CS350 Handout - Multithreaded Programming Basics (POSIX)

PART I - Programming with Threads

1. Create a Thread

pthread create()

Thread creation adds a new thread of control to the current Process. The procedure main(), itself, is a single thread of control. Each thread executes simultaneously with all the other threads within the calling process, and with other threads from other active processes.

A newly created thread shares all of the calling process' global data with the other threads in this process; however, it has its own set of attributes and private execution stack. The new thread inherits the calling thread's signal mask, possibly, and scheduling priority. Pending signals for a new thread are not inherited and will be empty.

<code>arg</code> - the argument passed to <code>start_routine</code>. Only one argument can be passed. If more than one argument needs to be passed to <code>start_routine</code>, the arguments can be packed into a structure and the address of that structure can be passed to <code>arg</code>.

2. Get the Thread Identifier

```
pthread_self() - Gets the ID of the calling thread.
-----
#include <pthread.h>
pthread_t pthread_self( void );
```

3. Yield Thread Execution

pthread_yield() - Causes the current thread to yield its execution in favor of another thread
with the same or greater priority.

```
#include <pthread.h>
void pthread_yield( void );
```

5. Send a Signal to a Thread

```
pthread_kill() - sends a signal to a thread.
```

```
#include <sys/types.h>
#include <signal.h>
int kill( pid_t target, int sig );
int tkill( pid_t tid, int sig );
```

kill() sends the signal sig to the thread specified by target. target must be a process. The sig argument must be from the list given in signal.

tkill() sends the signal sig to the thread specified by tid. tid must be a thread within the same process as the calling thread. The sig argument must be from the list given in signal. tkill() is **Linux specific**. Do not use tkill() if you plan on making your program portable.

6. Terminate a Thread

pthread_exit() - terminates a thread.

```
#include <pthread.h>
void pthread_exit( void *status );
```

The pthread_exit() function terminates the calling thread. All thread-specific data bindings are released. If the calling thread is not detached, the thread's ID and the exit status specified by status are retained until the thread is waited for. Otherwise, status is ignored and the thread's ID can be reclaimed immediately.

7. Wait for Thread Termination

pthread join() - waits for a thread to terminate.

```
#include <pthread.h>
int pthread_join( pthread_t wait_for, void **status );
```

The pthread_join() function blocks the calling thread until the thread specified by wait_for terminates. The specified thread must be in the current process and must not be detached. When status is not NULL, it points to a location that is set to the exit status of the terminated thread when pthread join() returns successfully.

*** Even though **pthread_join()** can be used to wait for a specific thread; there is no way to wait for "any" thread by using the Pthreads library in Linux.

8. Related Functions

thr_sigsetmask() - accesses the signal mask of the calling thread
pthread_key_create(), thr_setspecific(), thr_getspecific() - maintain thread-specific
data
thr_getconcurrency(), thr_setconcurrency() - get and set thread concurrency level
thr getprio(), thr setprio() - get and set thread priority

-- See manual pages for details.

PART II - Programming with Synchronization Objects

2.1 Mutual Exclusion Locks

1. Lock a Mutex

```
pthread_mutex_lock()
-----
#include <pthread.h>
int pthead_mutex_lock( pthread_mutex_t *mp );
```

Use pthead_mutex_lock() to lock the mutex pointed by *mp*. When the mutex is already locked, the calling thread blocks until the mutex becomes available (blocked threads wait on a prioritized queue). When pthread_mutex_lock() returns, the mutex is locked and the calling thread is the owner.

2. Lock with a Nonblocking Mutex

```
pthread_mutex_trylock()
-----
#include <pthread.h>
int pthread mutex trylock( pthread mutex t *mp );
```

Use pthread_mutex_trylock() to attempt to lock the mutex pointed to by mp. This function is a nonblocking version of pthread_mutex_lock(). When the mutex is already locked, this call returns with an error number (a positive integer). Otherwise, the mutex is locked (the calling thread becomes the owner of the lock) and a '0' (zero value) is returned.

3. Unlock a Mutex

Use the pthread_mutex_unlock() to unlock the mutex pointed to by mp. The mutex must be locked and the calling thread must be the one that last locked the mutex (the owner). When any other threads are waiting for the mutex to become available, the thread at the head of the queue is unblocked.

4. Code Example

```
pthread_mutex_t count_mutex;
int count;

increment_count()
{
    pthread_mutex_lock( &count_mutex );
    count = count + 1;
    pthread_mutex_unlock( &count_mutex );
}

int get_count()
{
    int c;

    pthread_mutex_lock( &count_mutex );
    c = count;
    pthread_mutex_unlock( &count_mutex );
    return c;
}
```

In this example, increment_count() uses the mutex lock simply to assure an atomic update of the shared variable. get_count() uses the mutex lock to guarantee the synchronization when it refers to count.

5. Related Functions

```
pthead_mutex_init() - initializes a mutual exclusion lock
pthread_mutex_destroy() - destroy mutex state
```

see manual pages for details

2.2 Condition Variables

Use condition variables to atomically block threads until a particular condition is true. Always use condition variables together with a mutex lock.

1. Block on a Condition Variable

Use pthread_cond_wait() to atomically release the mutex pointed by mp and to cause the calling thread to block on the condition variable pointed to by cvp. The blocked thread can be awakened by pthread_cond_signal(), pthread_cond_broadcast(), or when interrupted by delivery of a signal.

In typical use, a condition expression is evaluated under the protection of a mutex lock. When the condition expression is false, the thread blocks on the condition variable. The condition variable is then signaled by another thread when it changes the condition value. This causes one or all of the threads waiting on the condition to unblock and to try to reacquire the mutex lock.

Because the condition can change before an awakened thread returns from pthread_cond_wait(), the condition that caused the wait must be retested before the mutex lock is acquired. The recommended test method is to write the condition check as a while loop that calls pthread cond wait().

```
pthread_mutex_lock();
  while( condition_is_false )
    pthread_cond_wait();
pthread_mutex_unlock();
```

2. Unblock a Specific Thread

```
pthread_cond_signal()
-----
#include <pthread.h>
int pthread_cond_signal( pthread_cond_t *cvp );
```

Use pthead_cond_signal() to unblock one thread that is blocked on the condition variable pointed to by <code>cvp</code>. Call <code>pthread_cond_signal()</code> under the protection of the same mutex used with the condition variable being signaled. Otherwise, the condition variable could be signaled between the test of the associated condition and blocking in <code>pthread_cond_wait()</code>, which can cause an infinite wait.

When no threads are blocked on the condition variable, the pthread cond signal() has no effect.

3. Code Example

```
______
pthread mutex t count mutex;
pthread cond t count nonzero;
unsigned int count;
decrement count()
  pthread mutex lock( &count lock );
  while( count == 0 )
    pthread cond wait ( &count nonzero, &count lock );
  count = count - 1;
  pthread mutex unlock( &count lock );
}
increment count()
  pthread mutex lock( &count lock );
  if(count == 0)
    pthread_cond_signal( &count_nonzero );
  count = count + 1;
  pthread mutex unlock( &count lock );
```

4. Related Functions

```
pthread_cond_init() - initializes a condition variable
pthread_cond_timedwait() - blocks until a specified event
pthread cond broadcast() - unblocks all threads
```

pthread_cond_destroy() - destroys condition variable state

--See manual pages for details

References
1. POSIX online manual pages