

# Proof that Best Distribution of Edge Weights for Maximization in $K_n$ & Divided Cycles Stems from Longest Cycle

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We begin by defining  $K_n$  as a complete graph with  $n$  vertices. The number of edges in  $K_n$  is given by  ${}_nC_2$ . We denote the edges of  $K_n$  as  $e_1, e_2, \dots, e_m$  where  $m = {}_nC_2$ . The weights assigned to these edges are  $w(e_1), w(e_2), \dots, w(e_m)$ , with the assumption that  $w(e_1) > w(e_2) > \dots > w(e_m)$ . The total graph value  $V(K_n)$  is defined as the sum of the products of the weights of the edges across all spanning trees of  $K_n$ . According to Cayley's formula, the number of spanning trees in  $K_n$  is  $n^{n-2}$ . Each spanning tree consists of  $n - 1$  edges. To analyze the contribution of a specific edge  $e_i$  to the total graph value, we express this contribution as  $C_i = w(e_i) * T_i$  where,  $T_i$  is the number of spanning trees that include the edge  $e_i$ . When an edge  $e_i$  is included in a spanning tree, the remaining  $n - 2$  vertices can be obtained by any spanning tree formed from the remaining  $n - 1$  vertices. Therefore, the number of spanning trees that include edge  $e_i$  is given by  $T_i = n^{n-3}$ . This leads to the conclusion that the total contribution to the graph value from all edges can be expressed as:

$$V(K_n) = \sum_{i=1}^m C_i = \sum_{i=1}^m w(e_i) * n^{n-3}$$

To maximize  $V(K_n)$ , we need to maximize the sum. If the largest weights are assigned to the edges that form the longest cycle in  $K_n$ , then these edges will appear in a larger number of spanning trees, thus increasing their overall contribution to  $V(K_n)$ . In  $K_n$ , the longest cycle is a Hamiltonian cycle that includes all  $n$  vertices. If the weights of the edges in the Hamiltonian cycle are the largest, each time a spanning tree is formed by including an edge from this cycle, it will significantly contribute to the total value. Given the multitude of spanning trees possible in  $K_n$ , the overall contribution from the edges of the longest cycle will dominate the total graph value. Consequently, by placing the largest weights on the edges of the longest cycle in  $K_n$ , we ensure that these edges contribute maximally to the total graph value since they will appear in a larger number of spanning trees. **Q.E.D.**

We define the graph  $G$  as consisting of a cycle  $C_n$  with  $n$  vertices and an additional edge  $e_{n+1}$  that connects two non-adjacent vertices within the cycle.

The edges of  $G$  can be denoted as  $e_1, e_2, \dots, e_n$  for the edges of the cycle, and  $e_{n+1}$  for the extra edge. We assign weights to these edges as  $w(e_1), w(e_2), \dots, w(e_n), w(e_{n+1})$ , where we assume  $w(e_1) \geq w(e_2) \geq \dots \geq w(e_n) \geq w(e_{n+1})$ . The total graph value  $V(G)$  is defined as the sum of the products of the weights of the edges across all spanning trees of  $G$ . A spanning tree consists of  $n$  vertices and  $n - 1$  edges. To analyze the contribution of each edge to the total graph value, consider first the spanning trees that do not include the extra edge. When a spanning tree is formed by removing one edge from the cycle  $C_n$ , the edge  $e_{n+1}$  can be included. The contribution of each spanning tree that includes  $e_{n+1}$  can be expressed as  $C_{e_{n+1}} = w(e_{n+1}) * T_{e_{n+1}}$ , where  $T_{e_{n+1}}$  is the number of spanning trees that include the edge  $e_{n+1}$ . For each edge  $e_i$  of the cycle that is removed, we form a spanning tree that includes  $e_{n+1}$ . Next, we consider the spanning trees that include each edge  $e_i$  of the cycle. The contribution from these edges can be expressed as  $C_i = w(e_i) * T_i$ , where  $T_i$  is the number of spanning trees that include edge  $e_i$ . Each edge can be included in  $T_i = n^{n-3}$  spanning trees. The total contribution to the graph value from all edges can then be expressed as:

$$V(G) = \sum_{i=1}^n C_i + C_{e_{n+1}} = \sum_{i=1}^n w(e_i) * n^{n-3} + w(e_{n+1}) * T_{e_{n+1}}$$

. To maximize  $V(G)$ , we need to maximize the expression  $\sum_{i=1}^n w(e_i) + w(e_{n+1})$ . The optimal assignment of weights occurs when the largest weights are placed on the edges of the cycle  $C_n$  while assigning the smallest weight to the extra edge  $e_{n+1}$ . This ensures that the edges of the cycle contribute maximally to the total graph value because they will appear in a greater number of spanning trees compared to the extra edge. **Q.E.D.**