

习题课 3

郭海林

2014.5.27

4.9

```
for (i=0;i<300;i++) {  
    c_re[i] = a_re[i] * b_re[i] - a_im[i] * b_im[i];  
    c_im[i] = a_re[i] * b_im[i] + a_im[i] * b_re[i];  
}
```

a. 运算密度 = 浮点操作个数 / 内存访问字节数

六次浮点数操作

4 次读操作: a_re, a_im, b_re, b_im + 2 次写操作: c_re, c_im

运算密度 = $6/(6*4) = 1/4$

4.9 cont.

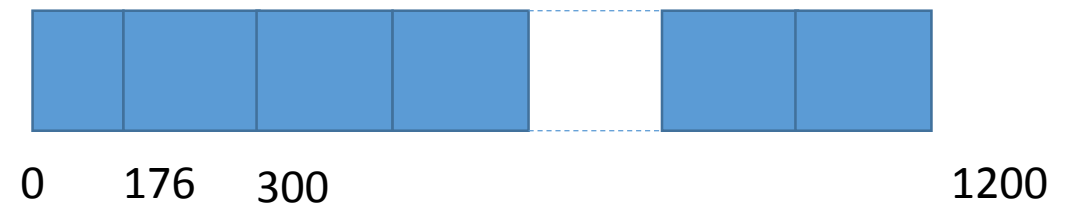
- b.

```
li    $VL,44          # perform the first 44 ops
li    $r1,0            # initialize index

loop:
    lv    $v1,a_re+$r1  # load a_re
    lv    $v3,b_re+$r1  # load b_re
    mulvv.s $v5,$v1,$v3  # a+re*b_re
    lv    $v2,a_im+$r1  # load a_im
    lv    $v4,b_im+$r1  # load b_im
    mulvv.s $v6,$v2,$v4  # a+im*b_im
    subvv.s $v5,$v5,$v6  # a+re*b_re - a+im*b_im
    sv    $v5,c_re+$r1  # store c_re
    mulvv.s $v5,$v1,$v4  # a+re*b_im
    mulvv.s $v6,$v2,$v3  # a+im*b_re
    addvv.s $v5,$v5,$v6  # a+re*b_im + a+im*b_re
    sv    $v5,c_im+$r1  # store c_im
```

```
        bne    $r1,0,else  # check if first iteration
        addi    $r1,$r1,#176 # first iteration,
                                increment by 176
j loop      # guaranteed next iteration

else:
        addi    $r1,$r1,#256 # not first iteration,
                                increment by 256
skip:
        blt     $r1,1200,loop # next iteration?
```



4.9 cont.

- c.

1. mulvv.s lv # a_re * b_re (assume already loaded), load a_im
2. lv mulvv.s # load b_im, a_im*b_im
3. subvv.s sv # subtract and store c_re
4. mulvv.s lv # a_re*b_im, load next a_re vector
5. mulvv.s lv # a_im*b_re, load next b_re vector
6. addvv.s sv # add and store c_im

6 次钟鸣

d. 每次迭代所用的时钟数: $6 \times 64 + 15 \times 6 + 8 \times 4 + 5 \times 2 = 516$

产生64 个结果: $516 / 64 = 8$

4.9 cont.

- e. 三条内存流水线

1. mulvv.s	# a_re*b_re
2. mulvv.s	# a_im*b_im
3. subvv.s sv	# subtract and store c_re
4. mulvv.s	# a_re*b_im
5. mulvv.s lv lv	# a_im*b_re, load next a_re, load next b_re
6. addvv.s sv lv lv	# add, store c_im, a_im,b_im

仍然是 6 次钟鸣，所需时钟数不受影响。

4.10

- Vector :
 - execution time: 400 ms
 - mem access: $(200\text{MB}+100\text{MB})/(30\text{GB/s}) = 10\text{ms}$
 - total time: 410ms
- Hybrid:
 - execution time: 400 ms
 - mem access: $(200\text{MB}+100\text{MB})/(150\text{GB/s}) = 2\text{ms}$
 - host IO: $(200\text{MB}+100\text{MB})/(10\text{GB/s}) = 30\text{ms}$
 - latency: 10 ms
 - total time: 442ms

4.12

```
for (int x=0; x<NX-1; x++) {  
  for (int y=0; y<NY-1; y++) {  
    for (int z=0; z<NZ-1; z++) {  
      int index = x*NY*NZ + y*NZ + z;  
      if (y>0 && x >0) {  
        material = IDx[index];  
        dH1 = (Hz[index] - Hz[index-incrementY])/dy[y];  
        dH2 = (Hy[index] - Hy[index-incrementZ])/dz[z];  
        Ex[index] = Ca[material]*Ex[index]+Cb[material]*(dH2-dH1);  
      }  
    }  
  }  
}
```

- a. 浮点操作个数：共 8 FLOPS
内存访问字节数：9 次reads + 1 次write
运算密度：8 / 40
- b. 可以
- c. $30 \text{ GB/s} * 8/40 \text{ FLOPS/B} = 6 \text{ GFLOPS/s}$
如果峰值性能大于 6 GFLPOS/s，则访存受限；否则，计算受限。
- d. 单个浮点数运算密度： $85/4 = 21.25 \text{ GFLOPS/s}$

4.13

- a. $1.5 \text{ GHz} \times .80 \times .85 \times 0.70 \times 10 \text{ cores} \times 32/4 = 57.12 \text{ GFLOPs/s}$
- b.
 - (1) Speedup = $16/8 = 2$
 - (2) Speedup = $15/10 = 1.5$
 - (3) Speedup = $0.95/0.85 = 1.11$

4.14

- a.

```
for (i=0; i<100; i++) {  
    A[i] = B[2*i+4];  
    B[4*i+5] = A[i];  
}
```

GCD 测试:

$\gcd(4, 2) \mid (4-5)$?

所以, 不存在相关。

4.14 cont.

• b.

```
for (i=0;i<100;i++) {  
    A[i] = A[i] * B[i]; /* S1 */  
    B[i] = A[i] + c;    /* S2 */  
    A[i] = C[i] * c;    /* S3 */  
    C[i] = D[i] * A[i]; /* S4 */  
}
```

重命名消除反相关和输出相关:

```
for (i=0;i<100;i++) {  
    A[i] = A[i] * B[i]; /* S1 */  
    B1[i] = A[i] + c;   /* S2 */  
    A1[i] = C[i] * c;   /* S3 */  
    C1[i] = D[i] * A1[i]; /* S4 */  
}
```

真相关:

S2-S1 A[i]; S4-S3 A[i];

反相关:

S1-S2 B[i]; S2-S3 A[i];

S1-S3 A[i]; S3-S4 C[i];

输出相关:

S1-S3 A[i]

(消除 s1-s2 关于 B[i] 的反相关)

(消除 s1-s3 关于 A[i] 的输出相关;

同时消除 s1-s3,s2-s3 关于A[i] 的反相关)

(消除 s3-s4 关于 C[i] 的反相关)

4.14 cont.

- C.

```
for (i=0; i < 100; i++) {  
    A[i] = A[i] + B[i];    /* S1 */  
    B[i+1] = C[i] + D[i]; /* S2 */  
}
```

关于 **B** 存在循环间依赖，因此不能并行。
修改为：

```
A[0] = A[0] + B[0];  
for (i=1; i<100; i++) {  
    B[i] = C[i-1] + D[i-1];  
    A[i] = A[i] + B[i];  
}  
B[100] = C[99] + D[99];
```

$$A[0] = A[0] + B[0]$$

$$B[1] = C[0] + D[0]$$

$$A[1] = A[1] + B[1]$$

$$B[2] = C[1] + D[1]$$

$$A[2] = A[2] + B[2]$$

$$B[3] = C[2] + D[2]$$

$$A[3] = A[3] + B[3]$$

$$B[4] = C[3] + D[3]$$

...