Homework 2

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problem 1

```
(\lambda p.\lambda q.\lambda r.p \ q \ r)(\lambda p.\lambda q.p \ q \ r) \rightarrow (\lambda s.\lambda t.\lambda u.s \ t \ u)(\lambda p.\lambda q.p \ q \ r) \\ \rightarrow \lambda t.\lambda u.(\lambda p.\lambda q.p \ q \ r)t \ u \rightarrow \lambda t.\lambda u.(\lambda q.t \ q \ r)u \rightarrow \lambda t.\lambda u.t \ u \ r \\ \rightarrow r
```

If we don't rename bound variables, free variable r be bound. So rename bound variables so that all bound variables are different from each other and different from all of the free variables.

problem 5

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\begin{array}{c} ({\bf a}){<}\; x+y,\sigma> \to {<}\; 2+y,\sigma> \to {<}\; 2+3,\sigma> \to {<}\; 5,\sigma> \\ ({\bf b}){<}\; x=x+3,\sigma> \to {<}\; x=1+3,\sigma> \to {<}\; x=4,\sigma> \to {<}\; 4, Put(\sigma,x,4)> \\ ({\bf c}){<}\; (x=3)+x,\sigma> \to {<}\; 3+x, Put(\sigma,x,3)> \to {<}\; 3+3,\sigma'> \to {<}\; 6,\sigma'> \\ ({\bf d}){<}\; x=(x=x+3)+(x=x+5),\sigma> \to {<}\; x=(x=1+3)+(x=x+5),\sigma> \\ \to {<}\; x=(x=4)+(x=x+5)\sigma> \to {<}\; x=4+(x=x+5), Put(\sigma,x,4)> \to {<}\; x=4+(x=4+5),\sigma'> \to {<}\; x=4+(x=9),\sigma'> \to {<}\; x=4+9, Put(\sigma',x,9)> \to {<}\; x=13,\sigma''> \to {<}\; 13, Put(\sigma'',x,13)> \end{array}
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problem 6

- (a) 2
- (b) 4
- (c) 3

problem 7

- (a) $\omega = 8$
- (b)

Activation Records				
(1)	access link	(0)		
	f			
	h			
(2)	access link	()		
	x			
(3) f(1)	access link	()		α
	X			Closure
(4)	access link	()		<(), ·>
	h			<(), ·
$\overline{(5)}$	access link	()	•	
	w			
(6) h(3)	access link	()	•	
	z			
(7) f(3)	access link	()	•	
	У			
			•	

$\begin{array}{c} Closures & Compiled \\ <(\),\ > \\ <(\),\ > & |codeforf| \\ |codeforh| \end{array}$

problem 8

(a)					
Activation Records					
(1)	access link	(0)			
	f				
	h				
(2) h(2)	access link	()			
	x				
(3) f(1)	access link	()			
	X				

 $\begin{array}{ll} \textit{Closures} & \textit{Compiled Codes} \\ A<(\), \ > & |\textit{function}, line 1| \\ B<(\), \ > & |\textit{function}, line 3| \end{array}$

(b) B

(c) 第二行之后调用 h(2) 会调用 $f_1(2)$ 然后再调用 $f_1(1)$; 而第四行调用 h(2) 时, f 已经被修改并指向 $f_3()$, 所以会先调用 $f_1(2)$ 再调用 $f_3(1)$.

l)

$Activation \ Records$				
(1)	access link	(0)		
	f			
	h			
(2) h(2)	access link	()		
	x			
(3) g(1)	access link	()		
	X			
/ \ - / \ - \ \ \ \ \ \ \ \ \ \ \ \ \ \				

f(e)h(2) = 2. 先调用 g(2) 再调用 g(1).