

# Time Warp Simulation on Multi-core Processors and Clusters

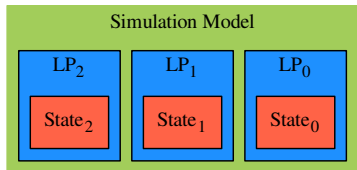
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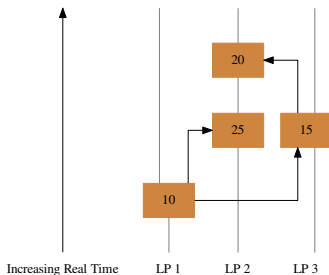
March 8, 2016

- ▶ Three main components
  - ▶ State variables
  - ▶ Simulation clock
  - ▶ Pending event set
- ▶ Unprocessed events stored in pending event set
- ▶ Events processed in time stamp order
- ▶ Simulation clock and state variables updated only when event occurs

- ▶ Model system as a set of Logical Processes (LPs)
- ▶ Events exchanged between LPs

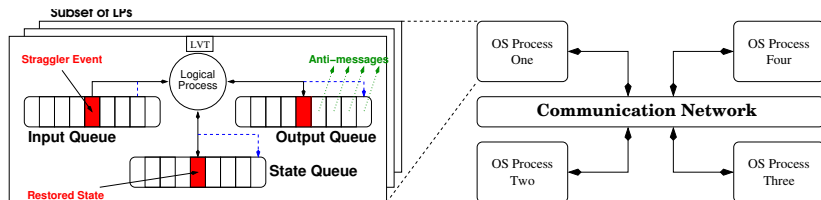


## Possible Causality Violations

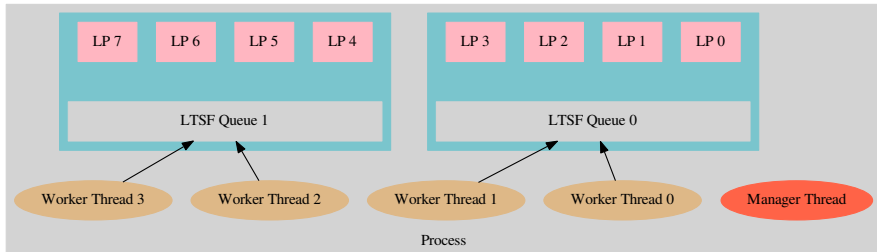


- ▶ Events can be received *and* processed out of order
- ▶ Two solutions
  - ▶ Conservative
  - ▶ Optimistic

## Optimistic Mechanism



- ▶ Rollback Mechanism
  - ▶ State Restoration, Anti-Messages
- ▶ Local Virtual Time (LVT) & Global Virtual Time (GVT)
- ▶ Fossil Collection



- ▶ Worker Threads and Manager Thread
- ▶ LTSF Queues/Worker Thread: Contention vs Rollbacks
- ▶ LP Partitioning

## Worker Threads

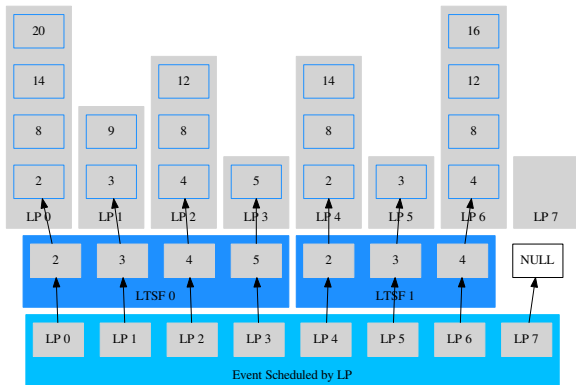
```
while termination not detected do
  e ← getNextEvent()
  lp ← receiver of e

  if e < last processed event for lp
  then
    └ rollback lp

  if e is an anti-message then
    └ cancel event with e
    └ schedule new event for lp
    continue

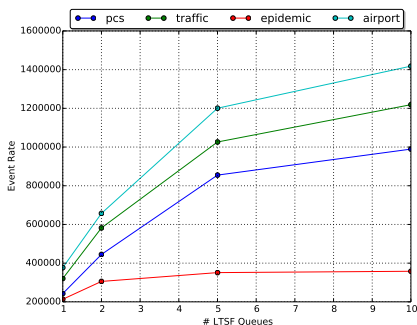
  process event e
  save state of lp
  send new events

  move e to processed queue
  replace scheduled event for lp
```

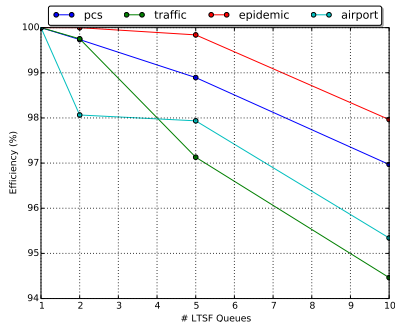


- Unprocessed and Processed Queues
- Scheduling to LTSF
- Rollbacks

Intel® Xeon® X5675 (6 cores/10 worker threads)



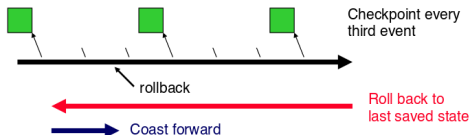
$$EventRate = \frac{CommittedEvents}{Runtime}$$



$$Efficiency = \frac{CommittedEvents}{ProcessedEvents} * 100\%$$

Reducing contention more important than reducing rollbacks

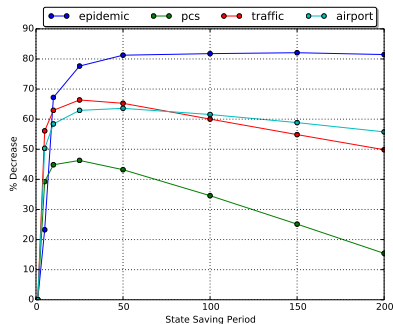
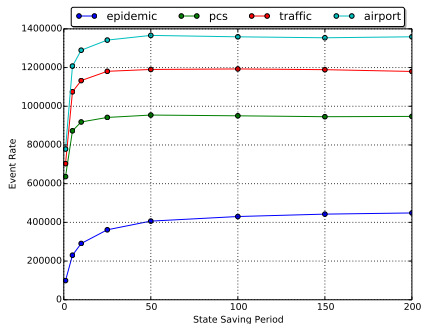
## Save state only once every $N$ events



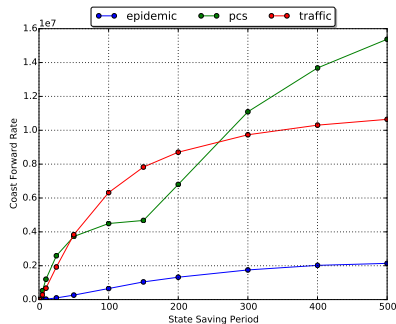
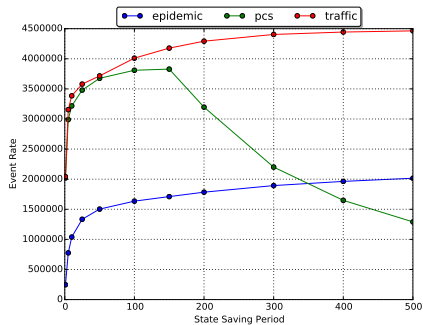
- ▶ Not all states available to roll back to
- ▶ Must “Coast Forward” to reproduce state
- ▶ Increases time to rollback
- ▶ Decrease time to save states and reduce memory footprint



Intel® Xeon® X5675 (6 cores/10 worker threads)



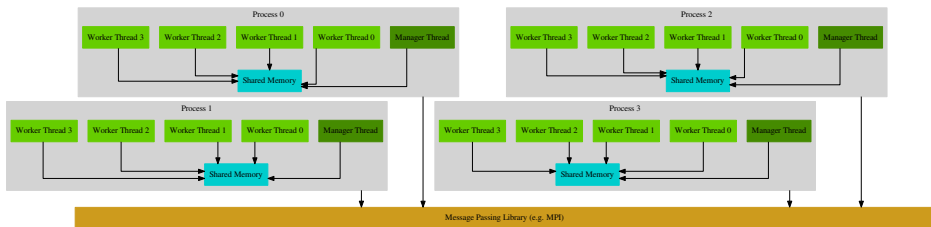
## Intel® Xeon® E5410 (8 nodes)



$$\text{CoastForwardRate} = \frac{\text{CoastForwardEvents}}{\text{Rollbacks}} * \text{RollbackRate}$$

## Shared Memory vs Message Passing

- ▶ Sending Events, GVT, Termination

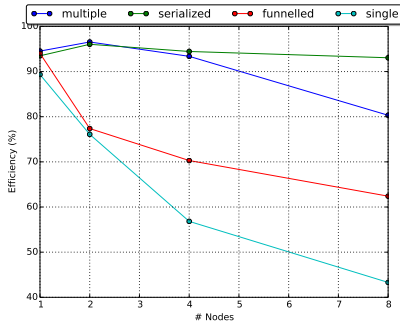
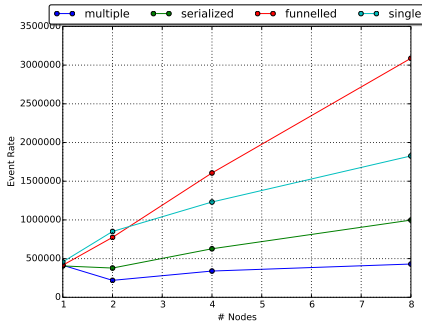


## Thread Message Passing Models

- ▶ Single threaded
- ▶ Funnelled
- ▶ Serialized
- ▶ Multiple

## Intel® Xeon® E5410 (8 nodes)

- ▶ Single: 1 thread (combining worker/manager threads)
- ▶ Others: 7 worker threads and manager thread



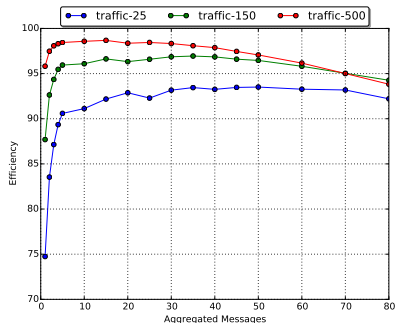
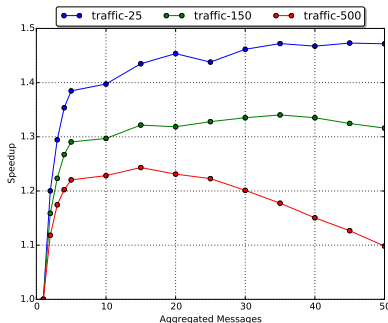
- ▶ Funneled/Single: Minimal Synchronization, High Latency
- ▶ Serialized/Multiple: Lots of Synchronization, Low Latency
- ▶ Performance will vary with model and partitioning.

- ▶ Wait for  $N$  messages with same **destination process** and send together

Num Msgs 4 bytes	Length 0 (L[0]) 4 bytes	Length 1 (L[1]) 4 bytes	Length 2 (L[2]) 4 bytes	...	Message 0 L[0]	Message 1 L[1]	Message 2 L[2]	...
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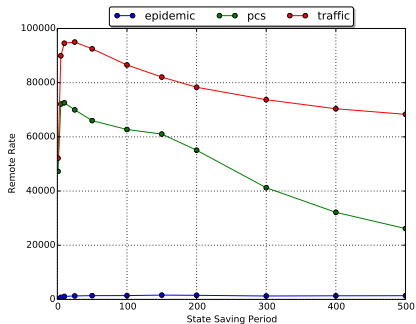
- ▶ Pros
  - ▶ Single buffer allocated for multiple messages
  - ▶ Latency shared by multiple messages (TCP/IP/Ethernet/...)
  - ▶ Events destined for same LP may be sent together
- ▶ Cons
  - ▶ May increase event latency and increase receiver rollbacks

# Message Aggregation: Traffic Model

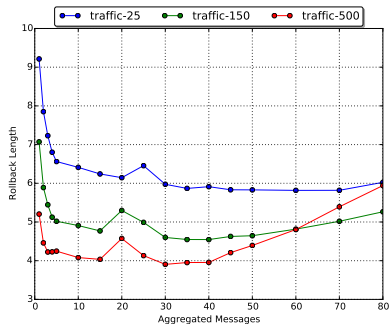


- ▶ Best with small number of aggregated events
- ▶ Different behavior with different state saving periods

# Message Aggregation: More Observations



► Faster rate of remote events better

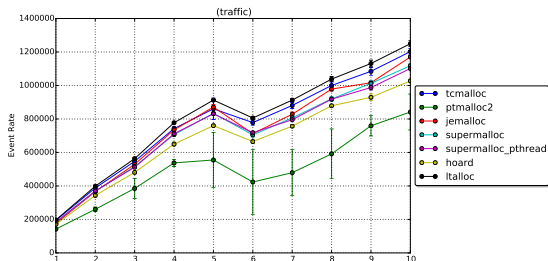


► Can lower rollback length?

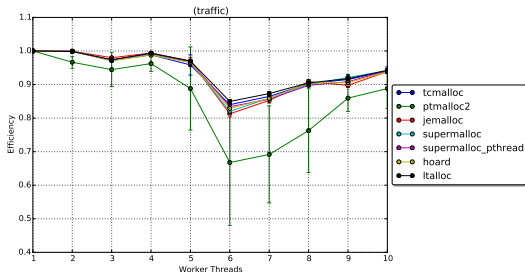
$$\text{AverageRollbackLength} = \text{RolledBackEvents} / \text{Rollbacks}$$

# Memory Allocation: Traffic Model

Intel® Xeon® X5675 (6 cores/up to 12 threads)



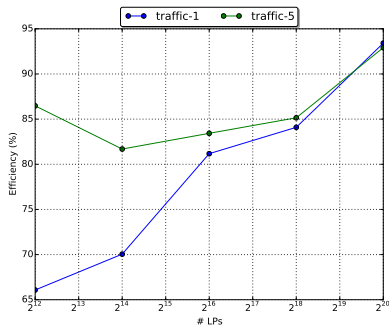
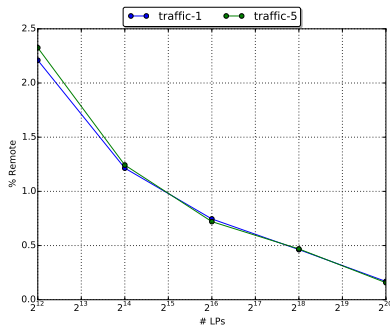
► Imbalance at 6 worker threads



► ptmalloc2 (the default in GLIBC) is inefficient and unpredictable

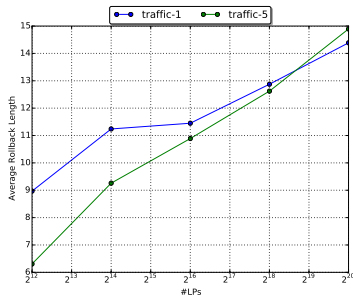
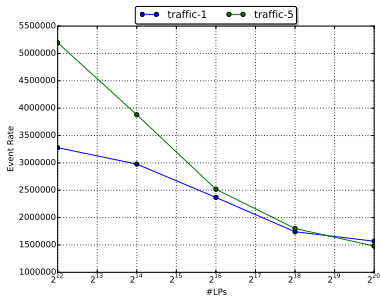


Intel® Xeon® E5410 (8 nodes/7 worker threads)



- ▶ Coarser Granularity
- ▶ Higher Efficiency

Intel® Xeon® E5410 (8 nodes/7 worker threads)



- ▶ More LPs to fossil collect
- ▶ Longer Rollbacks

## ▶ Conclusions

- ▶ Synchronization should always be avoided
  - ▶ Interprocess communication: Higher latency over synchronization
  - ▶ 1 LTSF queue per worker thread
- ▶ Avoid communication ... If possible
  - ▶ If it cannot be avoided, aggregating messages may help
  - ▶ Hard to calculate GVT efficiently
  - ▶ Memory consumption: May need flow control
- ▶ Alternative multi-threaded memory allocators better than GLIBC default

## ▶ Future Research

- ▶ Optimistic Fossil Collection
- ▶ Scaling up