

```

//AmicablePair sum computing method
private int factorsSum(int n) {
    long startTime = System.nanoTime();//Time count
    int j, sum = 0, num = 0;
    int[] sums = new int [n+1];
    for (int i = 1; i <=n/2; i++) {
        j = i*2;
        while (j <= n) {
            sums[j] = sums[j]+i; // add factor i to every sums in the list
            j = j+i;
        }
    }
    for (int i = 2; i <= n; i++) {
        sum = sums[i];
        if (sum > n || sum <= i)// avoid sum out of n and delete repeating such as
        "284-220" from "220-284 "
            continue;
        else {
            if (sums[sum] == i) { // Judge Amicable Pair
                System.out.println(num+": "+i+" and "+sum); //output Amicable Pair
                num++;
            }
        }
    }
    long endTime = System.nanoTime();// Time count
    double d2 = u.timeInSec(endTime,startTime) ;// Time count
    System.out.println("AmicablePair " + " CPU time = " + d2 + " seconds"); // Time count
    return num;
}

```

By regular methods, we need try every number from 2 to \sqrt{n} to compute remainder (%) to judge whether it is one of the factors of n. And the final Time Complexity would be $O(n\sqrt{n})$, that will cost so long time.

So I used this optimized algorithm as followed:

	n	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	...
			1	1	1	1	1	1	1	1	1	1	1	1	1	1	...
i=2	j=4				2		2		2		2		2		2		...
i=3	j=6						3			3			3			3	...
i=4	j=8								4				4				...
i=5	j=10										5						...
i=6	j=12												6				...

.....

Sums[]: Sums[1] sums[2].....

We created a Sums[] Array to add up the potential factors except themselves as the table above. In all ,we only need to count:

$n/2 + n/3 + n/4 + n/5 + \dots + 1/n + n = n \log n$ times. So that we could get Time

Complexity of $O(n \log n)$, which is about 300 times faster than the

$O(n \sqrt{n})$ when n equals 100 million. Finally I got the result within 20 seconds(17 s).