

Exercise 10.1

Date

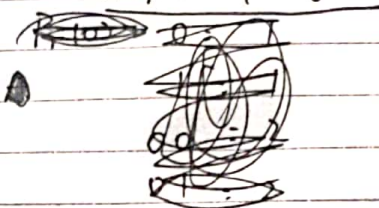
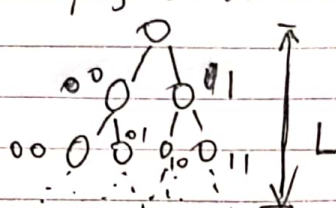
No.

Following the statement of this ~~the~~ binary huffman code of this sequence.

$$Pr(X_i) = \frac{1}{n}$$

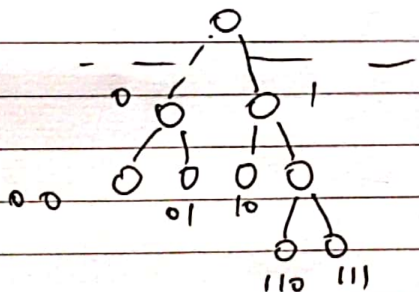
Then we use Huffman algorithm,

① When $n = 2^L$, we can know that this huffman tree structure can construct a perfect distribution



Therefore, every symbol has same length L and the average length is L

② When $n = 2^L + 1$, it means one node was added in the previous huffman tree, but for the existence of huffman tree, one node must be combined ~~with~~ this new node.



For $(2^L - 1)$ symbol, their length are all L
For 2 symbol, their length are $L+1$

$$\begin{aligned} \text{Average length} &= (2^L - 1) \cdot \frac{1}{2^L + 1} \cdot L + 2 \cdot \frac{1}{2^L + 1} \cdot (L+1) \\ &= \frac{2^L \cdot L + L + 2}{2^L + 1} \\ &= L + \frac{2}{2^L + 1} \end{aligned}$$



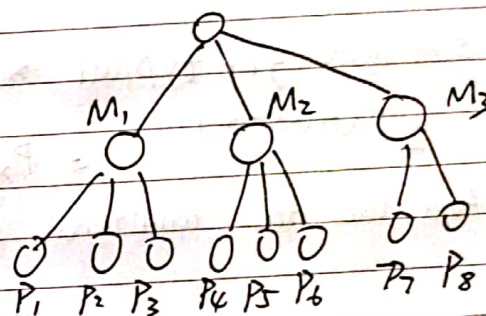
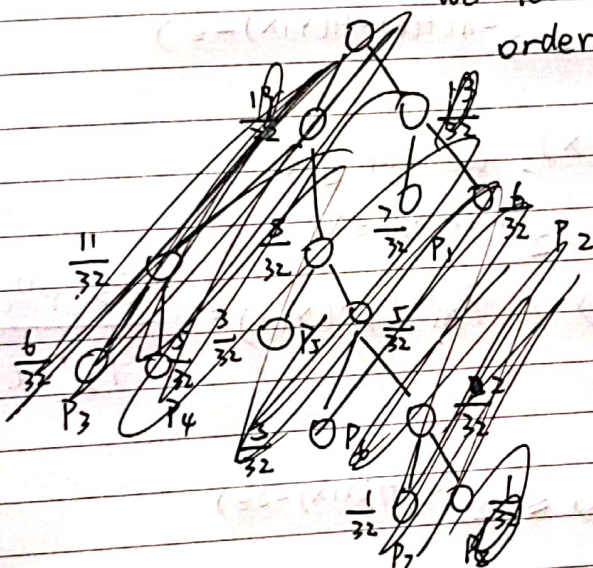
Exercise 10.3

a) Without mixing the water, the minimum expected number of taste is 1 because you ~~can~~ select the sugared water at the first time

In this order: $(\frac{7}{32}, \frac{6}{32}, \frac{6}{32}, \frac{5}{32}, \frac{3}{32}, \frac{3}{32}, \frac{1}{32}, \frac{1}{32})$

$P_1, P_2, P_3, P_4, P_5, P_6, P_7, P_8$

b). Allowing mixing, ~~we use Huffman code strategy.~~ we let ~~this~~ 8 bottles of water in this order to be mixed



After tasting (M_1, M_2, M_3) , we can output the sugar mixture, then tasting each part in this mixture

Expected number of ~~taste~~ taste: 5.9375
Maximal number of taste: 6.

In this strategy, $[P_1, P_2, P_3]$ mixture should be tasted first.



Exercise 10.4

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No.

a). From the statement of joint typical sequence, we know

$$A_{\epsilon}^{(n)}(XY) := \left\{ (x, y) \in X^n \times Y^n : \left| \frac{1}{n} \log \frac{1}{P_X^n Y^n(x, y)} - H(X, Y) \right| \leq \epsilon \right\}$$

$$\Downarrow$$

$$H(X, Y) - \epsilon \leq \frac{1}{n} \log \frac{1}{P_X^n Y^n(x, y)} \leq H(X, Y) + \epsilon, \text{ for all sequence } (x, y) \in A_{\epsilon}^{(n)}$$

$$\Downarrow$$

$$H(X) + H(Y|X) - \epsilon \leq \frac{1}{n} \log \frac{1}{P_X^n Y^n(x, y)} \leq H(X) + H(Y|X) + \epsilon$$

$$\Downarrow$$

$$2^{-n(H(X) + H(Y|X) + \epsilon)} \leq P_X^n Y^n(x, y) \leq 2^{-n(H(X) + H(Y|X) - \epsilon)}$$

b). Similarly, for typical $x \in A_{\epsilon}^{(n)}(X)$, we can get

$$2^{-n(H(X) + \epsilon)} \leq P_X^n(x) \leq 2^{-n(H(X) - \epsilon)}$$

And use the conditional probability, $P_{Y^n|X^n}(y|x) = \frac{P_X^n Y^n(x, y)}{P_X^n(x)}$

$$\Downarrow$$

$$2^{-n(H(Y|X) + 2\epsilon)} \leq P_{Y^n|X^n}(y|x) \leq 2^{-n(H(Y|X) - 2\epsilon)}$$

c). Since every sequence in $A_{\epsilon}^{(n)}(X, Y)$ has probability at least $2^{-n(H(X, Y) + \epsilon)}$ by definition, there can be at most $2^{n(H(X, Y) + \epsilon)}$

such sequence in the typical set, otherwise the total probability of

(c). suppose the new random variables $Z_i = \log \frac{1}{P_X^n Y^n(x, y)} - H(X, Y)$

$$P[X^n, Y^n \in A_{\epsilon}^{(n)}(X, Y)] = P\left[\left|\frac{1}{n} \log \frac{1}{P_X^n Y^n(x, y)} - H(X, Y)\right| \leq \epsilon\right]$$

$$= P\left[\left|\frac{1}{n} \log \frac{1}{\prod_{i=1}^n P_X^n Y^n(x_i, y_i)} - H(X, Y)\right| \leq \epsilon\right]$$



$$= P \left[\left| \frac{1}{n} \sum_{i=1}^n \left(\log \frac{1}{P_{X,Y}(x,y)} - H(X,Y) \right) \right| \leq \varepsilon \right]$$

$$= P \left[\left| \frac{1}{n} \sum_{i=1}^n Z_i \right| \leq \varepsilon \right]$$

$$= 1 - P \left[\left| \frac{1}{n} \sum_{i=1}^n Z_i \right| > \varepsilon \right]$$

we can apply the weak law of large numbers (because Z_i are i.i.d and zero mean)

$$\lim_{n \rightarrow \infty} P \left[\left| \frac{1}{n} \sum_{i=1}^n Z_i \right| > \varepsilon \right] = 0$$

$$\lim_{n \rightarrow \infty} P \left[x^n, y^n \in A_{\varepsilon}^{(n)}(x, y) \right] = 1$$

$$P_{x^n, y^n} \left(\{x, y\} \in \mathcal{X}^n \times \mathcal{Y}^n : x \in A_{\varepsilon}^{(n)}(x), (x, y) \in A_{\varepsilon}^{(n)}(x, y)\} \right) \rightarrow 1$$

