

DERS :

TARİH :

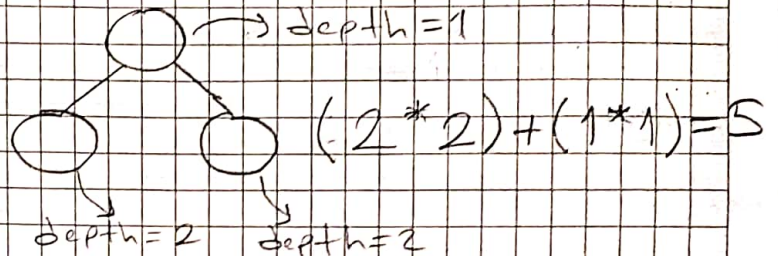
KONU :

GÜN :

1)

a) Calculate the total depth of nodes in a complete binary tree of height  $h$ .

if height is 2;



$$\text{depth in level-1} = (\text{height of level-1}) * (\text{number of nodes at level-1})$$

$$\Rightarrow h * 2^{h-1} = 1 * 2^0 = 1$$

$$\text{depth in level-2} = h * 2^{h-1} = 2 * 2^1 = 2$$

$$\text{depth in level-3} = h * 2^{h-1} = 3 * 2^2 = 12$$

So, total number of depth in a tree of height 3 =  $12 + 2 + 1$

According to the above calculations, we can generalize formula:

$$(\text{Total depth of } h \text{ height tree}) = (2^{h-1} \cdot h) + (2^{h-2} \cdot (h-1)) + \dots + (2^0 \cdot 1)$$

We can find total depth of nodes in complete binary tree of height  $h$  recursively:

In function `int total_depth(int h);`

1- if  $h=0$  } its our base case  
return 1

2- Else, return  $(2^{h-1} \cdot h) + \text{total\_depth}(h-1)$ .

\* if we generalize the formula more with using geometric series, we'll find this;  $a(h) = (h-1) \cdot 2^h + 1$



DERS : Worst case =  $\log n$

KONU : Best case = 1

TARİH : .....

GÜN : .....

b) Average number of comparisons in a BST

$$\text{Average number of Comparison in Binary Search Tree is} = \frac{(\text{Total number of Comparison})}{(\text{Number of Nodes})}$$

→ According to my research, total number of comparison has a same value with total depth of nodes

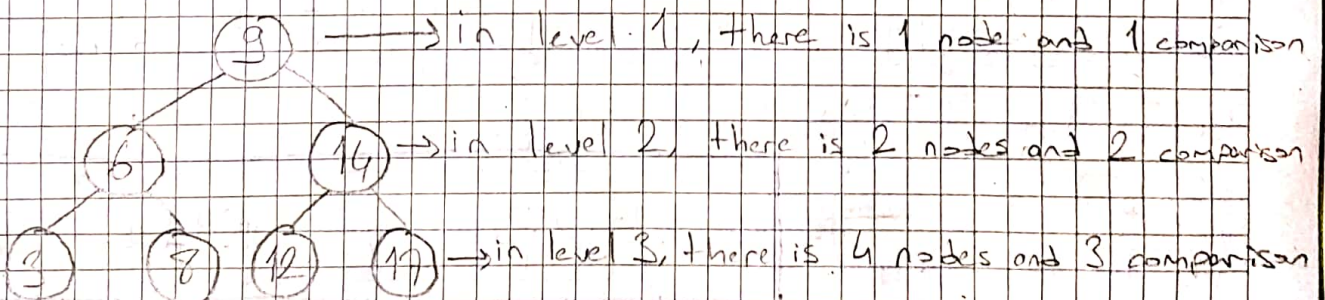
$$\text{For height } h; \underline{n = 2^h - 1}$$

$$\text{Total Number of Comparison} \Rightarrow a(h) = (h-1) \cdot 2^h + 1$$

$$\text{Number of Nodes} \Rightarrow n = 2^h - 1$$

$$\text{Average number of comparison is} \Rightarrow \frac{(h-1) \cdot 2^h + 1}{2^h - 1}$$

Proof :



$$\text{So average number of comparison} = \frac{1 \cdot 1 + 2 \cdot 2 + 4 \cdot 3}{7}$$

$$= 2.42$$

With our formula ( $h=3$ );

$$\text{Average num. of comparison} = \frac{(3-1) \cdot 2^3 + 1}{7} = \frac{2 \cdot 8 + 1}{7} = 2.42$$

$2.4 = 2.4$ , So formula is true.



C)

→ No, there is no restriction on the number of nodes in a full binary tree. Rule of being full binary tree is that each node has 2 or 0 children. It doesn't matter how many descendant they have.

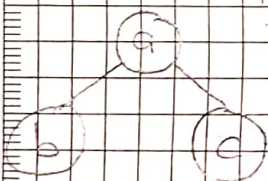
→ A leaf node is node that has no children

→ An internal node is node that has at least 1 child.

→ In full binary tree, nodes can have 0 or 2 children.

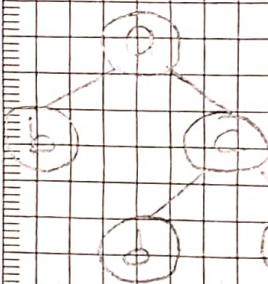
Lets start with 1 heighted binary tree.

(a) This tree has 1 leaf node and 0 internal node  
If we add 2 children to node a;



→ Node a is no more leaf node but it created 2 more leaf node.

→ This tree has 2 leaf node and 1 internal node now. Lets add 2 children to node c.



→ Now, c become internal node, d and e become leaf node.

→ So, number of internal node and number of leaf node increased by 1. The difference between them doesn't changed.

→ Leaf node = 3, internal node = 2

→ Num of nodes = Num of leaf nodes + Num of internal nodes

⇒ We know that the number of internal nodes is always 1 less than number of leaf nodes.

⇒ So, if we say num of internal nodes to k, k+1 became num of leaf nodes.

$$n = k + (k+1) \Rightarrow n = 2k+1 \Rightarrow \boxed{k = \frac{n-1}{2}}, \boxed{k+1 = \frac{n+1}{2}}$$

$$\text{Num of internal nodes} = \frac{n-1}{2}$$

$$\text{Num of leaf nodes} = \frac{n+1}{2} //$$



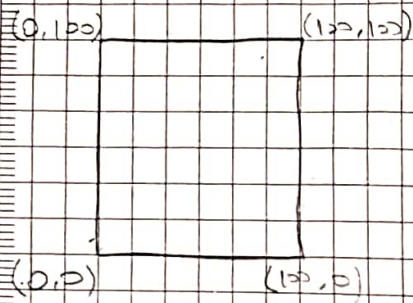
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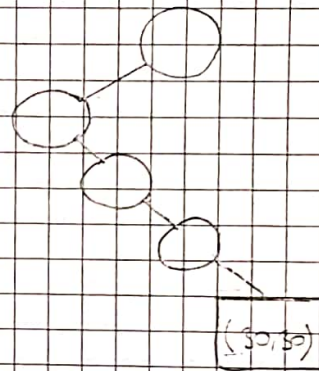
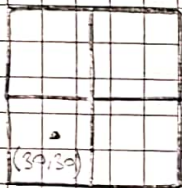
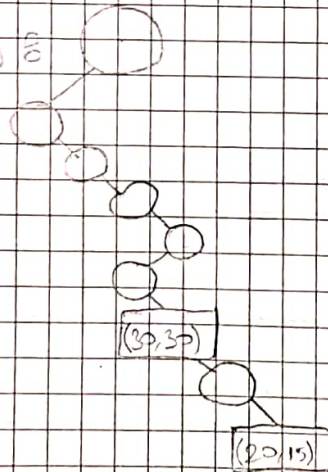
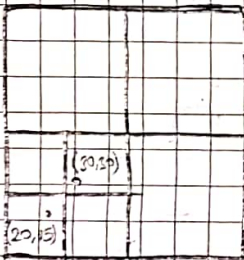
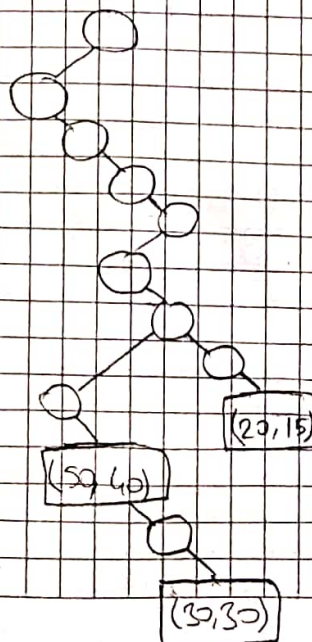
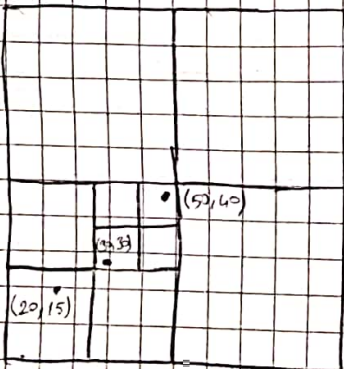
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GÜN : .....

2- Initial?



No Tree

Inserting  $(30,30)$ ?Inserting  $(20,15)$ ?Inserting  $(50,40)$ ?



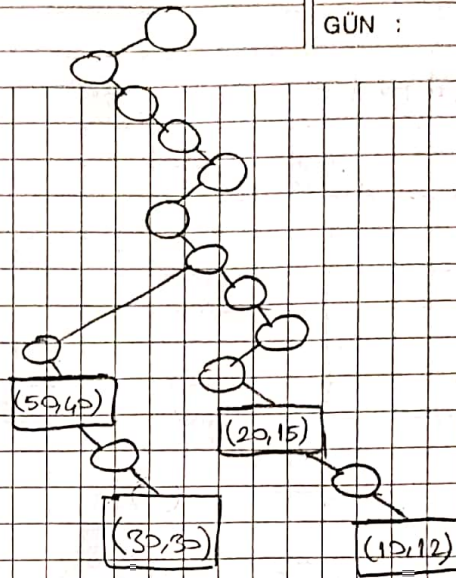
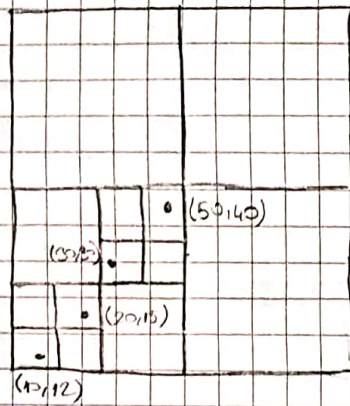
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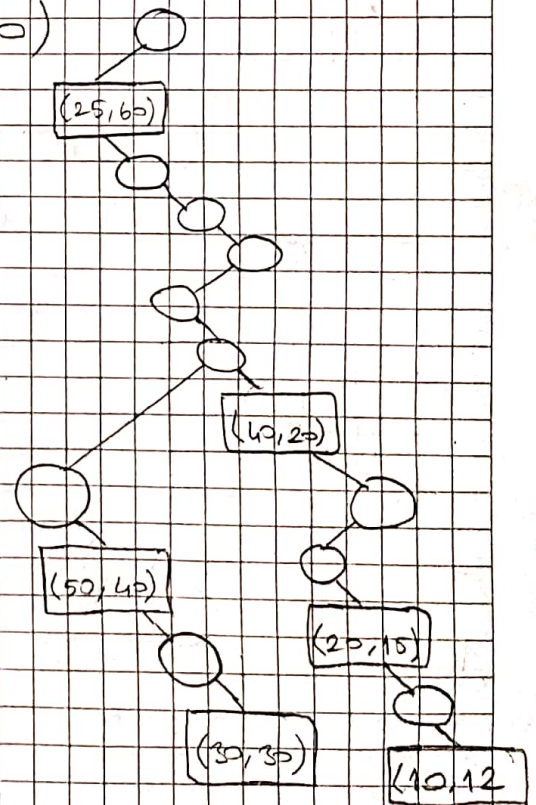
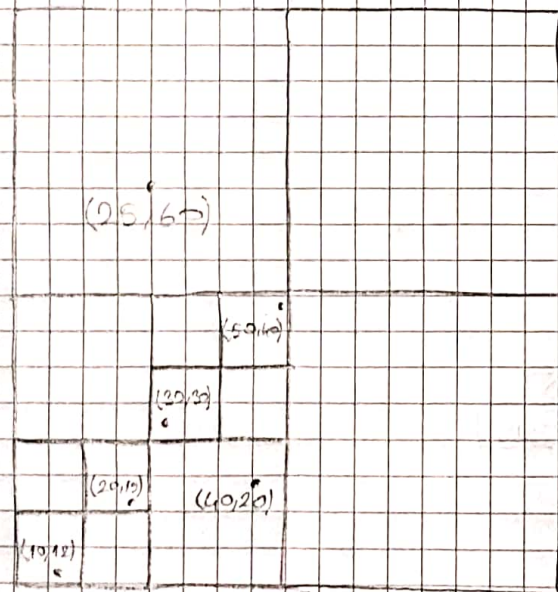
TARİH : .....

GÜN : .....

Inserting (10, 12) B

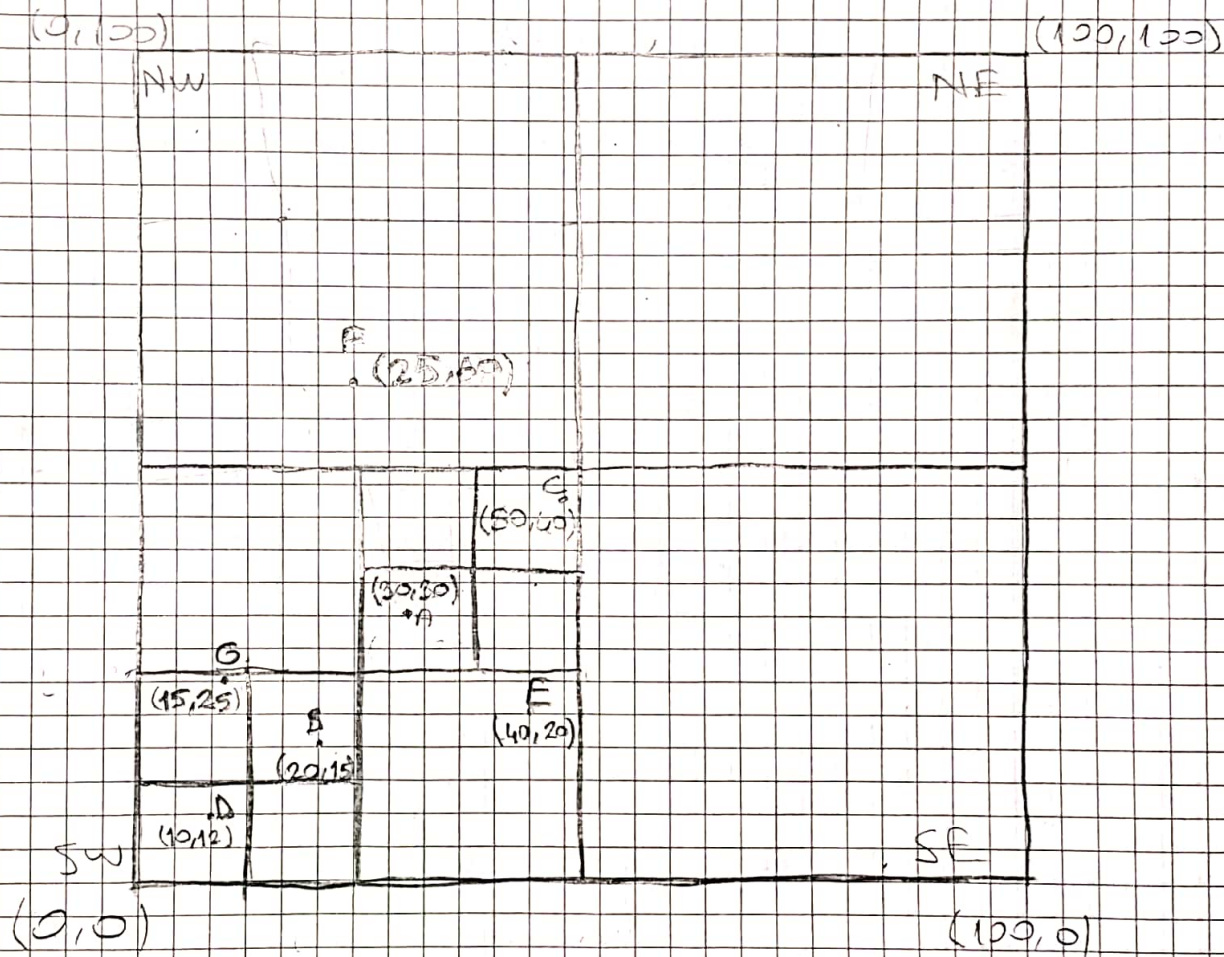


Inserting (40, 20) and (25, 60)





Inserting (15, 25) (Table)



DERS :

KONU :

TARİH :

GÜN :



DERS :

KONU :

Inserting (15, 25) (Binary Tree)

TARİH: / /

GÜN : / /

