

## Date \_\_\_\_\_

Figure 1

2. Host

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graph LR; H1[1] -.-> S[Server]; H2[2] -.-> S; H12[12] -.-> S;
```

For switch,

the throughput is 1200 Mbps.

3.  $AL_5 - AL_4 - AL_3 - AL_2 - AL_1 - RL_1 - RL_2 - RL_3 - RL_2 - RL_1 - BL_1 - BL_2 - BL_3 - BL_4 - BL_5$

4. The configuration message is  $\langle Y, 7, 5 \rangle$

5. 1 ms When the link ~~is~~ bandwidth is double, the total ~~time~~ throughput is no changed, so the time became  $\frac{1}{2} \times 2 \text{ ms} = 1 \text{ ms}$

## Section 2

6. (a) For the list of the outstanding frames

$$\{14, 15, 0, 1, 2, 3\}$$
~~10 - 4 = 6~~

So the 6 sequence of frames has been send.

(b) When receiving RR2, means  $\{14, 15, 0, 1, 2\}$  has been received, so sender window updated  $\{4, 5, 6, 7, 8, 9, 10, 11, 12\}$

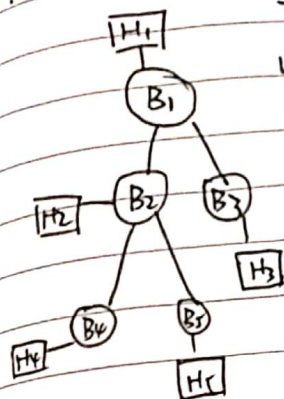
7. 
$$\begin{array}{r} 10010110 \\ 101 \overline{) 101010011000} \\ \underline{101} \phantom{00} \\ 00100 \phantom{00} \\ \underline{101} \phantom{00} \\ 111 \phantom{00} \\ \underline{101} \phantom{00} \\ 100 \phantom{00} \\ \underline{101} \phantom{00} \\ 100 \phantom{00} \\ \underline{101} \phantom{00} \end{array}$$

$E[x] = 1$ , means there is an error

Therefore, the receiver can detect an error.



8. (a) When  $H_1$  sends to  $H_4$ , every bridge forwards the packet through all the interfaces as none of the bridges knows where  $H_4$  is; all bridges learn where  $H_1$  is.



(b) When  $H_5$  sends to  $H_1$ , the packet use the path  $H_5 - B_5 - B_3 - B_1 - H_1$ , as every bridge knows where  $H_1$  is;  $B_5, B_3, B_1$  learn where  $H_5$  is, Bridge  $B_2, B_4$  does not learn where  $H_5$  is.

The ~~following~~ forwarding tables are built as follows:

Bridge  $B_1$

Host	Port
$H_1$	$B_1 \rightarrow H_1$
$H_4$	$B_1 \rightarrow B_2$

$B_2$

Host	Port
$H_1$	$B_2 \rightarrow B_1$
$H_5$	$B_2 \rightarrow B_3$

$B_3$

Host	Port
$H_1$	$B_3 \rightarrow B_1$

$B_4$

Host	Port
$H_1$	$B_4 \rightarrow B_1$

$B_5$

Host	Port
$H_1$	$B_5 \rightarrow B_3$
$H_5$	$B_5 \rightarrow H_5$

9. I don't think, each host has 0.5 probability of winning the race?

No.

Reasons: For both A and B, this collision is the second collision.  
means: A's slot is  $[0, 3]$  B's slot is  $[0, 3]$

For A, if A win the race, when  $A=0, B=1, 2, 3$   
when  $A=1, B=2, 3$   
when  $A=2, B=3$   
when  $A=3, B=\phi$

$$\text{Therefore, } P(A \text{ win the race}) = \frac{3+2+1}{4 \times 4} = \frac{3}{8}$$

Using the similar analyse, we also can get  $P(B \text{ win the race}) = \frac{3}{8}$

In a word, EACH host has  $\frac{3}{8}$  probability of winning the race.





$$T_f = \frac{10 \cdot 10^{12}}{3 \times 60 \times 60} = 10^8 \text{ bps}$$

$$T_{P_1} = \frac{2000 \times 10^3}{3 \times 10^8} \approx 66 \text{ ms}$$

$$T_{P_2} = 5 \times 2000 = 10000 \text{ us} = 10 \text{ ms}$$

