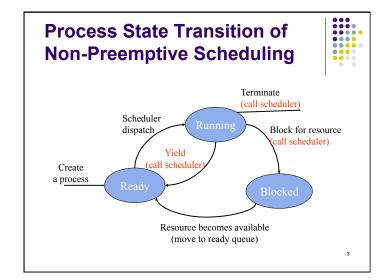
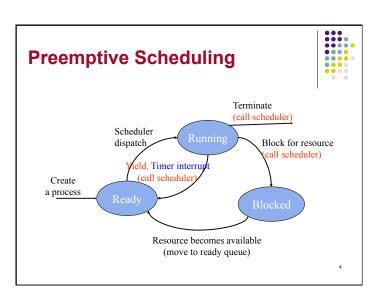


### **Timesharing Systems**



- Timesharing systems support interactive use
  - each user feels he/she has the entire machine
    - CPU, memory, I/O devices
- CPU How?
  - optimize response time
  - based on time-slicing





### **Review on CPU scheduling**



- · Mechanism is easy, policy is hard
  - Jobs have diverse characteristics
  - 4 performance metrics
  - Don't know about future
- Hard to analyze even when narrowing down metric/ job nature
  - FIFO vs. Round Robin in average turnaround time
    - 10 procs, 100 sec each, who wins?
    - 10 procs, 100 sec for first, 10 sec other 9, who wins?
  - · Adding I/O to the diversity

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#### STCF vs. SRTCF



- Shortest time to completion first (shortest job first)
  - Non-preemptive
- Shortest remaining time to completion first
  - Preemptive
- (Under what job arrival scenario are they the same?)
- Advantage
  - Minimal average turnaround time
- Disadvantage
  - Can cause starvation; difficult to know the future

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### **Deep thinking**



- Can we do better than Shortest Job First in terms of average turnaround time?
  - · Assume all jobs arrive at the beginning
- Can we prove it?
  - Proof by contradiction?
  - · Proof by induction?

### Real challenge in using STCF



- Used in long-term scheduling batch system
  - e.g., explicit-length job queues (for repeated jobs)
- What about short-term scheduling (e.g. time-sharing system)?
  - Try prediction for repeated jobs
  - Exponential average:Pred(n+1) = alpha\*t\_n + (1-alpha)\*Pred(n)
- What about starvation?
  - · "Aging"

#### Observations so far

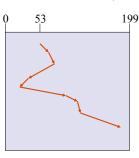


- Need to accommodate interactive jobs
  - Need some kind of RR
- Diversity in jobs job length, I/O mix
  - RR also appears to help
- SJF also has virtue
  - · Reduce avg. turnaround time
- Can we accommodate all?

## [week12] SSTF (Shortest Seek Time First)



- Method
  - · Pick the one closest on disk
- Pros?
- Cons?
- Question
  - Is SSTF optimal?
  - How can we avoid starvation?



98, 183, 37, 122, 14, 124, 65, 67 (65, 67, 37, 14, 98, 122, 124, 183)<sub>2</sub>

### **Scheduling policies**



FIFO Response time

Throughput

RR

Avg. turnaround time

SJF

Fairness

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## **Priority Scheduling**



- To incorporate the spirits of SJF/RR/FIFO
- The method
  - Assign each process a priority
  - Run the process with highest priority in ready queue first
  - Use FIFO for processes with equal priority
  - Adjust priority dynamically
    - To deal with all issues: e.g. aging, I/O wait raises priority
- Advantage
  - Flexibility: Not all processes are "born" equal
- · Challenge?

### **Priority Scheduling (cont)**

- Who sets the priorities
  - Internally by OS
    - I/O to computation ratio (can be dynamic)
    - Memory requirement (can be dynamic)
    - Time constraints (e.g. real-time systems)
  - Externally by users/sysadm
    - Importance
    - Funds paid for
    - e.g. nice
- How -- Dynamically adjustment is tricky

# Approach 1: Multiple Queue Scheduling



- Motivation: processes may be of different nature and can be easily classified
  - e.g. foreground jobs vs. background jobs
- The method:
  - Processes permanently assigned to one queue, based on processes priority / type
  - · Each queue has its own scheduling algo.
    - e.g. RR for foreground queue, FCFS for background queue
  - Need a scheduling among the queues
    - e.g. fixed priority preemptive scheduling (high-pri queue trumps other)
    - e.g. time-slice between queues

.

# Pros/Cons of Multiple Queue Scheduling



- Pros:
  - Low scheduling overhead
    - Jobs do not move across queues
- Cons:
  - Processes permanently assigned to one queue not flexible
    - Program behavior may change
    - E.g. can switch between I/O bound and CPU bound
    - ightarrow Need some learning/adaptation at runtime
  - Starvation cannot be easily handled
    - → Need some learning/adaptation at runtime

### Approach 2:

### **Multilevel Feedback Queue**



Priority	Time slic
0	1
1	2
2	4
3	8

- · Jobs start at high priority queue
- Feedback
  - If timeout expires, drop one level
  - If waiting too long, push up one level (aging)
- Leave I/O bound and interactive processes in higher-priority queue

### **UNIX System V**

- One of the first commercial versions of the Unix
- Originally developed by AT&T
- First released in 1983
- 4 major versions, termed R1 R2, R3, and R4
- SVR4, commercially most successful version, the result of Unix System Unification (with collaborations of major UNIX vendors)
  - IBM's AIX is based on SVR3
  - Sun 's Solaris and HP's HP-UX are based on SVR4
- 80s-early 90s, System V and BSD were the two major "flavors" of UNIX
  - Large multi-user systems vs. Desktop workstations

(SVR3, 4.3 BSD)



Multilevel Feedback Queue with 1 sec preemption

**Traditional Unix Scheduling** 

- "1 sec" preemption
- Preempt if a process doesn't block or complete within 1 sec
- Priority is recomputed every sec (low # is higher)  $P_i$  = base +  $CPU_i$  / 2 + nice, where  $CPU_i$  =  $(U_i + CPU_{i-1})$  / 2

base is the base priority of the process U<sub>i</sub> is process utilization in interval i, nice in [-20,20]

- Base priorities
  - Virtual memory Swapper
  - Block I/O device control
  - Character I/O device control
  - User processes

## **Scheduling Algorithms in OSes**



Operating System	Preemption	Algorithm	
Windows 3.1x	None	Cooperative Scheduler	
Windows 95, 98, Me	Half	Preemptive for 32-bit processes, <u>Cooperative Scheduler</u> for 16-bit processes	
Windows NT (2000, XP, Vista, 7, and Server)	Yes	Multilevel feedback queue	
Mac OS pre-9	None	Cooperative Scheduler	
Mac OS 9	Some	Preemptive for MP tasks, Cooperative Scheduler for processes and threads	
Mac OS X	Yes	Multilevel feedback queue	
Linux pre-2.6	Yes	Multilevel feedback queue	
Linux 2.6-2.6.23	Yes	O(1) scheduler	
Linux post-2.6.23	Yes	Completely Fair Scheduler	
Solaris	Yes	Multilevel feedback queue	
NetBSD	Yes	Multilevel feedback queue 22	
FreeBSD	Yes	Multilevel feedback queue	

### **Cooperative Scheduling**



- Early multitasking systems used applications that voluntarily ceded time to one another.
- This approach, which was eventually supported by many OSes, is known today as cooperative multitasking.
- Cooperative multitasking was once the scheduling scheme employed by Microsoft Windows (prior to Windows 95 and Windows NT) and Mac OS (prior to Mac OS X) ito enable multiple applications to be run simultaneously
- · Now rarely used in larger systems

## **Example Details in Chapter 5**



- Solaris 2
- Windows XP
- Linux
- All are variations of Multilevel Feedback Queue

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## Mechanism vs. policy: Lottery Scheduling (1994)



- Motivation
  - Priority scheduling is implemented by adjusting priorities
  - Adjusting priority is a bit ad hoc (e.g. what rate?)
- Lottery method
  - Give each job a number of tickets
  - High priority jobs get more tickets
  - Randomly pick a winning ticket
  - To avoid starvation, give each job at least one ticket

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## **Reading assignment**



OSC, Chapter 5