

南京大学 ACM-ICPC 集训队代码模版库



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1 General

1.1 Code library checksum

```
ab14 #!/usr/bin/python3
c502 import re, sys, hashlib
427e
f7db for line in sys.stdin.read().strip().split("\n") :
ddf5     print(hashlib.md5(re.sub(r'\s|//[.]*', '', line).encode('utf8')).hexdigest()
        [-4:], line)
```

1.2 Makefile

```
dab2 .PHONY : run
427e
207e $(t) : $(t).cpp
2d16     g++ --std=c++14 -Wall -D__LOCAL_DEBUG__ -fsanitize=undefined -fsanitize=
        address -ggdb -pipe -o $@ $<
427e
5f25 run : $(t)
bf3e     ./$$(t) < $(t).in
```

1.3 .vimrc

```
914c set nocompatible
733d syntax on
6bbc colorscheme slate
7db5 set number
b0e3 set cursorline
061b set shiftwidth=2
8011 set softtabstop=2
a66d set tabstop=2
d23a set expandtab
5245 set magic
740c set smartindent
bee8 set backspace=indent,eol,start
815d set cmdheight=1
0a40 set laststatus=2
1c67 set whichwrap=b,s,<,>,[,]
```

1.4 Stack

```
const int STK_SZ = 2000000;
char STK[STK_SZ * sizeof(void)];
void *STK_BAK;

#if defined(__i386__)
#define SP "%esp"
#elif defined(__x86_64__)
#define SP "%rsp"
#endif

int main() {
    asm volatile("movl SP, %0; movl %1, SP: "=g"(STK_BAK):"g"(STK+sizeof(STK)):");
    ;

    // main program

    asm volatile("movl %0, SP:="g"(STK_BAK));
    return 0;
}
```

1.5 Template

```
#include <bits/stdc++.h>
using namespace std;

#ifdef __LOCAL_DEBUG__
# define _debug(fmt, ...) fprintf(stderr, "[%s] " fmt "\n", \
    __func__, ##_VA_ARGS_)
#else
# define _debug(...) ((void) 0)
#endif

#define rep(i, n) for (int i=0; i<(n); i++)
#define Rep(i, n) for (int i=1; i<=(n); i++)
#define range(x) begin(x), end(x)
typedef long long LL;
typedef unsigned long long ULL;
```

2 Miscellaneous Algorithms

2.1 2-SAT

Usage:

init(n) Initialize the solver with n variables.
add_clause(x, xval, y, yval) Add a clause $(x == xval) \rightarrow (y == yval)$.
solve() Solve the problem. Return **true** if SAT, or **false** if UN-SAT.
operator[] (i) Get the value of i -th variable.

```

0f42 const int MAXN = 100005;
03a9 struct twoSAT {
5c83     int n;
8f72     vector<int> G[MAXN*2];
d060     bool mark[MAXN*2];
b42d     int S[MAXN*2], c;
427e
d34f     void init(int n) {
b985         this->n = n;
f9ec         for (int i=0; i < n*2; i++) G[i].clear();
0609         memset(mark, 0, sizeof(mark));
95cf     }
427e
3bd5     bool dfs(int x) {
bd70         if (mark[x^1]) return false;
c96a         if (mark[x]) return true;
fd23         mark[x] = true;
4bea         S[c++] = x;
bd55         for (int u : G[x]) if (!dfs(u)) return false;
3361         return true;
95cf     }
427e
5894     void add_clause(int x, bool xval, int y, bool yval) {
6afe         x = x * 2 + xval;
e680         y = y * 2 + yval;
81cc         G[x^1].push_back(y);
95cf     }
427e
d0cb     bool solve() {
7c39         for (int i=0; i<n*2; i+=2) {
e63f             if (!mark[i] && !mark[i+1]) {
88fb                 c = 0;

```

```

        if (!dfs(i)) {
            while (c > 0) mark[S[--c]] = false;
            if (!dfs(i+1)) return false;
        }
    }
    return true;
}

bool operator[] (int x) { return mark[2*x+1]; }
};

```

f4b9
 3f03
 86c5
 95cf
 95cf
 95cf
 3361
 95cf
 427e
 fb3b
 329b

2.2 Matroid Intersection

Find the maximum cardinality common independent set of two matroids. Matroids are given by independence oracle.

Usage:

MatroidOracle The independence oracle maintaining an independent set.
Note that the default constructor must properly initialize inner state to an empty set.
insert(x) Insert element labeled x to the independent set.
test(x) Test whether the set is still independent if x is inserted.
MatroidIntersection<MT1, MT2>(n) Construct the matroid intersection solver with n elements labeled from 0 and matroid oracles MT1 and MT2.
run() Run the algorithm and return the matroid intersection.

```

struct MatroidOracle {
    MatroidOracle() { /* TODO */ }
    void insert(int x) { /* TODO */ }
    bool test(int x) const { /* TODO */ }
};

```

0935
 297b
 53e5
 ff18
 329b
 427e
 a015
 94cc
 3288
 5c83
 5550
 fe84
 0b32
 93d2
 427e
 c152

```

const int MAXN = 8192;
template <typename MT1, typename MT2>
struct MatroidIntersection {
    int n;
    bool in[MAXN] = {}, t[MAXN], vis[MAXN];
    int pre[MAXN];
    vector<int> adj[MAXN];
    queue<int> q;

    MatroidIntersection(int n) : n(n) { }

```

```

427e vector<int> getcur() {
2ed1     vector<int> ret;
995a     rep (i, n) if (in[i]) ret.push_back(i);
a585     return ret;
ee0f }
95cf
427e
ca2b void enqueue(int x, int p) {
e5da     if (vis[x]) return;
f4a6     vis[x] = true; pre[x] = p; q.push(x);
ff59     if (t[x]) throw x;
329b };
427e
9081 vector<int> run() {
1026     while (true) {
c40f         vector<int> cur = getcur();
6f47         fill(vis, vis + n, 0);
943b         rep (i, n) adj[i].clear();
0e02         MT2 mt2;
3e54         for (int i : cur) mt2.insert(i);
191d         rep (i, n) t[i] = mt2.test(i);
e167         vector<MT1> mt1s(cur.size());
46d2         vector<MT2> mt2s(cur.size());
660b         rep (i, cur.size()) rep (j, cur.size()) if (i != j) {
3cd7             mt1s[i].insert(cur[j]);
9680             mt2s[i].insert(cur[j]);
95cf         }
e8d7         rep (i, n) if (!in[i]) rep (j, cur.size()) {
3fe9             if (mt1s[j].test(i)) adj[cur[j]].push_back(i);
645e             if (mt2s[j].test(i)) adj[i].push_back(cur[j]);
95cf         }
cf76         q = {};
85eb         try {
2f4f             MT1 mt1;
2f34             for (int i : cur) mt1.insert(i);
4053             rep (i, n) if (mt1.test(i)) enqueue(i, -1);
1c7d             while (q.size()) {
c048                 int u = q.front(); q.pop();
a697                 for (int v : adj[u]) enqueue(v, u);
95cf             }
5a9a         } catch (int v) {
a8f3             while (v >= 0) { in[v] ^= 1; v = pre[v]; }
b333             continue;
95cf         }

```

```

        break;
    };
    return getcur();
}
};

```

```

6173
329b
f2de
95cf
329b

```

3 String

3.1 Knuth-Morris-Pratt algorithm

```

const int SIZE = 10005;

struct kmp_matcher {
    char p[SIZE];
    int fail[SIZE];
    int len;

    void construct(const char* needle) {
        len = strlen(p);
        strcpy(p, needle);
        fail[0] = fail[1] = 0;
        for (int i = 1; i < len; i++) {
            int j = fail[i];
            while (j && p[i] != p[j]) j = fail[j];
            fail[i + 1] = p[i] == p[j] ? j + 1 : 0;
        }

        inline void found(int pos) {
            // ! add codes for having found at pos
        }

        void match(const char* haystack) { // must be called after construct
            const char* t = haystack;
            int n = strlen(t);
            int j = 0;
            rep(i, n) {
                while (j && p[j] != t[i]) j = fail[j];
                if (p[j] == t[i]) j++;
                if (j == len) found(i - len + 1);
            }
        }
    }
}

```

```

2836
427e
d02b
2d81
9847
57b7
427e
60cf
aaa1
3a87
3dd4
d8a8
147f
3c79
4643
95cf
95cf
427e
c464
427e
95cf
427e
2daf
700f
8482
8fd0
be8e
4e19
b5d5
f024
95cf

```

```
95cf }
329b };
```

3.2 Manacher algorithm

```
81d4 struct Manacher {
cd09     int Len;
9255     vector<int> lc;
b301     string s;
427e
ec07     void work() {
c033         lc[1] = 1;
6bef         int k = 1;
427e
491f         for (int i = 2; i <= Len; i++) {
7957             int p = k + lc[k] - 1;
5e04             if (i <= p) {
24a1                 lc[i] = min(lc[2 * k - i], p - i + 1);
8e2e             } else {
e0e5                 lc[i] = 1;
95cf             }
74ff             while (s[i + lc[i]] == s[i - lc[i]]) lc[i]++;
2b9a             if (i + lc[i] > k + lc[k]) k = i;
95cf         }
95cf     }
427e
bfd5     void init(const char *tt) {
aaaf         int len = strlen(tt);
f701         s.resize(len * 2 + 10);
7045         lc.resize(len * 2 + 10);
8e13         s[0] = '*';
ae54         s[1] = '#';
1321         for (int i = 0; i < len; i++) {
e995             s[i * 2 + 2] = tt[i];
69fd             s[i * 2 + 1] = '#';
95cf         }
43fd         s[len * 2 + 1] = '#';
75d1         s[len * 2 + 2] = '\0';
61f7         Len = len * 2 + 2;
3e7a         work();
95cf     }
427e }
```

```
pair<int, int> maxpal(int l, int r) {
    int center = l + r + 1;
    int rad = lc[center] / 2;
    int rmid = (l + r + 1) / 2;
    int rl = rmid - rad, rr = rmid + rad - 1;
    if ((r ^ l) & 1) {
        } else rr++;
    return {max(1, rl), min(r, rr)};
}
};
```

```
b194
901a
ffb2
ab54
17e4
3908
69f3
69dc
95cf
329b
```

3.3 Aho-corasick automaton

```
struct AC : Trie {
    int fail[MAXN];
    int last[MAXN];

    void construct() {
        queue<int> q;
        fail[0] = 0;
        rep(c, CHARN) {
            if (int u = tr[0][c]) {
                fail[u] = 0;
                q.push(u);
                last[u] = 0;
            }
        }
        while (!q.empty()) {
            int r = q.front();
            q.pop();
            rep(c, CHARN) {
                int u = tr[r][c];
                if (!u) {
                    tr[r][c] = tr[fail[r]][c];
                    continue;
                }
                q.push(u);
                int v = fail[r];
                while (v && !tr[v][c]) v = fail[v];
                fail[u] = tr[v][c];
                last[u] = tag[fail[u]] ? fail[u] : last[fail[u]];
            }
        }
    }
};
```

```
a1ad
9143
daca
427e
8690
93d2
a7a6
ce3c
b1c6
a506
3e14
f689
95cf
95cf
cc78
31f0
15dd
ce3c
ab59
0ef5
9d58
b333
95cf
3e14
b3ff
d2ea
c275
654c
95cf
```

```

95cf     }
95cf     }
427e
7752 void found(int pos, int j) {
043e     if (j) {
427e         // ! add codes for having found word with tag[j]
4a96         found(pos, last[j]);
95cf     }
95cf }
427e
9785 void find(const char* text) { // must be called after construct()
80a4     int p = 0, c, len = strlen(text);
9c94     rep(i, len) {
b3db         c = id(text[i]);
f119         p = tr[p][c];
f08e         if (tag[p])
389b             found(i, p);
1e67         else if (last[p])
299e             found(i, last[p]);
95cf     }
95cf }
329b };

```

3.4 Trie

```

e6f1 const int MAXN = 12000;
dd87 const int CHARN = 26;
427e
8ff5 inline int id(char c) { return c - 'a'; }
427e
a281 struct Trie {
5c83     int n;
f4f5     int tr[MAXN][CHARN]; // Trie tree, 0 denotes fail
35a5     int tag[MAXN];
427e
4fee     Trie() {
3ccc         memset(tr[0], 0, sizeof(tr[0]));
4d52         tag[0] = 0;
46bf         n = 1;
95cf     }
427e
427e // tag should not be 0

```

```

void add(const char* s, int t) {
    int p = 0, c, len = strlen(s);
    rep(i, len) {
        c = id(s[i]);
        if (!tr[p][c]) {
            memset(tr[n], 0, sizeof(tr[n]));
            tag[n] = 0;
            tr[p][c] = n++;
        }
        p = tr[p][c];
    }
    tag[p] = t;
}

// returns 0 if not found
// AC automaton does not need this function
int search(const char* s) {
    int p = 0, c, len = strlen(s);
    rep(i, len) {
        c = id(s[i]);
        if (!tr[p][c]) return 0;
        p = tr[p][c];
    }
    return tag[p];
}
};

```

```

30b0
d50a
9c94
3140
d6c8
26dd
2e5c
73bb
95cf
f119
95cf
35ef
95cf
427e
427e
427e
216c
d50a
9c94
3140
f339
f119
95cf
840e
95cf
329b

```

3.5 Suffix array

The character immediately after the end of the string **MUST** be set to the **UNIQUE SMALLEST** element.

Usage:

s[]	the source string
sa[i]	the index of starting position of i -th suffix
rk[i]	the number of suffixes less than the suffix starting from i
h[i]	the longest common prefix between the i -th and $(i-1)$ -th lexicographically smallest suffixes
n	size of source string
m	size of character set

```

void radix_sort(int x[], int y[], int sa[], int n, int m) {
    static int cnt[1000005]; // size > max(n, m)
    fill(cnt, cnt + m, 0);

```

```

de09
ec00
6066

```

```

93b7     rep (i, n) cnt[x[y[i]]]++;
9154     partial_sum(cnt, cnt + m, cnt);
acac     for (int i = n - 1; i >= 0; i--) sa[--cnt[x[y[i]]]] = y[i];
95cf }
427e
c939 void suffix_array(int s[], int sa[], int rk[], int n, int m) {
a69a     static int y[1000005]; // size > n
7306     copy(s, s + n, rk);
afbb     iota(y, y + n, 0);
7b42     radix_sort(rk, y, sa, n, m);
c8c2     for (int j = 1, p = 0; j <= n; j <= 1, m = p, p = 0) {
8c3a         for (int i = n - j; i < n; i++) y[p++] = i;
9323         rep (i, n) if (sa[i] >= j) y[p++] = sa[i] - j;
9e9d         radix_sort(rk, y, sa, n, m + 1);
ae41         swap_ranges(rk, rk + n, y);
ffd2         rk[sa[0]] = p = 1;
445e         for (int i = 1; i < n; i++)
f8dc             rk[sa[i]] = ((y[sa[i]] == y[sa[i-1]] and y[sa[i]+j] == y[sa[i-1]+j])
                    ? p : ++p);
02f0         if (p == n) break;
95cf     }
97d9     rep (i, n) rk[sa[i]] = i;
95cf }
427e
1715 void calc_height(int s[], int sa[], int rk[], int h[], int n) {
c41f     int k = 0;
f313     h[0] = 0;
be8e     rep (i, n) {
0883         k = max(k - 1, 0);
527d         if (rk[i]) while (s[i+k] == s[sa[rk[i]-1]+k]) ++k;
56b7         h[rk[i]] = k;
95cf     }
95cf }

```

3.6 Rolling hash

PLEASE call `init_hash()` in `int main()`!

Usage:

`build(str)` Construct the hasher with given string.
`operator()(l, r)` Get hash value of substring $[l, r)$.

```

1e42 const LL mod = 1006658951440146419, g = 967;
9f60 const int MAXN = 200005;

```

```

LL pg[MAXN];

inline LL mul(LL x, LL y) { return __int128_t(x) * y % mod; }

void init_hash() { // must be called in `int main()`
    pg[0] = 1;
    for (int i = 1; i < MAXN; i++) pg[i] = mul(pg[i-1], g);
}

struct hasher {
    LL val[MAXN];

    void build(const char *str) { // assume lower-case letter only
        for (int i = 0; str[i]; i++)
            val[i+1] = (mul(val[i], g) + str[i]) % mod;
    }

    LL operator() (int l, int r) { // [l, r)
        return (val[r] - mul(val[l], pg[r-l]) + mod) % mod;
    }
};

```

```

0291
427e
dfe7
427e
599a
286f
4af8
95cf
427e
7e62
534a
427e
4554
f937
9645
95cf
427e
19f8
9986
95cf
329b

```

4 Math

4.1 Extended Euclidean algorithm and Chinese remainder theorem

```

void exgcd(LL a, LL b, LL &g, LL &x, LL &y) {
    if (!b) g = a, x = 1, y = 0;
    else {
        exgcd(b, a % b, g, y, x);
        y -= x * (a / b);
    }
}

LL crt(LL r[], LL p[], int n) {
    LL q = 1, ret = 0;
    rep (i, n) q *= p[i];
    rep (i, n) {
        LL m = q / p[i];
        LL d, x, y;
        exgcd(p[i], m, d, x, y);
    }
}

```

```

4fba
7db6
037f
ffca
d798
95cf
95cf
427e
e491
84e6
00d9
be8e
98b4
9f4f
b082

```



```

3cd3     ret = (ret + y * m * r[i]) % q;
95cf     }
2e47     return (q + ret) % q;
95cf     }

```

4.2 Linear basis

```

8b44     const int MAXD = 30;
03a6     struct linearbasis {
3558         ULL b[MAXD] = {};
427e
1566         bool insert(LL v) {
9b2b             for (int j = MAXD - 1; j >= 0; j--) {
de36                 if (!(v & (1ll << j))) continue;
ee78                 if (b[j] & v) b[j] ^= v;
037f                 else {
7836                     for (int k = 0; k < j; k++)
f0b4                         if (v & (1ll << k)) v ^= b[k];
b0aa                     for (int k = j + 1; k < MAXD; k++)
46c9                         if (b[k] & (1ll << j)) b[k] ^= v;
8295                     b[j] = v;
3361                     return true;
95cf                 }
95cf             }
438e             return false;
95cf         }
329b     };

```

4.3 Gauss elimination over finite field

```

b784     const LL p = 1000000007;
427e
2a2c     LL powmod(LL b, LL e) {
95a2         LL r = 1;
3e90         while (e) {
1783             if (e & 1) r = r * b % p;
5549             b = b * b % p;
16fc             e >>= 1;
95cf         }
547e         return r;
95cf     }

```

```

typedef vector<LL> VLL;
typedef vector<VLL> VLLL;

```

```

LL gauss(VLLL &a, VLLL &b) {
    const int n = a.size(), m = b[0].size();
    vector<int> irow(n), icol(n), ipiv(n);
    LL det = 1;

```

```

    rep (i, n) {
        int pj = -1, pk = -1;
        rep (j, n) if (!ipiv[j])
            rep (k, n) if (!ipiv[k])
                if (pj == -1 || a[j][k] > a[pj][pk]) {
                    pj = j;
                    pk = k;
                }
        if (a[pj][pk] == 0) return 0;
        ipiv[pk]++;
        swap(a[pj], a[pk]);
        swap(b[pj], b[pk]);
        if (pj != pk) det = (p - det) % p;
        irow[i] = pj;
        icol[i] = pk;

```

```

        LL c = powmod(a[pk][pk], p - 2);
        det = det * a[pk][pk] % p;
        a[pk][pk] = 1;
        rep (j, n) a[pk][j] = a[pk][j] * c % p;
        rep (j, m) b[pk][j] = b[pk][j] * c % p;
        rep (j, n) if (j != pk) {
            c = a[j][pk];
            a[j][pk] = 0;
            rep (k, n) a[j][k] = (a[j][k] + p - a[pk][k] * c % p) % p;
            rep (k, m) b[j][k] = (b[j][k] + p - b[pk][k] * c % p) % p;
        }
    }

```

```

    for (int j = n - 1; j >= 0; j--) if (irow[j] != icol[j]) {
        for (int k = 0; k < n; k++) swap(a[k][irow[j]], a[k][icol[j]]);
    }
    return det;
}

```

427e
c130
42ac
427e
2c62
561b
a25e
2976
427e
be8e
d2b5
6b4a
e582
6112
a905
657b
95cf
d480
0305
8dad
aad8
be4d
d080
f156
427e
4ecd
865b
c36a
dd36
1b23
f8f3
e97f
c449
820b
f039
95cf
95cf
427e
37e1
50dc
95cf
f27f
95cf

4.4 Berlekamp-Massey algorithm

Call `berlekamp()` with input sequence $(x_0, x_1, \dots, x_{n-1})$. Return a vector of coefficients $(c_0 = 1, c_1, \dots, c_{m-1})$ with minimum m , such that $\sum_{i=0}^m c_i x_{j-i} = 0$ for all possible j .

```
6e50 LL mod = 1000000007;
97db vector<LL> berlekamp(const vector<LL>& a) {
8904     vector<LL> p = {1}, r = {1};
075b     LL dif = 1;
8bc9     rep (i, a.size()) {
1b35         LL u = 0;
bd0b         rep (j, p.size()) u = (u + p[j] * a[i-j]) % mod;
eae9         if (u == 0) {
b14c             r.insert(r.begin(), 0);
8e2e         } else {
0c78             auto op = p;
02f6             p.resize(max(p.size(), r.size() + 1));
0a2e             LL idif = powmod(dif, mod - 2);
9b57             rep (j, r.size())
dacc                 p[j+1] = (p[j+1] - r[j] * idif % mod * u % mod + mod) % mod;
bcd1             dif = u; r = op;
95cf         }
95cf     }
e149     return p;
95cf }
```

4.5 Fast Walsh-Hadamard transform

```
061e void fwt(int* a, int n){
5595     for (int d = 1; d < n; d <= 1)
05f2         for (int i = 0; i < n; i += d << 1)
b833             rep (j, d){
7796                 int x = a[i+j], y = a[i+j+d];
427e                 // a[i+j] = x+y, a[i+j+d] = x-y; // xor
427e                 // a[i+j] = x+y; // and
427e                 // a[i+j+d] = x+y; // or
95cf             }
95cf }
427e
4db1 void ifwt(int* a, int n){
5595     for (int d = 1; d < n; d <= 1)
05f2         for (int i = 0; i < n; i += d << 1)
b833             rep (j, d){
```

```
int x = a[i+j], y = a[i+j+d];
// a[i+j] = (x+y)/2, a[i+j+d] = (x-y)/2; // xor
// a[i+j] = x-y; // and
// a[i+j+d] = y-x; // or
    }
}

void conv(int* a, int* b, int n){
    fwt(a, n);
    fwt(b, n);
    rep(i, n) a[i] *= b[i];
    ifwt(a, n);
}
```

7796
427e
427e
427e
95cf
95cf
427e
2ab6
950a
e427
8a42
430f
95cf

4.6 Fast fourier transform

```
const int NMAX = 1<<20;

typedef complex<double> cplx;

const double PI = 2*acos(0.0);
struct FFT{
    int rev[NMAX];
    cplx omega[NMAX], oinv[NMAX];
    int K, N;
```

4e09
427e
3fbf
427e
abd1
12af
c47c
27d7
9827

```
FFT(int k){
    K = k; N = 1 << k;
    rep (i, N){
        rev[i] = (rev[i>>1]>>1) | ((i&1)<<(K-1));
        omega[i] = polar(1.0, 2.0 * PI / N * i);
        oinv[i] = conj(omega[i]);
    }
}
```

427e
1442
e209
b393
7ba3
1908
a166
95cf
95cf

```
void dft(cplx* a, cplx* w){
    rep (i, N) if (i < rev[i]) swap(a[i], a[rev[i]]);
    for (int l = 2; l <= N; l *= 2){
        int m = l/2;
        for (cplx* p = a; p != a + N; p += l)
            rep (k, m){
                cplx t = w[N/l*k] * p[k+m];
```

427e
b941
a215
ac6e
2969
b3cf
c24f
fe06

```

ecbf         p[k+m] = p[k] - t; p[k] += t;
95cf     }
95cf     }
95cf }
427e
617b void fft(cplx* a){dft(a, omega);}
a123 void ifft(cplx* a){
3b2f     dft(a, oinv);
57fc     rep (i, N) a[i] /= N;
95cf }
427e
bdc0 void conv(cplx* a, cplx* b){
6497     fft(a); fft(b);
12a5     rep (i, N) a[i] *= b[i];
f84e     ifft(a);
95cf }
329b };

```

4.7 Number theoretic transform

```

4ab9 const int NMAX = 1<<21;
427e
427e // 998244353 = 7*17*2^23+1, G = 3
fb9a const int P = 1004535809, G = 3; // = 479*2^21+1
427e
87ab struct NTT{
c47c     int rev[NMAX];
0eda     LL omega[NMAX], oinv[NMAX];
81af     int g, g_inv; // g: g_n = G^((P-1)/n)
9827     int K, N;
427e
2a2c     LL powmod(LL b, LL e){
95a2         LL r = 1;
3e90         while (e){
6624             if (e&1) r = r * b % P;
489e             b = b * b % P;
16fc             e >>= 1;
95cf         }
547e         return r;
95cf     }
427e
f420 NTT(int k){

```

```

K = k; N = 1 << k;
g = powmod(G, (P-1)/N);
g_inv = powmod(g, N-1);
omega[0] = oinv[0] = 1;
rep (i, N){
    rev[i] = (rev[i>>1]>>1) | ((i&1)<<(K-1));
    if (i){
        omega[i] = omega[i-1] * g % P;
        oinv[i] = oinv[i-1] * g_inv % P;
    }
}

void _ntt(LL* a, LL* w){
    rep (i, N) if (i < rev[i]) swap(a[i], a[rev[i]]);
    for (int l = 2; l <= N; l *= 2){
        int m = l/2;
        for (LL* p = a; p != a + N; p += l)
            rep (k, m){
                LL t = w[N/l*k] * p[k+m] % P;
                p[k+m] = (p[k] - t + P) % P;
                p[k] = (p[k] + t) % P;
            }
    }

    void ntt(LL* a){_ntt(a, omega);}
    void intt(LL* a){
        LL inv = powmod(N, P-2);
        _ntt(a, oinv);
        rep (i, N) a[i] = a[i] * inv % P;
    }

    void conv(LL* a, LL* b){
        ntt(a); ntt(b);
        rep (i, N) a[i] = a[i] * b[i] % P;
        intt(a);
    }
};

```

4.8 Sieve of Euler

```

cfc3 const int MAXX = 1e7+5;
5861 bool p[MAXX];
73ae int prime[MAXX], sz;
427e
9bc6 void sieve(){
9628     p[0] = p[1] = 1;
1ec8     for (int i = 2; i < MAXX; i++){
bf28         if (!p[i]) prime[sz++] = i;
e82c         for (int j = 0; j < sz && i*prime[j] < MAXX; j++){
b6a9             p[i*prime[j]] = 1;
5f51             if (i % prime[j] == 0) break;
95cf         }
95cf     }
95cf }

```

```

int x = i * prime[j]; p[x] = 1;
if (i % prime[j] == 0) {
    pval[x] = pval[i] * prime[j];
    pcnt[x] = pcnt[i] + 1;
} else {
    pval[x] = prime[j];
    pcnt[x] = 1;
}
if (x != pval[x]) {
    f[x] = f[x / pval[x]] * f[pval[x]]
}
if (i % prime[j] == 0) break;
}
}
}
}
}

```

f87a
20cc
9985
3f93
8e2e
cc91
6322
95cf
6191
d614
95cf
5f51
95cf
95cf
95cf

4.9 Sieve of Euler (General)

```

b62e namespace sieve {
6589     constexpr int MAXN = 10000007;
e982     bool p[MAXN]; // true if not prime
6ae8     int prime[MAXN], sz;
cbf7     int pval[MAXN], pcnt[MAXN];
6030     int f[MAXN];
427e
76f6     void exec(int N = MAXN) {
9628         p[0] = p[1] = 1;
427e
8a8a         pval[1] = 1;
bdda         pcnt[1] = 0;
c6b9         f[1] = 1;
427e
a643         for (int i = 2; i < N; i++) {
01d6             if (!p[i]) {
b2b2                 prime[sz++] = i;
37d9                 for (LL j = i; j < N; j *= i) {
758c                     int b = j / i;
81fd                     pval[j] = i * pval[b];
e0f3                     pcnt[j] = pcnt[b] + 1;
a96c                     f[j] = _____; // f[j] = f(i^pcnt[j])
95cf                 }
95cf             }
34c0         }
for (int j = 0; i * prime[j] < N; j++) {

```

4.10 Miller-Rabin primality test

The array `a[]` (excluding sentinel, i.e. `LLONG_MAX`) should be

{2}	when $n < 2,047$.
{2, 7, 61}	when $n < 4,759,123,141$ (2^{32}).
{2, 3, 5, 7, 11}	when $n < 2.1 \times 10^{12}$.
{2, 325, 9375, 28178, 450775, 9780504, 1795265022}	when $n < 2^{64}$.

```

bool test(LL n){
    if (n < 3) return n==2;
    // ! The array a[] should be modified if the range of x changes.
    const LL a[] = {2LL, 7LL, 61LL, LLONG_MAX};
    LL r = 0, d = n-1, x;
    while (~d & 1) d >>= 1, r++;
    for (int i=0; a[i] < n; i++){
        x = powmod(a[i], d, n); // ! powmod must use for 64bit mulmod
        if (x == 1 || x == n-1) goto next;
        rep (i, r) {
            x = mulmod(x, x, n);
            if (x == n-1) goto next;
        }
        return false;
    }
next:;
}
return true;

```

f16f
59f2
427e
3f11
c320
f410
2975
ece1
7f99
e257
d7ff
8d2e
95cf
438e
d490
95cf
3361

95cf

}

4.11 Integer factorization (Pollard's rho)

```

2e6b ULL gcd(ULL a, ULL b) {return b ? gcd(b, a % b) : a;}
427e
54a5 ULL PollardRho(ULL n){
45eb     ULL c, x, y, d = n;
d3e5     if (~n&1) return 2;
3c69     while (d == n){
0964         x = y = 2;
4753         d = 1;
5952         c = rand() % (n - 1) + 1;
9e5b         while (d == 1){
33d5             x = (mulmod(x, x, n) + c) % n;
e1bf             y = (mulmod(y, y, n) + c) % n;
e1bf             y = (mulmod(y, y, n) + c) % n;
a313             d = gcd(x>y ? x-y : y-x, n);
95cf         }
95cf     }
5d89     return d;
95cf }
```

4.12 Adaptive Simpson's Method

The Simpson's formula has order 3 algebraic precision.

Usage:

integrate(l, r, eps, fn) Integrate the function fn on interval $[l, r]$. eps is the estimated precision, while est is the current estimation, which can be set to arbitrary value initially.

```

b7ec template <typename T>
9c6c double simpson(double l, double r, T&& f) {
38f4     double mid = (l + r) / 2;
2075     return (f(l) + 4 * f(mid) + f(r)) * (r - l) / 6.0;
95cf }
427e
b7ec template <typename T>
9cbb double integrate(double l, double r, double eps, double est, T&& f) {
38f4     double mid = (l + r) / 2;
5d09     double lv = simpson(l, mid, f), rv = simpson(mid, r, f);
```

```

if (fabs(lv + rv - est) <= 15.0 * eps)
    return lv + rv + (lv + rv - est) / 15.0;
return integrate(l, mid, eps, lv, f) + integrate(mid, r, eps, rv, f);
}
```

d589
036c
13c4
95cf

4.13 Linear Programming (Simplex)

This function solves the following linear program

$$\begin{array}{ll} \max & c^T x \\ \text{s.t.} & Ax \leq b \\ & x \geq 0 \end{array}$$

If the program is infeasible, NAN is returned; if the program is unbounded, DBL_MAX is returned; otherwise, the optimal target is returned and the arguments are stored in x.

```

typedef vector<double> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const double EPS = 1e-9;

double LPSolve(VVD A, VD b, VD c, VD& x) {
    int m = b.size(), n = c.size();
    VI B(m), N(n+1);
    VVD D(m+2, VD(n+2));
    rep (i, m) rep (j, n) D[i][j] = A[i][j];
    rep (i, m) { B[i] = n + i; D[i][n] = -1; D[i][n+1] = b[i]; }
    rep (j, n) { N[j] = j; D[m][j] = -c[j]; }
    N[n] = -1; D[m+1][n] = 1;

    auto pivot = [&] (int r, int s) {
        double inv = 1.0 / D[r][s];
        rep (i, m+2) if (i != r) rep (j, n+2) if (j != s)
            D[i][j] -= D[r][j] * D[i][s] * inv;
        rep (j, n+2) if (j != s) D[r][j] *= inv;
        rep (i, m+2) if (i != r) D[i][s] *= -inv;
        D[r][s] = inv; swap(B[r], N[s]);
    };

    auto simplex = [&](int phase) {
        int x = m + (phase == 1);
        while (true) {
```

db00
9952
89a3
05b7
427e
5eb7
f1f6
1684
319d
7f8f
6b6c
9166
0def
427e
e0f7
3c4b
e090
48ea
79f3
73cf
82f1
329b
427e
3f89
adb8
1026

```

0676     int s = -1;
7e4d     for (int j = 0; j <= n; j++) {
30f5         if (phase == 2 and N[j] == -1) continue;
537c         if (s == -1 or D[x][j] < D[x][s] or
3262             D[x][j] == D[x][s] and N[j] < N[s]) s = j;
95cf     }
083a     if (s < 0 or D[x][s] > -EPS) return true;
bfc5     int r = -1;
356f     for (int i = 0; i < m; i++) {
691d         if (D[i][s] < EPS) continue;
6855         if (r == -1 or D[i][n+1] / D[i][s] < D[r][n+1] / D[r][s] or
26b3             D[i][n+1] / D[i][s] == D[r][n+1] / D[r][s] and
412f             B[i] < B[r]) r = i;
95cf     }
d829     if (r == -1) return false; else pivot(r, s);
95cf }
329b };
427e
7c08     int r = 0;
468b     for (int i = 1; i < m; i++) if (D[i][n+1] < D[r][n+1]) r = i;
8257     if (D[r][n+1] <= -EPS) {
d48d         pivot(r, n);
0175         if (!simplex(1) or D[m+1][n+1] < -EPS) return NAN;
fc91         rep (i, m) if (B[i] == -1) {
0676             int s = -1;
1e86             for (int j = 0; j <= n; j++) if (s == -1 or D[i][j] < D[i][s]
a48f                 or D[i][j] == D[i][s] and N[j] < N[s]) s = j;
c4cd             pivot(i, s);
95cf         }
95cf     }
e566     if (!simplex(2)) return DBL_MAX;
8720     x = VD(n);
3232     rep (i, m) if (B[i] < n) x[B[i]] = D[i][n+1];
bbe4     return D[m][n+1];
95cf }

```

5 Graph Theory

5.1 Vertex biconnected components

```

0f42 const int MAXN = 100005;

```

```

struct graph {
    int pre[MAXN], iscut[MAXN], bccno[MAXN], dfs_clock, bcc_cnt;
    vector<int> adj[MAXN], bcc[MAXN];
    set<pair<int, int>> bcce[MAXN];

    stack<pair<int, int>> s;

    void add_edge(int u, int v) {
        adj[u].push_back(v);
        adj[v].push_back(u);
    }

    int dfs(int u, int fa) {
        int lowu = pre[u] = ++dfs_clock;
        int child = 0;
        for (int v : adj[u]) {
            if (!pre[v]) {
                s.push({u, v});
                child++;
                int lowv = dfs(v, u);
                lowu = min(lowu, lowv);
                if (lowv >= pre[u]) {
                    iscut[u] = 1;
                    bcc[bcc_cnt].clear();
                    bcce[bcc_cnt].clear();
                    while (1) {
                        int xu, xv;
                        tie(xu, xv) = s.top(); s.pop();
                        bcce[bcc_cnt].insert({min(xu, xv), max(xu, xv)});
                        if (bccno[xu] != bcc_cnt) {
                            bcc[bcc_cnt].push_back(xu);
                            bccno[xu] = bcc_cnt;
                        }
                        if (bccno[xv] != bcc_cnt) {
                            bcc[bcc_cnt].push_back(xv);
                            bccno[xv] = bcc_cnt;
                        }
                    }
                    if (xu == u && xv == v) break;
                }
                bcc_cnt++;
            }
            else if (pre[v] < pre[u] && v != fa) {
                s.push({u, v});
                lowu = min(lowu, pre[v]);
            }
        }
    }
}

```

```

2ea0
33ae
848f
6b06
427e
76f7
427e
bfab
c71a
a717
95cf
427e
7d3c
9fe6
ec14
18f6
173e
e7f8
fdcf
f851
189c
b687
6323
57eb
90b8
a147
a6a3
a0c3
0ef5
3db2
e0db
d27f
95cf
f357
752b
57c9
95cf
7096
95cf
03f5
95cf
7470
e7f8
f115

```

```

95cf    }
95cf    }
e104    if (fa < 0 && child == 1) iscut[u] = 0;
1160    return lowu;
95cf    }
427e
17be    void find_bcc(int n) {
8c2f        memset(pre, 0, sizeof pre);
e2d2        memset(iscut, 0, sizeof iscut);
40d3        memset(bccno, -1, sizeof bccno);
fae2        dfs_clock = bcc_cnt = 0;
5c63        rep (i, n) if (!pre[i]) dfs(i, -1);
95cf    }
329b    };

```

5.2 Cut vertices

If the graph is unconnected, the algorithm should be run on each component. One may run `Rep (i, n) if (!dfn[i]) tarjan(i, i)` for unconnected graph.

Usage:

`add_edge(u, v)` Add an undirected edge (u, v) .
`tarjan(u, fa)` Run Tarjan's algorithm on tree rooted at `fa`. Please call with identical `u` and `fa`.
`cut[v]` Whether v is a cut vertex.

```

9f60    const int MAXN = 200005;
0b32    vector<int> adj[MAXN];
18e4    int dfn[MAXN], low[MAXN], idx;
d39d    bool cut[MAXN];
427e
bfab    void add_edge(int u, int v) {
c71a        adj[u].push_back(v);
a717        adj[v].push_back(u);
95cf    }
427e
50aa    void tarjan(int u, int fa) {
9891        dfn[u] = low[u] = ++idx;
ec14        int child = 0;
18f6        for (int v : adj[u]) {
3c64            if (!dfn[v]) {
9636                tarjan(v, fa); low[u] = min(low[u], low[v]);
f368                if (low[v] >= dfn[u] && u != fa) cut[u] = true;
7923                child += u == fa;

```

```

    }
    low[u] = min(low[u], dfn[v]);
}
if (u == fa && child > 1) cut[u] = true;
}

```

5.3 Minimum spanning arborescence, faster

All vertices are 1-based. Clear the fields when reuse the struct.

Usage:

`add_edge(u, v, w)` Add an edge from u to v with weight w .
`run(n, rt)` Compute the total weight of MSA rooted at `rt`. If not exist, return `LLONG_MIN`.

Time Complexity: $O(|E| \log^2 |V|)$

```

const int MAXN = 300005;
typedef pair<LL, int> pii;
struct MDST {
    priority_queue<pii, vector<pii>, greater<pii>> heap[MAXN];
    LL shift[MAXN];
    int fa[MAXN], vis[MAXN];

    int find(int x) { return fa[x] == x ? x : fa[x] = find(fa[x]); }

    void unite(int x, int y) {
        x = find(x); y = find(y); fa[y] = x; if (x == y) return;
        if (heap[x].size() < heap[y].size()) {
            swap(heap[x], heap[y]);
            swap(shift[x], shift[y]);
        }
        while (heap[y].size()) {
            auto p = heap[y].top(); heap[y].pop();
            heap[x].emplace(p.first - shift[y] + shift[x], p.second);
        }
    }

    void add_edge(int u, int v, LL w) { heap[v].emplace(w, u); }

    LL run(int n, int rt) {
        LL ans = 0;
        iota(fa, fa + n + 1, 0);
        Rep (i, n) if (find(i) != find(rt)) {

```

```

a7b1     int u = find(i);
010e     stack<int, vector<int>> s;
eff5     while (find(u) != find(rt)) {
0dda         if (vis[u] while (s.top() != u) {
c593             vis[s.top()] = 0; unite(u, s.top()); s.pop();
83c4         } else { vis[u] = 1; s.push(u); }
c76e         while (heap[u].size()) {
b385             ans += heap[u].top().first - shift[u];
dde2             shift[u] = heap[u].top().first;
da47             if (find(heap[u].top().second) != u) break;
9fbb             heap[u].pop();
95cf         }
6961         if (heap[u].empty()) return LLONG_MIN;
87e6         u = find(heap[u].top().second);
95cf     }
2d46     while (s.size()) { vis[s.top()] = 0; unite(rt, s.top()); s.pop(); }
95cf }
4206     return ans;
95cf }
329b };

```

5.4 Maximum flow (Dinic)

Usage:

add_edge(u, v, c) Add an edge from u to v with capacity c .
max_flow(s, t) Compute maximum flow from s to t .

Time Complexity: For general graph, $O(V^2E)$; for network with unit capacity, $O(\min\{V^{2/3}, \sqrt{E}\}E)$; for bipartite network, $O(\sqrt{VE})$.

```

bcf8 struct edge{
60e2     int from, to;
5e6d     LL cap, flow;
329b };
427e
e2cd const int MAXN = 1005;
9062 struct Dinic {
4dbf     int n, m, s, t;
9f0c     vector<edge> edges;
b891     vector<int> G[MAXN];
bbb6     bool vis[MAXN];
b40a     int d[MAXN];
ddc     int cur[MAXN];
427e

```

```

void add_edge(int from, int to, LL cap) {
    edges.push_back(edge{from, to, cap, 0});
    edges.push_back(edge{to, from, 0, 0});
    m = edges.size();
    G[from].push_back(m-2);
    G[to].push_back(m-1);
}

bool bfs() {
    memset(vis, 0, sizeof(vis));
    queue<int> q;
    q.push(s);
    vis[s] = 1;
    d[s] = 0;
    while (!q.empty()) {
        int x = q.front(); q.pop();
        for (int i = 0; i < G[x].size(); i++) {
            edge& e = edges[G[x][i]];
            if (!vis[e.to] && e.cap > e.flow) {
                vis[e.to] = 1;
                d[e.to] = d[x] + 1;
                q.push(e.to);
            }
        }
    }
    return vis[t];
}

LL dfs(int x, LL a) {
    if (x == t || a == 0) return a;
    LL flow = 0, f;
    for (int& i = cur[x]; i < G[x].size(); i++) {
        edge& e = edges[G[x][i]];
        if (d[x] + 1 == d[e.to] && (f = dfs(e.to, min(a, e.cap-e.flow))) > 0)
        {
            e.flow += f;
            edges[G[x][i]^1].flow -= f;
            flow += f;
            a -= f;
            if(a == 0) break;
        }
    }
    return flow;
}

```



```

427e LL max_flow(int s, int t) {
5bf2     this->s = s; this->t = t;
590d     LL flow = 0;
62e2     while (bfs()) {
ed58         memset(cur, 0, sizeof(cur));
f326         flow += dfs(s, LLONG_MAX);
fb3a     }
95cf     return flow;
84fb }
95cf }
427e
c72e vector<int> min_cut() { // call this after maxflow
1df9     vector<int> ans;
df9a     for (int i = 0; i < edges.size(); i++) {
56d8         edge& e = edges[i];
46a2         if(vis[e.from] && !vis[e.to] && e.cap > 0) ans.push_back(i);
95cf     }
4206     return ans;
95cf }
329b };

```

5.5 Maximum cardinality bipartite matching (Hungarian)

```

302f #include <bits/stdc++.h>
421c using namespace std;
427e
0d6c #define rep(i, n) for (int i = 0; i < (n); i++)
cfe3 #define Rep(i, n) for (int i = 1; i <= (n); i++)
8843 #define range(x) (x).begin(), (x).end()
5cad typedef long long LL;
427e
84ee struct Hungarian{
fbf6     int nx, ny;
9ec6     vector<int> mx, my;
9d4c     vector<vector<int>> > e;
edec     vector<bool> mark;
427e
8324     void init(int nx, int ny){
c1d1         this->nx = nx;
f9c1         this->ny = ny;
ac92         mx.resize(nx); my.resize(ny);
3f11         e.clear(); e.resize(nx);

```

```

        mark.resize(nx);
    }

    inline void add(int a, int b){
        e[a].push_back(b);
    }

    bool augment(int i){
        if (!mark[i]) {
            mark[i] = true;
            for (int j : e[i]){
                if (my[j] == -1 || augment(my[j])){
                    mx[i] = j; my[j] = i;
                    return true;
                }
            }
        }
        return false;
    }

    int match(){
        int ret = 0;
        fill(range(mx), -1);
        fill(range(my), -1);
        rep (i, nx){
            fill(range(mark), false);
            if (augment(i)) ret++;
        }
        return ret;
    }
};

```

5.6 Maximum matching of general graph (Edmond's blossom)

Usage:

init(n)	Initialize the template with n vertices, numbered from 1.
add_edge(u, v)	Add an undirected edge uv .
solve()	Find the maximum matching. Return the number of matched edges.
mate[]	The mate of a matched vertex. If it is not matched, then the value is 0.

Time Complexity: $O(|V|^3)$, but extremely fast in practice.

1023
95cf
427e
4589
486c
95cf
427e
0c2b
207c
dae4
6a1e
0892
9ca3
3361
95cf
95cf
438e
95cf
427e
3fac
5b57
b0f1
b957
4ed1
13a5
cc89
95cf
ee0f
95cf
329b

```

c041 const int MAXN = 1024;
6ab1 struct Blossom {
0b32     vector<int> adj[MAXN];
93d2     queue<int> q;
5c83     int n;
0de2     int label[MAXN], mate[MAXN], save[MAXN], used[MAXN];
427e
2186     void init(int nv) {
3728         n = nv; for (auto& v : adj) v.clear();
477d         fill(range(label), 0); fill(range(mate), 0);
bb35         fill(range(save), 0); fill(range(used), 0);
95cf     }
427e
c2dd     void add_edge(int u, int v) { adj[u].push_back(v); adj[v].push_back(u); }
427e
2a48     void rematch(int x, int y) {
8af8         int m = mate[x]; mate[x] = y;
1aa4         if (mate[m] == x) {
f4ba             if (label[x] <= n) {
740a                 mate[m] = label[x]; rematch(label[x], m);
8e2e             } else {
3341                 int a = 1 + (label[x] - n - 1) / n;
2885                 int b = 1 + (label[x] - n - 1) % n;
ef33                 rematch(a, b); rematch(b, a);
95cf             }
95cf         }
95cf     }
427e
8a50     void traverse(int x) {
43c0         Rep (i, n) save[i] = mate[i];
2ef7         rematch(x, x);
34d7         Rep (i, n) {
62c5             if (mate[i] != save[i]) used[i] ++;
97ef             mate[i] = save[i];
95cf         }
95cf     }
427e
8bf8     void relabel(int x, int y) {
d101         Rep (i, n) used[i] = 0;
c4ea         traverse(x); traverse(y);
34d7         Rep (i, n) {
dee9             if (used[i] == 1 and label[i] < 0) {
1c22                 label[i] = n + x + (y - 1) * n;

```

```

        q.push(i);
    }
}

int solve() {
    Rep (i, n) {
        if (mate[i]) continue;
        Rep (j, n) label[j] = -1;
        label[i] = 0; q = queue<int>(); q.push(i);
        while (q.size()) {
            int x = q.front(); q.pop();
            for (int y : adj[x]) {
                if (mate[y] == 0 and i != y) {
                    mate[y] = x; rematch(x, y); q = queue<int>(); break;
                }
                if (label[y] >= 0) { relabel(x, y); continue; }
                if (label[mate[y]] < 0) {
                    label[mate[y]] = x; q.push(mate[y]);
                }
            }
        }
        int cnt = 0;
        Rep (i, n) cnt += (mate[i] > i);
        return cnt;
    }
};

```

5.7 Minimum cost maximum flow

```

struct edge{
    int from, to;
    int cap, flow;
    LL cost;
};

const LL INF = LLONG_MAX / 2;
const int MAXN = 5005;
struct MCMF {
    int s, t, n, m;
    vector<edge> edges;

```

```

b891 vector<int> G[MAXN];
f74f bool inq[MAXN]; // queue
8f67 LL d[MAXN]; // distance
9524 int p[MAXN]; // previous
b330 int a[MAXN]; // improvement
427e
f7f2 void add_edge(int from, int to, int cap, LL cost) {
24f0     edges.push_back(edge{from, to, cap, 0, cost});
95f0     edges.push_back(edge{to, from, 0, 0, -cost});
fe77     m = edges.size();
dff5     G[from].push_back(m-2);
8f2d     G[to].push_back(m-1);
95cf }
427e
3c52 bool spfa(){
93d2     queue<int> q;
8494     fill(d, d + MAXN, INF); d[s] = 0;
fd48     memset(inq, 0, sizeof(inq));
5e7c     q.push(s); inq[s] = true;
2dae     p[s] = 0; a[s] = INT_MAX;
cc78     while (!q.empty()){
b0aa         int u = q.front(); q.pop(); inq[u] = false;
3bba         for (int i : G[u]) {
56d8             edge& e = edges[i];
3601             if (e.cap > e.flow && d[e.to] > d[u] + e.cost){
55bc                 d[e.to] = d[u] + e.cost;
ddf5                 p[e.to] = i;
8249                 a[e.to] = min(a[u], e.cap - e.flow);
e5d3                 if (!inq[e.to]) q.push(e.to), inq[e.to] = true;
95cf             }
95cf         }
95cf     }
6d7c     return d[t] != INF;
95cf }
427e
71a4 void augment(){
06f1     int u = t;
b19d     while (u != s){
db09         edges[p[u]].flow += a[t];
25a9         edges[p[u]^1].flow -= a[t];
e6c9         u = edges[p[u]].from;
95cf     }
95cf }
427e

```

```

#ifdef GIVEN_FLOW
bool min_cost(int s, int t, int f, LL& cost) {
    this->s = s; this->t = t;
    int flow = 0;
    cost = 0;
    while (spfa()) {
        augment();
        if (flow + a[t] >= f){
            cost += (f - flow) * d[t]; flow = f;
            return true;
        } else {
            flow += a[t]; cost += a[t] * d[t];
        }
    }
    return false;
}
#else
int min_cost(int s, int t, LL& cost) {
    this->s = s; this->t = t;
    int flow = 0;
    cost = 0;
    while (spfa()) {
        augment();
        flow += a[t]; cost += a[t] * d[t];
    }
    return flow;
}
#endif
};

```

5.8 Fast LCA

All indices of the tree are 1-based.

Usage:

```

preprocess(root)    Initialize with tree rooted at root.
lca(u, v)           Query the lowest common ancestor of u and v.

```

```

const int MAXN = 500005;
vector<int> adj[MAXN];
int id[MAXN], nid;
pair<int, int> st[MAXN << 1][33 - __builtin_clz(MAXN)];

void dfs(int u, int p, int d) {

```

```

6e20
5972
590d
21d4
23cb
22dc
bcd b
a671
b14d
3361
8e2e
2a83
95cf
95cf
438e
95cf
a8cb
f9a9
590d
21d4
23cb
22dc
bcd b
2a83
95cf
84fb
95cf
1937
329b

```

```

0e34
0b32
fccb
1356
427e
e16d

```

```

0df2     st[id[u] = nid++][0] = {d, u};
18f6     for (int v : adj[u]) {
bd87         if (v == p) continue;
f58c         dfs(v, u, d + 1);
08ad         st[nid++][0] = {d, u};
95cf     }
95cf }
427e
3d1b void preprocess(int root) {
3269     nid = 0;
91e1     dfs(root, 0, 1);
5e98     int l = 31 - __builtin_clz(nid);
213b     rep (j, l) rep (i, 1+nid-(1<<j))
1131         st[i][j+1] = min(st[i][j], st[i+(1<<j)][j]);
95cf }
427e
0f0b int lca(int u, int v) {
cfc4     tie(u, v) = minmax(id[u], id[v]);
be9b     int k = 31 - __builtin_clz(v-u+1);
8ebc     return min(st[u][k], st[v-(1<<k)+1][k]).second;
95cf }

```

5.9 Heavy-light decomposition

Time Complexity: The decomposition itself takes linear time. Each query takes $O(\log n)$ operations.

```

0f42 const int MAXN = 100005;
0b32 vector<int> adj[MAXN];
42f2 int sz[MAXN], top[MAXN], fa[MAXN], son[MAXN], depth[MAXN], id[MAXN];
427e
be5c void dfs1(int x, int dep, int par){
7489     depth[x] = dep;
2ee7     sz[x] = 1;
adb4     fa[x] = par;
b79d     int maxn = 0, s = 0;
c861     for (int c: adj[x]){
fe45         if (c == par) continue;
fd2f         dfs1(c, dep + 1, x);
b790         sz[x] += sz[c];
f0f1         if (sz[c] > maxn){
c749             maxn = sz[c];
fe19             s = c;

```

```

    }
    }
    son[x] = s;
}

int cid = 0;
void dfs2(int x, int t){
    top[x] = t;
    id[x] = ++cid;
    if (son[x]) dfs2(son[x], t);
    for (int c: adj[x]){
        if (c == fa[x]) continue;
        if (c == son[x]) continue;
        else dfs2(c, c);
    }
}

void decomp(int root){
    dfs1(root, 1, 0);
    dfs2(root, root);
}

void query(int u, int v){
    while (top[u] != top[v]){
        if (depth[top[u]] < depth[top[v]]) swap(u, v);
        // id[top[u]] to id[u]
        u = fa[top[u]];
    }
    if (depth[u] > depth[v]) swap(u, v);
    // id[u] to id[v]
}

```

95cf
95cf
0e08
95cf
427e
ba54
3644
8d96
d314
c4a1
c861
9881
5518
13f9
95cf
95cf
427e
0f04
9fa4
1c88
95cf
427e
2c98
03a1
45ec
427e
005b
95cf
6083
427e
95cf

5.10 Centroid decomposition

Note that the centroid here is not the exact centroid of the graph. It only guarantees that the size of each subtree does not exceed half of that of the original tree. This is enough to guarantee the correct time complexity. All vertices are numbered from 1. Call `decomp(root)` to use.

Usage:

`decomp(u, p)` Decompose the tree rooted at u with parent p .

Time Complexity: The decomposition itself takes $O(n \log n)$ time.

```

1fb6 vector<int> adj[100005];
88e0 int sz[100005], sum;
427e
f93d void getsz(int u, int p) {
5b36     sz[u] = 1; sum++;
18f6     for (int v : adj[u]) {
bd87         if (v == p) continue;
e3cb         getsz(v, u);
8449         sz[u] += sz[v];
95cf     }
95cf }
427e
67f9 int getcent(int u, int p) {
d51f     for (int v : adj[u])
76e4         if (v != p and sz[v] > sum / 2)
18e3             return getcent(v, u);
81b0     return u;
95cf }
427e
4662 void decompose(int u) {
618e     sum = 0; getsz(u, 0);
303c     u = getcent(u, 0); // update u to the centroid
427e
18f6     for (int v : adj[u]) {
427e         // get answer for subtree v
95cf     }
427e     // get answer for the whole tree
427e     // don't forget to count the centroid itself
427e
18f6     for (int v : adj[u]) { // divide and conquer
c375         adj[v].erase(find(range(adj[v]), u));
fa6b         decompose(v);
a717         adj[v].push_back(u); // restore deleted edge
95cf     }
95cf }

```

5.11 DSU on tree

This implementation avoids parallel existence of multiple data structures but requires that the data structure is invertible. To use this template, implement merge, enter, leave as needed; first call decomp(root, 0), then call work(root, 0, false). Labels of vertices start from 1.

Usage:

decomp(u, p) Decompose the tree *u*.
work(u, p, keep) Work for subtree *u*. When keep is set, information is not cleared.

Time Complexity: $O(n \log n)$ times the complexity for merge, enter, leave.

```

vector<int> adj[100005];
int sz[100005], son[100005];

void decomp(int u, int p) {
    sz[u] = 1;
    for (int v : adj[u]) {
        if (v == p) continue;
        decomp(v, u);
        sz[u] += sz[v];
        if (sz[v] > sz[son[u]]) son[u] = v;
    }
}

template <typename T>
void trav(T fn, int u, int p) {
    fn(u);
    for (int v : adj[u]) if (v != p) trav(fn, v, u);
}

#define for_light(v) for (int v : adj[u]) if (v != p and v != son[u])
void work(int u, int p, bool keep) {
    for_light(v) work(v, u, 0); // process light children

    // process heavy child
    // current data structure contains info of heavy child
    if (son[u]) work(son[u], u, 1);

    auto merge = [u] (int c) { /* count contribution of c */ };
    auto enter = [] (int c) { /* add vertex c */ };
    auto leave = [] (int c) { /* remove vertex c */ };

    for_light(v) {
        trav(merge, v, u);
        trav(enter, v, u);
    }

    // count answer for root and add it
    // Warning: special check may apply to root!

```

```

c54f     merge(u);
9dec     enter(u);
427e
427e     // Leave current tree
4e3e     if (!keep) trav(leave, u, p);
95cf }

```

```

}

void add(int n, LL x) {
    while (n) { tr[n] += x; n &= n - 1; }
}
};

```

```

95cf
427e
f4bd
0a2b
95cf
329b

```

6 Data Structures

6.1 Fenwick tree (point update range query)

```

9976 struct bit_purq { // point update, range query
d7af     int N;
99ff     vector<LL> tr;
427e
2d99     void init(int n) { tr.assign(N = n + 5, 0); }
427e
63d0     LL sum(int n) {
f7ff         LL ans = 0;
6770         while (n) { ans += tr[n]; n &= n - 1; }
4206         return ans;
95cf     }
427e
f4bd     void add(int n, LL x){
968e         while (n < N) { tr[n] += x; n += n & -n; }
95cf     }
329b };

```

6.2 Fenwick tree (range update point query)

```

3d03 struct bit_rupq{ // range update, point query
d7af     int N;
99ff     vector<LL> tr;
427e
2d99     void init(int n) { tr.assign(N = n + 5, 0); }
427e
38d4     LL query(int n) {
f7ff         LL ans = 0;
3667         while (n < N) { ans += tr[n]; n += n & -n; }
4206         return ans;

```

6.3 Segment tree

```

LL p;
const int MAXN = 4 * 100006;
struct segtree {
    int l[MAXN], m[MAXN], r[MAXN];
    LL val[MAXN], tadd[MAXN], tmul[MAXN];

#define lson (o<<1)
#define rson (o<<1|1)

    void pull(int o) {
        val[o] = (val[lson] + val[rson]) % p;
    }

    void push_add(int o, LL x) {
        val[o] = (val[o] + x * (r[o] - l[o])) % p;
        tadd[o] = (tadd[o] + x) % p;
    }

    void push_mul(int o, LL x) {
        val[o] = val[o] * x % p;
        tadd[o] = tadd[o] * x % p;
        tmul[o] = tmul[o] * x % p;
    }

    void push(int o) {
        if (l[o] == m[o]) return;
        if (tmul[o] != 1) {
            push_mul(lson, tmul[o]);
            push_mul(rson, tmul[o]);
            tmul[o] = 1;
        }
        if (tadd[o]) {
            push_add(lson, tadd[o]);

```

```

3942
1ebb
451a
27be
4510
427e
ac35
1294
427e
1344
bbe9
95cf
427e
e4bc
5dd6
6eff
95cf
427e
d658
b82c
aa86
649f
95cf
427e
b149
3159
0a90
0f4a
045e
ac0a
95cf
1b82
9547

```

```

0e73     push_add(rson, tadd[o]);
6234     tadd[o] = 0;
95cf     }
95cf     }
427e
471c void build(int o, int ll, int rr) {
0e87     int mm = (ll + rr) / 2;
9d27     l[o] = ll; r[o] = rr; m[o] = mm;
ac0a     tmul[o] = 1;
5c92     if (ll == mm) {
001f         scanf("%lld", val + o);
e5b6         val[o] %= p;
8e2e     } else {
7293         build(lson, ll, mm);
5e67         build(rson, mm, rr);
ba26         pull(o);
95cf     }
95cf }
427e
4406 void add(int o, int ll, int rr, LL x) {
3c16     if (ll <= l[o] && r[o] <= rr) {
db32         push_add(o, x);
8e2e     } else {
c4b0         push(o);
4305         if (m[o] > ll) add(lson, ll, rr, x);
d5a6         if (m[o] < rr) add(rson, ll, rr, x);
ba26         pull(o);
95cf     }
95cf }
427e
48cd void mul(int o, int ll, int rr, LL x) {
3c16     if (ll <= l[o] && r[o] <= rr) {
e7d0         push_mul(o, x);
8e2e     } else {
c4b0         push(o);
d1ba         if (ll < m[o]) mul(lson, ll, rr, x);
67f3         if (m[o] < rr) mul(rson, ll, rr, x);
ba26         pull(o);
95cf     }
95cf }
427e
0f62 LL query(int o, int ll, int rr) {
3c16     if (ll <= l[o] && r[o] <= rr) {
6dfe         return val[o];

```

```

    } else {
        push(o);
        if (rr <= m[o]) return query(lson, ll, rr);
        if (ll >= m[o]) return query(rson, ll, rr);
        return query(lson, ll, rr) + query(rson, ll, rr);
    }
}
}
} seg;

```

```

8e2e
c4b0
462a
5cca
bbf9
95cf
95cf
4d99

```

6.4 Mo's algorithm

All intervals are closed on both sides. When running functions `enter()` and `leave()`, the global `l` and `r` has not changed yet. **Assume the data structure is initialized for empty interval.**

Usage:

<code>add_query(id, l, r)</code>	Add id-th query $[l, r]$.
<code>run()</code>	Run Mo's algorithm.
<code>yield(id)</code>	TODO. Yield answer for id-th query.
<code>enter(o)</code>	TODO. Add o-th element.
<code>leave(o)</code>	TODO. Remove o-th element.

```

constexpr int BLOCK_SZ = 300;

struct query { int l, r, id; };
vector<query> queries;

void add_query(int id, int l, int r) {
    queries.push_back(query{l, r, id});
}

int l, r;

// ----- functions to implement -----
inline void yield(int id);
inline void enter(int o);
inline void leave(int o);

void run() {
    if (queries.empty()) return;
    sort(range(queries), [](query lhs, query rhs) {
        int lb = lhs.l / BLOCK_SZ, rb = rhs.l / BLOCK_SZ;
        if (lb != rb) return lb < rb;
        return lhs.r < rhs.r;
    });
}

```

```

5194
427e
3ec4
d26a
427e
1e30
54c9
95cf
427e
9f6b
427e
427e
50e1
b20d
13af
427e
37f0
ab0b
8508
c7f8
03e7
0780

```

```

b251     });
6196     l = queries[0].l;
9644     r = queries[0].r;
38e6     for (int i = l; i <= r; i++) enter(i);
5bc9     for (query q : queries) {
f422         while (l > q.l) enter(--l);
39fb         while (r < q.r) enter(++r);
46b3         while (l < q.l) leave(l++);
6234         while (r > q.r) leave(r--);
82f5         yield(q.id);
95cf     }
95cf }

```

6.5 Mo's algorithm on tree

Numbers of vertices are 1-based. Implement `deal(int u)` and `query::yield()`.

```

ed86 const int MAXN = 200005, BLOCK = 300;
35b8 int n, m;
0b32 vector<int> adj[MAXN];
a292 int en[MAXN], edx;
ebcd int dep[MAXN], fa[MAXN];
7744 bool in[MAXN];
427e
e1b1 inline void deal(int u) {
c672     if (in[u] ^= 1) {
427e         // enter
8e2e     } else {
427e         // leave
95cf     }
95cf }
427e
6c2e void moveto(int a, int b) {
e53f     if (a == b) return;
460b     int cross = in[b] ? b : 0;
ebc8     auto moveup = [&] (int &x) {
139d         if (!cross) {
ad52             if (in[x] and !in[fa[x]]) cross = x;
ed4e             else if (in[fa[x]] and !in[x]) cross = fa[x];
95cf         }
82fb         deal(x); x = fa[x];
329b     };
893a     while (dep[a] > dep[b]) moveup(a);

```

```

while (dep[b] > dep[a]) moveup(b);
while (a != b) moveup(a), moveup(b);
deal(a); if (cross) deal(cross);
}

```

```

void dfs(int u, int p) {
    en[u] = edx++; fa[u] = p;
    for (int v : adj[u]) if (v != p) {
        dep[v] = dep[u] + 1;
        dfs(v, u); edx++;
    }
}

```

```

struct query {
    int l, r, id;
    void yield() { /* TODO */}
};
vector<query> qs;

```

```

void run() {
    dfs(1, 0);

    sort(range(qs), [] (query lhs, query rhs) {
        int u0 = en[lhs.l], v0 = en[rhs.l];
        int b1 = u0 / BLOCK, br = v0 / BLOCK;
        if (b1 != br) return b1 < br;
        int u1 = en[lhs.r], v1 = en[rhs.r];
        return b1 & 1 ? u1 < v1 : u1 > v1;
    });

    int l = 1, r = 1; deal(1);
    for (auto& q : qs) {
        moveto(l, q.l); l = q.l;
        moveto(r, q.r); r = q.r;
        q.yield();
    }
}

```

6.6 Treap

Self-balanced binary search tree which supports split and merge.

Usage:

b334
9d99
d1d9
95cf
427e
e1a2
b00c
79e0
bbda
f624
95cf
95cf
427e
457a
7551
fa1f
329b
6b35
427e
37f0
99d6
427e
199c
28dc
adcc
6fbd
708c
ae17
b251
427e
5314
8b5c
09d4
ce55
1412
95cf
95cf

<code>push(x)</code>	Push lazy tags to children.
<code>pull(x)</code>	Update statistics of node x .
<code>Init(x, v)</code>	Initialize node x with value v .
<code>Add(x, v)</code>	Apply addition to subtree x .
<code>Reverse(x)</code>	Apply reversion to subtree x .
<code>Merge(x, y)</code>	Merge trees rooted at x and y . Return the root of new tree.
<code>Split(t, k, x, y)</code>	Split out the left k elements of tree t . The roots of left part and right part are stored in x and y , respectively.
<code>init(n)</code>	Initialize the treap with array of size n .
<code>work(op, l, r)</code>	Range operation over $[l, r)$.
Time Complexity: Expected $O(\log n)$ per operation.	

```

const int MAXN = 200005;
mt19937 gen(time(NULL));

struct Treap {
    int ch[MAXN][2];
    int sz[MAXN], key[MAXN], val[MAXN];
    int add[MAXN], rev[MAXN];
    LL sum[MAXN] = {0};
    int maxv[MAXN] = {INT_MIN}, minv[MAXN] = {INT_MAX};

    void Init(int x, int v) {
        ch[x][0] = ch[x][1] = 0;
        key[x] = gen(); val[x] = v; pull(x);
    }

    void pull(int x) {
        sz[x] = 1 + sz[ch[x][0]] + sz[ch[x][1]];
        sum[x] = val[x] + sum[ch[x][0]] + sum[ch[x][1]];
        maxv[x] = max({val[x], maxv[ch[x][0]], maxv[ch[x][1]]});
        minv[x] = min({val[x], minv[ch[x][0]], minv[ch[x][1]]});
    }

    void Add(int x, int a) {
        val[x] += a; add[x] += a;
        sum[x] += LL(sz[x]) * a; maxv[x] += a; minv[x] += a;
    }

    void Reverse(int x) {
        rev[x] ^= 1;
        swap(ch[x][0], ch[x][1]);
    }
}

```

```

void push(int x) {
    for (int c : ch[x]) if (c) {
        Add(c, add[x]);
        if (rev[x]) Reverse(c);
    }
    add[x] = 0; rev[x] = 0;
}

int Merge(int x, int y) {
    if (!x || !y) return x | y;
    push(x); push(y);
    if (key[x] > key[y]) {
        ch[x][1] = Merge(ch[x][1], y); pull(x); return x;
    } else {
        ch[y][0] = Merge(x, ch[y][0]); pull(y); return y;
    }
}

void Split(int t, int k, int &x, int &y) {
    if (t == 0) { x = y = 0; return; }
    push(t);
    if (sz[ch[t][0]] < k) {
        x = t; Split(ch[t][1], k - sz[ch[t][0]] - 1, ch[t][1], y);
    } else {
        y = t; Split(ch[t][0], k, x, ch[t][0]);
    }
    if (x) pull(x); if (y) pull(y);
}

} treap;

int root;

void init(int n) {
    Rep (i, n) {
        int x; scanf("%d", &x);
        treap.Init(i, x);
        root = (i == 1) ? 1 : treap.Merge(root, i);
    }
}

void work(int op, int l, int r) {
    int tl, tm, tr;
    treap.Split(root, l, tl, tm);
    treap.Split(tm, r - l, tm, tr);
}

```

1a53
5fe5
fd76
7a53
95cf
49ee
95cf
427e
9d2c
1b09
cd7e
bffa
a3df
8e2e
bf9e
95cf
95cf
427e
dc7e
6303
f26b
3465
ffdf
8e2e
8a23
95cf
89e3
95cf
b1f4
427e
24b6
427e
d34f
34d7
7681
0ed8
bcc8
95cf
95cf
427e
d030
6639
b6c4
8de3

```

3658     if (op == 1) {
c039         int x; scanf("%d", &x); treap.Add(tm, x);
1dcb     } else if (op == 2) {
ae78         treap.Reverse(tm);
581d     } else if (op == 3) {
e092         printf("%lld_%d_%d\n",
867f             treap.sum[tm], treap.minv[tm], treap.maxv[tm]);
95cf     }
6188     root = treap.Merge(treap.Merge(tl, tm), tr);
95cf }

```

6.7 Link/cut tree

Dynamic connectivity of undirected acyclic graph. Support single-vertex update, path aggregation and relative LCA query. Vertices are numbered from 1. Zero initialization is enough except for the statistic information.

Usage:

<code>pull(x)</code>	Update statistics of node x .
<code>Root(u)</code>	Get the root of tree where vertex u is in.
<code>Link(u, v)</code>	Link two unconnected trees.
<code>Cut(u, v)</code>	Cut an existent edge.
<code>Query(u, v)</code>	Path aggregation.
<code>Update(u, x)</code>	Single point modification.
<code>LCA(u, v, root)</code>	Get the lowest common ancestor of u and v in tree rooted at root.

Time Complexity: $O(\log n)$ per operation

```

2e73 const int MAXN = 1000005;
ca06 struct LCT {
6a6d     int fa[MAXN], ch[MAXN][2], val[MAXN], sum[MAXN];
c6e1     bool rev[MAXN];
427e
eba3     bool isroot(int x) { return ch[fa[x]][0] == x || ch[fa[x]][1] == x; }
f19f     void pull(int x) { sum[x] = val[x] ^ sum[ch[x][0]] ^ sum[ch[x][1]]; }
1c4d     void reverse(int x) { swap(ch[x][0], ch[x][1]); rev[x] ^= 1; }
1a53     void push(int x) {
89a0         if (rev[x]) rep (i, 2) if (ch[x][i]) reverse(ch[x][i]); rev[x] = 0;
95cf     }
425f     void rotate(int x) {
51af         int y = fa[x], z = fa[y], k = ch[y][1] == x, w = ch[x][!k];
e1fe         if (isroot(y)) ch[z][ch[z][1] == y] = x;
1e6f         ch[x][!k] = y; ch[y][k] = w; if (w) fa[w] = y;

```

```

        fa[y] = x; fa[x] = z; pull(y);
    }
    void pushall(int x) { if (isroot(x)) pushall(fa[x]); push(x); }
    void splay(int x) {
        int y = x, z = 0;
        for (pushall(y); isroot(x); rotate(x)) {
            y = fa[x]; z = fa[y];
            if (isroot(y)) rotate((ch[y][0] == x) ^ (ch[z][0] == y) ? x : y);
        }
        pull(x);
    }
    void access(int x) {
        int z = x;
        for (int y = 0; x; x = fa[y = x]) { splay(x); ch[x][1] = y; pull(x); }
        splay(z);
    }
    void chroot(int x) { access(x); reverse(x); }
    void split(int x, int y) { chroot(x); access(y); }

    int Root(int x) {
        for (access(x); ch[x][0]; x = ch[x][0]) push(x);
        splay(x); return x;
    }
    void Link(int u, int v) { chroot(u); fa[u] = v; }
    void Cut(int u, int v) { split(u, v); fa[u] = ch[v][0] = 0; pull(v); }
    int Query(int u, int v) { split(u, v); return sum[v]; }
    void Update(int u, int x) { splay(u); val[u] = x; }
    int LCA(int x, int y, int root) {
        chroot(root); access(x); splay(y);
        while (fa[y]) splay(y = fa[y]);
        return y;
    }
};

```

6d09
95cf
52c6
f69c
d095
c494
ceef
4449
95cf
78a0
95cf
6229
1548
8854
7afd
95cf
a067
126d
427e
d87a
f4f1
0d77
95cf
9e46
7c10
0691
a999
1f42
6cb2
02e5
c218
95cf
329b

6.8 Balanced binary search tree from pb_ds

```

#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;

tree<int, null_type, less<int>, rb_tree_tag, tree_order_statistics_node_update>
    rkt;
// null_tree_node_update

```

0475
332d
427e
43a7
427e

```

427e // SAMPLE USAGE
190e rkt.insert(x);      // insert element
05d4 rkt.erase(x);      // erase element
add5 rkt.order_of_key(x); // obtain the number of elements less than x
b064 rkt.find_by_order(i); // iterator to i-th (numbered from 0) smallest element
c103 rkt.lower_bound(x);
4ff4 rkt.upper_bound(x);
b19b rkt.join(rkt2);    // merge tree (only if their ranges do not intersect)
cb47 rkt.split(x, rkt2); // split all elements greater than x to rkt2

```

6.9 Persistent segment tree, range k-th query

```

f1a7 struct node {
2ff6     static int n, pos;
427e
7cec     int value;
70e2     node *left, *right;
427e
20b0     void* operator new(size_t size);
427e
3dc0     static node* Build(int l, int r) {
b6c5         node* a = new node;
ce96         if (r > l + 1) {
181e             int mid = (l + r) / 2;
3ba2             a->left = Build(l, mid);
8aaf             a->right = Build(mid, r);
8e2e         } else {
bfc4             a->value = 0;
95cf         }
5ffd         return a;
95cf     }
427e
5a45     static node* init(int size) {
2c46         n = size;
7ee3         pos = 0;
be52         return Build(0, n);
95cf     }
427e
93c0     static int Query(node* lt, node *rt, int l, int r, int k) {
d30c         if (r == l + 1) return l;
181e         int mid = (l + r) / 2;

```

```

if (rt->left->value - lt->left->value < k) {
    k -= rt->left->value - lt->left->value;
    return Query(lt->right, rt->right, mid, r, k);
} else {
    return Query(lt->left, rt->left, l, mid, k);
}
}

```

```

static int query(node* lt, node *rt, int k) {
    return Query(lt, rt, 0, n, k);
}

```

```

node *Inc(int l, int r, int pos) const {
    node* a = new node(*this);
    if (r > l + 1) {
        int mid = (l + r) / 2;
        if (pos < mid)
            a->left = left->Inc(l, mid, pos);
        else
            a->right = right->Inc(mid, r, pos);
    }
    a->value++;
    return a;
}

```

```

node *inc(int index) {
    return Inc(0, n, index);
}
nodes[8000000];

```

```

int node::n, node::pos;
inline void* node::operator new(size_t size) {
    return nodes + (pos++);
}

```

6.10 Block list

All indices are 0-based. All ranges are left-closed right-open.

Usage:

cb5a
8edb
2412
8e2e
0119
95cf
95cf
427e
c9ad
9e27
95cf
427e
b19c
5794
ce96
181e
203d
f44a
649a
1024
95cf
2b3e
5ffd
95cf
427e
e80f
c246
95cf
865a
427e
99ce
1987
bb3c
95cf

block::fix() Apply tags to the current block.
 Init(l, r) Range initializer.
 Reverse(l, r) Reverse the range.
 Add(l, r, x) Add x to the range.
 Query(l, r) Range aggregation.

```

fd9e  const int BLOCK = 800;
76b3  typedef vector<int> vi;
427e
a771  struct block {
8fbc      vi data;
e3b5      LL sum; int minv, maxv;
41db      int add; bool rev;
427e
d7eb      block(vi&& vec) : data(move(vec)),
1f0c          sum(accumulate(range(data), 0ll)),
8216          minv(*min_element(range(data))),
527d          maxv(*max_element(range(data))),
6437          add(0), rev(0) { }
427e
b919      void fix() {
0694          if (rev) reverse(range(data));      rev = 0;
0527          if (add) for (int& x : data) x += add;  add = 0;
95cf      }
427e
8bc4      void merge(block& another) {
b895          fix(); another.fix();
f516          vi temp(move(data));
d02c          temp.insert(temp.end(), range(another.data));
88ea          *this = block(move(temp));
95cf      }
427e
42e8      block split(int pos) {
3e79          fix();
ccab          block result(vi(data.begin() + pos, data.end()));
861a          data.resize(pos); *this = block(move(data));
56b0          return result;
95cf      }
329b  };
427e
2a18  typedef list<block>::iterator lit;
427e
ce14  struct blocklist {
5540      list<block> blk;
  
```

```

void maintain() {
    lit it = blk.begin();
    while (it != blk.end() && next(it) != blk.end()) {
        lit it2 = it;
        while (next(it2) != blk.end() &&
            it2->data.size() + next(it2)->data.size() <= BLOCK) {
            it2->merge(*next(it2));
            blk.erase(next(it2));
        }
        ++it;
    }
}

lit split(int pos) {
    for (lit it = blk.begin(); ; it++) {
        if (pos == 0) return it;
        while (it->data.size() > pos)
            blk.insert(next(it), it->split(pos));
        pos -= it->data.size();
    }
}

void Init(int *l, int *r) {
    for (int *cur = l; cur < r; cur += BLOCK)
        blk.emplace_back(vi(cur, min(cur + BLOCK, r)));
}

void Reverse(int l, int r) {
    lit it = split(l), it2 = split(r);
    reverse(it, it2);
    while (it != it2) {
        it->rev ^= 1;
        it++;
    }
    maintain();
}

void Add(int l, int r, int x) {
    lit it = split(l), it2 = split(r);
    while (it != it2) {
        it->sum += LL(x) * it->data.size();
        it->minv += x; it->maxv += x;
    }
}
  
```

427e
 7b8e
 3131
 4628
 852d
 188c
 3600
 93e1
 e1fa
 95cf
 5771
 95cf
 95cf
 427e
 b7b3
 2273
 5502
 8e85
 2099
 a5a1
 427e
 95cf
 95cf
 427e
 1c7b
 9919
 8950
 95cf
 427e
 a22f
 997b
 dfd0
 8f89
 6a06
 5283
 95cf
 b204
 95cf
 427e
 3cce
 997b
 8f89
 e927
 03d3

```

4511         it->add += x; it++;
95cf     }
b204     maintain();
95cf }
427e
3ad3 void Query(int l, int r) {
997b     lit it = split(l), it2 = split(r);
c33d     LL sum = 0; int minv = INT_MAX, maxv = INT_MIN;
8f89     while (it != it2) {
e472         sum += it->sum;
72c4         minv = min(minv, it->minv);
e1c4         maxv = max(maxv, it->maxv);
5283         it++;
95cf     }
b204     maintain();
8792     printf("%lld_%d_%d\n", sum, minv, maxv);
95cf }
958e } lst;

```

6.11 Persistent block list

Block list that supports persistence. All indices are 0-based. All ranges are left-closed right-open. `std::shared_ptr` is used to ease memory management. One should modify the constructor of `block` to maintain extra information. Here we use this policy that the size of each block does not exceed `BLOCK`, while the sum of sizes of two adjacent blocks does not less than `BLOCK`.

When some operation that breaks block list property, please call `maintain` in time to restore the property.

Usage:

<code>maintain()</code>	Maintain the block list property.
<code>split(pos)</code>	Split the block list at position <code>pos</code> . Returns an iterator to a block starting at <code>pos</code> .
<code>sum(l, r)</code>	An example function of list traversal between $[l, r)$.

Time Complexity: When `BLOCK` is properly selected, the time complexity is $O(\sqrt{n})$ per operation.

```

a19e constexpr int BLOCK = 800;
76b3 typedef vector<int> vi;
0563 typedef shared_ptr<vi> pvi;
013b typedef shared_ptr<const vi> pcvi;
427e
a771 struct block {

```

```

pcvi data;
LL sum;

// add information to maintain
block(pcvi ptr) :
    data(ptr),
    sum(accumulate(ptr->begin(), ptr->end(), 0ll))
{ }

void merge(const block& another) {
    pvi temp = make_shared<vi>(data->begin(), data->end());
    temp->insert(temp->end(), another.data->begin(), another.data->end());
    *this = block(temp);
}

block split(int pos) {
    block result(make_shared<vi>(data->begin() + pos, data->end()));
    *this = block(make_shared<vi>(data->begin(), data->begin() + pos));
    return result;
}

};

typedef list<block>::iterator lit;

struct blocklist {
    list<block> blk;

    void maintain() {
        lit it = blk.begin();
        while (it != blk.end() and next(it) != blk.end()) {
            lit it2 = it;
            while (next(it2) != blk.end() and
                    it2->data->size() + next(it2)->data->size() <= BLOCK) {
                it2->merge(*next(it2));
                blk.erase(next(it2));
            }
            ++it;
        }
    }

    lit split(int pos) {
        for (lit it = blk.begin(); ; it++) {
            if (pos == 0) return it;
            while (it->data->size() > pos) {

```

2989
8fd0
427e
427e
a613
24b5
0cf0
e93b
427e
5c0f
0b18
ac21
6467
95cf
427e
42e8
dac1
01db
56b0
95cf
329b
427e
2a18
427e
ce14
5540
427e
7b8e
3131
5e44
852d
0b03
029f
93e1
e1fa
95cf
5771
95cf
95cf
427e
b7b3
2273
5502
d480

```

2099         blk.insert(next(it), it->split(pos));
95cf     }
a1c8     pos -= it->data->size();
95cf     }
95cf     }
427e
fd38     LL sum(int l, int r) { // traverse
48b4         lit it1 = split(l), it2 = split(r);
ac09         LL res = 0;
9f1d         while (it1 != it2) {
8284             res += it1->sum;
61fd             it1++;
95cf         }
b204         maintain();
244d         return res;
95cf     }
329b };

```

6.12 Sparse table, range minimum query

The array is 0-based and the range is left-closed right-open.

```

db63 const int MAXN = 100007;
cefd int a[MAXN], st[MAXN][30];
427e
d34f void init(int n){
c73d     int l = log2(n);
cf75     rep (i, n) st[i][0] = a[i];
426b     rep (j, l) rep (i, 1+n-(1<<j))
1131         st[i][j+1] = min(st[i][j], st[i+(1<<j)][j]);
95cf }
427e
c863 int rmq(int l, int r){
f089     int k = log2(r - l);
6117     return min(st[l][k], st[r-(1<<k)][k]);
95cf }

```

7 Geometrics

7.1 2D geometric template

```

#include <bits/stdc++.h>
using namespace std;

typedef int T;
typedef struct pt {
    T x, y;
    T operator , (pt a) { return x*a.x + y*a.y; } // inner product
    T operator * (pt a) { return x*a.y - y*a.x; } // outer product
    pt operator + (pt a) { return {x+a.x, y+a.y}; }
    pt operator - (pt a) { return {x-a.x, y-a.y}; }

    pt operator * (T k) { return {x*k, y*k}; }
    pt operator - () { return {-x, -y}; }
} vec;

typedef pair<pt, pt> seg;

bool ptOnSeg(pt& p, seg& s){
    vec v1 = s.first - p, v2 = s.second - p;
    return (v1, v2) <= 0 && v1 * v2 == 0;
}

// 0 not on segment
// 1 on segment except vertices
// 2 on vertices
int ptOnSeg2(pt& p, seg& s){
    vec v1 = s.first - p, v2 = s.second - p;
    T ip = (v1, v2);
    if (v1 * v2 != 0 || ip > 0) return 0;
    return (v1, v2) ? 1 : 2;
}

// if two orthogonal rectangles do not touch, return true
inline bool nIntRectRect(seg a, seg b){
    return min(a.first.x, a.second.x) > max(b.first.x, b.second.x) ||
           min(a.first.y, a.second.y) > max(b.first.y, b.second.y) ||
           min(b.first.x, b.second.x) > max(a.first.x, a.second.x) ||
           min(b.first.y, b.second.y) > max(a.first.y, a.second.y);
}

// >0 in order
// <0 out of order
// =0 not standard

```

```

7538 inline double rotOrder(vec a, vec b, vec c){return double(a*b)*(b*c);}
427e
31ed inline bool intersect(seg a, seg b){
427e // ! if (nIntRectRect(a, b)) return false; // if commented, assume that a
      and b are non-collinear
cb52 return rotOrder(b.first-a.first, a.second-a.first, b.second-a.first) >= 0 &&
059e rotOrder(a.first-b.first, b.second-b.first, a.second-b.first) >= 0;
95cf }
427e
427e // 0 not intersect
427e // 1 standard intersection
427e // 2 vertex-line intersection
427e // 3 vertex-vertex intersection
427e // 4 collinear and have common point(s)
4d19 int intersect2(seg& a, seg& b){
5dc4 if (nIntRectRect(a, b)) return 0;
42c0 vec va = a.second - a.first, vb = b.second - b.first;
2096 double j1 = rotOrder(b.first-a.first, va, b.second-a.first),
72fe j2 = rotOrder(a.first-b.first, vb, a.second-b.first);
5ac6 if (j1 < 0 || j2 < 0) return 0;
9400 if (j1 != 0 && j2 != 0) return 1;
83db if (j1 == 0 && j2 == 0){
6b0c if (va * vb == 0) return 4; else return 3;
fb17 } else return 2;
95cf }
427e
2c68 template <typename Tp = T>
5894 inline pt getIntersection(pt P, vec v, pt Q, vec w){
6850 static_assert(is_same<Tp, double>::value, "must_be_double!");
7c9a return P + v * (w*(P-Q)/(v*w));
95cf }
427e
427e // -1 outside the polygon
427e // 0 on the border of the polygon
427e // 1 inside the polygon
cbdd int ptOnPoly(pt p, pt* poly, int n){
5fb4 int wn = 0;
1294 for (int i = 0; i < n; i++) {
427e
3cae T k, d1 = poly[i].y - p.y, d2 = poly[(i+1)%n].y - p.y;
b957 if (k = (poly[(i+1)%n] - poly[i])*(p - poly[i])){
8c40 if (k > 0 && d1 <= 0 && d2 > 0) wn++;
3c4d if (k < 0 && d2 <= 0 && d1 > 0) wn--;
aad3 } else return 0;

```

```

    }
    return wn ? 1 : -1;
}

istream& operator >> (istream& lhs, pt& rhs){
    lhs >> rhs.x >> rhs.y;
    return lhs;
}

istream& operator >> (istream& lhs, seg& rhs){
    lhs >> rhs.first >> rhs.second;
    return lhs;
}

```

```

95cf
0a5f
95cf
427e
d4a3
fa86
331a
95cf
427e
07ae
5cab
331a
95cf

```

8 Appendices

8.1 Primes

8.1.1 First primes

p	$g(p)$	p	$g(p)$	p	$g(p)$	p	$g(p)$	p	$g(p)$
2	1	3	2	5	2	7	3	11	2
13	2	17	3	19	2	23	5	29	2
31	3	37	2	41	6	43	3	47	5
53	2	59	2	61	2	67	2	71	7
73	5	79	3	83	2	89	3	97	5
101	2	103	5	107	2	109	6	113	3
127	3	131	2	137	3	139	2	149	2
151	6	157	5	163	2	167	5	173	2
179	2	181	2	191	19	193	5	197	2
199	3	211	2	223	3	227	2	229	6

8.1.2 Arbitrary length primes

$\lg p$	p	$g(p)$	p	$g(p)$
3	967	5	1031	14
4	9859	2	10273	10
5	96331	10	102931	3
6	958543	6	1031137	5
7	9594539	2	10169651	2
8	96243449	3	103211039	7
9	980483981	2	1042484357	2
10	9858935453	2	10261276009	7
11	95748666809	3	101759940101	2
12	950781833849	3	1012797784423	5
13	9739822952371	7	10037217092377	7
14	96181051140397	5	104974966380359	11
15	981030138360889	13	1029038416465403	2
16	9655206098080843	3	10116299875820773	2
17	97687777921994419	3	101506415998163437	2

8.1.3 $\sim 1 \times 10^9$

p	$g(p)$	p	$g(p)$	p	$g(p)$
954854573	3	967607731	2	973215833	3
975831713	3	978949117	2	980766497	3
983879921	3	985918807	3	986608921	29
991136977	5	991752599	13	997137961	11
1003911991	3	1009775293	2	1012423549	6
1021000537	5	1023976897	7	1024153643	2
1037027287	3	1038812881	11	1044754639	3
1045125617	3	1047411427	3	1047753349	6

8.1.4 $\sim 1 \times 10^{18}$

p	$g(p)$	p	$g(p)$
951970612352230049	3	963284339889659609	3
967495386904694119	3	969751761517096213	2
983238274281901499	2	984647442475101409	23
989286107138674069	11	1002507954383424641	3
1006658951440146419	2	1020152326159075903	3
1034876265966119449	7	1042753851435034019	2
1043609016597371563	2	1045571042176595707	2
1048364250160580293	2	1049495624119026949	2

8.2 Pell's equation

$x^2 - ny^2 = 1$, where n is a positive nonsquare integer.

Let (x_0, y_0) be the smallest positive solution of the equation, then the k -th solution is:

$$\begin{pmatrix} x_k \\ y_k \end{pmatrix} = \begin{pmatrix} x_0 & ny_0 \\ y_0 & x_0 \end{pmatrix}^k \begin{pmatrix} x_0 \\ y_0 \end{pmatrix}$$

Some smallest solutions to Pell's equation:

n	2	3	5	6	7	8	10	11	12	13	14	15	17	18	19	20
x	3	2	9	5	8	3	19	10	7	649	15	4	33	17	170	9
y	2	1	4	2	3	1	6	3	2	180	4	1	8	4	39	2

8.3 Burnside's lemma and Polya's enumeration theorem

The Burnside's lemma says that

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

where G is a group acting on X , X^g is the set of elements in X that are fixed by g , i.e. $X^g = \{x \in X : gx = x\}$.

The unweighted version of Pólya enumeration theorem says that

$$|Y^X/G| = \frac{1}{|G|} \sum_{g \in G} m^{c_g}$$

where $m = |X|$ is the number of colors, c_g is the number of the cycles of permutation g .

8.4 Supnick TSP

Given f and $x_1 \leq x_2 \leq \dots \leq x_n$, if f is Supnick, then

$$\sum_{i=1}^n f(x_{\pi(i)}, x_{\pi(i+1)})$$

1. is minimized when $\pi = (1, 3, 5, 7, \dots, 8, 6, 4, 2)$.
2. is maximized when $\pi = (n, 2, n-2, 4, \dots, 5, n-3, 3, n-1, 1)$.

8.5 Lagrange's interpolation

For sample points $(x_0, y_0), \dots, (x_k, y_k)$, define

$$l_j(x) = \prod_{0 \leq m \leq k, m \neq j} \frac{x - x_m}{x_j - x_m}$$

then the Lagrange polynomial is

$$L(x) = \sum_{j=0}^k y_j l_j(x).$$

To use the script below, type two lines

```
x0 x1 x2 ... xn
y0 y1 y2 ... yn
```

the script will print the fractional coefficient of the polynomial in ascending exponent order.

```
#!/usr/bin/python2
from fractions import *

def polymul(a, b) :
    p = [0] * (len(a)+len(b)-1)
    for e1, c1 in enumerate(a) :
        for e2, c2 in enumerate(b) :
            p[e1+e2] += c1*c2
    return p

x, y = [map(Fraction, raw_input().split()) for _ in 0,0]
n = len(x)
lj = [reduce(polymul, [[-x[m]/(x[j]-x[m]), 1/(x[j]-x[m])]
    for m in range(n) if m != j]] for j in range(n)]
print '_'.join(map(str, map(sum, zip(*map(
    lambda a, b : [x*a for x in b], y, lj)))))
```

```
6dc9
4b2b
427e
796b
83e4
f697
156c
dfce
5849
427e
f06d
e80a
a649
9dfa
3cae
7c0d
```