

WATSUP: Web Authentication without Sending or Storing User Passwords

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Abstract

User passwords are regularly compromised through failures of trusted web servers. Most approaches to solve password problems involve improving server security, improving client practices, or replacing passwords altogether. However, these approaches are not always practical or achievable. In many cases, the undetectable negligence of server owners impedes security improvements. We propose Web Authentication without Transmitting or Storing User Passwords (WATSUP), a new application-layer web protocol to replace standard website authentication, without requiring trust in any server or third party. WATSUP is designed to be auditable, in that a client can always detect if the server incorrectly implements the protocol. WATSUP is secure against eavesdropping and spoofing, even if the communication or server database is compromised. We provide a proof-of-concept implementation of WATSUP, and we discuss the weaknesses and necessary improvements prior to deployment.

1 Introduction

Internet users must entrust private information to many different companies, but online security is often overlooked. Even large and technically advanced companies have lost sensitive information to data breaches. For example, Yahoo suffered two major breaches in the past few years. First, in 2013 attackers perpetrated the largest recorded data breach when they stole data from roughly one billion Yahoo accounts [19]. Then in 2014, a separate attack compromised roughly 500 million Yahoo accounts

[8]. Neither of these breaches were discovered until 2016, meaning sensitive user information such as hashed passwords and security questions were stolen years before any user became aware.

These types of breaches are not unique to Yahoo [13, 14]. They exemplify a major issue with internet security: users must entrust their personal information to servers and companies that are not transparent and often fail to implement standard security methods. Most companies do not publish their full security practices, and even users who are responsible or knowledgeable about online security have no means to verify that their data is being handled correctly. Furthermore, the companies themselves are often unable to detect breaches promptly. As users register for an increasing number of services, these risks grow in magnitude.

In spite of the difficulties and flaws of passwords, such as reused, weak, and difficult-to-remember passwords, as well as constant breaches of improperly stored passwords, passwords are still unavoidable for users. Despite a push from both users and security experts, alternatives rarely catch on, especially in web authentication [12]. Given that passwords have not been eliminated, we focus on improving the security of user passwords rather than eliminating them entirely. We propose Web Authentication without Transmitting or Storing User Passwords (WATSUP), a new application-layer protocol that provides users with an auditable, consistent, and easy way to log into web services without trusting any service to properly handle their login credentials. In this paper, we first discuss security risks, their solutions and drawbacks, and related work. Next, we propose the WATSUP protocol; then we discuss a proof-of-

concept implementation; and finally, we discuss the advantages and disadvantages of WATSUP and propose future work.

2 Background

2.1 Risks of using passwords

Verifying a user's identity is an important challenge in internet security. Passwords, despite their many drawbacks, are a critical and ubiquitous solution to the problem. Today, anyone who shops, banks, communicates, or socializes online creates multiple user accounts with a number of websites. But password usage has a number of risks, and existing solutions have disadvantages. The WATSUP protocol is designed to eliminate or mitigate the common risks mentioned below.

2.1.1 Eavesdropping

Solution: *HTTPS*

Complication: *Not always supported*

A common risk is that a user's password is captured in transit as plaintext. The most popular solution to eavesdropping is to use HTTPS, which tunnels an HTTP connection through a TLS connection. Unfortunately, while HTTPS is supported by most large companies, it is still not universally used [6]. In addition, many users may not realize that they should not use websites that do not support HTTPS.

2.1.2 Data breaches

Solution: *Salting and hashing*

Complication: *Not always implemented*

A preventable, hidden risk users face is that their passwords are improperly stored on a company's web server. As discussed in the introduction, data breaches have happened or been discovered as recently as last year, often at staggering scales. The most common methods to mitigate the effects of losing passwords is to

salt and hash them. A cryptographic hash function is a one-way function for which it is easy to verify that an input is correct, but hard to extract the input from the output. A company should hash a new password before saving it in a database; to authenticate a user, it hashes the candidate password and compares it to the hashed password on record. Salting passwords adds fixed random information to each password. Combining these methods helps prevent rainbow table attacks in which precomputed tables of common passwords hashed with typical hashing algorithms are used to decode hashed passwords. They also help prevent an attack reusing a compromised password on another website. Unfortunately, many popular websites do not properly handle user data. For example, in its official statement, Yahoo only confirmed that the "vast majority" of passwords had been securely hashed before they were stolen in the 2013 attack [16], implying some (out of a billion) were not. And in a 2012 LinkedIn data breach, LinkedIn failed to salt users' passwords and only hashed them with the SHA-1 algorithm [10, 17], which has been known to be insufficiently secure as early as 2005 [20].

2.1.3 Phishing

Solution: *Client-side hashing on hostname*

Complication: *Susceptible to rainbow table attacks*

Another common means of compromising a user's account is called "phishing." In phishing, a user is asked to submit their username and password to a malicious website that poses as a legitimate one, typically through an email that encourages the user to log in to a domain posing as the legitimate service. One solution to this problem is to hash the password with the website's hostname as a salt before sending it. If this is done consistently, phishing is prevented because the malicious website will have a different hostname. This allows for improved portability, since no secret data is stored, and no trusted third party is required [11]. However this method is still vulnerable to weak pass-

words and poor server salting. For example, an attacker who has compromised a database of passwords that have been salted with the hostname and then hashed can still create a rainbow table for that website.

2.1.4 Replay attacks

Solution: *Nonces*

Complication: *Not typically implemented*

Another common attack vector is called a "replay attack." An adversary intercepts the user's encrypted authentication credentials and re-transmits them without modification in a subsequent login to masquerade as the user. In a replay attack, the adversary does not need to decrypt the data to pose as the user. To prevent this, each encrypted communication must include a nonce, which is a number that can only ever be used once. This scrambles the encrypted data so that a re-transmitted sequence will be invalid. HTTPS includes this feature.

2.1.5 Password reuse

Solution: *Strong, unique passwords*

Complication: *Cognitive overhead*

Finally, passwords are insecure because remembering many strong passwords is hard, and users have a tendency to reuse passwords. Many systems have been designed to lessen the cognitive load for users. For example, password managers are popular tools to generate and save unique user passwords, reducing password reuse without increasing a user's cognitive load. However, these password managers are frequently insecure [15]. Importantly, they often either rely on a trusted third-party or are not portable. Additionally, a password manager may not protect against phishing or replay attacks.

2.2 Related work

One solution to many of the above problems is Password Multiplier, which derives a site-specific password from a base password and

the site's hostname. This uses SHA-1 as a key derivation function, and the hostname as a salt. Key derivation functions take the user's base password and produce multiple distinct high entropy keys by using a distinct salt. If one derived key is compromised, the base password and all other keys remain secure [11]. This does not protect against replay attacks on the same site. In addition, if everyone were to use this and the server failed to adequately salt and hash the derived keys, then everyone would be using the same salt. In this scenario, a rainbow table or dictionary attack on a per-hostname basis would be reasonable, particularly against large databases.

Another similar solution is a client-server web authentication protocol known as Secure Quick Reliable Login (SQRL). It performs web authentication on a "something you have" model. SQRL uses a randomly generated master password to provide strong, derived logins on a per-site basis, almost completely eliminating user passwords and making phishing difficult. On a login request, a link or QR code is provided for input to the SQRL app, which may be on a different device. The SQRL app then completes the login. As a side effect, this also protects against phishing [9]. However, SQRL has a few issues. First, it requires its own protocol, and does not run over HTTP or HTTPS. This is a significant barrier to deployment. Second, it requires the rendering of QR codes. Finally, the code is implemented in assembly. While this prevents a compiler from eliminating security-critical code, it makes incremental deployment difficult.

2.3 Implications

We argue that any solution to the web authentication problem must resolve the following issues:

- Users should not have to trust servers.
- The solution should be portable. It should not require any data be stored or transmitted before it can be used on a new device. This means using passwords.

- Users should be protected against replay and reuse attacks, even without HTTPS.
- The solution should be as simple as possible for servers to implement, so that any developer can safely add it to their server, without risking user security.
- The solution should allow for phased deployment by not interfering with standard login.

We designed WATSUP to address these issues.

3 Protocol

3.1 Design

In the standard web login process, a user inputs their cleartext password into a form and submits it to the server, ideally via HTTPS. With the WATSUP protocol, a user still inputs their cleartext password, but it is never transmitted or stored. Instead, it is salted with the username and hostname, hashed, and then passed to a key derivation function to produce a seed for a pseudorandom number generator (PRNG). While this information is deterministically generated, it still serves a purpose similar to that of a salt by preventing against rainbow table and dictionary attacks. The PRNG is used to generate an asymmetric key pair. Next, the client requests a nonce from the server; the server generates the nonce and encrypts it with the user's public key. This public key was provided by the user during registration. Finally, the client receives the nonce, decrypts it with the deterministically generated private key, and submits the nonce for final authentication.

This design has a few benefits. It is easy to use and it is auditable, meaning a client can always detect if the server incorrectly implements the protocol. If a browser and server were to implement this protocol, users would not need to do anything differently. In addition, the method addresses the attacks mentioned previously.

- **Eavesdropping:** The user's password is never sent.

- **Data breaches:** Servers never see passwords and therefore cannot store them. Salting on the username protects against site-wide dictionary and rainbow table attacks.
- **Phishing:** The protocol makes phishing more difficult in two ways: first, we use different keys for different hostnames; and second, we never provide the server with a reusable key.
- **Replay attack:** Since nonces are good for only one use, replay attacks are prevented.
- **Password reuse:** The protocol primarily mitigates the effects of password reuse by salting with the hostname. While we do not recommend that users rely on a single password, we are realistic in our assumption that some people will reuse passwords.

3.2 Specification

Here we describe our procedure for generating an asymmetric key pair, which is the main cryptographic basis for the security of the WATSUP Protocol. We use the following definitions:

- `hash(x)` is the hash of `x`.
- `|` is the concatenation operator.
- `KDF(key, salt, iterations, hash_function)` derives an array of bits, with the specified hash function, salt, number of iterations, and key.
- `seed_PRNG(bit_array)` provides a cryptographic PRNG seeded from `bit_array`.
- `key_gen(prng)` provides an asymmetric key pair generated with the provided PRNG.

The procedure is:

```
salt = hash(hash(hostname)
              | hash(username))
bits[size] = KDF(root_pass, salt,
```

```

        iterations,
        hash_function)
prng = seed_PRNG(bits)
key_pair = key_gen(prng)

```

Using a high iteration value as input for KDF is called "key stretching" and makes the base password more difficult to crack.

3.3 Proof-of-concept

The proof-of-concept implementation has two components, a back-end web server and an browser extension. The web server is a stand-in for any web application that implements the WATSUP protocol. The browser extension sandboxes the login form from the web page and implements the client-side of the WATSUP protocol. The code for both the server and client are available here: <https://github.com/watsup-protocol>.

3.3.1 Server

The back-end is written in Python 2.7 and uses the Flask web framework. The server is an instance of Gunicorn, a Python WSGI HTTP server. The server runs on a Heroku instance running Ubuntu 14.04.

To generate the nonce, the server uses Python's `random` library function `SystemRandom` which is designed for cryptographic use. The function generates random numbers from information seeded by the operating system. Since the server runs on a UNIX-like operating system, `SystemRandom` queries `/dev/urandom` [4, 5].

To generate the RSA key pair the servers uses `cryptography`, an open-source Python package developed by the Python Cryptographic Authority [3]. The package primarily delegates to "backends" such as `OpenSSL` and `CommonCrypto` for the actual for platform-specific cryptographic algorithm implementations. The server uses the SHA-1 hash function for optimal asymmetric encryption padding; SHA-1 is acceptable here because adding padding does not hide sensitive information.

3.3.2 Client

Ideally a fully deployed implementation of the WATSUP protocol would be integrated into browsers by their vendors. However, for our proof-of-concept implementation, the client is implemented as a Google Chrome extension. It is written in JavaScript, HTML, and CSS.

Critically, Chrome sandboxes all browser extensions to prevent malicious websites from accessing privileged, extension-only operations [1]. The WATSUP client utilizes this security policy to keep user passwords safe. No other JavaScript program can access the extension's state, which contains references to the user's plaintext password.

For the protocol, the extension reads the username and password from their input fields and the hostname from Chrome's Tabs API [2]. The extension follows the key pair generation procedure defined in Specifications. We use SHA-256 as our hash function. We use PBKDF2 as our KDF, with 1000 iterations and SHA256 as the KDF hash function. We use ARC4 as our PRNG. We use `Cryptico`'s `generateRSAKey` to create a keypair [18].

3.4 Deployment

Deployment of new technologies and protocols to the internet can be a complicated process, and adoption tends to be slow and spotty. We believe the first step would be to recruit a major browser to implement client-side WATSUP, while providing open source browser plugins for all major browsers. Then, we could recruit some major websites and web frameworks to adopt WATSUP. Due to the minimal work required to implement server-side WATSUP and cross-compatibility with existing login services, any web service should be able to support WATSUP with ease.

4 Discussion

4.1 Weaknesses

The WATSUP protocol as described in this paper still suffers from some existing security issues. Current common implementations of asymmetric encryption schemes are vulnerable to offline attacks [7]. Additionally, if the user chooses a weak base password, the base password may be susceptible to offline dictionary attacks. Both of these attacks are only possible if an attacker obtains a nonce encrypted with a client's public key, or the attacker obtains the public key itself. Even with alternate asymmetric encryption algorithms, the attacker can still perform an offline dictionary attack if they also capture the plaintext nonce; the attack would perform a standard dictionary attack but then generate a public key for each candidate password using the WATSUP's protocol. If any public key decrypts the encrypted nonce and results in the unencrypted nonce, the attack knows the correct password. Alternately, with current encryption schemes, man-in-the-middle attacks are still possible with the WATSUP protocol. These weaknesses can be eliminated by using HTTPS, but the user still has some guarantees using HTTP.

Another weakness is a denial-of-service attack against the login protocol. An attacker could rapidly request nonces for a target user. Depending on the server implementation, a request may invalidate previous nonces, so the target may be unable to log in.

The WATSUP protocol also does not protect a user from themselves. In order for the WATSUP protocol to provide portable access to online content, a user's private key is reproducible using the user's username and password for a given host. Therefore, if a user loses their password through some social engineering attack for example, their accounts will still be vulnerable. This also means that online dictionary attacks are still a threat if users employ weak base passwords. We do not propose WATSUP as a replacement for good password hygiene. Users should still use good practices, such as different

passwords on each site and using high entropy passwords.

4.2 Future Work

Before WATSUP is ready for public use, it needs to be improved and thoroughly vetted by security experts. A major concern is that the protocol relies on the server to produce a cryptographic nonce. Our threat model does not allow for trust of the server. Therefore, we propose using POSIX time in milliseconds as a nonce. As long as the round trip time between client and server is greater than one millisecond, and the server is correctly providing the time, this nonce will be unique, and hence adequate for cryptographic use. This allows the client to audit the server's compliance.

Another direction that could be explored is that of mutual authentication based on a trust-on-first-use system. This would provide the user with server authenticity guarantees, which could be critical in the event of HTTPS public key infrastructure compromise. To solve this, one model is that the user saves the server's public key, but this means the user would need to save the server's public key on each of their devices. Another model is that the user signs the server's public key, which the server saves and reissues to the user on each connection, but this makes revocation in the event of server compromise difficult or impossible. In either model, the server would perform authentication similar to the client authentication.

4.3 Conclusion

The WATSUP protocol provides web authentication with guarantees for users and reduces the exposure of their passwords. Because the protocol never stores or transmits passwords, users are guaranteed that their cleartext passwords will not be stolen during transit or authentication or from the server. Perhaps most importantly, this protocol does not require a user to trust any server.

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