

# Problem Set 1

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## Question 1

Response

## Question 1.32

Response

## Question 1.40

Response

## Question 1.41

For languages  $A$  and  $B$ , let the **perfect shuffle** of  $A$  and  $B$  be the language

$$\{w : w = a_1b_1 \cdots a_kb_k, \text{ where } a_1 \cdots a_k \in A, b_1 \cdots b_k \in B, \text{ each } a_i, b_i \in \Sigma\}$$

Show that the class of regular languages is closed under perfect shuffle.

## Response

Let  $D_1 = (Q_1, \Sigma, \delta_1, q_1, F_1)$ ,  $D_2 = (Q_2, \Sigma, \delta_2, q_2, F_2)$  be two DFA's that recognize  $A$  and  $B$  respectively. Now, let  $N = (Q, \Sigma, \delta, q_0, F)$  be the NFA recognizing **perfect shuffle** and be constructed as:

$$Q = Q_1 \times Q_2 \times \{odd, even\}$$

$$\Sigma = \Sigma$$

$$q_0 = \{(q_1, q_2, odd)\}$$

$$F = F_1 \times F_2 \times \{odd\}$$

$$\delta = \begin{cases} \{\delta_1(q, \sigma)\} & q \in Q_2 \setminus F_2, \sigma \in \Sigma \\ \{\delta_2(q, \sigma)\} & q \in Q_1 \setminus F_1, \sigma \in \Sigma \\ \emptyset & otherwise \end{cases}$$

## Question 1.62

Response

## Question 1.66

Response

## Question 1.69

Response