#### CS1010S Tutorial 3

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## Today's Agenda

- Recap
- Question One
  - Standard solution
- Question 2
  - Discussion
  - Standard solution
  - Standard solution
- Question 3
  - Discussion
  - Standard solutions
- Question 4
  - Discussion
  - Standard solutions
- 6 Extra stuff: one more past midterm question

## Recap - Re-emphasize Recursion & Iteration

One way to convert between recursion and iteration!

- Recursion
  - base case (end condition) e.g. if n == 1: return something
  - recurrence relation e.g. f(n) = g(f(n-1))
  - starting status e.g. the first recursive function call
- Iteration
  - starting status (initialization) e.g n = 0
  - what you do in each loop e.g. the indented statements
  - end condition e.g. for i in range(1, n+1):
- How to convert between them? Notice their similarities

## Recap - Higher order functions

What have you learnt?

What are higher order functions? Explain non-abstractly

- Taking functions as input
  - Generalize common patterns
  - Functions taken in describe the varying features
- Or their return values are functions (closure)
  - Generate functions of **similar logic** but different internal parameters
  - Often return the function(lambda) defined in function
- Specifically written higher-order functions in this course
  - sum functions as input
  - product functions as input
  - fold functions as input
  - make\_something functions as output
  - accumulate not introduced yet

## Recap - Higher order function examples

```
def sum(term, a, next, b):
    if a > b:
        return 0
    else:
        return term(a) +
        sum(term, next(a), next, b)
```

Practice: find the sum of squares of integers?

## Recap - Higher order function examples

```
def product(term, a, next, b):
   if a > b:
      return 1
   else:
      return term(a) * product(term, next(a), next, b)
```

Practice: How would you define factorial in terms of product? Practice: How would you define exponential in terms of product?

## Recap - Higher order function examples

```
def fold(op, f, n):
   if n==0:
     return f(0)
   else:
     return op(f(n), fold(op, f, n-1))
```

Fold is much more general than sum and product. Practice: Can you mimic sum or product using fold?

## Recap - Order of Growth

- Method on lecture slides (may not work so well!)
  - Identify the basic computation steps
  - Try a few small values of n
  - Extrapolate for really large n
  - Look for worst case scenario
- Time complexity
  - Iteration : k nested loops means  $O(n^k)$
  - Recursion : k functions call of f(n-1) means  $O(k^n)$
  - In the future, look out for string/list concaternation, ...
- Space complexity
  - ullet Iteration : k nested loops means nothing, usually O(1)
  - Recursion : k functions call of f(n-1) means  $\mathcal{O}(n)$  (memory release)
  - In the future, look out for storage space taken by lists, tuples, dictionaries, . . .

## Question 1 Discussion

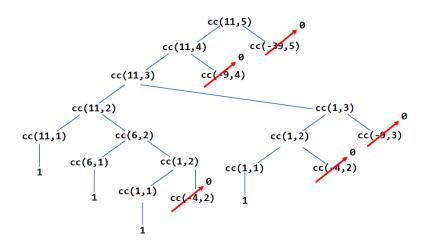
Draw the tree illustrating the process generated by the cc(amount, d) function given in the lecture, in making change for 11 cents. What are the orders of growth of the space and number of steps used by this process as the amount to be changed increases?

## Question 1 function cc code

```
def cc(amount, kinds of coins):
  if amount == 0:
    return 1
  elif (amount < 0) or (kinds of coins == 0):</pre>
                                                                Using 1 coin for
    return 0
                                                                   first kind
  else:
    return cc(amount - first_denomination(kinds_of_coins),
               kinds of coins) +
           cc(amount, kinds of coins-1)
                                                   Without using first
def first denomination(kinds of coins):
                                                      kind of coin
  ... <left as an exercise>
def count_change(amount)
  return cc(amount,5)
```

Takeaway: you can also try this top-down approach

## Question 1 Prof's solution



Order of growth of the space is O(n) (assuming amount is n) Order of growth of number of steps is  $O(2^n)$ 

## Question 2 Part (a) Discussion

A function f is defined by the rule that f(n) = n if n < 3 and f(n) = f(n-1) + 2f(n-2) + 3f(n-3) if n < 3.

 $\bullet$  Write a function that computes f by means of a recursive process.

## Question 2 Part (b) Discussion

A function f is defined by the rule that f(n) = n if n < 3 and f(n) = f(n-1) + 2f(n-2) + 3f(n-3) if n < 3.

 $\bullet$  Write a function that computes f by means of an iterative process.

# Question 2 (a) Zexin's solution

```
def f(n):
    if n < 3:
        return n
    else:
        return f(n-1) + 2 * f(n-2) + 3 * f(n-3)</pre>
```

- Order of growth for time complexity is  $O(2^n)$
- Order of growth for space complexity is O(n)
- Why is it so? Let's analyse it mathematically.

# Question 2 (b) Zexin's solution

```
def f(n):
    if n < 3:
        return n
    a = 0
    b = 1
    c = 2
    for i in range (3, n+1):
        d = 3*a + b*2 + c
        a = b
        b = c
        c = d
    return d
```

- Order of growth for time complexity is O(n)
- ullet Order of growth for space complexity is O(1)
- The storage space required remains constant as n increases.

## Question 3 Discussion

Write a function  $is\_fib(n)$  that returns True if n is a Fibonacci number, and False otherwise. What is the order of growth in terms of time and space for the function that you wrote? Explain.

There are a few things which we can be sure:

- There is no closed form formula to determine if n is Fibonacci
- If n is not Fibonacci, it must lie between two Fibonacci numbers
- If n is Fibonacci, it must appear in the Fibonacci series (increasing)

Here we form our algorithm:

- By monotonicity, if there is Fibonacci number greater than n, n is not.
- If any number in the series appear to be the same as n, n is Fibonacci.

## Question 3 Zexin's solutions

```
def is_fib(n):
    a = 0
    b = 1
    while n > b:
        c = a + b
        a = b
        b = c
    return n==b
```

- Order of growth for time complexity is  $O(\log n)$
- ullet Order of growth for space complexity is O(1)
- The storage space required remains constant as n increases.
- From the exponential growing rate of fib, we know that it will approximately take log n steps to reach n, so the time taken is logarithmically related to n.

## Question 4 Discussion

Define a function <code>make\_fare</code> that takes as arguments stage1, stage2, start\_fare, increment, block1, block2 and returns a function that calculates the taxi\_fare using those values.

Some obvious features of this function make\_fare

- Its return value should be a function.
- Its input parameters determine how the function returned will be.
- Also, the internal logic of the return function should be the same.

## Question 4 Other source to refer to

The function your *make\_fare* returns should possess same internal logic.

## Question 4 Zexin's solution

How to figure out this solution? Let's do live programming.

If time permits, we will go through this. If you all have understood, we shall skip this question and go to the second one.

```
def boo(x):
    return lambda y: x(x(y))
print(boo(boo)(lambda x:x+1)(5))
```

- What is the output of the print statement?
- Don't be intimidated by the complexity of problem first.
- Let's tackle this problem part by part.

This question is of the same type as the previous Question 3.

```
def boo(x):
    return lambda y: x(x(y))
print(boo(boo)(lambda x:x+1)(5))
```

Let's tackle the boo(boo) part first.

- boo(boo) takes in boo as parameter.
  - It will return lambda y : x(x(y)) with x replaced by boo.
  - So boo(boo) is equivalent to lambda y : boo(boo(y)),
  - which is just repeating boo twice for one single argument.

```
def boo(x):
    return lambda y: x(x(y))
print(boo(boo)(lambda x:x+1)(5))
```

Let's tackle the boo(boo)(lambda x : x + 1) part then.

- boo(boo) basically repeats boo twice for one single argument.
  - Let's define a function to simplify lambda x : x + 1, addOne.
  - Now we need to evaluate boo(boo(addOne)) layer by layer.
- boo(addOne) will return lambda y : addOne(addOne(y)),
- which is just lambda y: y + 2, and let's define it to be addTwo
- Similarly, boo(addTwo) will be addFour (lambda y: y + 4)

```
def boo(x):
    return lambda y: x(x(y))
print(boo(boo)(lambda x:x+1)(5))
```

Let's tackle the boo(boo)(lambda x : x + 1)(5) part then.

- boo(boo)(lambda x : x + 1) basically add four to the input.
- addFour(5) will return 9.
- Is this answer correct? Let's test it out!

```
x = 5
y = 10

def add_5(y):
    return lambda: y + 5

add_5(x)()
```

- This question is definitely easier than the previous one.
- Takeaway: they may test you on something never taught
- Fret not, most of you should be able to solve it.

```
x = 5
y = 10

def add_5(y):
    return lambda: y + 5

add_5(x)()
```

- What does lambda : y + 5 mean?
- Let's guess!

```
x = 5
y = 10

def add_5(y):
    return lambda: y + 5

add_5(x)()
```

- Now that we have reached concensus on the lambda
- Notice that this question tests on scoping as well.
- Look out for the difference between local y and global y

```
x = 5
y = 10

def add_5(y):
    return lambda: y + 5

add_5(x)()
```

- $add_{-}5$  takes x in as local y, and returns lambda: local y + 5
- Since x is always 5, it returns *lambda* : 10
- And so  $add_{-}5(x)()$  returns 10.

#### Feedback & more

• Slides + relevant material available at:

https://github.com/wangzexin/Teaching

• After the tutorial, if you have further questions:

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# Thank You

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