

## Week 1 Summary Sheet

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Before you start, make sure you understand the basic operation between sets including the **Union ( $\cup$ )**, **Intersection ( $\cap$ )**, **Difference ( $-$ )** and **Complement ( $\bar{A}$ )**

### Basic definitions

#### 1. Alphabets

- finite **set** of symbols
- denoted as  $\Sigma$ , read as sigma
- *examples*  
the binary alphabet,  $\Sigma = \{0,1\}$   
a set of all lower-case letters,  $\Sigma = \{a, b, c, d, e, \dots, x, y, z\}$

#### 2. Strings

- finite **sequence** of symbols
- $w, x, y, z$
- *examples*  
given an alphabet  $\Sigma = \{0,1\}$ , define a string as  $w = 010110$   
given an alphabet  $\Sigma = \{a, b, c\}$ , define a string as  $x = aabbcc$   
all symbols must taken from a given alphabet

#### 3. Empty strings

- string with zero symbols, means nothing in the string
- denoted as  $\varepsilon$ , read as epsilon

#### 4. Length of a string

- number of symbols in the string
- length of  $w$  is  $|w|$ , read as the modulus of  $w$
- *examples*  
a given alphabet  $\Sigma = \{0,1\}$ , a string  $w = 010110$ , then  $|w| = 6$   
for empty string,  $|\varepsilon| = 0$

#### 5. Power of an alphabet

- $\Sigma^*$  is a set of **all strings** over an alphabet  $\Sigma$ , also known as **kleene start operation**
- $\Sigma^* = \Sigma^0 \cup \Sigma^1 \cup \Sigma^2 \cup \dots \cup \Sigma^{infinite}$
- more specifically,  $\Sigma^k$  is a **set** of **all** strings of length  $k$  chosen from  $\Sigma$
- read as  $\Sigma$  **to the power of  $k$**
- $\Sigma^0 = \{\varepsilon\}$
- given  $\Sigma = \{0,1\}$ , then we have  $\Sigma^1 = \{0,1\}$  and  $\Sigma^2 = \{00,01,10,11\}$
- $\Sigma^+$  is a set of nonempty strings from  $\Sigma$
- $\Sigma^+ = \Sigma^1 \cup \Sigma^2 \cup \dots \cup \Sigma^{infinite}$
- $\Sigma^* = \Sigma^+ \cup \Sigma^0 = \Sigma^+ \cup \{\varepsilon\}$
- **powerset**

a set of **all possible subsets** of the original set, including  $\varepsilon$  ( $\emptyset$ ) and **the set itself**

#### 6. Concatenation of strings

- $x = 1010$ , and  $y = 0111$

- $xy = 10100111$
- guess what the concatenation of  $yx$  is
- $w\varepsilon = \varepsilon w = w$

## 7. Language

- a **set of strings** chosen from  $\Sigma^*$
- denoted by  $L$
- *examples*

7.1 language of all strings consisting of  $n$  0s followed by  $n$  1s, for  $n \geq 0$

it can be denoted as  $L = \{\varepsilon, 01, 0011, 000111, \dots\}$

7.2 a set of strings of 0s and 1s with equal number of each

it can be denoted as  $L = \{\varepsilon, 01, 10, 0011, 1100, 000111, 111000, \dots\}$

7.3 a set of binary numbers whose value is a prime\*

It can be denoted as  $L = \{10, 11, 101, 111, 1011, \dots\}$

\* A prime number is a natural number greater than 1 that has no positive divisors other than 1 and itself. Number like **2, 3, 5, 7, 11, 13**

- empty language, denoted by  $\varphi$ , read as **phi**, means no strings

## Finite Automata (FA)

### 1. introduction

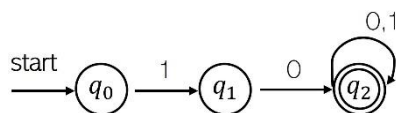
- a mathematical model, and **no memory**
- consists of **a set of states**, **transitions** among states based on input symbols
- deterministic finite automate (DFA)
- non-deterministic finite automata (NFA)

### 2. applications

- digital circuits, lexical analyzer, pattern matching (searching for a webpage)

### 3. design a DFA to accept the strings start with 10

#### 3.1 transition diagram



#### 3.2 list of transition functions

a DFA  $D = \{Q, \Sigma, \delta, q_0, F\}$

a set of states are denoted as  $Q = \{q_0, q_1, q_2\}$

an alphabet is  $\Sigma = \{0, 1\}$

transition functions are as follows,  $\delta$ , read as **delta**

- first transition is  $\delta(q_0, 1) = q_1$ , then  $\delta(q_1, 0) = q_2$ , then  $\delta(q_2, 0) = q_2$ , and  $\delta(q_2, 1) = q_2$
- start state is  $q_0$
- the final/ accept state is  $q_2$

#### 3.3 transition table

States	Input symbols	
	0	1
$\rightarrow q_0$	X	$q_1$
$q_1$	$q_2$	X
$* q_2$	$q_2$	$q_2$