

### 1.3 Algorithms and Convergence

**Algorithm:** procedure that unambiguously describes a finite sequence of steps in a specified order (can typically be written as pseudocode  $\Rightarrow$  no specific language required)

- **stable:** small changes to input  $\rightarrow$  small changes to output
- **conditionally stable:** stable for some input
- **unstable:** not stable for any input

Let  $E_0 > 0$  be error at initial step

$E_n$  be error at  $n^{\text{th}}$  step

Algorithm has:

- **Linear error growth:** if  $E_n \approx c \cdot n \cdot E_0$ ,  $c > 0$  constant
- **Exponential error growth:** if  $E_n \approx c^n \cdot E_0$ ,  $c > 1$  constant

**Rmk.** Linear error growth  $\rightarrow$  stable

Exponential error growth  $\rightarrow$  unstable

### Convergence Rate of Sequences

Let  $\{\alpha_n\}_{n=1}^{\infty}$  be a sequence such that  $\alpha_n \rightarrow \alpha$  as  $n \rightarrow \infty$ .

**Q.** How quickly is  $\alpha_n$  approaching  $\alpha$ ?

Use a second known sequence  $\{\beta_n\}$  to describe convergence behavior of  $\{\alpha_n\}$ .

**Def.** Let  $\alpha_n \rightarrow \alpha$  and  $\beta_n \rightarrow \beta$  as  $n \rightarrow \infty$ .

If there exists  $K > 0$  and integer  $n_0$  such that  $|\alpha_n - \alpha| \leq K |\beta_n - \beta|$  for all  $n \geq n_0$ ,

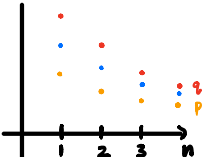
then  $a_n$  converges to  $\alpha$  with rate/order of  $O(\beta_n)$ , written as  $a_n = \alpha + O(\beta_n)$

Rmk. 1:  $\beta_n$  is usually chosen as  $n^{-p}$ ,  $p > 0$

(generally interested in largest possible  $p$  st  $a_n = \alpha + O(n^{-p})$ )

Rmk. 2: If  $0 < q < p$  and  $a_n = \alpha + O(n^{-p})$ , then  $a_n = \alpha + O(n^{-q})$

e.g.  $|a_n - \alpha| \leq K \underbrace{\left| \frac{1}{n^5} \right|}_{p=5} \leq K \underbrace{\left| \frac{1}{n^2} \right|}_{q=2}$



Ex. Let  $n \geq 1$ .

$$a_n = \frac{n+1}{n^2} \quad (\alpha = 0)$$

$$|a_n - 0| = \left| \frac{n+1}{n^2} \right| = \left| \frac{1}{n} + \frac{1}{n^2} \right| \leq 2 \underbrace{\left| \frac{1}{n} \right|}_{K|\beta_n|}$$

$$\rightarrow a_n = 0 + O\left(\frac{1}{n}\right)$$

$$\hat{a}_n = \frac{n+1}{n^3} \quad (\alpha = 0)$$

$$|\hat{a}_n - 0| = \left| \frac{n+1}{n^3} \right| = \left| \frac{1}{n^2} + \frac{1}{n^3} \right| \leq 2 \left| \frac{1}{n^2} \right|$$

$$\rightarrow \hat{a}_n = 0 + O\left(\frac{1}{n^2}\right)$$

$\{\hat{a}_n\}$  converges faster!

Similarly for functions:

$$\text{Let } \lim_{h \rightarrow 0} F(h) = L$$

Q. How fast is  $F$  approaching  $L$ ? (as  $h \rightarrow 0$ )

Use a known function  $G(h)$ , where  $\lim_{h \rightarrow 0} G(h) = 0$ .

Def. Let  $\lim_{h \rightarrow 0} F(h) = L$ ,  $\lim_{h \rightarrow 0} G(h) = 0$ . If there exists  $K > 0$ ,  $h_0 > 0$  such that

$$|F(h) - L| \leq K |G(h)| \text{ for } h < h_0, \text{ then we write } F(h) = L + O(G(h)).$$

**Rmk.**  $G(h)$  is usually chosen as  $h^p$ ,  $p > 0$

(generally interested in  $\max \{p: F(h) = L + O(h^p)\}$ )

**Ex.** Analyze the convergence rate of  $F(h) = \sin(h) - h\cos(h)$  as  $h \rightarrow 0$  ( $L=0$ )

Note: by Taylor's Thm,

$$\sin(h) = h - \frac{h^3}{6} \cos(\xi) \quad \text{where } 0 \leq \xi \leq h$$

$$\cos(h) = 1 - \frac{h^2}{2} \cos(\eta) \quad \text{where } 0 \leq \eta \leq h$$

$$\begin{aligned} |\sin(h) - h\cos(h)| &= \left| h - \frac{h^3}{6} \cos(\xi) - h + \frac{h^3}{2} \cos(\eta) \right| \\ &\leq \underbrace{\left| \frac{h^3}{6} \cos(\xi) \right|}_{\leq 1} + \underbrace{\left| \frac{h^3}{2} \cos(\eta) \right|}_{\leq 1} \\ &\leq \left( \frac{1}{6} + \frac{1}{2} \right) |h^3| \end{aligned}$$

$$\rightarrow \sin(h) - h\cos(h) = 0 + O(h^3)$$