

Appendix H

Supporting code for μ ML

H.1 Printing types and constraints

This code prints types without too many redundant parentheses. It uses infix notation for the function and tuple constructors.

```
693  <printing types 693>≡ (295c)
      <definition of separate 218d>
      local
        (* precedences *)
        val CONp   = 3
        val STARp  = 2
        val ARROWp = 1
        val NONEp  = 0

        fun parens s = "(" ^ s ^ ")"
        fun bracket (s, context, prec) = if prec <= context then parens s else s
        fun p (context, CONAPP (TYCON "tuple", l)) = bracket (ptuple l, context, STARp)
          | p (context, CONAPP (TYCON "function", [arg, ret])) =
              bracket (p (ARROWp, arg) ^ " -> " ^ p (ARROWp, ret), context, ARROWp)
          | p (context, CONAPP (n, [])) = p (context, n)
          | p (context, CONAPP (n, [t])) = p (CONp, t) ^ " " ^ p (CONp, n)
          | p (context, CONAPP (n, l)) =
              "(" ^ separate ("", " ", ") (map typeString l) ^ ")" ^ p (CONp, n)
          | p (context, TYCON n) = n
          | p (context, TYVAR v) = "'" ^ v
        and ptuple l = separate ("unit", " * ") (map (fn t => p (STARp, t)) l)
        and typeString ty = p (NONEp, ty)
      in
        val typeString = typeString
        val ptuple = ptuple
      end
```

A constraint can be printed in full, but it's easier to read if its first passed to `untriviate`, which removes as many TRIVIAL sub-constraints as possible.

```
694a  <printing constraints 694a>≡ (320c)
      fun untriviate (c /\ c') = (case (untriviate c, untriviate c')
                                   of (TRIVIAL, c) => c
                                    | (c, TRIVIAL) => c
                                    | (c, c') => c /\ c')
      | untriviate atomic = atomic

      fun constraintString (c /\ c') = constraintString c ^ " /\ " ^ constraintString c'
      | constraintString (t == t') = typeString t ^ " == " ^ typeString t'
      | constraintString TRIVIAL = "TRIVIAL"
```

H.2 Parsing

μ ML can use μ Scheme's lexical analysis, so all we have here is a parser. As with μ Scheme, we begin with error-detection functions.

```
694b  <parsing 694b>≡ (331c) 695a>
      fun letDups LETSTAR (loc, bindings) = OK bindings
      | letDups kind (loc, bindings) =
          let val names = map (fn (n, _) => n) bindings
              val kindName = case kind of LET => "let" | LETREC => "letrec" | _ => "??"
              in nodups ("bound name", kindName) (loc, names) >>=+ (fn _ => bindings)
              end
      fun noExp kind _ _ =
          ERROR ("uML does not include '" ^ kind ^ "' expressions")
```

```
== 320b
>>=+ 652a
ERROR 651a
LET 287a
LETREC 287a
LETSTAR 287a
nodups 666a
OK 651a
TRIVIAL 320b
typeString 693
```

695a <parsing 694b>+≡

(331c) <694b 695b>

```

val name      = (fn (NAME n) => SOME n | _ => NONE) <$>? token
val booltok   = (fn (SHARP b) => SOME b | _ => NONE) <$>? token
val int       = (fn (INT n) => SOME n | _ => NONE) <$>? token
val quote     = (fn (QUOTE) => SOME () | _ => NONE) <$>? token

fun exp tokens = (
  VAR          <$> name
  <|> (LITERAL o NUM) <$> int
  <|> (LITERAL o BOOL) <$> booltok
  <|> LITERAL   <$> (quote *> sexp)
  <|> bracket "if"      "(if e1 e2 e3)"      (curry3 IFX <$> exp <*> exp <*> exp)
  <|> bracket "while"   "(while e1 e2)"      (noExp "while" <$> exp <*>! exp)
  <|> bracket "set"     "(set x e)"          (noExp "set" <$> name <*>! exp)
  <|> bracket "begin"   ""                  ( BEGIN <$> many exp)
  <|> bracket "lambda"  "(lambda (names) body)" ( lambda <$> @@ formals <*>!
  <|> bracket "let"      "(let (bindings) body)" (letx LET <$> @@ bindings <*>!
  <|> bracket "letrec"   "(letrec (bindings) body)" (letx LETREC <$> @@ bindings <*>!
  <|> bracket "let*"     "(let* (bindings) body)" (letx LETSTAR <$> @@ bindings <*>!
  <|> "(" >-- literal ")" <|> "empty application"
  <|> curry APPLY <$> "(" >-- exp <*> many exp --> ")"
) tokens
and lambda xs exp =
  nodups ("formal parameter", "lambda") xs >==+ (fn xs => LAMBDA (xs, exp))
and letx kind bs exp = letDups kind bs >==+ (fn bs => LETX (kind, bs, exp))
and formals ts = "(" >-- many name --> ")" ts
and bindings ts = "(" >-- (many binding --> ")" <?> "(x e)...") ts
and binding ts = "(" >-- (pair <$> name <*> exp --> ")" <?> "(x e) in bindings")

and sexp tokens = (
  SYM <$> (notDot <$>! name)
  <|> NUM          <$> int
  <|> BOOL         <$> booltok
  <|> (fn v => embedList [SYM "quote", v]) <$> (quote *> sexp)
  <|> embedList    <$> "(" >-- many sexp --> ")"
) tokens
and notDot "." = ERROR "this interpreter cannot handle . in quoted S-expressions"
  | notDot s = OK s

```

695b <parsing 694b>+≡

(331c) <695a 696a>

```

fun define f formals body =
  nodups ("formal parameter", "definition of function " ^ f) formals >==+
  (fn xs => DEFINE (f, (xs, body)))

val def =
  bracket "define" "(define f (args) body)" (define <$> name <*> @@ formals <*>!
  <|> bracket "val"   "(val x e)"          (curry VAL <$> name <*> exp)
  <|> bracket "val-rec" "(val-rec x e)"      (curry VALREC <$> name <*> exp)
  <|> bracket "use"    "(use filename)"      (USE <$> name)
  <|> literal ")" <|> "unexpected right parenthesis"
  <|> EXP <$> exp
  <?> "definition"

```

696a $\langle \text{parsing } 694b \rangle + \equiv$ (331c) <695b
 val mlSyntax = (schemeToken, def)

H.3 Initial basis

These parts of the initial basis are identical to what we find in μ Scheme.

```
696b      (additions to the  $\mu$ ML initial basis 696b)≡
          (define caar (lambda (l) (car (car l))))
          (define cadr (lambda (l) (car (cdr l))))
          (define cdar (lambda (l) (cdr (car l))))
          (define length (lambda (l)
            (if (null? l) 0
                (+ 1 (length (cdr l))))))
          (define and (lambda (b c) (if b c b)))
          (define or (lambda (b c) (if b b c)))
          (define not (lambda (b) (if b #f #t)))
          696c>
```

```

696c      <additions to the  $\mu$ ML initial basis 696b>+≡                                     <696b 696d>
      (define append (xs ys)
        (if (null? xs)
            ys
            (cons (car xs) (append (cdr xs) ys))))
      (define revapp (xs ys)
        (if (null? xs)
            ys
            (revapp (cdr xs) (cons (car xs) ys))))

```

```
696d      <additions to the  $\mu ML$  initial basis 696b>+≡
          (define o (f g) (lambda (x) (f (g x)))))
          (define curry  (f) (lambda (x) (lambda (y) (f x y)))))
          (define uncurry (f) (lambda (x y) ((f x) y)))
          <696c 696e>
```

```

696e      <additions to the  $\mu$ ML initial basis 696b>+≡                                     <696d 696f>
          (define filter (p? l)
            (if (null? l)
                '()
                (if (p? (car l))
                    (cons (car l) (filter p? (cdr l)))
                    (filter p? (cdr l))))))
329b
329b
329b

```

[illegible]

697a	\langle additions to the μ ML initial basis 696b $\rangle + \equiv$ <pre> (define exists? (p? l) (if (null? l) #f (if (p? (car l)) #t (exists? p? (cdr l))))) (define all? (p? l) (if (null? l) #t (if (p? (car l)) (all? p? (cdr l)) #f))) </pre>	\langle 696f 697b \rangle	
697b	\langle additions to the μ ML initial basis 696b $\rangle + \equiv$ <pre> (define foldr (op zero l) (if (null? l) zero (op (car l) (foldr op zero (cdr l))))) (define foldl (op zero l) (if (null? l) zero (foldl op (op (car l) zero) (cdr l)))) </pre>	\langle 697a 697c \rangle	
697c	\langle additions to the μ ML initial basis 696b $\rangle + \equiv$ <pre> (define <= (x y) (not (> x y))) (define >= (x y) (not (< x y))) (define != (x y) (not (= x y))) </pre>	\langle 697b 697d \rangle	
697d	\langle additions to the μ ML initial basis 696b $\rangle + \equiv$ <pre> (define max (x y) (if (> x y) x y)) (define min (x y) (if (< x y) x y)) (define mod (m n) (- m (* n (/ m n)))) (define gcd (m n) (if (= n 0) m (gcd n (mod m n)))) (define lcm (m n) (* m (/ n (gcd m n)))) </pre>	\langle 697c 697e \rangle	
697e	\langle additions to the μ ML initial basis 696b $\rangle + \equiv$ <pre> (define min* (l) (foldr min (car l) (cdr l))) (define max* (l) (foldr max (car l) (cdr l))) (define gcd* (l) (foldr gcd (car l) (cdr l))) (define lcm* (l) (foldr lcm (car l) (cdr l))) </pre>	\langle 697d 697f \rangle	
697f	\langle additions to the μ ML initial basis 696b $\rangle + \equiv$ <pre> (define list1 (x) (cons x '())) (define list2 (x y) (cons x (list1 y))) (define list3 (x y z) (cons x (list2 y z))) (define list4 (x y z a) (cons x (list3 y z a))) (define list5 (x y z a b) (cons x (list4 y z a b))) (define list6 (x y z a b c) (cons x (list5 y z a b c))) (define list7 (x y z a b c d) (cons x (list6 y z a b c d))) (define list8 (x y z a b c d e) (cons x (list7 y z a b c d e))) </pre>	\langle 697e 698 \rangle	car \mathcal{P} 329b cdr \mathcal{P} 329b cons \mathcal{P} 329b not 696b null? \mathcal{P} 329a

```

698  <additions to the  $\mu$ ML initial basis 696b>+≡
      (define takewhile (p? l)
        (if (null? l)
            '()
            (if (p? (car l))
                (cons (car l) (takewhile p? (cdr l)))
                '()))))
      (define dropwhile (p? l)
        (if (null? l)
            '()
            (if (p? (car l))
                (dropwhile p? (cdr l))
                l)))

```

```

car      P 329b
cdr      P 329b
cons     P 329b
null?    P 329a

```