

# 南京大学 ACM-ICPC 集训队代码模版库



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## 1 General

### 1.1 Code library checksum

```
ab14 #!/usr/bin/python3
c502 import re, sys, hashlib
427e
f7db for line in sys.stdin.read().strip().split("\n") :
ddf5     print(hashlib.md5(re.sub(r'\s|//[.]*', '', line).encode('utf8')).hexdigest()
        [-4:], line)
```

### 1.2 Makefile

```
dab2 .PHONY : run
427e
207e $(t) : $(t).cpp
2d16     g++ --std=c++14 -Wall -D__LOCAL_DEBUG__ -fsanitize=undefined -fsanitize=
        address -ggdb -pipe -o $@ $<
427e
5f25 run : $(t)
bf3e     ./$$(t) < $$(t).in
```

### 1.3 .vimrc

```
914c set nocompatible
733d syntax on
6bbc colorscheme slate
7db5 set number
b0e3 set cursorline
061b set shiftwidth=2
8011 set softtabstop=2
a66d set tabstop=2
d23a set expandtab
5245 set magic
740c set smartindent
bee8 set backspace=indent,eol,start
815d set cmdheight=1
0a40 set laststatus=2
1c67 set whichwrap=b,s,<,>,[,]
```

### 1.4 Stack

```
const int STK_SZ = 2000000;
char STK[STK_SZ * sizeof(void)];
void *STK_BAK;

#if defined(__i386__)
#define SP "%esp"
#elif defined(__x86_64__)
#define SP "%rsp"
#endif

int main() {
    asm volatile("movl SP, %0; movl %1, SP: "=g"(STK_BAK):"g"(STK+sizeof(STK)):");
    ;

    // main program

    asm volatile("movl %0, SP: "=g"(STK_BAK));
    return 0;
}
```

### 1.5 Template

```
#include <bits/stdc++.h>
using namespace std;

#ifdef __LOCAL_DEBUG__
# define _debug(fmt, ...) fprintf(stderr, "[%s] " fmt "\n", \
    __func__, __VA_ARGS__)
#else
# define _debug(...) ((void) 0)
#endif

#define rep(i, n) for (int i=0; i<(n); i++)
#define Rep(i, n) for (int i=1; i<=(n); i++)
#define range(x) begin(x), end(x)
typedef long long LL;
typedef unsigned long long ULL;
```

## 2 Miscellaneous Algorithms

### 2.1 2-SAT

#### Usage:

init(n) Initialize the solver with  $n$  variables.  
 add\_clause(x, xval, y, yval) Add a clause  $(x == xval) \rightarrow (y == yval)$ .  
 solve() Solve the problem. Return **true** if SAT, or **false** if UN-SAT.  
 operator[] (i) Get the value of  $i$ -th variable.

```
0f42 const int MAXN = 100005;
03a9 struct twoSAT {
5c83     int n;
8f72     vector<int> G[MAXN*2];
d060     bool mark[MAXN*2];
b42d     int S[MAXN*2], c;
427e
d34f     void init(int n) {
b985         this->n = n;
f9ec         for (int i=0; i < n*2; i++) G[i].clear();
0609         memset(mark, 0, sizeof(mark));
95cf     }
427e
3bd5     bool dfs(int x) {
bd70         if (mark[x^1]) return false;
c96a         if (mark[x]) return true;
fd23         mark[x] = true;
4bea         S[c++] = x;
bd55         for (int u : G[x]) if (!dfs(u)) return false;
3361         return true;
95cf     }
427e
5894     void add_clause(int x, bool xval, int y, bool yval) {
6afe         x = x * 2 + xval;
e680         y = y * 2 + yval;
81cc         G[x^1].push_back(y);
95cf     }
427e
d0cb     bool solve() {
7c39         for (int i=0; i<n*2; i+=2) {
e63f             if (!mark[i] && !mark[i+1]) {
88fb                 c = 0;
```

```
        if (!dfs(i)) {
            while (c > 0) mark[S[--c]] = false;
            if (!dfs(i+1)) return false;
        }
    }
    return true;
}

bool operator[] (int x) { return mark[2*x+1]; }
};
```

f4b9  
3f03  
86c5  
95cf  
95cf  
95cf  
3361  
95cf  
427e  
fb3b  
329b

### 2.2 Mo's algorithm

All intervals are closed on both sides. When running functions enter() and leave(), the global  $l$  and  $r$  has not changed yet. **Assume the data structure is initialized for empty interval.**

#### Usage:

add\_query(id, l, r) Add id-th query  $[l, r]$ .  
 run() Run Mo's algorithm.  
 yield(id) **TODO.** Yield answer for id-th query.  
 enter(o) **TODO.** Add o-th element.  
 leave(o) **TODO.** Remove o-th element.

```
constexpr int BLOCK_SZ = 300;

struct query { int l, r, id; };
vector<query> queries;

void add_query(int id, int l, int r) {
    queries.push_back(query{l, r, id});
}

int l, r;

// ----- functions to implement -----
inline void yield(int id);
inline void enter(int o);
inline void leave(int o);

void run() {
    if (queries.empty()) return;
    sort(range(queries), [](query lhs, query rhs) {
```

5194  
427e  
3ec4  
d26a  
427e  
1e30  
54c9  
95cf  
427e  
9f6b  
427e  
427e  
50e1  
b20d  
13af  
427e  
37f0  
ab0b  
8508

```

c7f8     int lb = lhs.l / BLOCK_SZ, rb = rhs.l / BLOCK_SZ;
03e7     if (lb != rb) return lb < rb;
0780     return lhs.r < rhs.r;
b251     });
6196     l = queries[0].l;
9644     r = queries[0].r;
38e6     for (int i = l; i <= r; i++) enter(i);
5bc9     for (query q : queries) {
7bc7         while (l > q.l) enter(l - 1), l--;
d646         while (r < q.r) enter(r + 1), r++;
13f0         while (l < q.l) leave(l), l++;
e1c6         while (r > q.r) leave(r), r--;
82f5         yield(q.id);
95cf     }
95cf }
```

## 3 String

### 3.1 Knuth-Morris-Pratt algorithm

```

2836     const int SIZE = 10005;
427e
d02b     struct kmp_matcher {
2d81         char p[SIZE];
9847         int fail[SIZE];
57b7         int len;
427e
60cf     void construct(const char* needle) {
aaa1         len = strlen(p);
3a87         strcpy(p, needle);
3dd4         fail[0] = fail[1] = 0;
d8a8         for (int i = 1; i < len; i++) {
147f             int j = fail[i];
3c79             while (j && p[i] != p[j]) j = fail[j];
4643             fail[i + 1] = p[i] == p[j] ? j + 1 : 0;
95cf         }
95cf     }
427e
c464     inline void found(int pos) {
427e         // ! add codes for having found at pos
95cf     }
```

```

void match(const char* haystack) { // must be called after construct
    const char* t = haystack;
    int n = strlen(t);
    int j = 0;
    rep(i, n) {
        while (j && p[j] != t[i]) j = fail[j];
        if (p[j] == t[i]) j++;
        if (j == len) found(i - len + 1);
    }
}
};
```

```

427e
2daf
700f
8482
8fd0
be8e
4e19
b5d5
f024
95cf
95cf
329b
```

### 3.2 Manacher algorithm

```

struct Manacher {
    int Len;
    vector<int> lc;
    string s;

    void work() {
        lc[1] = 1;
        int k = 1;

        for (int i = 2; i <= Len; i++) {
            int p = k + lc[k] - 1;
            if (i <= p) {
                lc[i] = min(lc[2 * k - i], p - i + 1);
            } else {
                lc[i] = 1;
            }
            while (s[i + lc[i]] == s[i - lc[i]]) lc[i]++;
            if (i + lc[i] > k + lc[k]) k = i;
        }
    }

    void init(const char *tt) {
        int len = strlen(tt);
        s.resize(len * 2 + 10);
        lc.resize(len * 2 + 10);
        s[0] = '*';
        s[1] = '#';
```

```

81d4
cd09
9255
b301
427e
ec07
c033
6bef
427e
491f
7957
5e04
24a1
8e2e
e0e5
95cf
74ff
2b9a
95cf
95cf
427e
bfd5
aaaf
f701
7045
8e13
ae54
```

```

1321     for (int i = 0; i < len; i++) {
e995         s[i * 2 + 2] = tt[i];
69fd         s[i * 2 + 1] = '#';
95cf     }
43fd     s[len * 2 + 1] = '#';
75d1     s[len * 2 + 2] = '\0';
61f7     len = len * 2 + 2;
3e7a     work();
95cf }
427e
b194 pair<int, int> maxpal(int l, int r) {
901a     int center = l + r + 1;
ffb2     int rad = lc[center] / 2;
ab54     int rmid = (l + r + 1) / 2;
17e4     int rl = rmid - rad, rr = rmid + rad - 1;
3908     if ((r ^ 1) & 1) {
69f3     } else rr++;
69dc     return {max(l, rl), min(r, rr)};
95cf }
329b };

```

### 3.3 Aho-corasick automaton

```

a1ad struct AC : Trie {
9143     int fail[MAXN];
daca     int last[MAXN];
427e
8690 void construct() {
93d2     queue<int> q;
a7a6     fail[0] = 0;
ce3c     rep(c, CHARN) {
b1c6         if (int u = tr[0][c]) {
a506             fail[u] = 0;
3e14             q.push(u);
f689             last[u] = 0;
95cf         }
95cf     }
cc78     while (!q.empty()) {
31f0         int r = q.front();
15dd         q.pop();
ce3c         rep(c, CHARN) {
ab59             int u = tr[r][c];

```

```

        if (!u) {
            tr[r][c] = tr[fail[r]][c];
            continue;
        }
        q.push(u);
        int v = fail[r];
        while (v && !tr[v][c]) v = fail[v];
        fail[u] = tr[v][c];
        last[u] = tag[fail[u]] ? fail[u] : last[fail[u]];
    }
}

void found(int pos, int j) {
    if (j) {
        // ! add codes for having found word with tag[j]
        found(pos, last[j]);
    }
}

void find(const char* text) { // must be called after construct()
    int p = 0, c, len = strlen(text);
    rep(i, len) {
        c = id(text[i]);
        p = tr[p][c];
        if (tag[p])
            found(i, p);
        else if (last[p])
            found(i, last[p]);
    }
}
};

```

```

0ef5
9d58
b333
95cf
3e14
b3ff
d2ea
c275
654c
95cf
95cf
95cf
427e
7752
043e
427e
4a96
95cf
95cf
427e
9785
80a4
9c94
b3db
f119
f08e
389b
1e67
299e
95cf
95cf
329b

```

### 3.4 Trie

```

const int MAXN = 12000;
const int CHARN = 26;

inline int id(char c) { return c - 'a'; }

struct Trie {
    int n;

```

```

e6f1
dd87
427e
8ff5
427e
a281
5c83

```

```

f4f5  int tr[MAXN][CHARN]; // Trie tree, 0 denotes fail
35a5  int tag[MAXN];
427e
4fee  Trie() {
3ccc      memset(tr[0], 0, sizeof(tr[0]));
4d52      tag[0] = 0;
46bf      n = 1;
95cf  }
427e
427e  // tag should not be 0
30b0  void add(const char* s, int t) {
d50a      int p = 0, c, len = strlen(s);
9c94      rep(i, len) {
3140          c = id(s[i]);
d6c8          if (!tr[p][c]) {
26dd              memset(tr[p][c], 0, sizeof(tr[p][c]));
2e5c              tag[p][c] = 0;
73bb              tr[p][c] = n++;
95cf          }
f119          p = tr[p][c];
95cf      }
35ef      tag[p] = t;
95cf  }
427e
427e  // returns 0 if not found
427e  // AC automaton does not need this function
216c  int search(const char* s) {
d50a      int p = 0, c, len = strlen(s);
9c94      rep(i, len) {
3140          c = id(s[i]);
f339          if (!tr[p][c]) return 0;
f119          p = tr[p][c];
95cf      }
840e      return tag[p];
95cf  }
329b };

```

### 3.5 Suffix array

The character immediately after the end of the string **MUST** be set to the **UNIQUE SMALLEST** element.

**Usage:**

s[] the source string  
sa[i] the index of starting position of  $i$ -th suffix  
rk[i] the number of suffixes less than the suffix starting from  $i$   
h[i] the longest common prefix between the  $i$ -th and  $(i-1)$ -th  
lexicographically smallest suffixes  
n size of source string  
m size of character set

```

void radix_sort(int x[], int y[], int sa[], int n, int m) {
    static int cnt[1000005]; // size > max(n, m)
    fill(cnt, cnt + m, 0);
    rep(i, n) cnt[x[y[i]]]++;
    partial_sum(cnt, cnt + m, cnt);
    for (int i = n - 1; i >= 0; i--) sa[--cnt[x[y[i]]]] = y[i];
}

void suffix_array(int s[], int sa[], int rk[], int n, int m) {
    static int y[1000005]; // size > n
    copy(s, s + n, rk);
    iota(y, y + n, 0);
    radix_sort(rk, y, sa, n, m);
    for (int j = 1, p = 0; j <= n; j <= 1, m = p, p = 0) {
        for (int i = n - j; i < n; i++) y[p++] = i;
        rep(i, n) if (sa[i] >= j) y[p++] = sa[i] - j;
        radix_sort(rk, y, sa, n, m + 1);
        swap_ranges(rk, rk + n, y);
        rk[sa[0]] = p = 1;
        for (int i = 1; i < n; i++)
            rk[sa[i]] = ((y[sa[i]] == y[sa[i-1]] and y[sa[i]+j] == y[sa[i-1]+j])
                ? p : ++p);
        if (p == n) break;
    }
    rep(i, n) rk[sa[i]] = i;
}

void calc_height(int s[], int sa[], int rk[], int h[], int n) {
    int k = 0;
    h[0] = 0;
    rep(i, n) {
        k = max(k - 1, 0);
        if (rk[i] while (s[i+k] == s[sa[rk[i]-1]+k]) ++k;
        h[rk[i]] = k;
    }
}

```

de09  
ec00  
6066  
93b7  
9154  
acac  
95cf  
427e  
c939  
a69a  
7306  
afbb  
7b42  
c8c2  
8c3a  
9323  
9e9d  
ae41  
ffd2  
445e  
f8dc  
02f0  
95cf  
97d9  
95cf  
427e  
1715  
c41f  
f313  
be8e  
0883  
527d  
56b7  
95cf  
95cf

### 3.6 Rolling hash

**PLEASE** call `init_hash()` in `int main()`!

**Usage:**

`build(str)` Construct the hasher with given string.  
`operator()(l, r)` Get hash value of substring  $[l, r)$ .

```
1e42 const LL mod = 1006658951440146419, g = 967;
9f60 const int MAXN = 200005;
0291 LL pg[MAXN];
427e
dfe7 inline LL mul(LL x, LL y) { return __int128_t(x) * y % mod; }
427e
599a void init_hash() { // must be called in `int main()`
286f     pg[0] = 1;
4af8     for (int i = 1; i < MAXN; i++) pg[i] = mul(pg[i-1], g);
95cf }
427e
7e62 struct hasher {
534a     LL val[MAXN];
427e
4554     void build(const char *str) { // assume lower-case letter only
f937         for (int i = 0; str[i]; i++)
9645             val[i+1] = (mul(val[i], g) + str[i]) % mod;
95cf     }
427e
19f8     LL operator() (int l, int r) { // [l, r)
9986         return (val[r] - mul(val[l], pg[r-l]) + mod) % mod;
95cf     }
329b };
```

## 4 Math

### 4.1 Extended Euclidean algorithm and Chinese remainder theorem

```
4fba void exgcd(LL a, LL b, LL &g, LL &x, LL &y) {
7db6     if (!b) g = a, x = 1, y = 0;
037f     else {
ffca         exgcd(b, a % b, g, y, x);
d798         y -= x * (a / b);
95cf     }
95cf }
```

```
LL crt(LL r[], LL p[], int n) {
    LL q = 1, ret = 0;
    rep (i, n) q *= p[i];
    rep (i, n) {
        LL m = q / p[i];
        LL d, x, y;
        exgcd(p[i], m, d, x, y);
        ret = (ret + y * m * r[i]) % q;
    }
    return (q + ret) % q;
}
```

427e  
e491  
84e6  
00d9  
be8e  
98b4  
9f4f  
b082  
3cd3  
95cf  
2e47  
95cf

### 4.2 Linear basis

```
const int MAXD = 30;
struct linearbasis {
    ULL b[MAXD] = {};

    bool insert(LL v) {
        for (int j = MAXD - 1; j >= 0; j--) {
            if (!(v & (1ll << j))) continue;
            if (b[j]) v ^= b[j]
            else {
                for (int k = 0; k < j; k++)
                    if (v & (1ll << k)) v ^= b[k];
                for (int k = j + 1; k < MAXD; k++)
                    if (b[k] & (1ll << j)) b[k] ^= v;
                b[j] = v;
                return true;
            }
        }
        return false;
    }
};
```

8b44  
03a6  
3558  
427e  
1566  
9b2b  
de36  
ee78  
037f  
7836  
f0b4  
b0aa  
46c9  
8295  
3361  
95cf  
95cf  
438e  
95cf  
329b

### 4.3 Gauss elimination over finite field

```
const LL p = 1000000007;

LL powmod(LL b, LL e) {
```

b784  
427e  
2a2c



```

95a2 LL r = 1;
3e90 while (e) {
1783     if (e & 1) r = r * b % p;
5549     b = b * b % p;
16fc     e >>= 1;
95cf }
547e return r;
95cf }

427e
c130 typedef vector<LL> VLL;
42ac typedef vector<VLL> VVLL;
427e
2c62 LL gauss(VVLL &a, VVLL &b) {
561b     const int n = a.size(), m = b[0].size();
a25e     vector<int> irow(n), icol(n), ipiv(n);
2976     LL det = 1;
427e
be8e     rep (i, n) {
d2b5         int pj = -1, pk = -1;
6b4a         rep (j, n) if (!ipiv[j])
e582             rep (k, n) if (!ipiv[k])
6112                 if (pj == -1 || a[j][k] > a[pj][pk]) {
a905                     pj = j;
657b                     pk = k;
95cf                 }
d480         if (a[pj][pk] == 0) return 0;
0305         ipiv[pk]++;
8dad         swap(a[pj], a[pk]);
aad8         swap(b[pj], b[pk]);
be4d         if (pj != pk) det = (p - det) % p;
d080         irow[i] = pj;
f156         icol[i] = pk;
427e
4ecd         LL c = powmod(a[pk][pk], p - 2);
865b         det = det * a[pk][pk] % p;
c36a         a[pk][pk] = 1;
dd36         rep (j, n) a[pk][j] = a[pk][j] * c % p;
1b23         rep (j, m) b[pk][j] = b[pk][j] * c % p;
f8f3         rep (j, n) if (j != pk) {
e97f             c = a[j][pk];
c449             a[j][pk] = 0;
820b             rep (k, n) a[j][k] = (a[j][k] + p - a[pk][k] * c % p) % p;
f039             rep (k, m) b[j][k] = (b[j][k] + p - b[pk][k] * c % p) % p;
95cf         }

```

```

}

for (int j = n - 1; j >= 0; j--) if (irow[j] != icol[j]) {
    for (int k = 0; k < n; k++) swap(a[k][irow[j]], a[k][icol[j]]);
}
return det;
}

```

95cf  
427e  
37e1  
50dc  
95cf  
f27f  
95cf

## 4.4 Berlekamp-Massey algorithm

Call `berlekamp()` with input sequence  $(x_0, x_1, \dots, x_{n-1})$ . Return a vector of coefficients  $(c_0 = 1, c_1, \dots, c_{m-1})$  with minimum  $m$ , such that  $\sum_{i=0}^m c_i x_{j-i} = 0$  for all possible  $j$ .

```

LL mod = 1000000007;
vector<LL> berlekamp(const vector<LL>& a) {
    vector<LL> p = {1}, r = {1};
    LL dif = 1;
    rep (i, a.size()) {
        LL u = 0;
        rep (j, p.size()) u = (u + p[j] * a[i-j]) % mod;
        if (u == 0) {
            r.insert(r.begin(), 0);
        } else {
            auto op = p;
            p.resize(max(p.size(), r.size() + 1));
            LL idif = powmod(dif, mod - 2);
            rep (j, r.size())
                p[j+1] = (p[j+1] - r[j] * idif % mod * u % mod + mod) % mod;
            dif = u; r = op;
        }
    }
    return p;
}

```

6e50  
97db  
8904  
075b  
8bc9  
1b35  
bd0b  
eae9  
b14c  
8e2e  
0c78  
02f6  
0a2e  
9b57  
dacc  
bcd1  
95cf  
95cf  
e149  
95cf

## 4.5 Fast Walsh-Hadamard transform

```

void fwt(int* a, int n){
    for (int d = 1; d < n; d <= 1)
        for (int i = 0; i < n; i += d < 1)
            rep (j, d){
                int x = a[i+j], y = a[i+j+d];

```

061e  
5595  
05f2  
b833  
7796

```

427e          // a[i+j] = x+y, a[i+j+d] = x-y;    // xor
427e          // a[i+j] = x+y;                    // and
427e          // a[i+j+d] = x+y;                    // or
95cf      }
95cf  }
427e
4db1 void ifwt(int* a, int n){
5595     for (int d = 1; d < n; d <= 1)
05f2         for (int i = 0; i < n; i += d << 1)
b833             rep (j, d){
7796                 int x = a[i+j], y = a[i+j+d];
427e                 // a[i+j] = (x+y)/2, a[i+j+d] = (x-y)/2;    // xor
427e                 // a[i+j] = x-y;                            // and
427e                 // a[i+j+d] = y-x;                            // or
95cf             }
95cf }
427e
2ab6 void conv(int* a, int* b, int n){
950a     fwt(a, n);
e427     fwt(b, n);
8a42     rep(i, n) a[i] *= b[i];
430f     ifwt(a, n);
95cf }

```

## 4.6 Fast fourier transform

```

4e09 const int NMAX = 1<<20;
427e
3fbf typedef complex<double> cplx;
427e
abd1 const double PI = 2*acos(0.0);
12af struct FFT{
c47c     int rev[NMAX];
27d7     cplx omega[NMAX], oinv[NMAX];
9827     int K, N;
427e
1442     FFT(int k){
e209         K = k; N = 1 << k;
b393         rep (i, N){
7ba3             rev[i] = (rev[i>>1]>>1) | ((i&1)<<(K-1));
1908             omega[i] = polar(1.0, 2.0 * PI / N * i);
a166             oinv[i] = conj(omega[i]);

```

```

    }
}

void dft(cplx* a, cplx* w){
    rep (i, N) if (i < rev[i]) swap(a[i], a[rev[i]]);
    for (int l = 2; l <= N; l *= 2){
        int m = l/2;
        for (cplx* p = a; p != a + N; p += l)
            rep (k, m){
                cplx t = w[N/l*k] * p[k+m];
                p[k+m] = p[k] - t; p[k] += t;
            }
    }
}

void fft(cplx* a){dft(a, omega);}
void ifft(cplx* a){
    dft(a, oinv);
    rep (i, N) a[i] /= N;
}

void conv(cplx* a, cplx* b){
    fft(a); fft(b);
    rep (i, N) a[i] *= b[i];
    ifft(a);
}
};

```

## 4.7 Number theoretic transform

```

const int NMAX = 1<<21;

// 998244353 = 7*17*2^23+1, G = 3
const int P = 1004535809, G = 3; // = 479*2^21+1

struct NTT{
    int rev[NMAX];
    LL omega[NMAX], oinv[NMAX];
    int g, g_inv; // g: g_n = G^((P-1)/n)
    int K, N;

    LL powmod(LL b, LL e){

```

```

95a2     LL r = 1;
3e90     while (e){
6624         if (e&1) r = r * b % P;
489e         b = b * b % P;
16fc         e >>= 1;
95cf     }
547e     return r;
95cf }

427e
f420     NTT(int k){
e209         K = k; N = 1 << k;
7652         g = powmod(G, (P-1)/N);
4b3a         g_inv = powmod(g, N-1);
e04f         omega[0] = oinv[0] = 1;
b393         rep (i, N){
7ba3             rev[i] = (rev[i>>1]>>1) | ((i&1)<<(K-1));
ad4f             if (i){
8d8b                 omega[i] = omega[i-1] * g % P;
9e14                 oinv[i] = oinv[i-1] * g_inv % P;
95cf             }
95cf         }
427e     }

9668     void _ntt(LL* a, LL* w){
a215         rep (i, N) if (i < rev[i]) swap(a[i], a[rev[i]]);
ac6e         for (int l = 2; l <= N; l *= 2){
2969             int m = l/2;
7a1d             for (LL* p = a; p != a + N; p += l)
c24f                 rep (k, m){
0ad3                     LL t = w[N/l*k] * p[k+m] % P;
6209                     p[k+m] = (p[k] - t + P) % P;
fa1b                     p[k] = (p[k] + t) % P;
95cf                 }
95cf             }
95cf         }
427e
92ea     void ntt(LL* a){_ntt(a, omega);}
5daf     void intt(LL* a){
1f2a         LL inv = powmod(N, P-2);
9910         _ntt(a, oinv);
a873         rep (i, N) a[i] = a[i] * inv % P;
95cf     }
427e
3a5b     void conv(LL* a, LL* b){

```

```

        ntt(a); ntt(b);
        rep (i, N) a[i] = a[i] * b[i] % P;
        intt(a);
    }
};

```

```

ad16
e49e
5748
95cf
329b

```

## 4.8 Sieve of Euler

```

const int MAXX = 1e7+5;
bool p[MAXX];
int prime[MAXX], sz;

void sieve(){
    p[0] = p[1] = 1;
    for (int i = 2; i < MAXX; i++){
        if (!p[i]) prime[sz++] = i;
        for (int j = 0; j < sz && i*prime[j] < MAXX; j++){
            p[i*prime[j]] = 1;
            if (i % prime[j] == 0) break;
        }
    }
}

```

```

cfc3
5861
73ae
427e
9bc6
9628
1ec8
bf28
e82c
b6a9
5f51
95cf
95cf
95cf

```

## 4.9 Sieve of Euler (General)

```

namespace sieve {
    constexpr int MAXN = 10000007;
    bool p[MAXN]; // true if not prime
    int prime[MAXN], sz;
    int pval[MAXN], pcnt[MAXN];
    int f[MAXN];

    void exec(int N = MAXN) {
        p[0] = p[1] = 1;

        pval[1] = 1;
        pcnt[1] = 0;
        f[1] = 1;

        for (int i = 2; i < N; i++) {
            if (!p[i]) {

```

```

b62e
6589
e982
6ae8
cbf7
6030
427e
76f6
9628
427e
8a8a
bdda
c6b9
427e
a643
01d6

```

```

b2b2     prime[sz++] = i;
37d9     for (LL j = i; j < N; j *= i) {
758c         int b = j / i;
81fd         pval[j] = i * pval[b];
e0f3         pcnt[j] = pcnt[b] + 1;
a96c         f[j] = _____; // f[j] = f(i^pcnt[j])
95cf     }
95cf     }
34c0     for (int j = 0; i * prime[j] < N; j++) {
f87a         int x = i * prime[j]; p[x] = 1;
20cc         if (i % prime[j] == 0) {
9985             pval[x] = pval[i] * prime[j];
3f93             pcnt[x] = pcnt[i] + 1;
8e2e         } else {
cc91             pval[x] = prime[j];
6322             pcnt[x] = 1;
95cf         }
6191         if (x != pval[x]) {
d614             f[x] = f[x / pval[x]] * f[pval[x]]
95cf         }
5f51         if (i % prime[j] == 0) break;
95cf     }
95cf     }
95cf     }
95cf     }

```

#### 4.10 Miller-Rabin primality test

The array `a[]` (excluding sentinel, i.e. `LLONG_MAX`) should be

{2}	when $n < 2,047$ .
{2, 7, 61}	when $n < 4,759,123,141 (2^{32})$ .
{2, 3, 5, 7, 11}	when $n < 2.1 \times 10^{12}$ .
{2, 325, 9375, 28178, 450775, 9780504, 1795265022}	when $n < 2^{64}$ .

```

f16f bool test(LL n){
59f2     if (n < 3) return n==2;
427e     // ! The array a[] should be modified if the range of x changes.
3f11     const LL a[] = {2LL, 7LL, 61LL, LLONG_MAX};
c320     LL r = 0, d = n-1, x;
f410     while (~d & 1) d >>= 1, r++;
2975     for (int i=0; a[i] < n; i++){
ece1         x = powmod(a[i], d, n); // ! powmod must use for 64bit mulmod

```

```

        if (x == 1 || x == n-1) goto next;
        rep (i, r) {
            x = mulmod(x, x, n);
            if (x == n-1) goto next;
        }
        return false;
next;;
    }
    return true;
}

```

```

7f99
e257
d7ff
8d2e
95cf
438e
d490
95cf
3361
95cf

```

#### 4.11 Integer factorization (Pollard's rho)

```

ULL gcd(ULL a, ULL b) {return b ? gcd(b, a % b) : a;}

ULL PollardRho(ULL n){
    ULL c, x, y, d = n;
    if (~n&1) return 2;
    while (d == n){
        x = y = 2;
        d = 1;
        c = rand() % (n - 1) + 1;
        while (d == 1){
            x = (mulmod(x, x, n) + c) % n;
            y = (mulmod(y, y, n) + c) % n;
            y = (mulmod(y, y, n) + c) % n;
            d = gcd(x>y ? x-y : y-x, n);
        }
    }
    return d;
}

```

```

2e6b
427e
54a5
45eb
d3e5
3c69
0964
4753
5952
9e5b
33d5
e1bf
e1bf
a313
95cf
95cf
5d89
95cf

```

## 5 Graph Theory

### 5.1 Strongly connected component

```

const int MAXV = 100005;

struct graph{
    vector<int> adj[MAXV];

```

```

837c
427e
2ea0
88e3

```

```

9cad    stack<int> s;
3d02    int V; // number of vertices
8b6c    int pre[MAXV], lnk[MAXV], scc[MAXV];
27ee    int time, sccn;
427e
bfab    void add_edge(int u, int v){
c71a        adj[u].push_back(v);
95cf    }
427e
d714    void dfs(int u){
7e41        pre[u] = lnk[u] = ++time;
80f6        s.push(u);
18f6        for (int v : adj[u]){
173e            if (!pre[v]){
5f3c                dfs(v);
002c                lnk[u] = min(lnk[u], lnk[v]);
6068            } else if (!scc[v]){
d5df                lnk[u] = min(lnk[u], pre[v]);
95cf            }
95cf        }
8de2        if (lnk[u] == pre[u]){
660f            sccn++;
3c9e            int x;
a69f            do {
3834                x = s.top(); s.pop();
b0e9                scc[x] = sccn;
6757            } while (x != u);
95cf        }
95cf    }
427e
4c88    void find_scc(){
f4a2        time = sccn = 0;
8de7        memset(scc, 0, sizeof scc);
8c2f        memset(pre, 0, sizeof pre);
6901        Rep (i, V){
56d1            if (!pre[i]) dfs(i);
95cf        }
95cf    }
427e
27ce    vector<int> adjc[MAXV];
364d    void contract(){
1a1e        Rep (i, V)
21a2            rep (j, adj[i].size()){
b730                if (scc[i] != scc[adj[i][j]])

```

```

        adjc[scc[i]].push_back(scc[adj[i][j]]);
    }
};

```

b46e  
95cf  
95cf  
329b

## 5.2 Vertex biconnected component

```

const int MAXN = 100005;
struct graph {
    int pre[MAXN], iscut[MAXN], bccno[MAXN], dfs_clock, bcc_cnt;
    vector<int> adj[MAXN], bcc[MAXN];
    set<pair<int, int>> bcce[MAXN];

    stack<pair<int, int>> s;

    void add_edge(int u, int v) {
        adj[u].push_back(v);
        adj[v].push_back(u);
    }

    int dfs(int u, int fa) {
        int lowu = pre[u] = ++dfs_clock;
        int child = 0;
        for (int v : adj[u]) {
            if (!pre[v]) {
                s.push({u, v});
                child++;
                int lowv = dfs(v, u);
                lowu = min(lowu, lowv);
                if (lowv >= pre[u]) {
                    iscut[u] = 1;
                    bcc[bcc_cnt].clear();
                    bcce[bcc_cnt].clear();
                    while (1) {
                        int xu, xv;
                        tie(xu, xv) = s.top(); s.pop();
                        bcce[bcc_cnt].insert({min(xu, xv), max(xu, xv)});
                        if (bccno[xu] != bcc_cnt) {
                            bcc[bcc_cnt].push_back(xu);
                            bccno[xu] = bcc_cnt;
                        }
                    }
                    if (bccno[xv] != bcc_cnt) {

```

0f42  
2ea0  
33ae  
848f  
6b06  
427e  
76f7  
427e  
bfab  
c71a  
a717  
95cf  
427e  
7d3c  
9fe6  
ec14  
18f6  
173e  
e7f8  
fdcf  
f851  
189c  
b687  
6323  
57eb  
90b8  
a147  
a6a3  
a0c3  
0ef5  
3db2  
e0db  
d27f  
95cf  
f357

```

752b         bcc[bcc_cnt].push_back(xv);
57c9         bccno[xv] = bcc_cnt;
95cf     }
7096         if (xu == u && xv == v) break;
95cf     }
03f5         bcc_cnt++;
95cf     }
7470     } else if (pre[v] < pre[u] && v != fa) {
e7f8         s.push({u, v});
f115         lowu = min(lowu, pre[v]);
95cf     }
95cf }
e104     if (fa < 0 && child == 1) iscut[u] = 0;
1160     return lowu;
95cf }
427e
17be void find_bcc(int n) {
8c2f     memset(pre, 0, sizeof pre);
e2d2     memset(iscut, 0, sizeof iscut);
40d3     memset(bccno, -1, sizeof bccno);
fae2     dfs_clock = bcc_cnt = 0;
5c63     rep (i, n) if (!pre[i]) dfs(i, -1);
95cf }
329b };

```

### 5.3 Cut vertices

If the graph is unconnected, the algorithm should be run on each component. One may run `Rep (i, n) if (!dfn[i]) tarjan(i, i)` for unconnected graph.

**Usage:**

`add_edge(u, v)` Add an undirected edge  $(u, v)$ .  
`tarjan(u, fa)` Run Tarjan's algorithm on tree rooted at `fa`. Please call with identical `u` and `fa`.  
`cut[v]` Whether  $v$  is a cut vertex.

```

9f60 const int MAXN = 200005;
0b32 vector<int> adj[MAXN];
18e4 int dfn[MAXN], low[MAXN], idx;
d39d bool cut[MAXN];
427e
bfab void add_edge(int u, int v) {
c71a     adj[u].push_back(v);
a717     adj[v].push_back(u);

```

```

}

void tarjan(int u, int fa) {
    dfn[u] = low[u] = ++idx;
    int child = 0;
    for (int v : adj[u]) {
        if (!dfn[v]) {
            tarjan(v, fa); low[u] = min(low[u], low[v]);
            if (low[v] >= dfn[u] && u != fa) cut[u] = true;
            child += u == fa;
        }
        low[u] = min(low[u], dfn[v]);
    }
    if (u == fa && child > 1) cut[u] = true;
}

```

### 5.4 Minimum spanning arborescence, faster

All vertices are 1-based. Clear the fields when reuse the struct.

**Usage:**

`add_edge(u, v, w)` Add an edge from  $u$  to  $v$  with weight  $w$ .  
`run(n, rt)` Compute the total weight of MSA rooted at `rt`. If not exist, return `LLONG_MIN`.

**Time Complexity:**  $O((|E| + |V| \log |V|) \log |V|)$

```

const int MAXN = 300005;
typedef pair<LL, int> pii;
struct MDST {
    priority_queue<pii, vector<pii>, greater<pii>> heap[MAXN];
    LL shift[MAXN];
    int fa[MAXN], vis[MAXN];

    int find(int x) { return fa[x] == x ? x : fa[x] = find(fa[x]); }

    void unite(int x, int y) {
        x = find(x); y = find(y); fa[y] = x; if (x == y) return;
        if (heap[x].size() < heap[y].size()) {
            swap(heap[x], heap[y]);
            swap(shift[x], shift[y]);
        }
        while (heap[y].size()) {
            auto p = heap[y].top(); heap[y].pop();
            heap[x].emplace(p.first - shift[y] + shift[x], p.second);

```

```

95cf    }
95cf    }
427e
0bbd    void add_edge(int u, int v, LL w) { heap[v].emplace(w, u); }
427e
a526    LL run(int n, int rt) {
f7ff        LL ans = 0;
81f2        iota(fa, fa + n + 1, 0);
19b3        Rep (i, n) if (find(i) != find(rt)) {
a7b1            int u = find(i);
010e            stack<int, vector<int>> s;
eff5            while (find(u) != find(rt)) {
0dda                if (vis[u]) while (s.top() != u) {
c593                    vis[s.top()] = 0; unite(u, s.top()); s.pop();
83c4                } else { vis[u] = 1; s.push(u); }
c76e                while (heap[u].size()) {
b385                    ans += heap[u].top().first - shift[u];
dde2                    shift[u] = heap[u].top().first;
da47                    if (find(heap[u].top().second) != u) break;
9fbb                    heap[u].pop();
95cf                }
6961                if (heap[u].empty()) return LLONG_MIN;
87e6                u = find(heap[u].top().second);
95cf            }
2d46            while (s.size()) { vis[s.top()] = 0; unite(rt, s.top()); s.pop(); }
95cf        }
4206        return ans;
95cf    }
329b    };

```

## 5.5 Maximum flow (Dinic)

### Usage:

add\_edge(u, v, c)      Add an edge from *u* to *v* with capacity *c*.  
max\_flow(s, t)      Compute maximum flow from *s* to *t*.

**Time Complexity:** For general graph,  $O(V^2E)$ ; for network with unit capacity,  $O(\min\{V^{2/3}, \sqrt{E}\}E)$ ; for bipartite network,  $O(\sqrt{VE})$ .

```

bcf8    struct edge{
60e2        int from, to;
5e6d        LL cap, flow;
329b    };
427e

```

```

const int MAXN = 1005;
struct Dinic {
    int n, m, s, t;
    vector<edge> edges;
    vector<int> G[MAXN];
    bool vis[MAXN];
    int d[MAXN];
    int cur[MAXN];

    void add_edge(int from, int to, LL cap) {
        edges.push_back(edge{from, to, cap, 0});
        edges.push_back(edge{to, from, 0, 0});
        m = edges.size();
        G[from].push_back(m-2);
        G[to].push_back(m-1);
    }

    bool bfs() {
        memset(vis, 0, sizeof(vis));
        queue<int> q;
        q.push(s);
        vis[s] = 1;
        d[s] = 0;
        while (!q.empty()) {
            int x = q.front(); q.pop();
            for (int i = 0; i < G[x].size(); i++) {
                edge& e = edges[G[x][i]];
                if (!vis[e.to] && e.cap > e.flow) {
                    vis[e.to] = 1;
                    d[e.to] = d[x] + 1;
                    q.push(e.to);
                }
            }
        }
        return vis[t];
    }

    LL dfs(int x, LL a) {
        if (x == t || a == 0) return a;
        LL flow = 0, f;
        for (int& i = cur[x]; i < G[x].size(); i++) {
            edge& e = edges[G[x][i]];
            if (d[x] + 1 == d[e.to] && (f = dfs(e.to, min(a, e.cap-e.flow))) > 0)
                {

```

```

e2cd
9062
4dbf
9f0c
b891
bbb6
b40a
ddec
427e
5973
7b55
1db7
fe77
dff5
8f2d
95cf
427e
1836
3b73
93d2
5d13
2cd2
721d
cc78
66ba
3b61
b510
bba9
cd72
cf26
ca93
95cf
95cf
95cf
b23b
95cf
427e
9252
6904
8bf9
f515
b510
2374

```

```

1cce         e.flow += f;
e16d         edges[G[x][i]^1].flow -= f;
a74d         flow += f;
23e5         a -= f;
97ed         if(a == 0) break;
95cf     }
95cf     }
84fb     return flow;
95cf }
427e
5bf2 LL max_flow(int s, int t) {
590d     this->s = s; this->t = t;
62e2     LL flow = 0;
ed58     while (bfs()) {
f326         memset(cur, 0, sizeof(cur));
fb3a         flow += dfs(s, LLONG_MAX);
95cf     }
84fb     return flow;
95cf }
427e
c72e vector<int> min_cut() { // call this after maxflow
1df9     vector<int> ans;
df9a     for (int i = 0; i < edges.size(); i++) {
56d8         edge& e = edges[i];
46a2         if(vis[e.from] && !vis[e.to] && e.cap > 0) ans.push_back(i);
95cf     }
4206     return ans;
95cf }
329b };

```

## 5.6 Maximum cardinality bipartite matching (Hungarian)

```

302f #include <bits/stdc++.h>
421c using namespace std;
427e
0d6c #define rep(i, n) for (int i = 0; i < (n); i++)
cfe3 #define Rep(i, n) for (int i = 1; i <= (n); i++)
8843 #define range(x) (x).begin(), (x).end()
5cad typedef long long LL;
427e
84ee struct Hungarian{
fbf6     int nx, ny;

```

```

vector<int> mx, my;
vector<vector<int> > e;
vector<bool> mark;

void init(int nx, int ny){
    this->nx = nx;
    this->ny = ny;
    mx.resize(nx); my.resize(ny);
    e.clear(); e.resize(nx);
    mark.resize(nx);
}

inline void add(int a, int b){
    e[a].push_back(b);
}

bool augment(int i){
    if (!mark[i]) {
        mark[i] = true;
        for (int j : e[i]){
            if (my[j] == -1 || augment(my[j])){
                mx[i] = j; my[j] = i;
                return true;
            }
        }
    }
    return false;
}

int match(){
    int ret = 0;
    fill(range(mx), -1);
    fill(range(my), -1);
    rep (i, nx){
        fill(range(mark), false);
        if (augment(i)) ret++;
    }
    return ret;
}
};

```

```

9ec6
9d4c
edec
427e
8324
c1d1
f9c1
ac92
3f11
1023
95cf
427e
4589
486c
95cf
427e
0c2b
207c
dae4
6a1e
0892
9ca3
3361
95cf
95cf
95cf
438e
95cf
427e
3fac
5b57
b0f1
b957
4ed1
13a5
cc89
95cf
ee0f
95cf
329b

```



## 5.7 Maximum matching of general graph (Edmond's blossom)

### Usage:

init(n) Initialize the template with  $n$  vertices, numbered from 1.  
 add\_edge(u, v) Add an undirected edge  $uv$ .  
 solve() Find the maximum matching. Return the number of matched edges.  
 mate[] The mate of a matched vertex. If it is not matched, then the value is 0.

**Time Complexity:**  $O(|V|^3)$ , but extremely fast in practice.

```
c041 const int MAXN = 1024;
6ab1 struct Blossom {
0b32     vector<int> adj[MAXN];
93d2     queue<int> q;
5c83     int n;
0de2     int label[MAXN], mate[MAXN], save[MAXN], used[MAXN];

427e     void init(int nv) {
2186         n = nv; for (auto& v : adj) v.clear();
3728         fill(range(label), 0); fill(range(mate), 0);
477d         fill(range(save), 0); fill(range(used), 0);
bb35     }

95cf     void add_edge(int u, int v) { adj[u].push_back(v); adj[v].push_back(u); }

427e     void rematch(int x, int y) {
2a48         int m = mate[x]; mate[x] = y;
8af8         if (mate[m] == x) {
1aa4             if (label[x] <= n) {
f4ba                 mate[m] = label[x]; rematch(label[x], m);
740a             } else {
8e2e                 int a = 1 + (label[x] - n - 1) / n;
3341                 int b = 1 + (label[x] - n - 1) % n;
2885                 rematch(a, b); rematch(b, a);
ef33             }
95cf         }
95cf     }
95cf }

427e     void traverse(int x) {
8a50         Rep (i, n) save[i] = mate[i];
43c0         rematch(x, x);
2ef7         Rep (i, n) {
34d7             if (mate[i] != save[i]) used[i] ++;
62c5         }
```

```

        mate[i] = save[i];
    }
}

void relabel(int x, int y) {
    Rep (i, n) used[i] = 0;
    traverse(x); traverse(y);
    Rep (i, n) {
        if (used[i] == 1 and label[i] < 0) {
            label[i] = n + x + (y - 1) * n;
            q.push(i);
        }
    }
}

int solve() {
    Rep (i, n) {
        if (mate[i]) continue;
        Rep (j, n) label[j] = -1;
        label[i] = 0; q = queue<int>(); q.push(i);
        while (q.size()) {
            int x = q.front(); q.pop();
            for (int y : adj[x]) {
                if (mate[y] == 0 and i != y) {
                    mate[y] = x; rematch(x, y); q = queue<int>(); break;
                }
                if (label[y] >= 0) { relabel(x, y); continue; }
                if (label[mate[y]] < 0) {
                    label[mate[y]] = x; q.push(mate[y]);
                }
            }
        }
        int cnt = 0;
        Rep (i, n) cnt += (mate[i] > i);
        return cnt;
    }
};
```

97ef  
 95cf  
 95cf  
 427e  
 8bf8  
 d101  
 c4ea  
 34d7  
 dee9  
 1c22  
 eb31  
 95cf  
 95cf  
 95cf  
 427e  
 a0ce  
 34d7  
 a073  
 1fc0  
 7676  
 1c7d  
 66ba  
 b98c  
 c07f  
 7f36  
 95cf  
 d315  
 58ec  
 c9c4  
 95cf  
 95cf  
 95cf  
 95cf  
 8abb  
 b52f  
 6808  
 95cf  
 329b

## 5.8 Minimum cost maximum flow

```
struct edge{
```

bcf8

```

60e2     int from, to;
d698     int cap, flow;
32cc     LL cost;
329b };
427e
cc3e     const LL INF = LLONG_MAX / 2;
2aa8     const int MAXN = 5005;
c6cb     struct MCMF {
9ceb         int s, t, n, m;
9f0c         vector<edge> edges;
b891         vector<int> G[MAXN];
f74f         bool inq[MAXN]; // queue
8f67         LL d[MAXN];    // distance
9524         int p[MAXN];    // previous
b330         int a[MAXN];   // improvement
427e
f7f2     void add_edge(int from, int to, int cap, LL cost) {
24f0         edges.push_back(edge{from, to, cap, 0, cost});
95f0         edges.push_back(edge{to, from, 0, 0, -cost});
fe77         m = edges.size();
dff5         G[from].push_back(m-2);
8f2d         G[to].push_back(m-1);
95cf     }
427e
3c52     bool spfa(){
93d2         queue<int> q;
8494         fill(d, d + MAXN, INF); d[s] = 0;
fd48         memset(inq, 0, sizeof(inq));
5e7c         q.push(s); inq[s] = true;
2dae         p[s] = 0; a[s] = INT_MAX;
cc78         while (!q.empty()){
b0aa             int u = q.front(); q.pop(); inq[u] = false;
3bba             for (int i : G[u]) {
56d8                 edge& e = edges[i];
3601                 if (e.cap > e.flow && d[e.to] > d[u] + e.cost){
55bc                     d[e.to] = d[u] + e.cost;
0bea                     p[e.to] = G[u][i];
8249                     a[e.to] = min(a[u], e.cap - e.flow);
e5d3                     if (!inq[e.to]) q.push(e.to), inq[e.to] = true;
95cf                 }
95cf             }
95cf         }
6d7c         return d[t] != INF;
95cf     }

```

```

427e     void augment(){
71a4         int u = t;
06f1         while (u != s){
b19d             edges[p[u]].flow += a[t];
db09             edges[p[u]^1].flow -= a[t];
25a9             u = edges[p[u]].from;
e6c9         }
95cf     }
95cf
427e     #ifndef GIVEN_FLOW
6e20     bool min_cost(int s, int t, int f, LL& cost) {
5972         this->s = s; this->t = t;
590d         int flow = 0;
21d4         cost = 0;
23cb         while (spfa()) {
22dc             augment();
bcd8             if (flow + a[t] >= f){
a671                 cost += (f - flow) * d[t]; flow = f;
b14d                 return true;
3361             } else {
8e2e                 flow += a[t]; cost += a[t] * d[t];
2a83             }
95cf         }
95cf         return false;
438e     }
95cf
a8cb     #else
f9a9     int min_cost(int s, int t, LL& cost) {
590d         this->s = s; this->t = t;
21d4         int flow = 0;
23cb         cost = 0;
22dc         while (spfa()) {
bcd8             augment();
2a83             flow += a[t]; cost += a[t] * d[t];
95cf         }
84fb         return flow;
95cf     }
1937     #endif
329b };

```

## 5.9 Fast LCA

All indices of the tree are 1-based.

**Usage:**

preprocess(root)      Initialize with tree rooted at root.  
lca(u, v)              Query the lowest common ancestor of  $u$  and  $v$ .

```
0e34 const int MAXN = 500005;
0b32 vector<int> adj[MAXN];
fccb int id[MAXN], nid;
1356 pair<int, int> st[MAXN << 1][33 - __builtin_clz(MAXN)];
427e
e16d void dfs(int u, int p, int d) {
0df2     st[id[u] = nid++][0] = {d, u};
18f6     for (int v : adj[u]) {
bd87         if (v == p) continue;
f58c         dfs(v, u, d + 1);
08ad         st[nid++][0] = {d, u};
95cf     }
95cf }
427e
3d1b void preprocess(int root) {
3269     nid = 0;
91e1     dfs(root, 0, 1);
5e98     int l = 31 - __builtin_clz(nid);
213b     rep (j, l) rep (i, 1+nid-(1<<j))
1131         st[i][j+1] = min(st[i][j], st[i+(1<<j)][j]);
95cf }
427e
0f0b int lca(int u, int v) {
cfc4     tie(u, v) = minmax(id[u], id[v]);
be9b     int k = 31 - __builtin_clz(v-u+1);
8ebc     return min(st[u][k], st[v-(1<<k)+1][k]).second;
95cf }
```

## 5.10 Heavy-light decomposition

**Time Complexity:** The decomposition itself takes linear time. Each query takes  $O(\log n)$  operations.

```
0f42 const int MAXN = 100005;
0b32 vector<int> adj[MAXN];
42f2 int sz[MAXN], top[MAXN], fa[MAXN], son[MAXN], depth[MAXN], id[MAXN];
427e
be5c void dfs1(int x, int dep, int par){
7489     depth[x] = dep;
```

```
sz[x] = 1;
fa[x] = par;
int maxn = 0, s = 0;
for (int c: adj[x]){
    if (c == par) continue;
    dfs1(c, dep + 1, x);
    sz[x] += sz[c];
    if (sz[c] > maxn){
        maxn = sz[c];
        s = c;
    }
}
son[x] = s;
}

int cid = 0;
void dfs2(int x, int t){
    top[x] = t;
    id[x] = ++cid;
    if (son[x]) dfs2(son[x], t);
    for (int c: adj[x]){
        if (c == fa[x]) continue;
        if (c == son[x]) continue;
        else dfs2(c, c);
    }
}

void decomp(int root){
    dfs1(root, 1, 0);
    dfs2(root, root);
}

void query(int u, int v){
    while (top[u] != top[v]){
        if (depth[top[u]] < depth[top[v]]) swap(u, v);
        // id[top[u]] to id[u]
        u = fa[top[u]];
    }
    if (depth[u] > depth[v]) swap(u, v);
    // id[u] to id[v]
}
```

2ee7  
adb4  
b79d  
c861  
fe45  
fd2f  
b790  
f0f1  
c749  
fe19  
95cf  
95cf  
0e08  
95cf  
427e  
ba54  
3644  
8d96  
d314  
c4a1  
c861  
9881  
5518  
13f9  
95cf  
95cf  
427e  
0f04  
9fa4  
1c88  
95cf  
427e  
2c98  
03a1  
45ec  
427e  
005b  
95cf  
6083  
427e  
95cf

## 5.11 Centroid decomposition

Note that the centroid here is not the exact centroid of the graph. It only guarantees that the size of each subtree does not exceed half of that of the original tree. This is enough to guarantee the correct time complexity. All vertices are numbered from 1. Call `decomp(root)` to use.

### Usage:

`decomp(u, p)` Decompose the tree rooted at  $u$  with parent  $p$ .

**Time Complexity:** The decomposition itself takes  $O(n \log n)$  time.

```

1fb6 vector<int> adj[100005];
88e0 int sz[100005], sum;
427e
f93d void getsz(int u, int p) {
5b36     sz[u] = 1; sum++;
18f6     for (int v : adj[u]) {
bd87         if (v == p) continue;
e3cb         getsz(v, u);
8449         sz[u] += sz[v];
95cf     }
95cf }
427e
67f9 int getcent(int u, int p) {
d51f     for (int v : adj[u])
76e4         if (v != p and sz[v] > sum / 2)
18e3             return getcent(v, u);
81b0     return u;
95cf }
427e
4662 void decompose(int u) {
618e     sum = 0; getsz(u, 0);
303c     u = getcent(u, 0); // update u to the centroid
427e
18f6     for (int v : adj[u]) {
427e         // get answer for subtree v
95cf     }
427e     // get answer for the whole tree
427e     // don't forget to count the centroid itself
427e
18f6     for (int v : adj[u]) { // divide and conquer
c375         adj[v].erase(find(range(adj[v]), u));
fa6b         decompose(v);
a717         adj[v].push_back(u); // restore deleted edge
95cf     }

```

}

95cf

## 5.12 DSU on tree

This implementation avoids parallel existence of multiple data structures but requires that the data structure is invertible. To use this template, implement `merge`, `enter`, `leave` as needed; first call `decomp(root, 0)`, then call `work(root, 0, false)`. Labels of vertices start from 1.

### Usage:

`decomp(u, p)` Decompose the tree  $u$ .  
`work(u, p, keep)` Work for subtree  $u$ . When `keep` is set, information is not cleared.

**Time Complexity:**  $O(n \log n)$  times the complexity for `merge`, `enter`, `leave`.

```

vector<int> adj[100005];
int sz[100005], son[100005];

void decomp(int u, int p) {
    sz[u] = 1;
    for (int v : adj[u]) {
        if (v == p) continue;
        decomp(v, u);
        sz[u] += sz[v];
        if (sz[v] > sz[son[u]]) son[u] = v;
    }
}

template <typename T>
void trav(T fn, int u, int p) {
    fn(u);
    for (int v : adj[u]) if (v != p) trav(fn, v, u);
}

#define for_light(v) for (int v : adj[u]) if (v != p and v != son[u])
void work(int u, int p, bool keep) {
    for_light(v) work(v, u, 0); // process light children

    // process heavy child
    // current data structure contains info of heavy child
    if (son[u]) work(son[u], u, 1);

    auto merge = [u] (int c) { /* count contribution of c */ };

```

```

1fb6
901d
427e
5559
50c0
18f6
bd87
a851
8449
d28c
95cf
95cf
427e
b7ec
62f5
4412
30b3
95cf
427e
7467
33ff
72a2
427e
427e
427e
9866
427e
18a9

```

```

1ab0      auto enter = [] (int c) { /* add vertex c */ };
f241      auto leave = [] (int c) { /* remove vertex c */ };
427e
3d3b      for_light(v) {
74c6          trav(merge, v, u);
c13d          trav(enter, v, u);
95cf      }
427e
427e      // count answer for root and add it
427e      // Warning: special check may apply to root!
c54f      merge(u);
9dec      enter(u);
427e
427e      // Leave current tree
4e3e      if (!keep) trav(leave, u, p);
95cf  }

```

```

struct bit_rupq{ // range update, point query
    int N;
    vector<LL> tr;

    void init(int n) { tr.resize(N = n + 5);}

    LL query(int n) {
        LL ans = 0;
        while (n < N) { ans += tr[n]; n += n & -n; }
        return ans;
    }

    void add(int n, LL x) {
        while (n) { tr[n] += x; n &= n - 1; }
    }
};

```

```

3d03
d7af
99ff
427e
456d
427e
38d4
f7ff
3667
4206
95cf
427e
f4bd
0a2b
95cf
329b

```

## 6 Data Structures

### 6.1 Fenwick tree (point update range query)

```

9976 struct bit_purq { // point update, range query
d7af     int N;
99ff     vector<LL> tr;
427e
456d     void init(int n) { tr.resize(N = n + 5); }
427e
63d0     LL sum(int n) {
f7ff         LL ans = 0;
6770         while (n) { ans += tr[n]; n &= n - 1; }
4206         return ans;
95cf     }
427e
f4bd     void add(int n, LL x){
968e         while (n < N) { tr[n] += x; n += n & -n; }
95cf     }
329b };

```

### 6.2 Fenwick tree (range update point query)

### 6.3 Segment tree

```

LL p;
const int MAXN = 4 * 100006;
struct segtree {
    int l[MAXN], m[MAXN], r[MAXN];
    LL val[MAXN], tadd[MAXN], tmul[MAXN];

#define lson (o<<1)
#define rson (o<<1|1)

    void pull(int o) {
        val[o] = (val[lson] + val[rson]) % p;
    }

    void push_add(int o, LL x) {
        val[o] = (val[o] + x * (r[o] - l[o])) % p;
        tadd[o] = (tadd[o] + x) % p;
    }

    void push_mul(int o, LL x) {
        val[o] = val[o] * x % p;
        tadd[o] = tadd[o] * x % p;
        tmul[o] = tmul[o] * x % p;
    }
};

```

```

3942
1ebb
451a
27be
4510
427e
ac35
1294
427e
1344
bbe9
95cf
427e
e4bc
5dd6
6eff
95cf
427e
d658
b82c
aa86
649f

```

```

95cf    }
427e
b149    void push(int o) {
3159        if (l[o] == m[o]) return;
0a90        if (tmul[o] != 1) {
0f4a            push_mul(lson, tmul[o]);
045e            push_mul(rson, tmul[o]);
ac0a            tmul[o] = 1;
95cf        }
1b82        if (tadd[o]) {
9547            push_add(lson, tadd[o]);
0e73            push_add(rson, tadd[o]);
6234            tadd[o] = 0;
95cf        }
95cf    }
427e
471c    void build(int o, int ll, int rr) {
0e87        int mm = (ll + rr) / 2;
9d27        l[o] = ll; r[o] = rr; m[o] = mm;
ac0a        tmul[o] = 1;
5c92        if (ll == mm) {
001f            scanf("%lld", val + o);
e5b6            val[o] %= p;
8e2e        } else {
7293            build(lson, ll, mm);
5e67            build(rson, mm, rr);
ba26            pull(o);
95cf        }
95cf    }
427e
4406    void add(int o, int ll, int rr, LL x) {
3c16        if (ll <= l[o] && r[o] <= rr) {
db32            push_add(o, x);
8e2e        } else {
c4b0            push(o);
4305            if (m[o] > ll) add(lson, ll, rr, x);
d5a6            if (m[o] < rr) add(rson, ll, rr, x);
ba26            pull(o);
95cf        }
95cf    }
427e
48cd    void mul(int o, int ll, int rr, LL x) {
3c16        if (ll <= l[o] && r[o] <= rr) {
e7d0            push_mul(o, x);

```

```

    } else {
        push(o);
        if (ll < m[o]) mul(lson, ll, rr, x);
        if (m[o] < rr) mul(rson, ll, rr, x);
        pull(o);
    }
}

LL query(int o, int ll, int rr) {
    if (ll <= l[o] && r[o] <= rr) {
        return val[o];
    } else {
        push(o);
        if (rr <= m[o]) return query(lson, ll, rr);
        if (ll >= m[o]) return query(rson, ll, rr);
        return query(lson, ll, rr) + query(rson, ll, rr);
    }
}
} seg;

```

## 6.4 Treap

Self-balanced binary search tree which supports split and merge.

### Usage:

push(x)	Push lazy tags to children.
pull(x)	Update statistics of node $x$ .
Init(x, v)	Initialize node $x$ with value $v$ .
Add(x, v)	Apply addition to subtree $x$ .
Reverse(x)	Apply reversion to subtree $x$ .
Merge(x, y)	Merge trees rooted at $x$ and $y$ . Return the root of new tree.
Split(t, k, x, y)	Split out the left $k$ elements of tree $t$ . The roots of left part and right part are stored in $x$ and $y$ , respectively.
init(n)	Initialize the treap with array of size $n$ .
work(op, l, r)	Range operation over $[l, r)$ .

**Time Complexity:** Expected  $O(\log n)$  per operation.

```

const int MAXN = 200005;
mt19937 gen(time(NULL));
struct Treap {
    int ch[MAXN][2];
    int sz[MAXN], key[MAXN], val[MAXN];
    int add[MAXN], rev[MAXN];

```

```

2b1b    LL sum[MAXN] = {0};
a773    int maxv[MAXN] = {INT_MIN}, minv[MAXN] = {INT_MAX};
427e
a629    void Init(int x, int v) {
5a00        ch[x][0] = ch[x][1] = 0;
d8cd        key[x] = gen(); val[x] = v; pull(x);
95cf    }
427e
3bf9    void pull(int x) {
e1c3        sz[x] = 1 + sz[ch[x][0]] + sz[ch[x][1]];
99f8        sum[x] = val[x] + sum[ch[x][0]] + sum[ch[x][1]];
94e9        maxv[x] = max({val[x], maxv[ch[x][0]], maxv[ch[x][1]]});
6bb9        minv[x] = min({val[x], minv[ch[x][0]], minv[ch[x][1]]});
95cf    }
427e
8c8e    void Add(int x, int a) {
a7b1        val[x] += a; add[x] += a;
832a        sum[x] += LL(sz[x]) * a; maxv[x] += a; minv[x] += a;
95cf    }
427e
aaf6    void Reverse(int x) {
52c6        rev[x] ^= 1;
7850        swap(ch[x][0], ch[x][1]);
95cf    }
427e
1a53    void push(int x) {
5fe5        for (int c : ch[x]) if (c) {
fd76            Add(c, add[x]);
7a53            if (rev[x]) Reverse(c);
95cf        }
49ee        add[x] = 0; rev[x] = 0;
95cf    }
427e
9d2c    int Merge(int x, int y) {
1b09        if (!x || !y) return x | y;
cd7e        push(x); push(y);
bffa        if (key[x] > key[y]) {
a3df            ch[x][1] = Merge(ch[x][1], y); pull(x); return x;
8e2e        } else {
bf9e            ch[y][0] = Merge(x, ch[y][0]); pull(y); return y;
95cf        }
95cf    }
427e
dc7e    void Split(int t, int k, int &x, int &y) {

```

```

        if (t == 0) { x = y = 0; return; }
        push(t);
        if (sz[ch[t][0]] < k) {
            x = t; Split(ch[t][1], k - sz[ch[t][0]] - 1, ch[t][1], y);
        } else {
            y = t; Split(ch[t][0], k, x, ch[t][0]);
        }
        if (x) pull(x); if (y) pull(y);
    }
} treap;

int root;

void init(int n) {
    Rep (i, n) {
        int x; scanf("%d", &x);
        treap.Init(i, x);
        root = (i == 1) ? 1 : treap.Merge(root, i);
    }
}

void work(int op, int l, int r) {
    int tl, tm, tr;
    treap.Split(root, l, tl, tm);
    treap.Split(tm, r - l, tm, tr);
    if (op == 1) {
        int x; scanf("%d", &x); treap.Add(tm, x);
    } else if (op == 2) {
        treap.Reverse(tm);
    } else if (op == 3) {
        printf("%lld %d %d\n",
            treap.sum[tm], treap.minv[tm], treap.maxv[tm]);
    }
    root = treap.Merge(treap.Merge(tl, tm), tr);
}

```

```

6303
f26b
3465
ffd8
8e2e
8a23
95cf
89e3
95cf
b1f4
427e
24b6
427e
d34f
34d7
7681
0ed8
bcc8
95cf
95cf
427e
d030
6639
b6c4
8de3
3658
c039
1dcb
ae78
581d
e092
867f
95cf
6188
95cf

```

## 6.5 Link/cut tree

Dynamic connectivity of undirected acyclic graph. Support single-vertex update, path aggregation and relative LCA query. Vertices are numbered from 1. Zero initialization is enough except for the statistic information.

**Usage:**

<code>pull(x)</code>	Update statistics of node $x$ .
<code>Root(u)</code>	Get the root of tree where vertex $u$ is in.
<code>Link(u, v)</code>	Link two unconnected trees.
<code>Cut(u, v)</code>	Cut an existent edge.
<code>Query(u, v)</code>	Path aggregation.
<code>Update(u, x)</code>	Single point modification.
<code>LCA(u, v, root)</code>	Get the lowest common ancestor of $u$ and $v$ in tree rooted at root.

**Time Complexity:**  $O(\log n)$  per operation

```

2e73 const int MAXN = 1000005;
ca06 struct LCT {
6a6d     int fa[MAXN], ch[MAXN][2], val[MAXN], sum[MAXN];
c6e1     bool rev[MAXN];

427e
eba3     bool isroot(int x) { return ch[fa[x]][0] == x || ch[fa[x]][1] == x; }
f19f     void pull(int x) { sum[x] = val[x] ^ sum[ch[x][0]] ^ sum[ch[x][1]]; }
1c4d     void reverse(int x) { swap(ch[x][0], ch[x][1]); rev[x] ^= 1; }
1a53     void push(int x) {
89a0         if (rev[x]) rep (i, 2) if (ch[x][i]) reverse(ch[x][i]); rev[x] = 0;
95cf     }
425f     void rotate(int x) {
51af         int y = fa[x], z = fa[y], k = ch[y][1] == x, w = ch[x][!k];
e1fe         if (isroot(y)) ch[z][ch[z][1] == y] = x;
1e6f         ch[x][!k] = y; ch[y][k] = w; if (w) fa[w] = y;
6d09         fa[y] = x; fa[x] = z; pull(y);
95cf     }
52c6     void pushall(int x) { if (isroot(x)) pushall(fa[x]); push(x); }
f69c     void splay(int x) {
d095         int y = x, z = 0;
c494         for (pushall(y); isroot(x); rotate(x)) {
ceef             y = fa[x]; z = fa[y];
4449             if (isroot(y)) rotate((ch[y][0] == x) ^ (ch[z][0] == y) ? x : y);
95cf         }
78a0         pull(x);
95cf     }
6229     void access(int x) {
1548         int z = x;
8854         for (int y = 0; x; x = fa[y = x]) { splay(x); ch[x][1] = y; pull(x); }
7afd         splay(z);
95cf     }
a067     void chroot(int x) { access(x); reverse(x); }
126d     void split(int x, int y) { chroot(x); access(y); }
427e

```

```

int Root(int x) {
    for (access(x); ch[x][0]; x = ch[x][0]) push(x);
    splay(x); return x;
}
void Link(int u, int v) { chroot(u); fa[u] = v; }
void Cut(int u, int v) { split(u, v); fa[u] = ch[v][0] = 0; pull(v); }
int Query(int u, int v) { split(u, v); return sum[v]; }
void Update(int u, int x) { splay(u); val[u] = x; }
int LCA(int x, int y, int root) {
    chroot(root); access(x); splay(y);
    while (fa[y]) splay(y = fa[y]);
    return y;
}
};

```

## 6.6 Balanced binary search tree from pb\_ds

```

#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;

tree<int, null_type, less<int>, rb_tree_tag, tree_order_statistics_node_update>
rkt;
// null_tree_node_update

// SAMPLE USAGE
rkt.insert(x);           // insert element
rkt.erase(x);           // erase element
rkt.order_of_key(x);     // obtain the number of elements less than x
rkt.find_by_order(i);    // iterator to i-th (numbered from 0) smallest element
rkt.lower_bound(x);
rkt.upper_bound(x);
rkt.join(rkt2);          // merge tree (only if their ranges do not intersect)
rkt.split(x, rkt2);      // split all elements greater than x to rkt2

```

## 6.7 Persistent segment tree, range k-th query

```

struct node {
    static int n, pos;

    int value;
    node *left, *right;
}

```



```

427e void* operator new(size_t size);
20b0
427e
3dc0 static node* Build(int l, int r) {
b6c5     node* a = new node;
ce96     if (r > l + 1) {
181e         int mid = (l + r) / 2;
3ba2         a->left = Build(l, mid);
8aaf         a->right = Build(mid, r);
8e2e     } else {
bfc4         a->value = 0;
95cf     }
5ffd     return a;
95cf }
427e
5a45 static node* init(int size) {
2c46     n = size;
7ee3     pos = 0;
be52     return Build(0, n);
95cf }
427e
93c0 static int Query(node* lt, node *rt, int l, int r, int k) {
d30c     if (r == l + 1) return l;
181e     int mid = (l + r) / 2;
cb5a     if (rt->left->value - lt->left->value < k) {
8edb         k -= rt->left->value - lt->left->value;
2412         return Query(lt->right, rt->right, mid, r, k);
8e2e     } else {
0119         return Query(lt->left, rt->left, l, mid, k);
95cf     }
95cf }
427e
c9ad static int query(node* lt, node *rt, int k) {
9e27     return Query(lt, rt, 0, n, k);
95cf }
427e
b19c node *Inc(int l, int r, int pos) const {
5794     node* a = new node(*this);
ce96     if (r > l + 1) {
181e         int mid = (l + r) / 2;
203d         if (pos < mid)
f44a             a->left = left->Inc(l, mid, pos);
649a         else
1024             a->right = right->Inc(mid, r, pos);

```

```

    }
    a->value++;
    return a;
}

node *inc(int index) {
    return Inc(0, n, index);
}
} nodes[8000000];

int node::n, node::pos;
inline void* node::operator new(size_t size) {
    return nodes + (pos++);
}

```

```

95cf
2b3e
5ffd
95cf
427e
e80f
c246
95cf
865a
427e
99ce
1987
bb3c
95cf

```

## 6.8 Block list

All indices are 0-based. All ranges are left-closed right-open.

### Usage:

block::fix()	Apply tags to the current block.
Init(l, r)	Range initializer.
Reverse(l, r)	Reverse the range.
Add(l, r, x)	Add $x$ to the range.
Query(l, r)	Range aggregation.

```

const int BLOCK = 800;
typedef vector<int> vi;

```

```

struct block {
    vi data;
    LL sum; int minv, maxv;
    int add; bool rev;

```

```

    block(vi&& vec) : data(move(vec)),
        sum(accumulate(range(data), 0ll)),
        minv(*min_element(range(data))),
        maxv(*max_element(range(data))),
        add(0), rev(0) { }

```

```

    void fix() {
        if (rev) reverse(range(data));      rev = 0;
        if (add) for (int& x : data) x += add; add = 0;
    }

```

```

fd9e
76b3
427e
a771
8fbc
e3b5
41db
427e
d7eb
1f0c
8216
527d
6437
427e
b919
0694
0527
95cf

```

```

427e void merge(block& another) {
8bc4     fix(); another.fix();
b895     vi temp(move(data));
f516     temp.insert(temp.end(), range(another.data));
d02c     *this = block(move(temp));
88ea }
95cf
427e block split(int pos) {
42e8     fix();
3e79     block result(vi(data.begin() + pos, data.end()));
ccab     data.resize(pos); *this = block(move(data));
861a     return result;
56b0 }
95cf };
329b
427e
2a18 typedef list<block>::iterator lit;
427e
ce14 struct blocklist {
5540     list<block> blk;
427e
7b8e void maintain() {
3131     lit it = blk.begin();
4628     while (it != blk.end() && next(it) != blk.end()) {
852d         lit it2 = it;
188c         while (next(it2) != blk.end() &&
3600             it2->data.size() + next(it2)->data.size() <= BLOCK) {
93e1             it2->merge(*next(it2));
e1fa             blk.erase(next(it2));
95cf         }
5771         ++it;
95cf     }
95cf }
427e
b7b3 lit split(int pos) {
2273     for (lit it = blk.begin(); ; it++) {
5502         if (pos == 0) return it;
8e85         while (it->data.size() > pos)
2099             blk.insert(next(it), it->split(pos));
a5a1         pos -= it->data.size();
427e     }
95cf }
95cf
427e

```

```

void Init(int *l, int *r) {
    for (int *cur = l; cur < r; cur += BLOCK)
        blk.emplace_back(vi(cur, min(cur + BLOCK, r)));
}

void Reverse(int l, int r) {
    lit it = split(l), it2 = split(r);
    reverse(it, it2);
    while (it != it2) {
        it->rev ^= 1;
        it++;
    }
    maintain();
}

void Add(int l, int r, int x) {
    lit it = split(l), it2 = split(r);
    while (it != it2) {
        it->sum += LL(x) * it->data.size();
        it->minv += x; it->maxv += x;
        it->add += x; it++;
    }
    maintain();
}

void Query(int l, int r) {
    lit it = split(l), it2 = split(r);
    LL sum = 0; int minv = INT_MAX, maxv = INT_MIN;
    while (it != it2) {
        sum += it->sum;
        minv = min(minv, it->minv);
        maxv = max(maxv, it->maxv);
        it++;
    }
    maintain();
    printf("%lld_%d_%d\n", sum, minv, maxv);
}
} lst;

```

```

1c7b
9919
8950
95cf
427e
a22f
997b
dfd0
8f89
6a06
5283
95cf
b204
95cf
427e
3cce
997b
8f89
e927
03d3
4511
95cf
b204
95cf
427e
3ad3
997b
c33d
8f89
e472
72c4
e1c4
5283
95cf
b204
8792
95cf
958e

```

## 6.9 Persistent block list

Block list that supports persistence. All indices are 0-based. All ranges are left-closed right-open. `std::shared_ptr` is used to ease memory management. One should modify

the constructor of `block` to maintain extra information. Here we use this policy that the size of each block does not exceed `BLOCK`, while the sum of sizes of two adjacent blocks does not less than `BLOCK`.

When some operation that breaks block list property, please call `maintain` in time to restore the property.

#### Usage:

`maintain()` Maintain the block list property.  
`split(pos)` Split the block list at position `pos`. Returns an iterator to a block starting at `pos`.  
`sum(l, r)` An example function of list traversal between  $[l, r)$ .

**Time Complexity:** When `BLOCK` is properly selected, the time complexity is  $O(\sqrt{n})$  per operation.

```
a19e constexpr int BLOCK = 800;
76b3 typedef vector<int> vi;
0563 typedef shared_ptr<vi> pvi;
013b typedef shared_ptr<const vi> pcvi;
427e
a771 struct block {
2989     pcvi data;
8fd0     LL sum;
427e
427e     // add information to maintain
a613     block(pcvi ptr) :
24b5         data(ptr),
0cf0         sum(accumulate(ptr->begin(), ptr->end(), 0ll))
e93b     { }
427e
5c0f     void merge(const block& another) {
0b18         pvi temp = make_shared<vi>(data->begin(), data->end());
ac21         temp->insert(temp->end(), another.data->begin(), another.data->end());
6467         *this = block(temp);
95cf     }
427e
42e8     block split(int pos) {
dac1         block result(make_shared<vi>(data->begin() + pos, data->end()));
01db         *this = block(make_shared<vi>(data->begin(), data->begin() + pos));
56b0         return result;
95cf     }
329b };
427e
2a18 typedef list<block>::iterator lit;
427e
```

```
struct blocklist {
    list<block> blk;

    void maintain() {
        lit it = blk.begin();
        while (it != blk.end() and next(it) != blk.end()) {
            lit it2 = it;
            while (next(it2) != blk.end() and
                    it2->data->size() + next(it2)->data->size() <= BLOCK) {
                it2->merge(*next(it2));
                blk.erase(next(it2));
            }
            ++it;
        }
    }

    lit split(int pos) {
        for (lit it = blk.begin(); ; it++) {
            if (pos == 0) return it;
            while (it->data->size() > pos) {
                blk.insert(next(it), it->split(pos));
            }
            pos -= it->data->size();
        }
    }

    LL sum(int l, int r) { // traverse
        lit it1 = split(l), it2 = split(r);
        LL res = 0;
        while (it1 != it2) {
            res += it1->sum;
            it1++;
        }
        maintain();
        return res;
    }
};
```

ce14  
5540  
427e  
7b8e  
3131  
5e44  
852d  
0b03  
029f  
93e1  
e1fa  
95cf  
5771  
95cf  
95cf  
427e  
b7b3  
2273  
5502  
d480  
2099  
95cf  
a1c8  
95cf  
95cf  
427e  
fd38  
48b4  
ac09  
9f1d  
8284  
61fd  
95cf  
b204  
244d  
95cf  
329b

## 6.10 Sparse table, range minimum query

The array is 0-based and the range is left-closed right-open.

```
const int MAXN = 100007;
```

db63

```

cefd int a[MAXN], st[MAXN][30];
427e
d34f void init(int n){
c73d     int l = log2(n);
cf75     rep (i, n) st[i][0] = a[i];
426b     rep (j, l) rep (i, 1+n-(1<<j))
1131         st[i][j+1] = min(st[i][j], st[i+(1<<j)][j]);
95cf }
427e
c863 int rmq(int l, int r){
f089     int k = log2(r - l);
6117     return min(st[l][k], st[r-(1<<k)][k]);
95cf }

```

## 7 Geometrics

### 7.1 2D geometric template

```

302f #include <bits/stdc++.h>
421c using namespace std;
427e
4553 typedef int T;
c0ae typedef struct pt {
7a9d     T x, y;
ffa9     T operator , (pt a) { return x*a.x + y*a.y; } // inner product
3ec7     T operator * (pt a) { return x*a.y - y*a.x; } // outer product
221a     pt operator + (pt a) { return {x+a.x, y+a.y}; }
8b34     pt operator - (pt a) { return {x-a.x, y-a.y}; }
427e
368b     pt operator * (T k) { return {x*k, y*k}; }
90f4     pt operator - () { return {-x, -y}; }
ba8c } vec;
427e
0ea6 typedef pair<pt, pt> seg;
427e
8d6e bool ptOnSeg(pt& p, seg& s){
ce77     vec v1 = s.first - p, v2 = s.second - p;
de97     return (v1, v2) <= 0 && v1 * v2 == 0;
95cf }
427e
427e // 0 not on segment

```

```

// 1 on segment except vertices
// 2 on vertices
int ptOnSeg2(pt& p, seg& s){
    vec v1 = s.first - p, v2 = s.second - p;
    T ip = (v1, v2);
    if (v1 * v2 != 0 || ip > 0) return 0;
    return (v1, v2) ? 1 : 2;
}

// if two orthogonal rectangles do not touch, return true
inline bool nIntRectRect(seg a, seg b){
    return min(a.first.x, a.second.x) > max(b.first.x, b.second.x) ||
           min(a.first.y, a.second.y) > max(b.first.y, b.second.y) ||
           min(b.first.x, b.second.x) > max(a.first.x, a.second.x) ||
           min(b.first.y, b.second.y) > max(a.first.y, a.second.y);
}

// >0 in order
// <0 out of order
// =0 not standard
inline double rotOrder(vec a, vec b, vec c){return double(a*b)*(b*c);}

inline bool intersect(seg a, seg b){
    // ! if (nIntRectRect(a, b)) return false; // if commented, assume that a
    // and b are non-collinear
    return rotOrder(b.first-a.first, a.second-a.first, b.second-a.first) >= 0 &&
           rotOrder(a.first-b.first, b.second-b.first, a.second-b.first) >= 0;
}

// 0 not intersect
// 1 standard intersection
// 2 vertex-line intersection
// 3 vertex-vertex intersection
// 4 collinear and have common point(s)
int intersect2(seg& a, seg& b){
    if (nIntRectRect(a, b)) return 0;
    vec va = a.second - a.first, vb = b.second - b.first;
    double j1 = rotOrder(b.first-a.first, va, b.second-a.first),
           j2 = rotOrder(a.first-b.first, vb, a.second-b.first);
    if (j1 < 0 || j2 < 0) return 0;
    if (j1 != 0 && j2 != 0) return 1;
    if (j1 == 0 && j2 == 0){
        if (va * vb == 0) return 4; else return 3;
    } else return 2;
}

```

```

95cf }
427e
2c68 template <typename Tp = T>
5894 inline pt getIntersection(pt P, vec v, pt Q, vec w){
6850     static_assert(is_same<Tp, double>::value, "must_be_double!");
7c9a     return P + v * (w*(P-Q)/(v*w));
95cf }
427e
427e // -1 outside the polygon
427e // 0 on the border of the polygon
427e // 1 inside the polygon
cbdd int ptOnPoly(pt p, pt* poly, int n){
5fb4     int wn = 0;
1294     for (int i = 0; i < n; i++) {
427e
3cae         T k, d1 = poly[i].y - p.y, d2 = poly[(i+1)%n].y - p.y;
b957         if (k = (poly[(i+1)%n] - poly[i])*(p - poly[i])){
8c40             if (k > 0 && d1 <= 0 && d2 > 0) wn++;
3c4d             if (k < 0 && d2 <= 0 && d1 > 0) wn--;
aad3         } else return 0;
95cf     }
0a5f     return wn ? 1 : -1;
95cf }
427e
d4a3 istream& operator >> (istream& lhs, pt& rhs){
fa86     lhs >> rhs.x >> rhs.y;
331a     return lhs;
95cf }
427e
07ae istream& operator >> (istream& lhs, seg& rhs){
5cab     lhs >> rhs.first >> rhs.second;
331a     return lhs;
95cf }

```

## 8 Appendices

### 8.1 Primes

#### 8.1.1 First primes

$p$	$g(p)$	$p$	$g(p)$	$p$	$g(p)$	$p$	$g(p)$	$p$	$g(p)$
2	1	3	2	5	2	7	3	11	2
13	2	17	3	19	2	23	5	29	2
31	3	37	2	41	6	43	3	47	5
53	2	59	2	61	2	67	2	71	7
73	5	79	3	83	2	89	3	97	5
101	2	103	5	107	2	109	6	113	3
127	3	131	2	137	3	139	2	149	2
151	6	157	5	163	2	167	5	173	2
179	2	181	2	191	19	193	5	197	2
199	3	211	2	223	3	227	2	229	6

#### 8.1.2 Arbitrary length primes

$\lg p$	$p$	$g(p)$	$p$	$g(p)$
3	967	5	1031	14
4	9859	2	10273	10
5	96331	10	102931	3
6	958543	6	1031137	5
7	9594539	2	10169651	2
8	96243449	3	103211039	7
9	980483981	2	1042484357	2
10	9858935453	2	10261276009	7
11	95748666809	3	101759940101	2
12	950781833849	3	1012797784423	5
13	9739822952371	7	10037217092377	7
14	96181051140397	5	104974966380359	11
15	981030138360889	13	1029038416465403	2
16	9655206098080843	3	10116299875820773	2
17	97687777921994419	3	101506415998163437	2

#### 8.1.3 $\sim 1 \times 10^9$

$p$	$g(p)$	$p$	$g(p)$	$p$	$g(p)$
954854573	3	967607731	2	973215833	3
975831713	3	978949117	2	980766497	3
983879921	3	985918807	3	986608921	29
991136977	5	991752599	13	997137961	11
1003911991	3	1009775293	2	1012423549	6
1021000537	5	1023976897	7	1024153643	2
1037027287	3	1038812881	11	1044754639	3
1045125617	3	1047411427	3	1047753349	6

#### 8.1.4 $\sim 1 \times 10^{18}$

$p$	$g(p)$	$p$	$g(p)$
951970612352230049	3	963284339889659609	3
967495386904694119	3	969751761517096213	2
983238274281901499	2	984647442475101409	23
989286107138674069	11	1002507954383424641	3
1006658951440146419	2	1020152326159075903	3
1034876265966119449	7	1042753851435034019	2
1043609016597371563	2	1045571042176595707	2
1048364250160580293	2	1049495624119026949	2

### 8.2 Pell's equation

$x^2 - ny^2 = 1$ , where  $n$  is a positive nonsquare integer.

Let  $(x_0, y_0)$  be the smallest positive solution of the equation, then the  $k$ -th solution is:

$$\begin{pmatrix} x_k \\ y_k \end{pmatrix} = \begin{pmatrix} x_0 & ny_0 \\ y_0 & x_0 \end{pmatrix}^k \begin{pmatrix} x_0 \\ y_0 \end{pmatrix}$$

Some smallest solutions to Pell's equation:

$n$	2	3	5	6	7	8	10	11	12	13	14	15	17	18	19	20
$x$	3	2	9	5	8	3	19	10	7	649	15	4	33	17	170	9
$y$	2	1	4	2	3	1	6	3	2	180	4	1	8	4	39	2

### 8.3 Burnside's lemma and Polya's enumeration theorem

The Burnside's lemma says that

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

where  $G$  is a group acting on  $X$ ,  $X^g$  is the set of elements in  $X$  that are fixed by  $g$ , i.e.  $X^g = \{x \in X : gx = x\}$ .

The unweighted version of Pólya enumeration theorem says that

$$|Y^X/G| = \frac{1}{|G|} \sum_{g \in G} m^{c_g}$$

where  $m = |X|$  is the number of colors,  $c_g$  is the number of the cycles of permutation  $g$ .

### 8.4 Lagrange's interpolation

For sample points  $(x_0, y_0), \dots, (x_k, y_k)$ , define

$$l_j(x) = \prod_{0 \leq m \leq k, m \neq j} \frac{x - x_m}{x_j - x_m}$$

then the Lagrange polynomial is

$$L(x) = \sum_{j=0}^k y_j l_j(x).$$

To use the script below, type two lines

```
x0 x1 x2 ... xn
y0 y1 y2 ... yn
```

the script will print the fractional coefficient of the polynomial in ascending exponent order.

```
#!/usr/bin/python2
from fractions import *

def polymul(a, b) :
    p = [0] * (len(a)+len(b)-1)
    for e1, c1 in enumerate(a) :
        for e2, c2 in enumerate(b) :
            p[e1+e2] += c1*c2
    return p

x, y = [map(Fraction, raw_input().split()) for _ in 0,0]
n = len(x)
lj = [reduce(polymul, [[-x[m]/(x[j]-x[m]), 1/(x[j]-x[m])]
    for m in range(n) if m != j]]) for j in range(n)]
print '\n'.join(map(str, map(sum, zip(*map(
    lambda a, b : [x*a for x in b], y, lj)))))
```

6dc9  
4b2b  
427e  
796b  
83e4  
f697  
156c  
dfce  
5849  
427e  
f06d  
e80a  
a649  
9dfa  
3cae  
7c0d