OR Homework 1

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-0, -0.5

1 Problem 1

1.1 (a)

Set: Not italics, use \mathrm{}. -0.5

 $I \in \{1, \dots, 10\}, J \in \{\underline{S, L, H, M, O}\}.$

Variables:

 $x_i = 1$, if i is chosen, else $x_i = 0$. $\forall i \in I$.

 $y_{i,j} = 1$, if i is chosen as j, else $y_{i,j} = 0$. $\forall i \in I, \forall j \in J$.

 $w_{i,j}$ is a variable to linearize $x_i, y_{i,j}$. $w_j = \{0, 1\}. \forall i \in I, \forall j \in J$.

Constants:

 $z_{i,j}=1$, if i can be selected as j, else $z_{i,j}=0$. $\forall i\in I,\, \forall j\in J.$

 s_i, b_i, r_i, p_i are the rate of spiking, blocking, receiving, and passing abilities of player i. $\forall i \in I$.

Why don't you merely use y to represent player i's state? Variable w is redundant and may lead to some problem.

$$\begin{aligned} & \max \quad \sum_{i \in I} w_{i,H} s_i + \sum_{i \in I} w_{i,M} s_i + \sum_{i \in I} w_{i,O} s_i + 3 \sum_{i \in I} w_{i,S} p_i \\ & \text{s.t.} \quad w_{i,H} \leq x_i, w_{i,H} \leq y_{i,H} \quad \forall i \in I \\ & w_{i,M} \leq x_i, w_{i,M} \leq y_{i,M} \quad \forall i \in I \\ & w_{i,O} \leq x_i, w_{i,O} \leq y_{i,O} \quad \forall i \in I \\ & w_{i,S} \leq x_i, w_{i,S} \leq y_{i,S} \quad \forall i \in I \\ & w_{i,L} \leq x_i, w_{i,L} \leq y_{i,L} \quad \forall i \in I \\ & \sum_{i \in I} w_{i,S} = 1, \sum_{i \in I} w_{i,L} = 1, \sum_{i \in I} w_{i,H} = 2, \sum_{i \in I} w_{i,M} = 2, \sum_{i \in I} w_{i,O} = 1 \\ & \sum_{i \in I} x_i z_{i,S} \geq 2, \sum_{i \in I} x_i z_{i,O} \geq 3 \\ & 1 - x_1 \geq x_3 \\ & x_1 + x_3 + x_4 + x_6 + x_8 \geq 2 \\ & \sum_{j \in J} y_{i,j} \geq x_i \quad \forall i \in I \\ & 1 \geq \sum_{j \in J} y_{i,j} \quad \forall i \in I \\ & 1 \geq \sum_{j \in J} y_{i,j} \quad \forall i \in I, \forall j \in J \\ & w_{i,j}, x_i, y_{i,j}, z_{i,j} \in \{0,1\} \quad \forall i \in I, \forall j \in J \end{aligned} \end{aligned}$$

1.2 (b)

We only need to add two more constrains:

s.t.
$$\sum_{i \in I} w_{i,H} r_i + \sum_{i \in I} w_{i,O} r_i + \sum_{i \in I} w_{i,L} r_i \ge 4 \times 3.5$$
$$\sum_{i \in I} w_{i,M} b_i + \sum_{i \in I} w_{i,I} b_i \ge 3 \times 3.5$$
 (2)

2 Problem 2

2.1 (a)

Set:

$$I \in \{1, \dots, 5\}, J \in \{1, \dots, 8\}.$$

10 groups $G_{u,v}$ as the groups that students work on day u and day v. $\forall u, v \in I, u \neq v$. w_i as the maximum students needed in day i. $\forall i \in I$.

min
$$\sum_{u=1}^{4} \sum_{v=u+1}^{5} G_{u,v}$$
s.t.
$$\sum_{v \in I} G_{i,v} \ge w_i \quad \forall i \in I, v \ne i$$

$$w_i \ge D_{i,j} \quad \forall i \in I, j \in J$$

$$G_{u,v} \ge 0 \quad \forall u, v \in I, u \ne v$$

$$(3)$$

2.2 (b)

Set:

$$I \in \{1, \dots, 10\}.$$

define interval 1 to 10 as:

slot 1 to 4: 1, 2, 3, 4, 5 (day 1 to 5)

slot 5 to 8: 6, 7, 8, 9, 10 (day 1 to 5)

80 groups $G_{w,x,y,z}$ as the groups that students work on interval $w,x,y,z, \forall w,x,y,z \in I, w \neq x \neq y \neq z$.

 m_i as the maximum students needed in interval $i, \forall i \in I$.

min
$$\sum_{w=1}^{7} \sum_{x=w+1}^{8} \sum_{y=x+1}^{9} \sum_{z=y+1}^{10} G_{w,x,y,z}$$
s.t.
$$\sum_{x=w+1}^{8} \sum_{y=x+1}^{9} \sum_{z=y+1}^{10} G_{w,x,y,z} \ge m_w \quad \forall w \in I, w \ne x \ne y \ne z$$

$$m_w \ge D_{w,j} \quad \forall w \in \{1, ..., 5\}, j \in \{1, ..., 4\}$$

$$m_w \ge D_{w-5,j} \quad \forall w \in \{6, ..., 10\}, j \in \{5, ..., 8\}$$

$$G_{w,x,y,z} \ge 0 \quad \forall w, x, y, z \in I, w \ne x \ne y \ne z$$

$$(4)$$

3 Problem 3

Most notations are defined in the problem description, so we only define the ones that is not defined in it below.

Set:

 $I \in \{0,, 40\}$ as each CSR's number

 $x_{i,d,j}=1$ if CSR i is assigned to shift j in day d, else $x_{i,d,j}=0$. $\forall i\in I, \forall d\in D, \forall j\in J$.

 $y_{d,t}$ is the number of CSR shortage in period t of day t. $\forall d \in D, \forall j \in J$.

should be d (-0.5) should be t in T

$$\begin{array}{ll} \min & \sum_{d \in D} \sum_{t \in T} y_{d,t} \\ \mathrm{s.t.} & \sum_{j \in J} x_{i,d,j} = 1 \quad \forall i \in I, d \in D \\ & y_{d,t} \geq C_{t,d} - \sum_{j \in J} A_{j,t} \sum_{i \in I} x_{i,d,j} \quad \forall d \in D, \forall t \in T \\ & y_{d,t} \geq 0 \quad \forall d \in D, t \in T \\ & \text{ (-1) the predetermined number of days off per month may vary in each month.} \\ & \sum_{d \in D} x_{i,d,0} = \$ \quad \forall i \in I \quad (-1) \text{ should sum all the shifts belong to type leave} \\ & \sum_{d \in D} x_{i,d,0} = \$ \quad \forall i \in I \quad (-1) \text{ should sum all the shifts belong to type leave} \\ & \sum_{d \in D} x_{i,d,0} = \$ \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{j \in J} x_{i,d,j} B_{j,night} \leq 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{j \in J} x_{i,d,0} = 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{j \in J} x_{i,d,0} = 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{j \in J} x_{i,d,0} = 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{j \in J} x_{i,d,0} = 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{j \in J} x_{i,d,0} = 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{j \in J} x_{i,d,0} = 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{j \in J} x_{i,d,0} = 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{j \in J} x_{i,d,0} = 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{j \in J} x_{i,d,0} = 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{j \in J} x_{i,d,0} = 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{j \in J} x_{i,d,0} = 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{j \in J} x_{i,d,0} = 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{j \in J} x_{i,d,0} = 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{j \in J} x_{i,d,0} = 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{j \in J} x_{i,d,0} = 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{j \in J} x_{i,d,0} = 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{d \in E} \sum_{d \in E} x_{i,d,0} = 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} \sum_{d \in E} x_{i,d,0} = 1 \quad \forall i \in I, \forall k \in D - \{end \quad 6 \quad days\} \\ & \sum_{d \in E} x_{i,d,0} = 1$$

4 Problem 4

4.1 (a)

According to the answer solved by excel solver, an optimal solution is

$$G_{1,2} = 0.0, G_{1,3} = 0.0, G_{1,4} = 0.0, G_{1,5} = 8.0, G_{2,3} = 2.0,$$

$$G_{2,4} = 5.0, G_{2,5} = 0.0, G_{3,4} = 4.0, G_{3,5} = 0.0, G_{4,5} = 0.0$$

where the objective value equals 19.0.

So, the plan is to hire 8 students working in day 1 and 5, hire 2 students working in day 2 and 3, hire 5 students working in day 2 and 4, hire 4 students working in day 3 and 4. Then we can hire a minimum of 19 students to fulfill all the demands.

4.2 (b)

According to the answer solved by excel solver, an optimal solution is

$$G_{1,2} = 2.0, G_{1,3} = 2.0, G_{1,4} = 0.0, G_{1,5} = 1.0, G_{2,3} = 2.0,$$

$$G_{2,4} = 0.0, G_{2,5} = 0.0, G_{3,4} = 0.0, G_{3,5} = 0.0, G_{4,5} = 4.0$$

where the objective value equals 11.0.

So, the plan is to hire 2 students working in day 1 and 2, hire 2 students working in day 1 and 3, hire 1 student working in day 1 and 5, hire 2 students working in day 2 and 3, hire 4 students working in day 4 and 5. Then we can hire a minimum of 11 students to fulfill all the demands.

4.3 (c)

By comparing the two optimal solutions, there are differences between hiring 19 students and 11 students. Although the average numbers of students needed per time slot are identical in the above two parts, the second part is more well-distributed, which decreases the maximum students needed in each day in the constrain functions. Thus, causing a 8 student difference between when minimizing the student amounts to fulfill all the demands.