

Solving this set of $m + n$ equations in the $m + n$ variables x^* , ν^* gives the optimal primal and dual variables for (5.50).

Example 5.2 *Water-filling.* We consider the convex optimization problem

$$\begin{aligned} & \text{minimize} && -\sum_{i=1}^n \log(\alpha_i + x_i) \\ & \text{subject to} && x \succeq 0, \quad \mathbf{1}^T x = 1, \end{aligned}$$

where $\alpha_i > 0$. This problem arises in information theory, in allocating power to a set of n communication channels. The variable x_i represents the transmitter power allocated to the i th channel, and $\log(\alpha_i + x_i)$ gives the capacity or communication rate of the channel, so the problem is to allocate a total power of one to the channels, in order to maximize the total communication rate.

Introducing Lagrange multipliers $\lambda^* \in \mathbf{R}^n$ for the inequality constraints $x^* \succeq 0$, and a multiplier $\nu^* \in \mathbf{R}$ for the equality constraint $\mathbf{1}^T x = 1$, we obtain the KKT conditions

$$\begin{aligned} x^* \succeq 0, \quad \mathbf{1}^T x^* = 1, \quad \lambda^* \succeq 0, \quad \lambda_i^* x_i^* = 0, \quad i = 1, \dots, n, \\ -1/(\alpha_i + x_i^*) - \lambda_i^* + \nu^* = 0, \quad i = 1, \dots, n. \end{aligned}$$

We can directly solve these equations to find x^* , λ^* , and ν^* . We start by noting that λ^* acts as a slack variable in the last equation, so it can be eliminated, leaving

$$\begin{aligned} x^* \succeq 0, \quad \mathbf{1}^T x^* = 1, \quad x_i^* (\nu^* - 1/(\alpha_i + x_i^*)) = 0, \quad i = 1, \dots, n, \\ \nu^* \geq 1/(\alpha_i + x_i^*), \quad i = 1, \dots, n. \end{aligned}$$

If $\nu^* < 1/\alpha_i$, this last condition can only hold if $x_i^* > 0$, which by the third condition implies that $\nu^* = 1/(\alpha_i + x_i^*)$. Solving for x_i^* , we conclude that $x_i^* = 1/\nu^* - \alpha_i$ if $\nu^* < 1/\alpha_i$. If $\nu^* \geq 1/\alpha_i$, then $x_i^* > 0$ is impossible, because it would imply $\nu^* \geq 1/\alpha_i > 1/(\alpha_i + x_i^*)$, which violates the complementary slackness condition. Therefore, $x_i^* = 0$ if $\nu^* \geq 1/\alpha_i$. Thus we have

$$x_i^* = \begin{cases} 1/\nu^* - \alpha_i & \nu^* < 1/\alpha_i \\ 0 & \nu^* \geq 1/\alpha_i, \end{cases}$$

or, put more simply, $x_i^* = \max\{0, 1/\nu^* - \alpha_i\}$. Substituting this expression for x_i^* into the condition $\mathbf{1}^T x^* = 1$ we obtain

$$\sum_{i=1}^n \max\{0, 1/\nu^* - \alpha_i\} = 1.$$

The lefthand side is a piecewise-linear increasing function of $1/\nu^*$, with breakpoints at α_i , so the equation has a unique solution which is readily determined.

This solution method is called *water-filling* for the following reason. We think of α_i as the ground level above patch i , and then flood the region with water to a depth $1/\nu^*$, as illustrated in figure 5.7. The total amount of water used is then $\sum_{i=1}^n \max\{0, 1/\nu^* - \alpha_i\}$. We then increase the flood level until we have used a total amount of water equal to one. The depth of water above patch i is then the optimal value x_i^* .
