

Coupling Towards The Past: Local Sampling from Markov Chains

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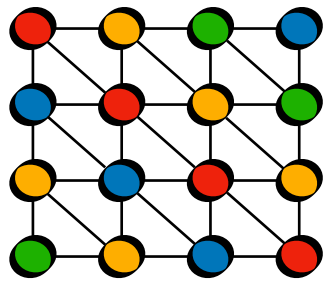
Based on joint works with Weiming Feng,² Heng Guo,³ Hongyang Liu,¹ Jingcheng Liu,¹ Jiaheng Wang,³ Yitong Yin¹ and Yixiao Yu¹

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Spin Systems and Gibbs sampling



graph $G = (V, E)$ $q \geq 2$ states

configuration $\sigma \in [q]^V$

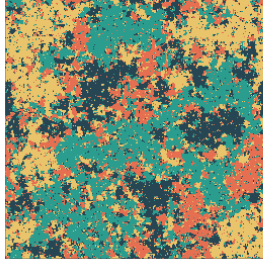
external fields $\lambda_v \in \mathbb{R}_{\geq 0}^q$ for each $v \in V$

interaction matrix $A_e \in \mathbb{R}_{\geq 0}^{q \times q}$ for each $e \in E$

weight: $w(\sigma) = \prod_{v \in V} \lambda_v(\sigma(v)) \prod_{(u,v) \in E} A_e(\sigma(u), \sigma(v))$

partition function: $Z = \sum_{\sigma \in [q]^V} w(\sigma)$

Gibbs distribution: $\mu(\sigma) = \frac{w(\sigma)}{Z}$



(Systematic Scan) Gibbs Sampling/Glauber Dynamics

start with arbitrary configuration σ with $w(\sigma) > 0$;

at each time $1 \leq t \leq T$:

pick the vertex $v = v_{t \bmod n}$; (assume $V = \{v_0, \dots, v_{n-1}\}$);

resample $\sigma_v \sim \mu_v(\cdot \mid \sigma_{V \setminus \{v\}})$;

return σ ;

converges to μ as $T \rightarrow \infty$!

Coupling Towards The Past (CTTP)

Idea: Imagine the chain runs from the infinite past to time 0, can we somehow deduce its final state (distributed as μ)?

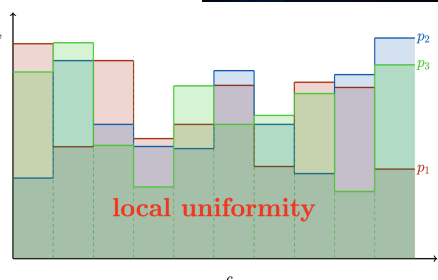


If only we can resolve an update without knowing the current configuration...

Local Uniformity (Marginal Lower Bound)

$\forall c \in [q], \mu_v^{\min}(c) \triangleq \min_{\mu(\sigma_{V \setminus \{v\}}) > 0} \mu_v(c \mid \sigma_{V \setminus \{v\}})$

θ -(locally) uniform: $\sum_{c \in [q]} \mu_v^{\min}(c) \geq \theta$

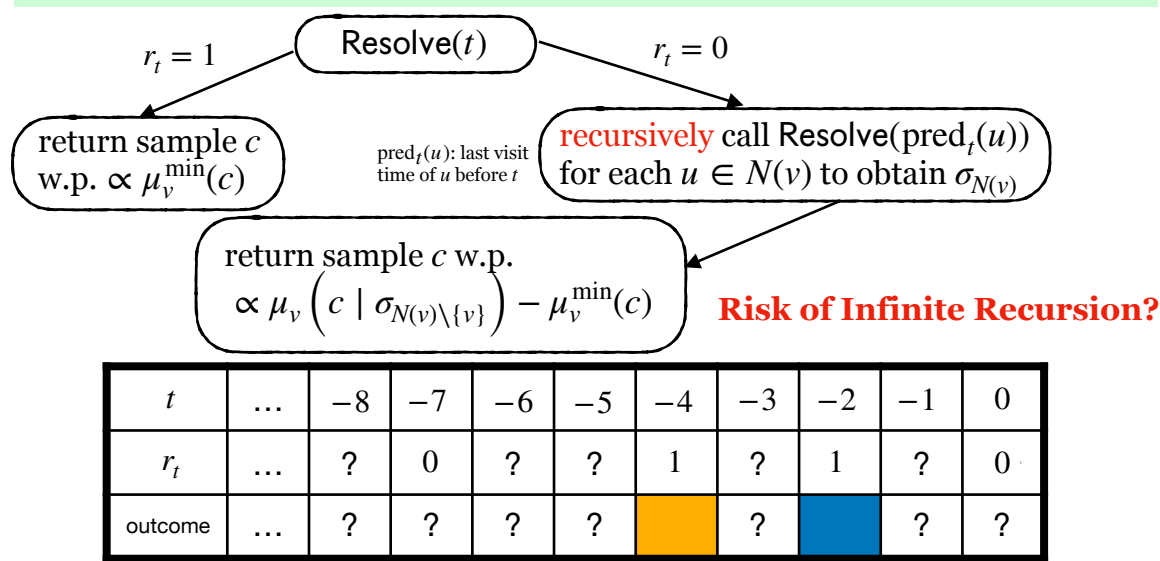


$\sigma_v \sim \mu_v(\cdot \mid \sigma_{V \setminus \{v\}}) \Rightarrow$ draw $r_t \sim \text{Ber}(\theta)$ and σ_v such that

$$\mu_v(\cdot \mid \sigma_{N(v) \setminus \{v\}}) \Pr[\sigma_v = c] \propto \begin{cases} \mu_v^{\min}(c) & r_t = 1 \\ \mu_v(c \mid \sigma_{N(v) \setminus \{v\}}) - \mu_v^{\min}(c) & r_t = 0 \end{cases}$$

With probability θ , an update can be directly resolved!

Otherwise, we need to know its neighbors' states to determine...



sufficient condition for termination: $(1 - \theta)\Delta \leq 1$ (Δ : maximum degree of graph)

a perfect (no bias) local (produces local samples within local time) sampler

a direct-sum style decomposition of Markov chains:

resolving a single update takes $O(1/n)$ time of learning the entire configuration

naming: hidden (default) grand coupling + backward deduction of states

Application: Deterministic Counting

CTTP often gives exponential tail bound: $\Pr[t_{\text{run}} \geq T] \leq \exp(-O(T))$

truncate up to $K = O_{\Delta} \left(\log \frac{n}{\epsilon} \right)$ random bits: ϵ -approximate marginals

brute force enumeration

→ efficient deterministic approximate counting matching MCMC bounds

also applies to problems with high-order constraints

Hypergraph Independent Sets (HIS)

Let $H = (V, \mathcal{E})$ be a hypergraph.

$S \subseteq V$ is a (weak) independent set if $S \cap e \neq e$ for all $e \in \mathcal{E}$.

We obtain optimal (on the exponent) deterministic approximate counting algorithms for k -uniform ($|e| = k$ for all $e \in \mathcal{E}$) HIS:

Hypergraph independent sets	Reference	Bound	Running time
Randomised counting / sampling	[BDK08, BDK06]	$\Delta \leq k - 2$	$\tilde{O}(n^2) / O(n \log n)$
	[HSZ19, QWZ22]	$\Delta \lesssim 2^{k/2}$	$\tilde{O}(n^2) / O(n \log n)$
Deterministic counting	[BGG ⁺ 19]	$\Delta \leq k$	$n^{O(\log(k\Delta))}$
	[JPV21b]	$\Delta \lesssim 2^{k/7}$	$n^{\text{poly}(k, \Delta)}$
	[HWY23]	$\Delta \lesssim 2^{k/5}$	$n^{\text{poly}(k, \Delta)}$
	Our result	$\Delta \lesssim 2^{k/2}$	$n^{\text{poly}(k, \Delta)}$
Hardness	[BGG ⁺ 19]	$\Delta \geq 5 \cdot 2^{k/2}$ assuming $\mathbf{P} \neq \mathbf{NP}$	

Application: Analytical Stability

CTTP provides a direct-sum style decomposition of the Gibbs measure.

lift to the complex plane

→ analytical stability of certain polynomials matching MCMC bounds

Hypergraph Independence Polynomial (HIP)

Let $H = (V, \mathcal{E})$ be a hypergraph and Ω collect its independent sets.

The (univariate) independence polynomial $Z_H: \mathbb{C} \rightarrow \mathbb{C}$ of H is given as:

$$Z_H(\lambda) = \sum_{S \subseteq \Omega} \lambda^{|S|}$$

Complex zeroes of $Z_H(\lambda)$ are often called Lee-Yang zeroes.

We obtain optimal (on the exponent) zero-free regions for k -uniform HIP:



Extending CTTP: Local Sampling near Criticality

The requirement of local uniformity may be restrictive for certain models.

improved grand coupling and deduction rules

→ efficient local samplers near criticality for Ising model and q -colorings

Instance	Tractable regimes		Our result (local sampling)
	Global sampling	Local sampling	
Ising model	$\beta \in \left(\frac{\Delta-2}{\Delta}, \frac{\Delta}{\Delta-2} \right)$	$\beta \in \left(1 - \frac{1}{\Theta(\Delta^2)}, 1 + \frac{1}{\Theta(\Delta^2)} \right)$	$\beta \in \left(\frac{\Delta-0.5}{\Delta}, \frac{\Delta}{\Delta-0.5} \right)$
q -colorings	$q \geq 1.809\Delta$	N/A	$q \geq 65\Delta$

- first local sampler for near-critical Ising model;
- first local sampler for q -coloring (also near-critical);
- perfect samplers;
- expected linear running time: $O(\Delta \cdot |\Lambda|)$ for Ising;
 $O(\Delta^2 q \cdot |\Lambda|)$ for q -colorings.

1. *Towards Derandomising Markov Chain Monte Carlo*. Weiming Feng, Heng Guo, Chunyang Wang, Jiaheng Wang, Yitong Yin. In **SICOMP '25** (preliminary version in **FOCS '23**).
2. *Phase Transitions via Complex Extensions of Markov Chains*. Jingcheng Liu, Chunyang Wang, Yitong Yin, Yixiao Yu. In **STOC '25**.
3. *Local Gibbs Sampling beyond Local Uniformity*. Hongyang Liu, Chunyang Wang, Yitong Yin. To appear in **SODA '26**.