# **Chapter S:VI**

#### VI. Relaxed Models

- Motivation
- $\Box$   $\varepsilon$ -Admissible Speedup Versions of A\*
- □ Using Information about Uncertainty of *h*
- □ Risk Measures
- Nonadditive Evaluation Functions
- □ Heuristics Provided by Simplified Models
- Mechanical Generation of Admissible Heuristics
- □ Probability-Based Heuristics

Using a non-admissible heuristic function

Idea [Harris]:

The heuristic function h estimates the cheapest remaining cost  $h^*$  mostly quite well, but sometimes overestimates  $h^*$  by no more than  $\varepsilon$ .

 $\rightarrow$  A\* using such a heuristic function h is  $\varepsilon$ -admissible.

The condition for  $\varepsilon$ -admissibility of A\* is satisfied because at termination it holds

$$h(n) - h^*(n) \le \varepsilon$$
 für alle  $n \in \mathsf{OPEN}$ .

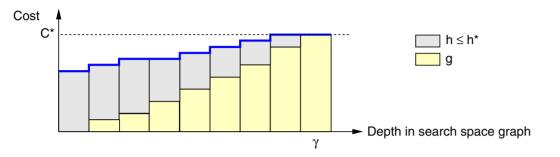
Also the weakened form of admissibility of h is often too restrictive.

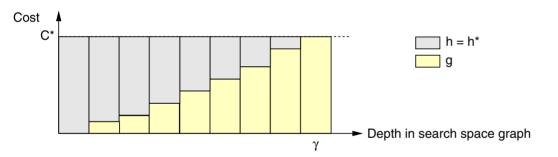
Often it is easier to find a heuristic estimate for  $h^*$  that mostly estimates precisely but sometimes overestimates  $h^*$  (by much more than any reasonable  $\varepsilon$ ).

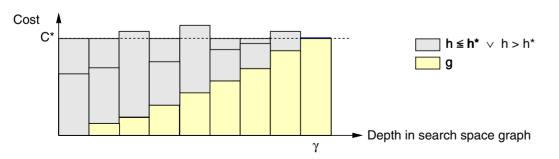
→ The error in the estimate is not limited, but a large error is unlikely.

- $\Box$  Heuristic functions h with  $h \leq (1+\varepsilon)h^*$  are called  $\varepsilon$ -admissible.
- Analogously to Lemma  $C^*$ -Bounded OPEN Node, it can be proven that, at any point in time before termination, there exists some node n in OPEN with  $f(n) \leq (1 + \varepsilon)C^*$ .
- The condition " $h(n) h^*(n) \le \varepsilon$  for all  $n \in \mathsf{OPEN}$ " is sufficient, but not necessary, for A\* being  $\varepsilon$ -admissible.

Illustration of Underestimating and Overestimating Estimation Functions







Example: Search in "Random" Graphs

Given is a graph with randomly drawn edge costs. The minimum number of edges to a target node is known in each node.

- $\Box$  Edge costs c(n, n') are known to be drawn independently from a common distribution function, uniform in interval [0; 1].
- $\supset$  For long paths with N edges from a node n to a goal node in  $\Gamma$  it is known that h\*(n) is most likely to be near  $\frac{N}{2}$ .
- $\Box$  The only *admissible* heuristic estimate for  $h^*$  is  $h_1(n) = 0$ .
- $\Box$  The most reasonable heuristic estimate for  $h^*$  is  $h_2(n) = \frac{N}{2}$ .

The heuristic estimate  $h_2$  leads to a worst-case cost overestimation of  $\frac{N}{2}$  and is therefore not  $(\varepsilon$ -)admissible. But the likelihood of this event is extremely small.

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## → Algorithm $R^*_{\delta}$ :

- Besides an estimation function h for  $h^*$  there is also knowledge about the uncertainty of the estimation process.
- Knowledge about the uncertainty of the estimation process is expressed in the form of a probability density function  $\rho_{h^*}(x)$ .

Describing the Estimation Uncertainty using Density Functions

Viewing cost functions as a random variables:

cost function	random variable
$h^*(n)$	$h_n^*$
$f^*(n) = g^*(n) + h^*(n)$	$f_n^*$
$f^+(n) = g(n) + h^*(n)$	$f_n^+$

Let  $\rho_{h_n^*}$  be a density function for the random variable  $h_n^*$ .

#### Semantics:

On the basis of  $\rho_{h_n^*}$  one can define the probability with which  $h^*(n)$  can be found in a neighborhood of x costs.

$$P(h_n^* = x) = \rho_{h_n^*}(x)$$

Describing the Estimation Uncertainty using Density Functions (continued)

Let  $\rho_{h_n^*}$  be a density function for the random variable  $h_n^*$ .

## Further applies:

1. From  $\rho_{h_n^*}(x)$  a density function  $\rho_{f_n^*}(y)$  can be derived for the random variable  $f_n^*$ , if  $g^*$  is known (e.g., when searching a tree):

$$\rho_{f_n^*}(y) := \rho_{h_n^*}(y - g^*)$$

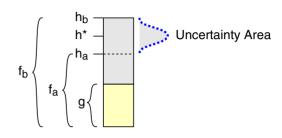
2. Let  $P_{s-n}$  be the cheapest known path from s to an OPEN node n. From  $\rho_{h_n^*}(x)$  a density function  $\rho_{f_n^+}(y)$  can be derived for the random variable  $f_n^+$ , which specifies the cost of an optimal solution path that continues  $P_{s-n}$ :

$$\rho_{f_n^+}(y) := \rho_{h_n^*}(y - g)$$

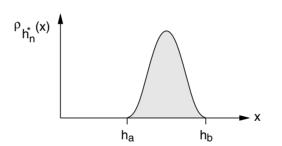
- The random variable  $f_n^+$  with associated density function  $\rho_{f_n^+}$  is given for each node n.
- The random variable  $f_n^+$  describes the possible costs of an optimal solution path that contains the pointer path  $P_{s-n}$  as a subpath.
- If goal nodes can be reached from s, the OPEN list always contains a node n, to which  $f^+(n) = f^*(n)$  applies. [Corollrary Shallowest OPEN Node on Optimum Path]

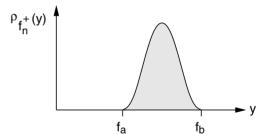
Describing the Estimation Uncertainty using Density Functions (continued)

Uncertainty area:

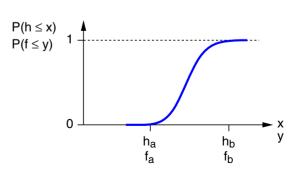


Density functions  $\rho$ :





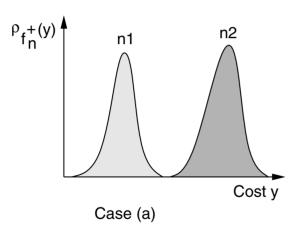
Related distribution function:

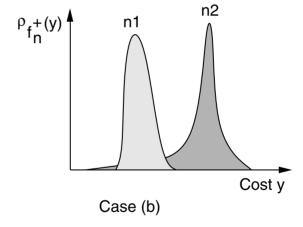


Describing the Estimation Uncertainty using Density Functions (continued)

How should an evaluation order be calculated from the density functions  $\rho_{f_n^+}$  for the nodes in the OPEN list?

Possible shapes of two density functions:





- (a) If the density functions do not overlap, the node for which the corresponding density function  $\rho_{f^+}$  has the lowest density value  $f_a^+$  with respect to all other nodes would be selected.
- (b)  $f_{n_1}^+$  has the lower expected value;  $n_2$  has the possibility that the cost  $f_{n_2}^+$  may be lower than  $n_1$ . An admissible algorithm would expand  $n_2$ . It would make more sense to expand  $n_1$  because the " $f^+(n_2) < f^+(n_1)$ " event is unlikely.
- → Due to uncertainty, costs can be overestimated or underestimated. I.e., not expanding a node in OPEN and terminating it too expensively as a result, represents a *risk*.
- → Quantification of the risk of terminating with too high costs (= terminating too early).

## Defining the Order of Node Evaluations

#### Idea:

Estimate the risk of terminating too early using a *risk measure* R for each node in the OPEN list.

- $\Box$  For a given cost value C (of a goal node), the risk measure evaluates for each node n in the OPEN list to what extent C can be improved by expanding n.
- $\square$  R = R(C). The risk measure is a nondecreasing function of the C cost. The greater the R(C) value of a node n, the greater the risk of missing an improvement of C if terminating with C without expanding n.
- $\ \square$  R(C) should use knowledge about the cost distribution for the node n, so it should be based on  $\rho_{f^+}$ .

Principle of the Algorithm R\*<sub>δ</sub>

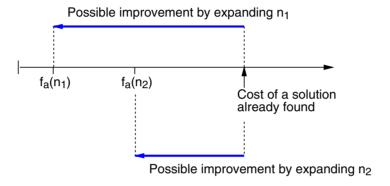
Search continues until the risk value R(C) of each node in the OPEN list is below a user accepted risk threshold  $\delta$ .

- → If a high risk is acceptable, nodes with a high risk value of R(C) (= high cost reduction potential) remain unexpanded. As a result, cost underestimation becomes less likely.
  If only a small risk is acceptable, even nodes with a low risk value of R(C) (=
  - lf only a small risk is acceptable, even nodes with a low risk value of R(C) (= low cost reduction potential) are expanded. As a result, cost underestimation becomes more likely.
- $\rightarrow$  I.e., depending on a risk threshold  $\delta$  the probability of a cost underestimation (= probability of admissibility or optimality) can be controlled.

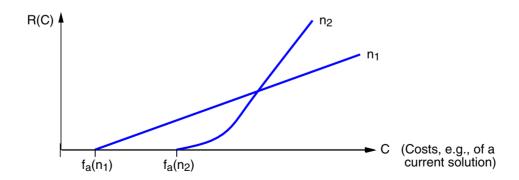
- □ The observance of this principle by the algorithm  $R^*_{\delta}$  is ensured as shown later by the use of certain risk functions R(C).
- When a goal node with cost C is selected from OPEN, its risk function must guarantee the property  $C \le C_{\delta}$ . Otherwise, one of the remaining nodes in OPEN can have a risk for C that is higher than  $\delta$ .

## Potential for Improvement to a Current Solution

Let  $n_1$ ,  $n_2$  be nodes of the OPEN list.



Example of risk functions R(C) for the nodes  $n_1, n_2$ :

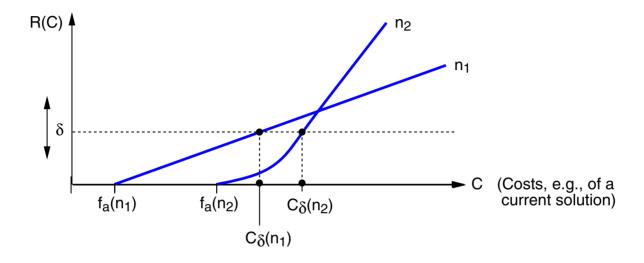


The nodes have different random variables  $f_{n_1}^+$  and  $f_{n_2}^+$  for the cost.

- $\Box$  The potential for improvement is a statistical quantity defined for a node n using  $f_n^+$ .
- $\Box$  The *evaluation* of the potential for improvement regarding given costs C is done with the help of a risk-measure R(C).

### Risk Threshold and Cost Treshold

The risk threshold  $\delta \geq 0$  defines for each node n in OPEN its cost threshold  $C_{\delta}(n)$ :



Let  $n_1$ ,  $n_2$  be nodes in OPEN. If the search was terminated with node  $n_2$  and cost  $C' = C_{\delta}(n_2)$ , the risk R(C') for  $n_1$  would be above the risk threshold  $\delta$ .

 $\Rightarrow$  R\* $_{\delta}$  chooses the node n with the lowest cost threshold  $C_{\delta}(n)$  in the OPEN list. In the above example, node  $n_1$  would be preferred to node  $n_2$ .

### **Definition 75 (Risk Measure)**

Let M be the ordered set of cost values. A risk measure R(C) for a node is a nondecreasing function  $R:M\to [0,\infty]$  measuring the risk associated with leaving that node unexplored when terminating with a solution with cost C.

### **Definition 76 (Cost Threshold)**

Let  $\delta$  be a nonnegative real number and let R(C) be the risk measure for a node n. The solution  $C_{\delta}(n)$  to the equation  $R(C) = \delta$  is called the cost threshold.

Assuming the cost of a solution path found is C, then for each node n in OPEN with  $C > C_{\delta}(n)$  the risk of missing a better solution path is higher than risk threshold  $\delta$ . These nodes should be expanded before termination.

- ☐ Risk measures and risk thresholds must be seen in context: not every risk threshold makes sense for a risk measure.
- $\Box$  Depending on the  $f_n^+$  cost random variable of a n node, the  $\delta$  risk threshold can lead to different sequences in the OPEN list.
- The cost-threshold  $C_{\delta}(n)$  indicates how high the cost of a solution may be without exceeding the  $\delta$ -risk-threshold for the node n.

## **Definition** 77 ( $\delta$ -Risk-Admissibility)

An algorithm is said to be  $\delta$ -risk-admissible if it always terminates with a solution cost C such that  $R(C) \leq \delta$  for each node left on OPEN.

The above version of the  $\delta$ -risk-admissible condition is equivalent to stating that at termination, the cost of the solution found is not greater than  $C_{\delta}(n)$  for each n on OPEN.

## **Definition** 78 (Algorithm $R^*_{\delta}$ )

 $\mathsf{R}^*_\delta$  is a search algorithm which is identical to  $\mathsf{A}^*$  except that it chooses for expansion that node n from OPEN with the lowest cost-threshold  $C_\delta(n)$ .

Note that with  $\delta = 0$ ,  $R^*_{\delta}$  is identical to  $A^*$  since it is guided by the (admissible) lowest tail of the density of f, namely by  $g + h_a$ .

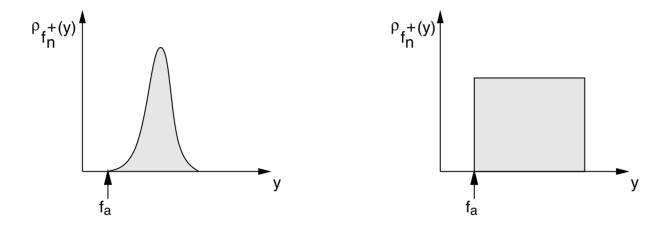
For  $\delta > 0$ ,  $R^*_{\delta}$  may prefer a node with high  $f_a$  and narrow distribution over a node with low  $f_a$  but highly diffused density.

- $\Box$  The first definition of  $\delta$ -risk-admissibility is from the perspective of risk, the second is from the perspective of cost.
- $\Box$  With  $\delta = 0$ ,  $R^*_{\delta}$  is identical to  $A^*$ . Justification:
  - 1. Computiong the cost-threshold  $C_{\delta}(n)$  for nodes in OPEN is solving the equation R(C)=0 for  $\delta=0$ .
  - 2. R(C) = 0 holds for the lowest point  $f_a$  on the tail of the density of  $f^+$ .
  - 3. Hence,  $f_a = g + h_a \le g(n) + h^*(n)$
  - 4.  $R^*_{\delta}$  is guided by an admissible heuristic function and, therefore,  $R^*_{\delta}$  is admissible.
- $\Box$  As the  $\delta$  increases,  $R^*_{\delta}$  tends to abandon admissibility.

Risk Measures of Type  $R(C) = \varrho[C - f^+]$ 

Starting point are density functions for the random variables  $f_n^+$  of nodes n in the OPEN list.

## Examples:



 $f_a$  (resp.  $h_a$ ) is the smallest positive preimage of the density function  $\rho_{f_n^+}$  (resp.  $\rho_{h_n^*}$ ).

Risk Measures of Type  $R(C) = \varrho[C - f^+]$  (continued)

1. Worst Case Risk  $R_1$ :

$$R_1(C) = \sup_{\{y \mid \rho_{f_n^+}(y) > 0\}} (C - y) = C - f_a = C - g - h_a$$

2. Probability of Suboptimal Termination  $R_2$ :

$$R_2(C) = P(C > f_n^+) = P(C - f_n^+ > 0) = \int_{y=-\infty}^{C} \rho_{f_n^+}(y) dy$$

3. Expected Risk  $R_3$ :

$$R_3(C) = E(\max\{C - f_n^+; 0\}) = \int_{y = -\infty}^{C} (C - y) \rho_{f_n^+}(y) dy$$

- $\Box$  The risk measures  $R_1$  and  $R_3$  describe costs, the risk measure  $R_2$  describes a probability.
  - $R_1$ : For the costs represented by the  $f_n^+$  random variable, the smallest possible value is assumed.  $R_1$  quantifies the maximum possible loss if a solution is satisfied with C costs.

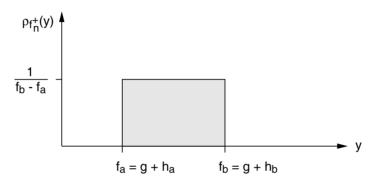
The lowest costs are the worst case because they represent the extreme case of a missed cost reduction. The probability that the remaining costs are lower than  $h_a$  is 0.

- $R_2$ : The probability for the occurrence of the event " $C > f_n^+$ " (i.e., event is a loss) is calculated if you are satisfied with a solution with C cost.
- $R_3$ : For the costs represented by the random variable  $f_n^+$ , the expected loss  $E(\max\{C-f_n^+;0\})$  is calculated if one is satisfied with a solution with costs C.

 $R_3$  weights the probability of the loss  $(R_2)$  with the amount of the occurring loss.

# Example

Let  $f_n^+$  be uniformly distributed between an optimistic estimate  $f_a$  and a pessimistic estimate  $f_b$  (The  $f_a$  and  $f_b$  estimates depend on n, where n is in OPEN.):



Density function:

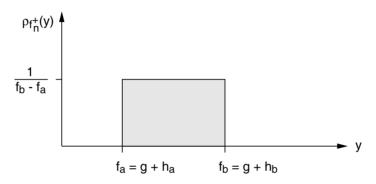
$$\rho_{f_n^+}(y) = \begin{cases} \frac{1}{f_b - f_a} & f_a \le y \le f_b \\ 0 & \text{else} \end{cases}$$

Risk measure:

$$R_1(C) = \sup_{\{y \mid \rho_{f_n^+}(y) > 0\}} (C - y) = C - f_a$$

Example (continued)

Let  $f_n^+$  be uniformly distributed between an optimistic estimate  $f_a$  and a pessimistic estimate  $f_b$  (The  $f_a$  and  $f_b$  estimates depend on n, where n is in OPEN.):



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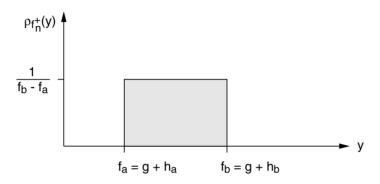
$$\rho_{f_n^+}(y) = \begin{cases} \frac{1}{f_b - f_a} & f_a \le y \le f_b \\ 0 & \text{else} \end{cases}$$

Risk measure:

$$R_2(C) = \int_{y=-\infty}^{C} \rho_{f_n^+}(y) dy = \begin{cases} 0 & C < f_a \\ \frac{(C-f_a)}{(f_b-f_a)} & f_a \le C \le f_b \\ 1 & f_b < C \end{cases}$$

Example (continued)

Let  $f_n^+$  be uniformly distributed between an optimistic estimate  $f_a$  and a pessimistic estimate  $f_b$  (The  $f_a$  and  $f_b$  estimates depend on n, where n is in OPEN.):



Density function:

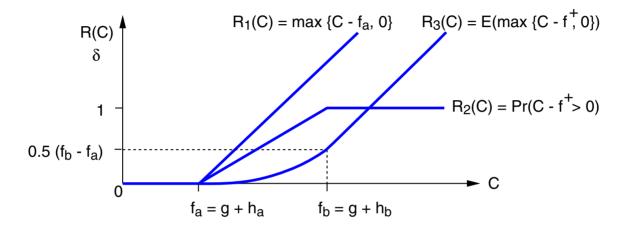
$$\rho_{f_n^+}(y) = \begin{cases} \frac{1}{f_b - f_a} & f_a \le y \le f_b \\ 0 & \text{else} \end{cases}$$

Risk measure:

$$R_3(C) = \int\limits_{y=-\infty}^C (C-y) \rho_{f_n^+}(y) dy = \begin{cases} 0 & C < f_a \\ \frac{(C-f_a)^2}{2(f_b-f_a)} & f_a \le C \le f_b \\ C - \frac{f_a+f_b}{2} & f_b < C \end{cases}$$
 Optimality Requirement

Example (continued)

Shape of risk measures  $R_1$ ,  $R_2$ , and  $R_3$  ( $f_n^+$  uniformly distributed):

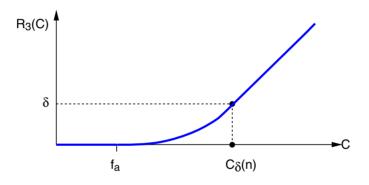


The vertical axis represents the functions R(C) for the three risk measures considered.

Example (continued)

Computing the cost threshold  $C_{\delta}$  for  $R_3$  ( $f_n^+$  uniformly distributed):

Let  $\delta$  be the user's risk tolerance of the user. For each node n in OPEN, it defines its cost threshold  $C_{\delta}(n)$  using the equation  $R(C) = \delta$ .



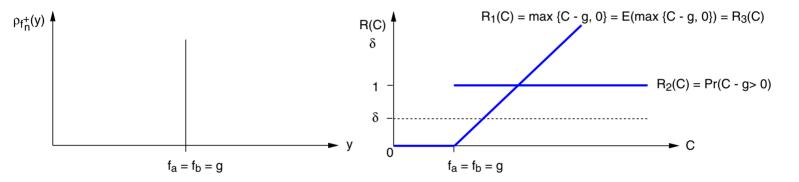
For a node n on OPEN,  $C_{\delta}(n)$  is computed by transforming  $R_3(C_{\delta}) = \delta$ :

$$C_{\delta}(n) = \begin{cases} f_a & \delta = 0 \\ f_a + \sqrt{2 \cdot (f_b - f_a) \cdot \delta} & 0 < \delta \le \frac{f_b - f_a}{2} \\ \delta + \frac{f_a + f_b}{2} & \frac{f_b - f_a}{2} < \delta \end{cases}$$

Example (continued)

For a goal node  $\gamma$  in OPEN, we have  $h(n) = h^*(n) = 0$ . Therefore, there remains no uncertainty regarding  $f_{\gamma}^+$ .

Graphs of random variable  $f_{\gamma}^+$  for solution cost and risk measures  $R_1$ ,  $R_2$ , and  $R_3$ :



#### Cost treshold:

$$C_{\delta}(\gamma) = \begin{cases} g(\gamma) + \delta & \text{for risk measure } R_1 \\ g(\gamma) & \text{for risk measure } R_2 \text{ and } \delta < 1 \\ g(\gamma) + \delta & \text{for risk measure } R_3 \end{cases}$$

### **Theorem** 79 ( $\delta$ -Risk-Admissibility of $R^*_{\delta}$ )

 $\mathsf{R}^*_\delta$  is  $\delta$ -risk-admissible with respect to risk measures  $R_1$ ,  $R_2$ , and  $R_3$  when G is a search space graph with  $\mathit{Prop}(G)$  and  $E(h^*) < \infty$  on solution paths.

### **Theorem** 79 ( $\delta$ -Risk-Admissibility of $R^*_{\delta}$ )

 $\mathsf{R}^*_\delta$  is  $\delta$ -risk-admissible with respect to risk measures  $R_1, R_2$ , and  $R_3$  when G is a search space graph with  $\mathit{Prop}(G)$  and  $E(h^*) < \infty$  on solution paths.

### **Proof** (sketch)

- 1.  $\delta$ -Risk-Admissibility:
  - (a) According to the previous example, it holds for the cost C of a solution path found by  $R^*_{\delta}$ :

$$C = g(\gamma) \le C_{\delta}(\gamma)$$
 for the risk measures  $R_1, R_2, R_3$ .

- (b) Since  $R^*_{\delta}$  chooses for expansion that node n from OPEN with the lowest cost-threshold  $C_{\delta}(n)$ ,  $\delta$ -risk-admissibility of  $R^*_{\delta}$  follows for risk measures  $R_1$ ,  $R_2$ , and  $R_3$ .
- 2. Completeness:
  - (a) At all times OPEN contains a node n on a solution path for which  $C_{\delta}(n)$  is finite. Obviously,  $R_1(C) = \delta$  and  $R_2(C) = \delta$  have a finite solution. If density  $\rho_{h^*}(x)$  possesses a finite expectation  $E(h^*) < \infty$  for any node on a solution path, for  $R_3$  we have

$$R_3(C) \ge C \cdot (1 - P(f^+ > C)) - E(f^+) \ge C - 2E(f^+) = C - 2g - 2E(h^*)$$

(b)  $C_{\delta}(n) \geq g(n)$  holds for each node n in OPEN since there is no risk in abandoning n after finding a solution path with cost  $\leq g(n)$ . A positive lower bound of the edge cost values guarantees that  $\mathsf{R}^*_{\delta}$  can neglect nodes on solution paths only for a limited number of node expansions.

- $\square$  Expectations can have the value  $\infty$ , e.g., for a random variable that returns values  $2^n$  with probability  $2^{-n}$ .
- In step 2(a) we use the fact  $R_3(0) = 0$  for graphs G with nonnegative edge cost values. As the lower bound for  $R_3(C)$  increases with C, there is a finite value C with  $R_3(C) > \delta$ . Hence,  $C_\delta(n) < \infty$ .
- $\Box$  The exact form of  $\rho_{h_n^*}$  is generally unknown. For this the edge costs must have been generated by a given probabilistic model.
- Generating a good estimate for  $C_{\delta}(n)$  is often possible. For this, the knowledge of upper and lower bounds of  $h_n^*$  together with the often reasonable assumption of a standardized distribution between them, such as an uniform distribution, an exponential distribution or a normal distribution, is sufficient.
- □ The principle of the  $\varepsilon$ -admissible acceleration in  $A^*_{\epsilon}$  for  $A^*$  can also be applied to  $R^*_{\delta}$  and leads to the algorithm  $R^*_{\delta,\epsilon}$ . The special version  $R^*_{\delta,\delta}$  is  $\delta$ -risk-admissible with respect to risk measures  $R_1$ ,  $R_2$ , and  $R_3$  under the preconditions of the previous theorem.