## MPI types, Scatter and Scatterv

## MPI types, Scatter and Scattery

| 0  |    | 2  | 3  | 4  | 5  |
|----|----|----|----|----|----|
| 6  | 7  | 8  | 9  | 10 |    |
| 12 | 13 | 14 | 15 | 16 | 17 |
| 18 | 19 | 20 | 21 | 22 | 23 |
| 24 | 25 | 26 | 27 | 28 | 29 |
| 30 | 31 | 32 | 33 | 34 | 35 |

Logical and physical layout of a C/C++ array in memory.

A = malloc(6\*6\*sizeof(int));



### MPI\_Scatter

sendbuf, sendcount, sendtype valid only at the sending process

### Equal number elements to all processors

A

| 0  | _  | 2  | 3  | 4  | 5  |
|----|----|----|----|----|----|
| 6  | 7  | 8  | 9  | 10 |    |
| 12 | 13 | 14 | 15 | 16 | 17 |
| 18 | 19 | 20 | 21 | 22 | 23 |
| 24 | 25 | 26 | 27 | 28 | 29 |
| 30 | 31 | 32 | 33 | 34 | 35 |

```
int MPI_Scatter(A, 9, MPI_Int, B, 9,
              MPI Int, 0,
              MPI COMM WORLD)
          11 12 13 14 15 16 17
   18 19 20 21 22 23 24 25 26
P<sub>3</sub> 27 28 29 30 31 32 33 34
```

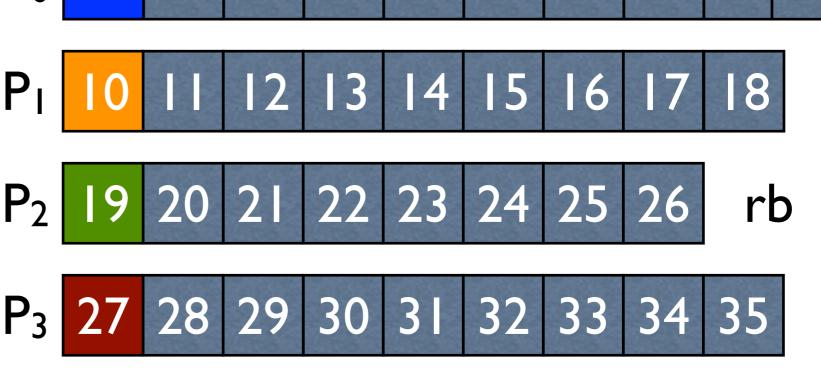
### MPI Scattery

```
int MPI Scatter(
    const void *sendbuf, // data to send
    const int *sendcounts, // sent to each process
                        // where in sendbuf
    const int* displ
                             // sent data is
    MPI_Datatype sendtype, // type of data sent
    void *recvbuf,  // where received
    int recvcount, // how much to receive
    MPI Datatype recytype, // type of data received
                             // sending process
    int root,
                             // communicator
    MPI Comm comm)
```

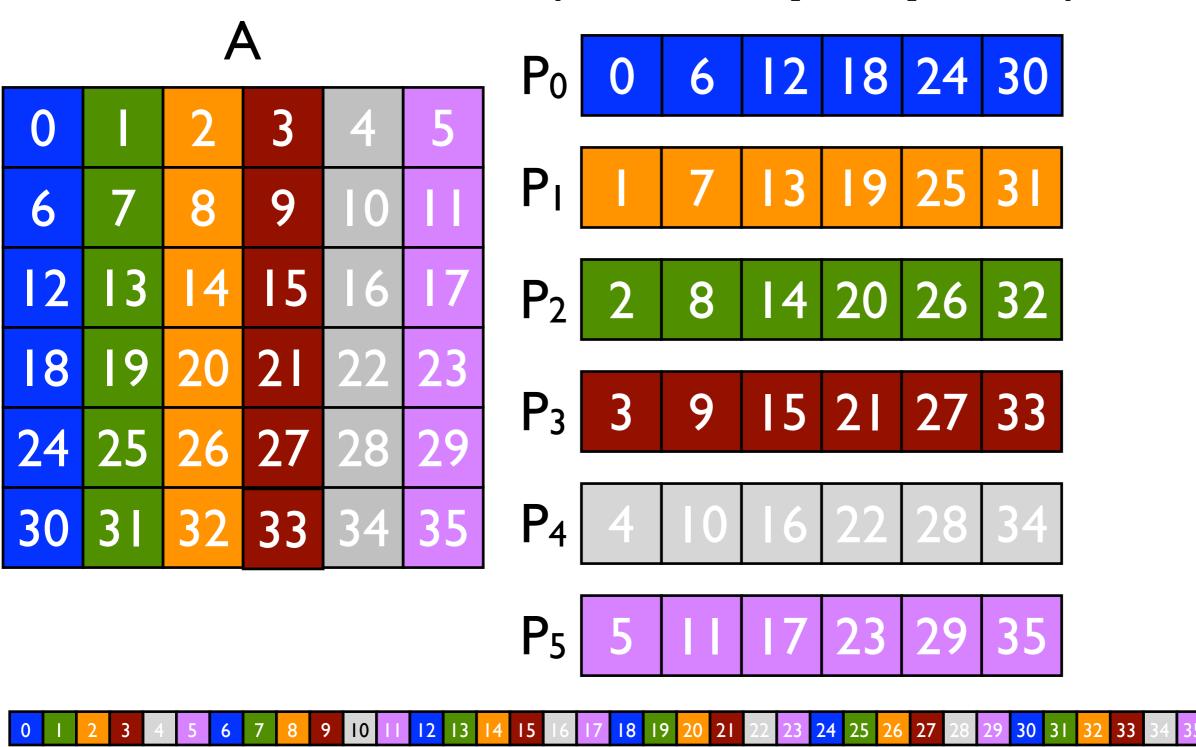
sendbuf, sendcount, sendtype valid only at the sending process

## Specify the number elements sent to each processor





## What if we want to scatter columns (C array layout)



### MPI\_Type\_vector

Allows a type to be created that puts together blocks of elements in a vector into another vector.

Note that a 2-D array in contiguous memory can be treated as a 1-D vector.

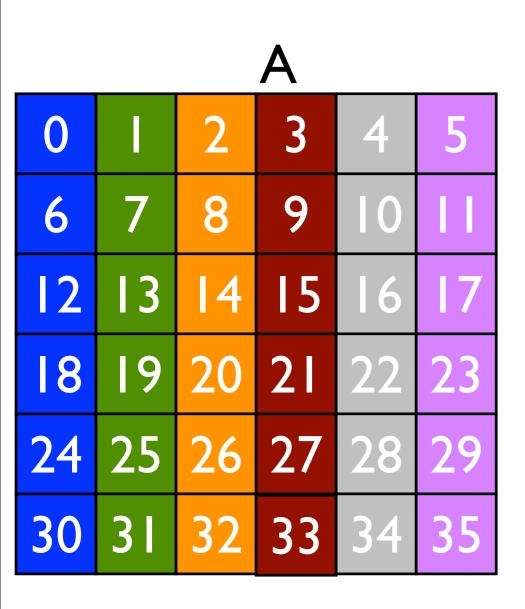
### MPI\_Type\_create\_resized

```
int MPI_Type_create_resized(
    MPI_Datatype oldtype, // type being resized
    MPI_Aint lb, // new lower bound
    MPI_Aint extent, // new extent ("length")
    MPI_Datatype *newtype) // resized type name
)
```

Allows a new size (or extent) to be assigned to an existing type.

Allows MPI to determine how far from an object O1 the next adjacent object O2 is. As we will see this is often necessitated because we treat a logically 2-D array as a 1-D vector.

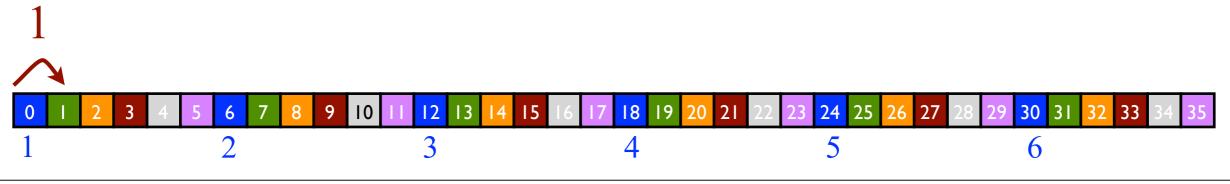
### Using MPI\_Type\_vector



## MPI\_Type\_vector: defining the type

```
9
        15
    14
25 26
       27
   32 | 33
```

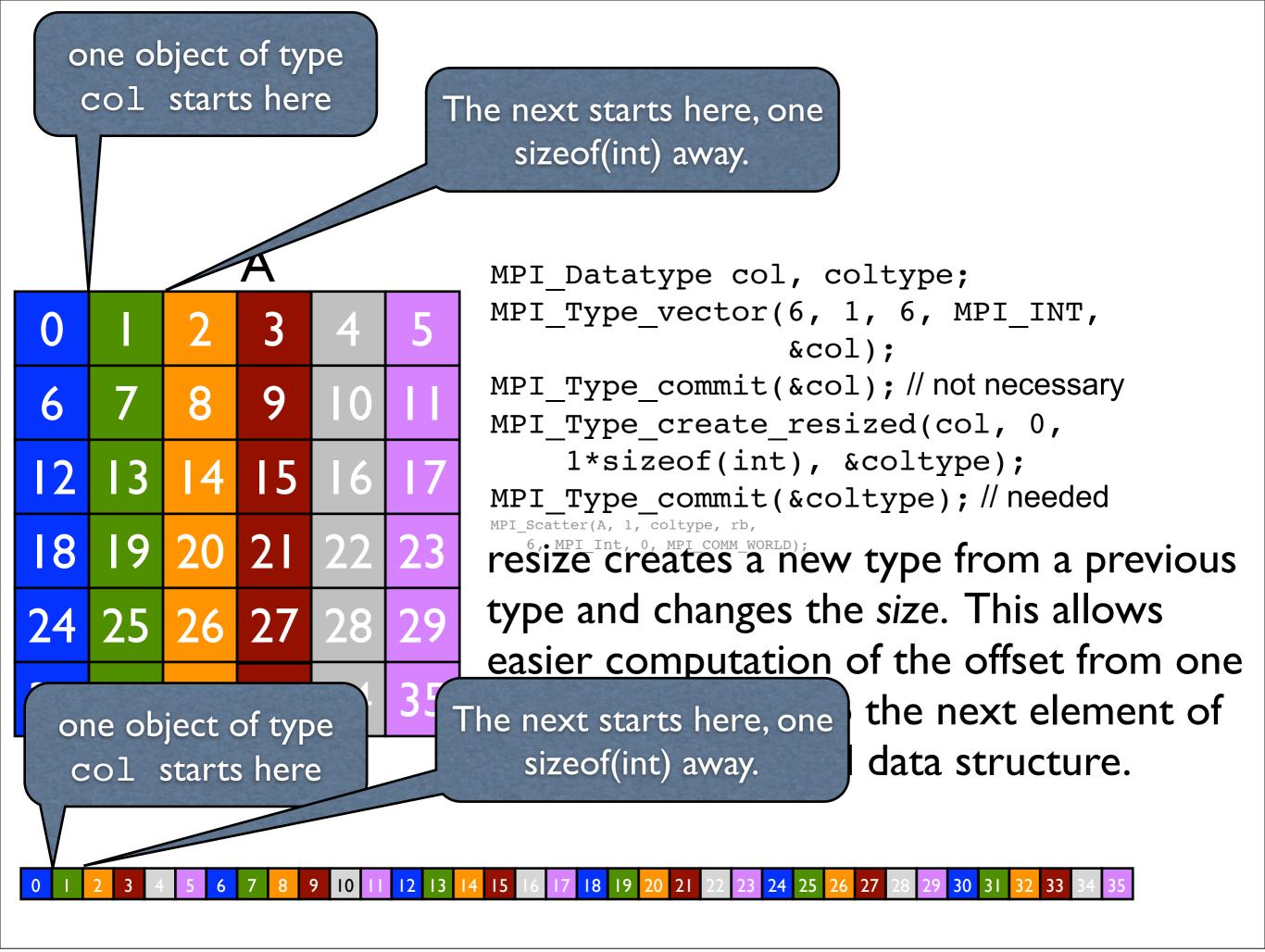
Each block is made of 6 ints of length one, and the new block starts I positions in the linearized array from the start of the previous block.



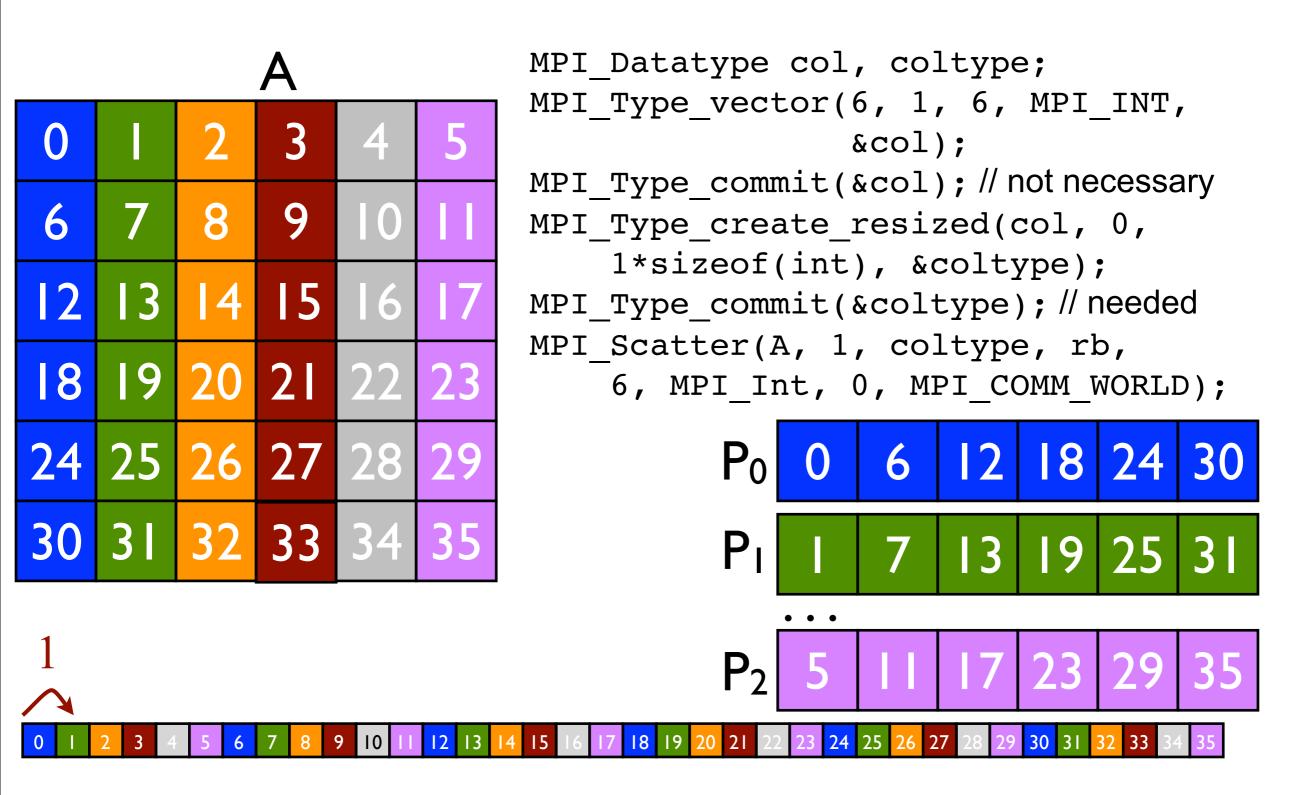
### Using MPI\_type\_create\_resized

9 4 33

resize creates a new type from a previous type and changes the *size*. This allows easier computation of the offset from one element of a type to the next element of a type in the original data structure.

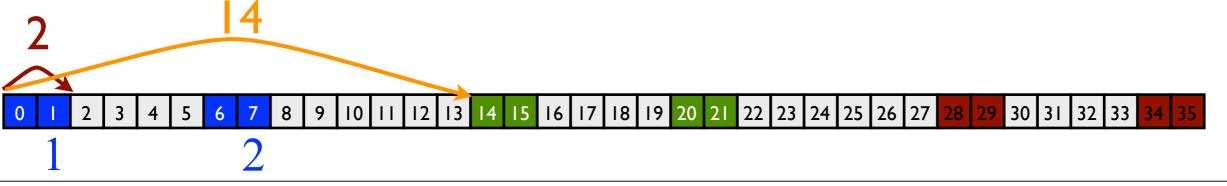


### The result of the communication



#### Scattering diagonal blocks

| A  |    |    |    |    |    |  |
|----|----|----|----|----|----|--|
| 0  | _  | 2  | 3  | 4  | 5  |  |
| 6  | 7  | 8  | 9  | 10 | П  |  |
| 12 | 13 | 14 | 15 | 16 | 17 |  |
| 18 | 19 | 20 | 21 | 22 | 23 |  |
| 24 | 25 | 26 | 27 | 28 | 29 |  |
| 30 | 31 | 32 | 33 | 34 | 35 |  |

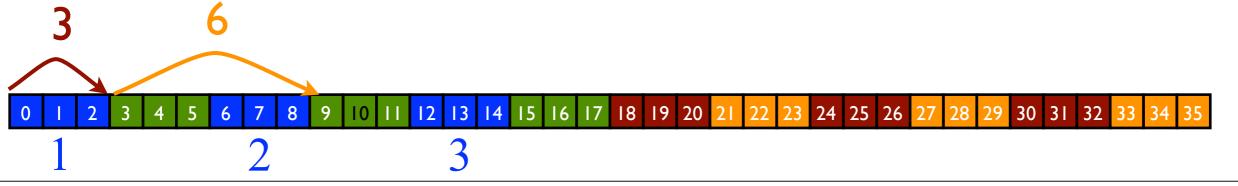


### Scattering the blocks

| A  |    |    |    |    |    |  |
|----|----|----|----|----|----|--|
| 0  | _  | 2  | 3  | 4  | 5  |  |
| 6  | 7  | 8  | 9  | 10 | П  |  |
| 12 | 13 | 14 | 15 | 16 | 17 |  |
| 18 | 19 | 20 | 21 | 22 | 23 |  |
| 24 | 25 | 26 | 27 | 28 | 29 |  |
| 30 | 31 | 32 | 33 | 34 | 35 |  |

```
MPI Datatype block, blocktype;
MPI Type vector(2, 2, 14, MPI INT,
                 &block);
MPI Type commit(&block); // not necessary
MPI Type create resized(block, 0,
    14*sizeof(int), &blocktype);
MPI Type commit(&blocktype); // needed
int MPI Scatter(A, 1, blocktype, B, 4,
                MPI Int, 0,
                MPI COMM WORLD)
          B
        15 20
```

## The Type\_vector statement describing this



### The create\_resize statement for this

```
Α
        9
   14
       15
   20
              23
       27
          28
              29
25 | 26 |
          34 35
   32 33
```

Distance between start of blocks varies, but are multiples of 3. Use MPI\_Scatterv



### Sending the data

```
Α
                      MPI Datatype block, blocktype;
                       int disp = \{0, 1, 6, 7\}
                       int scount = \{1, 1, 1, 1\}
                       int rcount = \{9, 9, 9, 9\}
         9
                      MPI Type vector(3, 3, 6, MPI INT,
                                        &block);
    14
                 17
                      MPI Type commit(&block); // not necessary
                      MPI Type create resized(block, 0,
    20
                           3*sizeof(int), &blocktype);
                      MPI Type commit(&blocktype); // needed
                 29
25 26
            28
                       int MPI_Scatterv(A, scount, displ,
                                       blocktype, rb, rcount,
    32 33
            34 35
                                       MPI Int, 0,
                                       MPI COMM WORLD)
```

displacement is sizeof(blockcol)

# Matrix Multiply Cannon's Algorithm

Useful for the small project

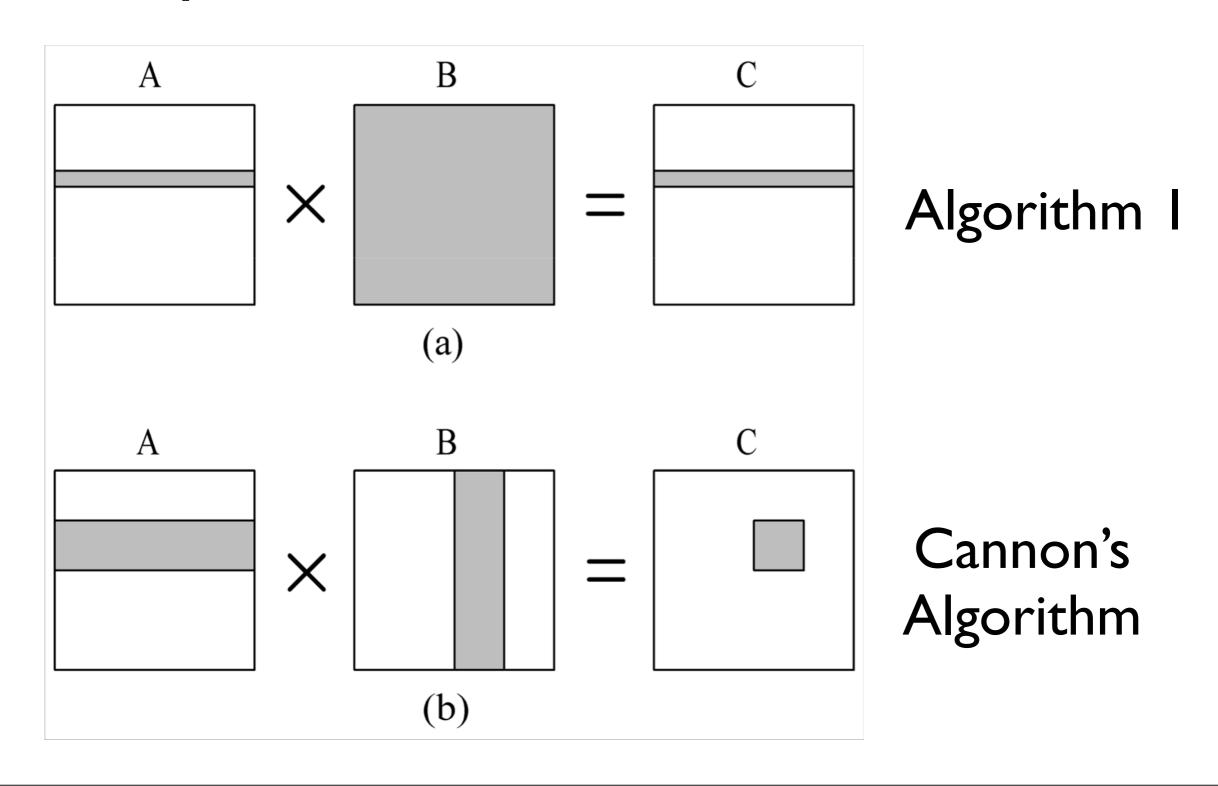
# Parallel Algorithm 2 (Cannon's Algorithm)

- Associate a primitive task with each matrix element
- Agglomerate tasks responsible for a square (or nearly square) block of C (the result matrix)
- Computation-to-communication ratio rises to  $n \mid \sqrt{p}$  (same total computation, more computation per communication)
  - $2n/p < n/\sqrt{p}$  when p > 4

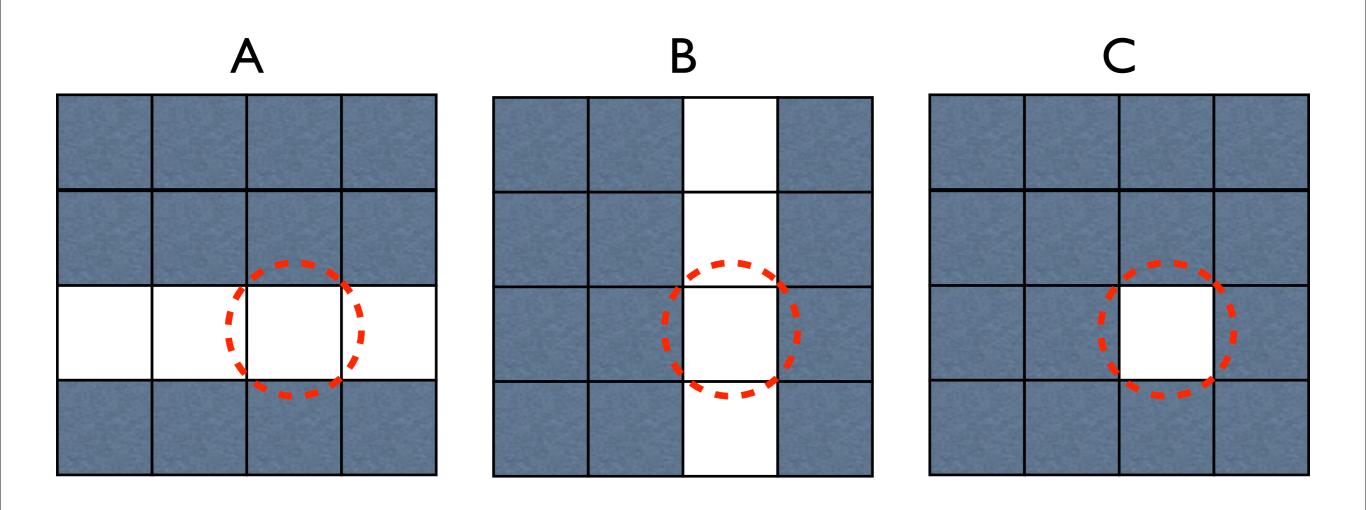
### A simplifying assumption

- Assume that
  - A, B and (consequently) C are n x n square matrices
  - $\sqrt{p}$  is an integer, and
  - $n = k \cdot \sqrt{p}$ , k an integer (i.e. n is a multiple of  $\sqrt{p}$

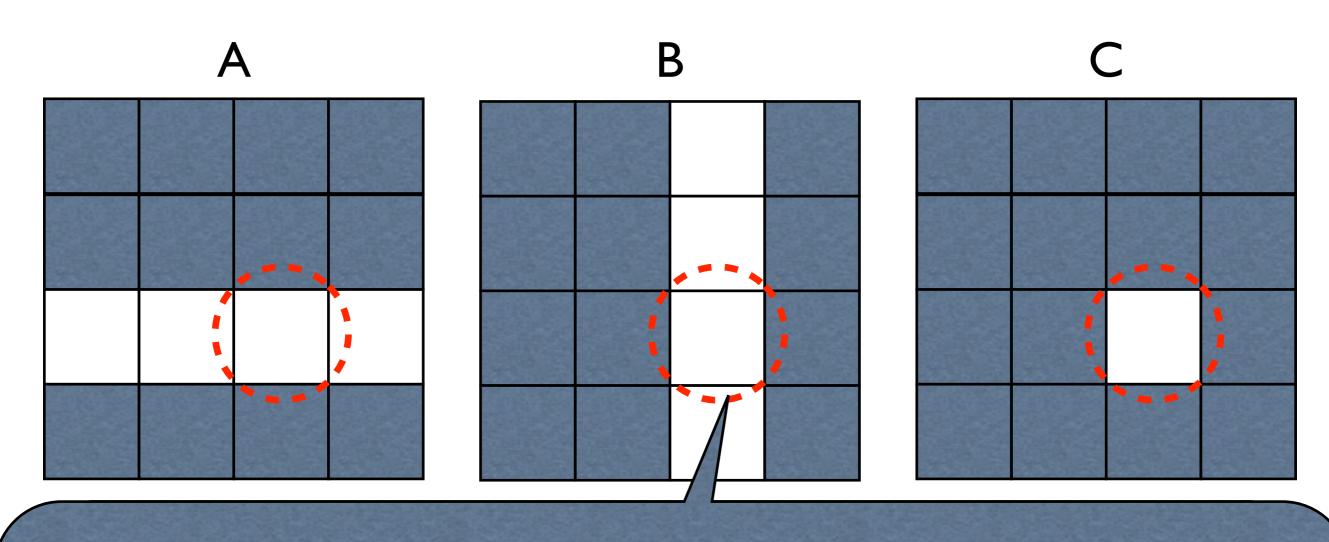
### Elements of A and B Needed to Compute a Process's Portion of C



## Blocks need to compute a C element



## Blocks need to compute a C element

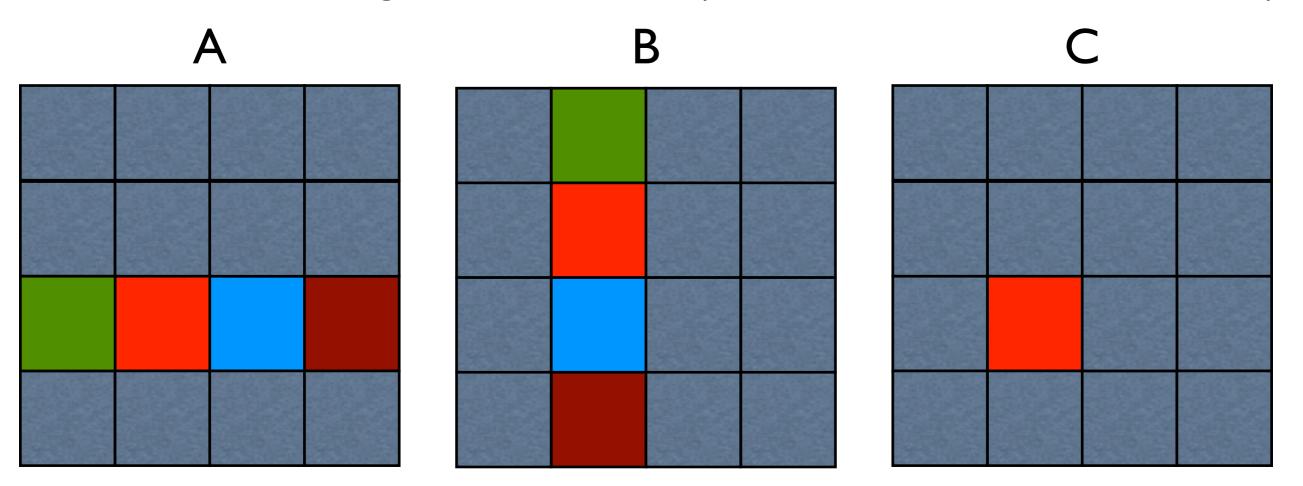


Processor that owns these blocks fully computes value of this C block (but needs more than these)

#### Blocks needed to compute a C element

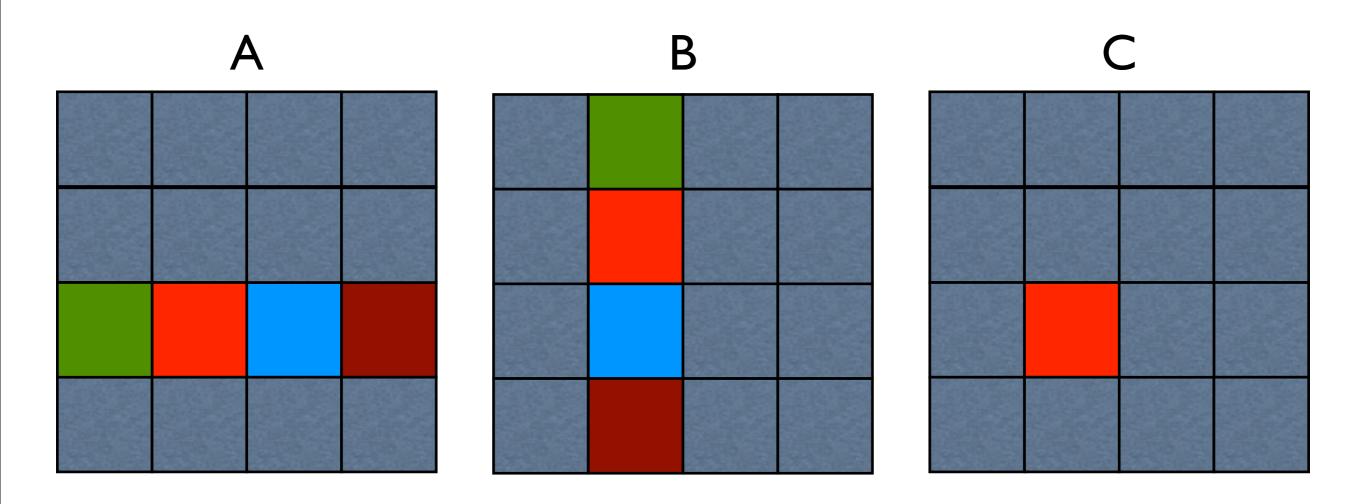
Processor  $P_{3,2}$  needs, at some point, to simultaneously hold the green A and B blocks, the red A and B blocks, the blue A and B blocks, and the cayenne A and B blocks.

With the current data layout it cannot do useful work because it does not contain matching A and B blocks (it has a red A and blue B block)



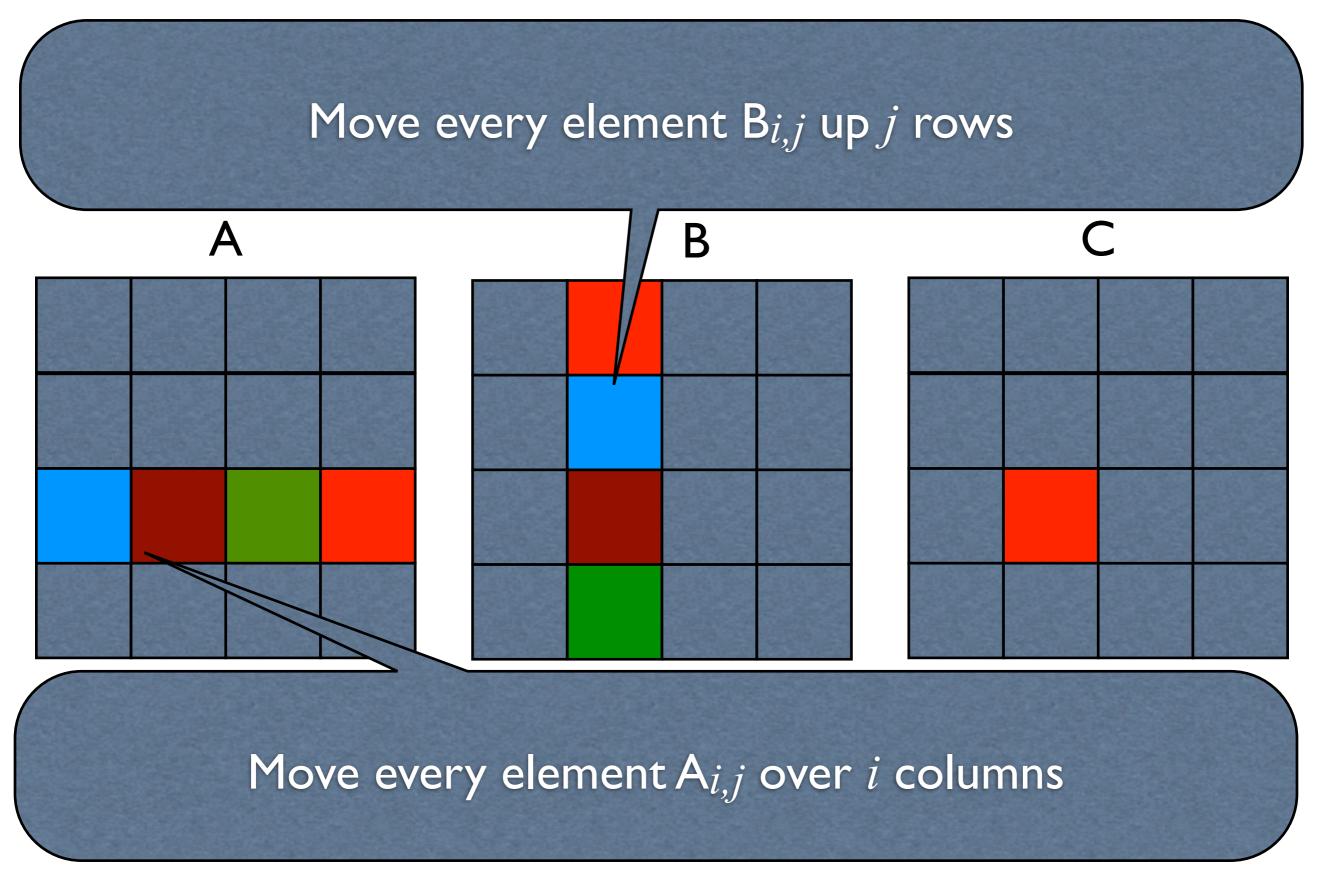
#### Blocks needed to compute a C element

We need to rearrange the data so that every block has useful work to do

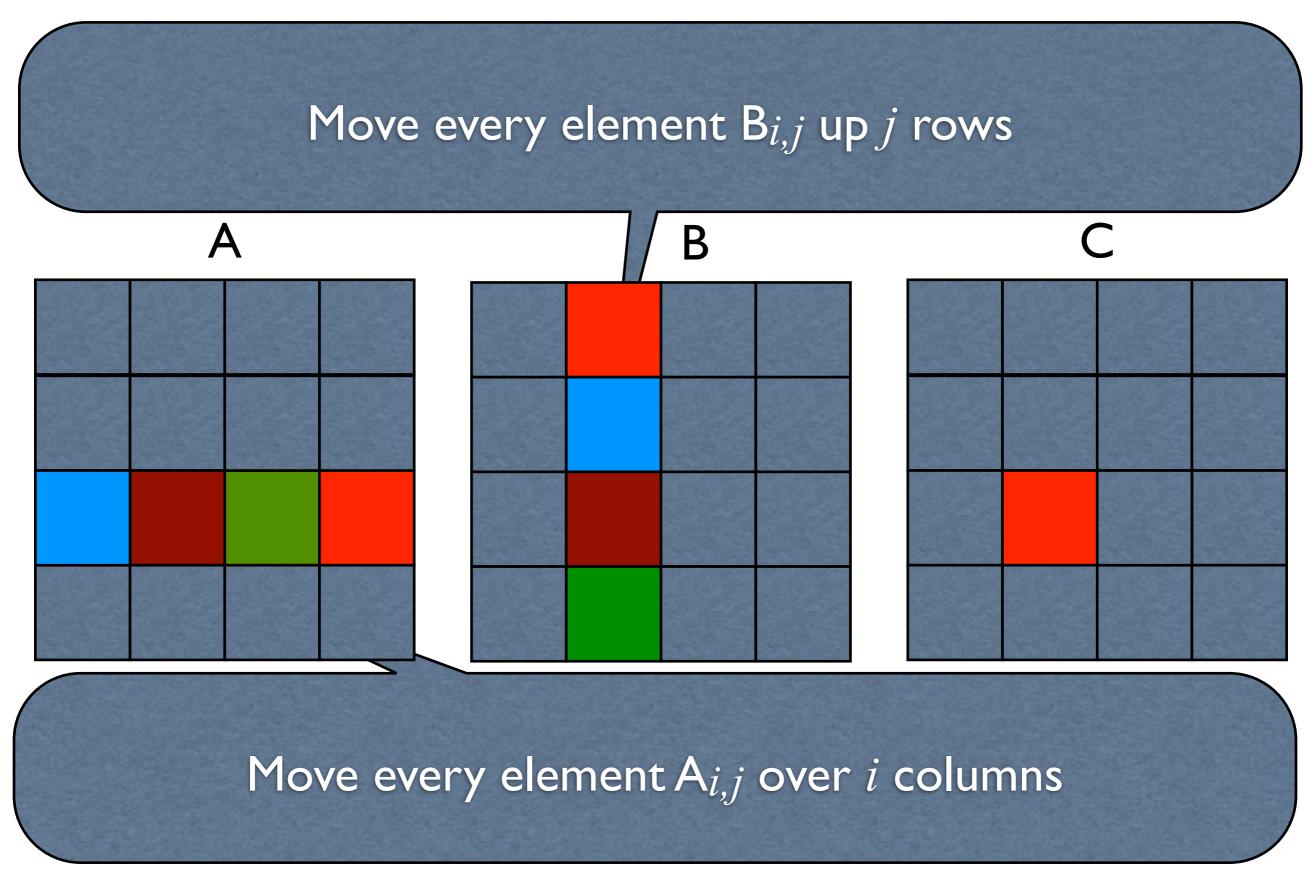


The initial data configuration does not provide for this

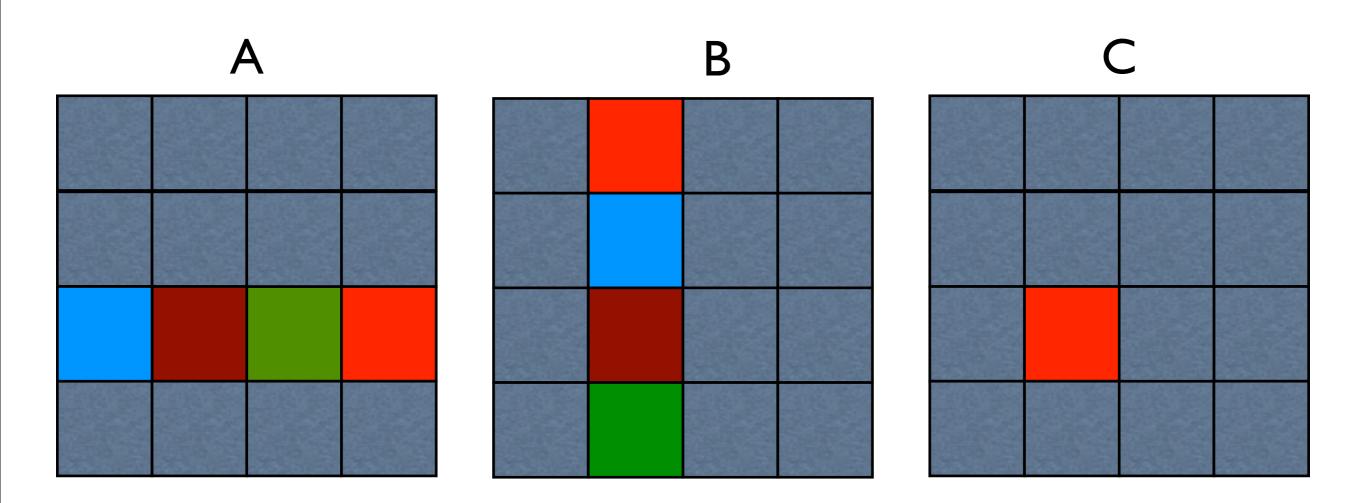
#### Change the initial data setup



#### Change the initial data setup

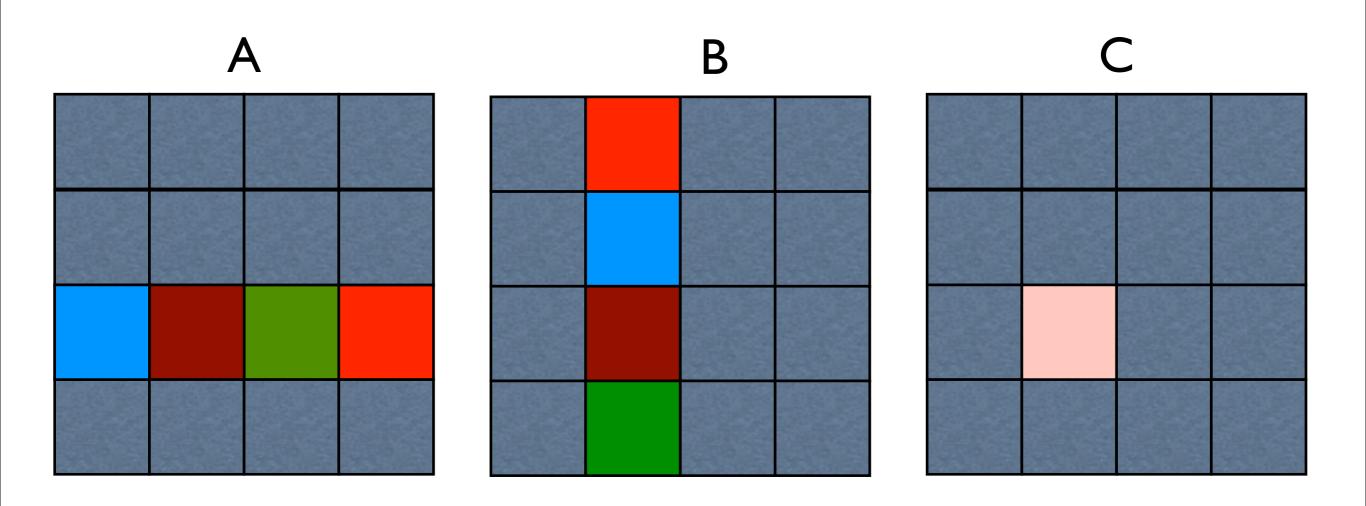


### Every processor now has useful work to do

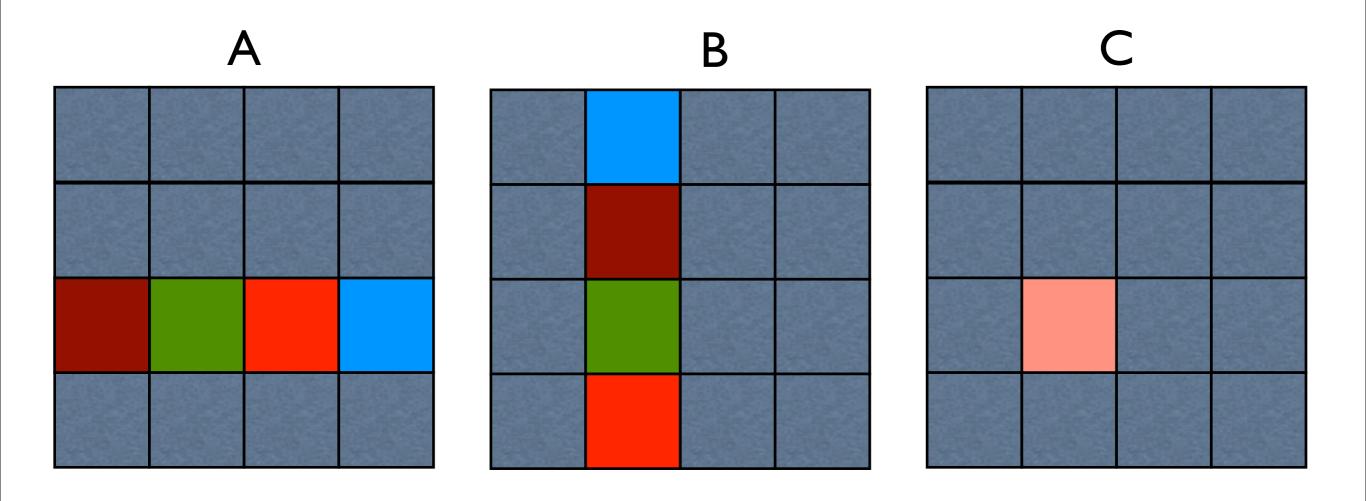


Note -- this only shows the full data layout for one processor

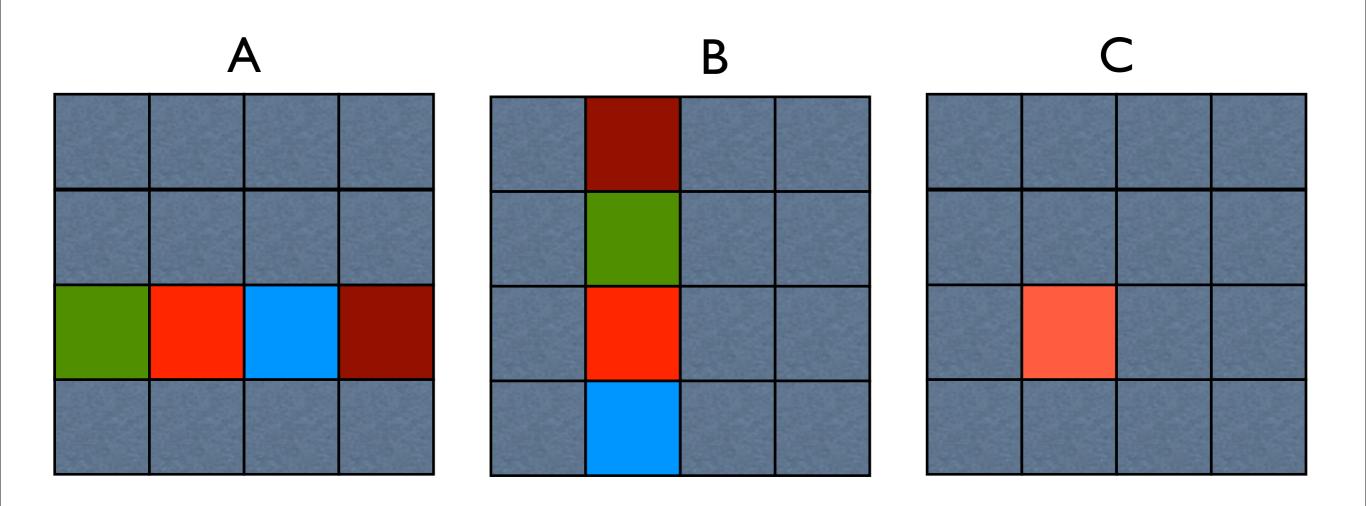
First partial sum



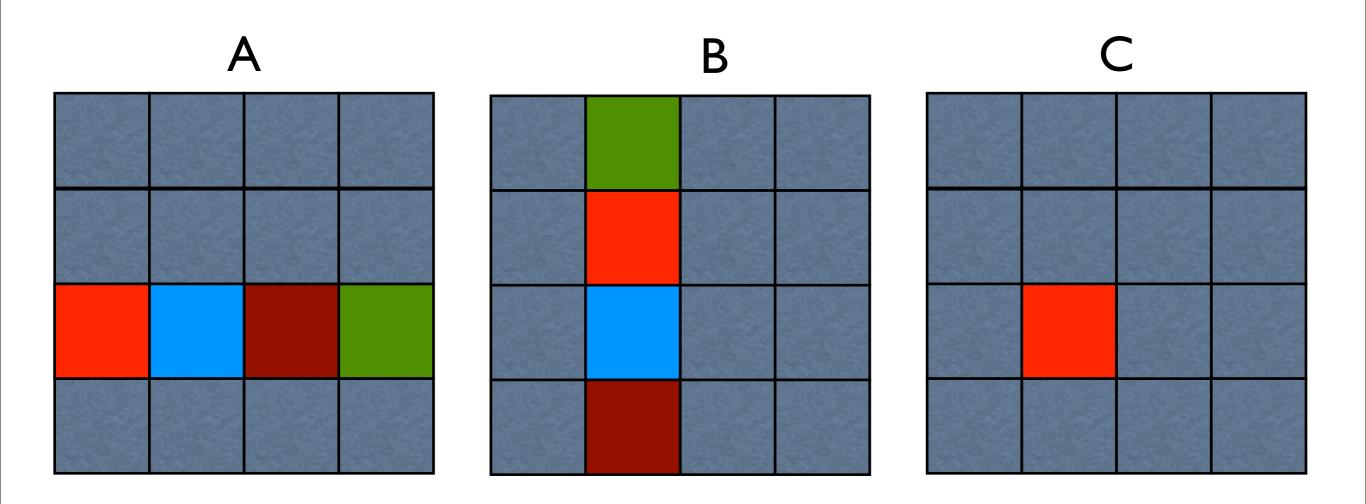
Second partial sum



Third partial sum



Fourth partial sum



### Another way to view this

$$\begin{bmatrix} A_{0,0} \\ B_{0,0} \end{bmatrix} \begin{bmatrix} A_{0,1} \\ B_{0,1} \end{bmatrix} \begin{bmatrix} A_{0,2} \\ B_{0,2} \end{bmatrix} \begin{bmatrix} A_{0,3} \\ B_{0,3} \end{bmatrix} \qquad \begin{bmatrix} A_{0,0} \\ B_{0,0} \end{bmatrix} \begin{bmatrix} A_{0,1} \\ B_{1,1} \end{bmatrix} \begin{bmatrix} A_{0,2} \\ B_{2,2} \end{bmatrix} \begin{bmatrix} A_{0,3} \\ B_{3,3} \end{bmatrix}$$

$$\begin{bmatrix} A_{1,0} \\ B_{1,0} \end{bmatrix} \begin{bmatrix} A_{1,1} \\ B_{1,1} \end{bmatrix} \begin{bmatrix} A_{1,2} \\ B_{1,2} \end{bmatrix} \begin{bmatrix} A_{1,3} \\ B_{1,3} \end{bmatrix} \qquad \begin{bmatrix} A_{1,1} \\ B_{1,0} \end{bmatrix} \begin{bmatrix} A_{1,2} \\ B_{2,1} \end{bmatrix} \begin{bmatrix} A_{1,3} \\ B_{3,2} \end{bmatrix} \begin{bmatrix} A_{1,0} \\ B_{0,3} \end{bmatrix}$$

$$\begin{bmatrix} A_{2,0} \\ B_{2,0} \end{bmatrix} \begin{bmatrix} A_{2,1} \\ B_{2,1} \end{bmatrix} \begin{bmatrix} A_{2,2} \\ B_{2,2} \end{bmatrix} \begin{bmatrix} A_{2,3} \\ B_{2,3} \end{bmatrix} \qquad \begin{bmatrix} A_{2,2} \\ B_{2,0} \end{bmatrix} \begin{bmatrix} A_{2,1} \\ B_{1,3} \end{bmatrix} \begin{bmatrix} A_{2,0} \\ B_{2,1} \end{bmatrix} \begin{bmatrix} A_{2,1} \\ B_{1,2} \end{bmatrix} \begin{bmatrix} A_{2,1} \\ B_{1,3} \end{bmatrix}$$

$$\begin{bmatrix} A_{3,0} \\ B_{3,0} \end{bmatrix} \begin{bmatrix} A_{3,1} \\ B_{3,1} \end{bmatrix} \begin{bmatrix} A_{3,2} \\ B_{3,2} \end{bmatrix} \begin{bmatrix} A_{3,3} \\ B_{3,3} \end{bmatrix} \begin{bmatrix} A_{3,0} \\ B_{0,1} \end{bmatrix} \begin{bmatrix} A_{3,1} \\ B_{1,2} \end{bmatrix} \begin{bmatrix} A_{3,2} \\ B_{2,3} \end{bmatrix}$$

**Before** 

**After** 

### Another to view this

B block goes here (up 1 (j) rows) 0,0 0,2 0,1 0,3 0,0 1,3 1,0 1,3 1,1 1,0  $A_{2,0}$ A 2,1 A<sub>2,1</sub>  $A_{2,0}$ A<sub>2,2</sub>  $A_{3,3}$  $A_{3,2}$  $A_{3,3}$ A<sub>3,1</sub>  $A_{3,0}$ 3,0 3,1 В 3,1 3,2 2,3

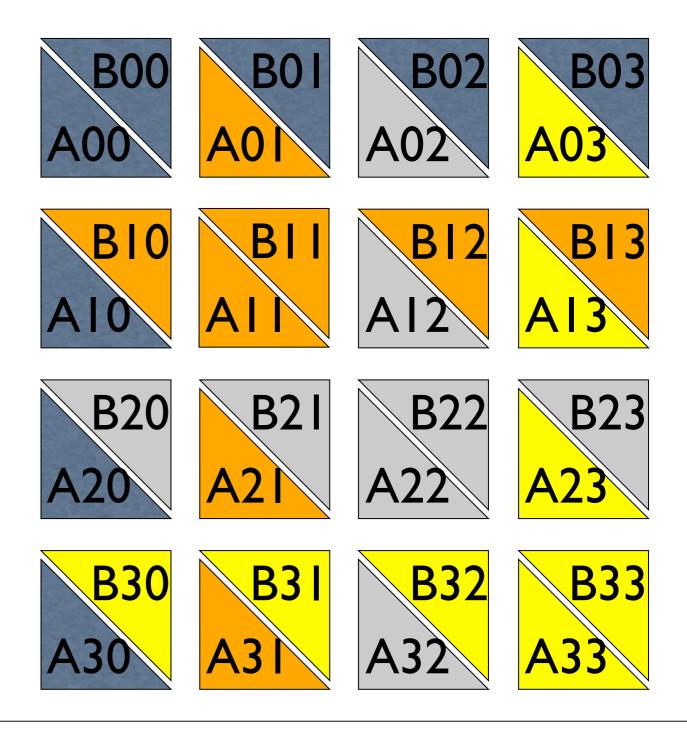
Before

A block goes here (over 2 (i) rows)

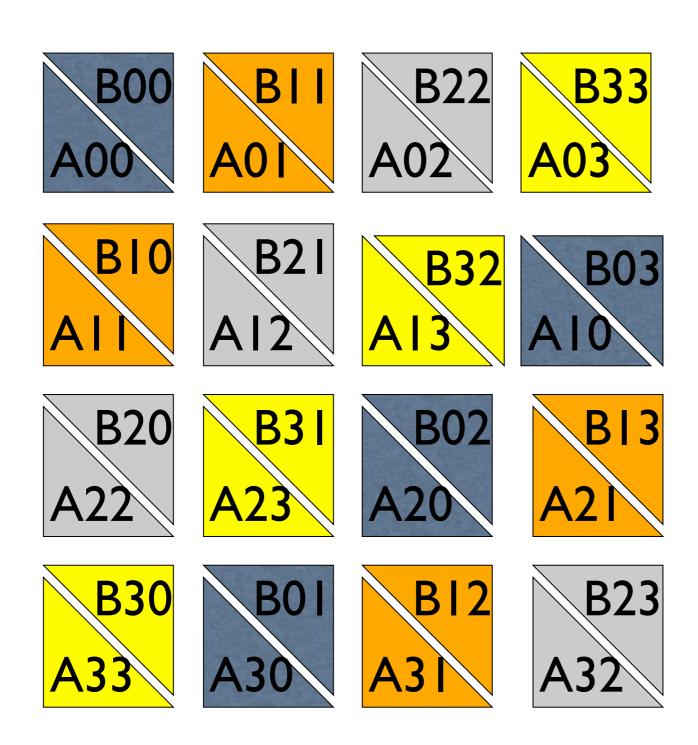
## Yet another way to view this

Each triangle represents a matrix block

Only same-color triangles should be multiplied

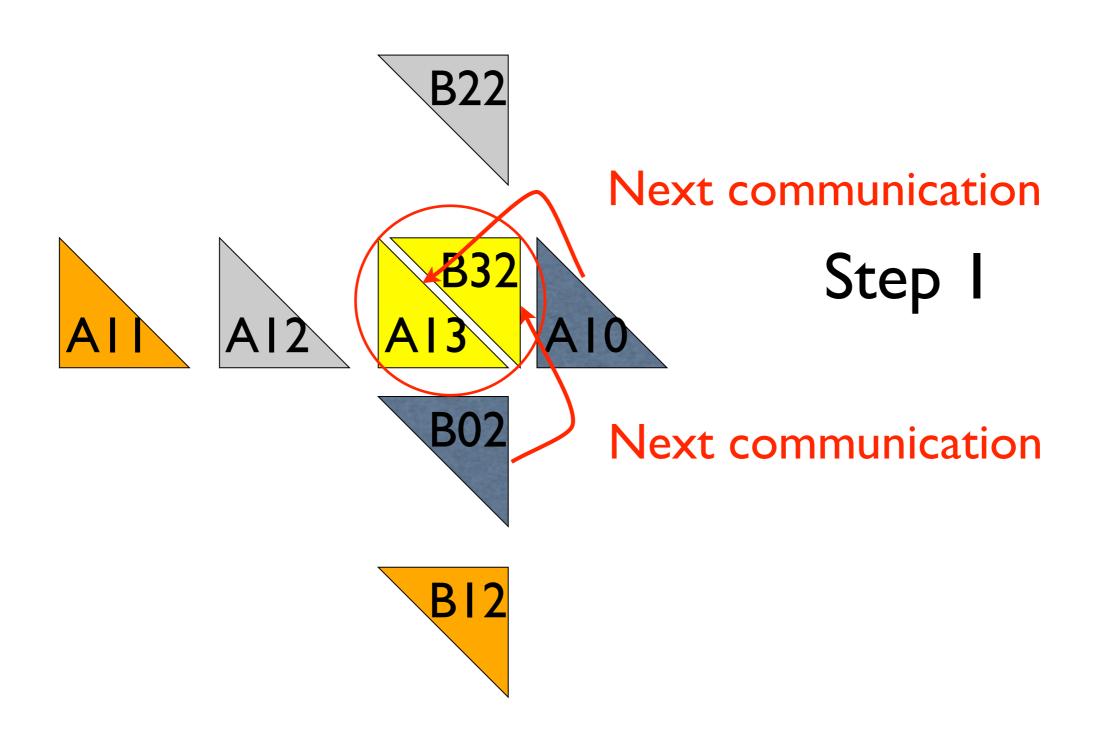


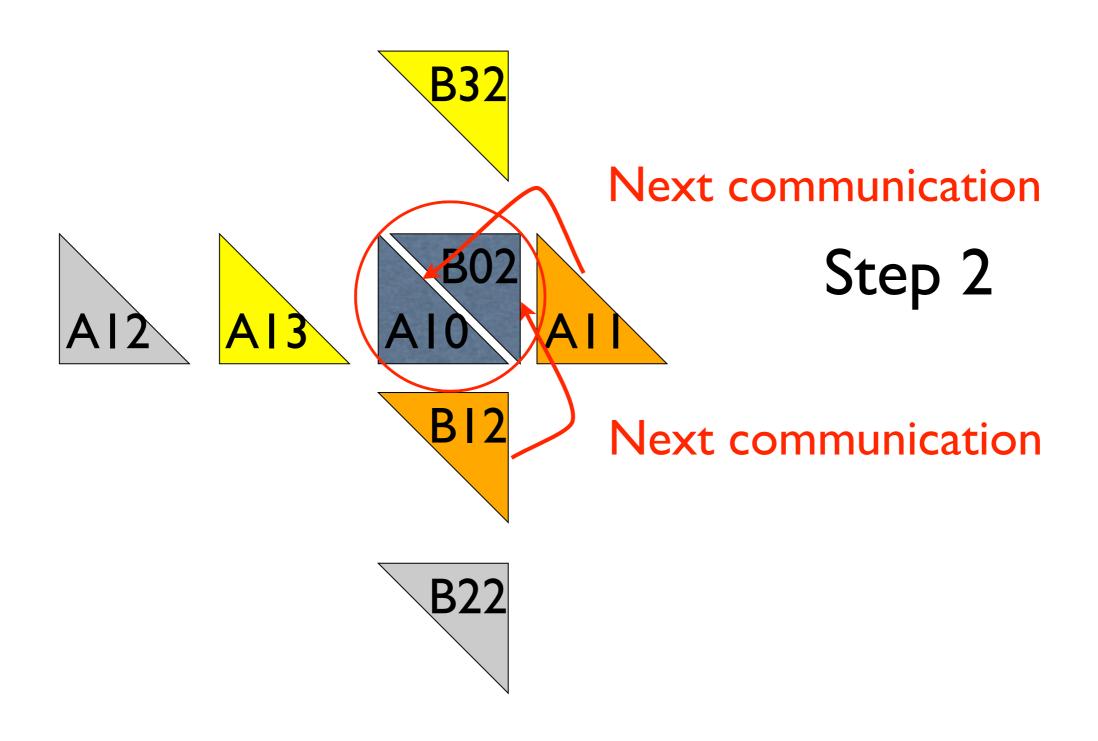
### Rearrange Blocks

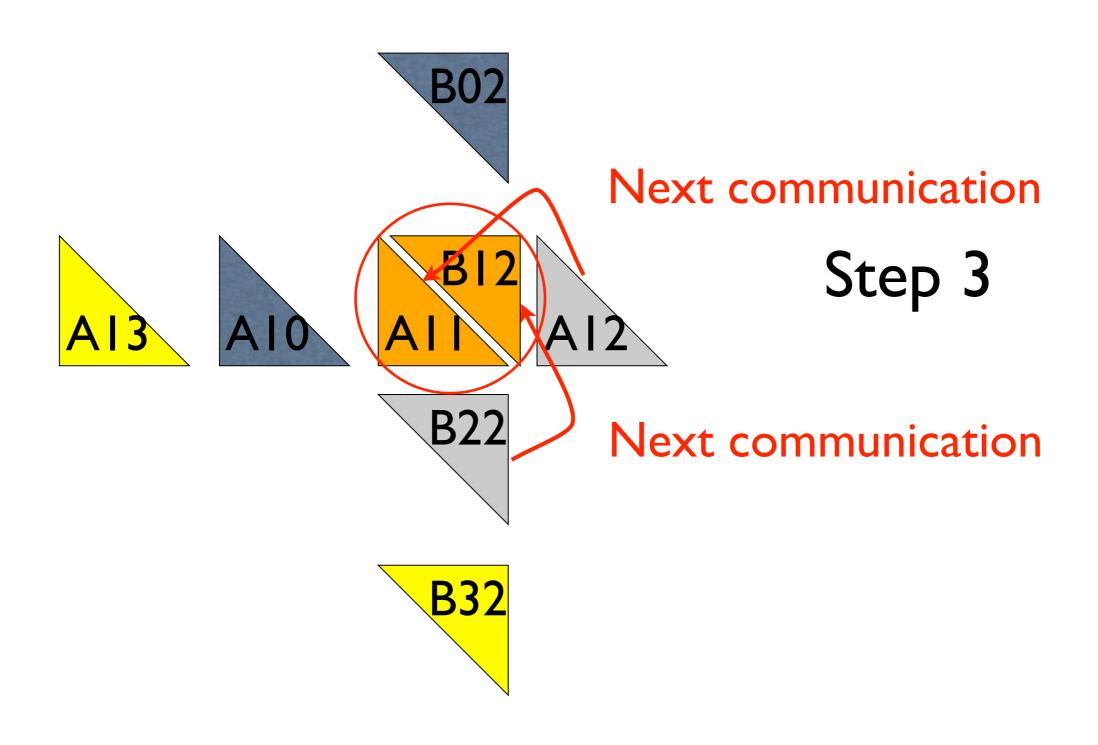


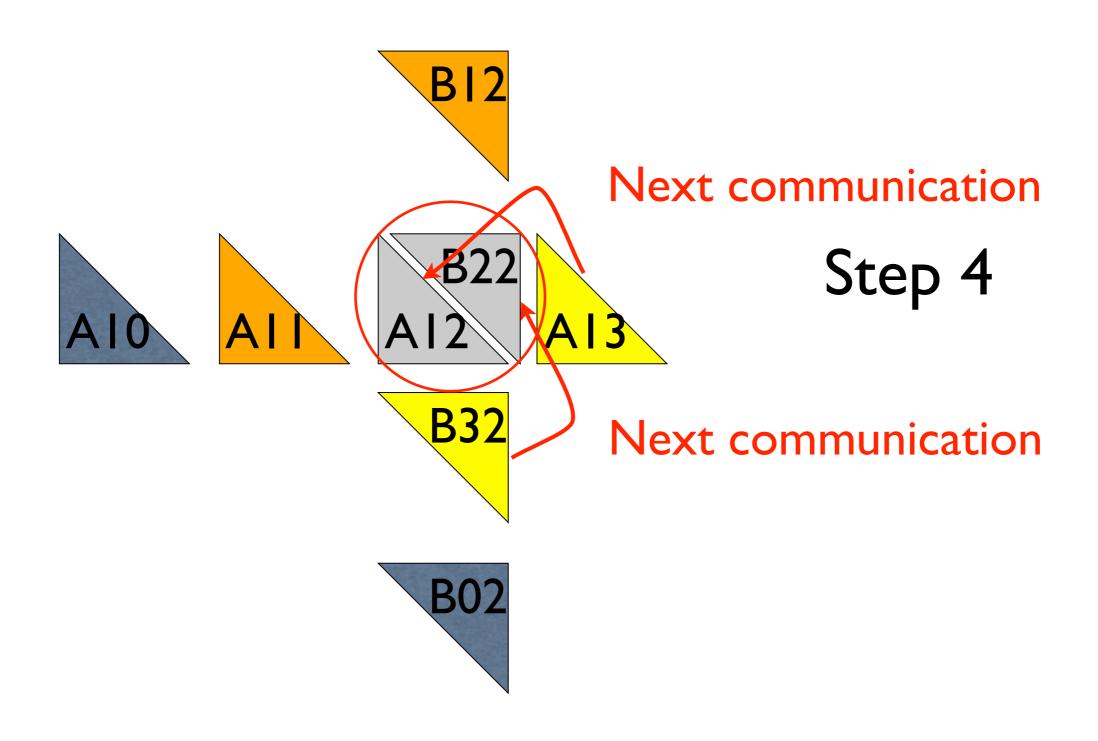
Block Ai,j shifts left i positions

Block Bi,j shifts up j positions









### Complexity Analysis

- Algorithm has  $\sqrt{p}$  iterations
  - During each iteration process multiplies two  $(n / \sqrt{p}) \times (n / \sqrt{p})$  matrices:  $\Theta(n / \sqrt{p})^3$  or  $\Theta(n^3 / p^{3/2})$
- Overall computational complexity:  $\sqrt{p} n^3/p^{3/2}$  or  $\Theta(n^3/p)$ 
  - During each  $\sqrt{p}$  iteration process sends and receives two blocks of size  $(n \mid \sqrt{p}) \times (n \mid \sqrt{p})$
- Overall communication complexity:  $\Theta(n^2/\sqrt{p})$