CS 301 High-Performance Computing

Lab 6

Problem 2: How does the interaction between charged particles influence their movement?

Group: 11

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1 Introduction

This problem includes calculating the final position of the charged particles under coulombic force. The problem's complexity is examined using LAB207 PC and HPC Cluster.

We can calculate the distance by,

$$\Delta d = v\Delta t + \frac{1}{2}a(\Delta t)^2 \tag{1}$$

2 Hardware Details

2.1 Hardware Details for LAB207 PCs

• Architecture: x86 64

• CPU op-mode(s): 32-bit, 64-bit

• Byte Order: Little Endian

• CPU(s): 4

• On-line CPU(s) list: 0-3

• Thread(s) per core: 1

• Core(s) per socket: 4

• Socket(s): 1

• NUMA node(s): 1

• Vendor ID: GenuineIntel

• CPU family: 6

• Model: 60

• Model name: Intel(R) Core(TM) i5-4590 CPU @ 3.30GHz

• Stepping: 3

• CPU MHz: 799.992

• CPU max MHz: 3700.0000

• CPU min MHz: 800.0000

• BogoMIPS: 6584.55

• Virtualization: VT-x

• L1d cache: 32K

• L1i cache: 32K

• L2 cache: 256K

• L3 cache: 6144K

• NUMA node0 CPU(s): 0-3

• Flags: fpu vme de pse tsc msr pae mce cx8 apic sep mtrr pge mca cmov pat pse36 clflush dts acpi mmx fxsr sse sse2 ss ht tm pbe syscall nx pdpe1gb rdtscp lm constant_tsc arch_perfmon pebs bts rep_good nopl xtopology nonstop_tsc aperfmperf eagerfpu pni pclmulqdq dtes64 monitor ds_cpl vmx smx est tm2 ssse3 fma cx16 xtpr pdcm pcid sse4_1 sse4_2 x2apic movbe popcnt tsc_deadline_timer aes xsave avx f16c rdrand lahf_lm abm epb invpcid_single tpr_shadow vnmi flexpriority ept vpid fsgsbase tsc_adjust bmi1 avx2 smep bmi2 erms invpcid xsaveopt dtherm ida arat pln pts

2.2 Hardware Details for HPC Cluster (Node gics1)

• Architecture: x86_64

• CPU op-mode(s): 32-bit, 64-bit

• Byte Order: Little Endian

• CPU(s): 16

• On-line CPU(s) list: 0-15

• Thread(s) per core: 1

• Core(s) per socket: 8

• Socket(s): 2

• NUMA node(s): 2

• Vendor ID: GenuineIntel

• CPU family: 6

• Model: 63

• Model name: Intel(R) Xeon(R) CPU E5-2640 v3 @ 2.60GHz

• Stepping: 2

• CPU MHz: 1288.726

• BogoMIPS: 5204.73

• Virtualization: VT-x

• L1d cache: 32K

• L1i cache: 32K

• L2 cache: 256K

• L3 cache: 20480K

• NUMA node0 CPU(s): 0-7

• NUMA node1 CPU(s): 8-15

3 Problem

3.1 Brief description of the problem

A space contains n charged particles with their charge and initial x and y coordinates given. We need to determine their final positions after 10 seconds. The movement of these particles is influenced by the interactions between them due to their charges. Assuming only coulomb's force is in action.

3.1.1 Algorithm description

We are given the initial postions (x, y) along with the charge q for n particles. We need to calculate the force and hence the valocity and hence the final postion of the particles after 10 seconds, with $\Delta t = 0.0001s$.

The formulas for the same are given below:

$$F_i = k \frac{|q1q2|}{r_i^2} \tag{2}$$

$$a_i = \frac{F_i}{m} \tag{3}$$

$$v_i = v_{i-1} + a_{i-1}t \tag{4}$$

$$d_i = v_{i-1}t + \frac{1}{2}a_{i-1}t^2 \tag{5}$$

3.2 The complexity of the algorithm (serial)

The number of subintervals used in the calculation for Problem is n, and the serial complexity is $O(n^2)$. This indicates that the time it takes for the method to finish will grow quadratically with n as the number of subintervals does. Because the method only must be run twice through the data to calculate the distance, this complexity is considered as quadratic. An algorithm's serial complexity gives an estimate of how long it will take to execute as the size of the input grows.

3.3 The complexity of the algorithm (parallel)

The parallelism reduces this complexity by dividing the work among t threads. Each thread computes a partial sum over a subset of the intervals. The time complexity of this parallel algorithm is $O(n^2/t)$, assuming that the work is evenly distributed among the threads and there are no overheads due to synchronization or communication.

3.4 Profiling information (using gprof)

Please note that profiling using gprof was done on codes without optimizing them as it was way faster and spontaneous and gprof was not showing any output.

3.4.1 Profiling using LAB207 PCs (with gprof)

The screenshots of profiling using the LAB207 PCs are given below for Q1 (both from the terminal and from the text file).

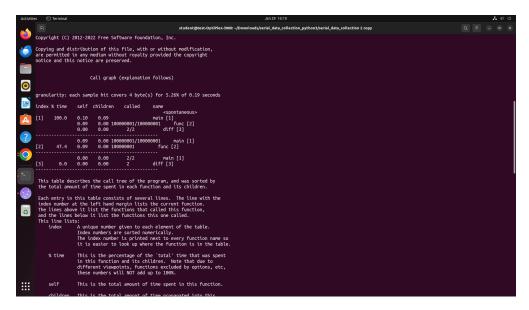


Figure 1: Screenshot of terminal from LAB207 PC

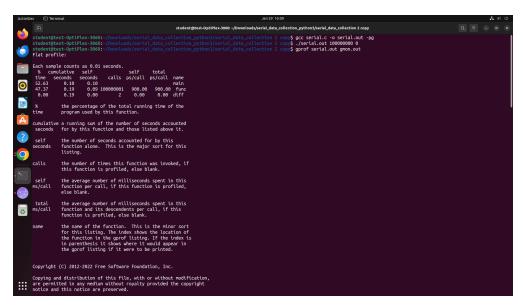


Figure 2: Screenshot of the gprof output(taken from a text file)

3.4.2 Profiling using HPC Cluster (with gprof)

The screenshots of profiling using the HPC Cluster are given below for Q2 (both from the terminal and from the text file).

```
Flat profile:

Each sample counts as 0.01 seconds.

% cumulative self self total
time seconds seconds calls Ts/call name
100.02 0.01 0.01 main
0.00 0.01 0.00 2 0.00 0.00 diff

% the percentage of the total running time of the
time program used by this function.

cumulative a running sum of the number of seconds accounted
seconds for by this function and those listed above it.

self the number of seconds accounted for by this
seconds function alone. This is the major sort for this
listing.

calls the number of times this function was invoked, if
this function is profiled, else blank.

self the average number of milliseconds spent in this
ms/call function per call, if this function is profiled,
else blank.

total the average number of milliseconds spent in this
```

function and its descendents per call, if this

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Call graph (explanation follows)

index	% time	self	children	called	name	
					<spontaneo< th=""><th></th></spontaneo<>	
[1]	100.0	0.91	0.00		main [1]	
		0.00	0.00	2/2	diff [2]	
		0.00	0.00	2/2	 main [1]	
[2]	0.0	0.00	0.00		diff [2]	

This table describes the call tree of the program, and was sorted by the total amount of time spent in each function and its children.

it is easier to look up where the function is in the table.

This is the percentage of the 'total' time that was spent

This is the percentage of the `total' time that was spent in this function and its children. Note that due to different viewpoints, functions excluded by options, etc, these numbers will NOT add up to 100%.

This is the total amount of time spent in this function.

If the function called itself recursively, the number only includes non-recursive calls, and is followed by

a '+' and the number of recursive calls.

The name of the current function. The index number is printed after it. If the function is a member of a cycle, the cycle number is printed between the

children This is the amount of time that was propagated from the function's children into this parent.

function `/' the total number of times the function was called. Recursive calls to the function are not

This is the name of the parent. The parent's index number is printed after it. If the parent is a member of a cycle, the cycle number is printed between the name and the index number. $\,$

If the parents of the function cannot be determined, the word `<spontaneous>' is printed in the `name' field, and all the other

```
For the function's children, the fields have the following meanings:

self Ihis is the amount of time that was propagated directly from the child into the function.

children Ihis is the amount of time that was propagated from the child's children to the function.

called Ihis is the number of times the function called this child '/' the total number of times the child was called. Recursive calls by the child are not listed in the number after the '/'.

name Ihis is the name of the child. The child's index number is printed after it. If the child is a member of a cycle, the cycle number is printed between the name and the index number.

If there are any cycles (circles) in the call graph, there is an entry for the cycle-as-a-whole. This entry shows who called the cycle (as parents) and the members of the cycle (as children.)

The '+' recursive calls entry shows the number of function calls that were internal to the cycle, and the calls entry for each member shows, for that member, how many times it was called from other members of the cycle.

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Index by function name
```

Figure 3: Screenshot of terminal from HPC Cluster

3.5 Optimization Strategy

The code uses OpenMP to parallelize the for loop that computes the sum of function evaluations:

```
\label{eq:pragma omp parallel for private(j)} \ensuremath{\text{for } (i=0;\ i<\text{PARTICLES};\ i++)\ \{} \\ \ensuremath{\text{forces}[i][0]} = 0.0; \\ \ensuremath{\text{for } (j=0;\ j<\text{PARTICLES};\ j++)\ \{} \\ \ensuremath{\text{if } (i!=j)\ \{} \\ \ensuremath{\text{double } dx=\text{particles}[j].x-\text{particles}[i].x; \\ \ensuremath{\text{double } dy=\text{particles}[j].y-\text{particles}[i].y; \\ \ensuremath{\text{double } r=\text{sqrt}(dx*dx+dy*dy); \\ \ensuremath{\text{double } forceMagnitude=(k*\text{particles}[i].charge*\text{particles}[j].charge)/(r*r); \\ \ensuremath{\text{forces}[i][0]+=\text{forceMagnitude*}(dx/r); //\text{Force in x-direction} \\ \ensuremath{\text{forces}[i][1]+=\text{forceMagnitude*}(dy/r); //\text{Force in y-direction} \\ \ensuremath{\}} \\ \ensuremath{} \ensuremat
```

The #pragma omp parallel for directive instructs the compiler to generate multi-threaded code for this loop. The private(i) clause specifies that the variable i should be private to each thread. This means that each thread has its own copy of this variable and can modify it independently without interfering with other threads.

3.6 Graph of Number of Processors vs Runtime

3.6.1 Graph of Number of Processors vs Algorithm Runtime for LAB207 PC

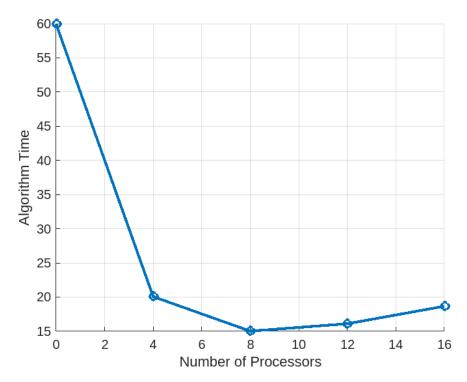


Figure 4: Algorithm execution time vs Number of Processors

3.6.2 Graph of Number of Processors vs Total Runtime for LAB207 PC

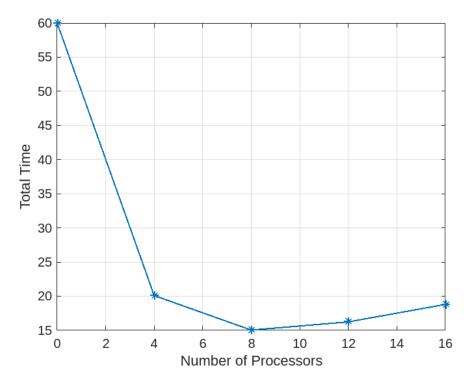


Figure 5: Total execution time vs Number of Processors

3.6.3 Graph of Number of Processors vs Algorithm Runtime for Cluster

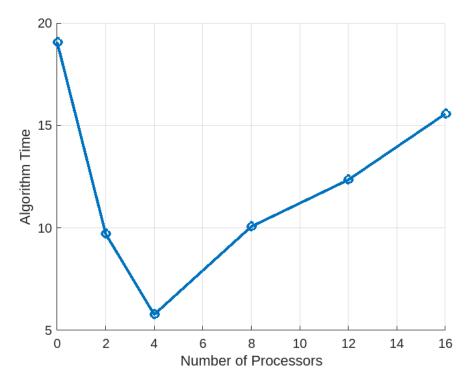


Figure 6: Algorithm execution time vs Number of Processors

3.6.4 Graph of Number of Processors vs Total Runtime for Cluster

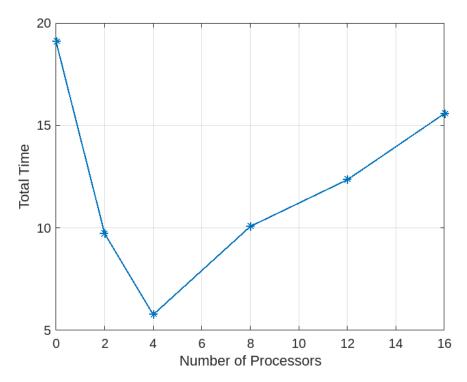


Figure 7: Total execution time vs Number of Processors

4 Conclusions

- We can observe that we get minimum time for 8 processors in Lab207PC.
- We can observe that we get minimum time for 4 processors in Cluster.