

# FOPL - Homework 1

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## Problem 4

(a) prob4.hs

```
data Tree a = Leaf a | Node (Tree a) (Tree a) deriving Show
maptree f (Leaf x) = Leaf (f x)
maptree f (Node x y) = Node (maptree f x) (maptree f y)
```

(b) Using pattern match to define function. When *maptree* function's second argument matches "Leaf x", *Leaf (f x)* is evaluated; When function's second argument matches "Node x y", *Node (maptree f x) (maptree f y)* is evaluated, which calls *maptree* recursively.

(c) Function type: *maptree* ::  $(t \rightarrow a) \rightarrow Tree\ t \rightarrow Tree\ a$ . Because if *x* is of type *t*, there is no guarantee that *f(x)* is also of type *t*, so *f(x)*'s type is represented by *a*.

## Problem 5

(a) prob5.hs

```
data Tree a = Leaf a | Node (Tree a) (Tree a)
reduce :: (a -> a -> a) -> Tree a -> a
reduce oper (Leaf x) = x
reduce oper (Node x y) = oper (reduce oper x) (reduce oper y)
```

(b) Using pattern match method to define function. When function's second argument matches "Leaf x", result is *x*; Otherwise, *oper (reduce oper x) (reduce oper y)* is evaluated, which calls *reduce* recursively.

## Problem 6

(a) prob6.hs

```
curry :: ((a, b) -> c) -> (a -> b -> c)
```

```

curry f a b = f (a, b)
uncurry :: (a -> (b -> c)) -> ((a, b) -> c)
uncurry f (a, b) = f a b

```

(b) let  $\mathbf{f}' = \mathbf{curry\ f}$ , so  $\mathbf{f}'$  is of type  $(a \rightarrow (b \rightarrow c))$ . Then  $\mathbf{f}'$  is taken as the first argument of function *uncurry*, so  $\mathbf{uncurry\ f'}$  is of type  $((a, b) \rightarrow c)$ , which is the same type of  $\mathbf{f}$ .

let  $\mathbf{g}' = \mathbf{uncurry\ g}$ , so  $\mathbf{g}'$  is of type  $((a, b) \rightarrow c)$ . Then  $\mathbf{g}'$  is taken as the first argument of function *curry*, so  $\mathbf{curry\ g'}$  is of type  $(a \rightarrow (b \rightarrow c))$ , which is the same type of  $\mathbf{g}$ .

## Problem 7

(a) prob7.hs

```

evens :: [Int]
evens = [x * 2 | x <- [1..]]
odds  :: [Int]
odds  = [ x * 2 - 1 | x <- [1..]]

```

(b) prob7.hs

```

mergeAux [] [] l = l
mergeAux x [] l = l ++ x
mergeAux [] x l = l ++ x
mergeAux x1@(h1:t1) x2@(h2:t2) h3 =
    if h1 < h2
    then mergeAux t1 x2 (h3 ++ [h1])
    else mergeAux x1 t2 (h3 ++ [h2])

merge :: [Int] -> [Int] -> [Int]
merge l1 l2 = mergeAux l1 l2 []

```

(c) No, because *evens* and *odds* are both infinite lists. *merge* function will call itself infinitely.