



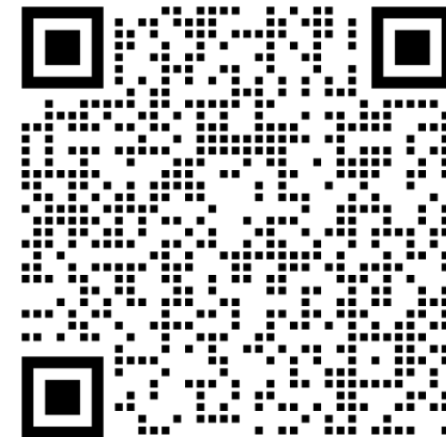
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Adversarial Search

Data Intelligence and Social Computing Lab (DISC)

October 26th, 2021



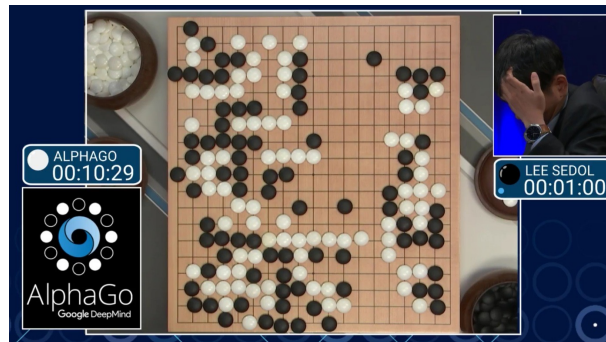
- Type of game

Types of Games

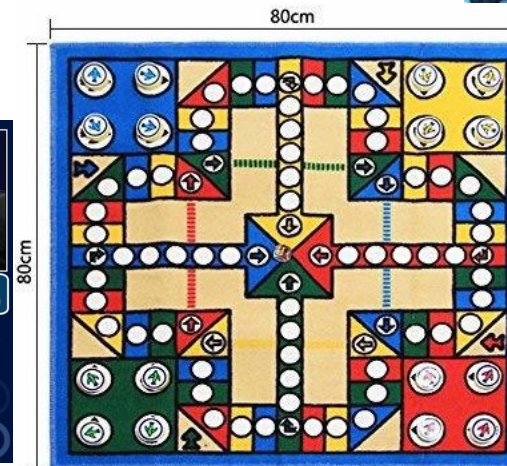
- **Deterministic or stochastic ?**
 - GO, Chess VS Aeroplane Chess, Monopoly
- **One, two, or more players ?**
- **Zero sum or not ?**
 - GO VS Contra
- **Perfect information ?**
 - GO VS DOTA



Chess



GO



Aeroplane Chess



Monopoly

- Zero-Sum Games
 - Agents have opposite utilities
 - Think of a single value that one maximizes and the other minimizes
 - **Adversarial**, pure competition

- General Games
 - Agents have independent utilities
 - Cooperation, competition, indifference and more are all possible

Example: Non-zero sum game

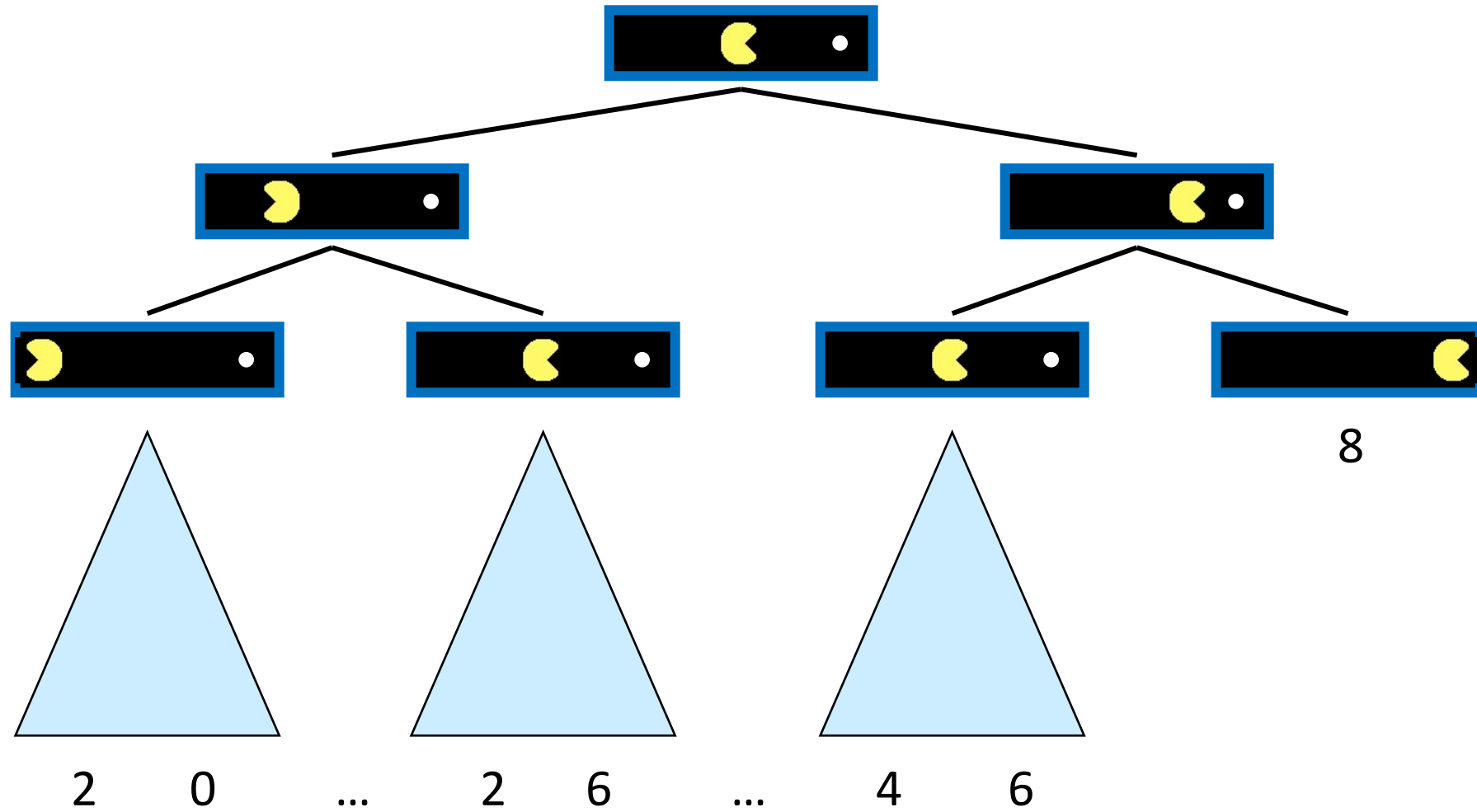
The Prisoners' Dilemma

		Prisoner A Choices	
		<i>Stay Silent</i>	<i>Confess and Betray</i>
Prisoner B Choices	<i>Stay Silent</i>	Each serves one month in jail 2 months in total	Prisoner A goes free 12 months in total Prisoner B serves full year in jail
	<i>Confess and Betray</i>	Prisoner A serves full year in jail 12 months in total Prisoner B goes free	Each serves three months in jail 6 months in total

- S : states (start at s_0)
 - **Player (s)**: the player has the move in this state
 - Actions (s): A set of legal moves in a state
 - Results (s, a): A transition model, return the results of a move
 - Terminal Test (s): {true, false} if s is the terminal state
 - **Terminal Utilities (s, p)**: A utility function gives the final numeric value of a game
-
- Zero Sum

- Type of Game
- Adversarial Search

Single-Agent Trees

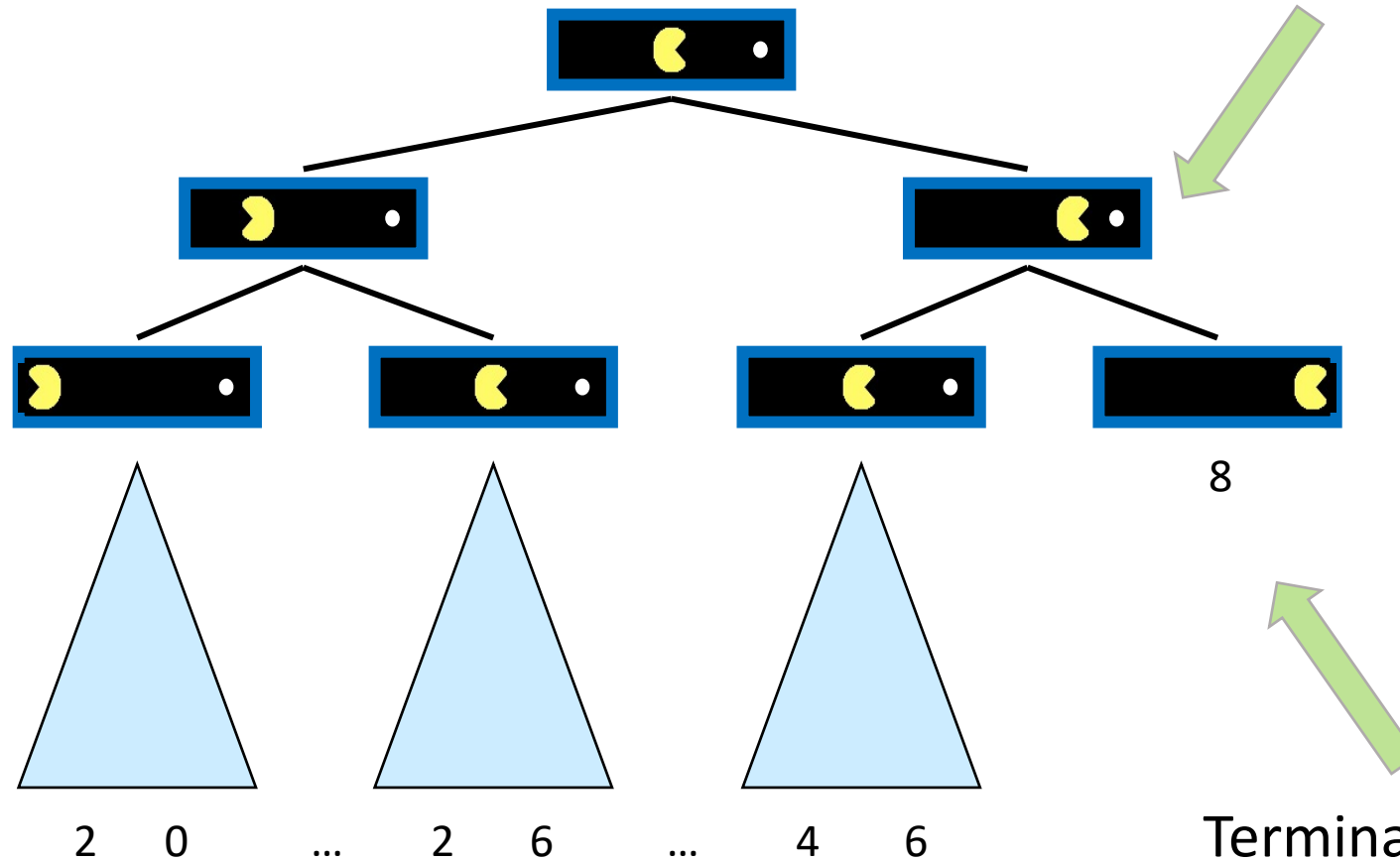


Value of a State

Value of a state: The **best achievable outcome** (utility) from that state

Non-Terminal States:

$$V(s) = \max_{s' \in \text{children}(s)} V(s')$$



Terminal States:

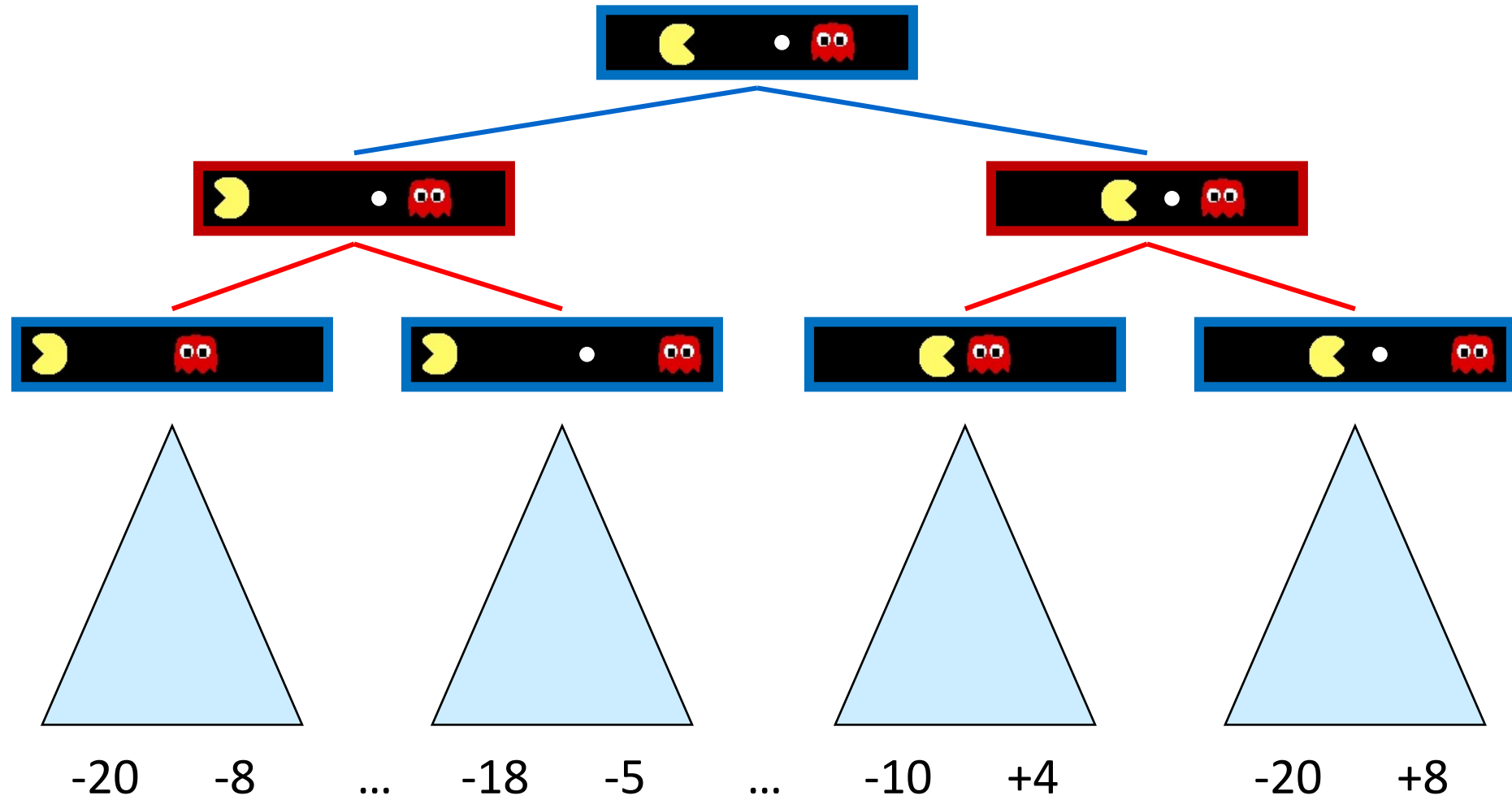
$$V(s) = \text{known}$$

Adversarial Game Trees

Xiao Huang

Ghost

Xiao Huang



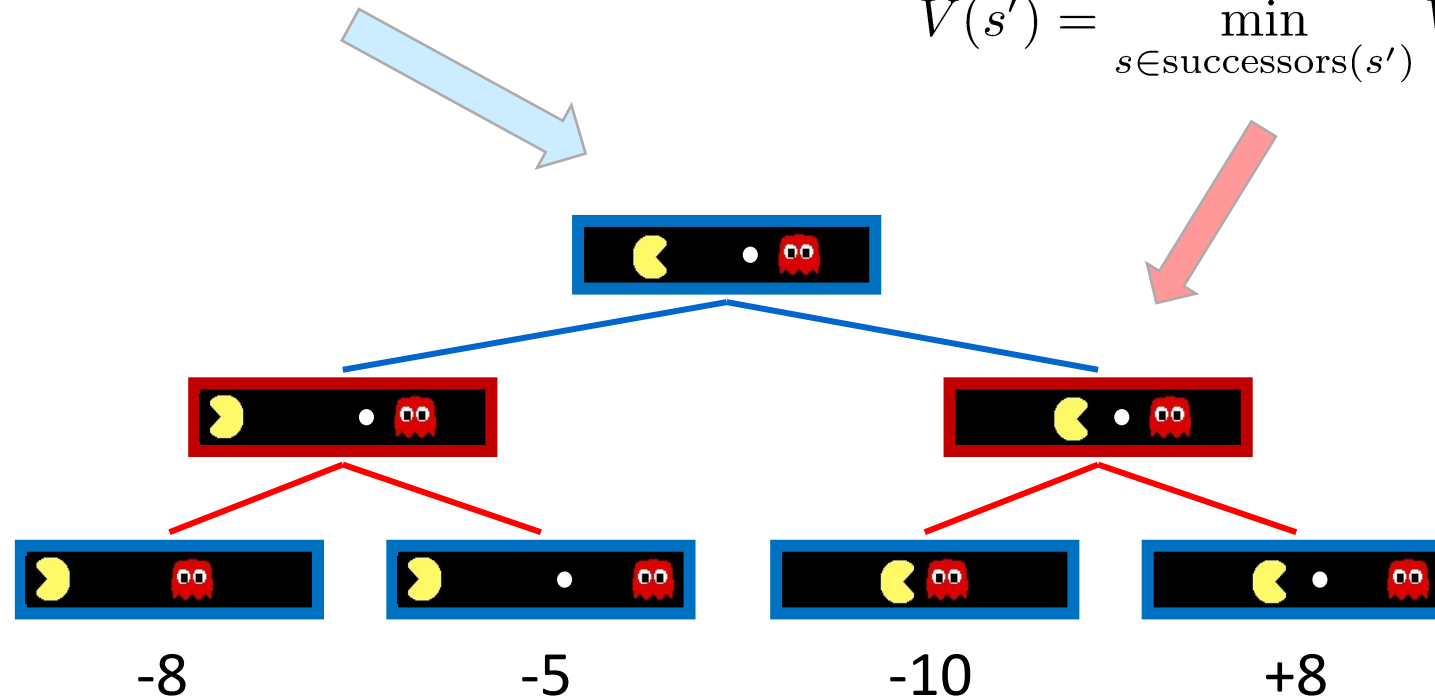
Minimax Values

States Under Agent's Control:

$$V(s) = \max_{s' \in \text{successors}(s)} V(s')$$

States Under Opponent's Control:

$$V(s') = \min_{s \in \text{successors}(s')} V(s)$$



Terminal States:

$$V(s) = \text{known}$$

Tic-Tac-Toe Game Tree



MAX (X)



MIN (O)



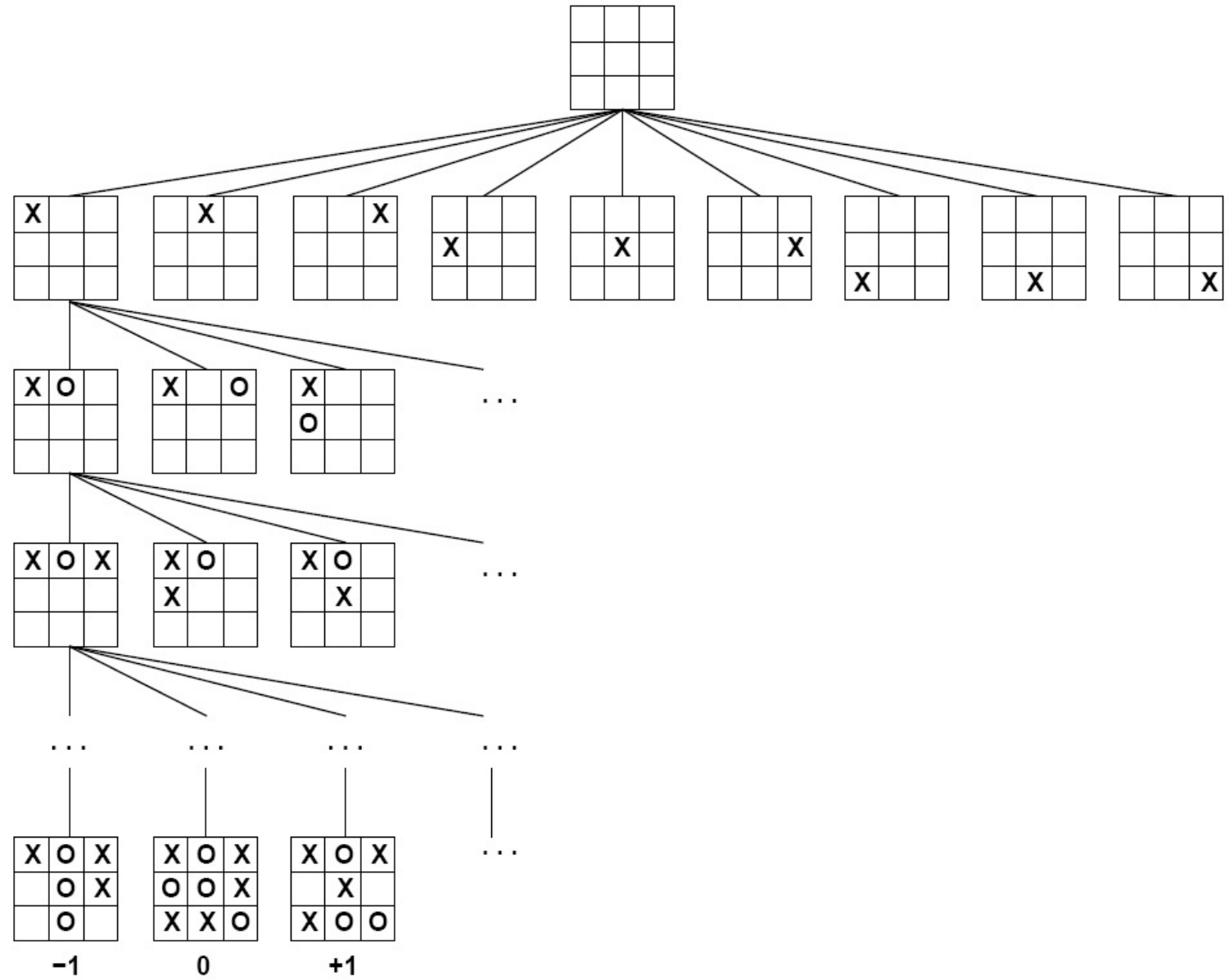
MAX (X)



MIN (O)

TERMINAL

Utility

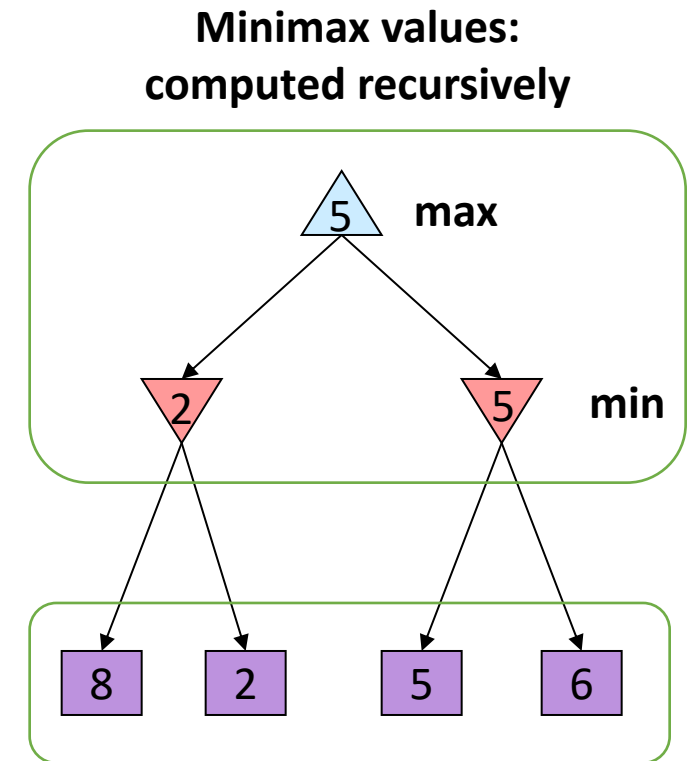


- Deterministic, zero-sum games:

- Tic-tac-toe, chess
- One player maximizes result
- The other minimizes result

- Minimax search:

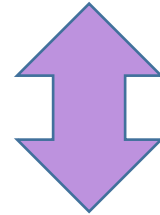
- A state-space search tree
- Players alternate turns
- Compute each node's **minimax value**: the best achievable utility against a rational (optimal) adversary



Minimax Implementation

```
def max-value(state):  
    initialize v =  $-\infty$   
    for each successor of state:  
        v = max(v, min-value(successor))  
    return v
```

$$V(s) = \max_{s' \in \text{successors}(s)} V(s')$$



```
def min-value(state):  
    initialize v =  $+\infty$   
    for each successor of state:  
        v = min(v, max-value(successor))  
    return v
```

$$V(s') = \min_{s \in \text{successors}(s')} V(s)$$

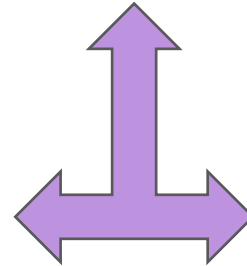
Minimax Implementation (Dispatch)

```
def value(state):
```

if the state is a terminal state: return the state's utility

if the next agent is MAX: return max-value(state)

if the next agent is MIN: return min-value(state)



```
def max-value(state):
```

initialize $v = -\infty$

for each successor of state:

$v = \max(v, \text{value}(\text{successor}))$

return v

```
def min-value(state):
```

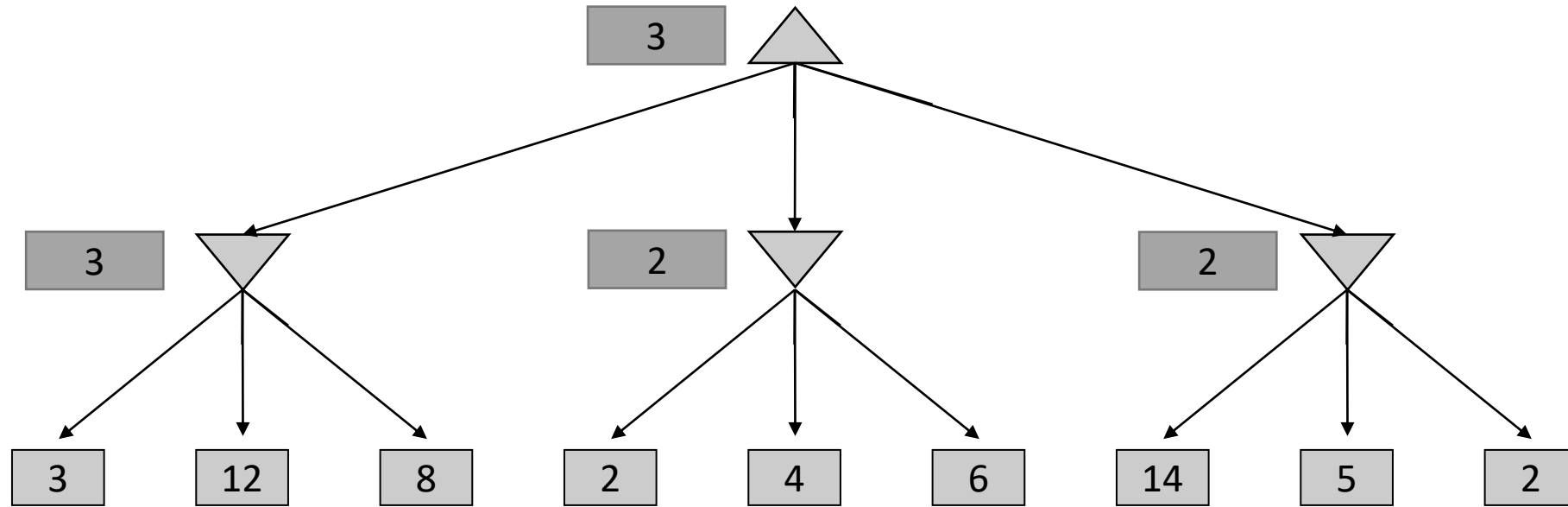
initialize $v = +\infty$

for each successor of state:

$v = \min(v, \text{value}(\text{successor}))$

return v

Minimax Example



1. Value inference : use minimax search to identify the value for each node
2. Action generation : generate the path based on the value

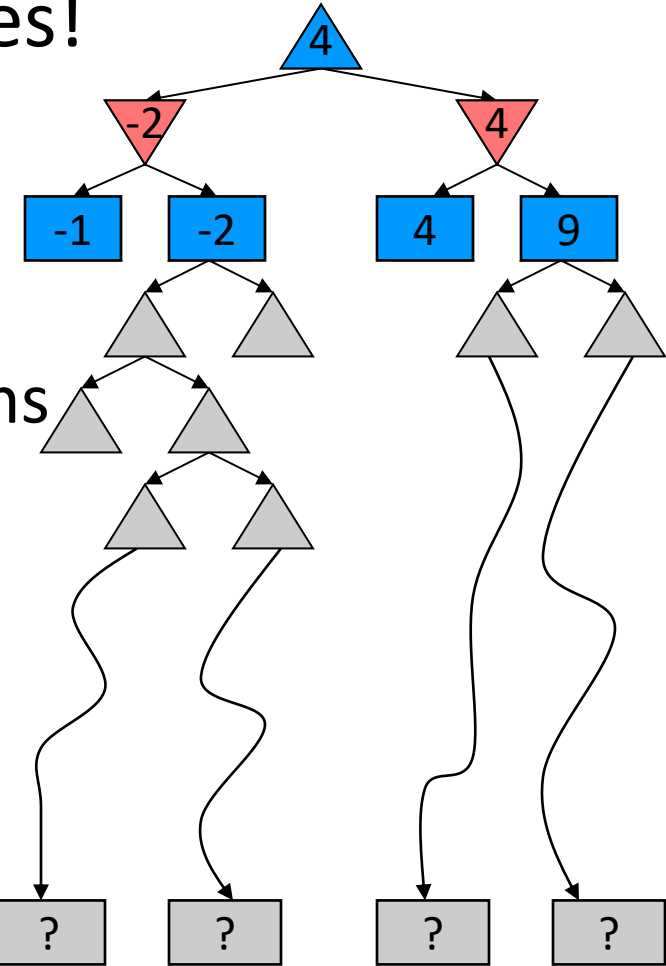
- Type of Game
- Adversarial Search
- Evaluation Function

- How efficient is minimax?
 - Just like (exhaustive) DFS
 - Time: $O(b^m)$
 - Space: $O(bm)$

- Example: For chess, $b \approx 35$, $m \approx 100$
 - Exact solution is completely infeasible
 - 8^{140} years for a computer that can process 1,000,000 nodes per second
 - But, do we need to explore the whole tree?

Depth-limited Search with Evaluation Function

- Problem: In realistic games, cannot search to leaves!
- Solution: Depth-limited search
 - Search only to a limited depth in the tree
 - Need an **evaluation function** for non-terminal positions
- Guarantee of optimal play is gone
- More steps forward makes a BIG difference
- How to determine an appropriate depth for search
 - Use iterative deepening



- Evaluation functions score non-terminals in depth-limited search
- Linear combination of various factors

$$Eval(s) = w_1 f_1(s) + w_2 f_2(s) + \dots + w_n f_n(s)$$

Eval(s) = material + mobility + king-safety + center-control

material = $10^{100}(K - K') + 9(Q - Q') + 5(R - R') + 3(B - B' + N - N') + 1(P - P')$

mobility = $0.1(\text{num-legal-moves} - \text{num-legal-moves}')$

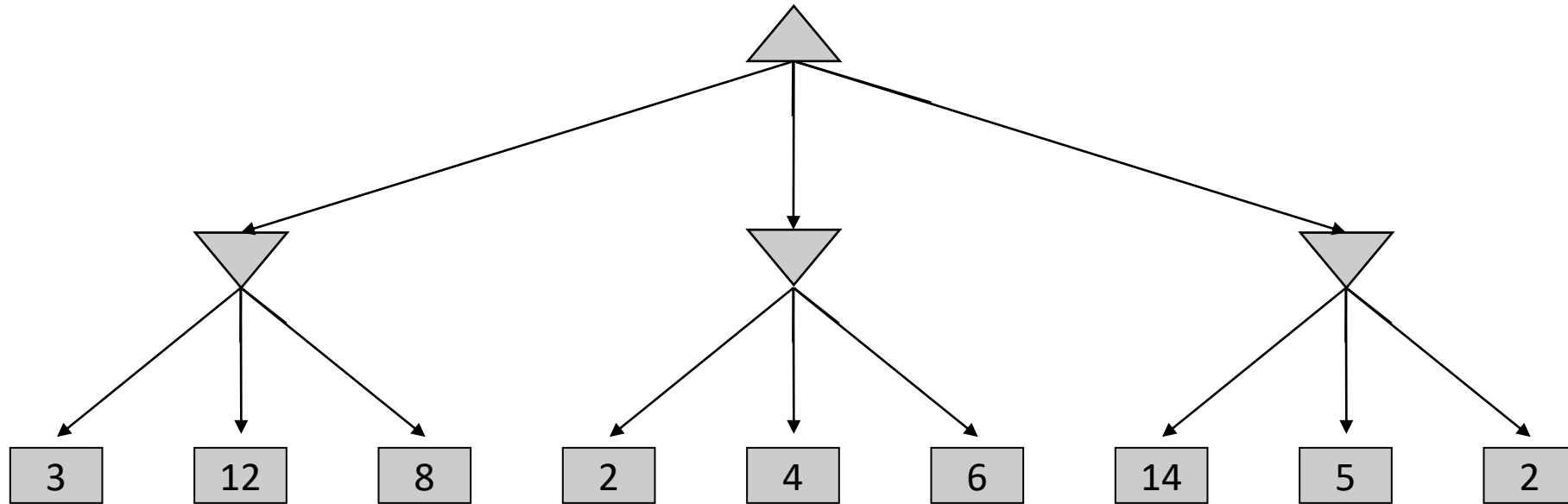
- Table-based evaluation function
 - Compute values for states in advance and store them in a table
 - Too many states to compute. Can be done for the starting and the ending.
- Machine-learning based evaluation function

- ✓ Principles for designing evaluation function
 - ✓ Utility for a win state should be higher than a tie. (correct)
 - ✓ The computation of evaluation function should be efficient. (quick)
 - ✓ It should be related to the chance of winning the game. (consistent)

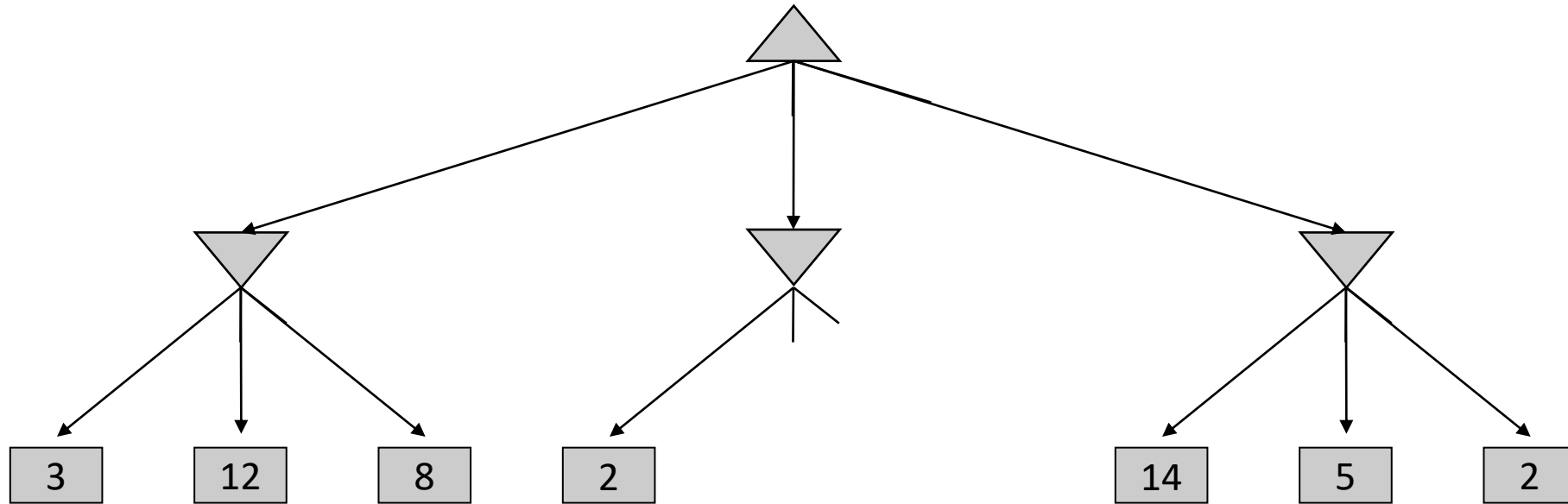
- **It takes time to compute the evaluation function.**
 - An important example of the tradeoff between **complexity of features** and **complexity of computation**
 - Ideal function: returns the actual minimax value of the position

- Type of Game
- Adversarial Search
- Evaluation Function
- Game Tree Pruning

Minimax Example

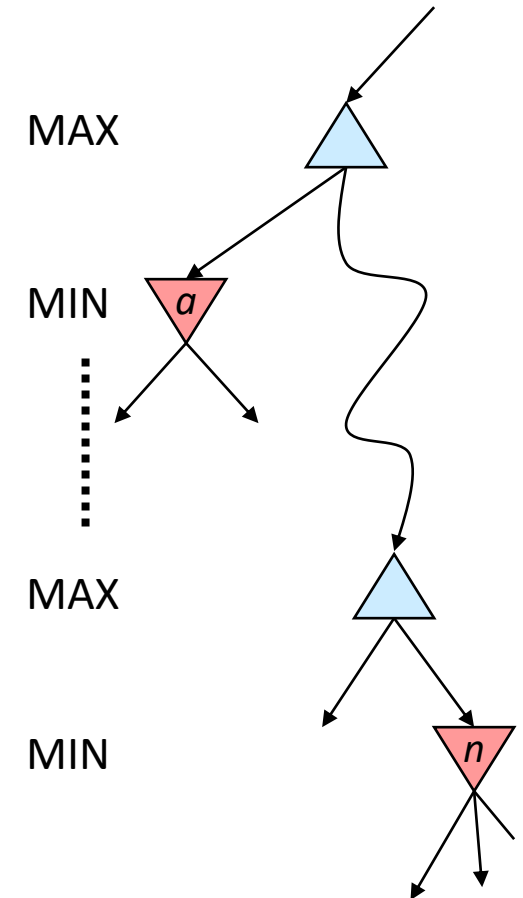


Minimax Pruning



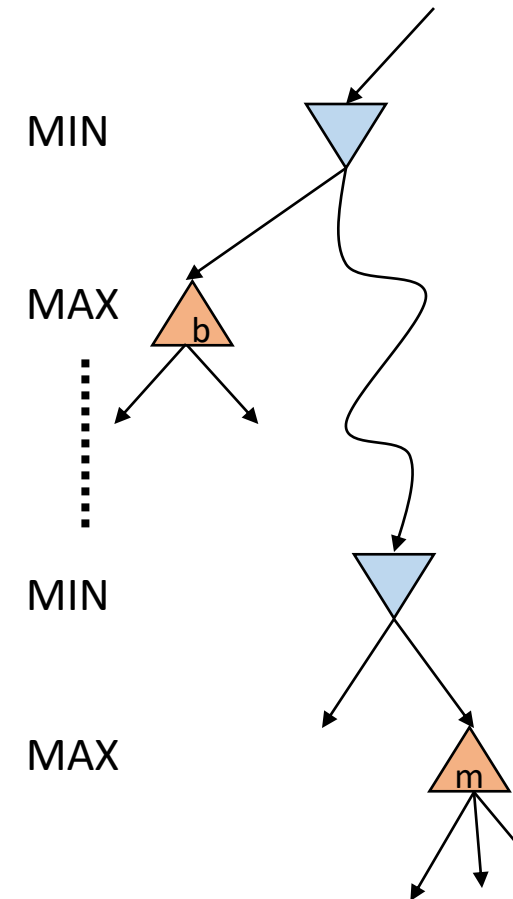
Alpha-Beta Pruning for MIN node

- Suppose we have a MAX node as root and it has **the best value α**
- We're computing the MIN-VALUE at some node n
- We're looping over n 's children
- **n 's estimate of the children's min is decreasing**
- **If n becomes worse than α , MAX will avoid it, so we can stop considering n 's other children**



Alpha-Beta Pruning for MAX node

- Suppose we have a MIN node as root and it has **the best value b**
- We're computing the MAX-VALUE at some node m
- We're looping over m 's children
- m 's estimate of the children's' max is increasing
- If m becomes bigger than b , MIN will avoid it, so we can stop considering m 's other children



- Set a range of value for each node with two parameters
 - $[\text{Alpha}, \text{Beta}]$
 - Obtain the initial value from parent node
- Update Process
 - Alpha: update in a max node, the best value it can achieve
 - Beta: update in a min node, the least value it can achieve
- Pruning Process
 - MAX node: if the current value is higher than beta, prune the subtree
 - MIN node: if the current value is lower than alpha, prune the subtree

α : MAX's best option on path to root
 β : MIN's best option on path to root

def max-value(state, α , β):

 initialize $v = -\infty$

 for each successor of state:

update of v $v = \max(v, \text{value}(\text{successor}, \alpha, \beta))$

update of alpha $\alpha = \max(\alpha, v)$

pruning if $v \geq \beta$ return v

 return v

def min-value(state, α , β):

 initialize $v = +\infty$

 for each successor of state:

update of v $v = \min(v, \text{value}(\text{successor}, \alpha, \beta))$

update of beta $\beta = \min(\beta, v)$

pruning if $v \leq \alpha$ return v

 return v

- Update alpha in MAX node, and prune subtree of a MAX node using beta
- Update beta in MIN node, and prune subtree of a MIN node using alpha

Adversarial Search (Minimax)

function ALPHA-BETA-SEARCH(*state*) **returns** an action
 $v \leftarrow \text{MAX-VALUE}(\text{state}, -\infty, +\infty)$
return the *action* in $\text{ACTIONS}(\text{state})$ with value v

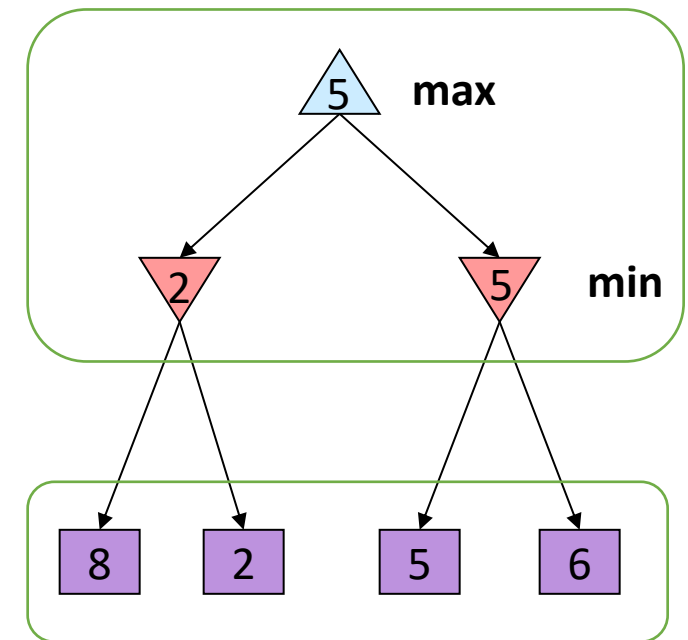
function MAX-VALUE(*state*, α , β) **returns** a utility value
if $\text{TERMINAL-TEST}(\text{state})$ **then return** $\text{UTILITY}(\text{state})$
 $v \leftarrow -\infty$
for each a **in** $\text{ACTIONS}(\text{state})$ **do**
 $v \leftarrow \text{MAX}(v, \text{MIN-VALUE}(\text{RESULT}(s, a), \alpha, \beta))$
if $v \geq \beta$ **then return** v
 $\alpha \leftarrow \text{MAX}(\alpha, v)$
return v

function MIN-VALUE(*state*, α , β) **returns** a utility value
if $\text{TERMINAL-TEST}(\text{state})$ **then return** $\text{UTILITY}(\text{state})$
 $v \leftarrow +\infty$
for each a **in** $\text{ACTIONS}(\text{state})$ **do**
 $v \leftarrow \text{MIN}(v, \text{MAX-VALUE}(\text{RESULT}(s, a), \alpha, \beta))$
if $v \leq \alpha$ **then return** v
 $\beta \leftarrow \text{MIN}(\beta, v)$
return v

Figure 5.7 The alpha-beta search algorithm. Notice that these routines are the same as the MINIMAX functions in Figure 5.3, except for the two lines in each of MIN-VALUE and MAX-VALUE that maintain α and β (and the bookkeeping to pass these parameters along).

- Alpha-Beta-Search (*state*): the entry of the search; *State* is the root of the search
- MAX-VALUE(*state*, α , β): compute the value for a MAX node; *state* is the current state; α is the MAX's best option on path to root; β is the MIN's best option on path to root.
- MIN-VALUE(*state*, α , β): compute the value for a MIN node
- ACTIONS (*state*): get all the possible actions for *state*
- RESULT (*state*, *action*): get the successor state for current state

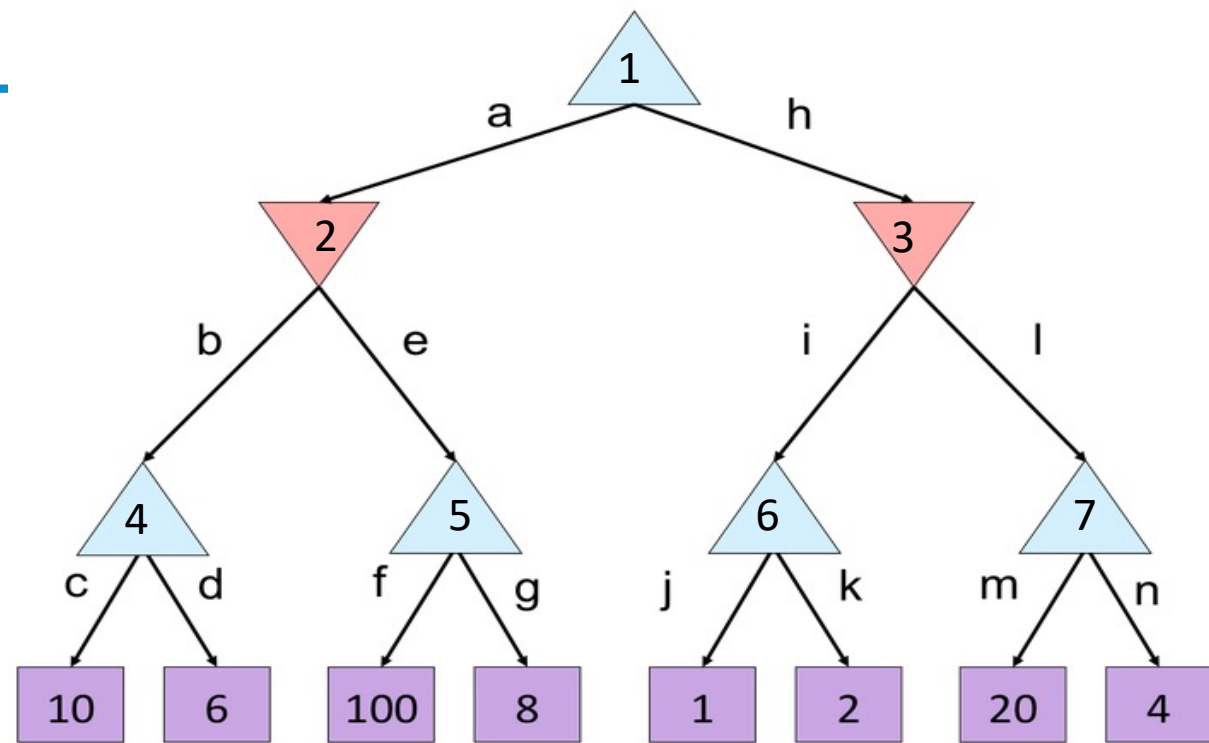
**Minimax values:
computed recursively**



**Terminal values:
part of the game**

In class exercise

node	v	Alpha	beta
1			
2			
3			
4			
5			
6			
7			



1. Which edges are pruned?
2. Write down the value, alpha and beta for each node after processing.

```
def max-value(state,  $\alpha$ ,  $\beta$ ):
```

 If terminal (state),

 return value

 initialize $v = -\infty$

 for each successor of state:

$v = \max(v, \text{value}(\text{successor}, \alpha, \beta))$

$\alpha = \max(\alpha, v)$

if $v \geq \beta$ return v

 return v

```
def min-value(state,  $\alpha$ ,  $\beta$ ):
```

 If terminal (state),

 return value

 initialize $v = +\infty$

 for each successor of state:

$v = \min(v, \text{value}(\text{successor}, \alpha, \beta))$

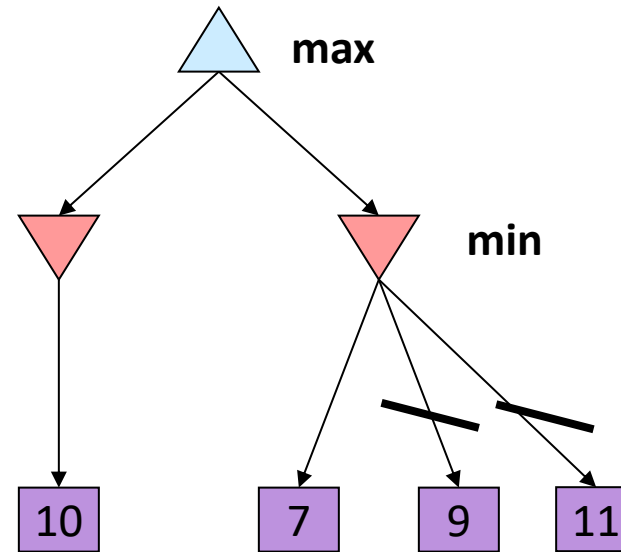
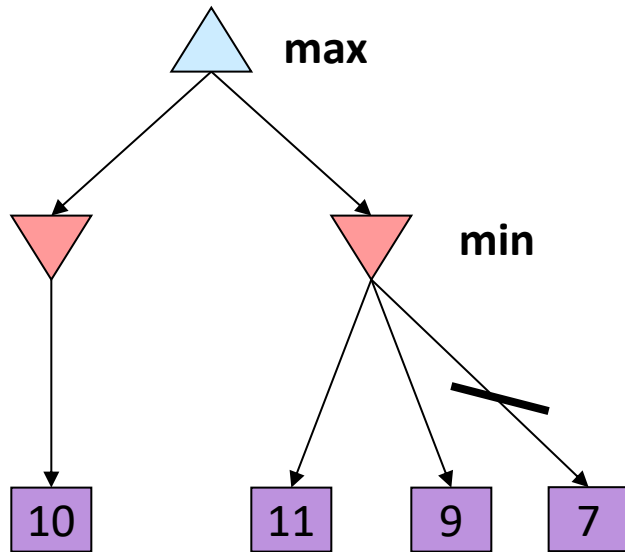
$\beta = \min(\beta, v)$

if $v \leq \alpha$ return v

 return v

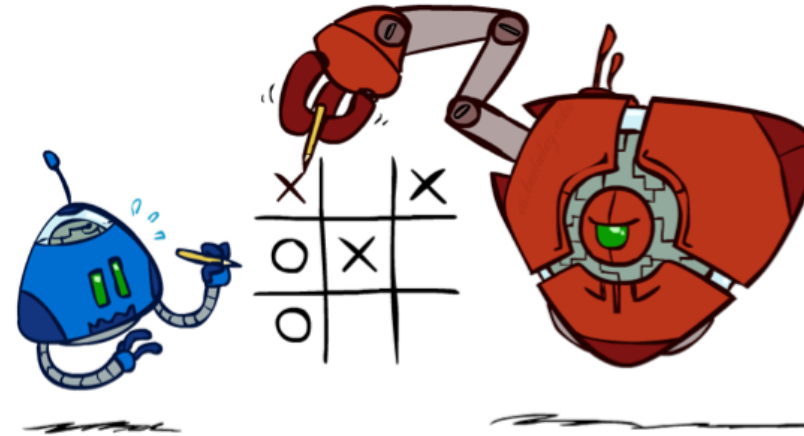
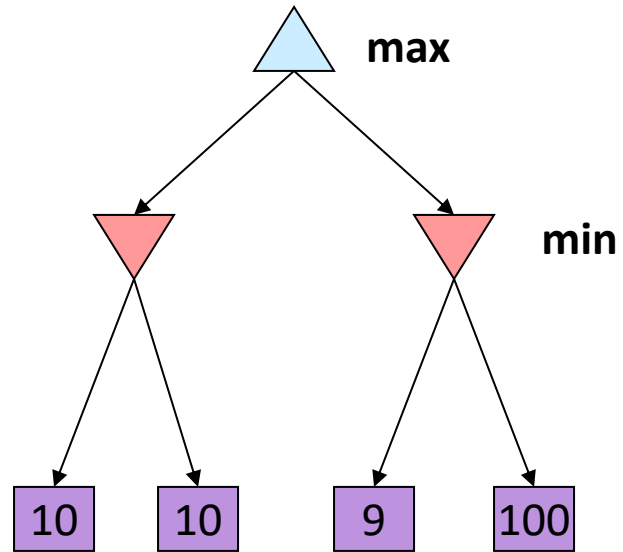
Alpha-Beta Pruning Properties

- This pruning has **no effect** on minimax value computed for the root!
- Minimax values of intermediate nodes might be wrong
- Ordering Matters
 - Good child ordering improves effectiveness of pruning
 - MAX node needs decreasing order of children nodes
 - MIN node needs increasing order of children nodes

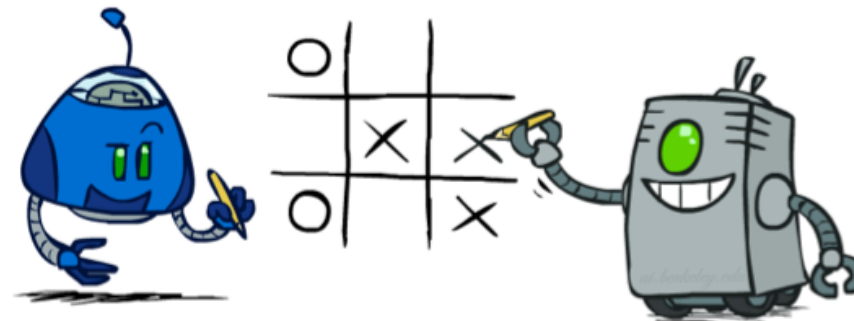


- Type of Game
- Adversarial Search
- Evaluation Function
- Game Tree Pruning
- Uncertain Outcomes

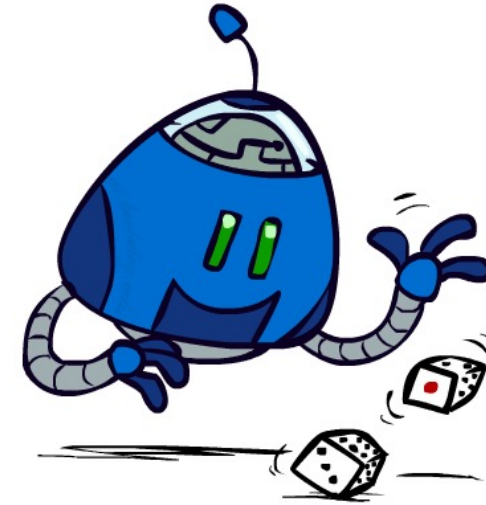
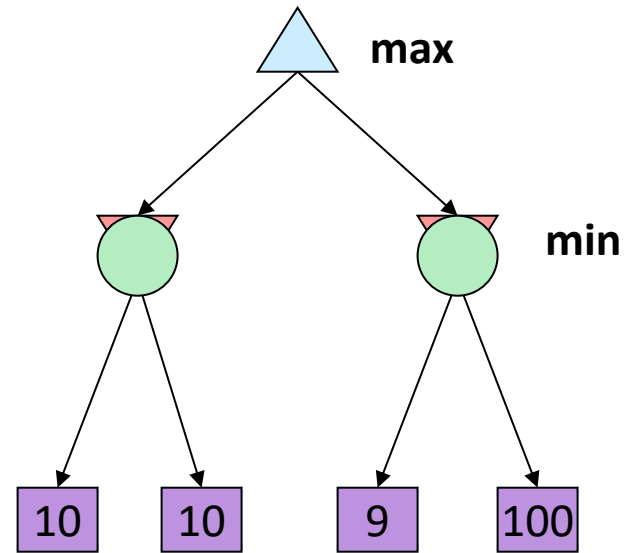
Minimax Properties



Optimal against a perfect player. Otherwise?



Worst-Case vs. Average Case



Idea: Uncertain outcomes controlled by chance.

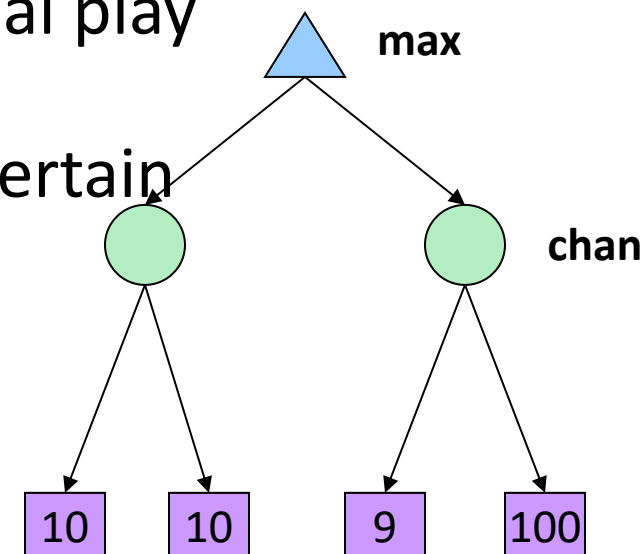
- A **random variable** represents an event whose outcome is unknown
- A **probability distribution** is an assignment of weights to outcomes
- Some laws of probability:
 - Probabilities are always non-negative
 - Probabilities over all possible outcomes sum to one
- Example: Traffic on freeway
 - Random variable: T = whether there's traffic
 - Outcomes: T in {none, light, heavy}
 - Distribution: $P(T=\text{none}) = 0.25$, $P(T=\text{light}) = 0.50$, $P(T=\text{heavy}) = 0.25$

Reminder: Expectations

- The expected value of a function of a random variable is the average, weighted by the probability distribution over outcomes
- Example: How long to get to the airport?
 - 20 mins if there **no traffic**.
 - 30 mins if there is **light traffic**.
 - 60 mins if there is **heavy traffic**.
 - $P(T=\text{none}) = 0.25$, $P(T=\text{light}) = 0.50$, $P(T=\text{heavy}) = 0.25$

$$\begin{array}{ccccccc} 20 \text{ min} & & 30 \text{ min} & & 60 \text{ min} & & \\ \times & & \times & & \times & & \\ 0.25 & + & 0.50 & + & 0.25 & \rightarrow & 35 \text{ min} \end{array}$$

- Why wouldn't we know what the result of an action will be?
 - Explicit randomness: rolling dice
 - Unpredictable opponents: the ghosts respond randomly
 - Actions can fail: when moving a robot, wheels might slip
- Values should now reflect average-case (expectimax) outcomes, not worst-case (minimax) outcomes
- **Expectimax search**: compute the average score under optimal play
 - Max nodes as in minimax search
 - Chance nodes are like min nodes but the outcome is uncertain
 - Calculate their **expected utilities**
 - I.e. take weighted average (expectation) of children



Expectimax Pseudocode



```
def value(state):
```

```
    if the state is a terminal state: return the state's utility
```

```
    if the next agent is MAX: return max-value(state)
```

```
    if the next agent is EXP: return exp-value(state)
```

```
def max-value(state):
```

```
    initialize  $v = -\infty$ 
```

```
    for each successor of state:
```

```
         $v = \max(v, \text{value}(\text{successor}))$ 
```

```
    return v
```

```
def exp-value(state):
```

```
    initialize  $v = 0$ 
```

```
    for each successor of state:
```

```
         $p = \text{probability}(\text{successor})$ 
```

```
         $v += p * \text{value}(\text{successor})$ 
```

```
    return v
```

Expectimax Pseudocode

```
def exp-value(state):
```

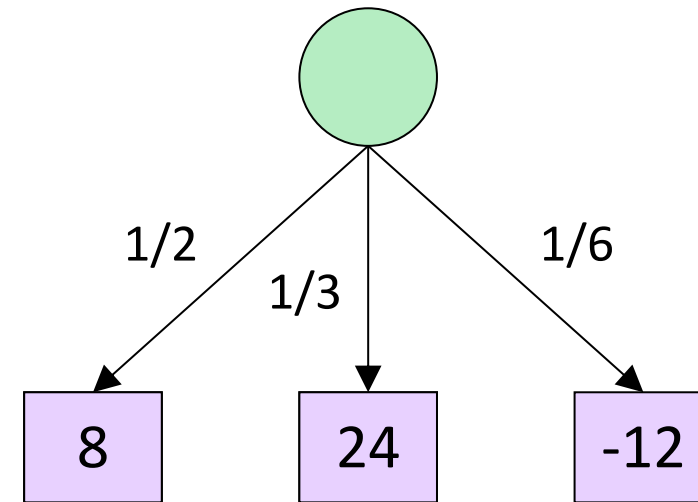
```
    initialize v = 0
```

```
    for each successor of state:
```

```
        p = probability(successor)
```

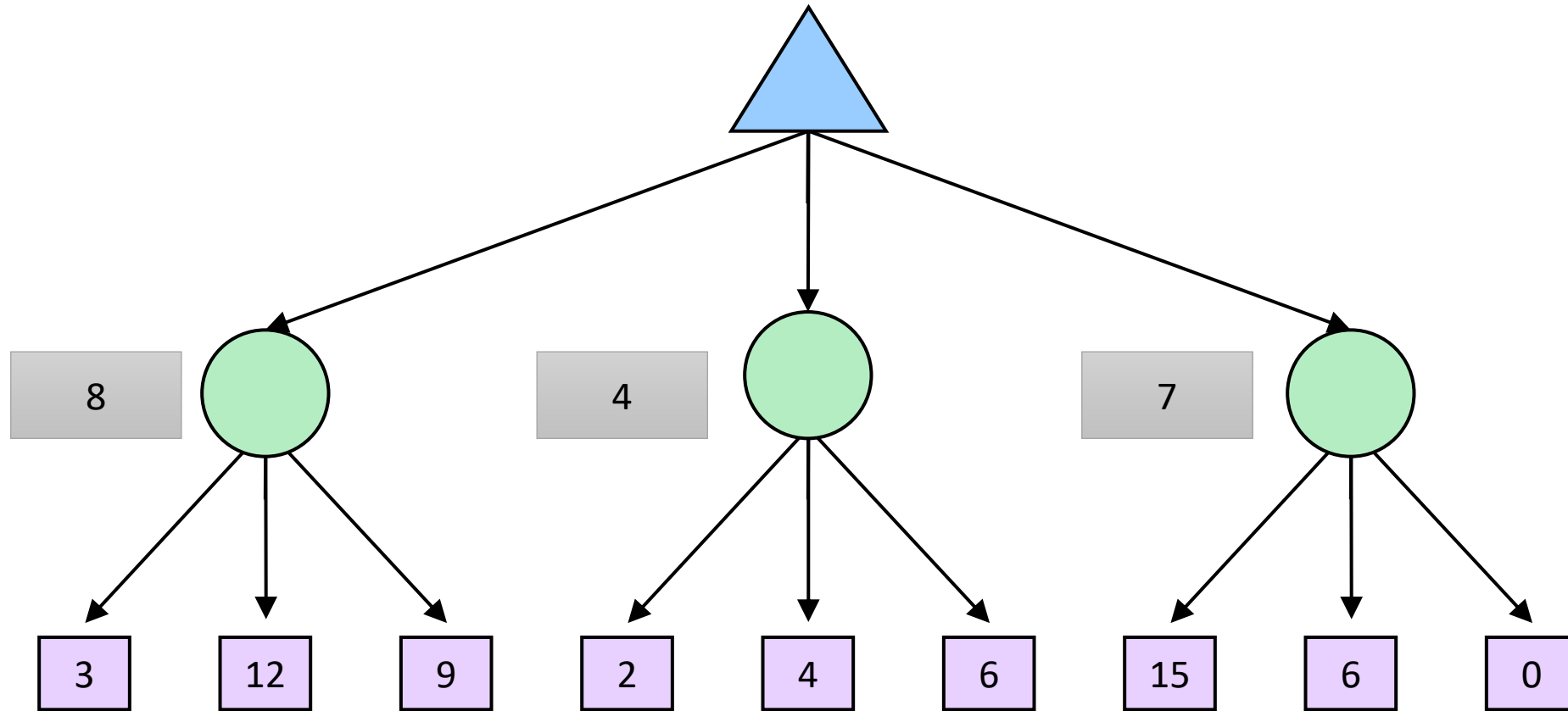
```
        v += p * value(successor)
```

```
    return v
```

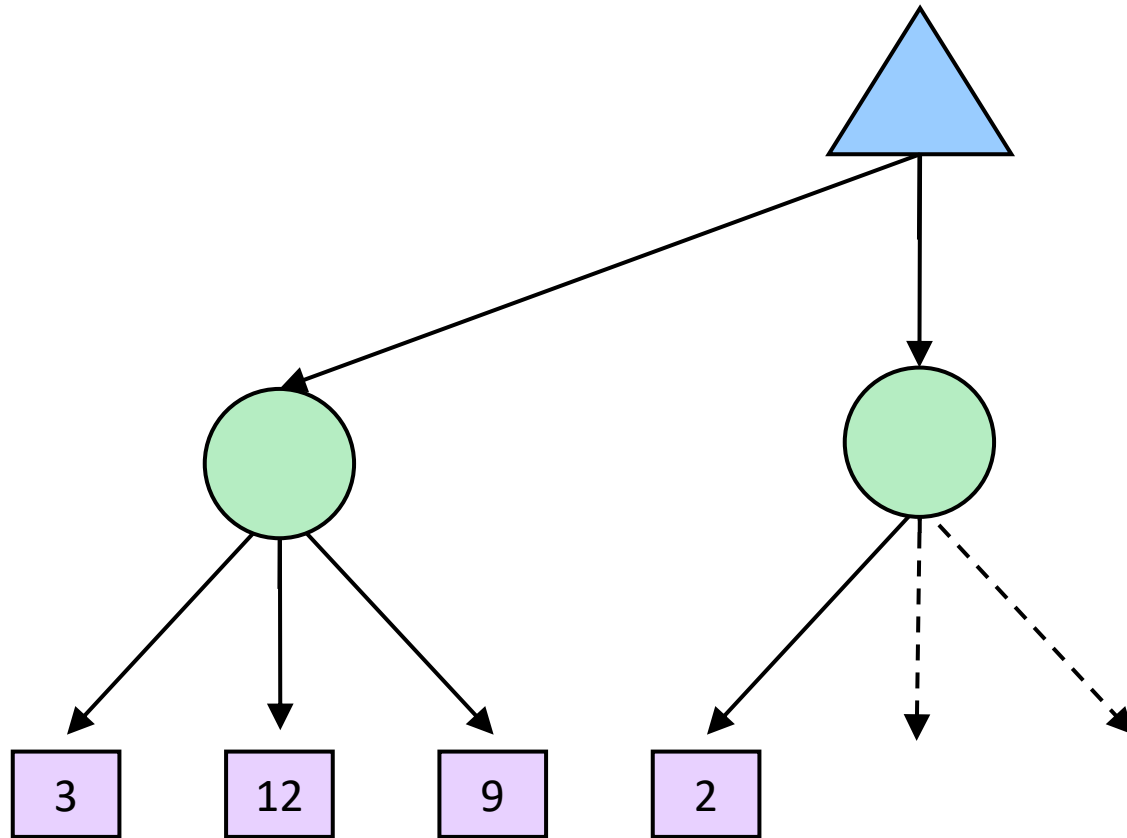


$$v = (1/2) (8) + (1/3) (24) + (1/6) (-12) = 10$$

Expectimax Example

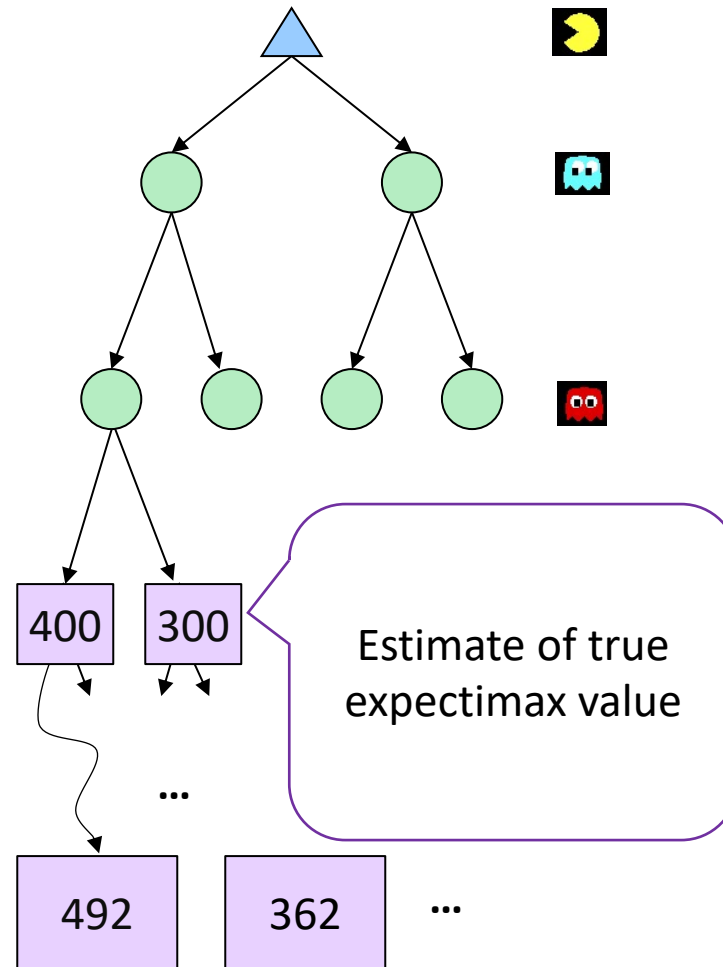


Expectimax Pruning?



No Pruning!
All Children nodes are involved.

Depth-Limited Expectimax



What Probabilities to Use?

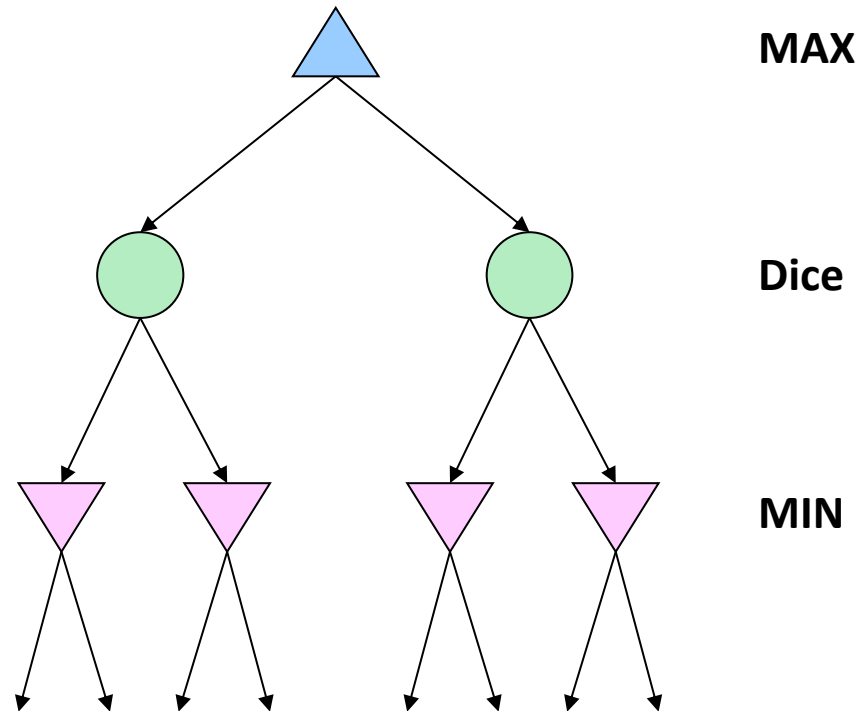


- We have a chance node for any outcome out of our control: opponent or environment
- A probabilistic model to describe how the opponent (or environment) will behave in any state
 - Model could be a simple uniform distribution (roll a die)
 - Model could be sophisticated and require a great deal of computation
 - Model the ghost based on its historical behaviors

- Type of Game
- Adversarial Search
- Evaluation Function
- Game Tree Pruning
- Uncertain Outcomes
- Other Game Types

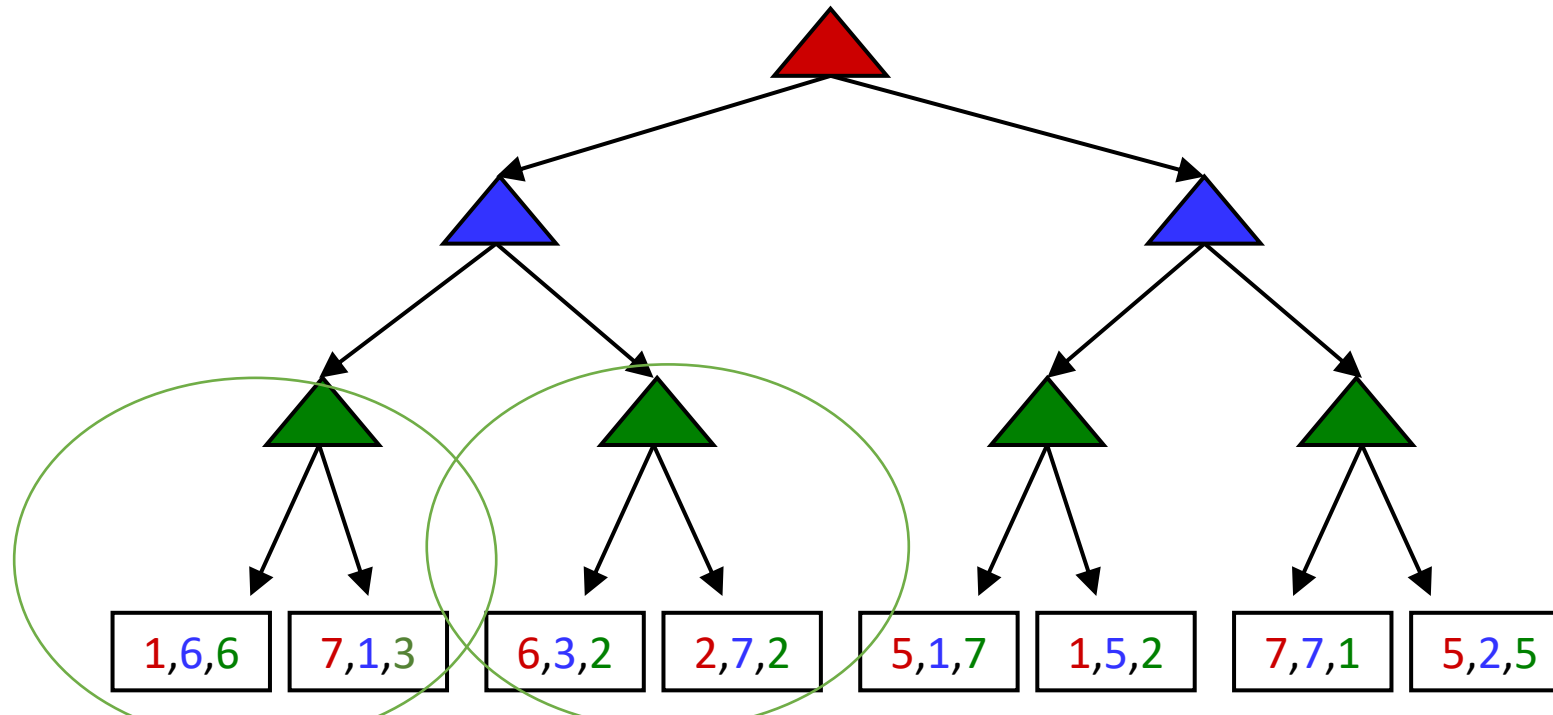
Mixed Layer Types

- E.g. Monopoly with two participants
- Expectiminimax
 - Environment is an extra “random agent” player that moves after each min/max agent
 - Each node computes the appropriate combination of its children



Multi-Agent Utilities

- What if the game is not zero-sum, or has multiple players?
- Generalization of minimax:
 - Terminals have utility tuples
 - Node values are also utility tuples
 - Each player maximizes its own component
 - Can give rise to cooperation and competition dynamically.



- Type of Game
- Adversarial Search
- Game Tree Pruning
- Uncertain Outcomes
- Other Game Types
- Utility

Maximum Expected Utility



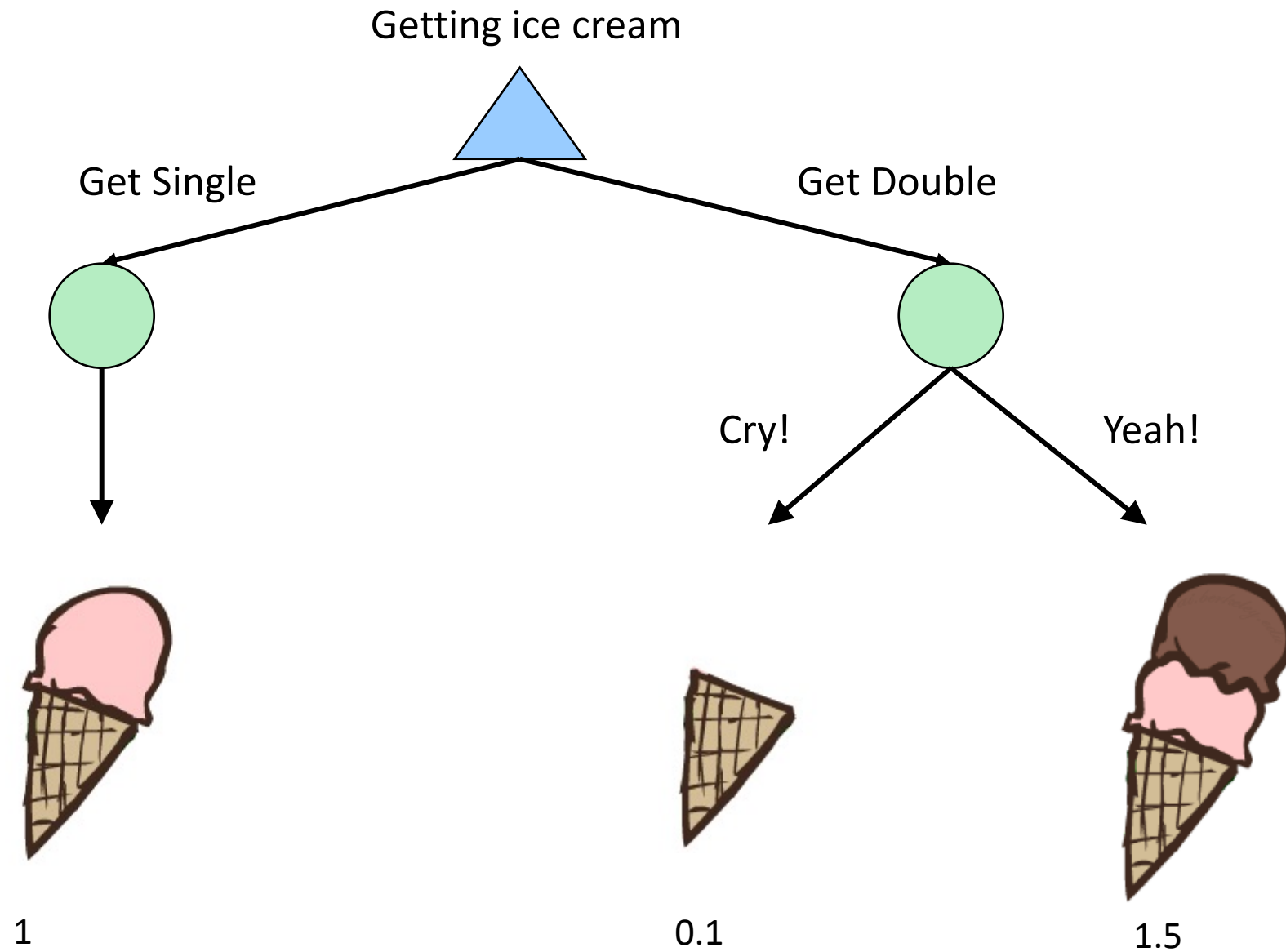
- Principle of maximum expected utility:
 - A rational agent should choose the action that **maximizes its expected utility**, given its knowledge

$$action = \operatorname{argmax} ExpectedUtility(a|e)$$

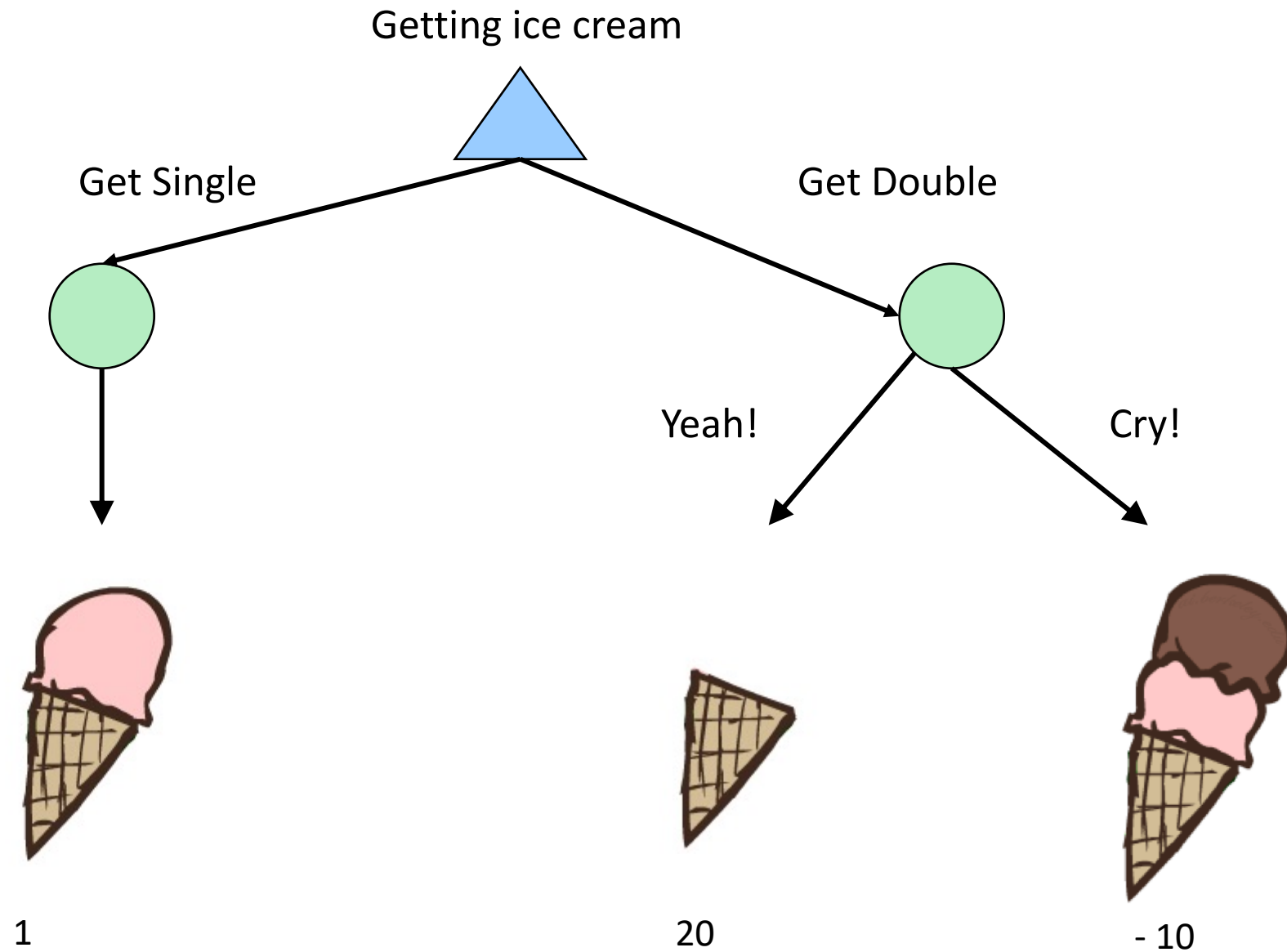
- Questions:
 - Where do utilities come from?
 - How do we ensure the agent is rational with such guidance?

- Utilities are functions **from outcomes (states of the world) to real numbers** that describe an agent's **preferences**
- Where do utilities come from?
 - In a game, may be simple (+1/-1)
 - In general, utilities function summarizes the agent's preferences
- In practice, we design a utility function and let behaviors emerge
- Why bother utility? Why not define behavior directly?
 - Note: an agent can be entirely rational (consistent with MEU) without representing or manipulating utilities and probabilities.
 - E.g., a lookup table for perfect tic-tac-toe, a reflex vacuum cleaner

Utilities: Uncertain Outcomes



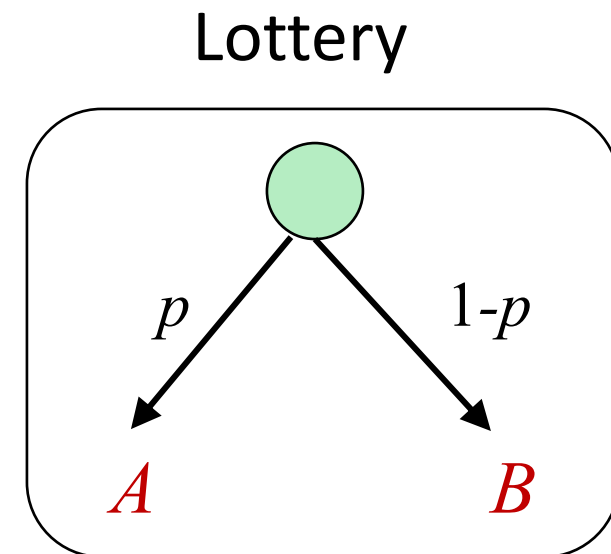
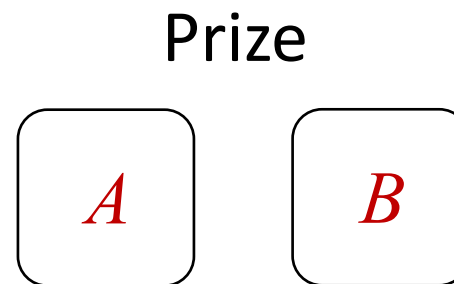
Utilities: Uncertain Outcomes



- An agent must have preferences among:
 - Prize: A , B , etc.
 - Lotteries: situations with uncertain prizes
 - Each uncertain decision can be interpreted as a lottery

$$L = [p, A; (1 - p), B]$$

- Notation:
 - Preference: $A \succ B$
 - Indifference: $A \sim B$

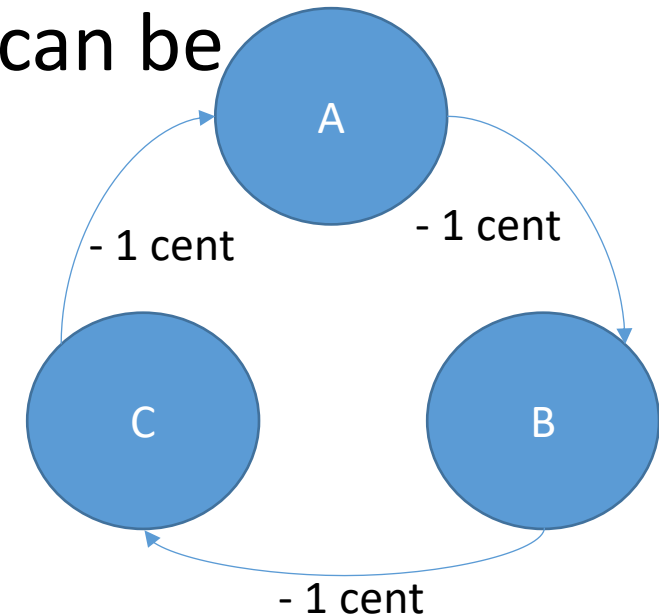


- We want some constraints on preferences before we call them rational, such as:

Axiom of Transitivity: $(A \succ B) \wedge (B \succ C) \Rightarrow (A \succ C)$

- For example: an agent with **intransitive preferences** can be induced to give away all its money

- Say, $C \succ B \succ A$ and $A \succ C$
- If $C \succ B$, then an agent with B would pay 1 cent to get C
- If $B \succ A$, then an agent with A would pay 1 cent to get B
- If $A \succ C$, then an agent with C would pay 1 cent to get A



The Axioms of Rationality

Orderability

$$(A \succ B) \vee (B \succ A) \vee (A \sim B)$$

Transitivity

$$(A \succ B) \wedge (B \succ C) \Rightarrow (A \succ C)$$

Continuity

$$A \succ B \succ C \Rightarrow \exists p [p, A; 1 - p, C] \sim B$$

Substitutability

$$A \sim B \Rightarrow [p, A; 1 - p, C] \sim [p, B; 1 - p, C]$$

Monotonicity

$$A \succ B \Rightarrow \\ (p \geq q \Leftrightarrow [p, A; 1 - p, B] \succeq [q, A; 1 - q, B])$$

Theorem: Rational preferences imply behavior describable as maximization of expected utility \rightarrow Rationality!

- Theorem [Ramsey, 1931; von Neumann & Morgenstern, 1944]
 - Given any preferences satisfying these constraints, there exists a real-valued function U such that:

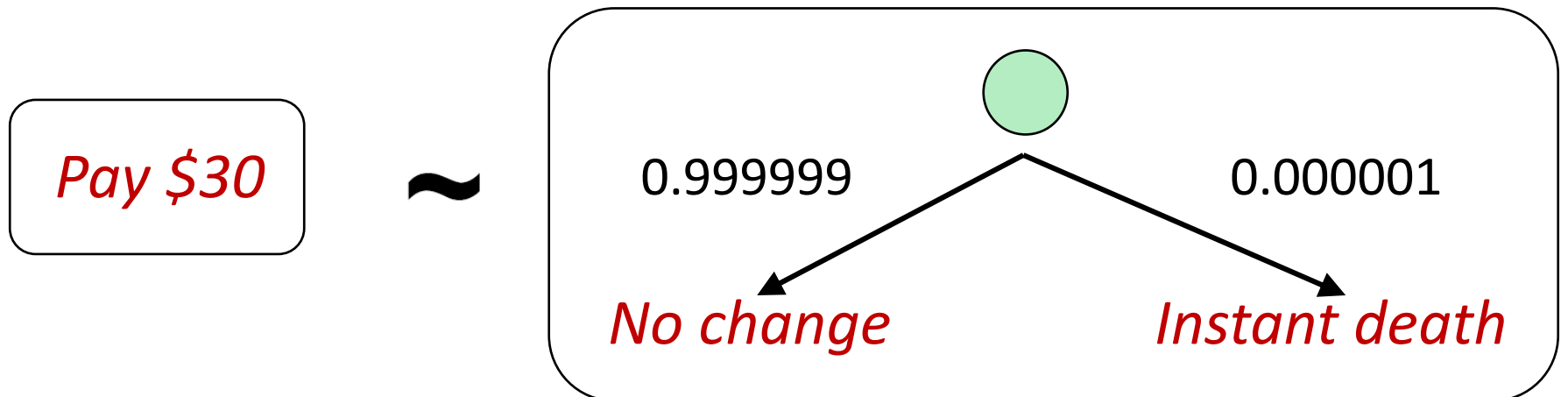
$$U(A) \geq U(B) \Leftrightarrow A \succeq B$$

$$U([p_1, S_1; \dots ; p_n, S_n]) = \sum_i p_i U(S_i)$$

- I.e. values assigned by U preserve preferences of both prizes and lotteries
- Maximum expected utility (MEU) principle:
 - Choose the action that maximizes expected utility
- **Rational Preferences \rightarrow Rational Utility \rightarrow Rational Agent**

- Type of Game
- Adversarial Search
- Game Tree Pruning
- Uncertain Outcomes
- Other Game Types
- Utility
- Human Utility

- **Normalized utilities:** $u_+ = 1.0$, $u_- = 0.0$
- Utilities map states to real numbers.
- Standard approach to assessment of human utilities:
 - Compare a prize A to a **standard lottery** L_p between
 - “best possible prize” u_+
 - “worst possible outcome” u_-
 - Adjust lottery probability p until indifference: $A \sim L_p$



- **Micromorts**: one-millionth chance of death, useful for paying to reduce product risks, etc.
 - 1000 dollar for safety airbag (reduce the **Micromorts from 1000 miles to 6000 miles**)

Micromort examples

- **QALYs** (quality adjusted life years) for medical decisions

QALYs (quality adjusted life year):

quality-adjusted life years, useful for medical decisions involving

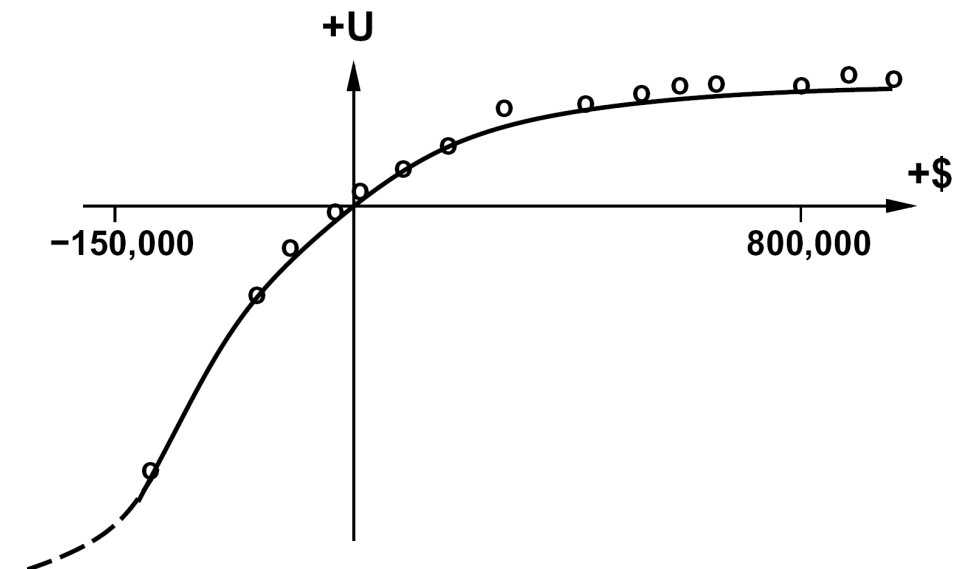
Death from	Micromorts per exposure
Scuba diving	5 per dive
Skydiving	7 per jump
Base-jumping	430 per jump
Climbing Mt. Everest	38,000 per ascent

1 Micromort	
Train travel	6000 miles
Jet	1000 miles
Car	230 miles
Walking	17 miles
Bicycle	10 miles
Motorbike	6 miles

rs, useful



- We can use having money (or being in debt) as the utility.
- Given a lottery $L = [p, \$X; (1-p), \$Y]$
 - The **expected monetary value** $EMV(L)$ is $p \cdot X + (1-p) \cdot Y$
 - $U(L) = p \cdot U(\$X) + (1-p) \cdot U(\$Y)$
- Typically, $U(L) < U(EMV(L))$
- In this sense, people are **risk-averse**
- When deep in debt, people are **risk-seeking**



Example: Insurance



- Consider the lottery $[0.5, \$1000; 0.5, \$0]$
 - What is its **expected monetary value**? (\$500)
 - What is its **certainty equivalent**?
 - \$400 for most people
- Difference of \$100 is the **insurance**
 - There's an insurance industry because people will pay to reduce their risk
 - If everyone were risk-neutral, no insurance needed!
- It's win-win: you'd rather have the \$400 and the insurance company would rather have the lottery (they have many lotteries)

Example: Human Rationality?



- Famous example of Allais (1953)
 - A: [0.8, \$4k; 0.2, \$0]
 - B: [1.0, \$3k; 0.0, \$0]
 - C: [0.2, \$4k; 0.8, \$0]
 - D: [0.25, \$3k; 0.75, \$0]

