Software Transactional Memory "for real"

Armin Rigo

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Introduction

This talk is about programming multi- or many-core machines

About myself

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- "Language implementation guy"
- PyPy project
 - Python in Python
 - includes a Just-in-Time Compiler "Generator" for Python and any other dynamic language

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Motivation

- A single-core program is getting exponentially slower than a multi-core one
- Using several processes exchanging data
 - works fine in some cases
 - but becomes a large mess in others
- Using several threads
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Common solution

- Organize your program in multiple threads
- Add synchronization when accessing shared, non-read-only data

Synchronization with locks

- Carefully place locks around every access to shared data
- How do you know if you missed a place?
 - hard to catch by writing tests
 - instead you get obscure rare run-time crashes
- Issues when scaling to a large program
 - order of acquisition
 - deadlocks

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Synchronization with TM

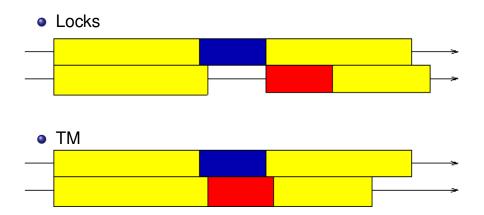
TM = Transactional Memory

```
mylock.acquire(); atomic {
x = list1.pop(); x = list1.pop();
list2.append(x);
mylock.release(); }
```

Synchronization with TM

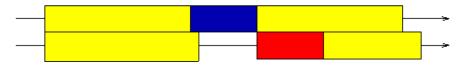
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Locks versus TM



Locks versus TM

Locks



TM in case of conflict



Synchronization with TM

- "Optimistic" approach:
 - no lock to protect shared data in memory
 - instead, track all memory accesses
 - detect actual conflicts
 - if conflict, restart the whole "transaction"
- Easier to use
 - no need to name locks
 - no deadlocks
 - "composability"

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HTM versus STM

- HTM = Hardware Transactional Memory
 - Intel Haswell CPU, 2013
 - and others
- STM = Software Transactional Memory
 - various approaches
 - large overhead (2x-10x), but getting faster
 - experimental in PyPy: read/write barriers, as with GC

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Analogy with Garbage Collection

- Explicit Memory Management:
 - messy, hard to debug rare leaks or corruptions
- Automatic GC solves it
 - common languages either have a GC or not
 - if they have a GC, it controls almost all objects
 - not just a small part of them

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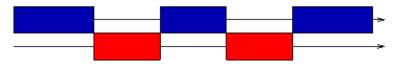
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Proposed solution

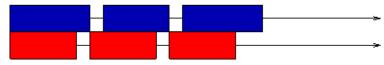
- Put atomic everywhere...
- in other words, Run Everything with TM

Proposed solution

Really needs TM. With locks, you'd get this:



With TM you can get this:



In a few words

- Longer transactions
- Corresponding to larger parts of the program
- The underlying multi-threaded model becomes implicit

Typical example

- You want to run f1 () and f2 () and f3 ()
- Assume they are "mostly independent"
 - i.e. we expect that we can run them in parallel
 - but we cannot prove it, we just hope that in the common case we can
- In case of conflicts, we don't want random behavior
 - i.e. we don't want thread-like non-determinism and crashes

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Pooling and atomic statements

- Solution: use a library that creates a pool of threads
- Each thread picks a function from the list and runs it with atomic

Results

- The behavior is "as if" we had run f1(), f2() and f3() sequentially
- The programmer chooses if he wants this fixed order, or if any order is fine
- Threads are hidden from the programmer

More generally

- This was an example only
- TM gives various new ways to hide threads under a nice interface

Not the Ultimate Solution

- Much easier for the programmer to get reproducible results
- But maybe too many conflicts
- "The right side" of the problem
 - start with a working program, and improve performance
 - as opposed to: with locks, start with a fast program, and debug crashes
 - we will need new debugging tools

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PyPy-STM

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 - in-progress, but basically working
 - solves the "GIL issue" but more importantly adds atomic
- http://pypy.org/
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