

### Problem 1

We define an  $\omega$ -word  $u$  as *universal* if, and only if, for any  $i \in \mathbb{N}$ , there exists a  $j > i$  such that  $u[i] \neq u[j]$ . Intuitively, this means that  $u$  contains an infinite subsequence of the form  $(ab)^\infty$ . Note that if  $u \in \Sigma^\omega$  is universal, then for any  $v \in \Sigma^\omega$ , it holds that  $v \sqsubseteq u$ . This is because we can remove finitely or infinitely many letters from  $u$  to obtain  $(ab)^\infty$ , and for each  $i$ -th pair of consecutive  $ab$  occurrences, we can remove either  $a$  or  $b$ , depending on the value of  $v[i]$ .

If an  $\omega$ -word  $u$  is not universal, then either  $\#_a(u)$  or  $\#_b(u)$  is finite. In the first case,  $u$  is of the form  $vb^\infty$ , and we will refer to such an  $\omega$ -word as *a-long*. In the second case,  $u$  is of the form  $va^\infty$ , and we will call it *b-long*. In both cases,  $v$  is a word of finite length.

Let  $u_1, u_2, \dots$  be any infinite sequence of  $\omega$ -words. We will prove that there exist indices  $i < j$  such that  $u_i \sqsubseteq u_j$ . If there are at least two universal  $\omega$ -words in this sequence, the proof is straightforward since  $u_1$  can embed into the second universal word in the sequence. Otherwise, either there are infinitely many *a-long* words in the sequence, or there are infinitely many *b-long* words. Without loss of generality, assume there are infinitely many *a-long* words in the sequence. Let  $k_1, k_2, \dots$  be the indices of the *a-long* words. Denote  $u_{k_l} = u'_{k_l}a^\infty$ , where  $u'_{k_l}$  is finite. By Higman's lemma, there exist indices  $i < j$  such that  $u'_{k_i}$  is a substring of  $u'_{k_j}$ , concluding the proof as  $u_{k_i} \sqsubseteq u_{k_j}$ .