

Section A

Question 1

13

- (i) a
- (ii) $\begin{matrix} 2 & 1 \\ 1 & 2 \end{matrix}$
- (iii) The pages will increase.
- (iv) a
- (v) 7
- (vi) the hitting ratio are h_1, h_2, h_3 for shared cache, RAM and disk.
time are t_1, t_2, t_3 for shared cache, RAM and disk.
 $T_{eff} = h_1 t_1 + (1-h_1) h_2 t_2 + (1-h_1)(1-h_2) h_3 t_3$
- (vii) (a). $\log_2(1024) = 10 \rightarrow$ access a block
 $\log_2(4) = 2 \rightarrow$ specific byte
(b). remaining bits are used for tag.
- (viii) $\sqrt{4m}$
- (ix) False
 $H_k = \frac{S_k}{K \cdot \tau} = \frac{1}{K \cdot \tau} \cdot S_k$ linear relationship
- (x) 1 False
- (xi) False
- (xii) the memory could be accessed concurrently.

Question 2

(i). For sequential:

Load y
 Load a
 Load b
 Load c
 Load x
 MUL b, b, x // $b = b * x$
 MUL x, x, x // $x = x * x$
 MUL a, a, x // $a = a * x$
 ADD y, a, b, c // $y = a + b + c$

For VLIW:

Load b Load x
 Load a Load c MUL b, b, x
 MUL x, x, x ADD b, b, c
 MUL a, a, x
 ADD y, a, b.

(ii). $\frac{5}{9}$

(iii) $y = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0$

~~$y = ax^2 + bx + c$~~

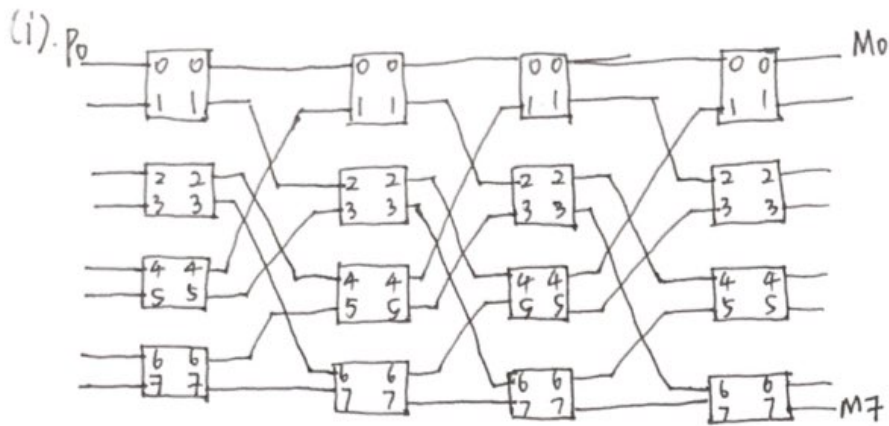
~~$O(\log(ax^2 + bx + c)) = O(\log t)$~~

Polynomial Time Complexity: $T(n) = O(n^k)$ for all positive k .

Here, we have the upper bound of the time complexity, which is $T(n) = O(x^n)$ for all positive n .

Question 4

(3)



(ii) P_0 to M_0 : 1

P_0 to M_7 : 1. Yes, it is hold for all pairs.

(iii) suppose switch 1 fails.
 (P_0, M_0) (P_4, M_2) fails.

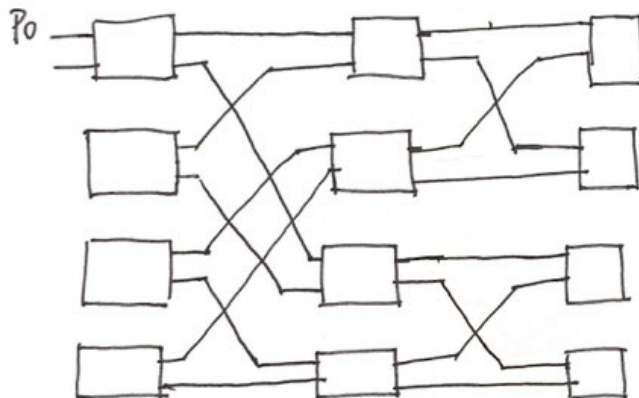
(iv) (P_0, M_0) fails.

(v) If only the top link ~~is~~ fails, the upper link still can use, and (P_4, M_2) will not fail.

(vi) Yes, it will increase.

The base-line network has self-routing ability.

Base Line.



Question 5

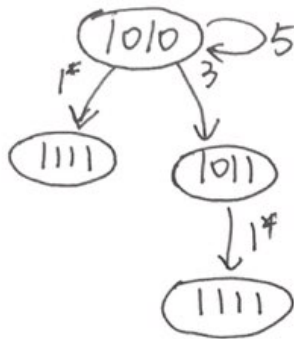
(4)

i) $FL = (4, 2)$

$P_L = (1, 3, 5, 6)$

(ii) (1010)

(iii)



(iv) (1)
 $(3, 1)$

(v) $MAL = 2$

(vi) True

(vii) True