Data Structures and Algorithms

Chapter 5

Learning Objectives

- Learn what a recursive method is
- Be able to implement a recursive method in Java
- Be able to analyze the runtime of a recursive method using asymptotic analysis

Recursion

- A recursive function is a function which is defined in terms of itself.
- A recursion, in programming, is a way of implementing repeated execution of statements (or a method), where a method invokes itself.

- Example: Factorial

$n! = 1$ if $n = 0$

$n * (n - 1)!$ if $n \geq 1$

Recursion
Factorial

- Java implementation

```
public static int factorial(int n) throws IllegalArgumentException {
  if (n < 0)
      throw new IllegalArgumentException( );   // argument must be
                                                // nonnegative
  else if (n == 0)
      return 1;                          // base case
  else
      return n * factorial(n-1);         // recursive case
}
```

Recursion
Factorial

- Recursion trace

factorial(4)

4 * factorial(3)

3 * factorial(2)

2 * factorial(1)

1 * factorial(0)

returns 1

returns 1*1 = 1

returns 2*1 = 2

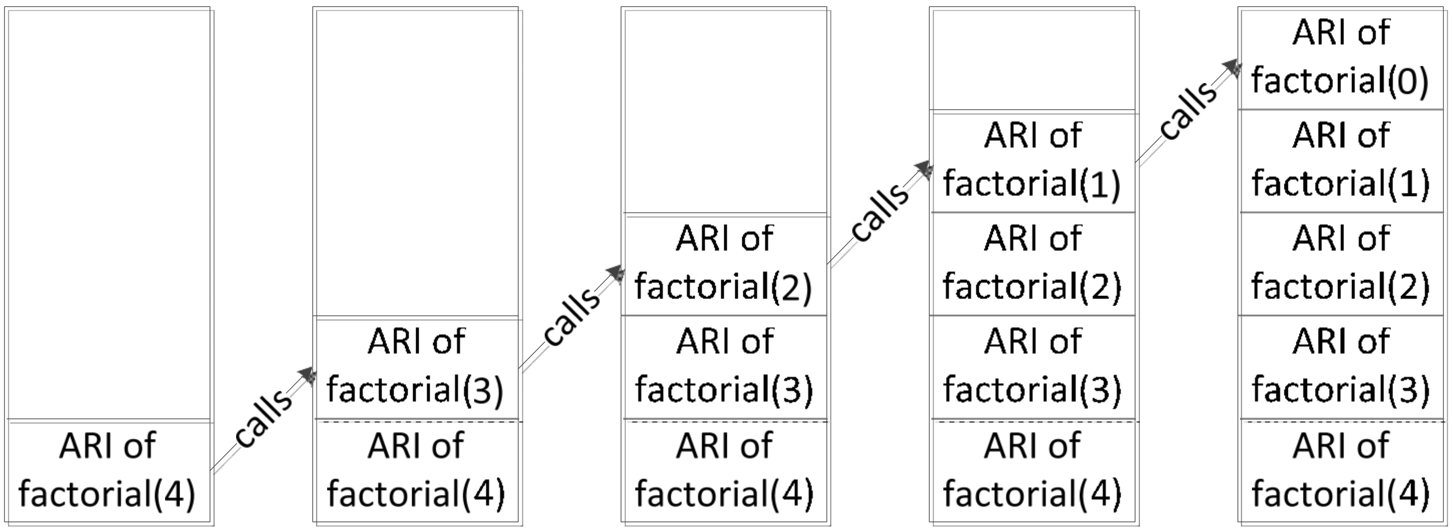
returns 3*2 = 6

returns 4*6 = 24

```
else if (n == 0)
   return 1;
else
   return n * factorial(n-1);
```

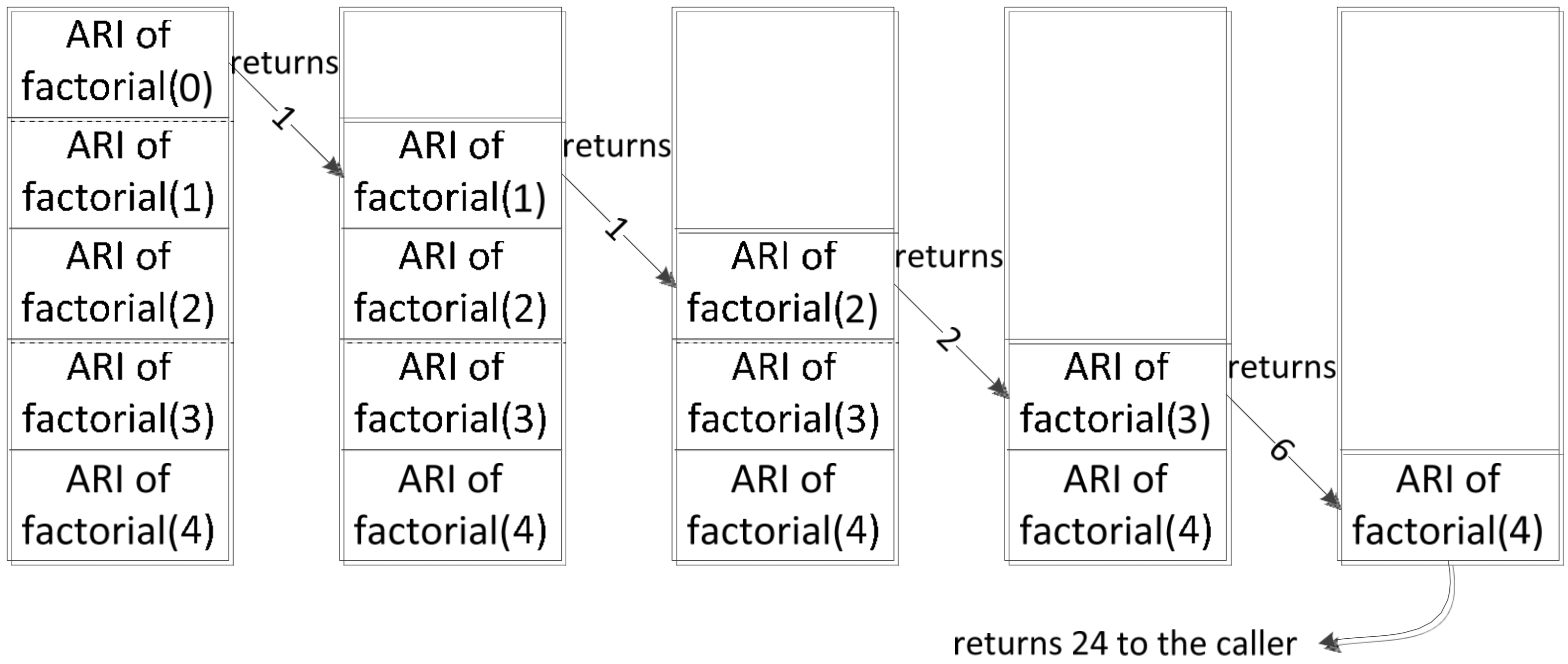
Recursion
Factorial

- Recursive calls



Recursion
Factorial

- Returning from calls



Recursion

Factorial

- Running time of *factorial*: $O(n)$

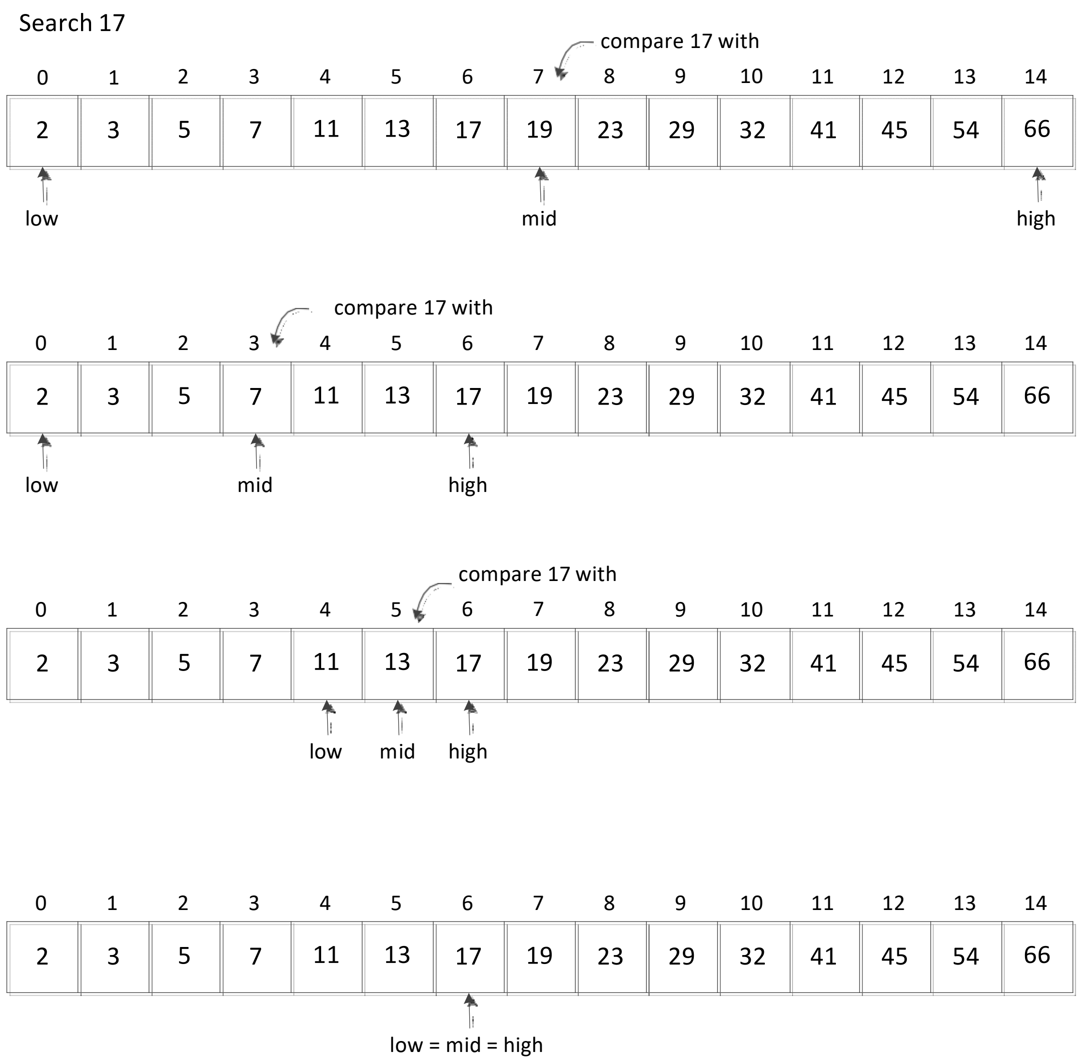
  - One execution of the method takes $O(1)$
  - It is invoked $n + 1$ times => $O(n)$
  - $O(1) \times O(n) = O(n)$

Recursion
Binary Search

- Search a sequence of n elements for a target element.
- Linear search
  - Examine each element while scanning the sequence
  - Best case: one comparison, or $O(1)$
  - Worst case: n comparisons, or $O(n)$
  - On average: $n/2$ comparisons, or $O(n)$
- Binary search
  - If the sequence is sorted, we can use binary search
  - Running time is $O(\log n)$

Recursion

Binary Search



Recursion
Binary Search

- Pseudocode

```
Algorithm binarySearch(int[ ] data, int target, int low, int high)
  If low > high                          // target is not found
     return false
  else
     mid = floor((low + high)/2) // median candidate
     if target = data[mid]              // target found
        return true
     else if target < data[mid]
        search data[low .. mid-1] recursively
     else
        search data[mid+1 .. high] recursively
```

Recursion
Binary Search

- Java implementation

```java
public static boolean binarySearch(int[ ] data, int target,
                                   int low, int high)
                                   {
if (low > high)
    return false;                  // interval empty; no match
else
{
    int mid = (low + high) / 2;
    if (target == data[mid])
        return true;               // found a match
    else if (target < data[mid]) // recurse left of the middle
        return binarySearch(data, target, low, mid - 1);
    else  // recurse right of the middle
        return binarySearch(data, target, mid + 1, high);
 }
}
```

Recursion
Binary Search

- Running time analysis
  - Execution of one call takes $O(1)$.
  - Each time binary search is (recursively) invoked, the number of elements to be searched is reduced to at most half.
  - Initially, there are n elements.
  - In the first recursive call, there are at most $n/2$ elements.
  - In the second recursive call, there are at most $n/4$ elements.
  - and so on …

Recursion
Binary Search

- Running time analysis (continued)
  - In the *j*-th recursive call, there are at most *n* / (2^j) elements.
  - In the worst case, the target is not in the sequence. In this case, recursion stops when there is no more elements to be searched.
  - The max. number of recursive calls is the smallest integer *r* such that $\frac{n}{2^r} < 1$
  - Or, *r* is the smallest integer such that *r* > log *n*
  - Therefore, $r = \lfloor \log n \rfloor + 1$
  - So, the total running time is $O(\log n)$

Recursion

More Examples

- Print array elements recursively - Pseudocode

```
Algorithm printArrayRecursively(data, i)
    if i = n, return
    else
        print data[i]
        i = i + 1
        printArrayRecursively(data, i)
```

Recursion
More Examples

- Print array elements recursively – Java code

```
public static void printArrayRecursive(int[ ] data, int i){
    if (i == data.length)
        return;
   else{
       System.out.print(data[i++] + " ");
      printArrayRecursive(data,i);
   }
}
```

Recursion

More Examples

- Reverse sequence recursively – Pseudocode

```
Algorithm reverseArray(data, low, high)
    if low >= high, return
    else
        swap data[low] with data[high]
        reverseArray(data, low+1, high-1)
```

Recursion
More Examples

- Reverse sequence recursively – Java code

```
public static void reverseArray(int[ ] data, int low,
                                int high) {
  if (low < high) {
    int temp = data[low];
    data[low] = data[high];
    data[high] = temp;
    reverseArray(data, low + 1, high - 1);
  }
}
```

Recursion

More Examples

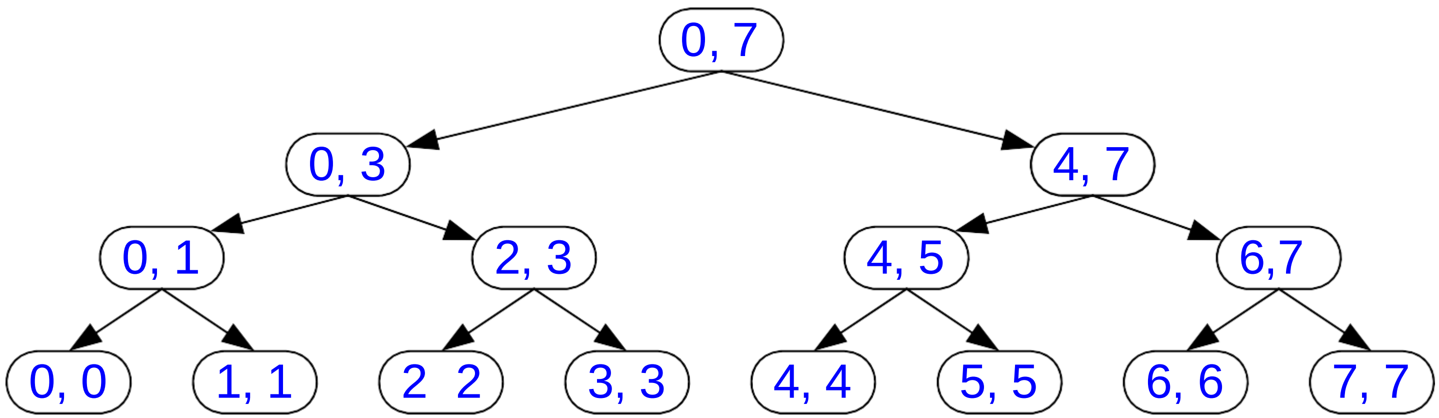
- Binary sum – Java code

```
public static int binarySum(int[ ] data, int low, int high) {
  if (low > high)              // zero elements in subarray
    return 0;
  else if (low == high)        // one element in subarray
    return data[low];
  else {
    int mid = (low + high) / 2;
    return binarySum(data, low, mid) +
                          binarySum(data, mid+1, high);
  }
}
```

Recursion

More Examples

- Binary sum – recursion trace



- Running time?

Recursion
Computing Powers

- Definition: $power(x, n) = x^n$
- Recursive definition

$$power(x, n) = \begin{cases} 1 & \text{if } n = 0 \\ x * power(x, n-1) & \text{otherwise} \end{cases}$$

- Direct implementation

```
public static double power (int x, int n)
{
    if(n==0)
       return 1;
    else{
        return x*power(x,n-1)
    }
}
```

- Execution of each method call takes $O(1)$.
- The method is invoked $(n + 1)$ times.
- Running time is $O(n)$

Recursion
Computing Powers

- There is an efficient method.
- Let $k = \left\lfloor \frac{n}{2} \right\rfloor$
- If *n* is even, $k = \frac{n}{2}$ and if *n* is odd, $k = \frac{n-1}{2}$
- So,

$$power(x, \left\lfloor \frac{n}{2} \right\rfloor)^2 \quad \text{if n is even}$$

$$power(x, \left\lfloor \frac{n}{2} \right\rfloor)^2 * x \quad \text{if n is odd}$$

Recursion

Computing Powers

- Then, we can redefine *power*(x, n) as follows:

1	If n == 0
$power(x, \left\lfloor \frac{n}{2} \right\rfloor)^2$	if n is even
$power(x, \left\lfloor \frac{n}{2} \right\rfloor)^2 * x$	if n is odd

Recursion
Computing Powers

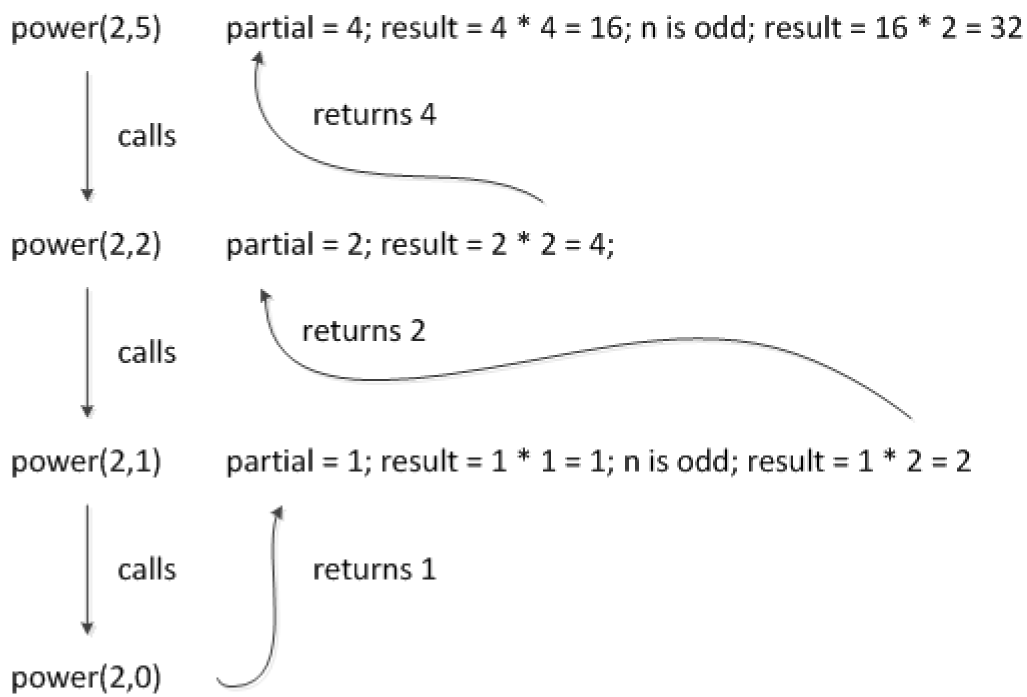
- Implementation

```
public static double power(double x, int n) {
  if (n == 0)
    return 1;
  else {
    double partial = power(x, n/2); // use integer division of n
    double result = partial * partial;
    if (n % 2 == 1)      // if n odd, include extra factor of x
      result *= x;
    return result;
  }
}
```

- Execution of one call takes $O(1)$.
- The method is invoked $O(\log n)$ times.
- Running time is $O(\log n)$

Recursion
Computing Powers

- Illustration



power(2,16) – power(2,8) – power(2,4) – power(2,2) – power(2,1) – power(2,0)
power(2,15) – power(2,7) – power(2,3) – power(2,1) – power(2,0)

Recursion

Designing Recursive Algorithms

- Two components: base case and recursion
- Base case:
  - Recursive call stops when a certain condition is met.
  - This is usually referred to as *base case*.
- Recursion: When the condition of the base case is not met, the algorithm invokes itself recursively.

- When poorly designed, very inefficient.
- Make sure the base case is always reached to avoid infinite recursion.

Recursion
Parameterizing Recursion

- Design of recursive algorithms sometimes requires the change of signature by adding more parameters.

- Natural signature of binary search:

  binarySearch (data, target)

- Recursive design requires additional parameters:

  binarySearch(data, target, low, high)

- Cleaner public interface:

```
public static boolean binarySearch(int[ ] data, int target) {
    return binarySearch(data, target, 0, data.length – 1);
}
```

Recursion
Tail Recursion

- Recursion allows exploitation of repetitive structure of a problem.
- Makes algorithm description more readable; avoids complex analyses and nested loops.
- Requires more memory.
- Tail recursion: A recursive call is the last operation.
- A tail recursion can be converted to a non-recursive algorithm (or implementation) that does not use additional memory.
- Example: binary search

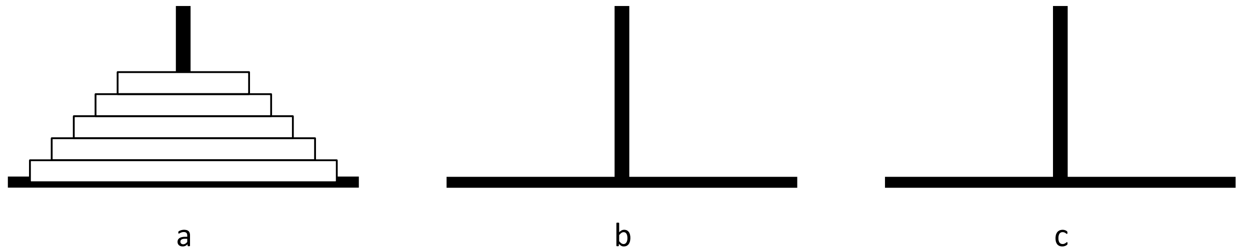
Recursion

Towers of Hanoi (Exercise)

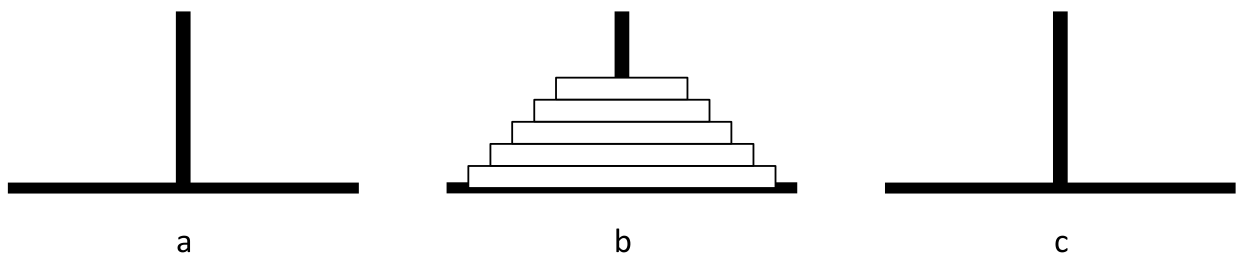
- Well known problem
- Given 3 pegs *a*, *b*, and *c*
- Peg *a* has *n* disks, the smallest on the top and the largest at the bottom
- Move all disks from *a* to *b*
- Use *c* as a temporary peg

Recursion
Towers of Hanoi (Exercise)

- Initial



- Final



Recursion

Towers of Hanoi (Exercise)

- When moving disks:
  - One disk at a time
  - Never place a larger disk on top of a smaller disk
- Each disk has a label
  - Label of the smallest disk is 1
  - Label of the next smallest disk is 2
  - . . .
  - Label of the largest disk is *n*

References

- M.T. Goodrich, R. Tamassia, and M.H. Goldwasser, “Data Structures and Algorithms in Java,” Sixth Edition, Wiley, 2014.