Tries

- R-way tries
- ternary search tries
- string symbol table API

Review: Summary of Performance of Symbol Table Implementations

Frequency of operations.

implementation		typical case	ordered	operations	
	search	insert	delete	operations	on keys
red-black BST	1.00 lg N	1.00 lg N	1.00 lg N	yes	compareTo()
hashing	1 †	1 †	1 †	no	equals() hashcode()

† under uniform hashing assumption

Q. Can we do better?

A. Yes, if we can avoid examining the entire key, as with string sorting.

String symbol table basic API

String symbol table. Symbol table specialized to string keys.

```
public class StringST<Value>

StringST()

create an empty symbol table

void put(String key, Value val)

put key-value pair into the symbol table

Value get(String key)

return value paired with given key

boolean contains(String key)

is there a value paired with the given key?
```

Goal. Faster than hashing, more flexible than binary search trees.

String symbol table implementations cost summary

	character accesses (typical case)				dedup	
implementation	search hit	search miss	insert	space (references)	moby.txt	actors.txt
red-black BST	L + c lg ² N	c lg ² N	c lg ² N	4N	1.40	97.4
hashing	L	L	L	4N to 16N	0.76	40.6

file	size	words	distinct	
moby.txt	1.2 MB	210 K	32 K	
actors.txt	82 MB	11.4 M	900 K	

Parameters

- N = number of strings
- L = length of string
- R = radix

Challenge. Efficient performance for string keys.

Tries

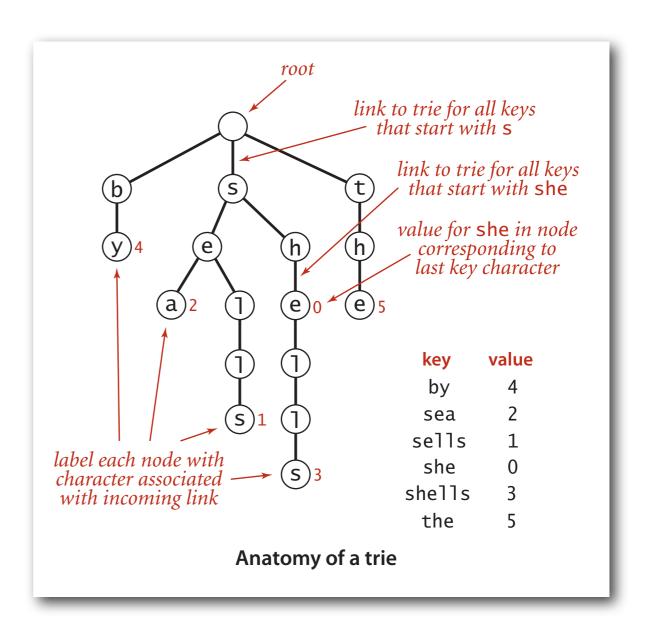
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Tries

Tries. [from retrieval, but pronounced "try"]

- Store characters and values in nodes (not keys).
- \bullet Each node has R children, one for each possible character.
- For now, we do not draw null links.

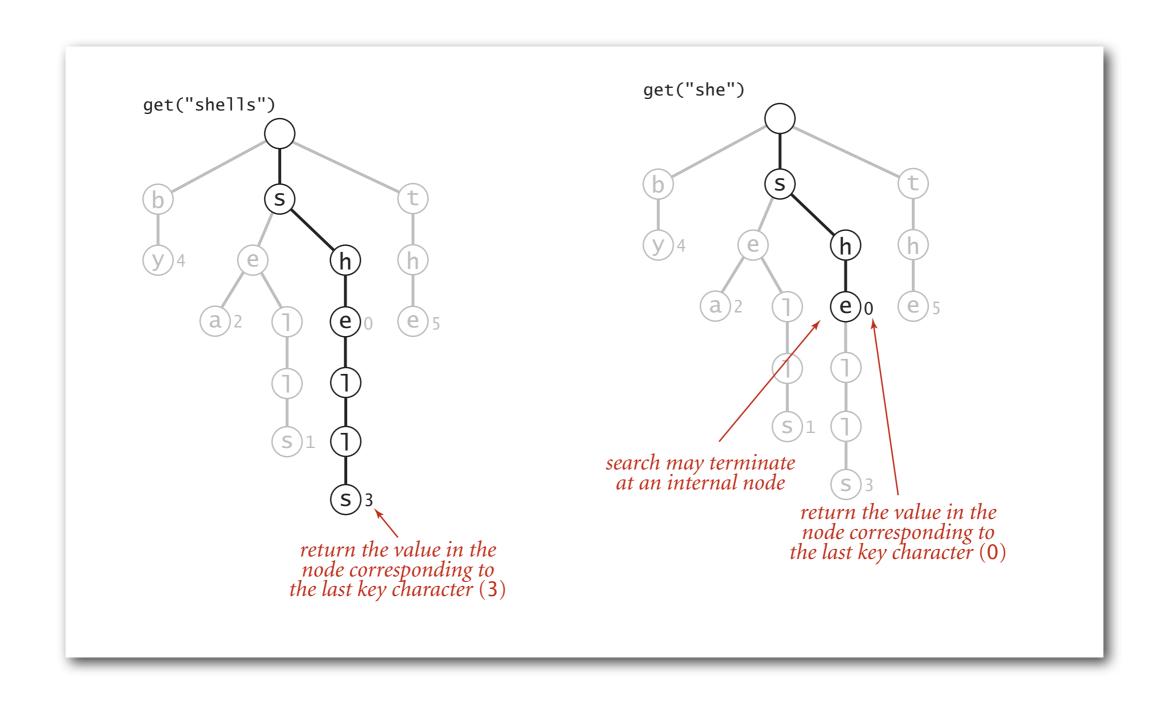
Ex. she sells sea shells by the



Search in a trie

Follow links corresponding to each character in the key.

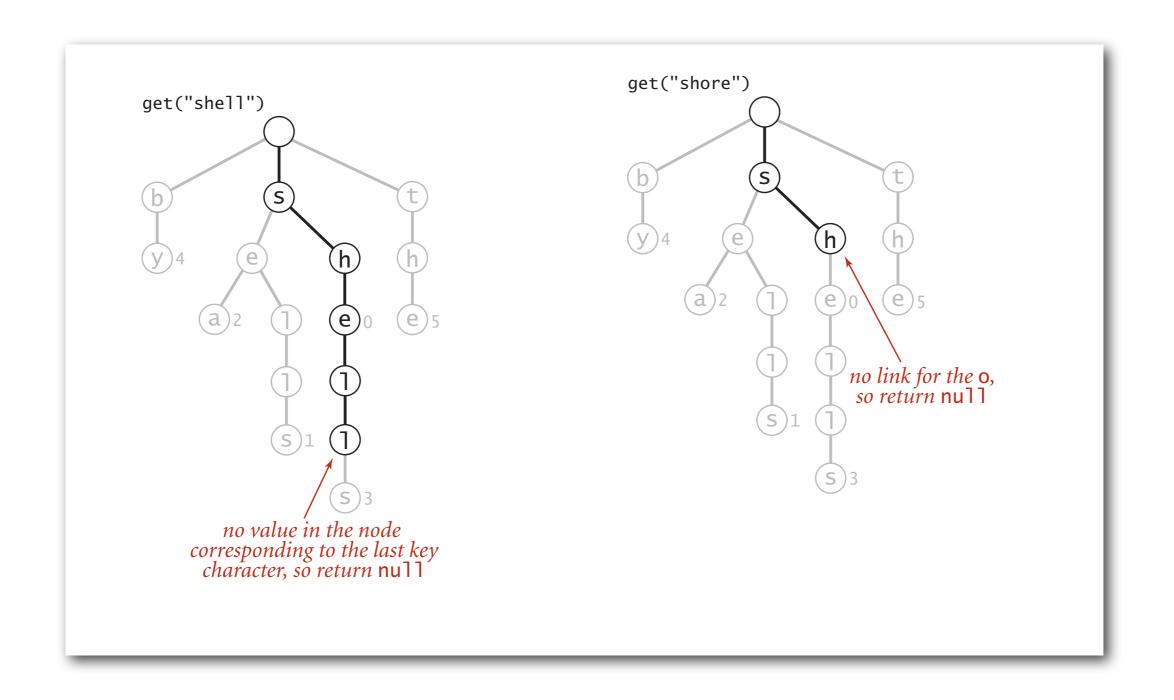
- Search hit: node where search ends has a non-null value.
- Search miss: reach a null link or node where search ends has null value.



Search in a trie

Follow links corresponding to each character in the key.

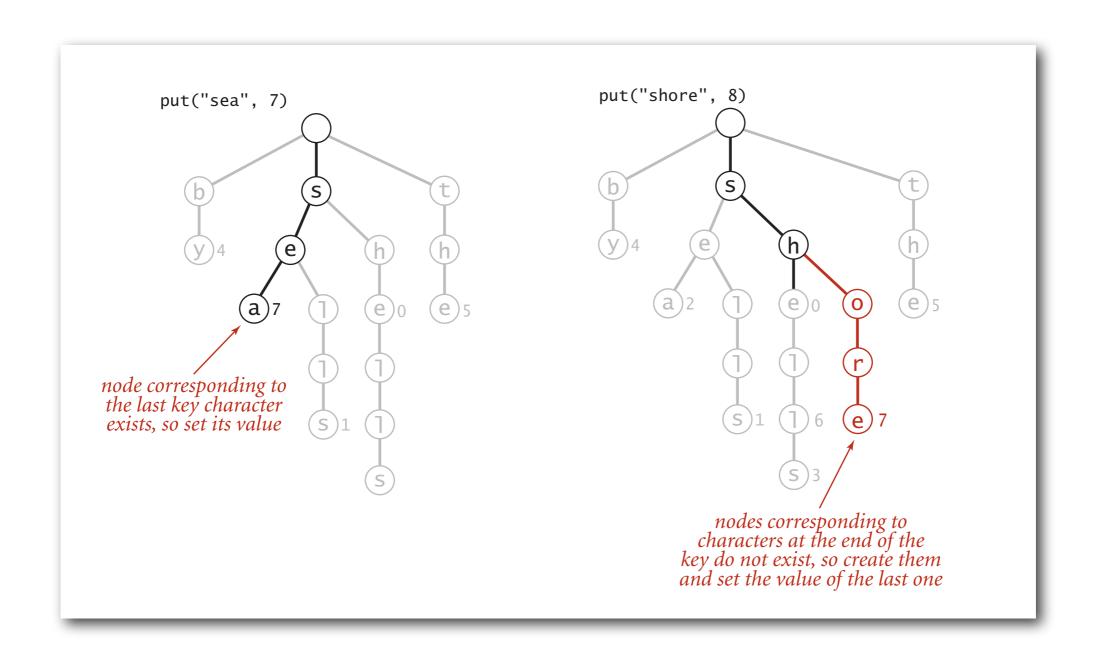
- Search hit: node where search ends has a non-null value.
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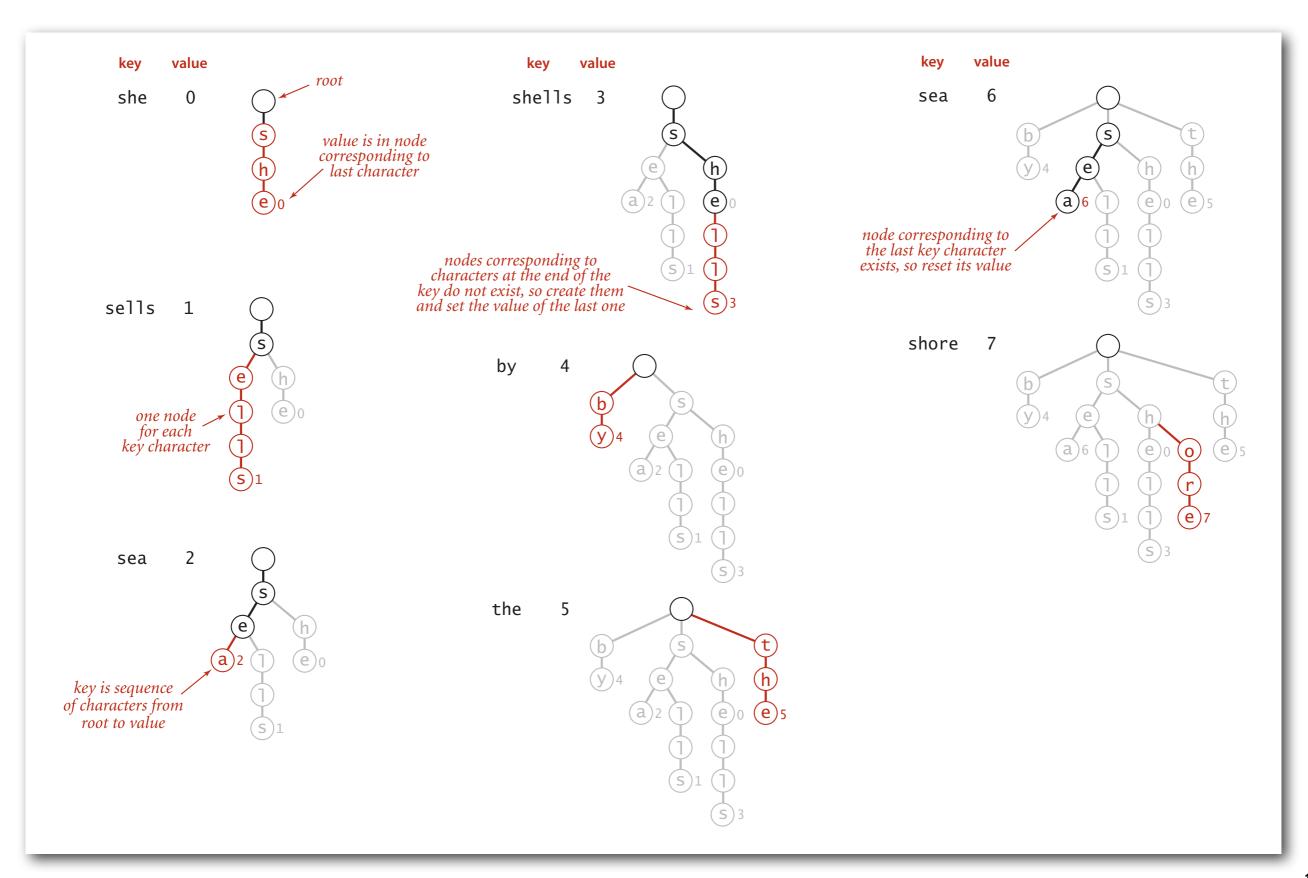
Insertion into a trie

Follow links corresponding to each character in the key.

- Encounter a null link: create new node.
- Encounter the last character of the key: set value in that node.



Trie construction example

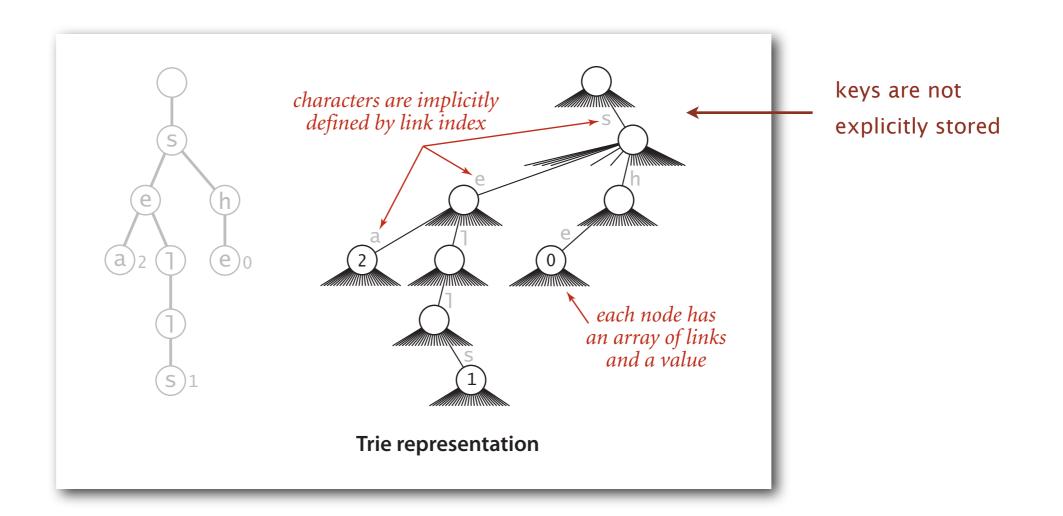


Trie representation: Java implementation

Node. A value, plus references to R nodes.

```
private static class Node
{
    private Object value;
    private Node[] next = new Node[R];
}

use Object instead of value since no generic array creation in Java
private Node[] next = new Node[R];
```



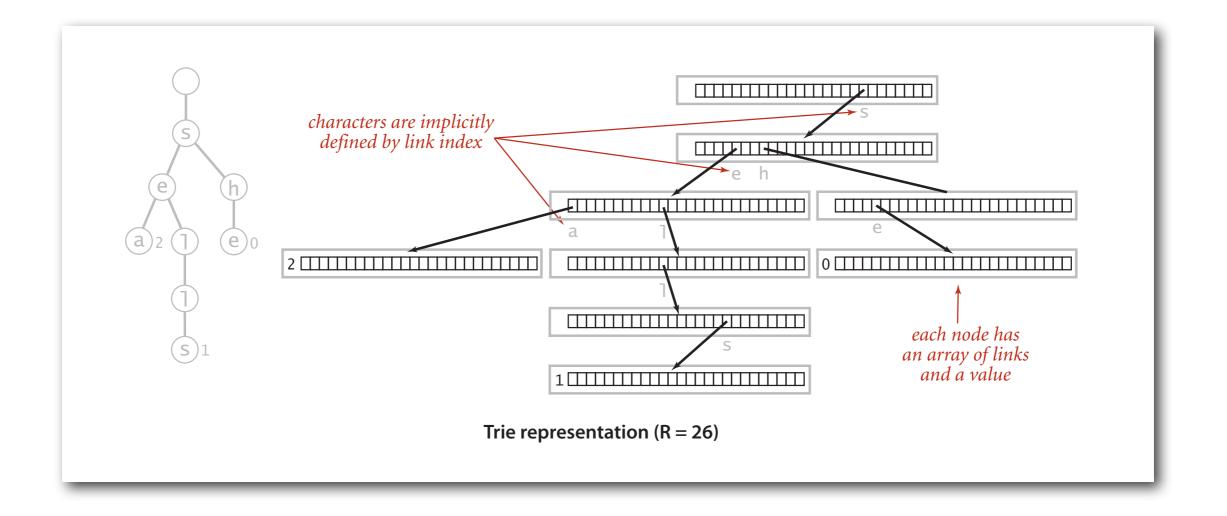
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use Object instead of Value since no generic array creation in Java

private Node[] next = new Node[R];
```



R-way trie: Java implementation

```
public class TrieST<Value>
{
   private static final int R = 256; ← extended ASCII
   private Node root;
   private static class Node
   { /* see previous slide */ }
   public void put(String key, Value val)
   { root = put(root, key, val, 0); }
   private Node put(Node x, String key, Value val, int d)
      if (x == null) x = new Node();
      if (d == key.length()) { x.val = val; return x; }
      char c = key.charAt(d);
      x.next[c] = put(x.next[c], key, val, d+1);
      return x;
```

R-way trie: Java implementation (continued)

```
public boolean contains(String key)
   return get(key) != null;
public Value get(String key)
   Node x = get(root, key, 0);
   if (x == null) return null;
   return (Value) x.val;
                                    cast needed
private Node get(Node x, String key, int d)
   if (x == null) return null;
   if (d == key.length()) return x;
   char c = key.charAt(d);
   return get(x.next[c], key, d+1);
```

Trie performance

Search miss.

- Could have mismatch on first character.
- Typical case: examine only a few characters (sublinear).

Search hit. Need to examine all L characters for equality.

Space. R null links at each leaf.

(but sublinear space possible if many short strings share common prefixes)

Bottom line. Fast search hit and even faster search miss, but wastes space.

Tries

- R-way tries
- ternary search tries
- string symbol table API

Ternary search tries

TST. [Bentley-Sedgewick, 1997]

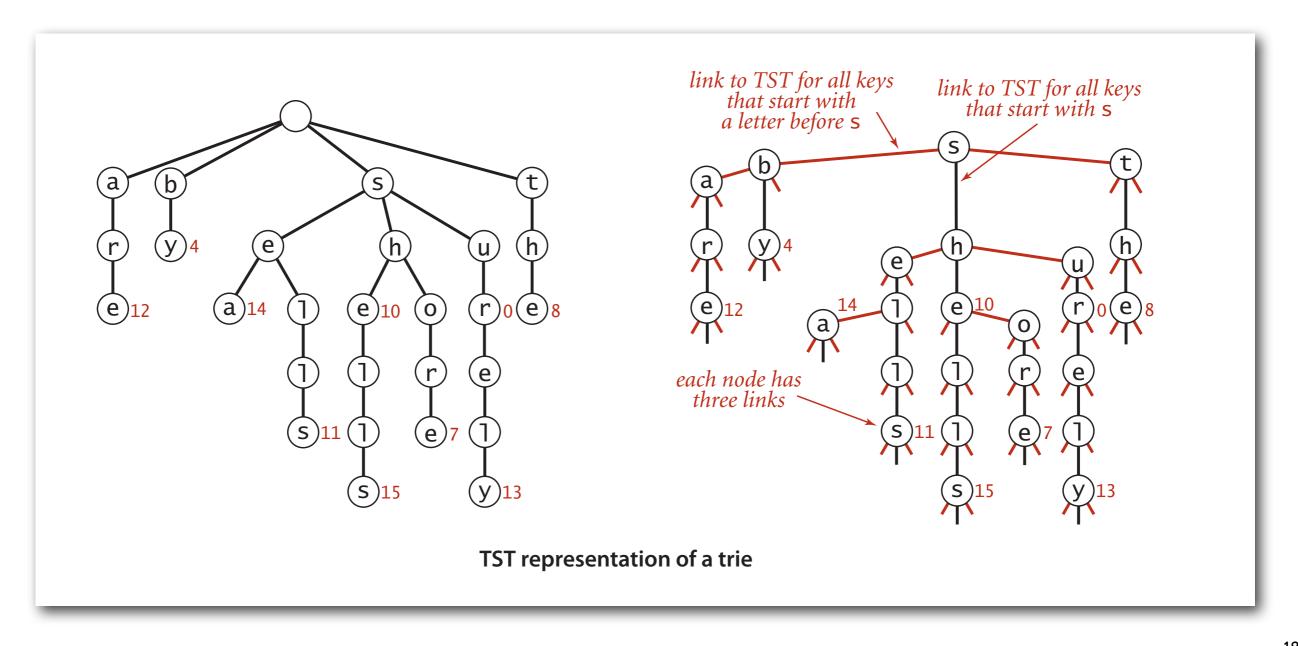
- Store characters and values in nodes (not keys).
- Each node has three children: smaller (left), equal (middle), larger (right).



Ternary search tries

TST. [Bentley-Sedgewick, 1997]

- Store characters and values in nodes (not keys).
- Each node has three children: smaller (left), equal (middle), larger (right).



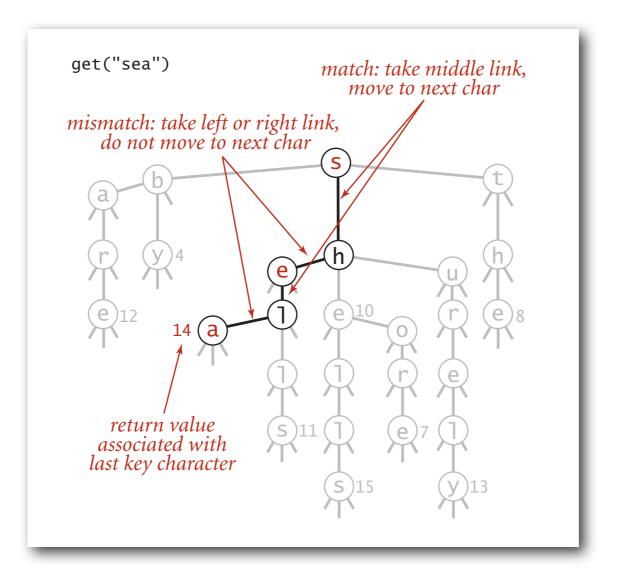
Search in a TST

Follow links corresponding to each character in the key.

- If less, take left link; if greater, take right link.
- If equal, take the middle link and move to the next key character.

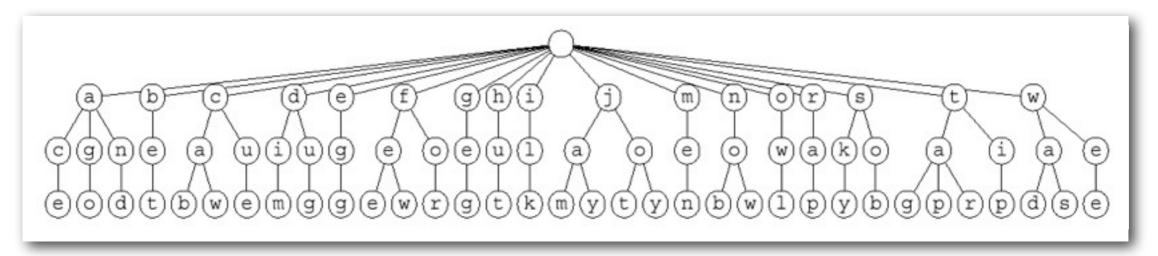
Search hit. Node where search ends has a non-null value.

Search miss. Reach a null link or node where search ends has null value.



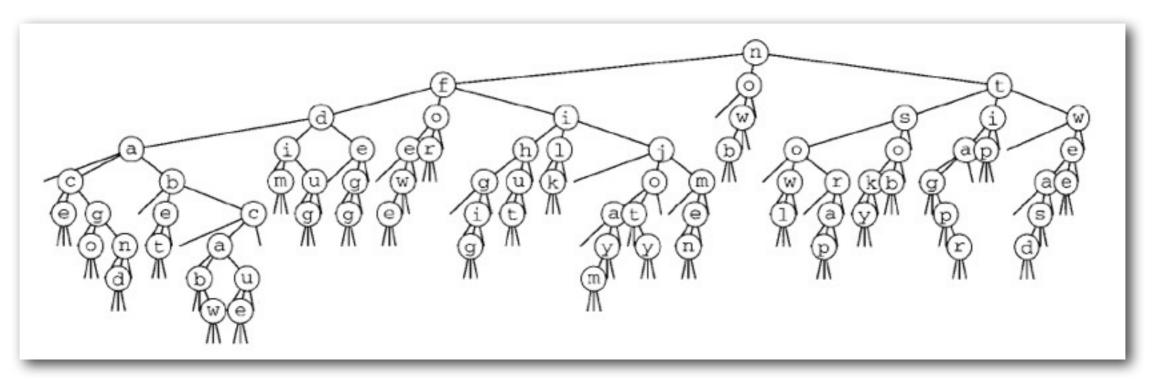
26-way trie vs.TST

26-way trie. 26 null links in each leaf.



26-way trie (1035 null links, not shown)

TST. 3 null links in each leaf.



TST (155 null links)

now for tip ilk dim tag jot sob nob sky hut ace bet men egg few jay owl joy rap gig wee was cab wad caw cue fee tap

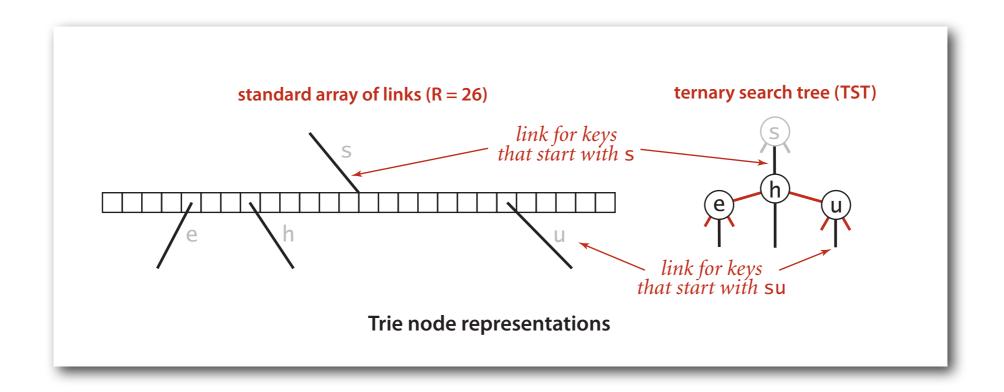
ago

TST representation in Java

ATST node is five fields:

- A value.
- A character c.
- A reference to a left TST.
- A reference to a middle TST.
- A reference to a right TST.

```
private class Node
{
   private Value val;
   private char c;
   private Node left, mid, right;
}
```



TST: Java implementation

```
public class TST<Value>
  private Node root;
  private class Node
   { /* see previous slide */ }
  public void put(String key, Value val)
   { root = put(root, key, val, 0); }
  private Node put(Node x, String key, Value val, int d)
     char c = key.charAt(d);
     if (x == null) \{ x = new Node(); x.c = c; \}
     if (c < x.c) x.left = put(x.left, key, val, d);
     else if (c > x.c)  x.right = put(x.right, key, val, d);
     else if (d < s.length() - 1) x.mid = put(x.mid, key, val, d+1);
     else
                                x.val = val;
     return x;
```

TST: Java implementation (continued)

```
public boolean contains(String key)
  return get(key) != null; }
public Value get(String key)
  Node x = get(root, key, 0);
  if (x == null) return null;
  return x.val;
private Node get(Node x, String key, int d)
  if (x == null) return null;
  char c = key.charAt(d);
  if (c < x.c) return get(x.left, key, d);
  else if (c > x.c) return get(x.right, key, d);
  else if (d < key.length() - 1) return get(x.mid, key, d+1);</pre>
  else
                                return x;
```

String symbol table implementation cost summary

	character accesses (typical case)				dedup	
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red-black BST	L + c lg ² N	c lg ² N	c lg ² N	4 N	1.40	97.4
hashing	L	L	L	4 N to 16 N	0.76	40.6
R-way trie	L	log _R N	L	(R + 1) N	1.12	out of memory
TST	L + In N	In N	L + In N	4 N	0.72	38.7

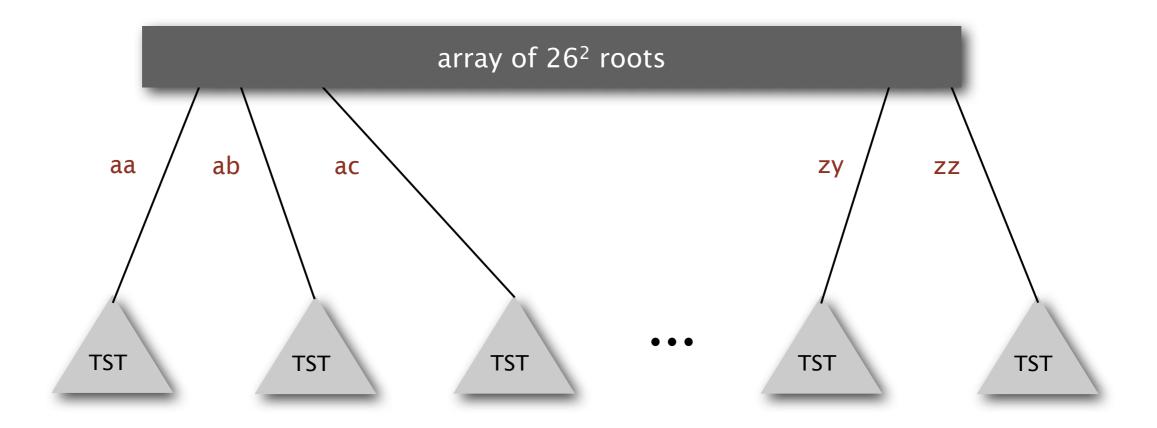
Remark. Can build balanced TSTs via rotations to achieve $L + \log N$ worst-case guarantees.

Bottom line. TST is as fast as hashing (for string keys), space efficient.

TST with R² branching at root

Hybrid of R-way trie and TST.

- Do R^2 -way branching at root.
- Each of R^2 root nodes points to a TST.



Q. What about one- and two-letter words?

String symbol table implementation cost summary

	character accesses (typical case)				dedup		
implementation	search hit	search miss	insert	space (references)	moby.txt	actors.txt	
red-black BST	L + c lg ² N	c lg ² N	c lg ² N	4 N	1.40	97.4	
hashing	L	L	L	4 N to 16 N	0.76	40.6	
R-way trie	L	log _R N	L	(R + 1) N	1.12	out of memory	
TST	L + In N	In N	L + In N	4 N	0.72	38.7	
TST with R ²	L + In N	In N	L + In N	4 N + R ²	0.51	32.7	

TST vs. hashing

Hashing.

- Need to examine entire key.
- Search hits and misses cost about the same.
- Need good hash function for every key type.
- No help for ordered symbol table operations.

TSTs.

- Works only for strings (or digital keys).
- Only examines just enough key characters.
- Search miss may only involve a few characters.
- Can handle ordered symbol table operations (plus others!).

Bottom line. TSTs are:

- Faster than hashing (especially for search misses).
- More flexible than red-black trees (next).

Tries

- R-way tries
- ternary search tries
- string symbol table API

String symbol table API

Character-based operations. The string symbol table API supports several useful character-based operations.

by sea sells she shells shore the

Prefix match. Keys with prefix "sh": "she", "shells", and "shore".

Longest prefix. Key that is the longest prefix of "shellsort": "shells".

Wildcard match. Keys that match ".he": "she" and "the".

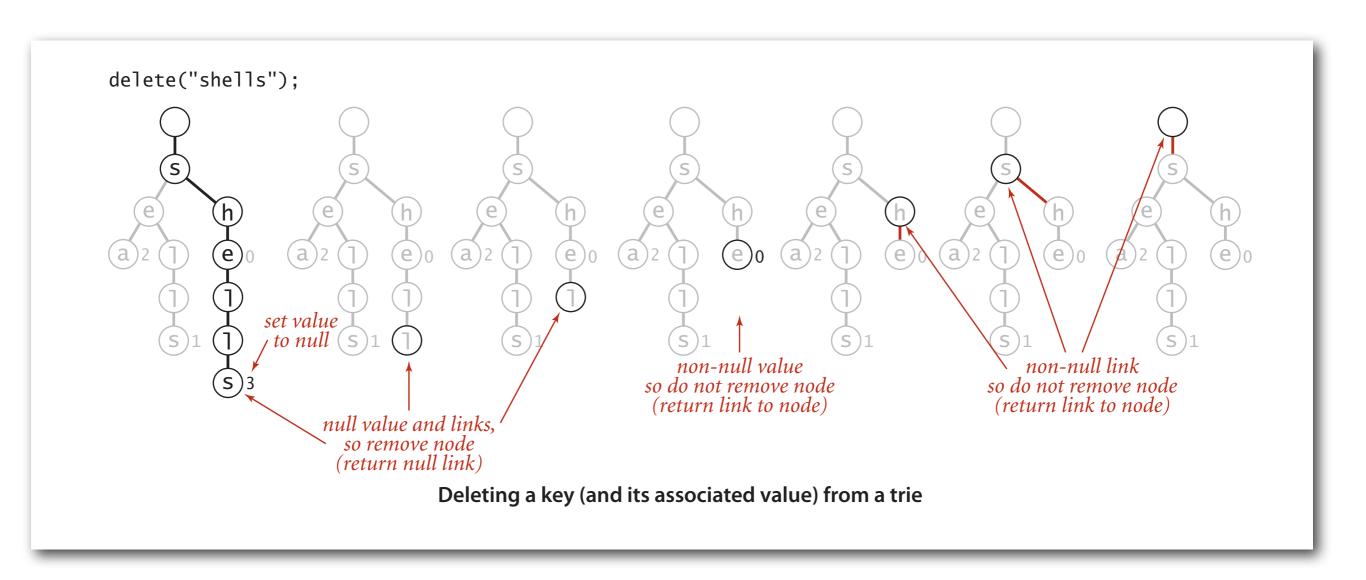
String symbol table API

```
public class StringST<Value>
                                                          create a symbol table with string keys
                       StringST()
                                                          create a symbol table with string keys
                       StringST(Alphabet alpha)
                                                          whose characters are taken from alpha.
                                                          put key-value pair into the symbol table
                void put(String key, Value val)
                                                          (remove key from table if value is null)
                                                          value paired with key
               Value get(String key)
                                                          (null if key is absent)
                void delete(String key)
                                                          remove key (and its value) from table
            boolean contains(String key)
                                                          is there a value paired with key?
            boolean isEmpty()
                                                          is the table empty?
             String longestPrefixOf(String s)
                                                          return the longest key that is a prefix of s
 Iterable<String> keysWithPrefix(String s)
                                                          all the keys having s as a prefix.
                                                          all the keys that match s (where.
 Iterable<String> keysThatMatch(String s)
                                                          matches any character).
                                                          number of key-value pairs in the table
                 int size()
 Iterable<String> keys()
                                                          all the keys in the symbol table
```

Deletion in an R-way trie

To delete a key-value pair:

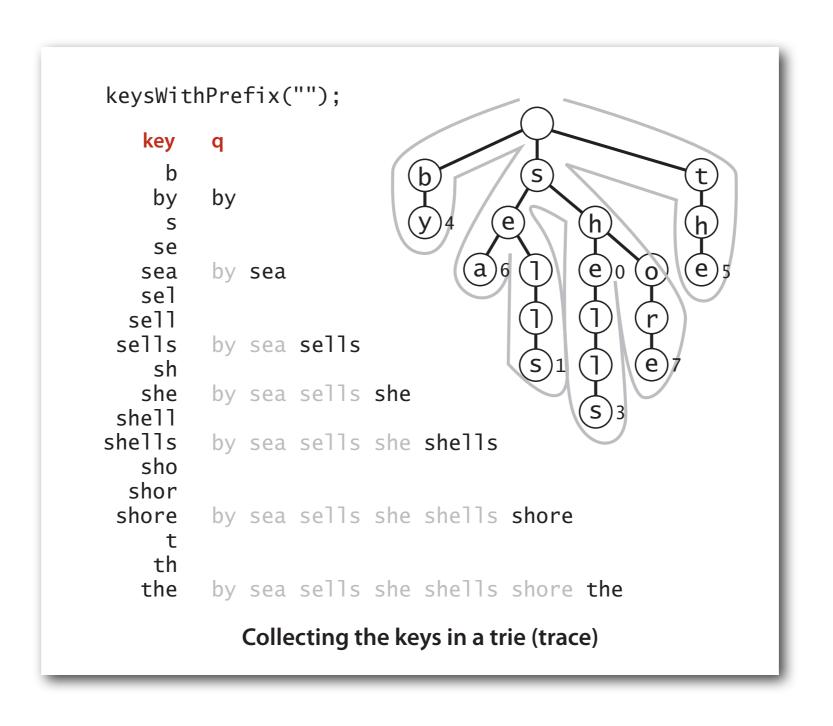
- Find the node corresponding to key and set value to null.
- If that node has all null links, remove that node (and recur).



Ordered iteration

To iterate through all keys in sorted order:

- Do inorder traversal of trie; add keys encountered to a queue.
- Maintain sequence of characters on path from root to node.



Ordered iteration: Java implementation

To iterate through all keys in sorted order:

- Do inorder traversal of trie; add keys encountered to a queue.
- Maintain sequence of characters on path from root to node.

```
public Iterable<String> keys()
   Queue<String> queue = new Queue<String>();
   collect(root, "", queue);
   return queue;
                                                         sequence of characters
                                                         on path from root to x
private void collect(Node x, String prefix, Queue<String> q)
   if (x == null) return;
   if (x.val != null) q.enqueue(prefix);
   for (char c = 0; c < R; c++)
      collect(x.next[c], prefix + c, q);
```

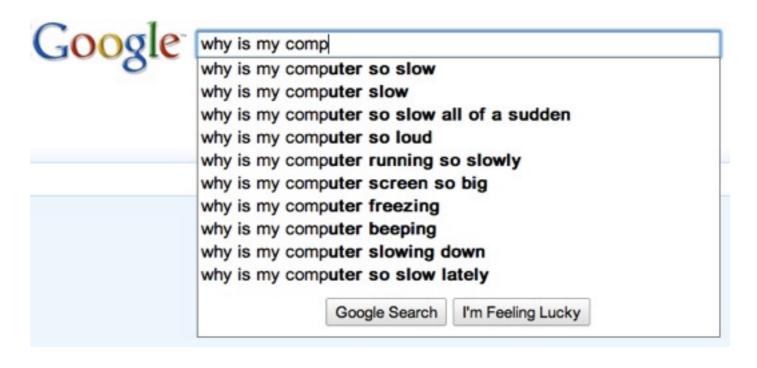
Prefix matches

Find all keys in symbol table starting with a given prefix.

Ex. Autocomplete in a cell phone, search bar, text editor, or shell.

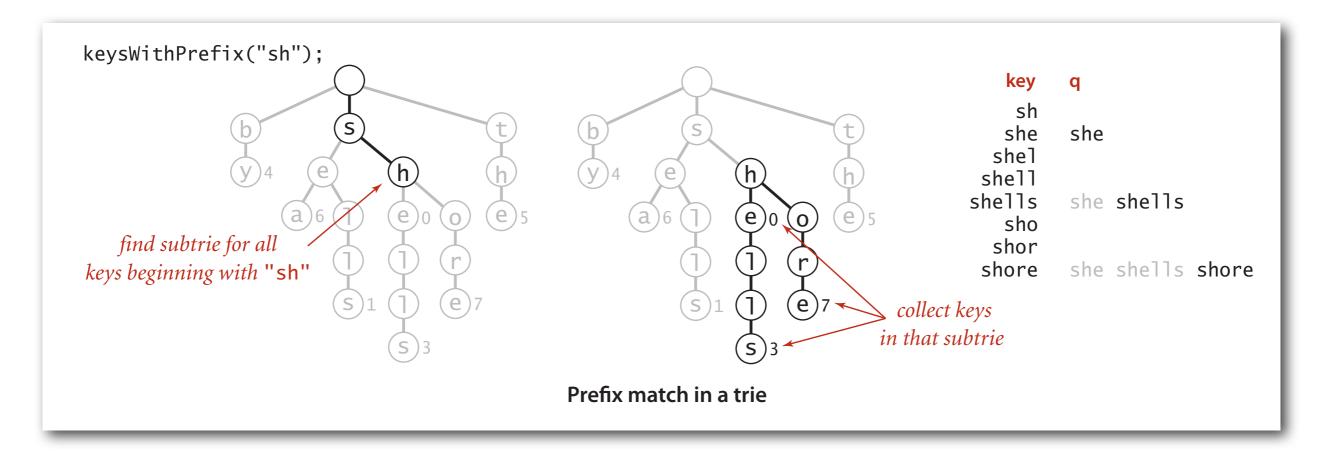
- User types characters one at a time.
- System reports all matching strings.





Prefix matches

Find all keys in symbol table starting with a given prefix.



```
public Iterable<String> keysWithPrefix(String prefix)
{
   Queue<String> queue = new Queue<String>();
   Node x = get(root, prefix, 0);
   collect(x, prefix, queue);
   root of subtrie for all strings
   return queue;
}
```

Wildcard matches

Use wildcard . to match any character in alphabet.

coalizer coberger codifier cofaster cofather cognizer cohelper colander coleader compiler composer computer cowkeper

co...er

acresce acroach acuracy octarch science scranch scratch scrauch screich scrinch scritch scrunch scudick scutock

.c...c.

Wildcard matches

Search as usual if character is not a period; go down all R branches if query character is a period.

```
public Iterable<String> keysThatMatch(String pat)
   Queue<String> queue = new Queue<String>();
   collect(root, "", 0, pat, queue);
   return queue;
private void collect(Node x, String prefix, String pat, Queue<String> q)
   if (x == null) return;
   int d = prefix.length();
   if (d == pat.length() && x.val != null) q.enqueue(prefix);
   if (d == pat.length()) return;
   char next = pat.charAt(d);
   for (char c = 0; c < R; c++)
       if (next == '.' || next == c)
           collect(x.next[c], prefix + c, pat, q);
```

Longest prefix

Find longest key in symbol table that is a prefix of query string.

Ex. Search IP database for longest prefix matching destination IP, and route packets accordingly.

```
"128"
"128.112"
"128.112.055"
"128.112.055.15"
"128.112.136"
"128.112.155.11"
"128.222"
"128.222"
"128.222.136"

prefix("128.112.136.11") = "128.112.136"
prefix("128.166.123.45") = "128"
```

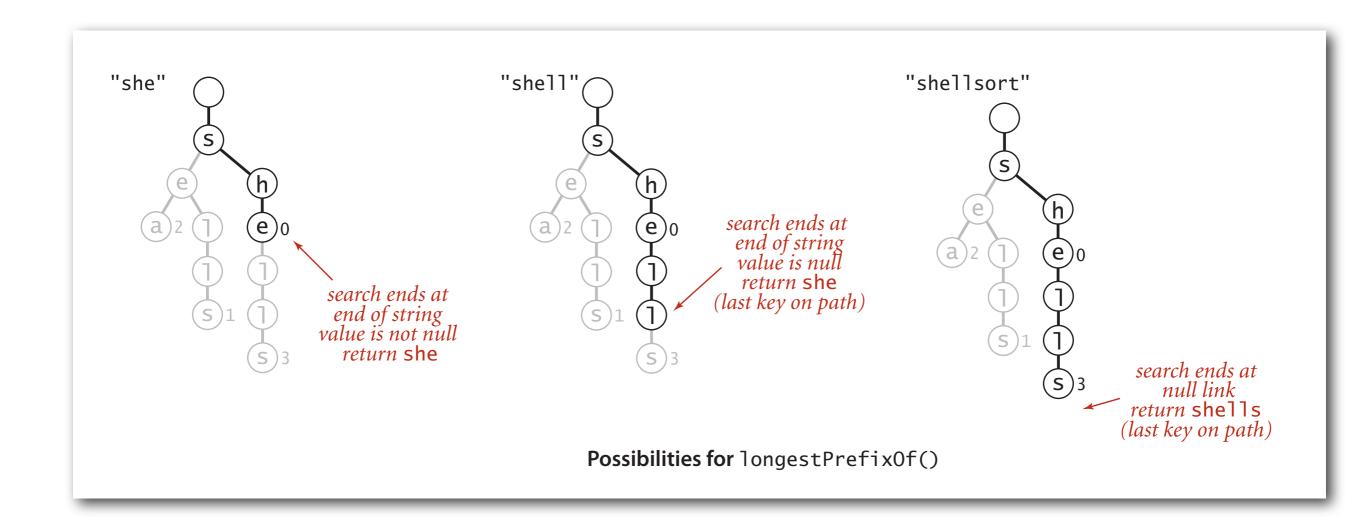
Note. Not the same as floor.

```
prefix("128.112.100.16") = "128.112"
floor("128.112.100.16") = "128.112.055.15"
```

Longest prefix

Find longest key in symbol table that is a prefix of query string.

- Search for query string.
- Keep track of longest key encountered.



Longest prefix: Java implementation

Find longest key in symbol table that is a prefix of query string.

- Search for query string.
- Keep track of longest key encountered.

```
public String longestPrefixOf(String query)
   int length = search(root, query, 0, 0);
   return query.substring(0, length);
private int search (Node x, String query, int d, int length)
   if (x == null) return length;
   if (x.val != null) length = d;
   if (d == query.length()) return length;
   char c = query.charAt(d);
   return search(x.next[c], query, d+1, length);
```

T9 texting

Goal. Type text messages on a phone keypad.

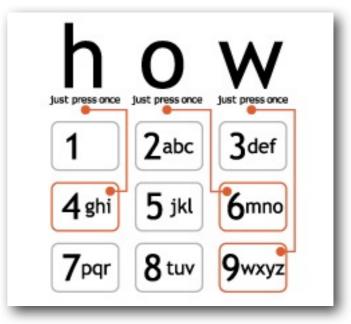
Multi-tap input. Enter a letter by repeatedly pressing a key until the desired letter appears.

T9 text input. ["A much faster and more fun way to enter text."]

- Find all words that correspond to given sequence of numbers.
- Press 0 to see all completion options.

Ex. hello

- Multi-tap: 4 4 3 3 5 5 5 5 6 6 6
- T9: **4 3 5 5 6**



www.t9.com

String symbol tables summary

A success story in algorithm design and analysis.

Red-black BST.

- ullet Performance guarantee: $\log N$ key compares.
- Supports ordered symbol table API.

Hash tables.

- Performance guarantee: constant number of probes.
- Requires good hash function for key type.

Tries. R-way, TST.

- ullet Performance guarantee: $\log N$ characters accessed.
- Supports character-based operations.

Bottom line. You can get at anything by examining 50-100 bits (!!!)