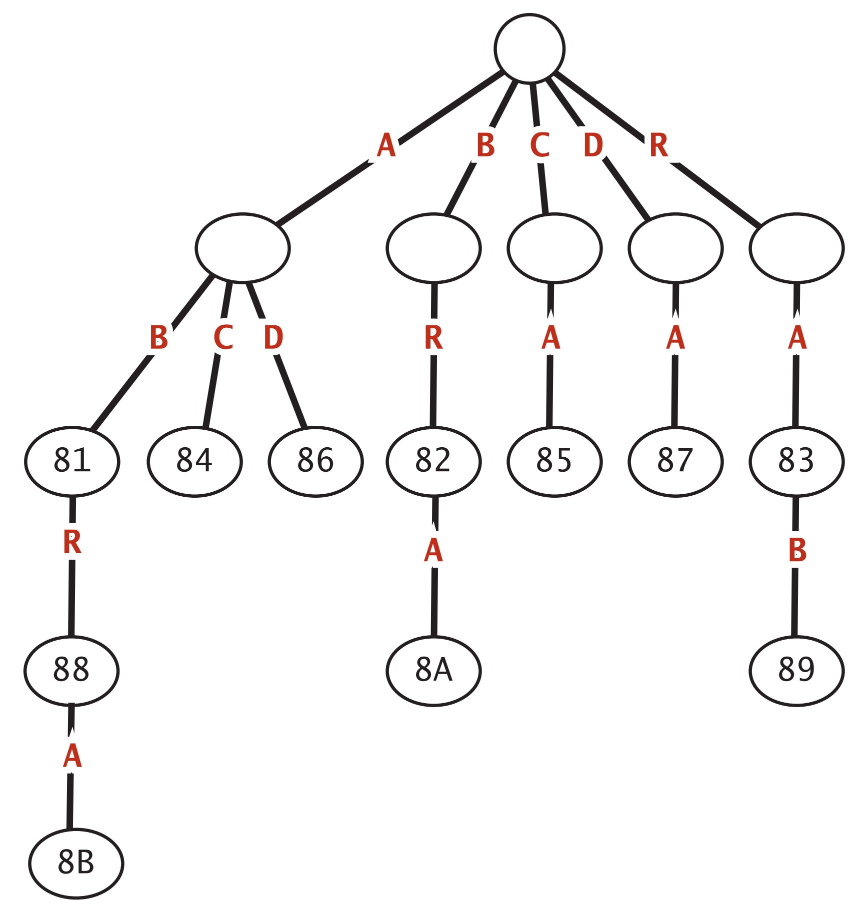
- Create ST associating W-bit codewords with string keys.
- Initialize ST with codewords for single-char keys.
- Find longest string **s** in ST that is a prefix of unscanned part of input.
- Write the W-bit codeword associated with **s**.
- Add **s + c** to ST, where **c** is next char in the input.

LZW compression for A B R A C A D A B R A B R A B R A

Representation of LZW code table

Q. How to represent LZW code table?

A. A trie: supports efficient longest prefix match.



Trie representation of LZW code table

Remark. Every prefix of a key in encoding table is also in encoding table.

Lempel-Ziv-Welch expansion example

value	41	42	52	41	43	41	44	81	83	82	88	41	80
output	A	B	R	A	C	A	D	A B	R A	B R	A B R	A	

LZW expansion for 41 42 52 41 43 41 44 81 83 82 88 41 80

value	key
…	…
41	A
42	B
43	C
44	D
…	…

value	key
81	AB
82	BR
83	RA
84	AC
85	CA
86	AD

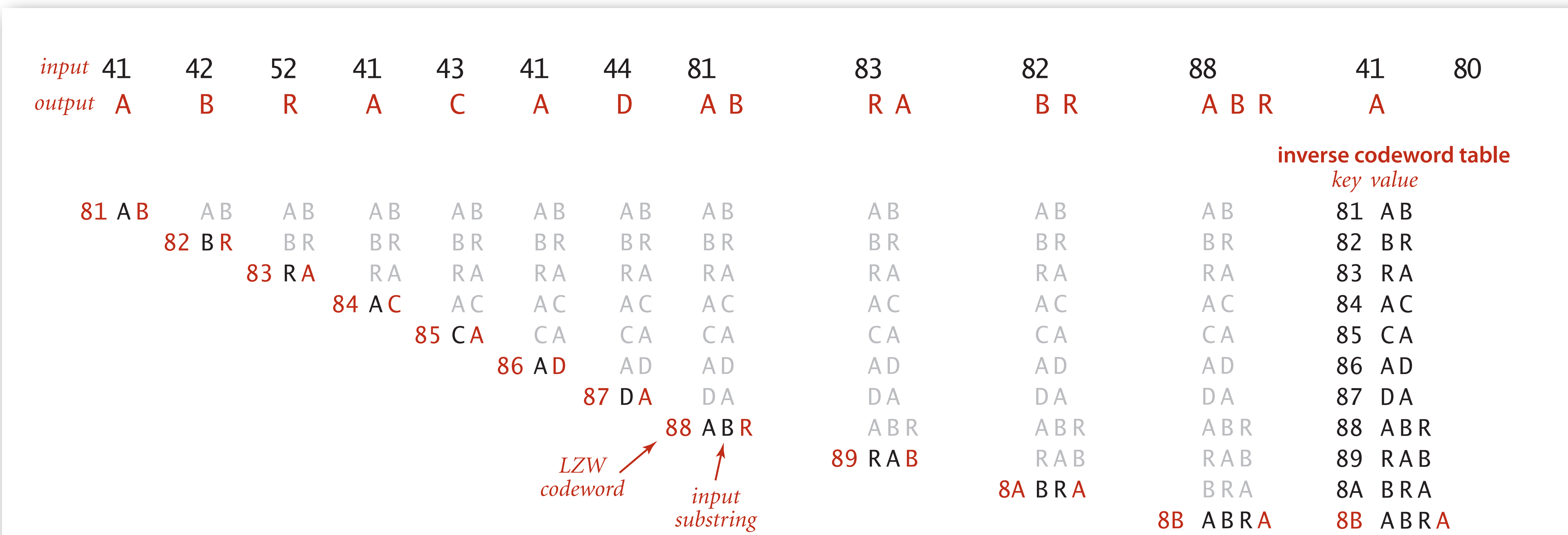
value	key
87	DA
88	ABR
89	RAB
8A	BRA
8B	ABRA

codeword table

LZW expansion

LZW expansion.

- Create ST associating string values with W-bit keys.
- Initialize ST to contain with single-char values.
- Read a W-bit key.
- Find associated string value in ST and write it out.
- Update ST.



LZW expansion for 41 42 52 41 43 41 44 81 83 82 88 41 80

LZW example: tricky situation

input	A	B	A	B	A	B	A	
matches	A	B	A B		A B A			
value	41	42	81		83			80

LZW compression for ABABABA

key	value
...	
A	41
B	42
C	43
D	44
...	

key	value
AB	81
BA	82
ABA	83

codeword table

LZW example: tricky situation

value	41	42	81	83	80
output	A	B	A B	A B A	

need to know which key has value 83 before it is in ST!

LZW expansion for 41 42 81 83 80

value	key
…	…
41	A
42	B
43	C
44	D
…	…

value	key
81	AB
82	BA
83	ABA

codeword table

LZW in the real world

Lempel-Ziv and friends.

- LZ77.
- LZ78.
- LZW.
- Deflate = LZ77 variant + Huffman.

PNG: LZ77.

7zip, gzip, jar, pdf, java.util.zip: deflate.

Unix compress: LZW.

Pkzip: LZW + Shannon-Fano.

GIF, TIFF, V.42bis modem: LZW.

Google: zlib which is based on deflate.

never expands a file

Lossless data compression benchmarks

year	scheme	bits / char
1967	ASCII	7.00
1950	Huffman	4.70
1977	LZ77	3.94
1984	LZMW	3.32
1987	LZH	3.30
1987	move-to-front	3.24
1987	LZB	3.18
1987	gzip	2.71
1988	PPMC	2.48
1994	SAKDC	2.47
1994	PPM	2.34
1995	Burrows-Wheeler	2.29
1997	BOA	1.99
1999	RK	1.89

final programming assignment

data compression using Calgary corpus

Data compression summary

Lossless compression.

- Represent fixed-length symbols with variable-length codes. [Huffman]
- Represent variable-length symbols with fixed-length codes. [LZW]

Lossy compression. [not covered in this course]

- JPEG, MPEG, MP3, …
- FFT, wavelets, fractals, …

Theoretical limits on compression. Shannon entropy.

Practical compression. Use extra knowledge whenever possible.