

COMP9101 Ass03

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4. You are on vacation for N days at a resort that has three possible activities 1, 2 and 3. For each day i , for each activity 1, 2 or 3, you've determined how

much enjoyment $e(i, j)$ ($1 \leq i \leq n$; $1 \leq j \leq 3$) you will get out of that activity if you do it on that particular day (the same activity might give you a different amounts of enjoyment at different days). However, you are not allowed to do the same activity two days in a row. Design an algorithm for determining the maximum total enjoyment possible over the entire stay of N days and the sequence of activities you should do at each day. (20 pts)

Q4.

This problem similar to the Q1, we can also use dynamic programming to solve the question.

We use $dp[i][j]$ to store the most enjoyment we get till the day i if we do activity j ;

$i \in [0, n]$, $j \in [1, 3]$

We know that the $e(i, j)$ to represent current day of the enjoyment we get if we do activity j .

Therefore, the base case is:

$dp[0][1] = e(0, 1)$

$dp[0][2] = e(0, 2)$

$dp[0][3] = e(0, 3)$

we know that we should not do the same activity two days in a row, therefore there are three situations for the day i :

Suppose we do activity 1 on day i :

$dp[i][1] = \max(dp[i-1][2] + e(i, 1), dp[i-1][3] + e(i, 1))$

Suppose we do activity 2 on day i :

$dp[i][2] = \max(dp[i-1][1] + e(i, 2), dp[i-1][3] + e(i, 2))$

Suppose we do activity 3 on day i :

$dp[i][3] = \max(dp[i-1][1] + e(i, 3), dp[i-1][2] + e(i, 3))$

for each algorithm, we should store each activity in a list called "res", if the result of enjoyment is the highest, the res should be the sequence of activity we should do at each day.

the recursion should be

therefore for day N , the max enjoyment we may get should be $\max(dp[N][1], dp[N][2], dp[N][3])$

the time complexity for this problem should be $O(3N)$, which is $O(N)$ at last.