习题课

第五次 第二次作业 陈航



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DADDIU R4, R1, #800

foo: L.D F2, 0(R1)

MUL.D F4, F2, F0

L.D F6, O(R2)

ADD.D F6, F4, F6

S.D F6, 0(R2)

DADDIU R1, R1, #8 2 stalls

DADDIU R2, R2, #8

DSLTU R3, R1, R4 2 stalls

BNEZ R3, foo

1 stall

2 stalls

6 stalls

2 stalls

4 stalls



Clock cycle	Unschedu	ıled code	R4,R1,#800 F2,0(R1) F6,0(R2) F4,F2,F0 R1,R1,#8 R2,R2,#8 R3,R1,R4 F6,F4,F6	
1	DADDIU	R4,R1,#800	DADDIU	R4,R1,#800
2	L.D	F2,0(R1)	L.D	F2,0(R1)
3	stall	2 stalls	L.D	F6,0(R2)
4	MUL.D	F4,F2,F0	MUL.D	F4,F2,F0
5	L.D	F6,0(R2)	DADDIU	R1,R1,#8
6	stall	6 stalls	DADDIU	R2,R2,#8
	stall	Ostalis	DSLTU	R3,R1,R4
	stall	>2 stalls	stall	
1	stall		stall	
7	ADD.D	F6,F4,F6	ADD.D	F6,F4,F6
8	stall		stall	
9	stall	4 stalls	stall	
10	stall		BNEZ	R3,foo
11	S.D	F6,0(R2)	S.D	F6,-8(R2)
12	DADDIU	R1,R1,#8		
13	DADDIU	R2, R2, #8		
14	DSLTU	R3,R1,R4 2 stall	.8	
15	stall	2 stalls		
16	BNEZ	R3, foo		
17	stall	↓ 1 stall		

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Iteration	Instruction	Issues at	Executes/ Memory	Write CDB at	Comment
1	L.D F2,0(R1)	1	2	3	First issue
1	MUL.D F4,F2,F0	2	4	19	Wait for F2
_					Mult rs [3–4] Mult use [5–18]
1	L.D F6,0(R2)	3	4	5	Ldbuf [4]
1	ADD.D F6,F4,F6	4	20	30	Wait for F4 Add rs [5–20] Add use [21–29]
1	S.D F6,0(R2)	5	31		Wait for F6 Stbuf1 [6-31]
1	DADDIU R1,R1,#8	6	7	8	
1	DADDIU R2,R2,#8	7	8	9	
1	DSLTU R3,R1,R4	8	9	10	
1	BNEZ R3,foo	9	11		Wait for R3
2	L.D F2,0(R1)	10	12	13	Wait for BNEZ
2	MUL.D F4,F2,F0	11	14 19	34	Ldbuf [11–12] Wait for F2 Mult busy Mult rs [12–19]
2	L.D F6,0(R2)	12	13	14	Mult use [20–33] Ldbuf [13]
2	ADD.D F6,F4,F6	13	35	45	Wait for F4 Add rs [14–35] Add use [36–44]
2	S.D F6,0(R2)	14	46		Wait for F6 Stbuf [15-46]
2	DADDIU R1,R1,#8	15	16	17	
2	DADDIU R2,R2,#8	16	17	18	
2	DSLTU R3,R1,R4	17	18	20	
2	BNEZ R3, foo	18	20		Wait for R3
3	L.D F2,0(R1)	19	21	22	Wait for BNEZ Ldbuf [20–21]
3	MUL.D F4,F2,F0	20	23 34		Wait for F2 Mult busy
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					Mult rs [21–34] Mult use [35–48]	
3	L.D F6,0(R2)	21	22	23	Ldbuf [22]	
3	ADD.D F6,F4,F6	22	50	60	Wait for F4 Add rs [23–49] Add use [51–59]	
3	S.D F6,0(R2)	23	55		Wait for F6 Stbuf [24–55]	
3	DADDIU R1,R1,#8	24	25	26		
3	DADDIU R2,R2,#8	25	26	27		
3	DSLTU R3,R1,R4	26	27	28		
					Wait for R3	



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- a. 对于非条件分支指令,其下一条指令直接存储在BTB中,因此省去了通过地址取指令的过程,因此节省了1 cycle
- b. 考虑如下指令 jmp L1

• • •

L: mov r1, r0

在对应5% unconditional branch, 90% hit rate, 2 cycyles penalty for a buffer miss中, 其性能提升为:



$$5\%*(90\%*(-1) + 10\%*2) = -0.035$$

为了使得BTB帮助机器获得性能提升,即: buffer 中存储的指令所带来的收益要比因buffer miss所带来的penalty要大,设hit rate = x,那么:

$$x^*(-1) + (1-x)^*2 < 0$$

解得: x > 2/3 = 66.7%

