PD109-1 hw5-1

Given a two dimension array map, which store the number of people living on location (x, y)

Given (n, m) as the boundary of (x, y) in the map

Given r as the range of coverage of a built hospital

// main function

int findMaxCover($n,\ m,\ r,\ map$)

initialize the number of maximum cover as $maxCover$ to be 0

for u from 0 to n

for v from 0 to m

set $numCover$ = cover($n,\ m,\ u,\ v,\ r,\ map$)

if $numCover > maxCover$

update $maxCover$ to be $numCover$

return $maxCover$

For the function cover(), we have two algorithms as following:

Algorithm 1

int cover($n,\ m,\ u,\ v,\ r,\ map$)

initialize the number of people covered by the hospital on (u, v) as $numOfCovered$ to be 0

for i from 0 to n

for j from 0 to m

set $distance = |i - u| + |j - v|$

if $distance \leq r$

update $numOfCovered$ to $numOfCovered + map[i][j]$

return $numOfCovered$

Algorithm 2

int cover($n,\ m,\ u,\ v,\ r,\ map$)

initialize the number of people covered by the hospital on (u, v) as $numOfCovered$ to be 0

for i from $u - r$ to $u + r$

set $yRange$ as $r - |i - u|$

for j from $v - yRange$ to $v + yRange$

update $numOfCovered$ to $numOfCovered + map[i][j]$

return $numOfCovered$

In algorithm 1, for each location, we have to compute $(n + 1) * (m + 1)$ times to compute the number of covered people. In algorithm 2, if $r = 1$, then we have to compute 1+3+1 times; if $r = 2$, then we have to compute 1 + 3 + 5 + 3 + 1 times. Thus, for a given range r, we have to compute $2\sum_{i=1}^{r}(2i - 1) + (2r + 1) = 2r(r + 1) + 1$ times. Theoretically, we can formulate the inequality as following to find the condition such that one of the algorithm is more efficient than the other one:

$$2r(r + 1) + 1 \leq (n + 1)(m + 1)$$

However, in practice, we only need to consider those villages inside of the map. When we implement these two algorithms, obviously the range to compute in algorithm 2 is less or equal than the range to compute in algorithm 1. Thus, algorithm 2 is more efficient.