

For the recurrence equation  $T(N) = aT(N/b) + f(N)$ , if  $af(N/b) = f(N)$ , then  $T(N) = \Theta(f(N)\log_b N)$ .

☒ T ☐ F

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For the recurrence equation  $T(N) = 8T(N/2) + N^3 \log N$ , we obtain  $T(N) = O(N^3 \log N)$  according to the Master Theorem.

F

For the recurrence equation  $T(N) = aT(N/b) + f(N)$ , if  $af(N/b) = Kf(N)$  for some constant  $K > 1$ , then  $T(N) = \Theta(f(N))$ .

F

If divide-and-conquer strategy is used to find the closest pair of points in a plane, unless the points are sorted not only by their  $x$  coordinates but also by their  $y$  coordinates, it would be impossible to solve it in a time of  $O(N \log N)$ , where  $N$  is the number of points.

T

Suppose that the divide-and-conquer strategy is used to find the maximum and the minimum of  $N$  positive numbers. At each step, the problem is divided into 2 sub-problems of size  $N/2$ . Then the time recurrence is  $T(N) = 2T(N/2) + f(N)$ , where  $f(N)$  is \_\_\_\_.

- ☐ A.  $N/2$
- ☒ B.  $O(1)$
- ☐ C.  $\Omega(N)$
- ☐ D.  $\Theta(\log N)$

答案正确: 1 分 [创建提问](#)