

CS221 Fall 2018 Homework [scheduling]

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By turning in this assignment, I agree by the Stanford honor code and declare that all of this is my own work.

Problem 1

1.a:

Variables:

$$X = (X_1, X_2, \dots, X_m)$$

$$\text{Domain}_i = \{ 0, 1 \}$$

Factors:

$$F_i(x) = \sum_{j=1}^m K_j \bmod 2 \text{ where } K_j = x_j \text{ when } i \in T_j \text{ otherwise } 0, i = 1, 2, \dots, n$$

Further explanation:

We have m variables, each one (X_j) indicate button j is turned on or off. Domain of X_j is $\{0, 1\}$, where 0 means button j has been pressed $2n$ times (unpressed state), 1 means button j has been pressed $2n+1$ times (pressed state).

We have n factor functions, each one represents the constrain to one light bulb, x is an assignment $x = (x_1, x_2, \dots, x_m)$, the scope of the factor function is (x_1, x_2, \dots, x_m) , the constraint should be:

$$\text{Weight}(x) = \prod_{i=1}^n f_i(X) = 1$$

1.b:

i:

There are two consistent assignments: $\{0, 1, 0\}$ and $\{1, 0, 1\}$.

ii:

Call stacks: (number of '==>' indicates layers of call stack)

CSP without using any heuristics: (9 times backtrack calls)

backtrack(\emptyset , 1, ($X_1=\{0, 1\}$, $X_2=\{0, 1\}$, $X_3=\{0, 1\}$))

==> backtrack($\{X_1 : 0\}$, 1, ($X_2=\{0, 1\}$, $X_3=\{0, 1\}$))

==>==> backtrack($\{X_1 : 0, X_3 : 0\}$, 1, ($X_2=\{0, 1\}$))

$\Rightarrow \Rightarrow \Rightarrow \Rightarrow \Rightarrow \text{backtrack}(\{X_1 : 0, X_2 : 1, X_3 : 0\}, 1, ())$ – success complete assignment
 $\Rightarrow \Rightarrow \Rightarrow \Rightarrow \text{backtrack}(\{X_1 : 0, X_3 : 1\}, 1, (X_2 = \{0, 1\}))$ – failed with X_2
 $\Rightarrow \Rightarrow \text{backtrack}(\{X_1 : 1\}, 1, (X_2 = \{0, 1\}, X_3 = \{0, 1\}))$
 $\Rightarrow \Rightarrow \Rightarrow \Rightarrow \text{backtrack}(\{X_1 : 1, X_3 : 0\}, 1, (X_2 = \{0, 1\}))$ – failed with X_2
 $\Rightarrow \Rightarrow \Rightarrow \Rightarrow \text{backtrack}(\{X_1 : 1, X_3 : 1\}, 1, (X_2 = \{0, 1\}))$
 $\Rightarrow \Rightarrow \Rightarrow \Rightarrow \Rightarrow \text{backtrack}(\{X_1 : 1, X_2 : 0, X_3 : 1\}, 1, ())$ – success complete assignment

iii:

Call stacks: (number of ' \Rightarrow ' indicates layers of call stack)

CSP with AC-3: (7 times backtrack calls)

$\text{backtrack}(\emptyset, 1, (X_1 = \{0, 1\}, X_2 = \{0, 1\}, X_3 = \{0, 1\}))$
 $\Rightarrow \text{backtrack}(\{X_1 : 0\}, 1, (X_2 = \{1\}, X_3 = \{0\}))$
 $\Rightarrow \Rightarrow \Rightarrow \text{backtrack}(\{X_1 : 0, X_3 : 0\}, 1, (X_2 = \{1\}))$
 $\Rightarrow \Rightarrow \Rightarrow \Rightarrow \Rightarrow \text{backtrack}(\{X_1 : 0, X_2 : 1, X_3 : 0\}, 1, ())$ – success complete assignment
 $\Rightarrow \Rightarrow \text{backtrack}(\{X_1 : 1\}, 1, (X_2 = \{0\}, X_3 = \{1\}))$
 $\Rightarrow \Rightarrow \Rightarrow \Rightarrow \text{backtrack}(\{X_1 : 1, X_3 : 1\}, 1, (X_2 = \{0\}))$
 $\Rightarrow \Rightarrow \Rightarrow \Rightarrow \Rightarrow \text{backtrack}(\{X_1 : 1, X_2 : 0, X_3 : 1\}, 1, ())$ – success complete assignment

Problem 2

2.a:

Answer:

The straightforward method is to set a single auxiliary variable A, and construct binary factors $A[1] = X_1$, $A[2] = X_2 + A[1]$, $A[3] = X_3 + A[2]$ and a final unary constraint that A[3] has to meet $A[3] \leq K$.

Let's we can use the adder-circuit- like structure given in class to define this in a formal way:

We have three auxiliary variables: B_1, B_2, B_3 , B_i represents $\sum_{j=1}^i X_j$.

$\text{Domain}(B_i) = \{(m, n) | 0 \leq m \leq 6, 0 \leq n \leq 6\}$

The factors are:

Initialization (Unary factor)	$[B_1[1] = 0]$
Processing (Binary factor)	$[B_i[2] = X_i + B_i[1]]$
Consistency (Binary factor)	$[B_{i+1}[1] = B_i[2]]$
Constraint on final output (Unary factor)	$[B_3[2] \leq K]$

Problem 3

3.c:

profile.txt:

```
# Unit limit per quarter.
minUnits 3
maxUnits 5

# These are the quarters that I need to fill. It is assumed that
# the quarters are sorted in chronological order.
register Win2018
register Spr2019

# Courses I've already taken
taken CS221
taken PHYSICS41

# Courses that I'm requesting
request EARTHSCI400 weight 5 # Win
request EARTHSYS164 # Win
request CS224N # Aut
request CS224S # Spr

Here's the best schedule:
  Quarter  Units      Course
Win2018    4      EARTHSYS164
Spr2019    3      EARTHSCI400
```

Problem 4

4.a:

Answer: Let's consider a worst case that we have at least one $p \in P$ which contains all n variables, we need n -ary factor. As all these variables are all independent, when we eliminate one variable, we got a $n-1$ factor. According to Markov blanket definition, the treewidth for this worst case is $n - 1$.