

Preliminary Observations on the Happy Ending Problem

Warren Jin

July 2023

1 Introduction and Literature Review

Given N points in the plane, we wish to check whether there must exist a convex polygon with n vertices. The lowest number of points required is denoted $k(n)$, such that \exists convex n -gon $\forall N \geq k(n)$. For simplicity, we can assume that no 3 points are collinear. This problem was posed by Erdős and Szekeres in this 1935 paper, where they conjectured that $k(n) = 2^{n-2} + 1$ (which I will refer to as the ES conjecture).¹ They managed to prove that $k(n) \geq 2^{n-2} + 1$ in a later paper.² A related problem is to find the infimum of the number of convex n -gons in a set of N points, which I will denote by $f(N, n)$. The case $f(N, 4)$ was considered by Hu, Chen, and Zhu, who found exact values for $N \leq 9$ and a lower bound for $N > 9$.³ I present an alternative method to find the lower bound of $f(N, 4)$ that coheres with Hu et al's finding on the growth rate of the function, and generalize it to any choice of n .

2 Some observations

For this article I will introduce some notation. Noting that if the power set of N points is taken to be a poset (P, \subseteq) , then every subset of the same number of points is pairwise incomparable. Thus, the number of elements on height H equals $\binom{N}{H}$. Evidently, the height of this poset is N . We can also see that P is a lattice, since the joins and meets for any two arbitrary elements are their union and intersection respectively, which are both elements of the power set.

Let x_i be an element of P with i elements. Define a function $A : P \rightarrow \mathbb{R}$, where $A(x)$ is the supremum of the possible areas enclosed by a cycle over the elements of x . We let $A(x_1) = A(x_2) = 0$. A cycle which gives this supremum does not have any edge crossings, because if the edges v_1v_2 and v_3v_4 crossed, then either v_1v_3 and v_2v_4 form a non-crossing cycle of larger area, or v_1v_4 and

¹http://www.numdam.org/item/CM_1935__2__463_0.pdf

²https://www.renyi.hu/~p_erdos/1960-09.pdf

³http://archive.ymsc.tsinghua.edu.cn/pacm_download/21/86-20100n_the_Minimum_Number_of_Convex_Quadrilaterals_in_Point_Sets_of_Given_Numbers_of_Points.pdf

v_2v_3 will achieve this. This process can be iterated as many times as necessary to obtain various cycles over x , from which it is possible to select one with the largest area.

Notice that for any polygon X , any other polygon with vertices chosen from X must also be convex and has a smaller area than X if and only if X is convex. We can then define a poset (Q, \geq) such that $A(x) \in Q$ if $x \in P$. Clearly Q is a subset of \mathbb{R} and is thus a total order, but right now I am only interested in the relations in Q of corresponding comparable elements in P , which is to say that we only need to consider $(A(x), A(y))$ for $x \subseteq y$ or vice versa. By the observation above, consider an ideal of P formed by taking some element at height i and all the chains up to height i with this element as their join. The i -sided polygon corresponding to this element is convex if and only if the aforementioned ideal is isomorphic to the subposet of Q formed by applying A to every element. For an element x_i exhibiting these properties, we call it a convex element of P .

From this we have a restatement of the happy ending problem: given the posets P and Q , as well as the function A , which itself has many properties that are still unknown to me, we want to find an integer $k(n)$ such that a convex element x_n must exist in P if the height of $P \geq k(n)$.

3 Lower bound for $f(N, n)$

Consider the poset P for a given power set of N elements. In general, we only consider cases where $N \geq k(n)$. Take the elements on the heights $H = N, k(n)$, and n . For each height, the number of elements would be $\binom{1, N}{k(n)}$ and $\binom{N}{n}$ respectively. The number of elements on $H = k(n)$ that some element on $H = n$ is a subset of is equal to the number of ways that one can choose another $k(n) - n$ points out of $N - n$ available ones.

Now consider the number of elements on $H = n$ which must be convex. Since every $k(n)$ -element subset must be comparable to a convex element on $H = n$, and every such element is comparable to exactly $\binom{N-n}{k(n)-n}$ unique $k(n)$ -element subsets, then minimally we would need

$$\left\lceil \binom{N}{k(n)} \div \binom{N-n}{k(n)-n} \right\rceil = \left\lceil \frac{N!(k(n)-n)!}{k(n)!(N-n)!} \right\rceil$$

convex elements on $H = n$ to satisfy the requirements.

By the result in the original paper where $k(4) = 5$, we have

$$f(N, 4) = \left\lceil \frac{N!}{120(N-4)!} \right\rceil$$

The growth rate of $f(N, 4)$ by this calculation is given as such:

$$\frac{f(N, 4)}{f(N-1, 4)} \approx \frac{N}{N-4}$$

for large enough N . This is the same result found in Hu et al. However, this formula gives $f(7, 4) = 7$, which is worse than the actual bound $f(7, 4) = 9$. The disparity is due to the fact that the minimal number of convex elements required does not match the actual graph in the plane, as the additional point added from N to $N + 1$ creates more convex elements than expected by our method. Additionally, the possible graphs of 7 points containing only 9 convex quadrilaterals must contain a convex pentagon.

Regardless of the actual value of $f(N, 4)$, the best bounds currently are $O(N^4)$. At $H = 5$, there are $\binom{N}{5}$ elements, which is $O(N^5)$. Given that the number of (x_4, x_5) pairs in P is $(N - 4)f(N, 4) = O(N^5)$, it is not possible to use the pigeonhole principle alone to prove the happy ending problem for case $n = 5$. A similar argument applies for every n .

4 Some sketchy guesses

Assuming that $k(n)$ and $f(N, n)$ are accurately formulated:

1. For some set of $k(n) - 1$ points which does not contain a convex polygon with n sides, the convex hull of the entire set is an $n - 1$ sided polygon. This is a **very** sketchy guess.
2. For $N \geq k(n)$, a set of N points which does not contain a convex polygon with n sides contains more than $f(N, n)$ convex $n - 1$ sided polygons.
3. $f(N, n)$ is $O(N^n)$.