

Machine-level Programming II: Control

'20H2

송인식

- Control: Condition codes
- Conditional branches
- Loops
- Switch Statements

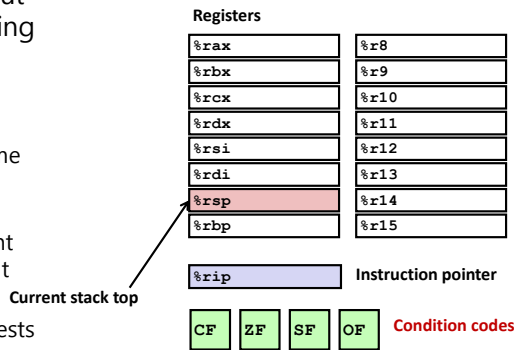
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Processor State (x86-64, Partial)

- Information about currently executing program

- Temporary data (`%rax`, ...)
- Location of runtime stack (`%rsp`)
- Location of current code control point (`%rip`, ...)
- Status of recent tests (`CF`, `ZF`, `SF`, `OF`)



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Condition Codes (Implicit Setting)

- Single bit registers
 - **CF** Carry Flag (for unsigned)
 - **SF** Sign Flag (for signed)
 - **ZF** Zero Flag
 - **OF** Overflow Flag (for signed)
- Implicitly set (think of it as side effect) by arithmetic operations

Example: `addq Src, Dest` ↔ `t = a+b`

 - CF set** if carry out from most significant bit (unsigned overflow)
 - ZF set** if `t == 0`
 - SF set** if `t < 0` (as signed)
 - OF set** if two's-complement (signed) overflow (`a > 0 && b > 0 && t < 0` || `a < 0 && b < 0 && t >= 0`)
- Not set by `leaq` instruction

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Condition Codes (Explicit Setting: Compare)

- Explicit Setting by Compare Instruction
 - `cmpq Src2, Src1`
 - `cmpq b, a` like computing `a-b` without setting destination
- **CF set** if carry out from most significant bit (used for unsigned comparisons)
- **ZF set** if `a == b`
- **SF set** if `(a-b) < 0` (as signed)
- **OF set** if two's-complement (signed) overflow (`a > 0 && b < 0 && (a-b) < 0` || `a < 0 && b > 0 && (a-b) > 0`)

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Condition Codes (Explicit Setting: Test)

- Explicit Setting by Test instruction
 - `testq Src2, Src1`
 - `testq b, a` like computing `a&b` without setting destination
- Sets condition codes based on value of `Src1` & `Src2`
- Useful to have one of the operands be a mask
- **ZF set** when `a&b == 0`
- **SF set** when `a&b < 0`

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Reading Condition Codes

- SetX Instructions

- Set low-order byte of destination to 0 or 1 based on combinations of condition codes
- Does not alter remaining 7 bytes

SetX	Condition	Description
sete	ZF	Equal / Zero
setne	~ZF	Not Equal / Not Zero
sets	SF	Negative
setns	~SF	Nonnegative
setg	~(SF^OF) & ~ZF	Greater (Signed)
setge	~(SF^OF)	Greater or Equal (Signed)
setl	(SF^OF)	Less (Signed)
setle	(SF^OF) ZF	Less or Equal (Signed)
seta	~CF & ~ZF	Above (unsigned)
setb	CF	Below (unsigned)

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x86-64 Integer Registers

%rax	%al	%r8	%r8b
%rbx	%bl	%r9	%r9b
%rcx	%cl	%r10	%r10b
%rdx	%dl	%r11	%r11b
%rsi	%sil	%r12	%r12b
%rdi	%dil	%r13	%r13b
%rsp	%spl	%r14	%r14b
%rbp	%bpl	%r15	%r15b

- Can reference low-order byte

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Reading Condition Codes (Cont.)

- SetX Instructions:

- Set single byte based on combination of condition codes

- One of addressable byte registers

- Does not alter remaining bytes
- Typically use `movzbl` to finish job
 - 32-bit instructions also set upper 32 bits to 0

```
int gt (long x, long y)
{
    return x > y;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

```
cmpq %rsi, %rdi # Compare x:y
setg %al         # Set when >
movzbl %al, %eax # Zero rest of %rax
ret
```

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Outline

- Control: Condition codes
- Conditional branches
- Loops
- Switch Statements

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Jumping

- jX Instructions

- Jump to different part of code depending on condition codes

jX	Condition	Description
jmp	1	Unconditional
je	ZF	Equal / Zero
jne	~ZF	Not Equal / Not Zero
js	SF	Negative
jns	~SF	Nonnegative
jg	~(SF^OF) & ~ZF	Greater (Signed)
jge	~(SF^OF)	Greater or Equal (Signed)
jl	(SF^OF)	Less (Signed)
jle	(SF^OF) ZF	Less or Equal (Signed)
ja	~CF & ~ZF	Above (unsigned)
jb	CF	Below (unsigned)

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Conditional Branch Example (Old Style)

- Generation

shark> gcc -Og -S -fno-if-conversion control.c

```
long absdiff
(long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

```
absdiff:
    cmpq    %rsi, %rdi # x:y
    jle     .L4
    movq    %rdi, %rax
    subq    %rsi, %rax
    ret
.L4:
    # x <= y
    movq    %rsi, %rax
    subq    %rdi, %rax
    ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

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Expressing with Goto Code

- C allows **goto** statement
- Jump to position designated by label

```
long absdiff
(long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

```
long absdiff_j
(long x, long y)
{
    long result;
    int ntest = x <= y;
    if (ntest) goto Else;
    result = x-y;
    goto Done;
Else:
    result = y-x;
Done:
    return result;
}
```

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General Conditional Expression Translation (Using Branches)

C Code

```
val = Test ? Then_Expr : Else_Expr;
```

```
val = x>y ? x-y : y-x;
```

Goto Version

```
ntest = !Test;
if (ntest) goto Else;
val = Then_Expr;
goto Done;
Else:
    val = Else_Expr;
Done:
    . . .
```

- Create separate code regions for then & else expressions
- Execute appropriate one

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Using Conditional Moves

- Conditional Move Instructions
 - Instruction supports: if (Test) Dest ← Src
 - Supported in post-1995 x86 processors
 - GCC tries to use them
 - But, only when known to be safe
- Why?
 - Branches are very disruptive to instruction flow through pipelines
 - Conditional moves do not require control transfer

C Code

```
val = Test
    ? Then_Expr
    : Else_Expr;
```

Goto Version

```
result = Then_Expr;
eval = Else_Expr;
nt = !Test;
if (nt) result = eval;
return result;
```

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Conditional Move Example

```
long absdiff
(long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

```
absdiff:
    movq    %rdi, %rax # x
    subq    %rsi, %rax # result = x-y
    movq    %rsi, %rdx
    subq    %rdi, %rdx # eval = y-x
    cmpq    %rsi, %rdi # x:y
    cmovle  %rdx, %rax # if <=, result = eval
    ret
```

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Bad Cases for Conditional Move

Expensive Computations

```
val = Test(x) ? Hard1(x) : Hard2(x);
```

- Both values get computed
- Only makes sense when computations are very simple

Risky Computations

```
val = p ? *p : 0;
```

- Both values get computed
- May have undesirable effects

Computations with side effects

```
val = x > 0 ? x*=7 : x+=3;
```

- Both values get computed
- Must be side-effect free

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Outline

- Control: Condition codes
- Conditional branches
- **Loops**
- Switch Statements

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"Do-While" Loop Example

C Code

```
long pcount_do
(unsigned long x) {
    long result = 0;
    do {
        result += x & 0x1;
        x >>= 1;
    } while (x);
    return result;
}
```

Goto Version

```
long pcount_goto
(unsigned long x) {
    long result = 0;
loop:
    result += x & 0x1;
    x >>= 1;
    if(x) goto loop;
    return result;
}
```

- Count number of 1's in argument x ("popcount")
- Use conditional branch to either continue looping or to exit loop

"Do-While" Loop Compilation

Goto Version

```
long pcount_goto
(unsigned long x) {
    long result = 0;
loop:
    result += x & 0x1;
    x >>= 1;
    if(x) goto loop;
    return result;
}
```

```
movl    $0, %eax        # result = 0
.L2:
movq    %rdi, %rdx
andl    $1, %edx
addq    %rdx, %rax       # result += t
shrq    %rdi             # x >>= 1
jne     .L2              # if (x) goto loop
rep; ret
```

Register	Use(s)
%rdi	Argument x
%rax	result

General "Do-While" Translation

C Code

```
do
    Body
while (Test);
```

Body:

```
Statement1;
Statement2;
...
Statementn;
}
```

Goto Version

```
loop:
    Body
    if (Test)
        goto loop
```

General "While" Translation #1

- "Jump-to-middle" translation
- Used with -Og

While version

```
while (Test)
    Body
```



Goto Version

```
goto test;
loop:
    Body
test:
    if (Test)
        goto loop;
done:
```

While Loop Example #1

C Code

```
long pcount_while
(unsigned long x) {
    long result = 0;
    while (x) {
        result += x & 0x1;
        x >>= 1;
    }
    return result;
}
```

Jump to Middle V

```
long pcount_goto_jtm
(unsigned long x) {
    long result = 0;
    goto test;
loop:
    result += x & 0x1;
    x >>= 1;
test:
    if(x) goto loop;
    return result;
}
```

- Compare to do-while version of function
- Initial goto starts loop at test

General "While" Translation #2

While version

```
while (Test)
    Body
```



Do-While Version

```
if (!Test)
    goto done;
do
    Body
while (Test);
done:
```



Goto Version

```
if (!Test)
    goto done;
loop:
    Body
    if (Test)
        goto loop;
done:
```

- "Do-while" conversion
- Used with -O1

While Loop Example #2

C Code

```
long pcount_while
(unsigned long x) {
    long result = 0;
    while (x) {
        result += x & 0x1;
        x >>= 1;
    }
    return result;
}
```

Do-While Version

```
long pcount_goto_dw
(unsigned long x) {
    long result = 0;
    if (!x) goto done;
loop:
    result += x & 0x1;
    x >>= 1;
    if (x) goto loop;
done:
    return result;
}
```

- Compare to do-while version of function
- Initial goto starts loop at test

"For" Loop Form

General Form

```
for (Init; Test; Update )
    Body
```

```
#define WSIZE 8*sizeof(int)
long pcount_for
(unsigned long x)
{
    size_t i;
    long result = 0;
    for (i = 0; i < WSIZE; i++)
    {
        unsigned bit =
            (x >> i) & 0x1;
        result += bit;
    }
    return result;
}
```

Init

```
i = 0
```

Test

```
i < WSIZE
```

Update

```
i++
```

Body

```
{
    unsigned bit =
        (x >> i) & 0x1;
    result += bit;
}
```

"For" Loop → While Loop

For Version

```
for (Init; Test; Update )
    Body
```

While Version

```
Init;
while (Test) {
    Body
    Update;
}
```

For-While Conversion

Init

```
i = 0
```

Test

```
i < WSIZE
```

Update

```
i++
```

Body

```
{
    unsigned bit =
        (x >> i) & 0x1;
    result += bit;
}
```

```
long pcount_for_while
(unsigned long x)
{
    size_t i;
    long result = 0;
    i = 0;
    while (i < WSIZE)
    {
        unsigned bit =
            (x >> i) & 0x1;
        result += bit;
        i++;
    }
    return result;
}
```

"For" Loop Do-While Conversion

Goto Version

C Code

```
long pcount_for
(unsigned long x)
{
    size_t i;
    long result = 0;
    for (i = 0; i < WSIZE; i++)
    {
        unsigned bit =
            (x >> i) & 0x1;
        result += bit;
    }
    return result;
}
```

```
long pcount_for_goto_dw
(unsigned long x) {
    size_t i;
    long result = 0;
    i = 0;
    if (!(i < WSIZE)) Init
        goto done; ! Test
loop:
    {
        unsigned bit =
            (x >> i) & 0x1; Body
        result += bit;
    }
    i++; Update
    if (i < WSIZE) Test
        goto loop;
done:
    return result;
}
```

- Initial test can be optimized away

Outline

- Control: Condition codes
- Conditional branches
- Loops
- Switch Statements

Switch Statement Example

```
long switch_eg
(long x, long y, long z)
{
    long w = 1;
    switch(x) {
        case 1:
            w = y*z;
            break;
        case 2:
            w = y/z;
            /* Fall Through */
        case 3:
            w += z;
            break;
        case 5:
        case 6:
            w -= z;
            break;
        default:
            w = 2;
    }
    return w;
}
```

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- Multiple case labels
 - Here: 5 & 6
- Fall through cases
 - Here: 2
- Missing cases
 - Here: 4

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Jump Table Structure

Switch Form

```
switch(x) {
    case val_0:
        Block 0
    case val_1:
        Block 1
    . . .
    case val_n-1:
        Block n-1
}
```

Translation (Extended C)

```
goto *JTab[x];
```

Jump Table

```
jtab: Targ0
      Targ1
      Targ2
      .
      .
      Targn-1
```

Jump Targets

```
Targ0: Code Block 0
Targ1: Code Block 1
Targ2: Code Block 2
.
.
Targn-1: Code Block n-1
```

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Switch Statement Example

```
long switch_eg(long x, long y, long z)
{
    long w = 1;
    switch(x) {
        . . .
    }
    return w;
}
```

Setup:

```
switch_eg:
    movq    %rdx, %rcx
    cmpq    $6, %rdi    # x:6
    ja      .L8
    jmp     *.L4(, %rdi, 8)
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

Note that **w** not initialized here

What range of values takes default?

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Switch Statement Example

```
long switch_eg(long x, long y, long z)
{
    long w = 1;
    switch(x) {
        . . .
    }
    return w;
}
```

Setup:

```
switch_eg:
    movq    %rdx, %rcx
    cmpq    $6, %rdi    # x:6
    ja      .L8          # Use default
    jmp     *.L4(, %rdi, 8) # goto *JTab[x]
```

Indirect jump

Jump table

```
.section .rodata
.align 8
.L4:
    .quad .L8 # x = 0
    .quad .L3 # x = 1
    .quad .L5 # x = 2
    .quad .L9 # x = 3
    .quad .L8 # x = 4
    .quad .L7 # x = 5
    .quad .L7 # x = 6
```

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Assembly Setup Explanation

- Table Structure
 - Each target requires 8 bytes
 - Base address at .L4
- Jumping
 - **Direct:** jmp .L8
 - Jump target is denoted by label .L8
 - **Indirect:** jmp *.L4(, %rdi, 8)
 - Start of jump table: .L4
 - Must scale by factor of 8 (addresses are 8 bytes)
 - Fetch target from effective Address .L4 + x*8
 - Only for 0 ≤ x ≤ 6

Jump table

```
.section .rodata
.align 8
.L4:
    .quad .L8 # x = 0
    .quad .L3 # x = 1
    .quad .L5 # x = 2
    .quad .L9 # x = 3
    .quad .L8 # x = 4
    .quad .L7 # x = 5
    .quad .L7 # x = 6
```

Jump Table

Jump table

```
.section .rodata
.align 8
.L4:
    .quad .L8 # x = 0
    .quad .L3 # x = 1
    .quad .L5 # x = 2
    .quad .L9 # x = 3
    .quad .L8 # x = 4
    .quad .L7 # x = 5
    .quad .L7 # x = 6
```

```
switch(x) {
    case 1: // .L3
        w = y*z;
        break;
    case 2: // .L5
        w = y/z;
        /* Fall Through */
    case 3: // .L9
        w += z;
        break;
    case 5:
    case 6: // .L7
        w -= z;
        break;
    default: // .L8
        w = 2;
}
```

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Code Blocks (x == 1)

```
switch(x) {
case 1:      // .L3
    w = y*z;
    break;
    . . .
}
```

```
.L3:
    movq    %rsi, %rax # y
    imulq   %rdx, %rax # y*z
    ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

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Handling Fall-Through

```
long w = 1;
. . .
switch(x) {
    . . .
case 2:
    w = y/z;
    /* Fall Through */
case 3:
    w += z;
    break;
    . . .
}
```

```
case 2:
    w = y/z;
    goto merge;
```

```
case 3:
    w = 1;
merge:
    w += z;
```

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Code Blocks (x == 2, x == 3)

```
long w = 1;
. . .
switch(x) {
    . . .
case 2:
    w = y/z;
    /* Fall Through */
case 3:
    w += z;
    break;
    . . .
}
```

```
.L5:                # Case 2
    movq    %rsi, %rax
    cqto
    idivq   %rcx     # y/z
    jmp     .L6      # goto merge
.L9:                # Case 3
    movl    $1, %eax # w = 1
.L6:                # merge:
    addq    %rcx, %rax # w += z
    ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

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Code Blocks (x == 5, x == 6, default)

```
switch(x) {
    . . .
case 5: // .L7
case 6: // .L7
    w -= z;
    break;
default: // .L8
    w = 2;
}
```

```
.L7:                # Case 5,6
    movl    $1, %eax # w = 1
    subq    %rdx, %rax # w -= z
    ret
.L8:                # Default:
    movl    $2, %eax # 2
    ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

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Summarizing

- C Control
 - if-then-else
 - do-while
 - while, for
 - switch
- Assembler Control
 - Conditional jump
 - Conditional move
 - Indirect jump (via jump tables)
 - Compiler generates code sequence to implement more complex control
- Standard Techniques
 - Loops converted to do-while or jump-to-middle form
 - Large switch statements use jump tables
 - Sparse switch statements may use decision trees (if-elseif-elseif-else)

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Questions?

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