

# ROUTING IN PURE OPENFLOW ENVIRONMENT

by can.

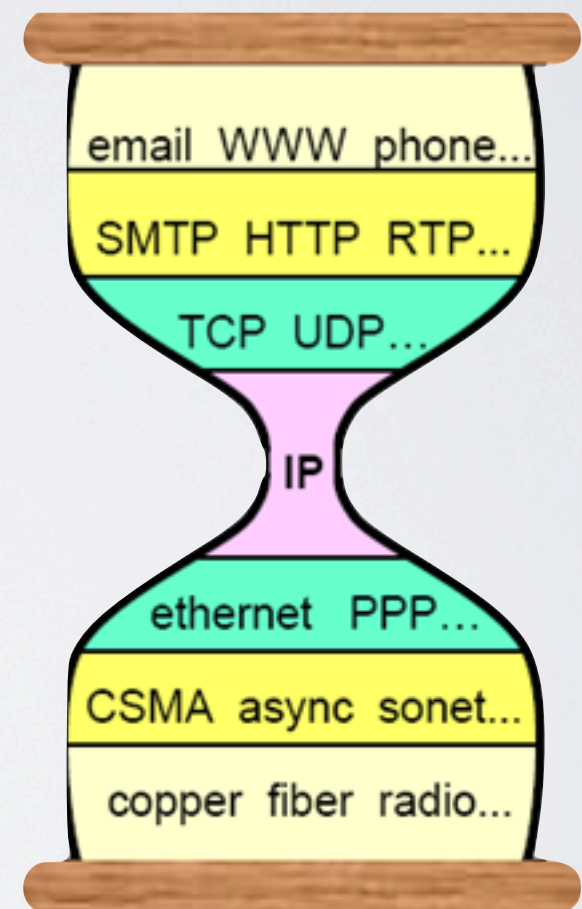
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some  
backgrounds



# WHY EXISTS IP?

- the *Internet* Protocol
- supposed to be globally unique
- “everything over IP”
- designed with routing in mind



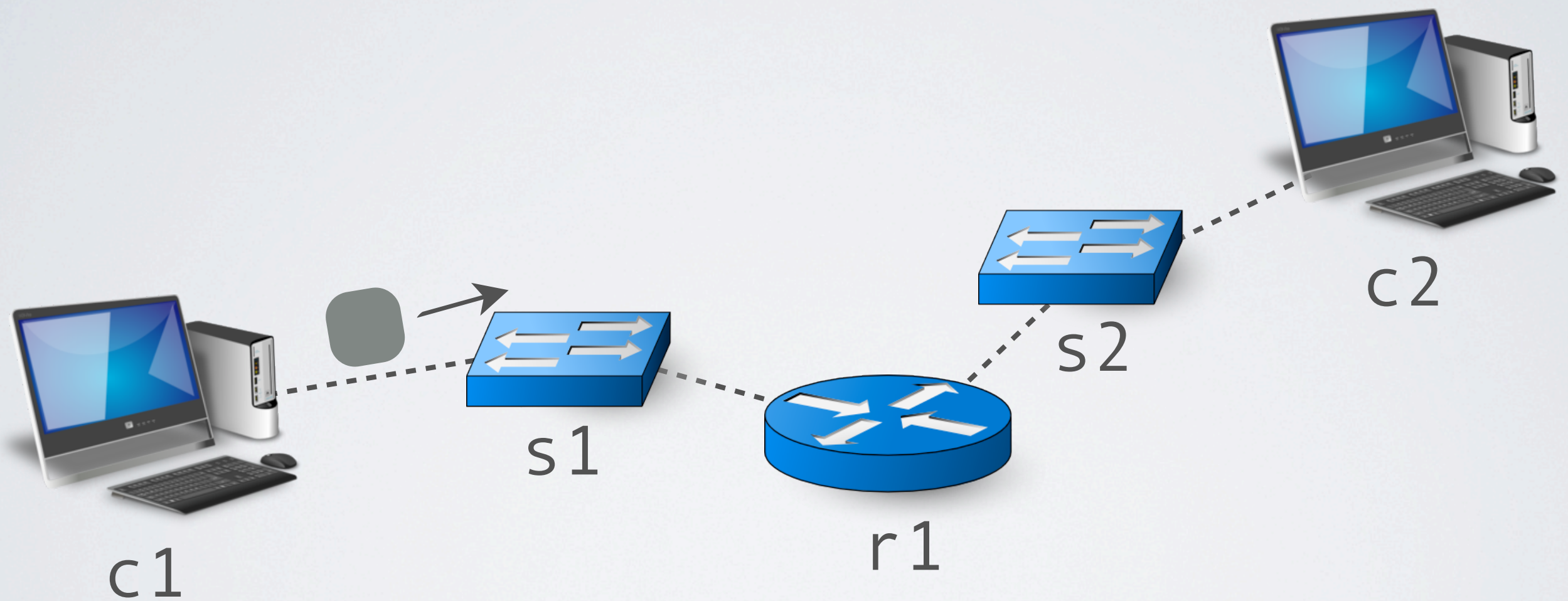


# WHY EXISTS MAC?

- by MAC I mean *Media Access Control* layer
- in a wired switching ethernet, the MAC layer is really doing nothing
- but in other types of networks, especially wireless ones, it means a lot

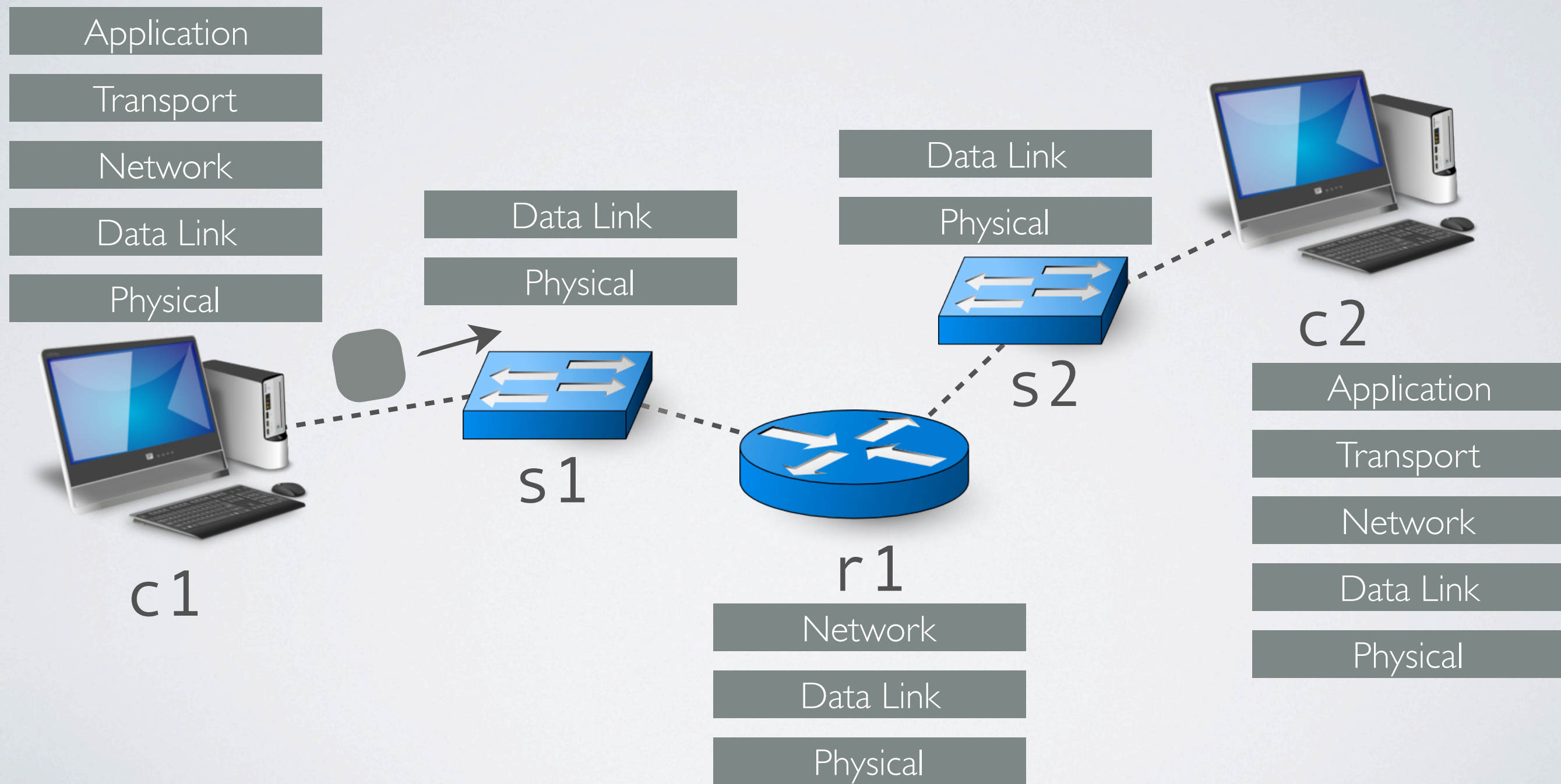


# HOW A PACKET TRAVELS?





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# HOW A PACKET TRAVELS?

Application

Transport

src[c1] dst[c2]

dst[unknown]

Physical

p1

packet p1: c1 => c2

gateway[c1] = r1.left

gateway[c2] = r1.right

every table of any devices are empty

Application

Transport

Network

Data Link

Physical

c1

Data Link

s1

Network

Data Link

Physical

r1

Data Link

Physical

s2

Application

Transport

Network

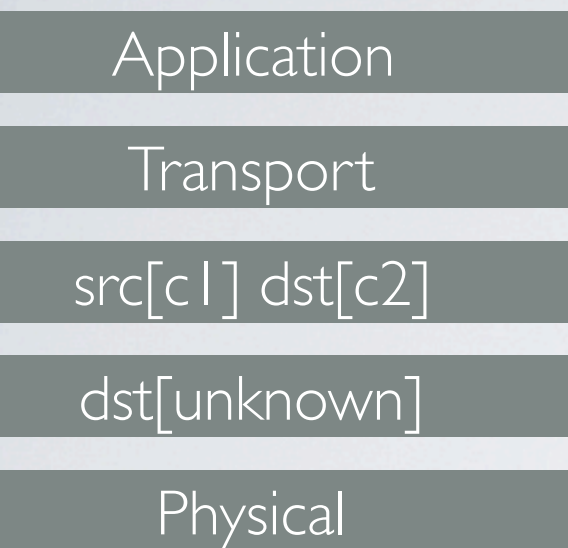
Data Link

Physical

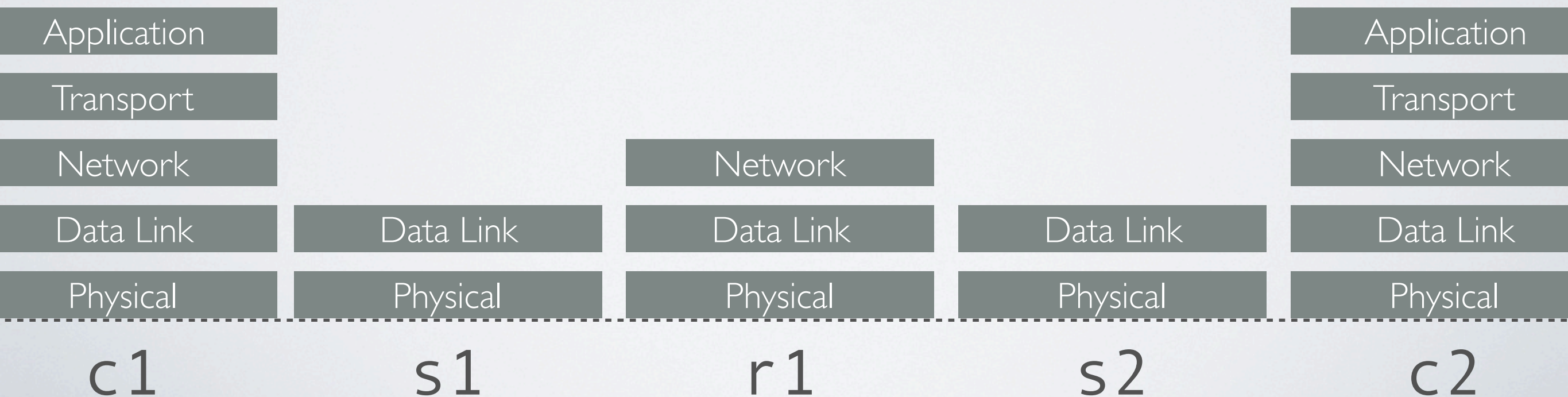
c2



# HOW A PACKET TRAVELS?

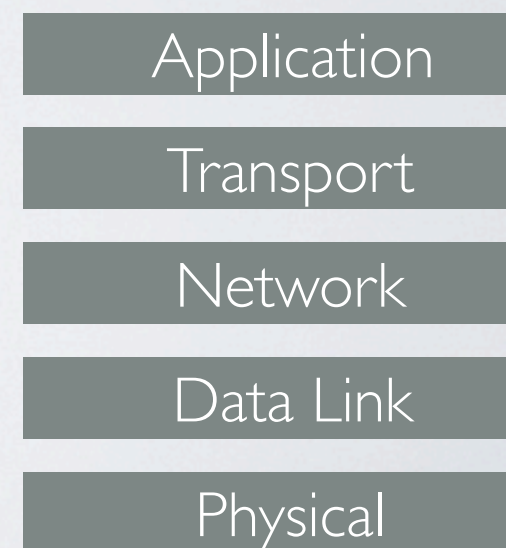
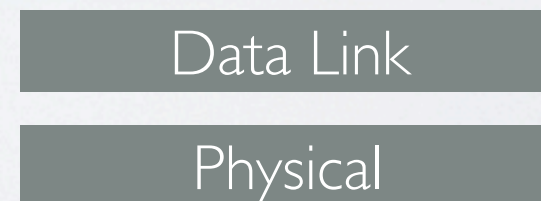
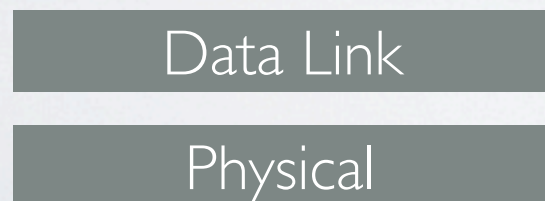
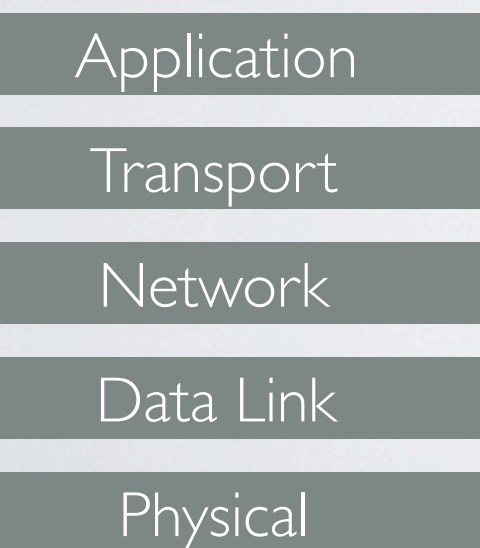
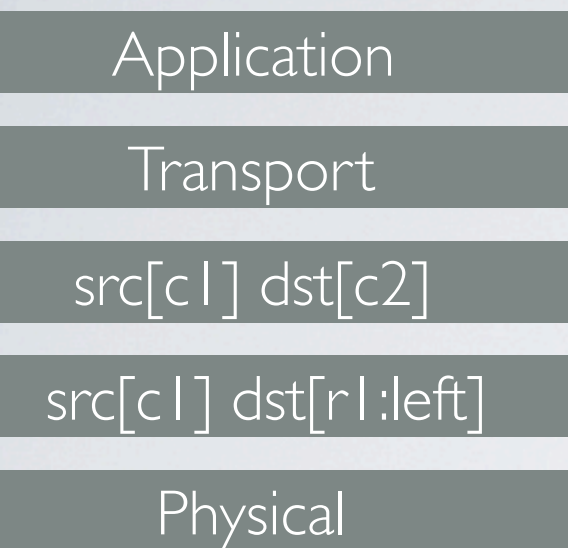


ARP :  
who's r1:left, tell c1





# HOW A PACKET TRAVELS?



**c1**

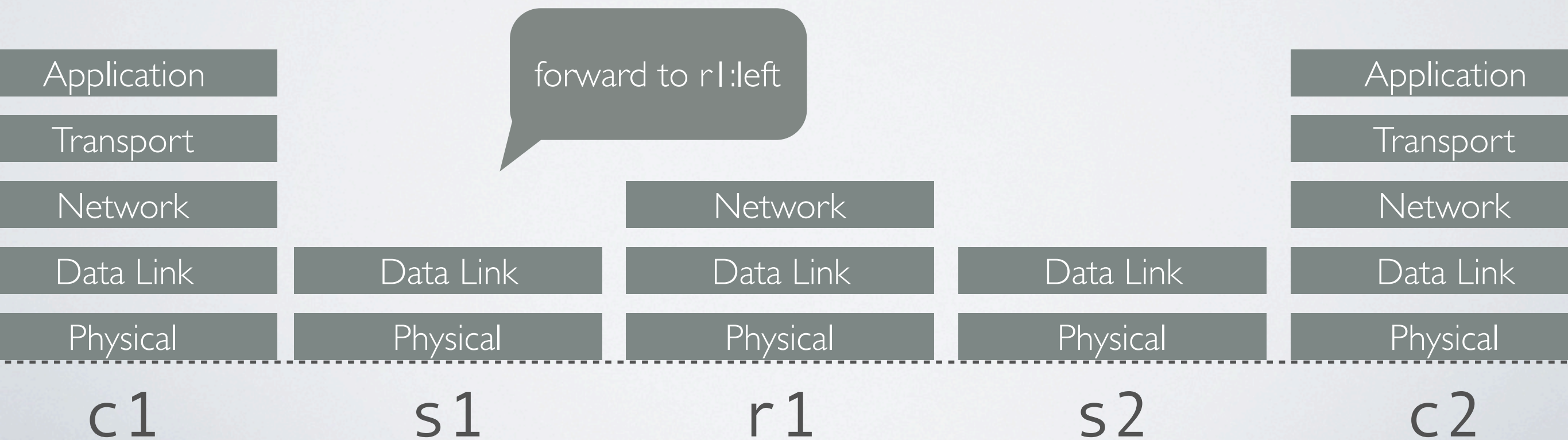
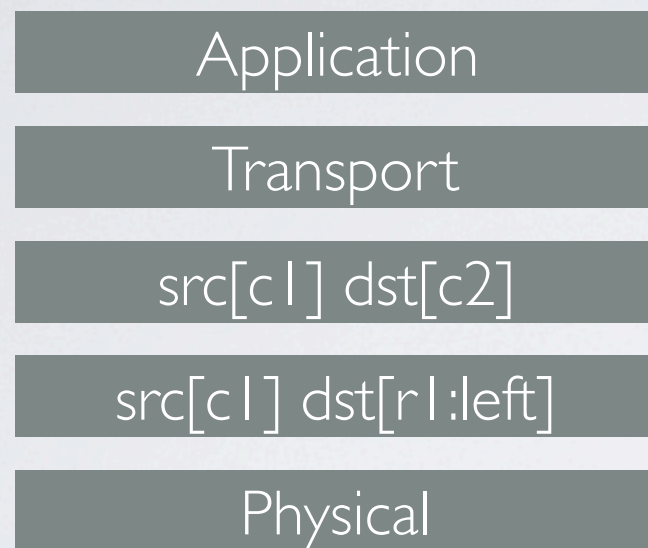
**s1**

**r1**

**s2**

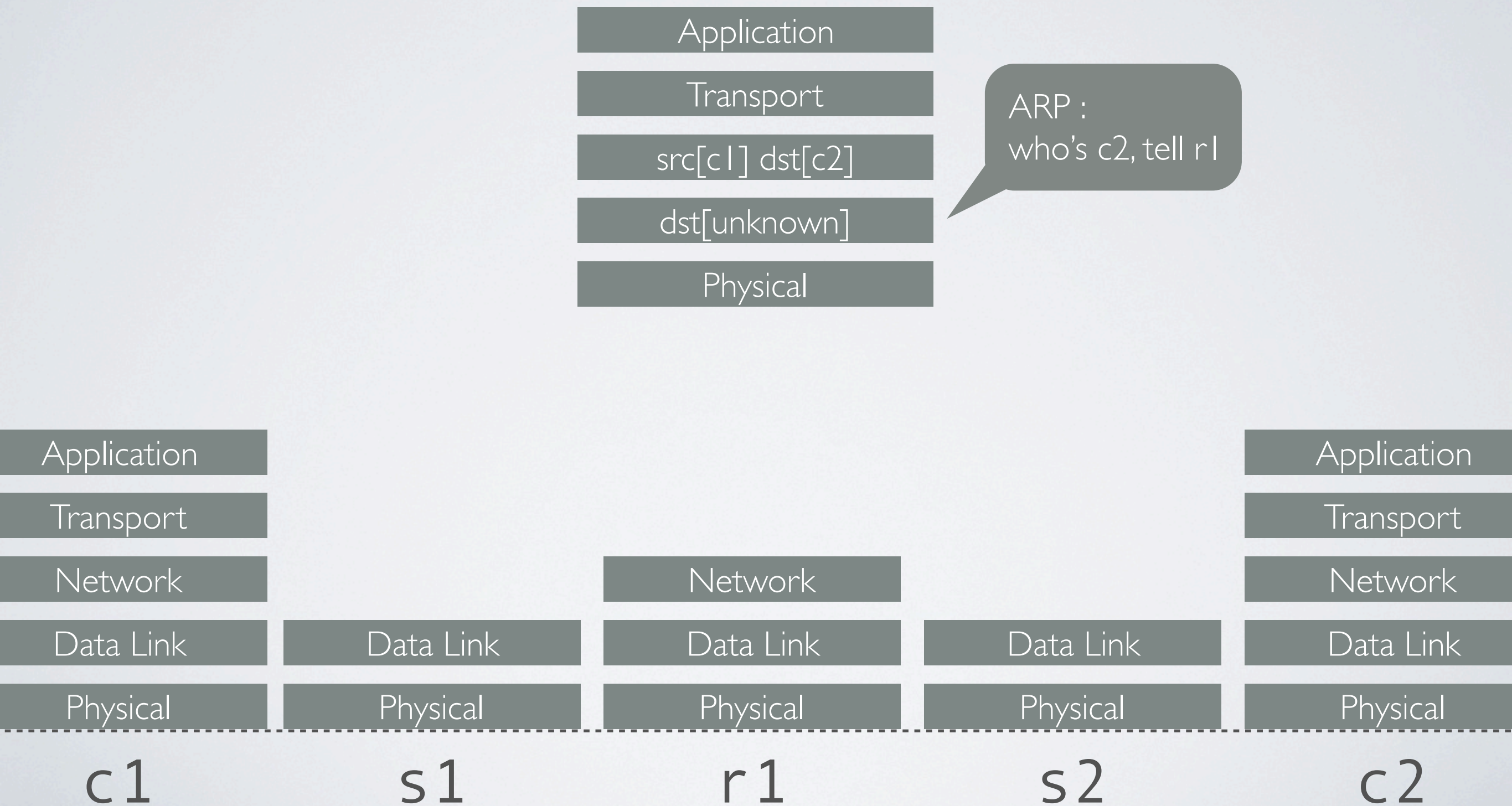
**c2**

# HOW A PACKET TRAVELS?

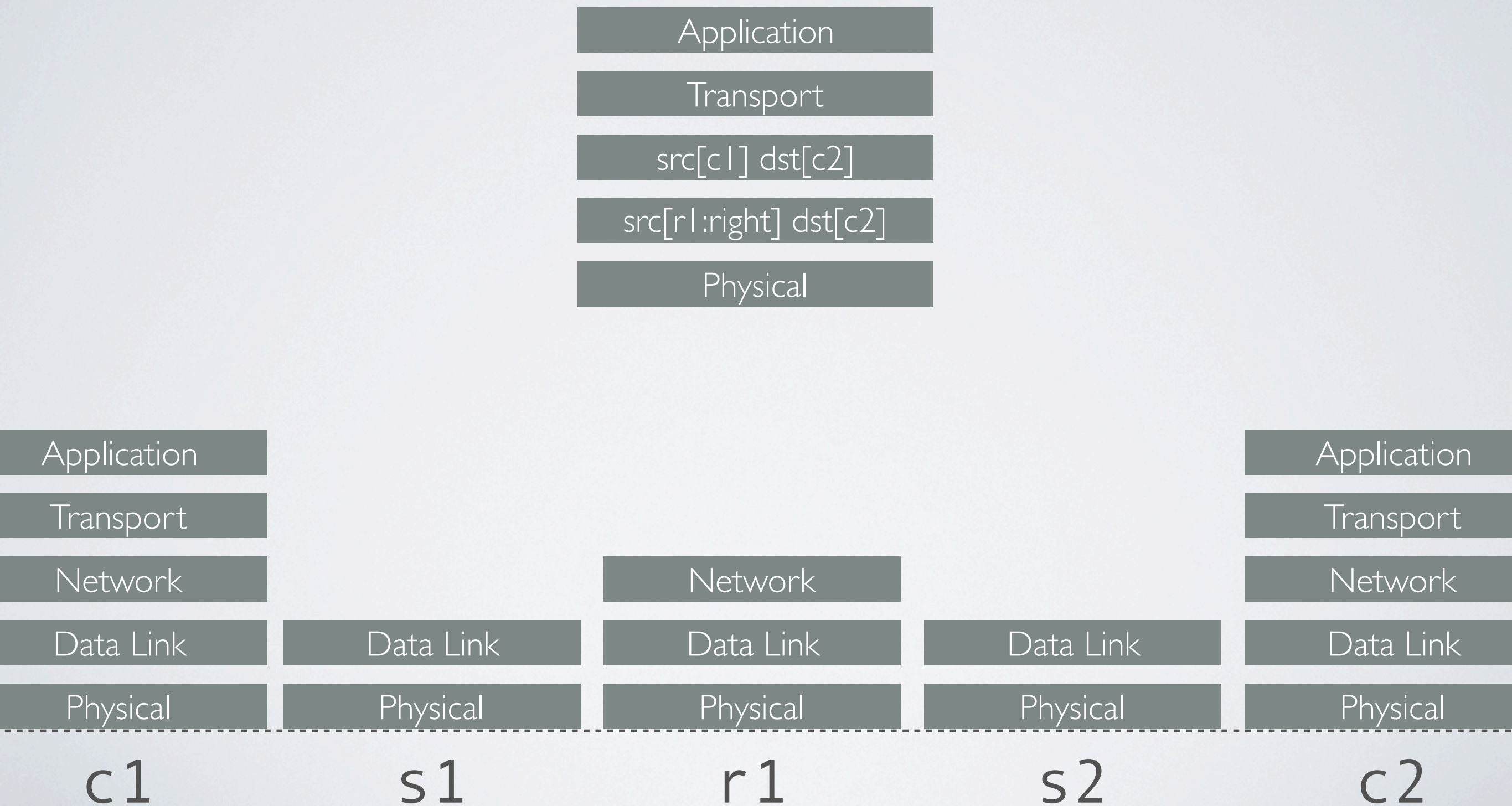




# HOW A PACKET TRAVELS?

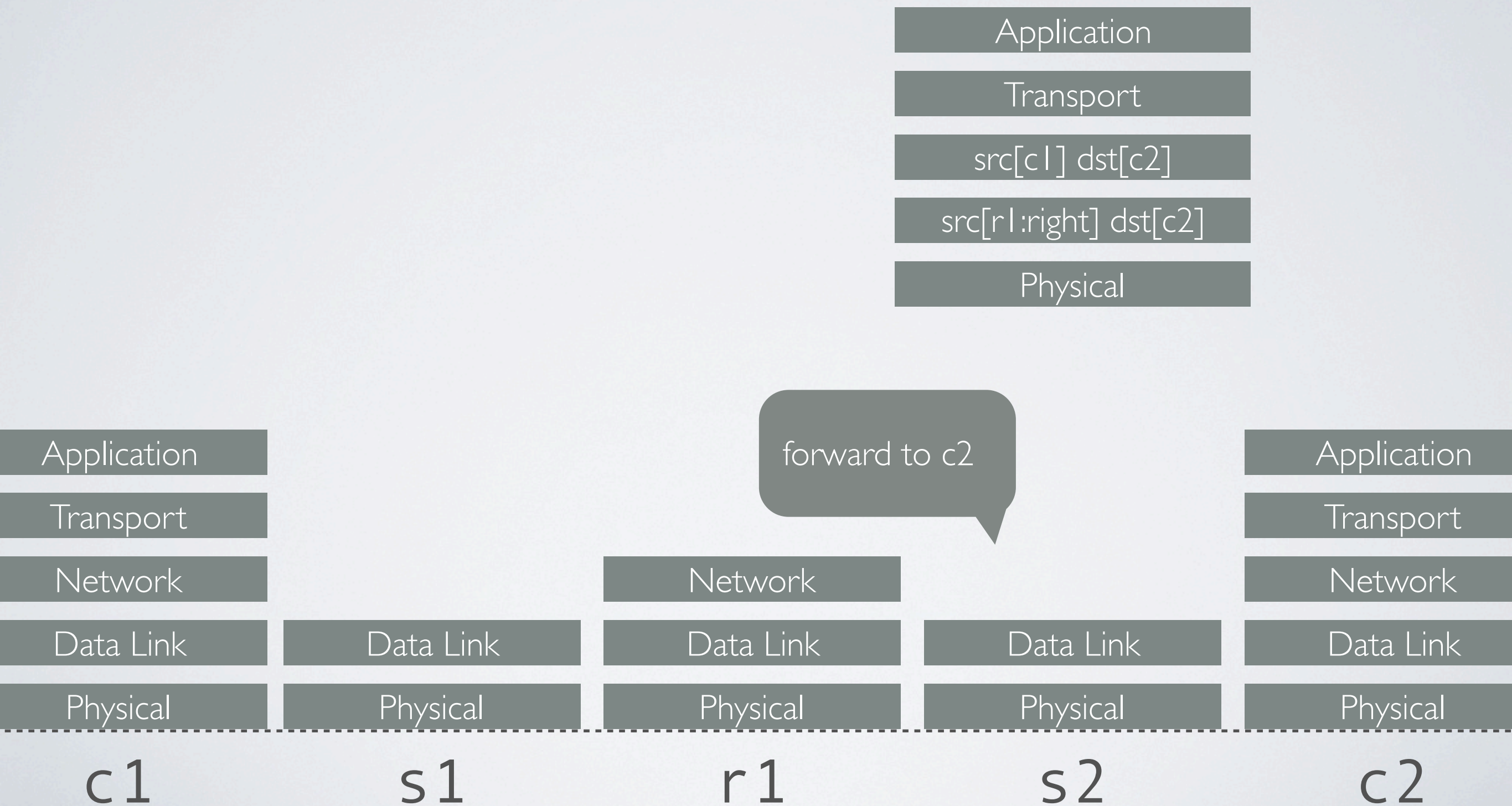


# HOW A PACKET TRAVELS?





# HOW A PACKET TRAVELS?



# HOW A PACKET TRAVELS?

finally there!

Application

Transport

src[c1] dst[c2]

src[r1:right] dst[c2]

Physical

Application

Transport

Network

Data Link

Physical

Data Link

Physical

Network

Data Link

Physical

Data Link

Physical

Application

Transport

Network

Data Link

Physical

c1

s1

r1

s2

c2



# NETWORK IN OPENFLOW'S EYE

Application

Transport

Network

Data Link

Physical

standard TCP/IP

powerful, yet less clean

Application

Ethernet

Network

Transport

Physical

OpenFlow

# DEFINE “ROUTE”

- Find a way from host A to host B
- More precisely, from network A to network B
- So,
  - there should be different subnets
  - hosts should be configured with gateways
  - network is well planned





design  
& reason

# A LEARNING ROUTER?

- sketch:

```
for every switch:  
    learn IP <--> mac of host  
    learn IP <--> switch:port(position)
```

```
when a packet comes in:  
    if we learned the IP:  
        forward to the switch:port  
        install flow entry  
    else:  
        flood
```

- plug-and-play, no configuration needed



# A LEARNING ROUTER?

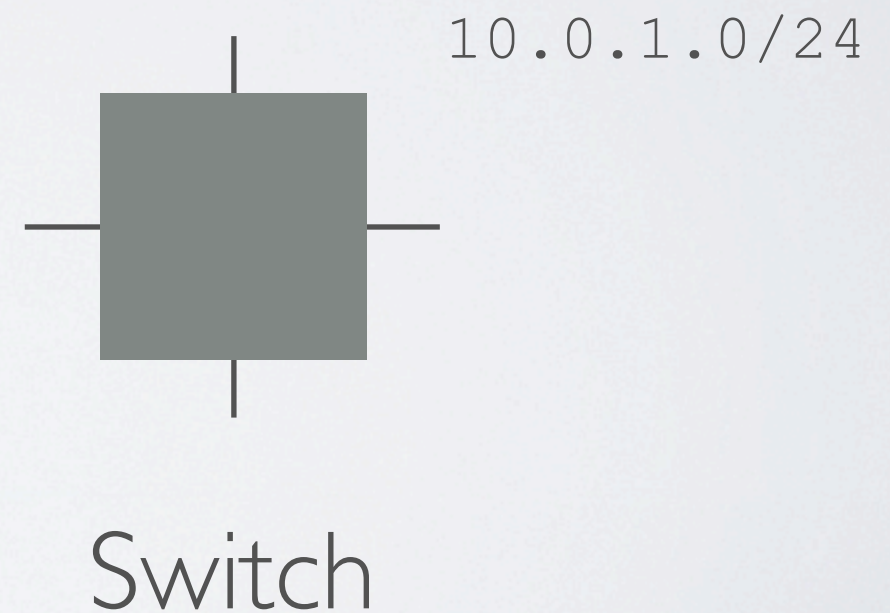
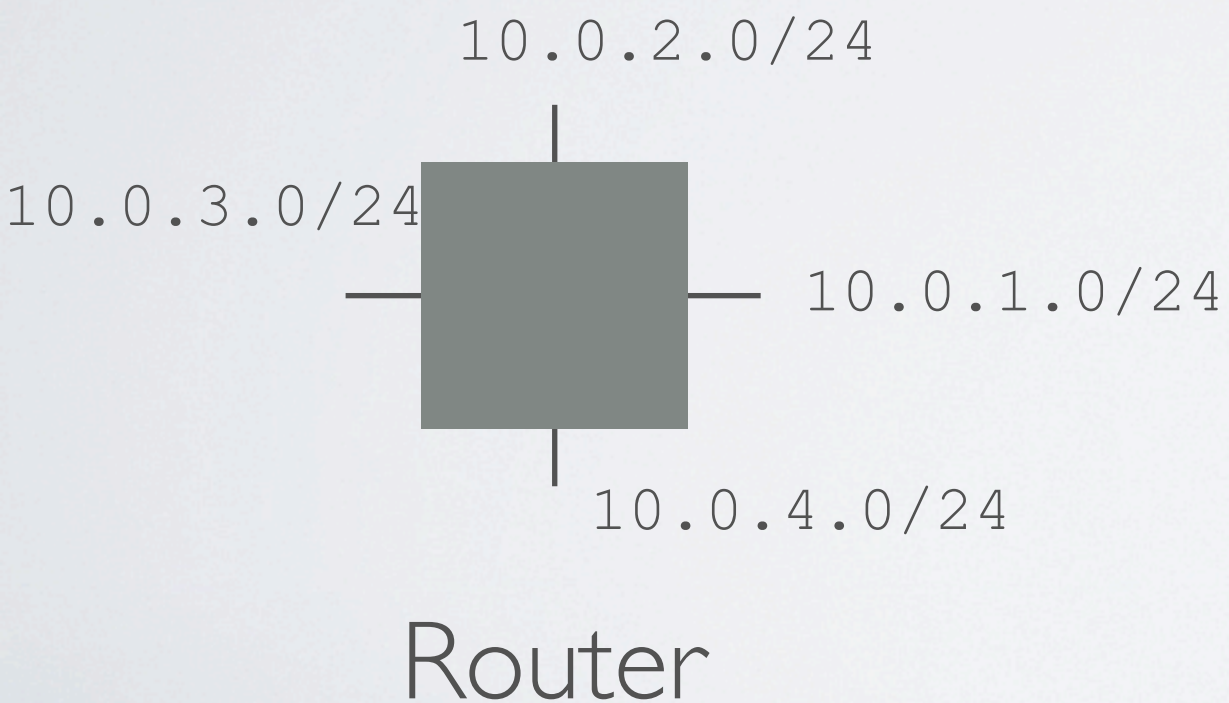
- But there's a problem:
  - We cannot handle loops
- That's why switches have *Spanning Tree Protocol(STP)*
- But STP doesn't fit for routing for lack of efficiency





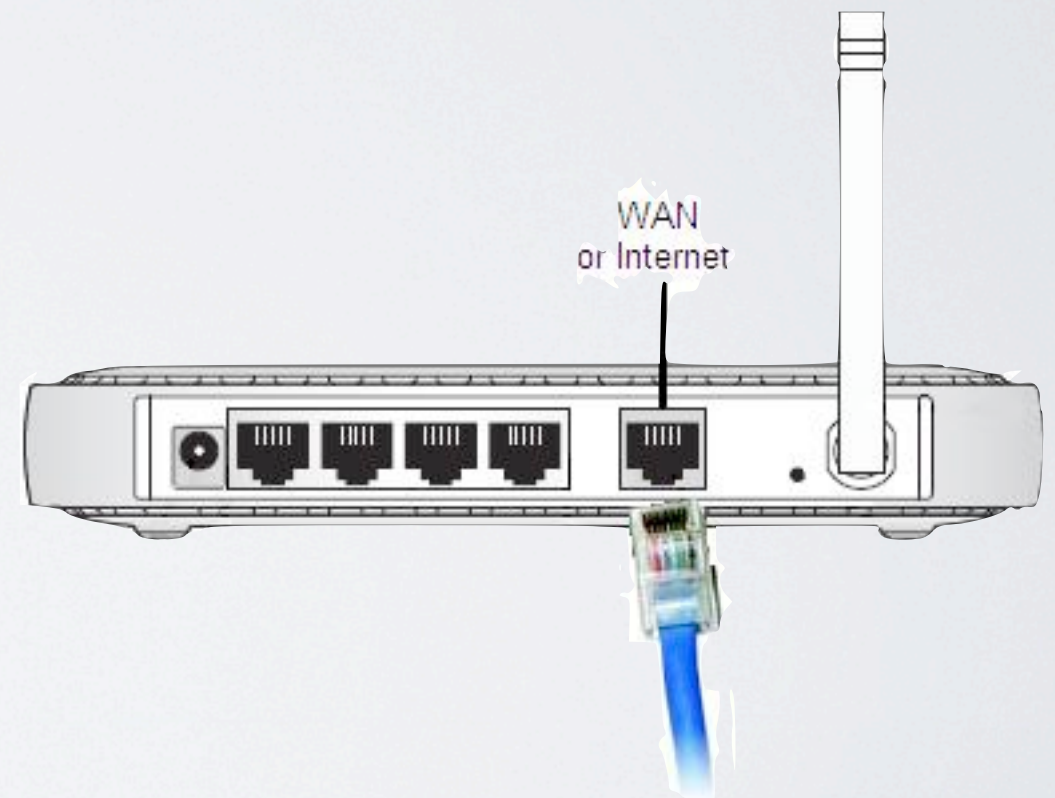
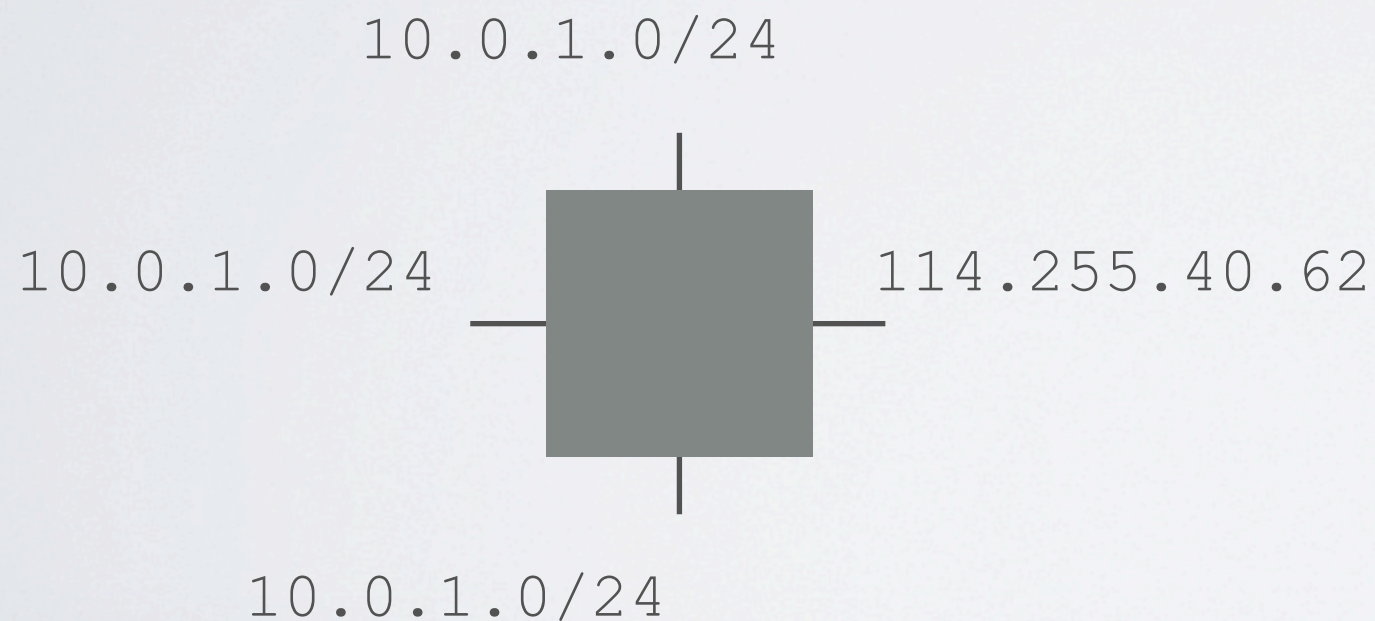
# SWITCH OR ROUTER?

- an OpenFlow switch could act like an L2 switch, or an L3 router



# SWITCH OR ROUTER?

- what if some ports are connected to same subnet?
- need *Network Address Translation(NAT)*





# DESIGN

- port-based mode configuration:
  - if the port is configured with an IP address, then the port is in ROUTER\_MODE
  - if the port has no IP address, then the port is in SWITCH\_MODE
- then the OpenFlow switch has 3 modes:
  - ROUTER\_MODE if all ports are in ROUTER\_MODE
  - SWITCH\_MODE if all ports are in SWITCH\_MODE
  - else HYBRID\_MODE



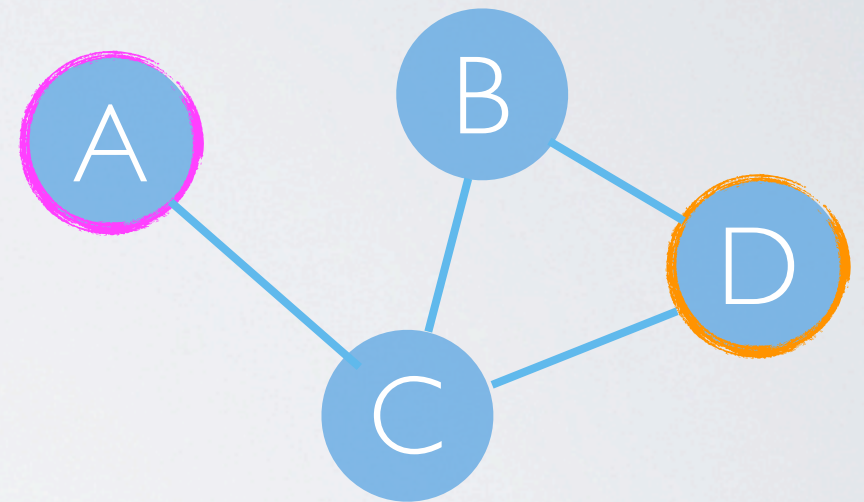
# DESIGN

- switch could be configured through configuration file, or a telnet command line interface
- then if nothing is configured, all ports are in SWITCH\_MODE and all OpenFlow switches act like normal L2 learning switches
- NAT is not supported, so the SWITCH\_MODE ports in a hybrid switch have no access to the outer world



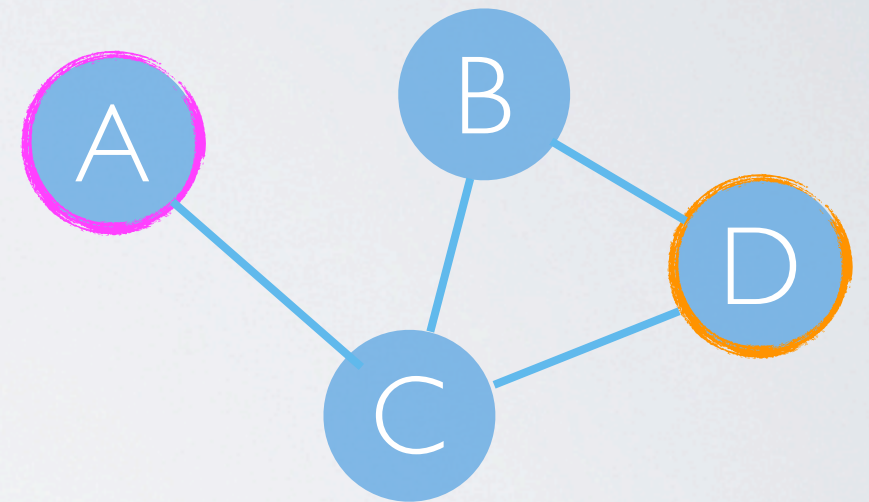
# HOW & WHEN TO INSTALL FLOW ENTRY?

- when a packet\_in, calculate the path to destination
- for every switch in the path, install flow entry in *reverse* order



# HOW & WHEN TO INSTALL FLOW ENTRY?

- but the flow entries also expire in reverse order
- this trick only puts off packet\_in messages
- so no real benefits





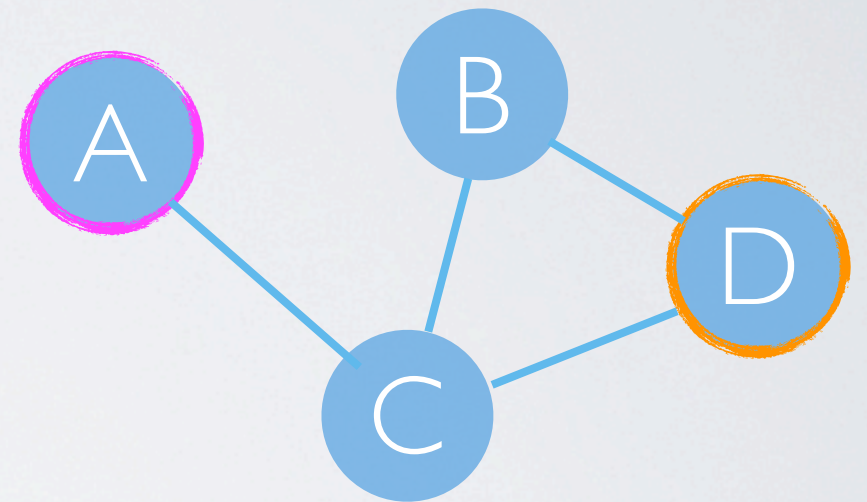
# DESIGN

- maintain a routing information data structure:

- if routing information is up-to-date, ROUTE\_CLEAN
- else ROUTE\_DIRTY

- when packet\_in:

- if ROUTE\_CLEAN, install the flow entry on this switch, for next hop
- else, re-calculate routing information, and install flow entry



# ACTING LIKE A GATEWAY

- response to ARP request
- response to ICMP Echo request(ping)
- using ARP to get MAC address of a host
- it turns out to implement a *protocol suite*, not just “routing”



# “改不动一个字”

- as you can see I tried many different ways to improve/simplify/blahblah the routing mechanism, but at last I just copied the current implementation
- that's why we call it “classic”

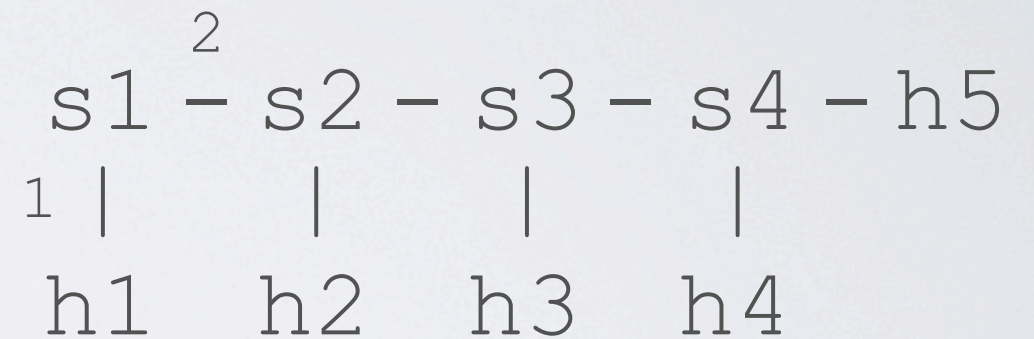
3

demo



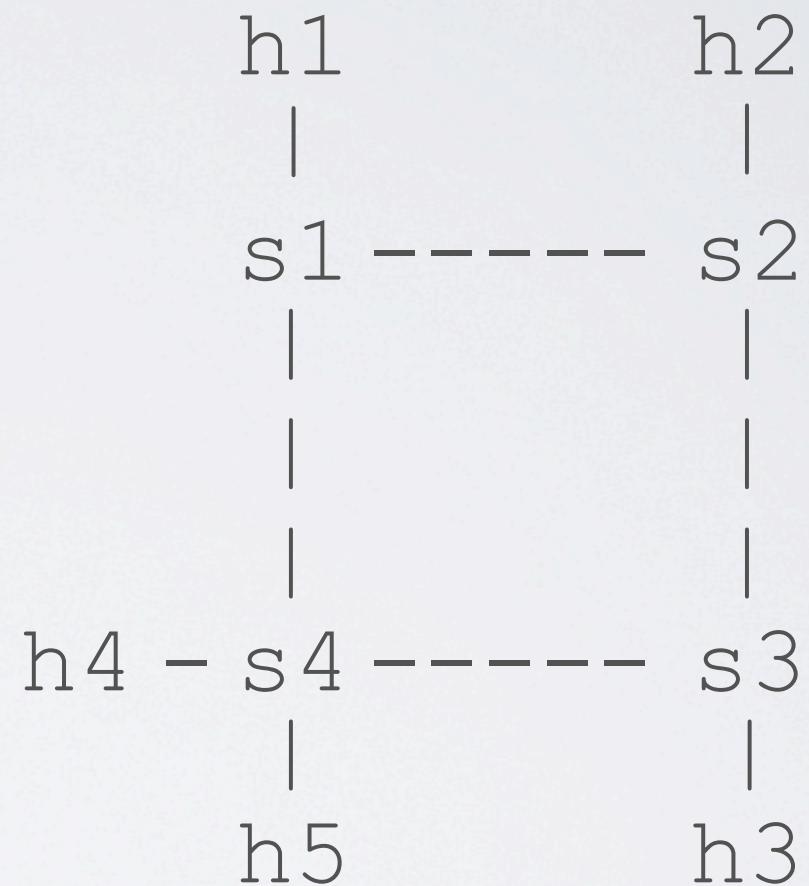
# A LINEAR TOPOLOGY

- s1: router
  - 1: 10.0.1.254, 10.0.1.0/24
  - 2: 10.0.0.254, 10.0.0.0/24
- h1: 10.0.1.1
- s2-s4: switch
- h2-h5: 10.0.0.2-10.0.0.5



# A MUCH COMPLICATED TOPOLOGY

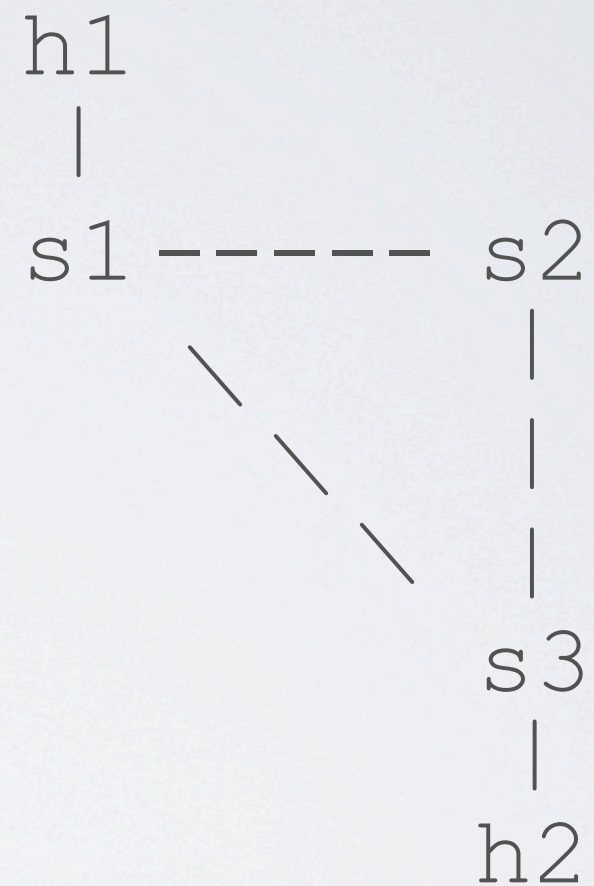
- h1: 10.0.1.1
- h2: 10.0.2.2
- h3: 10.0.3.3
- h4: 10.0.4.4
- h5: 10.0.4.5





# MANUALLY SET SWITCHES

- h1: 10.0.1.1
- h2: 10.0.2.2



4

open questions  
& future plan



# IS “PARTLY DIRTY” POSSIBLE?

- in the current implementation, a topology update causes the recalculation of whole network information

```
if topology_update:
    network.info_condition = ROUTE_DIRTY
... ..
if network.info_condition == ROUTE_DIRTY:
    network.calculate()
```

# IPv6

- as I have stated, it's actually implementing a protocol suite
- including NDP, ICMPv6, and of course IPv6

“ someone who did. I think the major parts of the problem are: 1) Adding the OpenFlow extension(s), 2) Proposing how to handle IPv6 addresses in `pox.lib.addresses` and implementing that, 3) Adding an `ipv6` class to `pox.lib.packet`. ”



# CODE

- [https://github.com/cannium/openflow\\_routing](https://github.com/cannium/openflow_routing)
- running with *POX*, test environment is *mininet*
- definitely need more tests !! so play with it !!

Q & A

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