

## Down Counters

All the examples considered so far have used counters that count up from binary zero to some maximum count value depending upon the Modulus value of the counter. When the counter reaches its maximum count value it is reset to binary zero and continues with the counting sequence. A down counter counts in a sequence which starts with some maximum count value and counts down to binary zero. It is then reset to the maximum count value and repeats the counting sequence. A Down-counter is implemented by connecting the  $\bar{Q}$  output instead of the Q output of all the flip-flops to the clock inputs of the next flip-flops. Figure 27.1

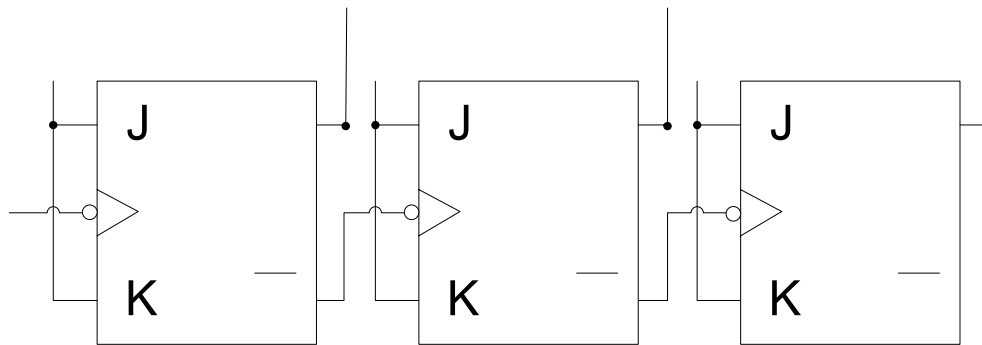


Figure 27.1a 3-bit Asynchronous Down-Counter

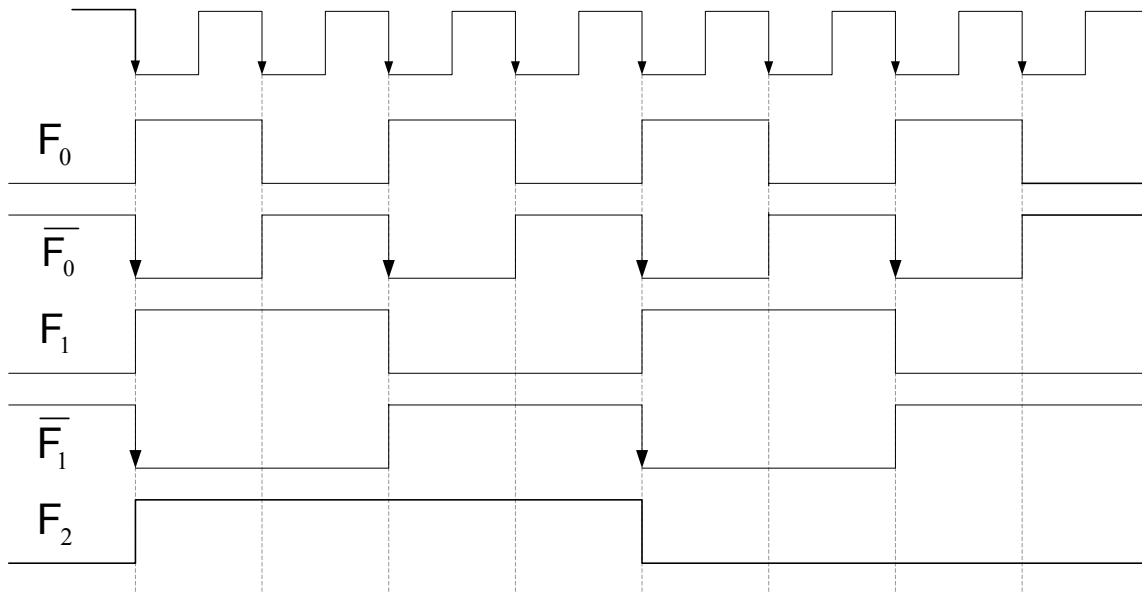


Figure 27.1b Timing diagram of a 3-bit Asynchronous Down-Counter

## Down Counter with truncated sequence

A down counter can be configured to count down a truncated sequence, similar to an up-counter which can count up to any truncated sequence. A down counter counts down from the maximum count value to some predefined count value which is the last count value in the truncated sequence. On reaching the last count value the down-counter is preset to the maximum count value instead of clearing the counter to zero count value as done in the case of an up-counter. The circuit shows a 3-bit down-counter configured to count down a truncated sequence from 111 to 011. On reaching the count value 011, the counter is preset to 111 when it is decremented to 010 on the negative clock transition. Figure 27.2

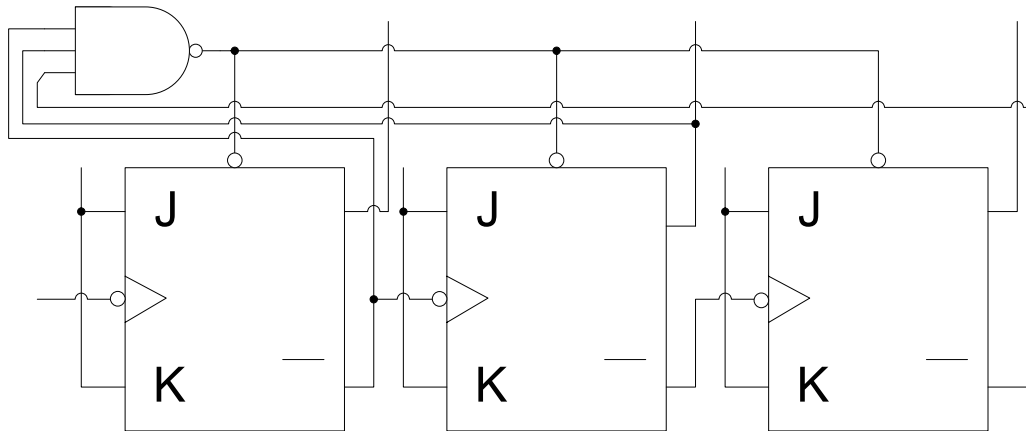


Figure 27.2a Down-counter configured to count a truncated sequence

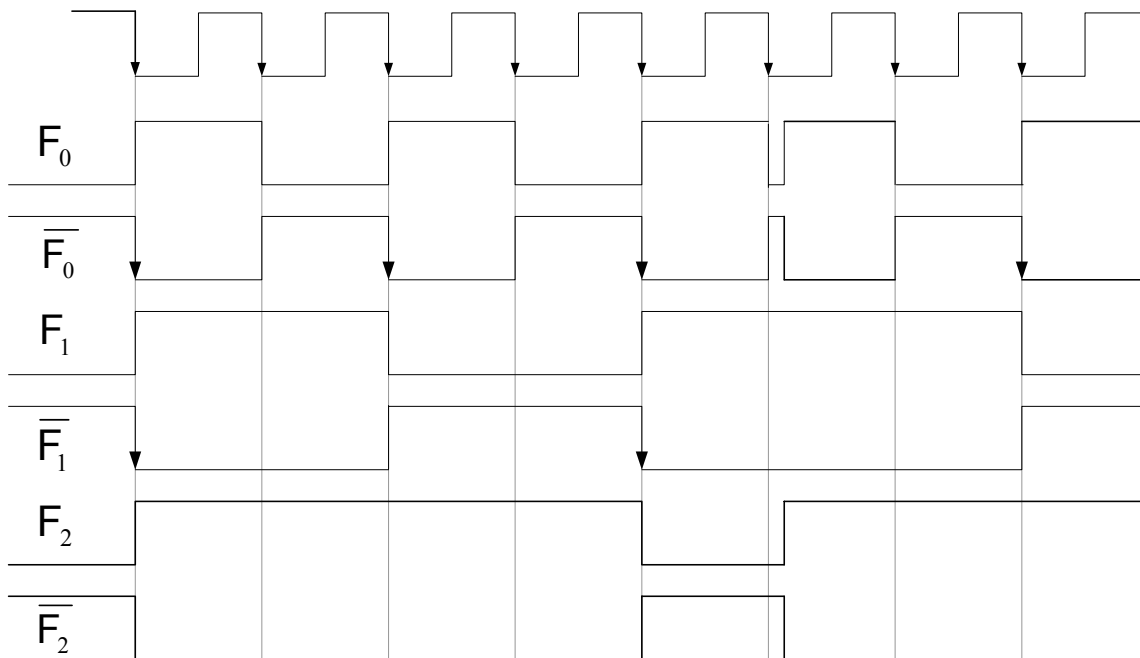


Figure 27.2b Timing diagram of a counter configured to count a truncated sequence

The counter counts down from 111 to 011 from interval  $t_1$  to interval  $t_5$ . At interval  $t_6$  the counter counts down to 010, the  $\overline{F_0}$ ,  $F_1$  and  $\overline{F_2}$  are set to logic 1, the output of the NAND gate is set to logic 0 which presets all the three flip-flops to state 111. The counter continues with its counting sequence and at the clock transition at interval  $t_7$  and  $t_8$  the counter count down to 110 and 101 respectively.

## Synchronous Counters

Asynchronous counters due to the delayed outputs caused by the rippling clock signal do not allow their operation with high frequency clock signals. Asynchronous counters having multiple bits also cause timing problems due to the excessive propagation delays.

Applications requiring 8, 16 and 32 counters and operating at high clock frequencies are implemented using Synchronous Counters. Synchronous counters use a common clock signal connected to the clock inputs of all the counter flip-flops. Therefore, on a clock transition all the flip-flops simultaneously change their output state. Figure 27.3.

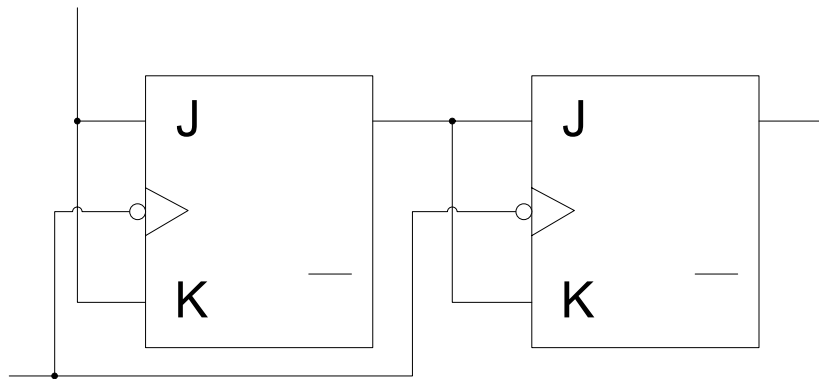


Figure 27.3a 2-bit Synchronous Counter

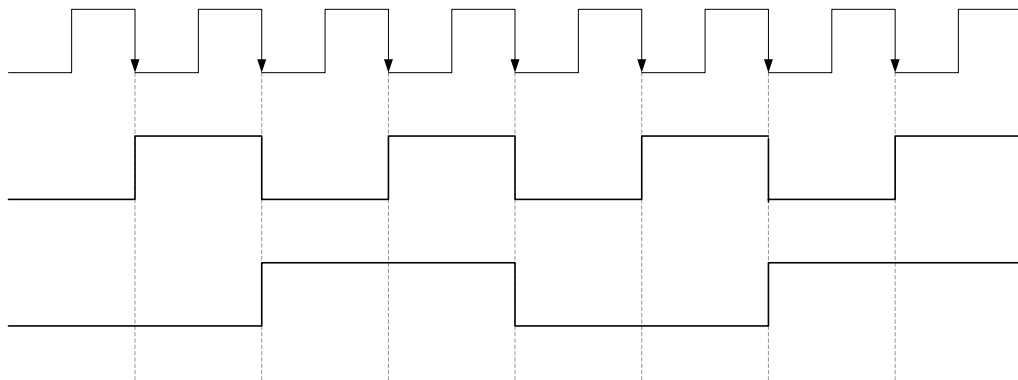


Figure 27.3b Timing diagram of a 2-bit Synchronous Counter

The 2-bit Synchronous counter has both its clock inputs connected to the clock signal. Both the flip-flops are reset to logic low states respectively. On a high to low clock transition at interval  $t_1$ , the  $F_0$  output of the first flip-flop toggles to logic high. Since the clock transition on the clock input of the second flip-flop also occurs at interval  $t_1$ , the J-K inputs of the second flip-flop are at interval  $t_1$  are at logic 0. The change at the inputs J-K to logic 1 of the second flip-flop occurs after a propagation delay  $t_{PLH}$  of the first flip-flop. Thus the output of the second flip-flop remains unchanged due to the input condition at the J-K inputs ( $J=0$ ,  $K=0$ ). At interval  $t_2$  the output  $F_0$  is at logic high (1) along with the J-K inputs as the three are connected together. On a clock transition at interval  $t_2$  the output  $F_0$  toggles to logic 0. At the very same instant the output  $F_1$  also toggles to logic 1. The inputs J-K of the second flip-flop is set to logic 0 after a propagation delay of  $t_{PHL}$  of the first flip-flop. At interval  $t_3$ , at the clock transition the output  $F_0$  toggles to logic 1. The inputs J-K of the second flip-flop at time interval  $t_3$  is logic 0 therefore at the clock transition the output  $F_1$  remains unchanged. The inputs J-K of the second flip-flop change after a propagation delay of  $t_{PLH}$ . Finally, at time interval  $t_4$ , the output  $F_0$  of the first flip-flop toggles to logic 0. The J-K inputs of the second flip-flop are at logic 1, therefore the output  $F_1$  of the second flip-flop is also set to logic 0.

### 3-bit & 4-bit Synchronous Counters

Multi-bit Synchronous Counters can easily be implemented by connecting together appropriate number of flip-flops together. The clock inputs of all the flip-flops are directly connected to a common clock signal. Implementing of Synchronous Counters larger than 2-bits requires the use of an AND gate. Figure 27.4

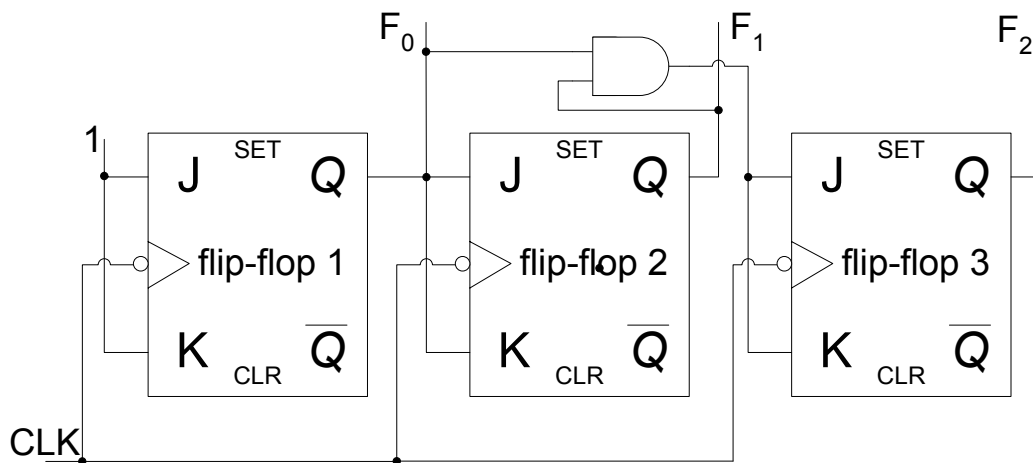


Figure 27.4a

A 3-bit Synchronous Counter

The operation of the 3-bit Synchronous Counter and the need for the AND gate can be understood by studying the timing diagram of the 3-bit counter. The timing of the first two flip-flops is identical to the timings of the 2-bit counter discussed earlier. The timing diagram shows that at interval  $t_4$ , the 3-bit counter should count from state 011 to 100. Similarly, at interval  $t_8$  the counter should count from state 111 to 000. At both the intervals the  $F_2$  output of the third flip-flop toggles to logic 1 and logic 0 respectively.

when the outputs  $F_0$  and  $F_1$  are both at logic 1. This is implemented by connecting the two outputs  $F_0$  and  $F_1$  to the inputs of a 2-input AND gate. The output of the AND gate is logic 1 when both its inputs ( $F_0$  and  $F_1$ ) are at logic 1. The output of the AND gate is connect to the J-K inputs of the third flip-flop. If the AND gate is not used and the J-K inputs of the third flip-flop are directly connected to the output  $F_1$  of the second flip-flop, the third flip-flop will change its state and set its output  $F_2$  to logic 1 at the time interval  $t_3$ . The count sequence is thus disturbed.

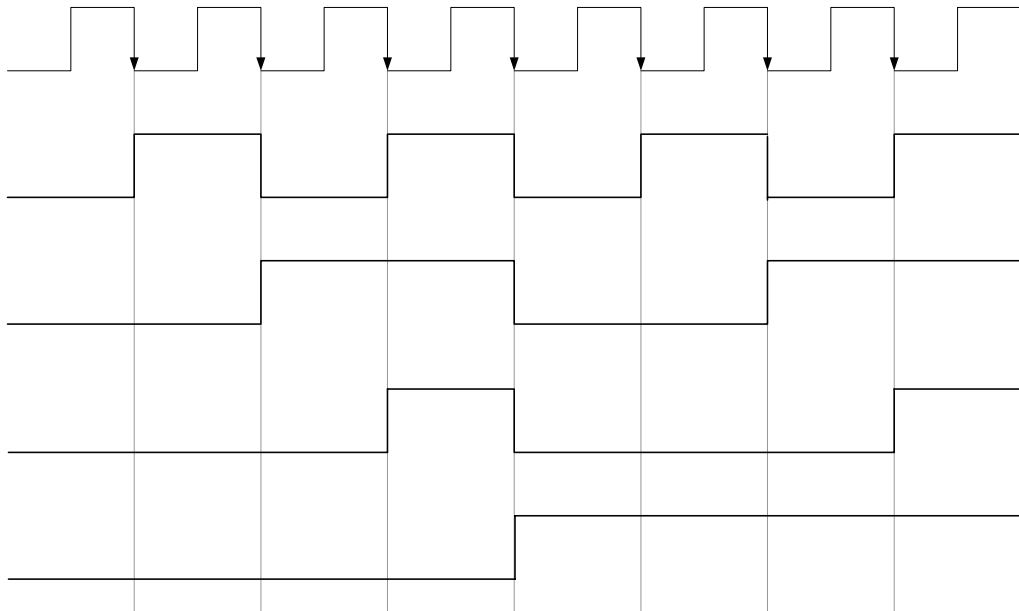


Figure 27.4b Timing diagram of a 3-bit Synchronous Counter

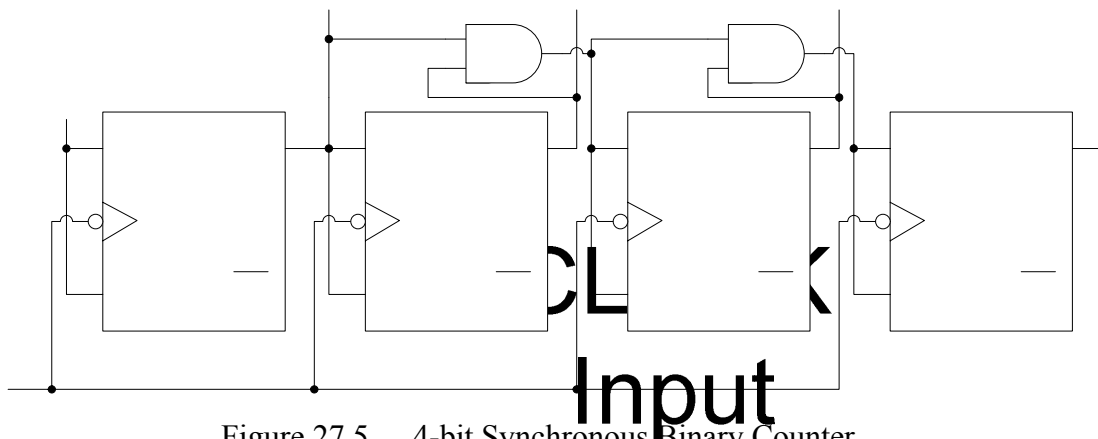


Figure 27.5 4-bit Synchronous Binary Counter

Larger counters can be implemented using similar AND gates. For example, a 4-bit counter uses four flip-flops. The counter circuit for the first three flip-flops is identical

to the 3-bit counter circuit. The input of the fourth flip-flop is connected through a 3-input AND gate with inputs  $F_0$ ,  $F_1$  and  $F_2$ . The fourth flip-flop changes its state when the outputs of the first three flip-flops are at logic 1. That is, the when the 4-bit counter is counting from 0111 to 1000 and 1111 to 0000. Figure 27.5

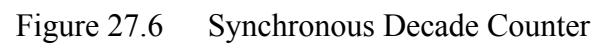
## 4-bit Synchronous Decade Counter

Earlier, an Asynchronous Decade counter has been discussed, which counts from state 0000 to 1001. The Asynchronous counter is cleared to state 0000 when the counter counts from 1001 to 1010. Synchronous counter can be implemented which counts from 0000 to 1001. In the synchronous counter, all the four flip-flops are connected to a common clock and are triggered simultaneously. However, instead of using the clear asynchronous inputs to clear the counter to the initial state, logic gates are used to reset the decade counter to state 0000 after it reaches state 1001. The implementation of the Synchronous Decade counter can be understood with the help of a function table that represents the operation of the Decade Counter. Table 27.1.

Input	Output			
Clock Pulses	$F_3$	$F_2$	$F_1$	$F_0$
1	0	0	0	0
2	0	0	0	1
3	0	0	1	0
4	0	0	1	1
5	0	1	0	0
6	0	1	0	1
7	0	1	1	0
8	0	1	1	1
9	1	0	0	0
10	1	0	0	1

Table 27.1 Output of a Synchronous Decade Counter

The output state of the first flip-flop  $F_0$  is shown to toggle between 1 and 0 on each clock transition. Therefore, the inputs J-K of the first flip-flop are connected to logic high. The output state of the second flip-flop  $F_1$  changes from logic 0 to logic 1 and vice-verse when  $F_0$  output is logic 1 and  $F_3$  output is logic 0. Therefore, the inputs J-K of the second flip-flop are connected to a function determined by the Boolean expression  $F_0 \overline{F_3}$ . The output state of the third flip-flop  $F_2$  changes from logic 0 to logic 1 and vice-versa when  $F_0$  and  $F_1$  outputs are both at logic 1. Therefore, the inputs J-K of the third flip-flop are connected to a function determined by the Boolean expression  $F_0 F_1$ . The output of the fourth flip-flop  $F_3$  changes its output state when outputs  $F_0$ ,  $F_1$  and  $F_2$  are at logic 1 or when outputs  $F_0$  and  $F_3$  are at logic 1. Therefore, the J-K inputs of the fourth flip-flop are connected to a function determined by a Boolean expression  $F_0 F_1 F_2 + F_0 F_3$ . The decade counter is shown in figure 27.6



# CLK