#### ECE 587 Midterm Review

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Concepts related to Entropy

The Asymptotic Equipartition Property (AEP)

Entropy Rates of a Stochastic Process

**Data Compression** 

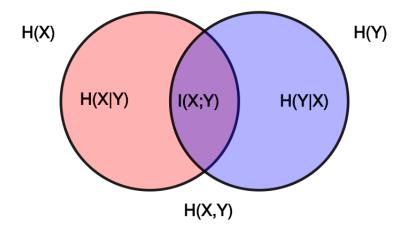


Figure: Venn diagram for entropy and mutual information

#### Entropy

- If the inequality is true prove it, otherwise, give a counter example
  - $\vdash$   $H(X|Y) \leq H(X)$
  - $\vdash$   $H(X, Y|Z) \geq H(X|Z)$
  - $\vdash$   $H(X|Z) \leq H(Z)$
  - ►  $H(X, Y, Z) H(X, Y) \ge H(X, Z) H(X)$
  - $\vdash$   $H(X|Z) \ge H(Z)$
- Relative entropy D(p||q)
  - ▶ show that  $D(p||q) \ge 0$
  - ▶ Is relative entropy symmetric, i.e D(p||q) = D(p||q)?

# Law of large numbers (LLN) for product of random variables

- ▶ Since  $X_i = \exp\{\ln X_i\}$
- we have  $\sqrt[n]{\prod_{i=1}^n X_i} = \exp\{\frac{1}{n}\sum_{i=1}^n \ln X_i\}.$
- Hence

$$\sqrt[n]{\prod_{i=1}^{n} X_i} = \exp\{\frac{1}{n} \sum_{i=1}^{n} \ln X_i\} 
\rightarrow \exp\{\mathbb{E} \ln X\} \quad (LLN)$$
(1)

$$\leq \exp\{\ln \mathbb{E}X\} = \mathbb{E}X$$
 (Jensen's Inequality) (3)

## The Asymptotic Equipartition Property (AEP)

▶ AEP. Let  $X_i$  be  $iid \sim p(x), x \in \{1, 2, \dots, m\}$ . Let  $\mu = \mathbb{E}X$  and  $H = -\sum p(x)log(x)$ . Let  $A^n = \{x^n \in \mathcal{X}^n : \left| -\frac{1}{n}\log p(x^n) - H \right| \le \epsilon$ . Let  $B^n = \{x^n \in \mathcal{X}^n : \left| -\frac{1}{n}\sum_{i=1}^n X_i - \mu \right| \le \epsilon \}$ 

- ▶  $Pr\{X^n \in A^n\} \rightarrow 1$ ?
- ▶ Does  $Pr\{X^n \in A^n \cap B^n\} \rightarrow 1$ ?
- ▶ Show that  $|A^n \cap B^n| \le 2^{n(H+\epsilon)}$  for all n.
- ▶ Show that  $|A^n \cap B^n| \le (\frac{1}{2})2^{n(H-\epsilon)}$  for n sufficiently large.

### Entropy rate of random walk on a weighted graphs

- ▶ The nodes of graphs are random variables, with state distribution  $\pi_t$  at time t and transition probability matrix T.
- T is determined by the weights over the edges.
- Stationary distribution is a distribution pi over states, such that  $\pi_t = \pi_{t+1}$  (shift invariant).
- ▶ The stationary distribution  $\pi(i)$ ,  $i = 1, \dots, |\mathcal{X}|$  satisfies

$$\pi = T\pi. \tag{4}$$

#### **Entropy Rates of a Stochastic Process**

Monotonicity of entropy per element. For a stationary stochastic process  $X_1, X_2, \dots, X_n$ , show that

 $\frac{H(X_1, X_2, \cdots, X_n)}{n} \le \frac{H(X_1, X_2, \cdots, X_{n-1})}{n-1}$  (5)

$$\frac{H(X_1, X_2, \cdots, X_n)}{n} \ge H(X_n | X_{n-1}, \cdots, X_1). \tag{6}$$

- Entropy rates of Markov Chains
  - What is the stationary distribution  $\pi$
  - Find the entropy rate the two-state Markov Chain with transition matrix

$$\begin{pmatrix} 1-p & p \\ q & 1-q \end{pmatrix}$$

#### Data compression

Huffman Coding Consider the random variable

$$X = \begin{pmatrix} x_1 & x_2 & x_3 & x_4 & x_5 & x_6 & X_7 \\ 0.49 & 0.26 & 0.12 & 0.04 & 0.04 & 0.03 & 0.02 \end{pmatrix}$$

- Find the binary Huffman code for X.
- Find the expected code length for this encoding.
- Find a ternary Huffman code for X.

# Summary

Good Luck!