

Planning (Chapter 2 from Lavalle book)

Abstraction of a planning problem

1. State space $s \in \mathcal{S}$. For example, 2D coordinate of a grid $s = (x, y)$.
2. Action space per state $u \in \mathcal{U}(s)$. For example, up, down, left right movement can be encoded as $\mathcal{U}(s_t) = \{(0, -1), (0, 1), (1, 0), (-1, 0)\}$.
3. State transition function $s_{t+1} = f(s_t, u_t)$. For example, the up-down-left-right action can be combined as addition to get the next state $s_{t+1} = s_t + u_t$.
4. Initial State $s_I \in \mathcal{S}$
5. Goal states $s_G \subseteq \mathcal{S}$

A Graph

A graph $\mathcal{G} = \{\mathcal{V}, \mathcal{E}\}$ is defined by a set of vertices \mathcal{V} and a set of edges \mathcal{E} such that each edge $e \in \mathcal{E}$ is formed by a pair of start and end vertices $e = (v_s, v_e), v_s \in \mathcal{V}, v_e \in \mathcal{V}$. The first vertex is called the start of the edge $v_s = \text{start}(e)$ and second vertex is called the end $v_e = \text{end}(e)$.

A discrete planning problem can be converted into a graph by defining

1. Vertices as the state space $\mathcal{V} = \mathcal{S}$.
2. The action space at each state as the edges connected to that vertex/state, $\mathcal{U}(s_t) = \{(s_t, s_j) \mid (s_t, s_j) \in \mathcal{E}\}$.
3. State transition function is the other end of the edge,

$$s_{t+1} = f(s_t, u_t) = \text{end}(u_t), \text{ where } s_t = \text{start}(u_t).$$

Representations of Graphs

Undirected graph

```
In [1]: # Programmatically you can represent a adjacency list as python lists  
# Python lists are not linked lists, they are arrays under the hood.  
G_adjacency_list = {  
    1 : [2, 5],  
    2 : [1, 5, 3, 4],  
    3 : [2, 4],  
    4 : [2, 5, 3],  
    5 : [4, 1, 2]  
}  
  
# Prefer to represent a matrix in python either as a list of lists or a numpy array  
import numpy as np  
G_adjacency_matrix = np.array([  
    [0, 1, 0, 0, 1],  
    [1, 0, 1, 1, 1],  
    [0, 1, 0, 1, 0],  
    [0, 1, 1, 0, 1],  
    [1, 1, 0, 1, 0]  
])  
  
# Edge list is another possible representation  
G_edge_list = [  
    (1, 2), (1, 5),  
    (2, 1), (2, 5), (2, 3), (2, 4),  
    (3, 2), (3, 4),  
    (4, 2), (4, 5), (4, 3),  
    (5, 4), (5, 1), (5, 2)  
]
```

Directed graph representation

```
In [2]: # Programmatically you can represent a adjacency list as python lists  
# Python lists are not linked lists, they are arrays under the hood.  
G_adjacency_list = {  
    1 : [2, 4],  
    2 : [5],  
    3 : [6, 5],  
    4 : [2],  
    5 : [4],  
    6 : [6]  
}  
  
# Prefer to represent a matrix in python either as a list of lists or a numpy array  
import numpy as np  
G_adjacency_matrix = np.array([  
    [0, 1, 0, 1, 0, 0],  
    [0, 0, 0, 0, 1, 0],  
    [0, 0, 0, 0, 1, 0],  
    [0, 0, 0, 0, 1, 0],  
    [0, 0, 0, 0, 1, 0],  
    [0, 0, 0, 0, 1, 0]  
])
```

```

    [0, 0, 0, 0, 1, 1],
    [0, 1, 0, 0, 0, 0],
    [0, 0, 0, 1, 0, 0],
    [0, 0, 0, 0, 0, 1]
])

# Edge list is another possible representation
G_edge_list = [
    (1, 2), (1, 4),
    (2, 5),
    (3, 6), (3, 5),
    (4, 2),
    (5, 6)
]

```

In [3]: # Exercise 1

```

# Write a function that converts a graph in adjacency list format to adjacency matrix
def adjacency_list_to_matrix(G_adj_list):
    G_adj_mat = None # TODO: Write code to convert to adj_mat
    return G_adj_mat

def adjacency_matrix_to_list(G_adj_mat):
    G_adj_list = None # TODO: Write code to convert to adj_mat
    return G_adj_list

# Use the above graphs to test
print(adjacency_list_to_matrix(G_adjacency_list))
print(adjacency_matrix_to_list(G_adjacency_matrix))

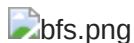
```

None

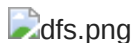
None

Graph Search algorithms

1. Breadth First Search



2. Depth First Search



Breadth first search (BFS)



In [4]: **from** queue **import** Queue, LifoQueue, PriorityQueue

```

graph = {
    's' : ['w', 'r'],
    'r' : ['v'],
    'w' : ['t', 'x'],
    'x' : ['y' ],

```

```

't' : ['u'],
'u' : ['y']
}

def bfs(graph, start, debug=False):
    seen = set() # Set for seen nodes (contains both frontier and dead state)
    # Frontier is the boundary between seen and unseen (Also called the alive)
    frontier = Queue() # Frontier of unvisited nodes as FIFO
    node2dist = {start : 0} # Keep track of distances
    search_order = []
    seen.add(start)
    frontier.put(start)

    i = 0 # step number
    while not frontier.empty(): # Creating loop to visit each node
        if debug: print("%d) Q = " % i, list(frontier.queue), end='; ')
        if debug: print("dists = " , [node2dist[n] for n in frontier.queue])
        m = frontier.get() # Get the oldest addition to frontier
        search_order.append(m)

        for neighbor in graph.get(m, []):
            if neighbor not in seen:
                seen.add(neighbor)
                frontier.put(neighbor)
                node2dist[neighbor] = node2dist[m] + 1
            else:
                assert node2dist[neighbor] <= node2dist[m] + 1, 'this should not happen'
                node2dist[neighbor] = min(node2dist[neighbor], node2dist[m] + 1)

        i += 1
    if debug: print("%d) Q = " % i, list(frontier.queue))
    return search_order, node2dist

```

In [5]: `print("Following is the Breadth-First Search order")`
`print(bfs(graph, 's', debug=True))` # function calling

```

Following is the Breadth-First Search order
0) Q = ['s']; dists = [0]
1) Q = ['w', 'r']; dists = [1, 1]
2) Q = ['r', 't', 'x']; dists = [1, 2, 2]
3) Q = ['t', 'x', 'v']; dists = [2, 2, 2]
4) Q = ['x', 'v', 'u']; dists = [2, 2, 3]
5) Q = ['v', 'u', 'y']; dists = [2, 3, 3]
6) Q = ['u', 'y']; dists = [3, 3]
7) Q = ['y']; dists = [3]
8) Q = []
(['s', 'w', 'r', 't', 'x', 'v', 'u', 'y'], {'s': 0, 'w': 1, 'r': 1, 't': 2,
'x': 2, 'v': 2, 'u': 3, 'y': 3})

```

Depth first search

 image.png

 bfs-states

```

In [6]: graph = {
    's' : ['w', 'r'],
    'r' : ['v'],
    'w' : ['t', 'x'],
    'x' : ['y'],
    't' : ['u'],
    'u' : ['y']
}

def dfs(graph, start, debug=False):
    seen = set([start]) # List for seen nodes (contains both frontier and de
    # Frontier is the boundary between seen and unseen (Also called the alive
    frontier = LifoQueue() # Frontier of unvisited nodes as FIFO
    node2dist = {start : 0} # Keep track of distances
    search_order = [] # Keep track of search order
    frontier.put(start)

    i = 0 # step number
    while not frontier.empty(): # Creating loop to visit each node
        if debug: print("%d) Q = " % i, list(frontier.queue), end='; ')
        if debug: print("dists = " , [node2dist[n] for n in frontier.queue])
        m = frontier.get() # Get the oldest addition to frontier
        search_order.append(m)

        for neighbor in graph.get(m, []):
            if neighbor not in seen:
                seen.add(neighbor)
                frontier.put(neighbor)
                node2dist[neighbor] = node2dist[m] + 1
            else:
                node2dist[neighbor] = min(node2dist[neighbor], node2dist[m])
        i += 1
    if debug: print("%d) Q = " % i, list(frontier.queue))
    return search_order, node2dist

```

```

In [7]: # Driver Code
print("Following is the Depth-First Search path")
print(dfs(graph, 's', debug=True)) # function calling

```

Following is the Depth-First Search path

```

0) Q = ['s']; dists = [0]
1) Q = ['w', 'r']; dists = [1, 1]
2) Q = ['w', 'v']; dists = [1, 2]
3) Q = ['w']; dists = [1]
4) Q = ['t', 'x']; dists = [2, 2]
5) Q = ['t', 'y']; dists = [2, 3]
6) Q = ['t']; dists = [2]
7) Q = ['u']; dists = [3]
8) Q = []
(['s', 'r', 'v', 'w', 'x', 'y', 't', 'u'], {'s': 0, 'w': 1, 'r': 1, 'v': 2,
't': 2, 'x': 2, 'y': 3, 'u': 3})

```

Converting a maze search to a graph search

In [8]: *# Skip these utilities for the class*

```
def batched(iterable, n):
    "Batch data into tuples of length n. The last batch may be shorter."
    # batched('ABCDEFGH', 3) --> ABC DEF G
    if n < 1:
        raise ValueError('n must be at least one')
    it = iter(iterable)
    while batch := tuple(islice(it, n)):
        yield batch

def draw_path(self, path, visited='*'):
    new_maze_lines = [list(l) for l in self.maze_lines]
    for (r, c) in path:
        new_maze_lines[r][c] = visited
    print('\n'.join([''.join(l) for l in new_maze_lines]))
    print('\n\n\n')

def init_plots(self, reinit=False):
    if self.fig is None or reinit:
        self.fig, self.ax = plt.subplots()

def plot_maze(self):
    self.init_plots()
    replace = { ' ' : 1, '+' : 0}
    maze_mat = np.array([[replace[c] for c in line]
                          for line in self.maze_lines])
    return [self.ax.imshow(maze_mat, cmap='gray')]

def plot_step(self, i_node):
    i, (r, c) = i_node
    return [self.ax.text(c, r, '%d' % (i+1))]

def plot_path(self, path):
    self.plot_maze()
    return [self.plot_step((i, (r,c)))
            for i, (r, c) in enumerate(path)]

def animate_search_path(maze, search_path, node2dist):
    maze.init_plots()
    return animation.FuncAnimation(maze.fig, maze.plot_step, frames=[(node2c for n in search_path)
                                                                    for n in range(len(search_path))],
                                   init_func=maze.plot_maze, blit=True, repeat=False)
```

In [9]: **import** matplotlib.pyplot **as** plt

import numpy **as** np

maze_str = \

"""

+++++++

+ +

+ + + +++

+ + + +

+ + + +

```

+ + +++ +
+      + +
+ +++ + +
+  +
+++++++
"""

```

```

class Maze:
    def __init__(self, maze_str, freepath=' '):
        self.maze_lines = [l for l in maze_str.split("\n")
                           if len(l)]
        self.FREEPATH = freepath
        self.fig = None

    def get(self, node, default):
        (r, c) = node
        m_row = self.maze_lines[r]
        nbrs = []
        if c-1 >= 0 and m_row[c-1] == self.FREEPATH:
            nbrs.append((r, c-1))
        if c+1 < len(m_row) and m_row[c+1] == self.FREEPATH:
            nbrs.append((r, c+1))
        if r-1 >= 0 and self.maze_lines[r-1][c] == self.FREEPATH:
            nbrs.append((r-1, c))
        if r+1 < len(self.maze_lines) and self.maze_lines[r+1][c] == self.FREEPATH:
            nbrs.append((r+1, c))
        return nbrs if len(nbrs) else default

    init_plots = init_plots
    plot_maze = plot_maze
    plot_step = plot_step
    plot_path = plot_path
    animate_search_path = animate_search_path

```

```

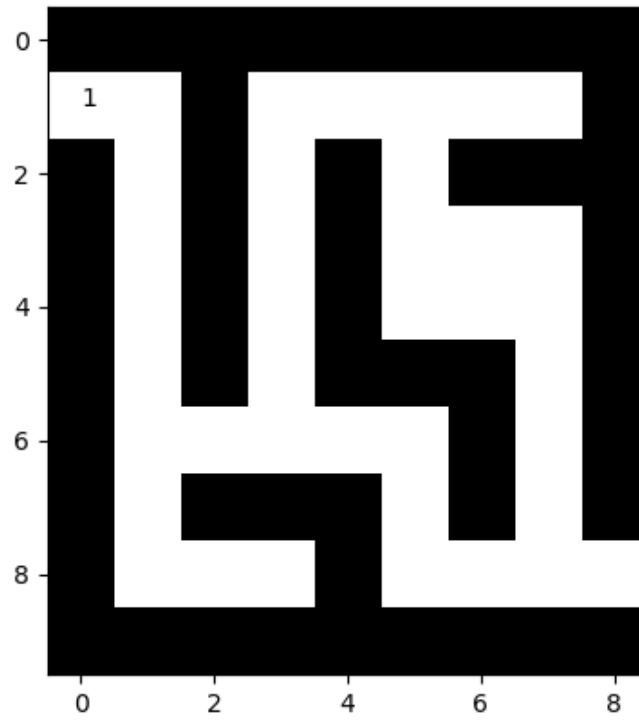
In [10]: import matplotlib.pyplot as plt
import matplotlib.animation as animation
import matplotlib as mpl
%matplotlib inline
mpl.rc('animation', html='jshtml')

maze = Maze(maze_str)
search_path, node2dist = bfs(maze, (1, 0)) # prints the order of search all
maze.plot_maze()

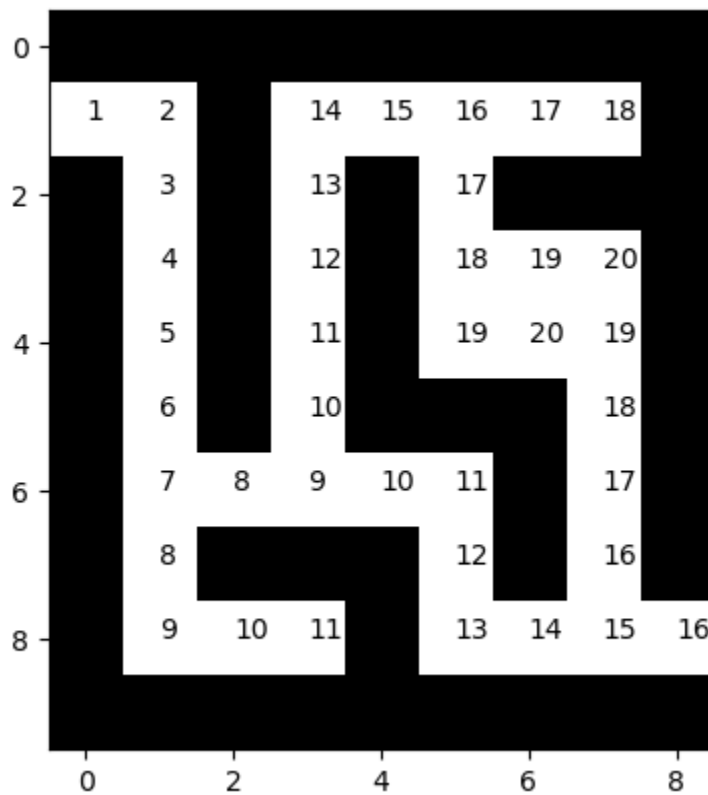
maze.animate_search_path(search_path, node2dist)

```

Out[10]:



☒ Once ☐ Loop ☐ Reflect




```
In [11]: def bfs_path(graph, start, goal):
    """
    Returns success and node2parent

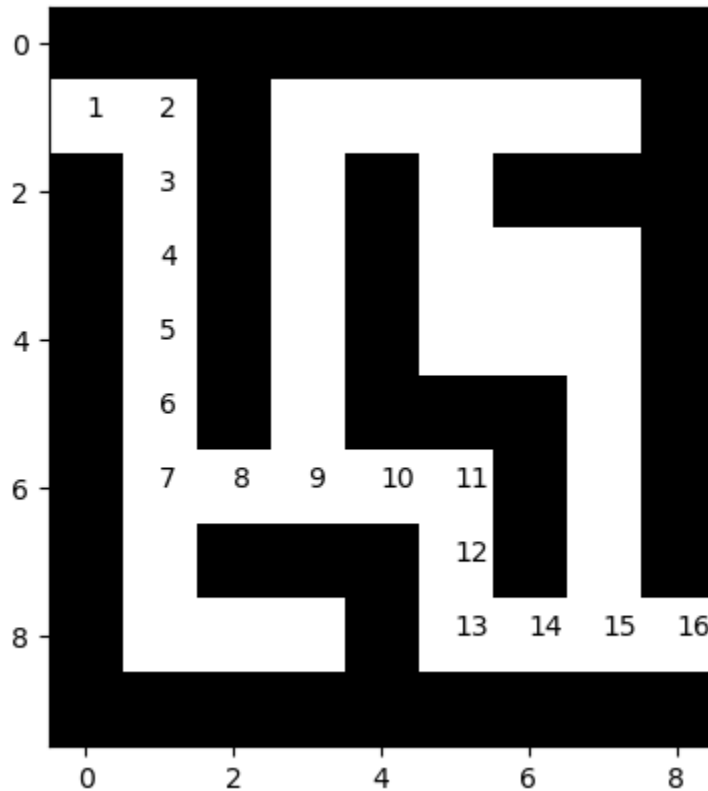
    success: True if goal is found otherwise False
    node2parent: A dictionary that contains the nearest parent for node
    """
    seen = [start] # List for seen nodes.
    # Frontier is the boundary between seen and unseen
    frontier = Queue() # Frontier of unvisited nodes as FIFO
    node2parent = dict() # Keep track of nearest parent for each node (required)
    frontier.put(start)

    while not frontier.empty(): # Creating loop to visit each node
        m = frontier.get() # Get the oldest addition to frontier
        if m == goal:
            return True, node2parent

        for neighbor in graph.get(m, []):
            if neighbor not in seen:
                seen.append(neighbor)
                frontier.put(neighbor)
                node2parent[neighbor] = m
    return False, []
```

```
In [12]: def backtrace_path(node2parent, start, goal):
    c = goal
    r_path = [c]
    parent = node2parent.get(c, None)
    while parent != start:
        r_path.append(parent)
        c = parent
        parent = node2parent.get(c, None) # Keep getting the parent until you reach start
        #print(parent)
    r_path.append(start)
    return reversed(r_path) # Reverses the path

maze = Maze(maze_str)
start = (1, 0)
goal = (8, 8)
success, node2parent = bfs_path(maze, (1, 0), (8, 8))
path = backtrace_path(node2parent, (1, 0), (8, 8))
#print(list(path))
maze.plot_path(path) # Draws all the searched nodes
plt.show()
#node2parent
```



Dijkstra algorithm



PriorityQueue

PriorityQueue returns the smallest (or the largest) item in the queue faster than other data structures

```
In [13]: #from queue import PriorityQueue
from hw2_solution import PriorityQueueUpdatable
from dataclasses import dataclass, field
from typing import Any

# https://docs.python.org/3/library/queue.html#queue.PriorityQueue
@dataclass(order=True)
class PItem:
    dist: int
    node: Any=field(compare=False)

    # Make the PItem hashable
    # https://docs.python.org/3/glossary.html#term-hashable
    def __hash__(self):
        return hash(self.node)

graph = {
    's' : [('x', 5), ('u', 10)],
```

```

'u' : [('v', 1), ('x', 2)],
'x' : [('u', 3), ('v', 9), ('y', 2)],
'y' : [('v', 6), ('s', 7)],
'v' : [('y', 4)]
}

def dijkstra(graph, start, goal, debug=False):
    """
    edgecost: cost of traversing each edge

    Returns success and node2parent

    success: True if goal is found otherwise False
    node2parent: A dictionary that contains the nearest parent for node
    """
    seen = set([start]) # Set for seen nodes.
    # Frontier is the boundary between seen and unseen
    frontier = PriorityQueueUpdatable() # Frontier of unvisited nodes as a P
    node2parent = {start : None} # Keep track of nearest parent for each node
    node2dist = {start: 0} # Keep track of cost to arrive at each node
    search_order = []
    frontier.put(PItem(0, start))
    i = 0
    while not frontier.empty():
        # Creating loop to visit each node
        dist_m = frontier.get() # Get the smallest addition to the frontier
        if debug: print("%d) Q = " % i, list(frontier.queue), end='; ')
        if debug: print("dists = " , [node2dist[n.node] for n in frontier.queue])
        m = dist_m.node
        m_dist = node2dist[m]
        search_order.append(m)
        if goal is not None and m == goal:
            return True, search_order, node2parent, node2dist

        for neighbor, edge_cost in graph.get(m, []):
            old_dist = node2dist.get(neighbor, float("inf"))
            new_dist = edge_cost + m_dist
            if neighbor not in seen:
                seen.add(neighbor)
                frontier.put(PItem(new_dist, neighbor))
                node2parent[neighbor] = m
                node2dist[neighbor] = new_dist
            elif new_dist < old_dist:
                node2parent[neighbor] = m
                node2dist[neighbor] = new_dist
                # ideally you would update the dist of this item in the priority queue
                # as well. But python priority queue does not support fast updates
                old_item = PItem(old_dist, neighbor)
                if old_item in frontier:
                    frontier.replace(old_item, PItem(new_dist, neighbor))

        i += 1
    if goal is not None:
        return False, [], {}, node2dist
    else:
        return True, search_order, node2parent, node2dist

```


[illegible]

```
In [17]: import math
          from itertools import islice

          class Maze8(MazeD):
              def get(self, node, default):
                  (r, c) = node
                  rmax = len(self.maze_lines)
                  cmax = len(self.maze_lines[0])
```

```

m_row = self.maze_lines[r]
possible_nbrs = [
    ((r, c-1), 1),
    ((r, c+1), 1),
    ((r-1, c), 1),
    ((r+1, c), 1),
    ((r-1, c-1), math.sqrt(2)),
    ((r-1, c+1), math.sqrt(2)),
    ((r+1, c-1), math.sqrt(2)),
    ((r+1, c+1), math.sqrt(2))
]
free_nbrs = []
for (ri, ci), dist in possible_nbrs:
    if (ri >= 0 and ci >= 0 and ri < rmax and ci < cmax
        and self.maze_lines[ri][ci] == self.FREEPATH):
        free_nbrs.append(((ri, ci), dist))
return free_nbrs if len(free_nbrs) else default

def _plot_path(self, path, char='+', color='c'):
    return [self.ax.text(c-0.5, r+0.5, char, color=color)
            for (r, c) in path]

def plot_path(self, path, **kw):
    self.plot_maze()
    return self._plot_path(path, **kw)

def animate(self, path, batch_size=200):
    self.init_plots()
    anim = animation.FuncAnimation(self.fig, self._plot_path,
                                   frames=batched(search_path, batch_size),
                                   init_func=self.plot_maze, blit=True, repeat=False)

    return anim

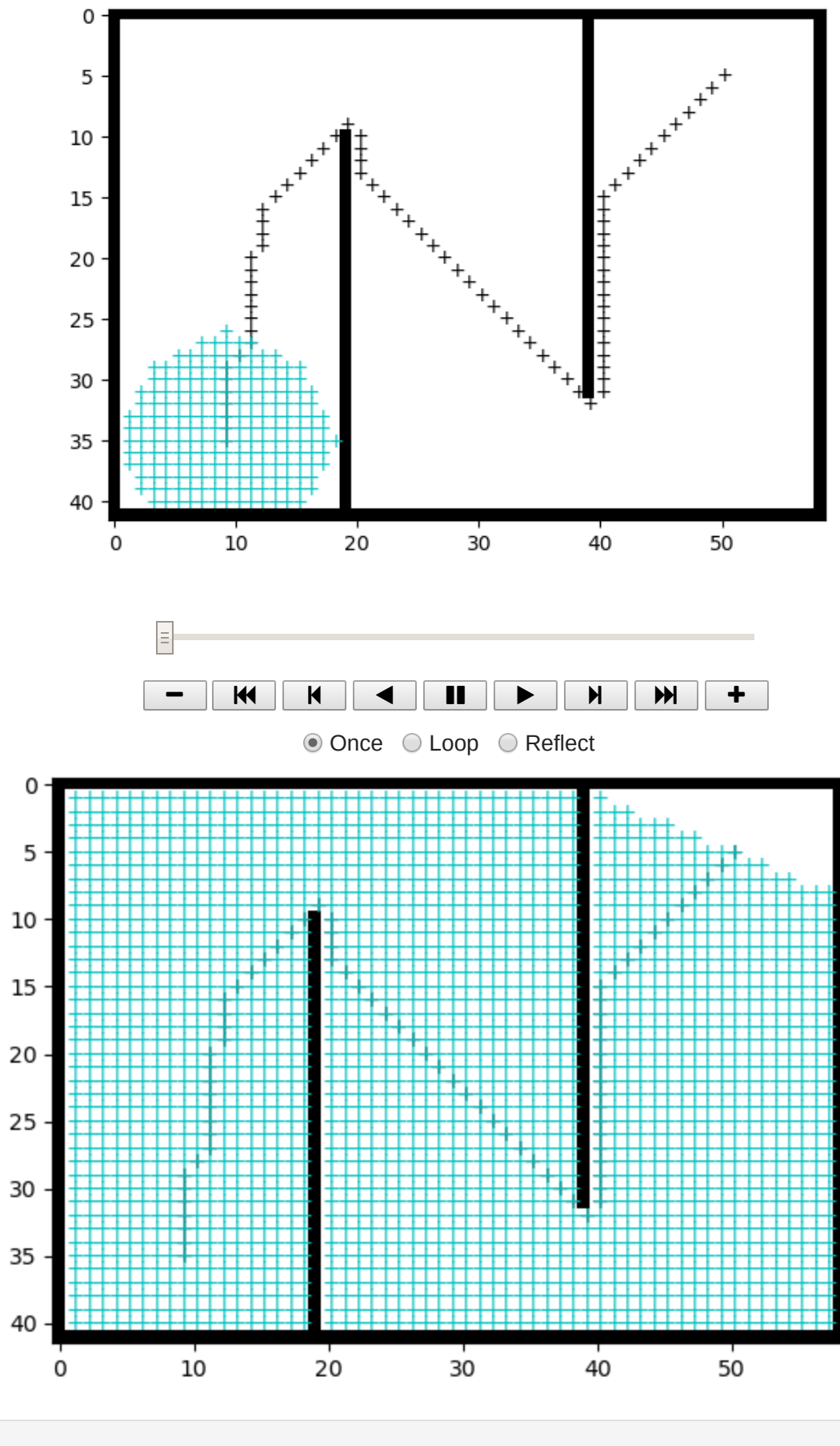
```

```

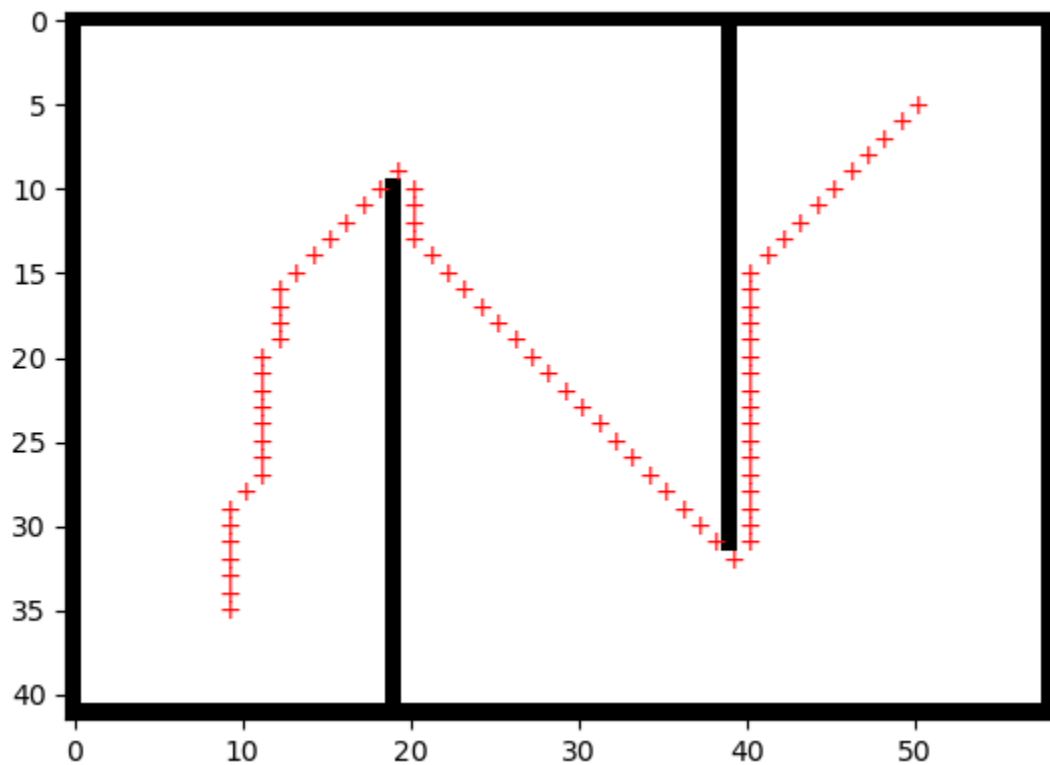
In [18]: maze = Maze8(maze_str)
success, search_path, node2parent, node2dist = dijkstra(maze, start_pos, goal_pos)
#print(success, search_path)
assert success
anim = maze.animate(search_path)
path = backtrace_path(node2parent, start_pos, goal_pos)
#maze.init_plots(reinit=True)
path_plot = maze.plot_path(path, color='k') # Draws the traced shortest path
anim

```

Out[18]:

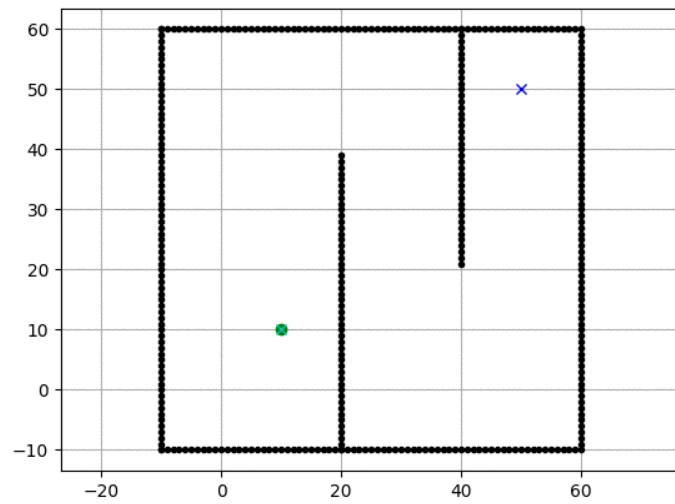
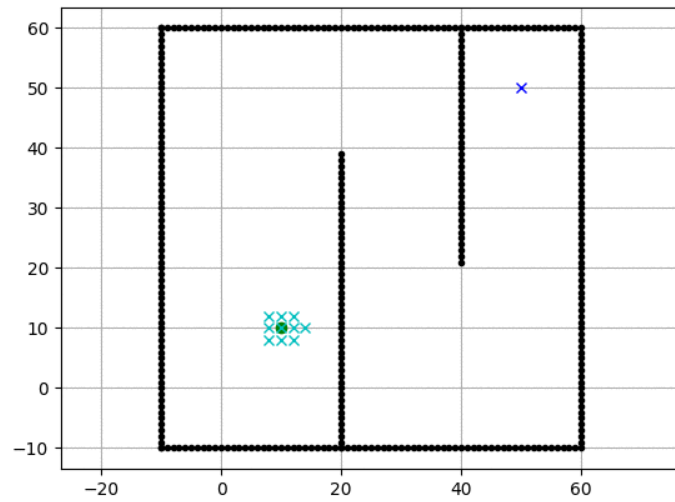


```
In [19]: path = backtrace_path(node2parent, (35, 9), (5, 50))
maze.init_plots(reinit=True)
maze.plot_path(path, color='r') # Draws the traced shortest path
plt.show()
```

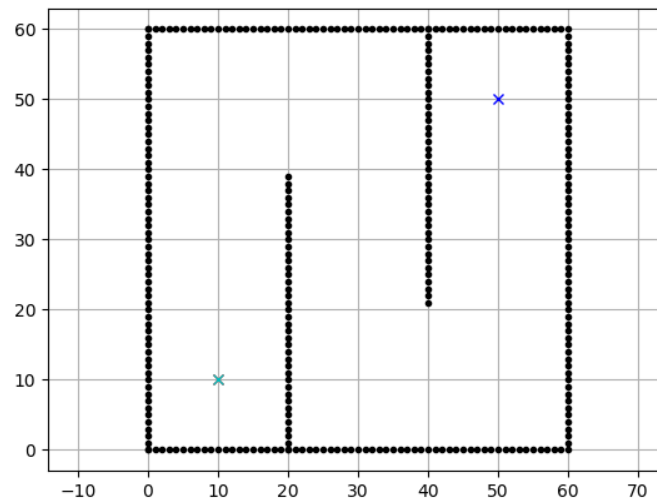
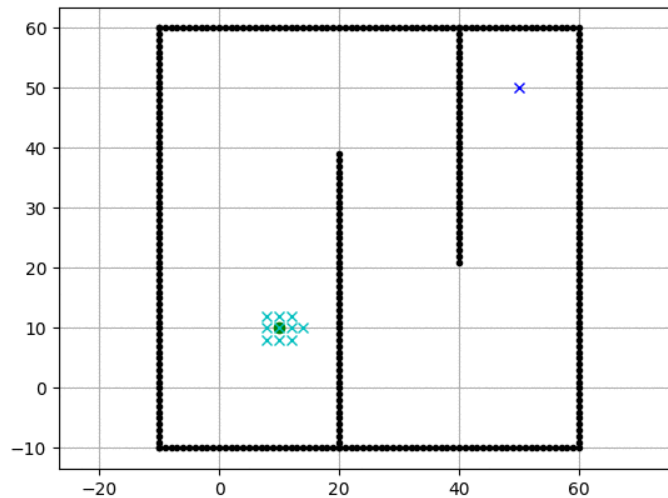


Search order in BFS vs DFS vs Dijkstra

Breadth first search vs Depth first search



Breadth first search vs Dijkstra



Computational complexity of BFS

```
In [20]: # Write down the computational complexity of each line in big-O notation  $O()$ 
# Assume the graph has  $|V|$  nodes and  $|E|$  edges
def bfs_barebones(graph, start):
    seen = {start} # Set for seen nodes (contains both frontier and dead states)
    # Frontier is the boundary between seen and unseen (Also called the alive frontier)
    frontier = Queue() # Frontier of unvisited nodes as FIFO #  $O(1)$ 
    frontier.put(start) #  $O(1)$ 

    while not frontier.empty(): # Creating loop to visit each node #  $O(|V|)$ 
        m = frontier.get() # Get the oldest addition to frontier #  $O(|V| * 1)$ 

        for neighbor in graph.get(m, []): #  $O(|V| * |E|/|V|) = O(|E|)$ 
            if neighbor not in seen: #  $O(|E| * 1)$ 
                seen.add(neighbor) #  $O(|E| * 1)$ 
                frontier.put(neighbor) #  $O(|E| * 1)$ 
```

The computational complexity of BFS is $O(|E|)$. Some books write it as $O(|V|)$
 # where $O(|V|)$ is the cost of initializing states of different nodes

Computational complexity of Dijkstra

In [21]: # Write down the computational complexity of each line in big-O notation $O()$
 # Assume the graph has $|V|$ nodes and $|E|$ edges
 def dijkstra_barebones(graph, start):
 seen = {start} # Set for seen nodes (contains both frontier and dead states)
 # Frontier is the boundary between seen and unseen (Also called the alive)
 frontier = PriorityQueue() # Frontier of unvisited nodes as PriorityQueue
 frontier.put(PItem(0, start)) # $O(1)$
 node2dist = {start: 0} # Keep track of cost to arrive at each node # $O(1)$

 while not frontier.empty(): # Creating loop to visit each node $O(1)$
 dist_and_node = frontier.get() # Get the smallest dist node # $O(|V|)$
 m_dist = dist_and_node.dist # $O(1)$
 m = dist_and_node.node # $O(1)$

 for neighbor, edge_dist in graph.get(m, []): # $O(|V| * |E|/|V|) = O(|E|)$
 if neighbor not in seen: # $O(|E| * 1)$
 seen.add(neighbor) # $O(|E| * 1)$
 frontier.put(neighbor) # $O(1)$
 node2dist[neighbor] = m_dist + edge_dist # $O(1)$
 elif node2dist[neighbor] > m_dist + edge_dist: # $O(1)$
 node2dist[neighbor] = m_dist + edge_dist # $O(1)$
 frontier.replace(neighbor, m_dist + edge_dist) # $O(1)$

 # The computational complexity of Dijkstra is $O(|V|\log(|V|) + |E|)$ when implemented using a Fibonacci heap based PriorityQueue

$n = \text{size of frontier}$
 worst case $= |V|$

$O(|V|)$

$O(|E|)$

$O(|V|\log(|V|) + |E|\log(|V|))$

PriorityQueue (Heaps Chapter 7 of Carmen's intro to algorithms)



Heap property

1. $H[\text{Parent}(i)] \geq H[i]$
2. $\text{Parent}(i) = \text{ceil}(i/2)$
3. $\text{LeftChild}(i) = 2i$
4. $\text{RightChild}(i) = 2i+1$

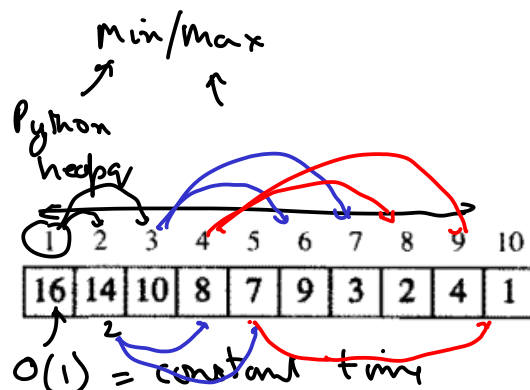
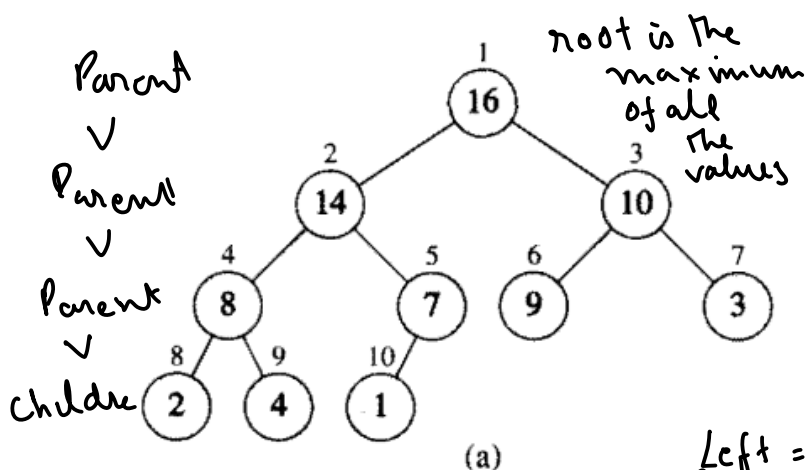
Heapify



Heapify pseudocode



Fibonacci heap
 insertion $O(1)$
 extract $O(\log n)$
 $O(|V|\log(|V|) + |E|)$



Left = $2 \times i = 2$ (b) Parent = $\text{floor}(i/2)$
Right = $2 \times i + 1 = 3$

Figure 7.1 A heap viewed as (a) a binary tree and (b) an array. The number within the circle at each node in the tree is the value stored at that node. The number next to a node is the corresponding index in the array.

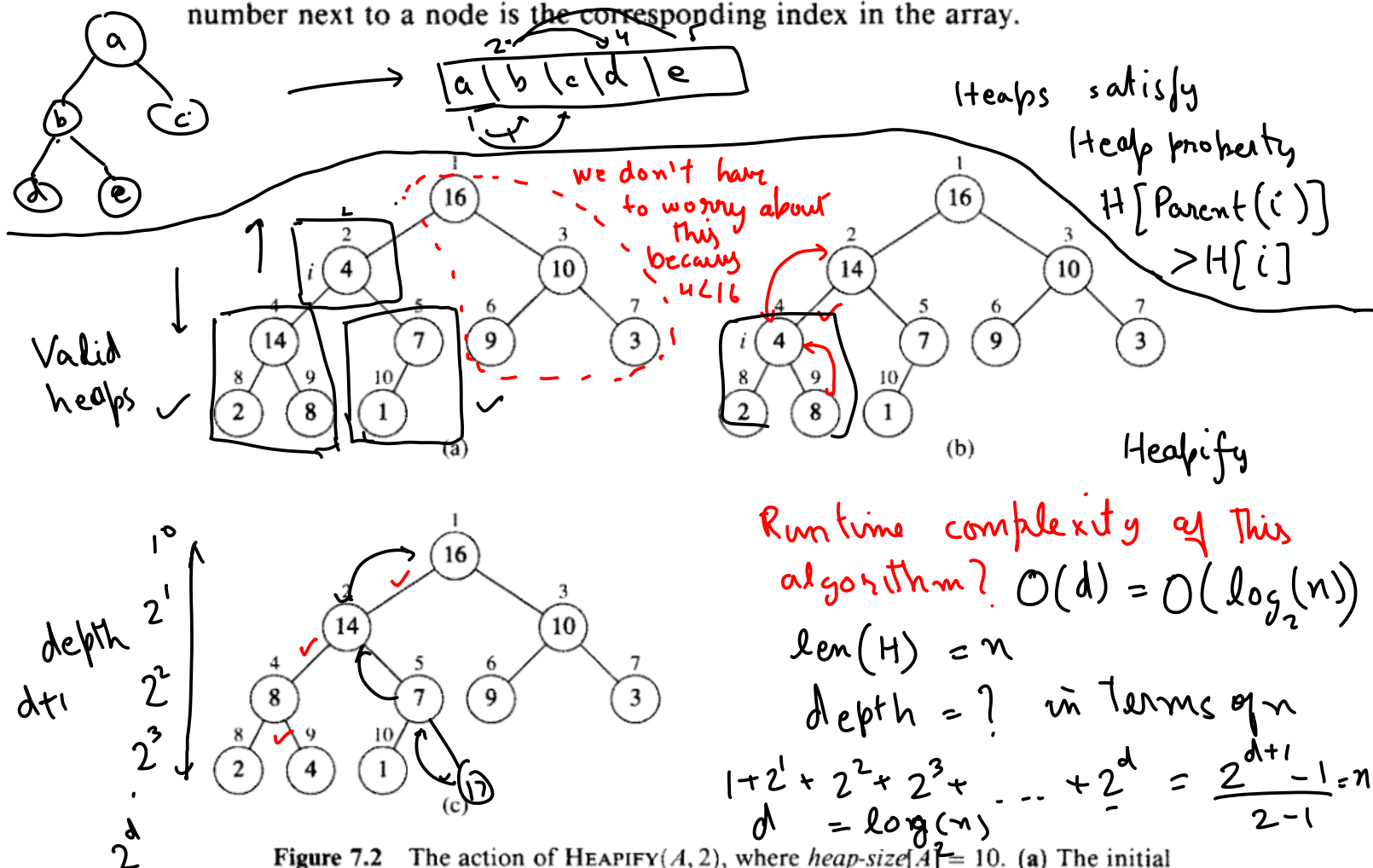


Figure 7.2 The action of $\text{HEAPIFY}(A, 2)$, where $\text{heap-size}[A] = 10$. (a) The initial configuration of the heap, with $A[2]$ at node $i = 2$ violating the heap property since it is not larger than both children. The heap property is restored for node 2 in (b) by exchanging $A[2]$ with $A[4]$, which destroys the heap property for node 4. The recursive call $\text{HEAPIFY}(A, 4)$ now sets $i = 4$. After swapping $A[4]$ with $A[9]$, as shown in (c), node 4 is fixed up, and the recursive call $\text{HEAPIFY}(A, 9)$ yields no further change to the data structure.

$$O(e^n) > O(n^3) > O(n) > O(\log(n)) > O(1)$$

```

HEAPIFY(A, i)
1  l ← LEFT(i) = 2i
2  r ← RIGHT(i) = 2i+1
3  if l ≤ heap-size[A] and A[l] > A[i]
4  then largest ← l
5  else largest ← i
6  if r ≤ heap-size[A] and A[r] > A[largest]
7  then largest ← r
8  if largest ≠ i
9  then exchange A[i] ↔ A[largest]
10 then HEAPIFY(A, largest)

```

recursive call
with l on n

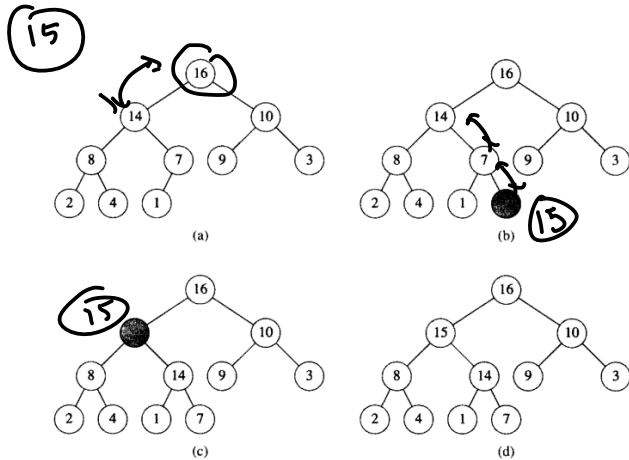
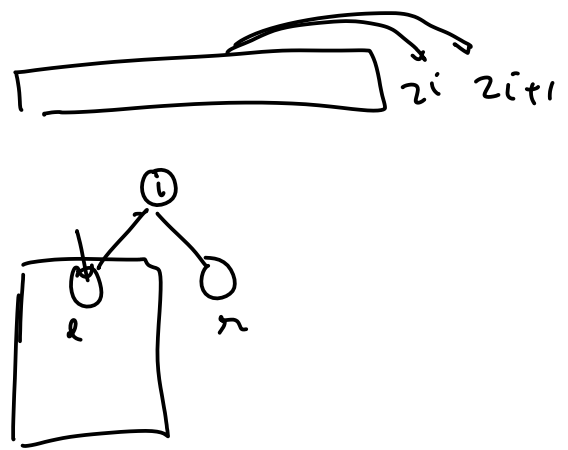
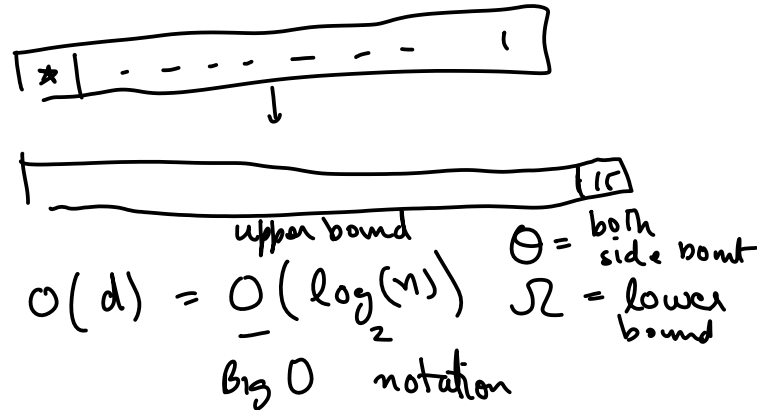


Figure 7.5 The operation of HEAP-INSERT. (a) The heap of Figure 7.4(a) before we insert a node with key 15. (b) A new leaf is added to the tree. (c) Values on the path from the new leaf to the root are copied down until a place for the key 15 is found. (d) The key 15 is inserted.

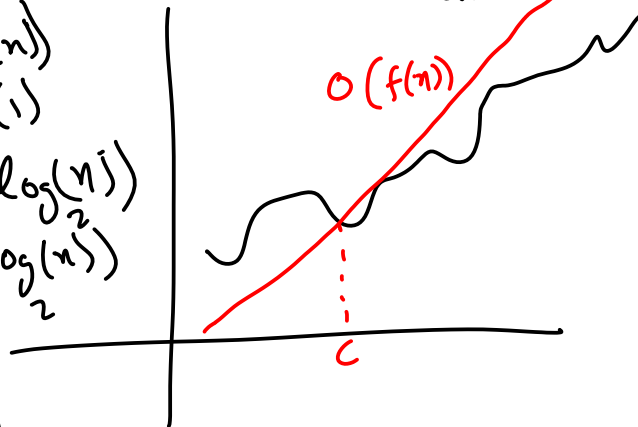


disregard constant factors and additions

Heap

Insert item in $O(\log(n))$
 Find maximum in $O(1)$
 Remove item in $O(\log(n))$
 Replace item in $O(\log^2(n))$

Replace does not have
to change the size of
array



Heap Insert



Heap runtimes



A-star (A*) algorithm

```
In [22]: from hw2_solution import PriorityQueueUpdatable
import sys
def astar(graph, heuristic_dist_fn, start, goal, debug=False, debugf=sys.stdout)
    """
    edgecost: cost of traversing each edge

    Returns success and node2parent

    success: True if goal is found otherwise False
    node2parent: A dictionary that contains the nearest parent for node
    """
    seen = set([start]) # Set for seen nodes.
    # Frontier is the boundary between seen and unseen
    frontier = PriorityQueueUpdatable() # Frontier of unvisited nodes as a P
    node2parent = {start : None} # Keep track of nearest parent for each node
    hfn = heuristic_dist_fn # make the name shorter
    node2dist = {start: 0 } # Keep track of cost to arrive at each node
    search_order = []
    frontier.put(PItem(0 + hfn(start, goal), start)) # <----- Diff

    if debug: debugf.write("goal = " + str(goal) + '\n')
    i = 0
    while not frontier.empty(): # Creating loop to visit each node
        dist_m = frontier.get() # Get the smallest addition to the frontier
        if debug: debugf.write("%d) Q = " % i + str(list(frontier.queue)) +
        if debug: debugf.write("%d) node = " % i + str(dist_m) + '\n')
        #if debug: print("dists = " , [node2dist[n.node] for n in frontier.queue])
        m = dist_m.node
        m_dist = node2dist[m]
        search_order.append(m)
        if goal is not None and m == goal:
            return True, search_order, node2parent, node2dist

        for neighbor, edge_cost in graph.get(m, []):
            old_dist = node2dist.get(neighbor, float("inf"))
            new_dist = edge_cost + m_dist
            if neighbor not in seen:
                seen.add(neighbor)
                frontier.put(PItem(new_dist + hfn(neighbor, goal), neighbor))
                node2parent[neighbor] = m
                node2dist[neighbor] = new_dist
            elif new_dist < old_dist:
                node2parent[neighbor] = m
                node2dist[neighbor] = new_dist
```

```

        # ideally you would update the dist of this item in the priority queue
        # as well. But python priority queue does not support fast updates
        # ----- Different from dijkstra -----
        old_item = PItem(old_dist + hfn(neighbor, goal), neighbor)
        if old_item in frontier:
            frontier.replace(
                old_item,
                PItem(new_dist + hfn(neighbor, goal), neighbor))

        i += 1
    if goal is not None:
        return False, [], {}, node2dist
    else:
        return True, search_order, node2parent, node2dist

```

```

In [23]: import math
        from functools import partial

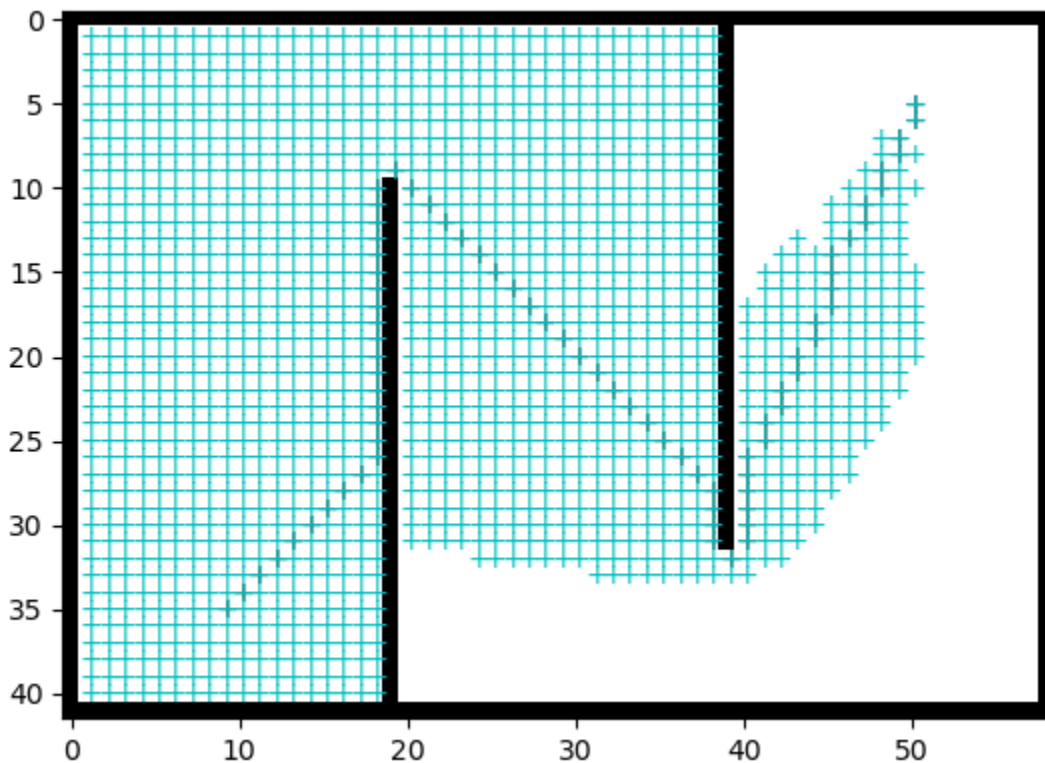
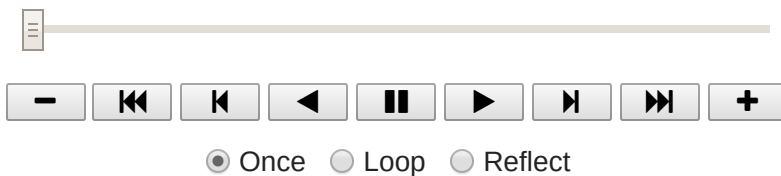
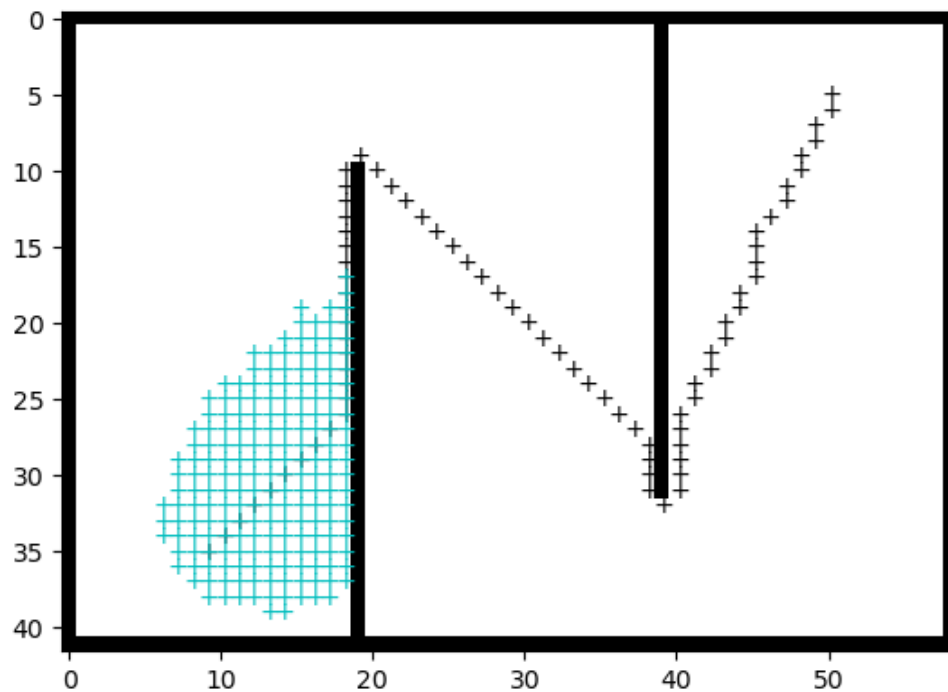
        def euclidean_heurist_dist(node, goal, scale=1):
            x_n, y_n = node
            x_g, y_g = goal
            return scale*math.sqrt((x_n-x_g)**2 + (y_n - y_g)**2)

        maze = Maze8(maze_str)
        debugf=open('log.txt', 'w')
        success, search_path, node2parent, node2dist = astar(
            maze, partial(euclidean_heurist_dist, scale=1),
            start_pos, goal_pos, debug=True, debugf=debugf)
        debugf.close()

        #print(success, search_path)
        assert success
        anim = maze.animate(search_path)
        path = backtrace_path(node2parent, start_pos, goal_pos)
        #maze.init_plots(reinit=True)
        path_plot = maze.plot_path(path, color='k') # Draws the traced shortest path
        anim

```

Out[23]:



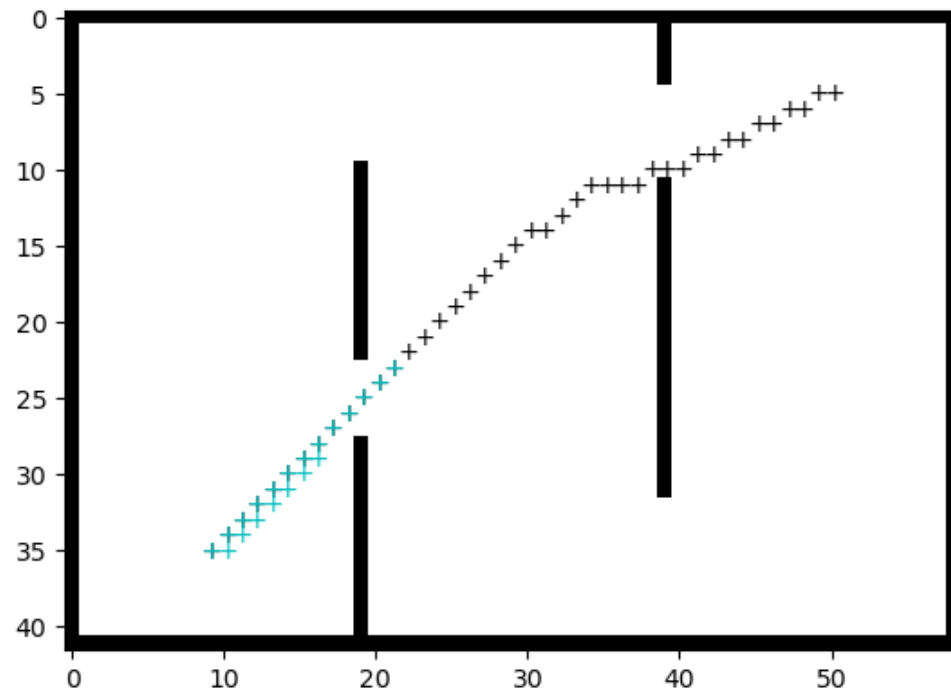

```
In [30]: maze_str = \
```

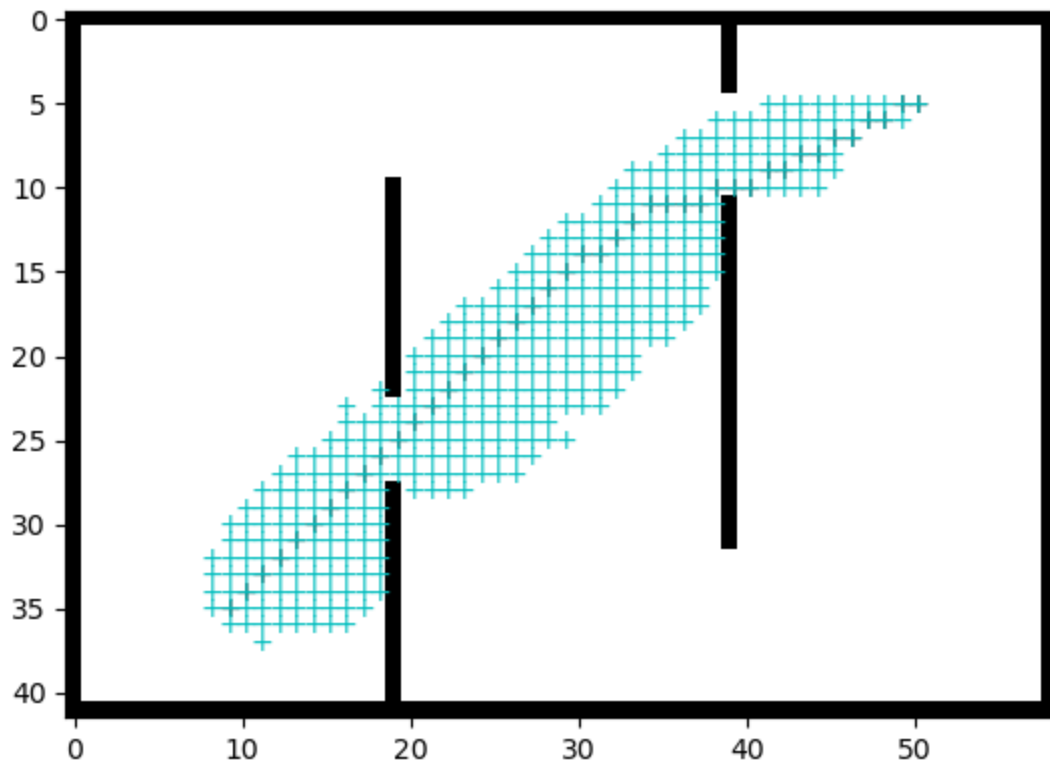
[illegible]

```
In [31]: maze = Maze8(maze_str)
          success, search_path, node2parent, node2dist = astar(
              maze, partial(euclidean_heurist_dist, scale=1),
              start_pos, goal_pos)
          #print(success, search_path)
          assert success
          anim = maze.animate(search_path, batch_size=20)
          path = backtrace_path(node2parent, start_pos, goal_pos)
```

```
#maze.init_plots(reinit=True)
path_plot = maze.plot_path(path, color='k') # Draws the traced shortest path
anim
```

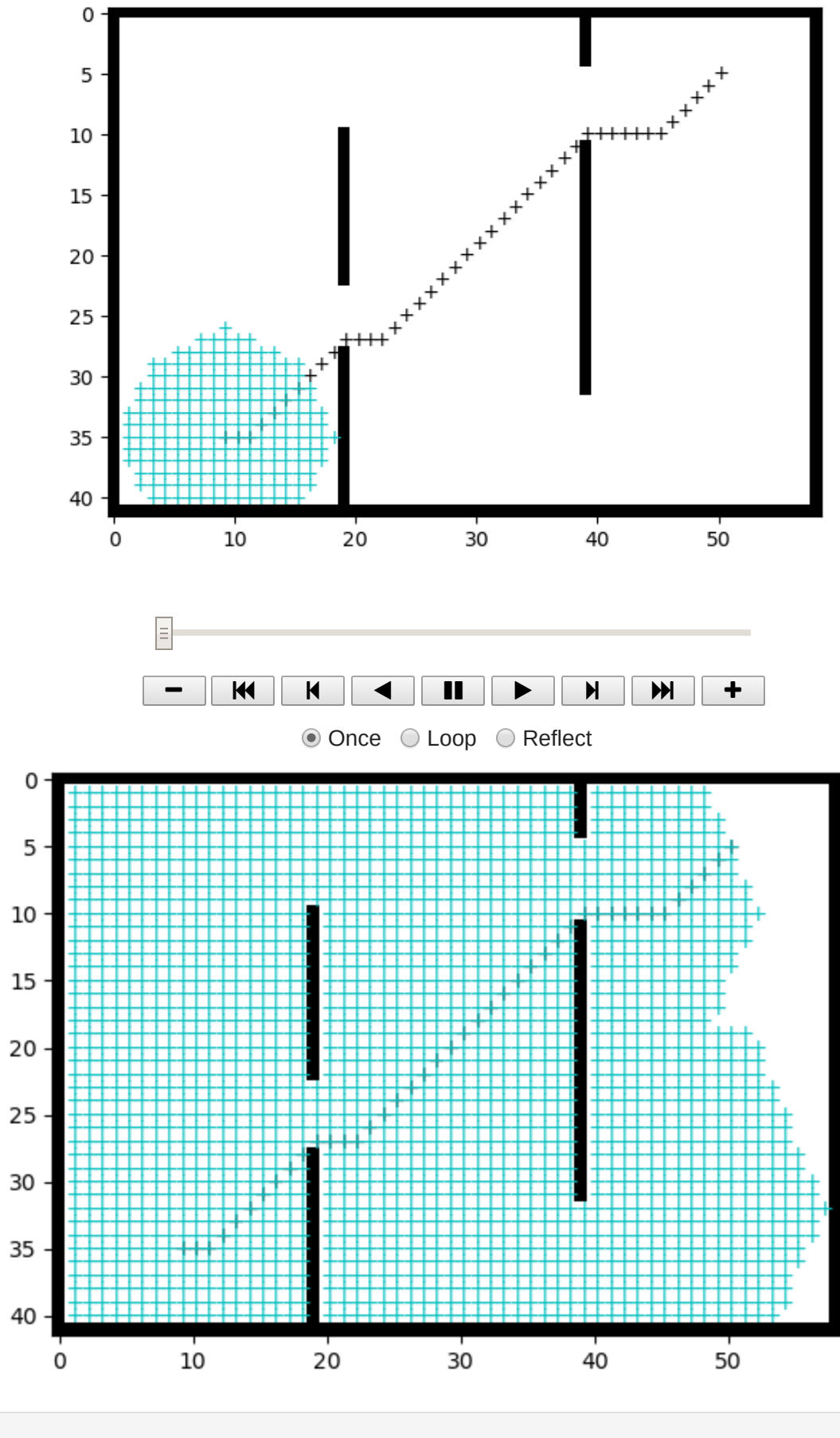
Out[31]:





```
In [32]: maze = Maze8(maze_str)
success, search_path, node2parent, node2dist = astar(
    maze, partial(euclidean_heurist_dist, scale=0),
    start_pos, goal_pos)
#print(success, search_path)
assert success
anim = maze.animate(search_path)
path = backtrace_path(node2parent, start_pos, goal_pos)
#maze.init_plots(reinit=True)
path_plot = maze.plot_path(path, color='k') # Draws the traced shortest path
anim
```

Out[32]:



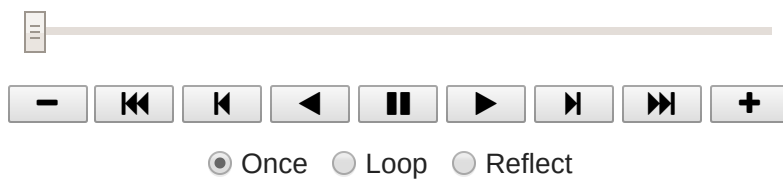
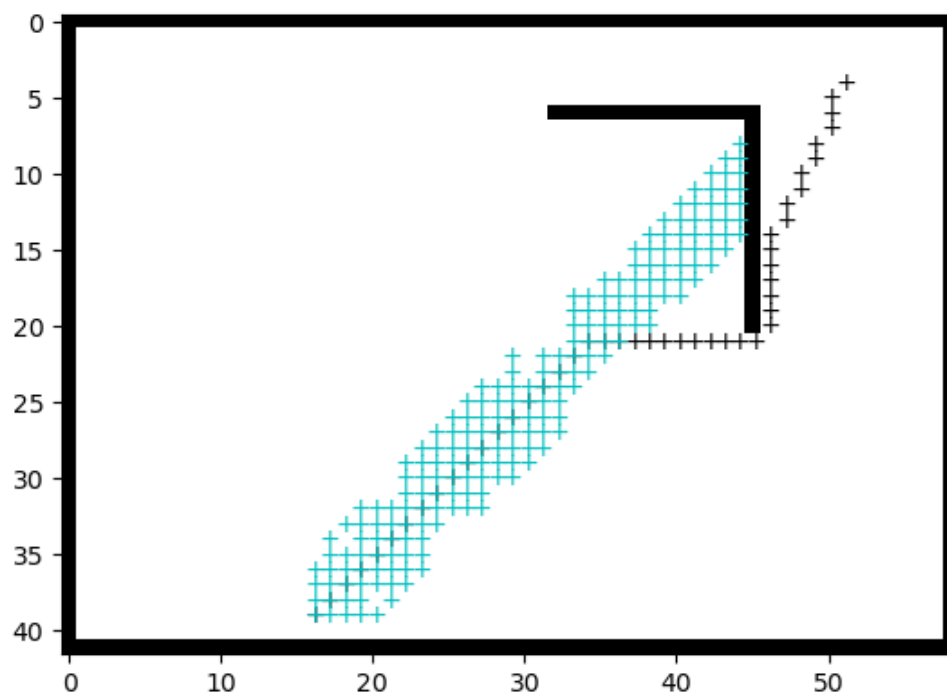
```
In [33]: maze_str = \
```

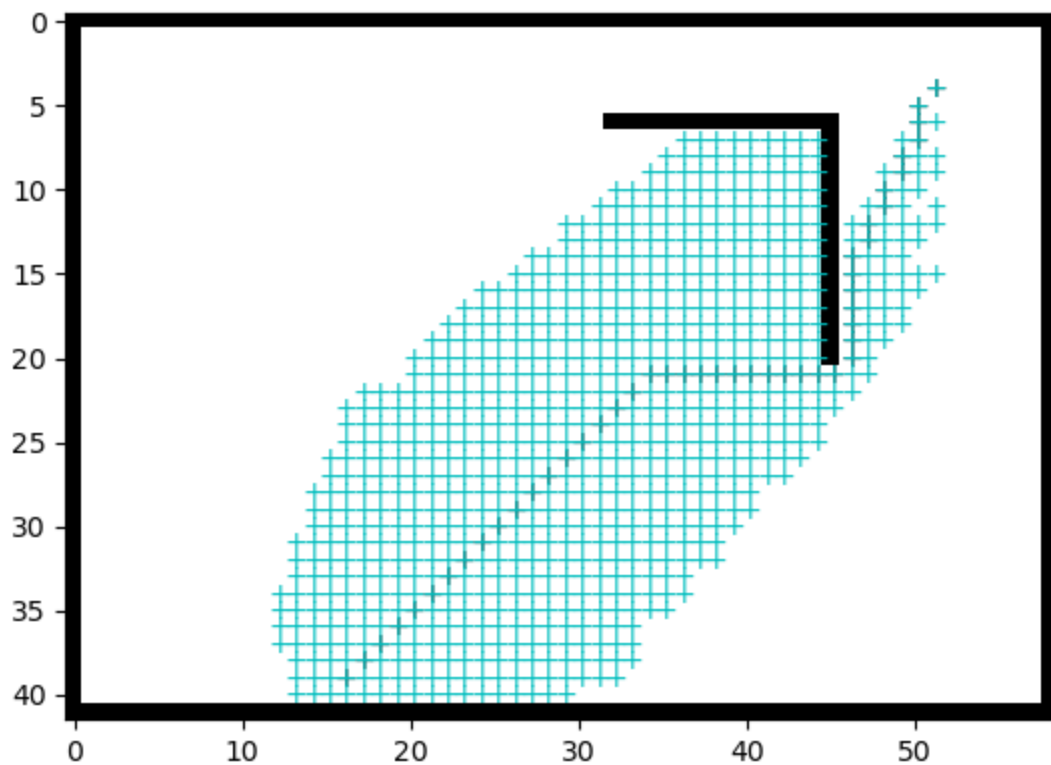
[illegible]

```
In [34]: maze = Maze8(maze_str)
          success, search_path, node2parent, node2dist = astar(
              maze, partial(euclidean_heurist_dist, scale=1),
              start_pos, goal_pos)
          #print(success, search_path)
          assert success
          anim = maze.animate(search_path)
          path = backtrace_path(node2parent, start_pos, goal_pos)
```

```
#maze.init_plots(reinit=True)
path_plot = maze.plot_path(path, color='k') # Draws the traced shortest path
anim
```

Out[34]:





In []: