

HOMEWORK 5 RISC-V CPU

Due date:

Overview

The goal of this homework is to help you understand **how a RISC-V work** and how to use Verilog hardware description language (Verilog HDL) to model electronic systems. In this homework, you need to implement ALU and decoder module and make your codes be able to **execute 17 RISC-V instructions**. You need to follow the instruction table in this homework and satisfy all the homework requirements. In addition, you need to verify your CPU by using Modelsim.

General rules for deliverables

- You need to complete this homework **INDIVIDUALLY**. You can discuss the homework with other students, but you need to do the homework by yourself. You should not **copy** anything from someone else, and you should not **distribute** your homework to someone else. If you violate any of these rules, you **will get NEGATIVE scores, or even fail this course directly**
- When submitting your homework, compress all files into a single **zip** file, and upload the compressed file to Moodle.
 - Please follow the file hierarchy shown in Figure 1.
F740XXXXX (your id) (folder)
src (folder) * Store your source code
report.docx (project report. The report template is already included. Follow the template to complete the report.)

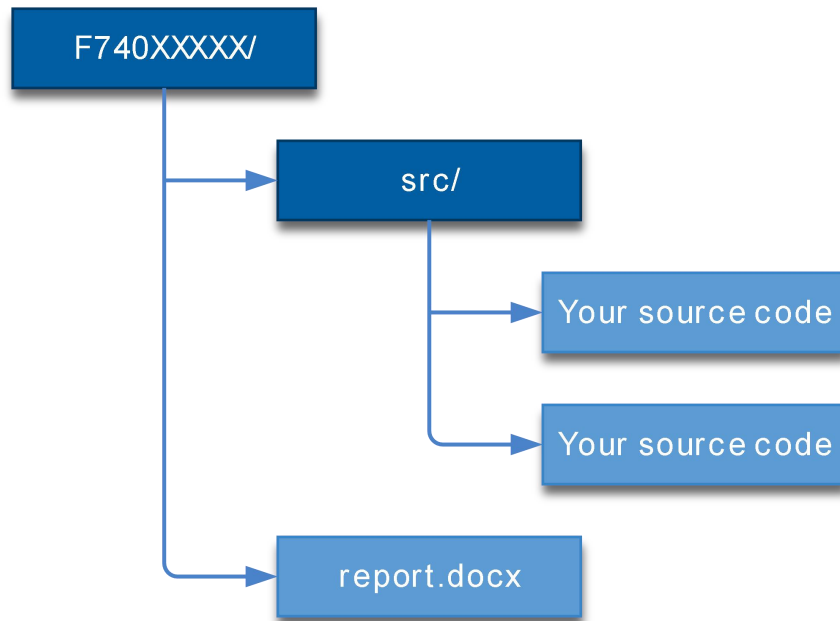


Figure 1. File hierarchy for homework submission

- **Important! DO NOT submit your homework in the last minute. Late submission is not accepted.**
- You should finish **all the requirements (shown below) in this homework** and Project report.
- **If your code can not be recompiled by TA successfully using modelsim, you will receive NO credit.**
- Verilog and SystemVerilog generators aren't allowed in this course.

Instruction format:

The highlighted instructions that have been completed in the previous job.

● R-type

31	25	24	20	19	15	14	12	11	7	6	0		
funct7		rs2		rs1		funct3		rd		opcode		Mnemonic	Description
0000000		rs2		rs1		000		rd		0110011		ADD	$rd = rs1 + rs2$
0100000		rs2		rs1		000		rd		0110011		SUB	$rd = rs1 - rs2$
0000000		rs2		rs1		100		rd		0110011		XOR	$rd = rs1 \wedge rs2$
0000000		rs2		rs1		110		rd		0110011		OR	$rd = rs1 \mid rs2$
0000000		rs2		rs1		111		rd		0110011		AND	$rd = rs1 \& rs2$

● I-type

31	20	19	15	14	12	11	7	6	0		
imm[11:0]		rs1		funct3		rd		opcode		Mnemonic	Description
imm[11:0]		rs1		010		rd		0000011		LW	$rd = M[rs1 + imm]$
imm[11:0]		rs1		000		rd		0010011		ADDI	$rd = rs1 + imm$
imm[11:0]		rs1		100		rd		0010011		XORI	$rd = rs1 \wedge imm$
imm[11:0]		rs1		110		rd		0010011		ORI	$rd = rs1 \mid imm$
imm[11:0]		rs1		111		rd		0010011		ANDI	$rd = rs1 \& imm$
imm[11:0]		rs1		000		rd		1100111		JALR	$rd = PC + 4$ $PC = imm + rs1$ (Set LSB of PC to 0)

● S-type

31	25	24	20	19	15	14	12	11	7	6	0		
imm[11:5]		rs2		rs1		funct3		imm[4:0]		opcode		Mnemonic	Description
imm[11:5]		rs2		rs1		010		imm[4:0]		0100011		SW	$M[rs1 + imm] = rs2$

● B-type

31	25	24	20	19	15	14	12	11	7	6	0		
imm[12 10:5]		rs2		rs1		funct3		imm[4:1 11]		opcode		Mnemonic	Description
imm[12 10:5]		rs2		rs1		000		imm[4:1 11]		1100011		BEQ	$PC = (rs1 == rs2) ?$ $PC + imm : PC + 4$
imm[12 10:5]		rs2		rs1		001		imm[4:1 11]		1100011		BNE	$PC = (rs1 != rs2) ?$ $PC + imm : PC + 4$

- **U-type**

31	12	11	7	6	0		
imm[31:12]		rd		opcode	Mnemonic	Description	
imm[31:12]		rd		0010111	AUIPC	rd = PC + imm	
imm[31:12]		rd		0110111	LUI	rd = imm	

- **J-type**

31	12	11	7	6	0		
imm[20 10:1 11 19:12]		rd		opcode	Mnemonic	Description	
imm[20 10:1 11 19:12]		rd		1101111	JAL	rd = PC + 4 PC = PC + imm	

Homework Description

- **Module**

- top_tb module**

- “top_tb” is not a part of CPU, it is a file that controls all the program and verify the correctness of our CPU. The main features are as follows: send periodical signal CLK to CPU, set the initial value of IM, print the value of DM, end the program.

✖You do not need to modify this module.

- top module**

“top” is the outmost module. It is responsible for connecting wires between CPU, IM and DM.

Here are the wires:

- *instr_read* represents the signal whether the instruction should be read in IM.
- *instr_addr* represents the instruction address in IM.
- *instr_out* represents the instruction send from IM .
- *data_read* represents the signal whether the data should be read in DM.
- *data_write* has four signal , and every signal represents the byte of the data whether should be wrote in DM.

Mem[0] =

{*Mem[0][31:24]*,*Mem[0][23:16]*,*Mem[0][15:8]*,*Mem[0][7:0]*}

data_write[3] => control *Mem[0][31:24]*

data_write[2] => control *Mem[0][23:16]*

data_write[1] => control *Mem[0][15:8]*

data_write[0] => control *Mem[0][7:0]*

- *data_addr* represents the data address in DM.
- *data_in* represents the data which will be wrote into DM .
- *data_out* represents the data send from DM .

※You do not need to modify this module.

d. **SRAM module**

“SRAM” is the abbreviation of “Instruction Memory” (or “Data Memory”). This module saves all the instructions (or data) and send instruction (or data) to CPU according to request.

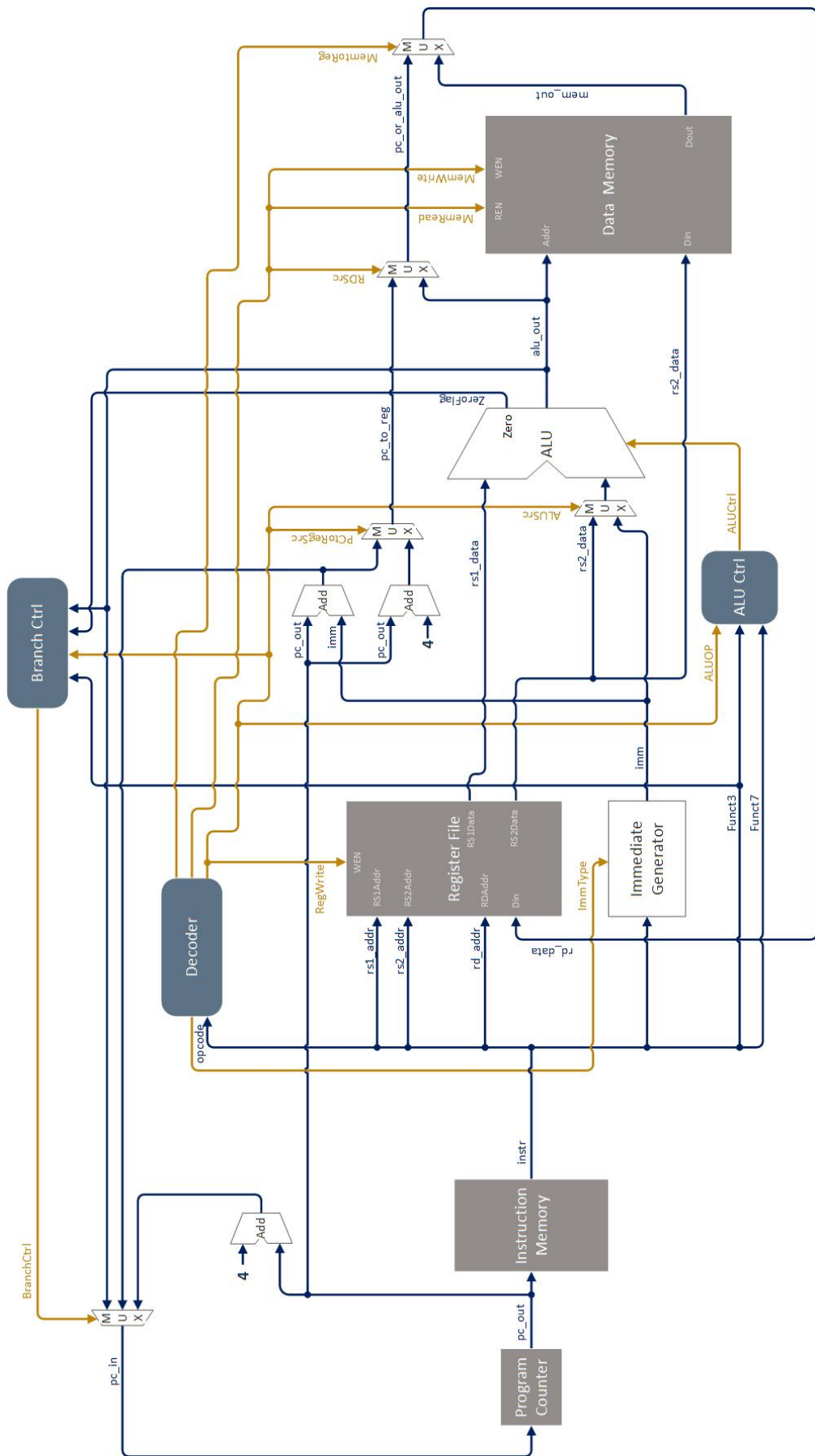
※You do not need to modify this module

e. **CPU module**

“CPU” is responsible for connecting wires between modules, please add your code to complete this module

※You should modify this module.

● Reference Block Diagram



- **Register File**

Register	ABI Name	Description	Saver
x0	zero	Hard-wired zero	---
x1	ra	Return address	Caller
x2	sp	Stack pointer	Callee
x3	gp	Global pointer	---
x4	tp	Thread pointer	---
x5	t0	Temporary / alternate link register	Caller
x6 - 7	t1 - 2	Temporaries	Caller
x8	s0/fp	Saved register / frame pointer	Callee
x9	s1	Saved register	Callee
x10 - 11	a0 - 1	Function arguments / return values	Caller
x12 - 17	a2 - 7	Function arguments	Caller
x18 - 27	s2 - 11	Saved registers	Callee
x28 - 31	t3 - 6	Temporaries	Caller

- **Test Instruction**

- a. **Memory layout**

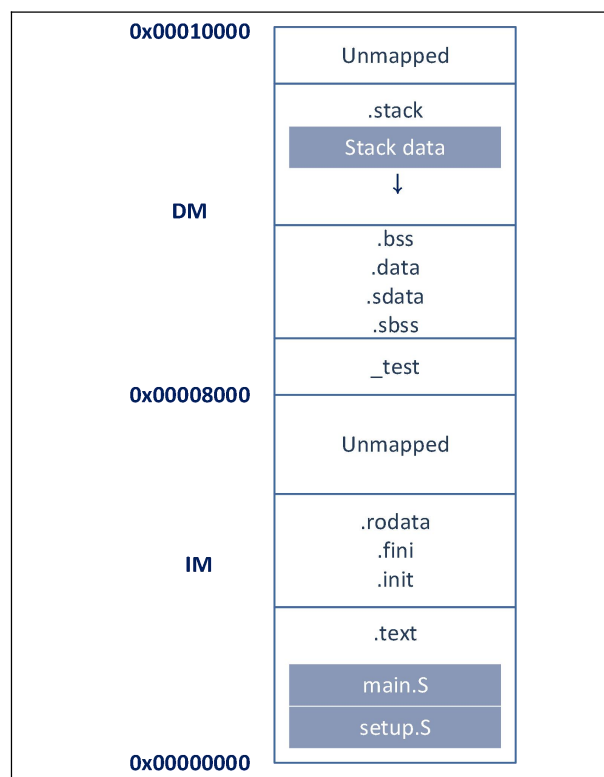


Figure 2. Memory layout

- .text: Store instruction code.
- .init & .fini: Store instruction code for entering & leaving the process.
- .rodata: Store constant global variable.
- .bss & .sbss: Store uninitiated global variable or global variable initiated as zero.
- .data & .sdata: Store global variable initiated as non-zero
- .stack: Store local variables

b. main.S

This will verify RISC-V instructions above (17 instructions).

c. main0.hex & main1.hex & main2.hex & main3.hex

Using the cross compiler of RISC-V to compile test program, and write result in verilog format. So you do not need to compile above program again.

Homework Requirements

1. Complete the CPU that can execute 17 instructions from the *RISC-V ISA* section.
2. Verify your CPU with the benchmark and take a snapshot (e.g. Figure 3)

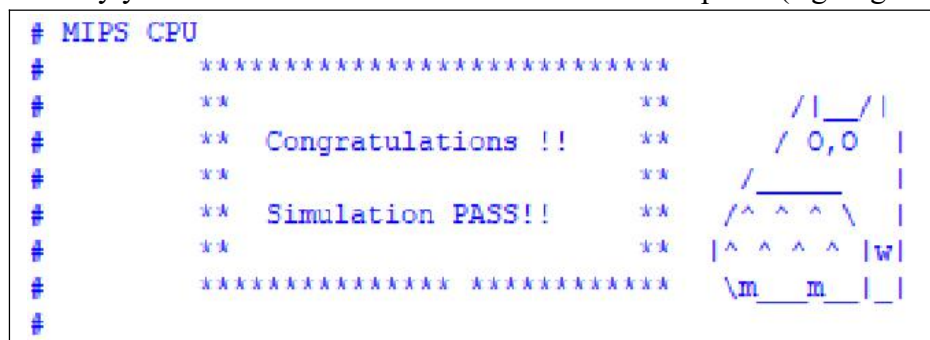


Figure 3. Snapshot of correct simulation

- a. You can verify the execution results by checking waveforms.
3. Finish the Project Report.
 - a. Complete the project report. The report template is provided “report.docx”.

Important

When you upload your file, please make sure you have satisfied all the homework requirements, including the **File hierarchy, Requirement file and Report format**.

If you have any questions, please contact us.

Score

Your score is divided into two parts:

- a. Functional Simulation (75%): TA will give you score based on the number of correct results. There are 15 test data, if you pass one of them, you will get 5 points.
- b. Report (25%): You should take a screenshot of your result, and write your report in “report.docx”.