Tight Last-Iterate Convergence of the Extragradient Method for Constrained Monotone Variational Inequalities

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Abstract

The monotone variational inequality is a central problem in mathematical programming that unifies and generalizes many important settings such as smooth convex optimization, two-player zero-sum games, convex-concave saddle point problems, etc. The extragradient method by Korpelevich [1976] is one of the most popular methods for solving monotone variational inequalities. Despite its long history and intensive attention from the optimization and machine learning community, the following major problem remains open. What is the last-iterate convergence rate of the extragradient method for monotone and Lipschitz variational inequalities with constraints? We resolve this open problem by showing a tight $O\left(\frac{1}{\sqrt{T}}\right)$ last-iterate convergence rate for arbitrary convex feasible sets, which matches the lower bound by Golowich et al. [2020b]. Our rate is measured in terms of the standard gap function. The technical core of our result is the monotonicity of a new performance measure – the tangent residual, which can be viewed as an adaptation of the norm of the operator that takes the local constraints into account. To establish the monotonicity, we develop a new approach that combines the power of the sum-of-squares programming with the low dimensionality of the update rule of the extragradient method. We believe our approach has many additional applications in the analysis of iterative methods.

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1 Introduction

The *monotone variational inequality (VI)* problem plays a crucial role in mathematical programming, providing a unifying setting for the study of optimization and equilibrium problems. It also serves as a computational framework for numerous important applications in fields such as Economics, Engineering, and Finance [Facchinei and Pang, 2007]. Monotone VIs have been studied since the 1960s [Hartman and Stampacchia, 1966, Browder, 1965, Lions and Stampacchia, 1967, Brezis and Sibony, 1968, Sibony, 1970]. Formally, a monotone VI is specified by a closed convex set $\mathcal{Z} \subseteq \mathbb{R}^n$ and a **monotone** operator $F: \mathcal{Z} \to \mathbb{R}^n$, with the goal of finding a $z^* \in \mathcal{Z}$ such that

$$\langle F(z^*), z^* - z \rangle \le 0 \quad \forall z \in \mathcal{Z}.$$
 (1)

We further assume the operator *F* to be **Lipschitz**, which is a natural assumption that is satisfied in most applications and is also made in the majority of algorithmic works concerning monotone VIs. An important special case of the monotone and Lipschitz VI is the convex-concave saddle point problem:

$$\min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} f(x, y), \tag{2}$$

where \mathcal{X} and \mathcal{Y} are closed convex sets in \mathbb{R}^n , and $f(\cdot,\cdot)$ is smooth, convex in x, and concave in y. Besides its central importance in Game Theory, Convex Optimization, and Online Learning, the convex-concave saddle point problem has recently received a lot of attention from the machine learning community due to several novel applications such as the generative adversarial networks (GANS) (e.g., [Goodfellow et al., 2014, Arjovsky et al., 2017]), adversarial examples (e.g., [Madry et al., 2018]), robust optimization (e.g., [Ben-Tal et al., 2009]), and reinforcement learning (e.g., [Du et al., 2017, Dai et al., 2018]).

The extragradient method (EG) by Korpelevich [1976] is one of the most popular methods for solving monotone and Lipschitz VIs. Despite the extensive effort from the optimization and machine learning community, the following major problem remains open.

What is the last-iterate convergence rate of the extragradient method for monotone
$$(\star)$$
 and Lipschitz variational inequalities with constraints?

Indeed, the same problem has not been answered even for two-player zero-sum games, arguably one of the most basic monotone and Lipschitz VIs.³ Korpelevich [1976] and Facchinei and Pang [2007] prove that the last iterate of EG converges to a solution of the monotone and Lipschitz VI, but do not provide any upper bound on the rate of convergence. Prior to our work, the only setting where an upper bound on the rate of convergence exists is when the problem is unconstrained, i.e., $\mathcal{Z} = \mathbb{R}^n$ [Golowich et al., 2020b, Gorbunov et al., 2021]. Determining the

¹*F* is monotone if $\langle F(z) - F(z'), z - z' \rangle \ge 0$ for all $z, z' \in \mathcal{Z}$.

²If we set $F(x,y) = \begin{pmatrix} \nabla_x f(x,y) \\ -\nabla_y f(x,y) \end{pmatrix}$ and $\mathcal{Z} = \mathcal{X} \times \mathcal{Y}$, then (i) F(x,y) is a monotone and Lipschitz operator, and (ii) the set of saddle points coincide with the solutions of the monotone VI for operator F and domain \mathcal{Z} .

³A two-player zero-sum game can be specified by its payoff matrix $A \in \mathbb{R}^{\ell \times m}$. It is a special case of the convex-concave saddle point problem, where $\mathcal{X} = \Delta^{\ell}$, $\mathcal{Y} = \Delta^{m}$ (Δ^{k} denotes the k-dimensional simplex), and the function $f(x,y) = x^{\top}Ay$.

last-iterate convergence rate of EG in the constrained setting has been posed as an open question in several recent works [Wei et al., 2021, Gorbunov et al., 2021, Golowich et al., 2020a, Hsieh et al., 2019]. We resolve this open problem by providing the tight last-iterate convergence rate of EG under arbitrary constraints.

Main Result: For any monotone and Lipschitz variational inequality problem with an **arbitrary convex constraint set** \mathcal{Z} , EG with **constant step size** achieves a last-iterate convergence of $O\left(\frac{1}{\sqrt{T}}\right)$ in terms of the standard convergence measures – the gap function (Definition 1) and the tangent residual (Definition 3).^a See Theorem 3 for the formal statement.

Our upper bound matches the lower bound of Golowich et al. [2020b] in a very strong sense, that is, they match in all of the following terms: T, the Lipschitz constant of F, and the distance between the starting point z_0 and the solution z^* .

The optimistic gradient descent-ascent (OGDA) algorithm by Popov [1980] is another widely used method for solving monotone VIs. The state of OGDA's last-iterate convergence guarantee is similar to EG's (prior to our result), that is, both the asymptotic last-iterate convergence [Popov, 1980, Hsieh et al., 2019] and the explicit convergence rate in the unconstrained setting (under a second-order smoothness assumption) are known [Golowich et al., 2020a], but no explicit convergence rate is known for the constrained setting, e.g., two-player zero-sum games. Indeed, to the best of our knowledge, our result is the first to provide a last-iterate convergence rate for solving monotone and Lipschitz VIs using any algorithm that belongs to the general class known as *p*-stationary canonical linear iterative algorithms (*p*-SCLI) [Arjevani and Shamir, 2016], which contains the EG, OGDA, and other well-known algorithms.

Why Last-Iterate Convergence? Both the EG and OGDA algorithms are known to have *average*iterate convergence. In particular, the average of the iterates of the algorithm converges at a rate of O(1/T) [Nemirovski, 2004, Auslender and Teboulle, 2005, Tseng, 2008, Monteiro and Svaiter, 2010, Mokhtari et al., 2020, Hsieh et al., 2019]. Nonetheless, there are several important reasons to study last-iterate convergence. First, not only is last-iterate convergence theoretically stronger and more appealing, it is also the only type of convergence that describes the trajectory of an algorithm. As demonstrated by Mertikopoulos et al. [2018], the trajectory of an algorithm may be cycling around in the space perpetually while still converges in the average-iterate sense. In game theory, we often view these algorithms as models of agents' behavior in a system/game. Thus, only lastiterate convergence provides a description of the evolution of the system. Additionally, EG and OGDA have been successfully applied to improve the training dynamics in GANs, as the training of GANs can be formulated as a saddle point problem [Daskalakis et al., 2018, Yadav et al., 2018, Liang and Stokes, 2019, Gidel et al., 2019a,b, Chavdarova et al., 2019]. On the one hand, in this formulation of GANs, the objective function f is usually non-convex and non-concave, making existing theoretical guarantees for the average iterate inapplicable. On the other hand, the last iterate typically has good performance in practice. Thus, it is crucial to develop machinery that

^aIn the unconstrained setting, the tangent residual is simply the ℓ_2 -norm of F. In the constrained setting, the tangent residual is the ℓ_2 -norm of F's projection to the tangent cone.

allows us to analyze the behavior of the last iterate of these algorithms.

1.1 A New Performance Measure: the Tangent Residual

A major challenge we face for establishing the last-iterate convergence for EG in the constrained setting is the choice of the convergence measure. In the unconstrained case, the central performance measure is the norm of the operator. The key component in both [Golowich et al., 2020b] and [Gorbunov et al., 2021] is to establish that the norm of the operator at the last iterate (also the T-th iterate) is upper bounded by $O(\frac{1}{\sqrt{T}})$, which implies a $O(\frac{1}{\sqrt{T}})$ last-iterate convergence rate for the gap function.

In the constrained setting, the norm of the operator is a poor choice to measure convergence, as it can be far away from 0 even in the limit, and is hence insufficient to guarantee convergence in terms of the gap function. A standard generalization of the norm of the operator in the constrained setting is the natural residual (Definition 4), which takes the constraints into account and is guaranteed to converge to 0 in the limit. Unfortunately, we observe that the natural residual is *not monotonically decreasing* even in basic bilinear games (see Appendix G), making it difficult to directly analyze. Similar non-monotonicity has been observed for several other natural performance measures such as the norm of the operator mapping introduced in [Diakonikolas, 2020] and the gap function, leaving all these performance measures unsuitable. See more discussion about these performance measures in Section 3 and Appendix G.

The first main contribution of this paper is the introduction of a new performance measure: the **tangent residual**, which can be viewed as the norm of the operator projected to the tangent cone of the current iterate (Definition 3). We believe that this new performance measure will be useful for the analysis of other iterative methods. Indeed, the tangent residual plays a crucial role in our analysis. Unlike the aforementioned performance measures, we show that the *tangent residual* is monotonically decreasing and has a last-iterate convergence rate of $O(\frac{1}{\sqrt{T}})$. Using the convergence rate of the tangent residual, we can easily derive the last-iterate convergence rate of other classical performance measures such as the natural residual or the gap function. However, we suspect these rates can be challenging to obtain directly without understanding the tangent residual first.

1.2 Sum-of-Squares based Analysis

To obtain our main result, we propose a novel approach to analyze iterative methods using the *Sum-of-Squares (SOS)* programming [Nesterov, 2000, Parrilo, 2000, 2003, Lasserre, 2001, Laurent, 2009]. In the past two decades, the SOS programming has found numerous applications in fields such as combinatorial optimization, control theory, quantum mechanics, etc. To the best of our knowledge, this is the first application of SOS programming in the analysis of iterative methods for solving VIs.

Very often, the convergence analysis of an iterative method builds on a potential function argument, which boils down to demonstrating the non-negativity of certain function of the iterates. This is usually done by chaining together a sequence of inequalities that are satisfied due to the update rule of the iterative method. Although easy to verify, such a proof could be extremely

challenging to discover, and is often found by much trial and error. We propose to use the SOS programming to significantly automate this process.

Sum-of-Squares (SOS) Programming. Suppose we want to prove that a polynomial $p(x) \in \mathbb{R}[x_1,\ldots,x_n]$ is non-negative over a semialgebraic set $\mathcal{S}=\{x\in\mathbb{R}^n:g_i(x)\leq 0, \forall i\in[m]\}$, where each $g_i(x)$ is also a polynomial. One way is to construct a *certificate of non-negativity*, for example, by providing a set of nonnegative coefficients $\{a_i\}_{i\in[m]}\in\mathbb{R}^m_{\geq 0}$ such that $p(x)+\sum_{i\in[m]}a_i\cdot g_i(x)$ is a sum-of-squares polynomial, that is, a polynomial that can be expressed as the sum of squares of further polynomials. Surprisingly, if p(x) is indeed non-negative over \mathcal{S} , a certificate of non-negativity always exists as guaranteed by a foundational result in real algebraic geometry – the Krivine-Stengle Positivestellensatz [Krivine, 1964, Stengle, 1974], a generalization of Artin's resolution of Hilbert's 17th problem [Artin, 1927]. Although the certificate in the example above is relatively simple, it is necessary to allow more sophisticated forms in general, e.g., replacing each coefficient a_i with a SOS polynomial $s_i(x)$, etc. The complexity of a certificate is parametrized by the highest degree of the polynomial involved. Such added complexity is useful and indeed needed in our analysis for the EG algorithm. The SOS programming consists of a hierarchy of algorithms, where the d-th hierarchy is an algorithm that searches for a *certificate of non-negativity* up to degree 2d based on semidefinite programming.

At the core of our analysis of the EG algorithm lies the monotonicity of the tangent residual, which can be formulated as the non-negativity of a **degree-6 polynomial** in the iterates. However, we cannot directly apply the SOS programming to search for such a certificate due to the following subtlety. As the EG algorithm applies to any dimension n, we need to prove the monotonicity of the corresponding polynomial for every positive integer n. Thus we must certify the non-negativity for an infinite family of polynomials, one for every n. We circumvent this difficulty by **exploiting the low-dimensionality of the EG update rule**, i.e., it only involves a constant number of points and projections per update. In particular, we manage to show that *the number of relevant dimensions in each update* is only a constant regardless of the dimension of the ambient space and the domain \mathbb{Z} . Taking advantage of this property, we show that it suffices to prove the non-negativity of a single polynomial, which we then use SOS programming to certify. The certificate is rather complex and involves a polynomial identity of a degree-8 polynomial in 27 variables, which we discover by solving a degree-8 SOS program. We think such a proof will be extremely difficult to discover using the traditional approach, if not impossible.

We believe our approach that combines the power of SOS programming with the low-dimensionality of the iterative update rule has many additional applications. Indeed, we can use it to recover the $O\left(\frac{1}{\sqrt{T}}\right)$ last-iterate convergence rate for the proximal point (PP) algo-

⁴The tangent residual is not a polynomial, but the squared tangent residual is a rational function, i.e., quotient of two polynomials. To establish the monotonicity of the tangent residual, we only need to prove that the difference between the squared tangent residual in two consecutive iterates is non-negative. Although, the difference is still only a rational function, proving its non-negativity turns out to be the same as proving the non-negativity of its numerator as the denominator is an SOS polynomial.

⁵Note that which constant number of dimensions are relevant change from iterate to iterate, so all n dimensions might be relevant if we look at the complete trajectory of the algorithm.

rithm [Bauschke et al., 2011, Cominetti et al., 2013].

1.3 Related Work

Last-Iterate Convergence Rate for EG-like Algorithms in the Unconstrained Setting. Golowich et al. [2020b,a] show a lower bound of $\Omega(\frac{1}{\sqrt{T}})$ for solving bilinear games using any p-SCLI algorithms, which include EG and OGDA. Matching upper bounds are provided by Golowich et al. [2020b], Gorbunov et al. [2021] for EG and by Golowich et al. [2020a] for OGDA. These upper bounds hold for all smooth and Lipschitz VIs. With the additional assumption that the operator F is cocoercive, Lin et al. [2020] show a $O(\frac{1}{\sqrt{T}})$ convergence rate for online gradient descent. If we further assume that either F is strongly monotone in VI or the payoff matrix A in a bilinear game has all singular values bounded away from 0, linear convergence rate is known for EG, OGDA, and several of their variants [Daskalakis et al., 2018, Gidel et al., 2019a, Liang and Stokes, 2019, Mokhtari et al., 2020, Peng et al., 2020, Zhang and Yu, 2020].

Last-Iterate Convergence Rate for EG-like Algorithms in the Constrained Setting. for the constrained setting are sparser. If the operator F is strongly monotone, we know that EG and some of its variants have linear convergence rate [Tseng, 1995, Malitsky, 2015]. Several papers establish the asymptotic convergence, i.e., converge in the limit, of the optimistic multiplicative weight updates in constrained convex-concave saddle point problems [Daskalakis and Panageas, 2019, Lei et al., 2021]. Finally, a recent paper by Wei et al. [2021] provides a linear rate convergence of OGDA for bilinear games when the domain is a polytope. They show that there is a problem dependent constant 0 < c < 1 that depends on the payoff matrix of the game as well as the constraint set, so that the error shrinks by a 1-c factor. However, c may be arbitrarily close to 0, even if we assume the corresponding operator to be L-Lipschitz. As a result, their convergence rate is slower than ours when T is not comparable to $\frac{1}{c}$, which may be exponentially large in the dimension n, though their rate will eventually catch up. Overall, their "instance-specific" bound is incomparable and complements the worst-case view taken in this paper, where we want to derive the worstcase convergence rate for all VIs with monotone and L-Lipschitz operator F. Our result is the first last-iterate convergence rate in this worst-case view and matches the lower bound by Golowich et al. [2020b].

Other Algorithms and Performance Measures. Other than the gap function, one can also measure the convergence using the norm of the operator if the setting is unconstrained, or the natural residual (Definition 4) or similar notions if the setting is constrained. In the unconstrained setting, Kim [2021] and Yoon and Ryu [2021] provide algorithms that obtain $O(\frac{1}{T})$ convergence rate in terms of the norm of the operator, which is shown to be optimal by Yoon and Ryu [2021] for Lipschitz and monotone VIs. In the constrained setting, Diakonikolas [2020] shows the same $O(\frac{1}{T})$ convergence rate under the extra assumption that the operator is cocoercive and loses an additional logarithmic factor when the operator is only monotone. Our result implies a $O(\frac{1}{\sqrt{T}})$ last-iterate convergence rate in terms of the natural residual. Although these algorithms achieve

faster rate than EG,⁶ they are based on different methods, i.e., accelerated proximal point or anchoring. As EG is arguably the most classical and one of the most popular methods for solving variational inequalities, we believe it is crucial to understand its last-iterate convergence rate.

Computer-Aided Proofs. A powerful computer-aided proof framework – the performance estimation problem (PEP) technique (e.g., [Drori and Teboulle, 2014, Taylor et al., 2017b]) is widely applied to analyze first-order iterative methods. Indeed, the last-iterate convergence rate of EG in the unconstrained setting by Gorbunov et al. [2021] is obtained via the PEP technique. While PEP can handle projections [Taylor et al., 2017a, Ryu et al., 2020, Goujaud et al., 2022, Dragomir et al., 2021], we believe that the PEP technique is insufficient to obtain our result. Our concern is that the PEP framework can only handle performance measures that are polynomials of degree 2 or less (see e.g., [Taylor et al., 2017a]). In fact, solving the PEP is equivalent to solving a degree-2 SOS program, which can be viewed as the dual of the PEP. In the unconstrained setting, the performance measure is a degree-2 polynomial – the squared norm of the operator, and that is why one can either use the PEP (as in [Gorbunov et al., 2021]) or a degree-2 SOS to certify its monotonicity (Theorem 1). In the constrained setting, we use the squared tangent residual to measure the algorithm's progress, which is a rational function with a degree 4 polynomial numerator and a quadratic denominator, making it unsuitable for the PEP framework. We were unable to derive a reformulation of our performance measure so that it is amenable to PEP and such a reformulation would be an interesting question for future work. Compared to PEP, our SOS approach is more flexible and can accommodate polynomial objectives and constraints of any degree, which allows us to directly apply it to certify the monotonicity of the tangent residual in the constrained setting.

Lessard et al. [2016] analyze first-order iterative algorithms for convex optimization using a technique inspired by the stability analysis from control theory. They model first-order iterative algorithms using discrete-time dynamical systems and search over quadratic potential functions that satisfy a set of Integral Quadratic Constraints (IQC). While the IQC is suitable for obtaining linear convergence rates of algorithms for smooth strongly convex optimization, it fails to recover the exact sublinear rates in the non-strongly convex case. Zhang et al. [2021] extend the IQC framework to study smooth and *strongly* monotone VIs in the unconstrained setting. However, since the IQC framework only concerns quadratic potential functions, it is unclear how to encode our potential function and constraints into the IQC framework.

SOS programming and Analysis of Iterative Methods. SOS programming has been employed in the design and analysis of algorithms in convex optimization. To the best of our knowledge, these results only concern minimization of smooth and strongly-convex functions in the unconstrained setting. Fazlyab et al. [2018] propose a framework to search the optimal parameters of

⁶In terms of the gap function, these algorithms have the same or slower convergence rate as the average-iterate of EG, i.e., $O(\frac{1}{T})$.

⁷More specifically, the PEP framework requires the performance measure as well as the constraints to be linear in (i) the function values at the iterates and (ii) the Gram matrix of a set of vectors consisting of the iterates and their gradients.

⁸The tangent residual is the square root of a rational function and can only be even harder to handle.

the algorithm, e.g., step size. They use SOS programming to search over quadratic potential functions and parameters of the algorithm with the goal of optimizing the exponential decay rate of the potential function. Tan et al. [2021] proposes to apply SOS programming to study the convergence rates of first-order methods. Their motivation is similar to ours, that is, it may be beneficial or even necessary to analyze performance measures that cannot be handled by the PEP framework, e.g., non-quadratic functions. Despite the motivation, their analyses only use degree-2 SOS programs, which can thus be derived within the PEP framework. Moreover, their approach is dimension-dependent for higher degree SOS proofs, that is, for every n, they need to run a separate SOS program for the dimension-n version of the problem. Compared to the literature, we are the first to analyse a high degree polynomial performance measure, and the first to apply SOS programming to analyze variational inequalities in both unconstrained and constrained settings. Moreover, unlike [Tan et al., 2021], our proof is dimension-free. We hope our result demonstrates the flexibility and power of the SOS programming based approach, and believe that there are many further applications.

2 Preliminaries

We consider the Euclidean Space $(\mathbb{R}^n, \|\cdot\|)$, where $\|\cdot\|$ is the ℓ_2 norm and $\langle\cdot,\cdot\rangle$ denotes inner product on \mathbb{R}^n . We use z[i] to denote the i-th coordinate of $z \in \mathbb{R}^n$ and e_i to denote the unit vector such that $e_i[j] := \mathbb{1}[i=j]$, the dimension of e_i is going to be clear from context.

Variational Inequality. Given a closed convex set $\mathcal{Z} \subseteq \mathbb{R}^n$ and an operator $F: \mathcal{Z} \to \mathbb{R}^n$, a variational inequality problem is defined as follows: find $z^* \in \mathcal{Z}$ such that

$$\langle F(z^*), z^* - z \rangle \le 0 \quad \forall z \in \mathcal{Z}.$$
 (3)

We say F is **monotone** if $\langle F(z) - F(z'), z - z' \rangle \ge 0$, for all $z, z' \in \mathcal{Z}$, and is L-Lipschitz if, $||F(z) - F(z')|| \le L||z - z'||$ for all $z, z' \in \mathcal{Z}$.

Remark 1. One sufficient condition for such a z^* to exist is when the set \mathcal{Z} is bounded, but there are also other sufficient conditions that apply to unbounded \mathcal{Z} . See [Facchinei and Pang, 2007] for more details. Throughout this paper, we only consider monotone VIs that have a solution.

Definition 1 (Gap Function). A standard way to measure the performance of $z \in \mathcal{Z}$ is by its gap function defined as $GAP_{\mathcal{Z},F,D}(z) = \max_{z' \in \mathcal{Z} \cap \mathcal{B}(z,D)} \langle F(z), z - z' \rangle$, where D > 0 is a fixed parameter and $\mathcal{B}(z,D)$ is a ball with radius D centered at z. When \mathcal{Z} , F and D are clear from context, we omit the subscripts and write the gap function at z as GAP(z).

⁹Sometimes the gap function is defined to allow z' to take value in $\mathcal{Z} \cap \mathcal{B}(z^*, \|z^* - z_0\|)$, where z_0 is the starting point of the EG algorithm, and z^* is the solution that the last iterate of the algorithm converges to. Due to Lemma 3, $\|z_k - z^*\| \le \|z_0 - z^*\|$ for every k, so $\mathcal{B}(z_k, 2\|z^* - z_0\|)$ contains $\mathcal{B}(z^*, \|z^* - z_0\|)$.

The Extragradient Method. Let z_k be the k-th iterate of the Extragradient (EG) method. The update rule of EG is as follows:

$$z_{k+\frac{1}{2}} = \Pi_{\mathcal{Z}} \left[z_k - \eta F(z_k) \right] = \arg \min_{z \in \mathcal{Z}} \| z - (z_k - \eta F(z_k)) \|, \tag{4}$$

$$z_{k+1} = \Pi_{\mathcal{Z}}\left[z_k - \eta F(z_{k+\frac{1}{2}})\right] = \arg\min_{z \in \mathcal{Z}} \left\|z - \left(z_k - \eta F(z_{k+\frac{1}{2}})\right)\right\|. \tag{5}$$

Sum-of-Squares (SOS) Polynomials. Let x be a set of variables. We denote the set of real polynomials in x as $\mathbb{R}[x]$. We say that polynomial $p(x) \in \mathbb{R}[x]$ is an SOS polynomial if there exist polynomials $\{q_i(x) \in \mathbb{R}[x]\}_{i \in [M]}$ such that $p(x) = \sum_{i \in [M]} q_i(x)^2$. We denote the set of SOS polynomials in x as SOS[x]. Note that any SOS polynomial is non-negative.

SOS Programs. In Figure 1 we present a generic formulation of a degree-d SOS program. The SOS program takes three kinds of input, a polynomial g(x), sets of polynomials $\{g_i(x)\}_{i\in[M]}$ and $\{h_i(x)\}_{i\in[N]}$. Each polynomial in $\{g(x)\}\cup\{g_i(x)\}_{i\in[M]}\cup\{h_i(x)\}_{i\in[N]}$ has degree of at most d. The SOS program searches for an SOS polynomial in the set of polynomials $\Sigma=\{g(x)+\sum_{i\in[M]}p_i(x)\cdot g_i(x)+\sum_{i\in[N]}q_i(x)\cdot h_i(x)\}$, where $\{p_i(x)\}_{i\in[M]}$ and $\{q_i(x)\}_{i\in[N]}$ are polynomials in x. More precisely for each $i\in[M]$, $p_i(x)$ is an SOS polynomial with degree at most $d-\deg(g_i(x))$. For each $i\in[N]$, $q_i(x)$ is a (not necessarily SOS) polynomial with degree at most $d-\deg(g_i(x))$. Note that any polynomial in set Σ is at most degree d. In our applications, we choose $\{g_i(x)\}_{i\in[M]}$ to be nonpositive polynomials and $\{h_i(x)\}_{i\in[N]}$ to be polynomials that are equal to 0. Any feasible solution to the program certifies the non-negativity of g(x).

Input Fixed Polynomials.

- Polynomial g(x)
- Polynomial $g_i(x) \in \mathbb{R}[x]$ for all $i \in [M]$.
- Polynomial $h_i(x) \in \mathbb{R}[x]$ for all $i \in [N]$.

Decision Variables of the SOS Program:

- $p_i(x) \in SOS[x]$ is an SOS polynomial with degree at most $d deg(g_i)$, for all $i \in [M]$.
- $q_i(x) \in \mathbb{R}[x]$ is a polynomial with degree at most $d \deg(h_i)$, for all $i \in [N]$.

Constraints of the SOS Program:

$$g(x) + \sum_{i \in [M]} p_i(x) \cdot g_i(x) + \sum_{i \in [N]} q_i(x) \cdot h_i(x) \in SOS[x]$$

Figure 1: Generic degree *d* SOS program.

3 Our Performance Measure: the Tangent Residual

We formally introduce our performance measure the tangent residual . As discussed in Section 1.1, many standard and natural performance measures, i.e., the natural residual, $||z_k - z_{k+1/2}||$, $||z_k - z_{k+1}||$, $\max_{z \in \mathcal{Z}} \langle F(z), z_k - z \rangle$ and $\max_{z \in \mathcal{Z}} \langle F(z_k), z_k - z \rangle$, are unfortunately non-decreasing for EG. See Appendix G for numerical examples.

Definition 2 (Unit Normal Cone). Given a closed convex set $\mathcal{Z} \subseteq \mathbb{R}^n$ and a point $z \in \mathcal{Z}$, we denote by $N_{\mathcal{Z}}(z) = \langle v \in \mathbb{R}^n : \langle v, z' - z \rangle \leq 0, \forall z' \in \mathcal{Z} \rangle$ the normal cone of \mathcal{Z} at point z and by $\widehat{N}_{\mathcal{Z}}(z) = \{v \in N_{\mathcal{Z}}(z) : ||v|| \leq 1\}$ the intersection of the unit ball with the the normal cone of \mathcal{Z} at z. Note that $\widehat{N}_{\mathcal{Z}}(z)$ is nonempty and compact for any $z \in \mathcal{Z}$, as $(0, \ldots, 0) \in \widehat{N}_{\mathcal{Z}}(z)$.

Definition 3 (Tangent Residual). Given an operator $F: \mathcal{Z} \to \mathbb{R}^n$ and a closed convex set \mathcal{Z} , let $T_{\mathcal{Z}}(z) := \{z' \in \mathbb{R}^n : \langle z', a \rangle \leq 0, \forall a \in \mathcal{C}_{\mathcal{Z}}(z)\}$ be the tangent cone of z, and define $J_{\mathcal{Z}}(z) := \{z\} + T_{\mathcal{Z}}(z)$. The tangent residual of F at $z \in \mathcal{Z}$ is defined as $r_{(F,\mathcal{Z})}^{tan}(z) := \|\Pi_{J_{\mathcal{Z}}(z)}[z - F(z)] - z\|$. An equivalent definition is $r_{(F,\mathcal{Z})}^{tan}(z) := \sqrt{\|F(z)\|^2 - \max_{a \in \widehat{N}_{\mathcal{Z}}(z), \langle a, F(z) \rangle^2}}$.

Remark 2. We show the equivalence of the two definitions of tangent residual in Lemma 7. For the rest of the paper, we may use either of the two equivalent definitions depending on which one is more convenient.

When the convex set \mathcal{Z} and the operator F are clear from context, we are going to omit the subscript and denote the unit normal cone as $\widehat{N}(z) = \widehat{N}_{\mathcal{Z}}(z)$ and the tangent residual as $r^{tan}(z) = r^{tan}_{(F,\mathcal{Z})}(z)$. Although the definition is slightly technical, one can think of the tangent residual as the norm of another operator \widehat{F} , which is F projected to all directions that are not "blocked" by the boundary of \mathcal{Z} if one takes an infinitesimally small step $\epsilon \cdot F(z)$, which is the same as projecting F to $J_{\mathcal{Z}}(z)$. Intuitively, if the tangent residual is small, then the next iterate will not be far away from the current one.

Next, we formally define the natural residual associated with the instance formally stated in Definition 4, and show how it is related to the tangent residual.

Definition 4. Consider an instance \mathcal{I} of the variational inequality problem on convex set $\mathcal{Z} \subseteq \mathbb{R}^n$ and monotone operator $F: \mathcal{Z} \to \mathbb{R}^n$. For $z \in \mathcal{Z}$, the natural map and natural residual associated with \mathcal{I} is defined as follows

$$F_K^{nat}(z) = z - \Pi_{\mathcal{Z}}(z - F(z)), \qquad r_{(F,\mathcal{Z})}^{nat}(z) = \|F_K^{nat}(z)\|.$$

Given an instance of the monotone VI constrained on convex set $\mathcal{Z} \subseteq \mathbb{R}^n$ and operator $F: \mathcal{Z} \to \mathbb{R}^n$, point z^* is a solution of the monotone VI iff $r_{(F,\mathcal{Z})}^{nat}(z^*) = 0$. In Lemma 1, we show that the tangent residual upper bounds the the natural residual. See Figure 2 for illustration of how the tangent residual relates to the natural residual.

Lemma 1. Consider an instance \mathcal{I} of the variational inequality problem on convex set $\mathcal{Z} \subseteq \mathbb{R}^n$ and monotone operator $F: \mathcal{Z} \to \mathbb{R}^n$. For any $z \in \mathcal{Z}$, $r_{(F,\mathcal{Z})}^{tan} \geq r_{(F,\mathcal{Z})}^{nat}(z)$.

 $^{^{10}}$ $\|z_k - z_{k+1/2}\|$ is proportional to the norm of the operator mapping introduced in [Diakonikolas, 2020].

Proof. Let $w = \Pi_{\mathcal{Z}}(z - F(z))$ and $a_1 = z - F(z) - w$. Observe that

$$||F(z)||^2 = ||z - w||^2 + ||a_1||^2 - 2\langle z - w, a_1 \rangle.$$

Since $r_{(F,Z)}^{nat}(z)^2 = \|z - w\|^2$ and $\langle z - w, a_1 \rangle \leq 0$, we have $r_{(F,Z)}^{nat}(z)^2 \leq \|F(z)\|^2 - \|a_1\|^2$.

According to Lemma 7, $r^{tan}(z)^2 = \|\Pi_{J_{\mathcal{Z}}(z)}[z - F(z)] - z\|^2$, where $J_{\mathcal{Z}}(z) := z + T_{\mathcal{Z}}(z)$ and $T_{\mathcal{Z}}(z) = \{z' \in \mathbb{R}^n : \langle z', a \rangle \leq 0, \forall a \in \mathcal{C}_{\mathcal{Z}}(z)\}$ is the tangent cone of z. Since $J_{\mathcal{Z}}(z)$ is a cone with origin z, we have $\langle z - F(z) - \Pi_{J_{\mathcal{Z}}(z)}[z - F(z)], z - \Pi_{J_{\mathcal{Z}}(z)}[z - F(z)] \rangle = 0$, and $\|\Pi_{J_{\mathcal{Z}}(z)}[z - F(z)] - z\|^2 = \|F(z)\|^2 - \|z - F(z) - \Pi_{J_{\mathcal{Z}}(z)}[z - F(z)]\|^2$. As $\mathcal{Z} \subseteq J_{\mathcal{Z}}(z)$, $\|a_1\|^2 \geq \|z - F(z) - \Pi_{J_{\mathcal{Z}}(z)}[z - F(z)]\|^2$, which implies that

$$r_{(F,\mathcal{Z})}^{nat}(z)^2 \le \left\| \Pi_{J_{\mathcal{Z}}(z)} \left[z - F(z) \right] - z \right\|^2 = r^{tan}(z)^2.$$

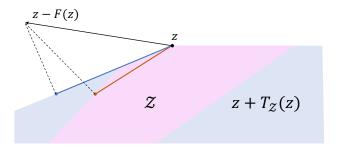


Figure 2: Illustration of the tangent residual and the natural residual. The blue line represents the tangent residual and the red line represents the natural residual. It is clear that the tangent residual upper bounds the natural residual.

Due to the above lemma, an upper bound of the tangent residual is also an upper bound of the natural residual. We show in Theorem 2 the monotonicity of the tangent residual of the EG updates, which is the technical core of our analysis and implies the $O(\frac{1}{\sqrt{T}})$ convergence rate of the tangent residual. As a result, we also show that the natural residual has a $O(\frac{1}{\sqrt{T}})$ convergence rate. One may be tempted to directly use the natural residual as the convergence measure. However, from our numerical experiments, the natural residual of the EG updates is not monotone, and we believe that it is very challenging to directly establish the convergence rate for the natural residue without using the tangent residual as a proxy.

In the next lemma, we argue why a small tangent residual implies a small gap function, hence an approximate solution of the variational inequality. The proof is postponed to Appendix A.

Lemma 2. [Adapted from the proof of Theorem 10 in [Golowich et al., 2020b].] Given a closed convex set $\mathcal{Z} \in \mathbb{R}^n$, an operator $F: \mathcal{Z} \to \mathbb{R}^n$ and $z \in \mathcal{Z}$, we have

$$\mathrm{Gap}_{\mathcal{Z},F,D}(z) := \max_{z' \in \mathcal{Z} \cap \mathcal{B}(z,D)} \left\langle F(z), z - z' \right\rangle \leq D \cdot r_{(F,\mathcal{Z})}^{tan}(z).$$

If we have a convex-concave function $f(z): \mathcal{Z} \to \mathbb{R}$ such that z = (x,y), $\mathcal{Z} = \mathcal{X} \times \mathcal{Y}$ where \mathcal{X} and \mathcal{Y} are closed convex sets, let $F(x,y) = \begin{pmatrix} \nabla_x f(x,y) \\ -\nabla_y f(x,y) \end{pmatrix}$, then the duality gap at z with respect to \mathcal{X}' and \mathcal{Y}' is $\deg_f^{\mathcal{X}',\mathcal{Y}'}(z) := \max_{y' \in \mathcal{Y}'} f(x,y') - \min_{x' \in \mathcal{X}'} f(x',y) \leq D\sqrt{2} \cdot r_{(F,\mathcal{Z})}^{tan}(z)$, if $z \in \mathcal{X}' \times \mathcal{Y}'$ and the diameters of \mathcal{X}' and \mathcal{Y}' are both upper bounded by $D.^{11}$

4 Best-Iterate Convergence of EG with Constant Step Size

En route to establish the last-iterate convergence of the EG algorithm, we first show a weaker guarantee known as the best-iterate convergence. Lemma 3 implies that after running EG for T steps, there exists an iteration $t^* \in [T]$ where $||z_{t^*} - z_{t^* + \frac{1}{2}}||^2 \le O(\frac{1}{T})$. The proof can be found in [Korpelevich, 1976] and [Facchinei and Pang, 2007] (included in Appendix B for completeness).

Lemma 3 ([Korpelevich, 1976, Facchinei and Pang, 2007]). Let \mathcal{Z} be a closed convex set in \mathbb{R}^n , $F(\cdot)$ be a monotone and L-Lipschitz operator mapping from \mathcal{Z} to \mathbb{R}^n . For any solution z^* of the monotone VI, that is, $\langle F(z^*), z^* - z \rangle \leq 0$ for all $z \in \mathcal{Z}$. For all k,

$$||z_k - z^*||^2 \ge ||z_{k+1} - z^*||^2 + (1 - \eta^2 L^2)||z_k - z_{k+\frac{1}{2}}||^2.$$
(6)

In Lemma 4, we relate $||z_k - z_{k+\frac{1}{2}}||$ with the tangent residual at z_{k+1} , and derive the best-iterate convergence guarantee in terms of the tangent residual in Lemma 5. The proofs of Lemma 4 and 5 are postponed to Appendix B.

Lemma 4. For all
$$k$$
, $r^{tan}(z_{k+1}) \leq (1 + \eta L + (\eta L)^2) \frac{||z_k - z_{k+1/2}||}{\eta}$.

In lemma 5, we argue the tangent residual has a best-iterate convergence with rate $O(\frac{1}{\sqrt{T}})$.

Lemma 5. Let \mathcal{Z} be a closed convex set in \mathbb{R}^n , $F(\cdot)$ be a monotone and L-Lipschitz operator mapping from \mathcal{Z} to \mathbb{R}^n . Suppose the step size of the EG algorithm $\eta \in (0, \frac{1}{L})$, then for any solution z^* of the monotone VI and any integer T > 0, there exists $t^* \in [T]$ such that:

$$\left\|z_{t^*} - z_{t^* + \frac{1}{2}}\right\|^2 \le \frac{1}{T} \frac{\|z_0 - z^*\|^2}{1 - (\eta L)^2}, \quad \text{AND} \quad r^{tan}(z_{t^* + 1}) \le \frac{1 + \eta L + (\eta L)^2}{\eta} \frac{1}{\sqrt{T}} \frac{\|z_0 - z^*\|}{\sqrt{1 - (\eta L)^2}}.$$

¹¹When \mathcal{X} and \mathcal{Y} are bounded, we choose $\mathcal{X}' = \mathcal{X}$ and $\mathcal{Y}' = Y$, otherwise the convention is to choose \mathcal{X}' and \mathcal{Y}' to be $\mathcal{X} \cap \mathcal{B}(x^*, 2\|x_0 - x^*\|)$ and $\mathcal{Y} \cap \mathcal{B}(y^*, 2\|y_0 - y^*\|)$ respectively, where (x^*, y^*) is a saddle point.

5 Last-Iterate Convergence of EG with Constant Step Size

In this section, we show that the last-iterate convergence rate is $O(\frac{1}{\sqrt{T}})$. In particular, we prove that the tangent residual is non-increasing, which, in combination with Lemma 5, implies the last-iterate convergence rate of EG. To establish the monotonicity of the tangent residual, we combine SOS programming with the low-dimensionality of the EG update rule. To better illustrate our approach, we first prove the result in the unconstrained setting (Section 5.1), then show how to generalize it to the constrained setting (Section 5.2).

5.1 Warm Up: Unconstrained Case

As a warm-up, we consider the unconstrained setting where $\mathcal{Z} = \mathbb{R}^n$. Although the last-iterate convergence rate is known in the unconstrained setting due to Golowich et al. [2020b] and Gorbunov et al. [2021], we provide a simpler proof, which also permits a larger step size. Our analysis holds for any step size $\eta \in (0, \frac{1}{L})$, while the previous analysis requires $\eta \leq \frac{1}{\sqrt{2L}}$ [Gorbunov et al., 2021].

Let z_k be the k-th iterate of the EG method. In Theorem 1, we show that the tangent residual is monotone in the unconstrained setting.¹² Our approach is to apply SOS programming to search for a certificate of non-negativity for $\|F(z_k)\|^2 - \|F(z_{k+1})\|^2$ for every k, over the semialgebraic set defined by the following polynomial constraints in variables $\{z_i[\ell], \eta F(z_i)[\ell]\}_{i \in \{k,k+\frac{1}{2},k+1\}, \ell \in [n]}$:

$$\begin{split} &z_{k+\frac{1}{2}}[\ell] - z_{k}[\ell] + \eta F(z_{k})[\ell] = 0, \quad z_{k+1}[\ell] - z_{k}[\ell] + \eta F(z_{k+\frac{1}{2}})[\ell] = 0, \ \forall \ell \in [n], \quad \text{(EG Update)} \\ & \left\| \eta F(z_{i}) - \eta F(z_{j}) \right\|^{2} - (\eta L)^{2} \left\| z_{i} - z_{j} \right\|^{2} \leq 0, \quad \forall i, j \in \{k, k + \frac{1}{2}, k + 1\}, \quad \text{(Lipschitzness)} \\ & \left\langle \eta F(z_{i}) - \eta F(z_{j}), z_{j} - z_{i} \right\rangle \leq 0, \quad \forall i, j \in \{k, k + \frac{1}{2}, k + 1\}. \quad \text{(Monotonicity)} \end{split}$$

We always multiply F with η in the constraints as it will be convenient later. We use K to denote the set $\{k, k+\frac{1}{2}, k+1\}$. To obtain a certificate of non-negativity, we apply SOS programming to search for a degree-2 SOS proof. More specifically, we want to find non-negative coefficients $\{\lambda_{i,j}^*, \mu_{i,j}^*\}_{i>j,i,j\in K}$ and degree-1 polynomials $\gamma_1^{(\ell)}(w)$ and $\gamma_2^{(\ell)}(w)$ in $\mathbb{R}[w]$ for each $\ell \in [n]$, where $w := \{z_i[\ell], \eta F(z_i)[\ell]\}_{i\in K, \ell \in [n]}$, such that the following is an SOS polynomial:

$$\|\eta F(z_{k})\|^{2} - \|\eta F(z_{k+1})\|^{2} + \sum_{i>j \text{ and } i,j\in K} \lambda_{i,j}^{*} \cdot \left(\|\eta F(z_{i}) - \eta F(z_{j})\|^{2} - (\eta L)^{2} \|z_{i} - z_{j}\|^{2}\right)$$

$$+ \sum_{i>j \text{ and } i,j\in K} \mu_{i,j}^{*} \cdot \left\langle \eta F(z_{i}) - \eta F(z_{j})\right\rangle, z_{j} - z_{i} \right\rangle + \sum_{\ell \in [n]} \gamma_{1}^{(\ell)}(\boldsymbol{w})(z_{k+\frac{1}{2}}[\ell] - z_{k}[\ell] + \eta F(z_{k})[\ell])$$

$$+ \sum_{\ell \in [n]} \gamma_{2}^{(\ell)}(\boldsymbol{w})(z_{k+1}[\ell] - z_{k}[\ell] + \eta F(z_{k+\frac{1}{2}})[\ell]). \tag{7}$$

 $[\]overline{}^{12}$ In the unconstrained setting, the tangent residual is simply the norm of the operator $r^{tan}_{(F,\mathbb{R}^n)}(z) = \|F(z)\|$.

Due to constraints satisfied by the EG iterates, the non-negativity of Expression (7) clearly implies that $\|F(z_k)\|^2 - \|F(z_{k+1})\|^2$ is non-negative. However, Expression (7) is in fact an infinite family of polynomials rather than a single one. Expression (7) corresponds to a different polynomial for every integer n. To directly search for the solution, we would need to solve an infinitely large SOS program, which is clearly infeasible. By exploring the symmetry in Expression (7), we show that it suffices to solve a constant size SOS program. Let us first expand Expression (7) as follows:

$$\sum_{\ell \in [n]} \left((\eta F(z_{k})[\ell])^{2} - (\eta F(z_{k+1})[\ell])^{2} + \sum_{i>j \text{ and } i,j \in K} \lambda_{i,j}^{*} \left((\eta F(z_{i})[\ell] - \eta F(z_{j})[\ell])^{2} - (\eta L)^{2} (z_{i}[\ell] - z_{j}[\ell])^{2} \right) \\
+ \sum_{i>j \text{ and } i,j \in K} \mu_{i,j}^{*} (\eta F(z_{i})[\ell] - \eta F(z_{j})[\ell]) \left(z_{j}[\ell] - z_{i}[\ell] \right) + \gamma_{1}^{(\ell)} (\boldsymbol{w}) (z_{k+\frac{1}{2}}[\ell] - z_{k}[\ell] + \eta F(z_{k})[\ell]) \\
+ \gamma_{2}^{(\ell)} (\boldsymbol{w}) (z_{k+1}[\ell] - z_{k}[\ell] + \eta F(z_{k+\frac{1}{2}})[\ell]) \right). \tag{8}$$

What we will argue next is that, due to the symmetry across coordinates, it suffices to directly search for a single SOS proof that shows that each of the n summands in Expression (8) is an SOS polynomial. More specifically, we make use of the following two key properties. (i) For any $\ell, \ell' \in [n]$, the ℓ -th summand and ℓ' -th summand are identical subject to a change of variable; ¹³ (ii) the ℓ -th summand only depends on the coordinate ℓ , i.e., variables in $\{z_i[\ell], \eta F(z_i)[\ell]\}_{i \in K}$ and does not involve any other coordinates.¹⁴ Both properties crucially rely on the assumption that we are in the unconstrained setting. As we will see in Section 5.2, the main source of difficulty in the constrained setting is the lack of these properties. We solve the following SOS program, whose solution can be used to construct $\{\lambda_{i,j}^*, \mu_{i,j}^*\}_{i>j,i,j\in K}$ and $\{\gamma_1^{(\ell)}(\boldsymbol{w}), \gamma_2^{(\ell)}(\boldsymbol{w})\}_{\ell\in [n]}$ so that each of the summands in Expression (8) is an SOS polynomial.

The proof of the following theorem is based on a feasible solution to the SOS program in Figure 3.

Theorem 1. Let $F: \mathbb{R}^n \to \mathbb{R}^n$ be a monotone and L-Lipschitz operator. Then for any $k \in \mathbb{N}$, the EG algorithm with step size $\eta \in (0, \frac{1}{L})$ satisfies $||F(z_k)||^2 \ge ||F(z_{k+1})||^2$.

Proof. Since *F* is monotone and *L*-Lipschitz, we have

$$\langle F(z_{k+1}) - F(z_k), z_k - z_{k+1} \rangle \le 0$$

and

$$\left\|F(z_{k+\frac{1}{2}}) - F(z_{k+1})\right\|^2 - L^2 \left\|z_{k+\frac{1}{2}} - z_{k+1}\right\|^2 \le 0.$$

We simplify them using the update rule of EG and $\eta L < 1$. In particular, we replace $z_k - z_{k+1}$ with $\eta F(z_{k+\frac{1}{2}})$ and $z_{k+\frac{1}{2}} - z_{k+1}$ with $\eta F(z_{k+\frac{1}{2}}) - \eta F(z_k)$.

Tissimply replace $\{z_i[\ell]\}_{i\in K}$ and $\{\eta F(z_i)[\ell]\}_{i\in K}$ with $\{z_i[\ell']\}_{i\in K}$ and $\{\eta F(z_i)[\ell']\}_{i\in K}$.

The mainly care about the polynomials arise from the constraints. Although $\gamma_1^{(\ell)}(w)$ and $\gamma_2^{(\ell)}(w)$ could depend on other coordinates, we show that it suffices to consider polynomials in $\{z_i[\ell], \eta F(z_i)[\ell]\}_{i\in K}$.

Input Fixed Polynomials. We use x to denote (x_0, x_1, x_2) and y to denote (y_0, y_1, y_2) . Interpret x_i as $z_{k+\frac{i}{2}}[\ell]$ and y_i as $\eta F(z_{k+\frac{i}{2}})[\ell]$ for $0 \le i \le 2$. Observe that $h_1(x,y)$ and $h_2(x,y)$ come from the EG update rule on coordinate ℓ . $g_{i,j}^L(x,y)$ and $g_{i,j}^m(x,y)$ come from the ℓ -th coordinate's contribution in the Lipschitzness and monotonicity constraints.

- $h_1(x, y) := x_1 x_0 + y_0$ and $h_2(x, y) := x_2 x_0 + y_1$.
- $g_{i,j}^L(x,y) := (y_i y_j)^2 C \cdot (x_i x_j)^2$ for any $0 \le j < i \le 2$.
- $g_{i,j}^m(x,y) := (y_i y_j)(x_j x_i)$ for any $0 \le j < i \le 2$.

Decision Variables of the SOS Program:

- $p_{i,j}^L \ge 0$, and $p_{i,j}^m \ge 0$, for all $0 \le j < i \le 2$.
- $q_1(x, y)$ and $q_2(x, y)$ are two degree 1 polynomials in $\mathbb{R}[x, y]$.

Constraints of the SOS Program:

s.t.
$$y_0^2 - y_2^2 + \sum_{2 \ge i > j \ge 0} p_{i,j}^L \cdot g_{i,j}^L(x, y) + \sum_{2 \ge i > j \ge 0} p_{i,j}^m \cdot g_{i,j}^m(x, y) + q_1(x, y) \cdot h_1(x, y) + q_2(x, y) \cdot h_2(x, y) \in SOS[x, y].$$
 (9)

Figure 3: Our SOS program in the unconstrained setting.

$$\left\langle F(z_{k+1}) - F(z_k), F(z_{k+\frac{1}{2}}) \right\rangle \le 0, \tag{10}$$

$$\left\| F(z_{k+\frac{1}{2}}) - F(z_{k+1}) \right\|^2 - \left\| F(z_{k+\frac{1}{2}}) - F(z_k) \right\|^2 \le 0.$$
 (11)

In Proposition 1 at Appendix C we verify the following identity.

$$||F(z_k)||^2 - ||F(z_{k+1})||^2 + 2 \cdot \text{LHS of Inequality}(10) + \text{LHS of Inequality}(11) = 0.$$

Thus,
$$||F(z_k)||^2 - ||F(z_{k+1})||^2 \ge 0.$$

Corollary 1 is implied by combing Lemma 2, Lemma 5, Theorem 1 and the fact that $\eta \in (0, \frac{1}{L})$.

Corollary 1. Let $F(\cdot): \mathbb{R}^n \to \mathbb{R}^n$ be a monotone and L-Lipshitz operator and $z^* \in \mathbb{R}^n$ be a solution to the variational inequality. For any $T \geq 1$, let z_T be T-th iterate of the EG algorithm with constant step size $\eta \in (0, \frac{1}{L})$, then $GAP_{\mathbb{R}^n, F, D}(z_T) \leq \frac{1}{\sqrt{T}} \frac{3D\|z_0 - z^*\|}{\eta \sqrt{1 - (\eta L)^2}}$.

^aC represents $(\eta L)^2$. Larger C corresponds to a larger step size and makes the SOS program harder to satisfy. Through binary search, we find that the largest possible value of C is 1 while maintaining the feasibility of the SOS program.

5.2 Last-Iterate Convergence of EG with Arbitrary Convex Constraints

We establish the last-iterate convergence rate of the EG algorithm in the constrained setting in this section. The plan is similar to the one in Section 5.1. First, we use the assistance of SOS programming to prove the monotonicity of the tangent residual (Theorem 2), then combine it with the best-iterate convergence guarantee from Lemma 5 to derive the last-iterate convergence rate (Theorem 3).

The crux is proving the monotonicity of the tangent residual. Due to the constraints introduced by the domain \mathcal{Z} , our approach in the unconstrained case no longer applies. More specifically, the coordinates are now entangled due to projection, ruining the two key properties – Property (i) and (ii) that we rely on in the unconstrained setting. We show how to reduce *both the number of constraints and the dimension by exploiting the low dimensionality of the EG update rule*. Equipped with the reduction, we convert the problem of proving the monotonicity of the tangent residual to solving a small constant size SOS program.

Reducing the Number of Constraints. Suppose we are not given the description of \mathcal{Z} , and we only observe one iteration of the EG algorithm. In other words, we know $z_k, z_{k+\frac{1}{2}}$, and z_{k+1} , as well as $F(z_k)$, $F(z_{k+\frac{1}{2}})$, and $F(z_{k+1})$. To compute the squared tangent residual at z_k , let us also assume that the unit vector $-a_k \in \widehat{N}(z_k)$ satisfies $r_{(F,\mathcal{Z})}^{tan}(z_k)^2 = \|F(z_k)\|^2 - \langle F(z_k), -a_k \rangle^2$. From this limited information, what can we learn about \mathcal{Z} ? We can conclude that \mathcal{Z} must lie in the intersection of the following halfspaces: (a) $\langle a_k, z \rangle \geq b_k$, where $b_k = \langle a_k, z_k \rangle$. This is true because $-a_k \in \widehat{N}(z_k)$. (b) $\langle a_{k+\frac{1}{2}}, z \rangle \geq b_{k+\frac{1}{2}}$, where $a_{k+\frac{1}{2}} = \frac{z_{k+\frac{1}{2}} - z_k + \eta F(z_k)}{\|z_{k+\frac{1}{2}} - z_k + \eta F(z_k)\|}$ and $b_{k+\frac{1}{2}} = \langle a_{k+\frac{1}{2}}, z_{k+\frac{1}{2}} \rangle$. This is true because $z_{k+\frac{1}{2}} = \prod_{\mathcal{Z}} (z_k - \eta F(z_k))$, so $\langle z_{k+\frac{1}{2}} - z_k + \eta F(z_k), z - z_{k+\frac{1}{2}} \rangle \geq 0$ for all $z \in \mathcal{Z}$. (c) $\langle a_{k+1}, z \rangle \geq b_{k+1}$, where $a_{k+1} = \frac{z_{k+1} - z_k + \eta F(z_{k+\frac{1}{2}})}{\|z_{k+1} - z_k + \eta F(z_{k+\frac{1}{2}})\|}$ and $b_{k+1} = \langle a_{k+1}, z_{k+1} \rangle$. This is true because $z_{k+1} = \prod_{\mathcal{Z}} (z_k - \eta F(z_{k+\frac{1}{2}}))$, so $\langle z_{k+1} - z_k + \eta F(z_{k+\frac{1}{2}}), z - z_{k+1} \rangle \geq 0$ for all $z \in \mathcal{Z}$. See Figure 4 for illustration.

The "hardest instance" of $\mathcal Z$ that is consistent with our knowledge of z_k , $z_{k+\frac12}$, and z_{k+1} is when $\mathcal Z$ is exactly the intersection of these three halfspaces. In such case, the squared tangent residual of z_{k+1} is $\|F(z_{k+1})\|^2 - \langle F(z_{k+1}), a_{k+1} \rangle^2 \cdot \mathbb{1}[\langle F(z_{k+1}), a_{k+1} \rangle \geq 0]$, and it is an upper bound of $r_{(F,\mathcal Z)}^{tan}(z_{k+1})^2$ for any other consistent $\mathcal Z$. Our goal is to prove the tangent residual is non-increasing even in the "hardest case", that is, to prove the non-negativity of

$$||F(z_k)||^2 - \langle F(z_k), a_k \rangle^2 - \left(||F(z_{k+1})||^2 - \langle F(z_{k+1}), a_{k+1} \rangle^2 \cdot \mathbb{1}[\langle F(z_{k+1}), a_{k+1} \rangle \ge 0] \right)$$
(12)

Low-dimensionality of an EG Update. As there are only three hyperplanes $\langle a_i, z \rangle \geq b_i$ for $i \in \{k, k+\frac{1}{2}, k+1\}$ involved, we can choose a new basis, so that $a_{k+1} = (1, 0, \ldots, 0), \ a_{k+\frac{1}{2}} = (\theta_1, \theta_2, 0, \ldots, 0), \ \text{and} \ a_k = (\sigma_1, \sigma_2, \sigma_3, 0, \ldots, 0).$ As $a_{k+\frac{1}{2}}$ and $z_{k+\frac{1}{2}} - z_k + \eta F(z_k)$ are co-directed, and a_{k+1} and $a_{k+1} - z_k + \eta F(z_{k+\frac{1}{2}})$ are co-directed, an important property of this change of basis is that the

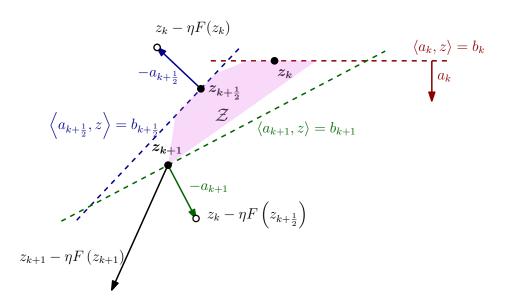


Figure 4: We illustrate how to reduce the number of constraints. Let $\mathcal{Z}' = \left\{z: \langle a_{k+\frac{i}{2}}, z \rangle \geq b_{k+\frac{i}{2}}, \text{ for } i \in \{0,1,2\}\right\}$. Pictorially $-F(z_{k+1}) \in N_{\mathcal{Z}}(z_{k+1})$, which further implies $r_{(F,\mathcal{Z})}^{tan}(z_{k+1}) = 0$ and $r_{(F,\mathcal{Z})}^{tan}(z_{k+1}) > 0$.

EG update from z_k to z_{k+1} is unconstrained in all coordinates $\ell \geq 4$. More specifically,

$$z_{k+\frac{1}{2}}[\ell] - z_k[\ell] + \eta F(z_k)[\ell] = 0, \quad z_{k+1}[\ell] - z_k[\ell] + \eta F(z_{k+\frac{1}{2}})[\ell] = 0, \ \forall \ell \geq 4.$$

Hence, we can represent all of the coordinates $\ell \geq 4$ with one coordinate in the SOS program similar to the unconstrained case. We still need to keep the first three dimensions, but now we only face a problem in dimension 4 rather than in dimension n, and we can form a constant size SOS program to search for a certificate of non-negativity for Expression (12).

In Lemma 6, we further simplify the instance that we need to consider. In particular, we argue that it is w.l.o.g. to assume that (1) a_{k+1} , $a_{k+\frac{1}{2}}$, and a_k are linear independent and (2) the intersection of the three halfspaces forms a cone, i.e., $b_k = b_{k+\frac{1}{2}} = b_{k+1} = 0$. Both assumption (1) and (2) reduce the number of variables we need to consider in the SOS program, so a low degree SOS proof is more likely to exist. To maximally reduce the number of variables, we only included the minimal number of constraints that suffice to derive an SOS proof.

The proof of Lemma 6 is postponed to Appendix E.

Lemma 6 (Simplification Procedure). Let \mathcal{I} be a variational inequality problem for a closed convex set $\mathcal{Z} \subseteq \mathbb{R}^n$ and a monotone and L-Lipschitz operator $F: \mathcal{Z} \to \mathbb{R}^n$. Suppose the EG algorithm has a constant step size η . Let z_k be the k-th iteration of the EG algorithm, $z_{k+\frac{1}{2}}$ be the $(k+\frac{1}{2})$ -th iteration as defined in (4), and z_{k+1} be the (k+1)-th iteration as defined in (5).

Then either $r_{(F,\mathcal{Z})}^{tan}(z_k) \geq r_{(F,\mathcal{Z})}^{tan}(z_{k+1})$, or there exist vectors $\overline{a}_k, \overline{a}_{k+\frac{1}{2}}, \overline{a}_{k+1}, \overline{z}_k, \overline{z}_{k+\frac{1}{2}}, \overline{z}_{k+1}, \overline{F}(\overline{z}_k)$, $\overline{F}(\overline{z}_{k+\frac{1}{2}}), \overline{F}(\overline{z}_{k+1}) \in \mathbb{R}^N$ with $N \leq n+5$ that satisfy the following conditions.

1. $\overline{a}_k = (\beta_1, \beta_2, 1, 0, \dots, 0), \overline{a}_{k+\frac{1}{2}} = (\alpha, 1, 0, \dots, 0), \text{ and } \overline{a}_{k+1} = (1, 0, \dots, 0) \text{ for some } \alpha, \beta_1, \beta_2 \in \mathbb{R}.$

$$2. \|\overline{F}(\overline{z}_k) - \frac{\langle F(\overline{z}_k), \overline{a}_k \rangle \cdot \overline{a}_k}{\|\overline{a}_k\|^2} \|^2 - \|\overline{F}(\overline{z}_{k+1}) - \frac{\langle \overline{F}(\overline{z}_{k+1}), \overline{a}_{k+1} \rangle \cdot \overline{a}_{k+1}}{\|\overline{a}_{k+1}\|^2} \mathbb{1}[\langle \overline{F}(\overline{z}_{k+1}), \overline{a}_{k+1} \rangle \ge 0] \|^2 < 0.$$

3. Additionally, $\langle \overline{a}_i, \overline{z}_j \rangle \geq 0$ and $\langle \overline{a}_i, \overline{z}_i \rangle = 0$ for all $i, j \in \{k, k + \frac{1}{2}, k + 1\}$. $\overline{a}_{k + \frac{1}{2}}$ and $\overline{z}_{k + \frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_k)$ are co-directed, i.e., they are colinear and have the same direction, and \overline{a}_{k+1} and $\overline{z}_{k+1} - \overline{z}_k + \eta \overline{F}(\overline{z}_{k + \frac{1}{2}})$ are co-directed.

4.

$$\left\| \overline{F}(\overline{z}_{k+1}) - \overline{F}(\overline{z}_{k+\frac{1}{2}}) \right\|^2 \le L^2 \left\| \overline{z}_{k+1} - \overline{z}_{k+\frac{1}{2}} \right\|^2$$
 (13)

$$\langle \overline{F}(\overline{z}_{k+1}) - \overline{F}(\overline{z}_k), \overline{z}_{k+1} - \overline{z}_k \rangle \ge 0$$
 (14)

$$\langle \overline{a}_k, \overline{F}(\overline{z}_k) \rangle \ge 0.$$
 (15)

In Theorem 2, we establish the monotonicity of the tangent residual. Our proof is based on the solution to the degree-8 SOS program concerning polynomials in 27 variables (Figure 5). We include a proof sketch of Theorem 2 in the main body, and the full proof in Appendix E.

Theorem 2. Let $\mathcal{Z} \subseteq \mathbb{R}^n$ be a closed convex set and $F: \mathcal{Z} \to \mathbb{R}^n$ be a monotone and L-Lipschitz operator. For any step size $\eta \in (0, \frac{1}{L})$ and any $z_k \in \mathcal{Z}$, the EG method update satisfies $r_{(F,\mathcal{Z})}^{tan}(z_k) \geq r_{(F,\mathcal{Z})}^{tan}(z_{k+1})$.

Proof Sketch Assume towards contradiction that $r^{tan}(z_k) < r^{tan}(z_{k+1})$, using Lemma 6 there exist numbers $\alpha, \beta_1, \beta_2 \in \mathbb{R}$ and vectors \overline{a}_k , $\overline{a}_{k+\frac{1}{2}}$, \overline{a}_{k+1} , \overline{z}_k , $\overline{z}_{k+\frac{1}{2}}$, \overline{z}_{k+1} , $\overline{F}(\overline{z}_k)$, $\overline{F}(\overline{z}_{k+\frac{1}{2}})$, $\overline{F}(\overline{z}_{k+1}) \in \mathbb{R}^N$ where N = n + 5 that satisfy the properties in the statement of Lemma 6 and

$$0 > \left\| \overline{F}(\overline{z}_k) - \frac{\langle F(\overline{z}_k), \overline{a}_k \rangle}{\left\| \overline{a}_k \right\|^2} \overline{a}_k \right\|^2 - \left\| \overline{F}(\overline{z}_{k+1}) - \frac{\langle \overline{F}(\overline{z}_{k+1}), \overline{a}_{k+1} \rangle \overline{a}_{k+1}}{\left\| \overline{a}_{k+1} \right\|^2} \mathbb{1} \left[\left\langle \overline{F}(\overline{z}_{k+1}), \overline{a}_{k+1} \right\rangle \ge 0 \right] \right\|^2.$$

We use TARGET to denote the RHS above. We first present 8 inequalities/equations, which directly follow from Lemma 6. We verify their correctness in the complete proof in Appendix E.

$$\langle \eta \overline{F}(\overline{z}_{k+1}) - \eta \overline{F}(\overline{z}_k), \overline{z}_k - \overline{z}_{k+1} \rangle \le 0,$$
 (16)

$$\left\| \eta \overline{F}(\overline{z}_{k+1}) - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}}) \right\|^2 - \left\| \overline{z}_{k+1} - \overline{z}_{k+\frac{1}{2}} \right\|^2 \le 0, \tag{17}$$

$$\overline{z}_{k+\frac{1}{2}}[1]\left(\left(\overline{z}_{k+\frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_k)\right)[1] - \alpha\left(\overline{z}_{k+\frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_k)\right)[2]\right) = 0, \tag{18}$$

$$\overline{z}_{k+1}[2]\left(\left(\overline{z}_{k+\frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_k)\right)[1] - \alpha\left(\overline{z}_{k+\frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_k)\right)[2]\right) = 0, \tag{19}$$

$$\left(\alpha\left(\overline{z}_{k}-\eta\overline{F}(\overline{z}_{k})\right)[1]+\left(\overline{z}_{k}-\eta\overline{F}(\overline{z}_{k})\right)[2]\right)\left(\alpha\overline{z}_{k+1}[1]+\overline{z}_{k+1}[2]\right)\leq0,\tag{20}$$

$$-\eta \left(\beta_1 \overline{F}(\overline{z}_k)[1] + \beta_2 \overline{F}(\overline{z}_k)[2] + \overline{F}(\overline{z}_k)[3]\right) \left(\beta_1 \overline{z}_{k+\frac{1}{2}}[1] + \beta_2 \overline{z}_{k+\frac{1}{2}}[2] + \overline{z}_{k+\frac{1}{2}}[3]\right) \le 0, \tag{21}$$

$$\overline{z}_{k}[1]\left(\overline{z}_{k}[1] - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[1]\right) \le 0, \tag{22}$$

$$-\eta \overline{F}(\overline{z}_{k+1})[1]\mathbb{1}\left[\overline{F}(\overline{z}_{k+1})[1] \le 0\right] \left(\overline{z}_{k}[1] - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[1]\right) \le 0. \tag{23}$$

We first add several non-positive terms to TARGET to derive Expression 24.

 $\eta^2 \cdot \text{TARGET} + 2 \cdot \text{LHS of Inequality (16)} + \text{LHS of Inequality (17)} + 2 \cdot \text{LHS of Equation (18)}$ $+ \frac{2\alpha}{1+\alpha^2} \cdot \text{LHS of Equation (19)} + \frac{2}{1+\alpha^2} \cdot \text{LHS of Inequality (20)} + \frac{2}{1+\beta_1^2+\beta_2^2} \cdot \text{LHS of Inequality (21)}$ $+ 2 \cdot \text{LHS of Inequality (22)} + 2 \cdot \text{LHS of Inequality (23)}$ (24)

After substituting the following six variables $\bar{z}_k[3]$, $\bar{z}_{k+\frac{1}{2}}[2]$, $\bar{z}_{k+\frac{1}{2}}[3]$, $\bar{z}_{k+1}[1]$, $\bar{z}_{k+1}[2]$, $\bar{z}_{k+1}[3]$ using Equation (63) to (68), Expression (24) equals to the following polynomial, which is clearly nonnegative. We verify the validity of the substitutions (Equation (63) to (68)) in the full proof.

$$\begin{split} & \left(\overline{z}_{k}[1] - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[1] + \eta \overline{F}(\overline{z}_{k+1})[1] \cdot \mathbb{1}[\overline{F}(\overline{z}_{k+1})[1] \geq 0] \right)^{2} \\ & + \frac{\left(\overline{z}_{k}[1] - \eta \overline{F}(\overline{z}_{k})[1] - \overline{z}_{k+\frac{1}{2}}[1] \right)^{2}}{1 + \beta_{1}^{2} + \beta_{2}^{2}} \\ & + \frac{\left(\eta \overline{F}(\overline{z}_{k})[3] + \beta_{1} \overline{z}_{k}[1] + \beta_{2} \overline{z}_{k}[2] + (\alpha \beta_{2} - \beta_{1}) \overline{z}_{k+\frac{1}{2}}[1] \right)^{2}}{1 + \beta_{1}^{2} + \beta_{2}^{2}} \\ & + \frac{\left(\overline{z}_{k}[2] - \eta \overline{F}(\overline{z}_{k})[2] + \alpha \overline{z}_{k+\frac{1}{2}}[1] \right)^{2}}{1 + \beta_{1}^{2} + \beta_{2}^{2}} \\ & + \frac{\left(\beta_{1}(\overline{z}_{k}[2] - \eta \overline{F}(\overline{z}_{k})[2] + \alpha \overline{z}_{k+\frac{1}{2}}[1] \right) - \beta_{2}(\overline{z}_{k}[1] - \eta \overline{F}(\overline{z}_{k})[1] - \overline{z}_{k+\frac{1}{2}}[1]) \right)^{2}}{1 + \beta_{1}^{2} + \beta_{2}^{2}} \end{split}$$

Theorem 3. Let $\mathcal{Z} \subseteq \mathbb{R}^n$ be a closed convex set, $F(\cdot): \mathcal{Z} \to \mathbb{R}^n$ be a monotone and L-Lipschitz operator and $z^* \in \mathcal{Z}$ be the solution to the variational inequality. Then for any $T \geq 1$, z_T produced by EG with any constant step size $\eta \in (0, \frac{1}{L})$ satisfies $GAP(z_T) \leq \frac{1}{\sqrt{T}} \frac{3D||z_0 - z^*||}{\eta \sqrt{1 - (\eta L)^2}}$.

Theorem 3 is implied by combing Lemma 2, Lemma 5, Theorem 2 and the fact that $\eta \in (0, \frac{1}{L})$. Choosing η to be $\frac{1}{2L}$ and $D = O(\|z_0 - z^*\|)$, then $\operatorname{GAP}(z_T) = O(\frac{D^2L}{\sqrt{T}})$ matching the $\Omega(\frac{D^2L}{\sqrt{T}})$ lower bound for EG, OGDA, and more generally all p-SCLI algorithms [Golowich et al., 2020b,a] in terms of the dependence on D, L, and T.

6 Conclusion

We propose a new method based on SOS programming to analyze iterative methods. Combining with the low-dimensionality of the EG update rule, we provide the first and tight last-iterate convergence rate of the EG algorithm for constrained monotone and Lipschitzs VIs, which, to the best of our knowledge, is also the first last-iterate convergence rate for any p-SCLI algorithm in the constrained setting. We believe many other iterative methods also have the low-dimensionality that make them amenable to our SOS programming based analysis.

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A Additional Preliminaries

For $z \in \mathbb{R}^n$ and D > 0, we use $\mathcal{B}(z, D) = \{z' \in \mathbb{R}^n : ||z' - z|| \le R\}$ to denote the ball of radius D, centered at z.

Min-Max Saddle Points. A special case of the variational inequality problem is the constrained min-max problem $\min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} f(x, y)$, where \mathcal{X} and \mathcal{Y} are closed convex sets in \mathbb{R}^n , and $f(\cdot, \cdot)$ is smooth, convex in x, and concave in y. It is well known that if one set $F(x,y) = \begin{pmatrix} \nabla_x f(x,y) \\ -\nabla_y f(x,y) \end{pmatrix}$, then F(x,y) is a monotone and Lipschitz operator [Facchinei and Pang, 2007].

Equilibria of Monotone Games. Monotone games are a large class of multi-player games that include many common and well-studied class of games such as bilinear games, λ -cocoercive games [Lin et al., 2020], zero-sum polymatrix games [Cai and Daskalakis, 2011, Cai et al., 2016], and zero-sum socially-concave games [Even-Dar et al., 2009]. Besides, the min-max saddle point problem is a special case of two-player monotone games. We include the definition of monotone games here and remind readers that finding a Nash Equilibrium of a monotone game is exactly the same as finding a solution to a monotone variational inequality.

A continuous game \mathcal{G} is denoted as $(\mathcal{N}, (\mathcal{X}_i)_{i \in [N]}, (f_i)_{i \in [N]})$ where there are N players $\mathcal{N} = \{1, \cdots, N\}$. Player $i \in \mathcal{N}$ chooses action from a closed convex set $\mathcal{X}_i \in \mathbb{R}^{n_i}$ such that $\mathcal{X} := \prod_{i \in \mathcal{N}} \mathcal{X}_i \in \mathbb{R}^n$ and wants to minimize its cost function $f_i : \mathcal{X} \to \mathbb{R}$. For each player i, we denote x_{-i} the vector of actions of all the other players. A *Nash Equilibrium* of game \mathcal{G} is an action profile $x^* \in \mathcal{X}$ such that $f_i(x^*) \leq f_i(x_i', x_{-i}^*)$ for any $x_i' \in \mathcal{X}_i$. Let $F(x) = (\nabla_{x_i} f_i(x), \cdots, \nabla_{x_N} f_N(x)) \in \mathbb{R}^n$. We say \mathcal{G} is *monotone* if $\langle F(x) - F(x'), x - x' \rangle \geq 0$ for any $x, x' \in \mathcal{X}$.

In Lemma 7 we present several equivalent formulations of the tangent residual.

Lemma 7. Let \mathcal{Z} be a closed convex set and $F: \mathcal{Z} \to \mathbb{R}^n$ be an operator. Denote $N_{\mathcal{Z}}(z)$ the normal cone of z and $J_{\mathcal{Z}}(z) := \{z\} + T_{\mathcal{Z}}(z)$, where $T_{\mathcal{Z}}(z) = \{z' \in \mathbb{R}^n : \langle z', a \rangle \leq 0, \forall a \in N_{\mathcal{Z}}(z)\}$ is the tangent cone of z. Then all of the following quantities are equivalent:

1.
$$\sqrt{\|F(z)\|^2 - \max_{\substack{a \in \widehat{N}_{\mathcal{Z}}(z), \ \langle F(z), a \rangle^2}} \frac{\langle F(z), a \rangle^2}{\langle F(z), a \rangle^2}}$$

2.
$$\min_{\substack{a \in \widehat{N}_{\mathcal{Z}}(z), \\ \langle F(z), a \rangle \leq 0}} \|F(z) - \langle F(z), a \rangle \cdot a\|$$

3.
$$\left\|\Pi_{T_{\mathcal{Z}}(z)}\right[-F(z)\right\|$$

4.
$$\left\|\Pi_{J_{\mathcal{Z}}(z)}\left[z-F(z)\right]-z\right\|$$

5.
$$\left\| -F(z) - \Pi_{N_{\mathcal{Z}}(z)} \right\| - F(z)$$

6.
$$\min_{a \in N_{\mathcal{Z}}(z)} ||F(z) + a||$$

Proof. (quantity 1 = quantity 2). Observe that

$$\min_{\substack{a \in \widehat{N}_{\mathcal{Z}}(z), \\ \langle F(z), a \rangle \leq 0}} \|F(z) - \langle F(z), a \rangle \cdot a\|^2 = \|F(z)\|^2 - \max_{\substack{a \in \widehat{N}_{\mathcal{Z}}(z), \\ \langle F(z), a \rangle \leq 0}} \langle F(z), a \rangle^2 \cdot \left(2 - \|a\|^2\right).$$

Therefore, it is enough to show that $\max_{\substack{a \in \widehat{N}_{\mathcal{Z}}(z), \\ \langle F(z), a \rangle \leq 0}} \langle F(z), a \rangle^2 \cdot (2 - \|a\|^2) = ax = (F(z), a)^2$ If $\widehat{N}_{\mathcal{Z}}(z) = \{(0, 0)\}$, then the equality holds trivially. Now we

 $\max_{\substack{a \in \widehat{N}_{\mathcal{Z}}(z), \\ \langle F(z), a \rangle \leq 0}} \langle F(z), a \rangle^2$. If $\widehat{N}_{\mathcal{Z}}(z) = \{(0, \dots, 0)\}$, then the equality holds trivially. Now we

assume that $\{(0,\ldots,0)\} \subseteq \widehat{N}_{\mathcal{Z}}(z)$ and consider any $a \in \widehat{N}_{\mathcal{Z}}(z) \setminus (0,\ldots,0)$. Let $c \in \left[1,\frac{1}{\|a\|}\right]$. By Definition 2, $\|a\| \le 1$, which implies that $c \cdot a \in \widehat{N}_{\mathcal{Z}}(z)$. We try to maximize the following objective

$$\langle F(z), c \cdot a \rangle^2 \cdot \left(2 - c^2 ||a||^2 \right) = \frac{\langle F(z), a \rangle^2}{||a||^2} \cdot c^2 ||a||^2 \cdot \left(2 - c^2 ||a||^2 \right).$$

One can easily verify that function $c^2\|a\|^2 \cdot (2-c^2\|a\|^2)$ is maximized when $c^2\|a\|^2 = 1 \Leftrightarrow c = \frac{1}{\|a\|}$. Thus when $\{(0,\ldots,0)\} \subsetneq \widehat{N}_{\mathcal{Z}}(z)$,

$$\max_{\substack{a \in \widehat{N}_{\mathcal{Z}}(z), \\ \langle F(z), a \rangle \leq 0}} \langle F(z), a \rangle^2 \cdot \left(2 - \|a\|^2 \right) = \max_{\substack{a \in \widehat{N}_{\mathcal{Z}}(z), \\ \langle F(z), a \rangle \leq 0, \\ \|a\| = 1}} \langle F(z), a \rangle^2 \cdot \left(2 - \|a\|^2 \right)$$

$$= \max_{\substack{a \in \widehat{N}_{\mathcal{Z}}(z), \\ \langle F(z), a \rangle \leq 0, \\ \|a\| = 1}} \langle F(z), a \rangle^2$$

$$= \max_{\substack{a \in \widehat{N}_{\mathcal{Z}}(z), \\ \langle F(z), a \rangle \leq 0, \\ \|a\| = 1}} \langle F(z), a \rangle^2,$$

$$= \max_{\substack{a \in \widehat{N}_{\mathcal{Z}}(z), \\ \langle F(z), a \rangle \leq 0}} \langle F(z), a \rangle^2,$$

which concludes the proof.

(quantity 3 = quantity 4). By definition, $J_{\mathcal{Z}}(z) = \{z\} + T_{\mathcal{Z}}(z)$. Thus we have

$$\left\|\Pi_{I_{\mathcal{Z}}(z)}\left[z-F(z)\right]-z\right\|=\left\|\Pi_{T_{\mathcal{Z}}(z)}\left[-F(z)\right]\right\|.$$

(quantity 4 = quantity 5). By definition, the tangent cone $T_{\mathcal{Z}}(z)$ is the polar cone of the normal cone $N_{\mathcal{Z}}(z)$. Since $N_{\mathcal{Z}}(z)$ is a closed convex cone, by Moreau's decomposition theorem, we have for any vector $x \in \mathbb{R}^n$,

$$x = \Pi_{N_{\mathcal{Z}}(z)}(x) + \Pi_{T_{\mathcal{Z}}(z)}(x), \qquad \left\langle \Pi_{N_{\mathcal{Z}}(z)}(x), \Pi_{T_{\mathcal{Z}}(z)}(x) \right\rangle = 0.$$

Thus it is clear that we have

$$\begin{split} \left\| \Pi_{J_{\mathcal{Z}}(z)} \Big[z - F(z) \Big] - z \right\| &= \left\| \Pi_{T_{\mathcal{Z}}(z)} \Big[- F(z) \Big] \right\| \\ &= \left\| - F(z) - \Pi_{N_{\mathcal{Z}}(z)} \Big[- F(z) \Big] \right\|. \end{split}$$

(quantity 5 = quantity 6). Denote $a^* := \Pi_{N_{\mathcal{Z}}(z)} \Big[-F(z) \Big]$. By definition of projection, we have

$$a^* = \operatorname*{argmin}_{a \in N_{\mathcal{Z}}(z)} ||F(z) + a||^2.$$

Thus

$$\left\| -F(z) - \Pi_{N_{\mathcal{Z}}(z)} \left[-F(z) \right] \right\|^2 = \|F(z) + a^*\|^2 = \min_{a \in N_{\mathcal{Z}}(z)} \|F(z) + a\|^2.$$

(quantity 6 = quantity 2). For any fix non-zero $a \in N_{\mathcal{Z}}(z)$, (i) if $\langle F(z), a \rangle > 0$, then $||F(z) + a||^2 \ge ||F(z)||^2$, and (ii) if $\langle F(z), a \rangle \le 0$, $||F(z) + a||^2 \ge ||F(z) - \langle F(z), \frac{a}{||a||} \rangle \cdot \frac{a}{||a||}||^2$, as

$$\min_{r \ge 0} ||F(z) + r \cdot a||^2 = \left| |F(z) - \left\langle F(z), \frac{a}{||a||} \right\rangle \cdot \frac{a}{||a||} \right|^2.$$

Hence,

$$\min_{a \in N_{\mathcal{Z}}(z)} \|F(z) + a\|^2 = \min_{\substack{a' \in \widehat{N}_{\mathcal{Z}}(z) \\ \langle F(z), a' \rangle \le 0}} \|F(z) - \langle F(z), a' \rangle \cdot a'\|^2$$

The first equality is because for any $a \in N_{\mathcal{Z}}(z)$, there exists $a' \in \widehat{N}_{\mathcal{Z}}(z)$ so that $||F(z) + a||^2 \ge ||F(z) - \langle F(z), a' \rangle||^2$.

Proof of Lemma 2: If $\langle a, F(z) \rangle \geq 0$ for all $a \in \widehat{N}(z)$, then we have $r^{tan}(z) = ||F(z)||$. Thus for any $z' \in \mathcal{Z}$, by Cauchy-Schwarz inequality, we have

$$\langle F(z), z - z' \rangle \le ||F(z)|| ||z - z'|| \le D \cdot r^{tan}(z).$$

Otherwise there exists $a \in \widehat{N}(z)$ such that ||a|| = 1, $\langle a, F(z) \rangle < 0$ and $r^{tan}(z) = \sqrt{||F(z)||^2 - \langle a, F(z) \rangle^2} = ||F(z) - \langle a, F(z) \rangle a||$. Then for any $z' \in \mathcal{Z}$, we have

$$\begin{split} \left\langle F(z), z - z' \right\rangle &= \left\langle F(z) - \left\langle a, F(z) \right\rangle a, z - z' \right\rangle + \left\langle a, F(z) \right\rangle \cdot \left\langle a, z - z' \right\rangle \\ &\leq \left\langle F(z) - \left\langle a, F(z) \right\rangle a, z - z' \right\rangle \\ &\leq \left\| F(z) - \left\langle a, F(z) \right\rangle a \right\| \left\| z - z' \right\| \\ &\leq D \cdot r^{tan}(z), \end{split}$$

where we use $\langle a, F(z) \rangle < 0$ and $\langle a, z - z' \rangle \ge 0$ in the first inequality and Cauchy-Schwarz inequality in the second inequality.

If
$$\mathcal{Z}' = \mathcal{X}' \times \mathcal{Y}'$$
 and $F(x,y) = \begin{pmatrix} \nabla_x f(x,y) \\ -\nabla_y f(x,y) \end{pmatrix}$ for a convex-concave function f then

$$\begin{split} \operatorname{dg}_{f}^{\mathcal{X}',\mathcal{Y}'}(z) &= \max_{y' \in \mathcal{Y}'} f(x,y') - \min_{x' \in \mathcal{X}'} f(x',y) \\ &= \max_{y' \in \mathcal{Y}'} (f(x,y') - f(x,y)) - \min_{x' \in \mathcal{X}'} (f(x',y) - f(x,y)), \\ &\leq \max_{y' \in \mathcal{Y}'} \left\langle \nabla_{y} f(x,y), y' - y \right\rangle + \max_{x' \in \mathcal{X}'} \left\langle \nabla_{x} f(x,y), x - x' \right\rangle, \\ &= \max_{z' \in \mathcal{Z}'} \left\langle F(z), z - z' \right\rangle, \\ &\leq D\sqrt{2} \cdot r^{tan}(z), \end{split}$$

where we use the fact that f is a convex-concave function in the first inequality and $||z-z'|| = \sqrt{||x-x'||^2 + ||y-y'||^2} \le \sqrt{2}D$ in the second inequality.

B Missing Proofs from Section 4

Proof of Lemma 3: By Pythagorean inequality,

$$||z_{k+1} - z^*||^2 \le ||z_k - \eta F(z_{k+\frac{1}{2}}) - z^*||^2 - ||z_k - \eta F(z_{k+\frac{1}{2}}) - z_{k+1}||^2$$

$$= ||z_k - z^*||^2 - ||z_k - z_{k+1}||^2 + 2\eta \langle F(z_{k+\frac{1}{2}}), z^* - z_{k+1} \rangle$$

$$= ||z_k - z^*||^2 - ||z_k - z_{k+1}||^2 + 2\eta \langle F(z_{k+\frac{1}{2}}), z^* - z_{k+\frac{1}{2}} \rangle + 2\eta \langle F(z_{k+\frac{1}{2}}), z_{k+\frac{1}{2}} - z_{k+1} \rangle.$$
(25)

We first use monotonicity of $F(\cdot)$ to argue that $\langle F(z_{k+\frac{1}{2}}), z^* - z_{k+\frac{1}{2}} \rangle \leq 0$.

Fact 1. For all
$$z \in \mathcal{Z}$$
, $\langle F(z), z^* - z \rangle \leq 0$.

Proof.

$$0 \le \langle F(z^*) - F(z), z^* - z \rangle$$
 (monotonicity of $F(\cdot)$)

$$= \langle F(z^*), z^* - z \rangle - \langle F(z), z^* - z \rangle$$
 (optimality of z^* and $z \in \mathcal{Z}$)

We can simplify Equation (25) using Fact 1:

$$\begin{split} \|z_{k+1} - z^*\|^2 &\leq \|z_k - z^*\|^2 - \|z_k - z_{k+1}\|^2 + 2\eta \langle F(z_{k+\frac{1}{2}}), z_{k+\frac{1}{2}} - z_{k+1} \rangle \\ &= \|z_k - z^*\|^2 - \|z_k - z_{k+\frac{1}{2}}\|^2 - \|z_{k+\frac{1}{2}} - z_{k+1}\|^2 - 2\langle z_k - \eta F(z_{k+\frac{1}{2}}) - z_{k+\frac{1}{2}}, z_{k+\frac{1}{2}} - z_{k+1} \rangle \\ &= \|z_k - z^*\|^2 - \|z_k - z_{k+\frac{1}{2}}\|^2 - \|z_{k+\frac{1}{2}} - z_{k+1}\|^2 \\ &\qquad \qquad - 2\langle z_k - \eta F(z_k) - z_{k+\frac{1}{2}}, z_{k+\frac{1}{2}} - z_{k+1} \rangle - 2\langle \eta F(z_k) - \eta F(z_{k+\frac{1}{2}}), z_{k+\frac{1}{2}} - z_{k+1} \rangle \\ &\leq \|z_k - z^*\|^2 - \|z_k - z_{k+\frac{1}{2}}\|^2 - \|z_{k+\frac{1}{2}} - z_{k+1}\|^2 - 2\eta \langle F(z_k) - F(z_{k+\frac{1}{2}}), z_{k+\frac{1}{2}} - z_{k+1} \rangle \end{split}$$

The last inequality is because $\langle z_k - \eta F(z_k) - z_{k+\frac{1}{2}}, z_{k+\frac{1}{2}} - z_{k+1} \rangle \geq 0$, which follows from the that fact that $z_{k+\frac{1}{2}} = \Pi_{\mathcal{Z}}[z_k - \eta F(z_k)]$ and $z_{k+1} \in \mathcal{Z}$.

Finally, since $F(\cdot)$ is L-Lipschitz, we know that

$$-\langle F(z_k) - F(z_{k+\frac{1}{2}}), z_{k+\frac{1}{2}} - z_{k+1} \rangle \leq \|F(z_k) - F(z_{k+\frac{1}{2}})\| \cdot \|z_{k+\frac{1}{2}} - z_{k+1}\| \leq L\|z_k - z_{k+\frac{1}{2}}\| \cdot \|z_{k+\frac{1}{2}} - z_{k+1}\|.$$

So we can further simplify the inequality as follows:

$$\begin{split} \|z_{k+1} - z^*\|^2 \leq & \|z_k - z^*\|^2 - \|z_k - z_{k+\frac{1}{2}}\|^2 - \|z_{k+\frac{1}{2}} - z_{k+1}\|^2 - 2\eta \langle F(z_k) - F(z_{k+\frac{1}{2}}), z_{k+\frac{1}{2}} - z_{k+1} \rangle \\ \leq & \|z_k - z^*\|^2 - \|z_k - z_{k+\frac{1}{2}}\|^2 - \|z_{k+\frac{1}{2}} - z_{k+1}\|^2 + 2\eta L \|z_k - z_{k+\frac{1}{2}}\| \cdot \|z_{k+\frac{1}{2}} - z_{k+1} \| \\ \leq & \|z_k - z^*\|^2 - (1 - \eta^2 L^2) \|z_k - z_{k+\frac{1}{2}}\|^2 \end{split}$$

Hence,

$$||z_k - z^*||^2 \ge ||z_{k+1} - z^*||^2 + (1 - \eta^2 L^2)||z_k - z_{k+\frac{1}{2}}||^2.$$

Proof of Lemma 4: We need the following fact for our proof.

Fact 2.
$$||z_{k+\frac{1}{2}} - z_{k+1}|| \le \eta L ||z_k - z_{k+\frac{1}{2}}||$$
. Moreover, when $\eta L < 1$, $||z_{k+\frac{1}{2}} - z_{k+1}|| \le \frac{||z_k - z_{k+1}||}{1 - \eta L}$.

Proof. Recall that $z_{k+\frac{1}{2}} = \Pi_{\mathcal{Z}}[z_k - \eta F(z_k)]$ and $z_{k+1} = \Pi_{\mathcal{Z}}[z_k - \eta F(z_{k+\frac{1}{2}})]$. By the non-expansiveness of the projection operator and the *L*-Lipschitzness of operator *F*, we have that $\|z_{k+\frac{1}{2}} - z_{k+1}\| \le \|\eta(F(z_{k+\frac{1}{2}}) - F(z_k))\| \le \eta L \|z_k - z_{k+\frac{1}{2}}\|$.

Finally, by the triangle inequality

$$||z_k - z_{k+1}|| \ge ||z_k - z_{k+\frac{1}{2}}|| - ||z_{k+\frac{1}{2}} - z_{k+1}|| \ge (1 - \eta L) ||z_k - z_{k+\frac{1}{2}}||.$$

Now we prove Lemma 4. Since $z_{k+1} = \Pi_{\mathcal{Z}}\left[z_k - \eta F(z_{k+\frac{1}{2}})\right]$, for every $z \in \mathcal{Z}$ we have:

$$\langle z_k - \eta F(z_{k+\frac{1}{2}}) - z_{k+1}, z_{k+1} \rangle \ge \langle z_k - \eta F(z_{k+\frac{1}{2}}) - z_{k+1}, z \rangle.$$

For the rest of the proof we assume that $z_k - \eta F(z_{k+\frac{1}{2}}) - z_{k+1} \neq 0$, we show how to remove this assumption at the end of the proof. Thus vector $a = \frac{z_k - \eta F(z_{k+\frac{1}{2}}) - z_{k+1}}{\|z_k - \eta F(z_{k+\frac{1}{2}}) - z_{k+1}\|} \in \widehat{N}(z_{k+1})$. Let

$$a_{\perp} = \begin{cases} \frac{F(z_{k+1}) - \langle a, F(z_{k+1}) \rangle \cdot a}{\|F(z_{k+1}) - \langle a, F(z_{k+1}) \rangle \cdot a\|} & \text{if } \|F(z_{k+1}) - \langle a, F(z_{k+1}) \rangle \cdot a\| \neq 0\\ (0, \dots, 0) & \text{otherwise} \end{cases}.$$

Observe that $\langle a, a \rangle = 1$, $\langle a_{\perp}, a \rangle = 0$ and $F(z_{k+1}) = \langle a, F(z_{k+1}) \rangle a + \langle a_{\perp}, F(z_{k+1}) \rangle a_{\perp}$. By Definition 3 we have

$$r^{tan}(z_{k+1}) \leq \sqrt{||F(z_{k+1})||^2 - \langle a, F(z_{k+1}) \rangle^2 \cdot \mathbb{1}[\langle a, F(z_{k+1}) \rangle \leq 0]}$$

$$= \sqrt{\langle a_{\perp}, F(z_{k+1}) \rangle^2 + \langle a, F(z_{k+1}) \rangle^2 \cdot \mathbb{1}[\langle a, F(z_{k+1}) \rangle > 0]}.$$
(26)

We first present a few inequalities that will help us to prove the claim. Let $\epsilon=F(z_{k+1})-F(z_{k+\frac{1}{2}})$. By Claim 2 we have $\|z_{k+\frac{1}{2}}-z_{k+1}\|\leq \eta L\|z_k-z_{k+\frac{1}{2}}\|$ and by the L-Lipschitzness of operator F we have

$$\|\epsilon\| = \|F(z_{k+1}) - F(z_{k+\frac{1}{2}})\| \le L\|z_{k+1} - z_{k+\frac{1}{2}}\| \le \eta L^2\|z_k - z_{k+\frac{1}{2}}\|.$$
(27)

Note that

$$0 = \langle a_{\perp}, a \rangle = \langle a_{\perp}, z_{k+1} - z_k + \eta F(z_{k+\frac{1}{2}}) \rangle$$

$$\Leftrightarrow \langle a_{\perp}, F(z_{k+\frac{1}{2}}) \rangle = \frac{\langle a_{\perp}, z_k - z_{k+1} \rangle}{\eta}.$$
(28)

Moreover, the fact that $0 \ge \langle a, z_{k+1} - z_k + \eta F(z_{k+\frac{1}{2}}) \rangle \Leftrightarrow \frac{\langle a, z_k - z_{k+1} \rangle}{\eta} \ge \langle a, F(z_{k+\frac{1}{2}}) \rangle$, which further imiplies that

$$\langle a, F(z_{k+1}) \rangle = \langle a, F(z_{k+\frac{1}{2}}) + \epsilon \rangle \le \langle a, \frac{z_k - z_{k+1}}{\eta} \rangle + \langle a, \epsilon \rangle.$$
 (29)

Combining Equation (26), Equation (28) and Equation (29) we have

$$r^{tan}(z_{k+1}) \leq \sqrt{\langle a_{\perp}, F(z_{k+1}) \rangle^{2} + \langle a, F(z_{k+1}) \rangle^{2} \cdot \mathbb{1}[\langle a, F(z_{k+1}) \rangle > 0]}$$

$$\leq \sqrt{\langle a_{\perp}, \frac{z_{k} - z_{k+1}}{\eta} + \epsilon \rangle^{2} + \langle a, \frac{z_{k} - z_{k+1}}{\eta} + \epsilon \rangle^{2}}$$

$$\leq \left\| \frac{z_{k} - z_{k+1}}{\eta} + \epsilon \right\|$$

$$\leq \frac{\|z_{k} - z_{k+1}\|}{\eta} + \|\epsilon\|$$

$$\leq \frac{\|z_{k} - z_{k+1}\| + (\eta L)^{2} \|z_{k} - z_{k+\frac{1}{2}}\|}{\eta}$$

$$\leq \frac{\|z_{k} - z_{k+1}\| + (\eta L)^{2} \|z_{k} - z_{k+\frac{1}{2}}\|}{\eta}$$

$$\leq \frac{\|z_{k} - z_{k+\frac{1}{2}}\| + \|z_{k+\frac{1}{2}} - z_{k+1}\| + (\eta L)^{2} \|z_{k} - z_{k+\frac{1}{2}}\|}{\eta}$$

$$\leq (1 + \eta L + (\eta L)^{2}) \frac{\|z_{k} - z_{k+\frac{1}{2}}\|}{\eta}.$$
(30)

The first inequality is because (i) $\langle a_{\perp}, F(z_{k+1}) \rangle = \langle a_{\perp}, \frac{z_k - z_{k+1}}{\eta} + \epsilon \rangle$ and (ii) $\langle a, F(z_{k+1}) \rangle^2 \cdot \mathbb{1}[\langle a, F(z_{k+1}) \rangle > 0] \leq \langle a, \frac{z_k - z_{k+1}}{\eta} + \epsilon \rangle^2$. The second inequality is because a and a_{\perp} are both unit vectors, and they are orthogonal to each other when $\|F(z_{k+1}) - \langle a, F(z_{k+1}) \rangle \cdot a\| \neq 0$, and because a_{\perp} is $(0, \ldots, 0)$ otherwise. The third and the fifth inequality follow from the triangle inequality. The fourth inequality follows from Equation (27). In the final inequality we use $||z_{k+\frac{1}{2}} - z_{k+1}|| \leq \eta L ||z_k - z_{k+\frac{1}{2}}||$ by Fact 2.

Proof of Lemma 5: By Lemma 3 we have

$$\|z_0 - z^*\|^2 \ge \|z_{T+1} - z^*\|^2 + (1 - \eta^2 L^2) \sum_{k=0}^{T} \|z_k - z_{k+\frac{1}{2}}\|^2 \ge (1 - \eta^2 L^2) \sum_{k=0}^{T} \|z_k - z_{k+\frac{1}{2}}\|^2$$

Thus there exists a $t^* \in [T]$ such that $||z_{t^*} - z_{t^* + \frac{1}{2}}||^2 \le \frac{||z_0 - z^*||^2}{T(1 - \eta^2 L^2)}$. We conclude the proof by applying Lemma 4.

C Missing Proofs from Section 5.1

Proposition 1.

$$||F(z_k)||^2 - ||F(z_{k+1})||^2 + 2 \cdot \left\langle F(z_{k+1}) - F(z_k), F(z_{k+\frac{1}{2}}) \right\rangle + \left(\left\| F(z_{k+\frac{1}{2}}) - F(z_{k+1}) \right\|^2 - \left\| F(z_{k+\frac{1}{2}}) - F(z_k) \right\|^2 \right) = 0.$$

Proof. Expanding the LHS of the equation in the statement we can verify that

$$\begin{split} \|F(z_k)\|^2 - \|F(z_{k+1})\|^2 + 2 \cdot \left\langle F(z_{k+1}), F(z_{k+\frac{1}{2}}) \right\rangle - 2 \cdot \left\langle F(z_k), F(z_{k+\frac{1}{2}}) \right\rangle \\ + \left\| F(z_{k+\frac{1}{2}}) \right\|^2 - 2 \cdot \left\langle F(z_{k+1}), F(z_{k+\frac{1}{2}}) \right\rangle + \left\| F(z_{k+1}) \right\|^2 \\ - \left\| F(z_{k+\frac{1}{2}}) \right\|^2 + 2 \cdot \left\langle F(z_k), F(z_{k+\frac{1}{2}}) \right\rangle - \left\| F(z_k) \right\|^2 = 0. \end{split}$$

D The SOS Program for the Constrained Setting

Interpretation and Notation.

- Interpret $\overline{\alpha}$ as α , $\overline{\beta}_1$ as β_1 , and $\overline{\beta}_2$ as β_2 . We use \overline{a}_0 to denote $(\overline{\beta}_1, \overline{\beta}_2, 1)$, \overline{a}_1 to denote $(\overline{\alpha}, 1, 0)$, and \overline{a}_2 to denote (1, 0, 0).
- Interpret $\overline{z}_i[\ell]$ as $\overline{z}_{k+\frac{i}{2}}[\ell]$ and $\overline{F}_i[\ell]$ as $\eta \overline{F}(\overline{z}_{k+\frac{i}{2}})[\ell]$ for $0 \le i \le 2$ and $\ell \le 3$. We use \overline{z}_i to denote $(\overline{z}_i[1], \overline{z}_i[2], \overline{z}_i[3])$ and \overline{F}_i to denote $(\overline{F}_i[1], \overline{F}_i[2], \overline{F}_i[3])$.
- Interpret x_i and y_i as the representatives for $\overline{z}_{k+\frac{i}{2}}[\ell]$ and $\eta \overline{F}(\overline{z}_{k+\frac{i}{2}})[\ell]$ for all $\ell \geq 4$, for $0 \leq i \leq 2$.
- We use x to denote $(\overline{z}_0, \overline{z}_1, \overline{z}_2, x_0, x_1, x_2, \overline{\alpha}, \overline{\beta}_1, \overline{\beta}_2)$ and y to denote $(F_0, F_1, F_2, y_0, y_1, y_2)$.

Input Fixed Polynomials.

- Polynomial $g(x,y) = \|\overline{a}_0\|^2 \cdot (\|\overline{F}_0 \frac{\langle \overline{F}_0, \overline{a}_0 \rangle \cdot \overline{a}_0}{\|\overline{a}_0\|^2}\|^2 + y_0^2 \|\overline{F}_2 \langle \overline{F}_2, \overline{a}_2 \rangle \cdot \overline{a}_2\|^2 y_2^2)$
- Set of polynomials

$$\begin{split} S^{(=0)} &= \left\{ \langle \overline{a}_0, \overline{z}_0 \rangle, \langle \overline{a}_1, \overline{z}_1 \rangle, \langle \overline{a}_2, \overline{z}_2 \rangle \right\} \\ &\quad \cup \left\{ \langle (1, -\overline{\alpha}, 0), \overline{z}_1 - \overline{z}_0 + \overline{F}_0 \rangle \right\} \\ &\quad \cup \left\{ \langle e_3, \overline{z}_1 - \overline{z}_0 + \overline{F}_0 \rangle, \langle e_2, \overline{z}_2 - \overline{z}_0 + \overline{F}_1 \rangle, \langle e_3, \overline{z}_2 - \overline{z}_0 + \overline{F}_1 \rangle \right\} \\ &\quad \cup \left\{ x_1 - x_0 + y_0, x_2 - x_0 + y_1 \right\}, \\ S^{(\geq 0)} &= \left\{ \langle \overline{a}_i, \overline{z}_j \rangle \right\}_{0 \leq i \neq j \leq 2} \\ &\quad \cup \left\{ \langle \overline{a}_0, \overline{F}_0 \rangle, \langle \overline{a}_1, \overline{z}_1 - \overline{z}_0 + \overline{F}_0 \rangle, \langle \overline{a}_2, \overline{z}_2 - \overline{z}_0 + \overline{F}_1 \rangle \right\}, \\ S^{(G, \leq 0)} &= \left\{ \langle \overline{F}_i - \overline{F}_j, \overline{z}_j - \overline{z}_i \rangle + (y_i - y_j) \cdot (x_j - x_i) \right\}_{i,j \in \{0,1,2\}} \\ &\quad \cup \left\{ \| \overline{F}_i - \overline{F}_j \|^2 - \| \overline{z}_i - \overline{z}_j \|^2 + (y_i - y_j)^2 - (x_i - x_j)^2 \right\}_{i,i \in \{0,1,2\}}. \end{split}$$

Inferred "Negative" Polynomials (Not part of Input).

• Let $S^{(\leq 0)} = \left\{ -h(x, y) \cdot h'(x, y) : h(x, y), h'(x, y) \in S^{(\geq 0)} \right\} \cup S^{(G, \leq 0)} \cup \{g(x, y)\}.$

Decision Variables of the SOS Program:

- $p_{(h(x,y))}(x,y) \in \mathbb{R}[x,y]$ is a polynomial with degree at most $d \deg(h(x,y))$, for all $h(x,y) \in S^{(=0)}$.
- $q_{(h(x,y))}(x,y) \in SOS[x,y]$ is an SOS polynomial with degree at most d deg(h(x,y)), for all $h(x,y) \in S(\le 0)$

Constraints of the SOS Program:

$$g(x,y) + \sum_{h(x,y) \in S^{(=0)}} p_{(h(x,y))}(x,y) \cdot h(x,y) + \sum_{h(x,y) \in S^{(\leq 0)}} q_{(h(x,y))}(x,y) \cdot h(x,y) \in SOS[x,y]$$

Figure 5: SOS program for the constrained case of degree d. We need d to be 8 to discover our **proof**. $S^{(=0)}$ contains all polynomials that are equal to 0 due to the constraints. $S^{(\geq 0)}$ contains all polynomials that only involve the first 3 coordinates and are non-negative due to the constraints. $S^{(G,\leq 0)}$ contains all polynomials that involve all coordinates and are non-negative. In our program, we need to include the inferred non-positive polynomials $S^{(\leq 0)}$.

E Missing Proofs from Section 5.2

Proof of Lemma 6: For our proof, it will be more convenient to work with terms $r^{tan}(z_k)^2$, $r^{tan}(z_{k+1})^2$ rather than $r^{tan}(z_k)$, $r^{tan}(z_{k+1})$. Since $r^{tan}(z_k)$, $r^{tan}(z_{k+1}) \ge 0$, then $r^{tan}(z_k) - r^{tan}(z_{k+1}) \ge 0$ iff $r^{tan}(z_k)^2 - r^{tan}(z_{k+1})^2 \ge 0$. For the rest of the proof, we refer to the property that $\overline{a}_k = (\beta_1, \beta_2, 1, 0, \dots, 0)$, $\overline{a}_{k+\frac{1}{2}} = (\alpha, 1, 0, \dots, 0)$, $\overline{a}_{k+1} = (1, 0, \dots, 0)$ as the form property.

Recall that the k-th update of EG is as follows $z_{k+\frac{1}{2}} = \Pi_{\mathcal{Z}}[z_k - \eta F(z_k)]$ and $z_{k+1} = \Pi_{\mathcal{Z}}[z_k - \eta F(z_{k+\frac{1}{2}})]$. We define the following vectors:

$$-a_k \in \underset{\substack{a \in \hat{N}_{\mathcal{Z}}(z_k), \\ \langle F(z_k), a \rangle \leq 0}}{\operatorname{argmin}} \|F(z_k) - \langle F(z_k), a \rangle \cdot a\|^2, \tag{31}$$

$$a_{k+\frac{1}{2}} = z_{k+\frac{1}{2}} - z_k + \eta F(z_k), \tag{32}$$

$$a_{k+1} = z_{k+1} - z_k + \eta F(z_{k+\frac{1}{2}}).$$
(33)

For now, let us assume that a_k , $a_{k+\frac{1}{2}}$, and a_{k+1} satisfy (i) the form property, and (ii) $\langle a_k, z_k \rangle = \langle a_{k+\frac{1}{2}}, z_{k+\frac{1}{2}} \rangle = \langle a_{k+1}, z_{k+1} \rangle = 0$. We use this simple case as the basis of our construction, and will remove these assumptions later.

We set $\overline{a}_k = a_k$, $\overline{a}_{k+\frac{1}{2}} = a_{k+\frac{1}{2}}$, $\overline{a}_{k+1} = a_{k+1}$, $\overline{z}_k = z_k$, $\overline{z}_{k+\frac{1}{2}} = z_{k+\frac{1}{2}}$, $\overline{z}_{k+1} = z_{k+1}$, $\overline{F}(\overline{z}_k) = F(z_k)$, $\overline{F}(\overline{z}_{k+\frac{1}{2}}) = F(z_{k+\frac{1}{2}})$ and $\overline{F}(\overline{z}_{k+1}) = F(z_{k+1})$. We first argue that Property 4 holds. Since $F(\cdot)$ is monotone and L-Lipschitz, Inequality (13) and (14) are satisfied. In addition, Inequality (15) is satisfied due to the definition of a_k .

Next, we show that property 3 holds. By the definition of $\overline{a}_{k+\frac{1}{2}}$ (or \overline{a}_{k+1}), it is clear that $\overline{a}_{k+\frac{1}{2}}$ (or \overline{a}_{k+1}) and $\overline{z}_{k+\frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_k)$ (or $\overline{z}_{k+1} - \overline{z}_k + \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})$) are co-directed. As $-a_k \in \widehat{N}_{\mathcal{Z}}(z_k)$,

$$\langle \overline{a}_k, \overline{z} \rangle \ge \langle \overline{a}_k, \overline{z}_k \rangle = 0, \quad \overline{z} \in \{\overline{z}_k, \overline{z}_{k+\frac{1}{2}}, \overline{z}_{k+1}\}.$$
 (34)

According to the update rule of the EG algorithm (Equation (4) and (5)), Equation (32), and Equation (33), we know that for all $z \in \mathcal{Z}$,

$$\left\langle a_{k+\frac{1}{2}}, z - z_{k+\frac{1}{2}} \right\rangle = \left\langle z_{k+\frac{1}{2}} - z_k + \eta F(z_k), z - z_{k+\frac{1}{2}} \right\rangle \ge 0,$$
 (35)

$$\langle a_{k+1}, z - z_{k+1} \rangle = \langle z_{k+1} - z_k + \eta F(z_{k+\frac{1}{2}}), z - z_{k+1} \rangle \ge 0,$$
 (36)

which implies that for any $i \in \{k + \frac{1}{2}, k + 1\}, j \in \{k, k + \frac{1}{2}, k + 1\}, \langle \overline{a}_i, \overline{z}_i \rangle \geq \langle \overline{a}_i, \overline{z}_i \rangle = 0.$

Finally, we verify Property 2. By Equation (36), $-\frac{a_{k+1}}{\|a_{k+1}\|} \in \widehat{N}(z_{k+1})$, which in combination with

Lemma 7 implies

$$\begin{split} & \left\| \overline{F}(\overline{z}_{k+1}) - \frac{\langle \overline{F}(\overline{z}_{k+1}), \overline{a}_{k+1} \rangle \cdot \overline{a}_{k+1}}{\left\| \overline{a}_{k+1} \right\|^2} \mathbb{1}[\langle \overline{F}(\overline{z}_{k+1}), \overline{a}_{k+1} \rangle \geq 0] \right\|^2 \\ & \geq \min_{\substack{a \in \widehat{N}_{\mathcal{Z}}(z_{k+1}), \\ \langle F(z_{k+1}), a \rangle \leq 0}} \|F(z_{k+1}) - \langle F(z_{k+1}), a \rangle \cdot a \|^2 \\ & = r^{tan}(z_{k+1})^2. \end{split}$$

According to Lemma 7 and Equation (31), we know that $r^{tan}(z_k)^2 = \|F(z_k) - \langle F(z_k), a_k \rangle \cdot a_k\|^2 = \|\overline{F}(\overline{z}_k) - \frac{\langle \overline{F}(\overline{z}_k), \overline{a}_k \rangle \cdot \overline{a}_k}{\|\overline{a}_k\|^2}\|^2$. If $r^{tan}(z_k)^2 - r^{tan}(z_{k+1})^2 < 0$, then

$$0 > r^{tan}(z_{k})^{2} - r^{tan}(z_{k+1})^{2}$$

$$\geq \left\| \overline{F}(\overline{z}_{k}) - \frac{\langle \overline{F}(\overline{z}_{k}), \overline{a}_{k} \rangle \cdot \overline{a}_{k}}{\|\overline{a}_{k}\|^{2}} \right\|^{2} - \left\| \overline{F}(\overline{z}_{k+1}) - \frac{\langle \overline{F}(\overline{z}_{k+1}), \overline{a}_{k+1} \rangle \cdot \overline{a}_{k+1}}{\|\overline{a}_{k+1}\|^{2}} \mathbb{1}[\langle \overline{F}(\overline{z}_{k+1}), \overline{a}_{k+1} \rangle \geq 0] \right\|^{2}.$$
(37)

Observe that when vectors a_k and a_{k+1} satisfy the form property, then $\|\overline{a}_k\| = 1$, and $\|\overline{a}_{k+1}\| \ge 1$ and Equation (37) is well-defined.

This completes the proof for the case, where (i) vectors a_k , $a_{k+\frac{1}{2}}$, and a_{k+1} satisfy the form property, and (ii) $\langle a_k, z_k \rangle = \langle a_{k+\frac{1}{2}}, z_{k+\frac{1}{2}} \rangle = \langle a_{k+1}, z_{k+1} \rangle = 0$. Our next step is to remove the assumptions. We first show how to modify the construction so that for any a_k , $a_{k+\frac{1}{2}}$, a_{k+1} , we can construct vectors \hat{a}_k , $\hat{a}_{k+\frac{1}{2}}$, \hat{a}_{k+1} , \hat{z}_k , $\hat{z}_{k+\frac{1}{2}}$, \hat{z}_{k+1} , $\hat{F}(\hat{z}_k)$, $\hat{F}(\hat{z}_{k+\frac{1}{2}})$, $\hat{F}(\hat{z}_{k+1}) \in \mathbb{R}^{n+5}$, that satisfy, among other properties, (a) $\langle \hat{a}_k, \hat{z}_k \rangle = \langle \hat{a}_{k+\frac{1}{2}}, \hat{z}_{k+\frac{1}{2}} \rangle = \langle \hat{a}_{k+1}, \hat{z}_{k+1} \rangle = 0$, and (b) vectors \hat{a}_k , $\hat{a}_{k+\frac{1}{2}}$, \hat{a}_{k+1} are linear independent. In our final step, we choose a proper basis to construct vectors \overline{a}_k , $\overline{a}_{k+\frac{1}{2}}$, \overline{a}_{k+1} , \overline{z}_k , $\overline{z}_{k+\frac{1}{2}}$, \overline{z}_{k+1} , $\overline{F}(\overline{z}_k)$, $\overline{F}(\overline{z}_{k+\frac{1}{2}})$, $\overline{F}(\overline{z}_{k+1}) \in \mathbb{R}^{n+5}$ that satisfy all four properties in the statement of Lemma 6.

We now present the construction of vectors \hat{a}_k , $\hat{a}_{k+\frac{1}{2}}$, \hat{a}_{k+1} , \hat{z}_k , $\hat{z}_{k+\frac{1}{2}}$, \hat{z}_{k+1} , $\widehat{F}(\hat{z}_k)$, $\widehat{F}(\hat{z}_{k+\frac{1}{2}})$, $\widehat{F}(\hat{z}_{k+1})$. High-levelly speaking, we introduce five dummy dimensions. The purpose of the first dummy dimension is to ensure property (a). We use the remaining four dummy dimensions to ensure that the newly created vectors \hat{a}_k , $\hat{a}_{k+\frac{1}{2}}$ and \hat{a}_{k+1} are linearly independent satisfying property (b). More specifically for parameters ℓ , $\epsilon > 0$ that we determine later, we define \hat{z}_k , $\hat{z}_{k+\frac{1}{2}}$, \hat{z}_{k+1} , $\widehat{F}(\hat{z}_k)$, $\widehat{F}(\hat{z}_{k+\frac{1}{2}})$,

 $\widehat{F}(\widehat{z}_{k+1})$, \widehat{a}_k , $\widehat{a}_{k+\frac{1}{2}}$, \widehat{a}_{k+1} as follows

$$\hat{z}_i := (-\epsilon^{-1}, 0, 0, 0, 0, z_i) \qquad \forall i \in \{k, k + \frac{1}{2}, k + 1\}$$
(38)

$$\widehat{F}(\widehat{z}_k) := \left(\frac{\epsilon}{\eta} \cdot \left\langle a_{k+\frac{1}{2}}, z_{k+\frac{1}{2}} \right\rangle, 0, \frac{\epsilon}{\eta}, 0, \frac{\ell \epsilon}{\eta}, F(z_k)\right) \tag{39}$$

$$\widehat{F}(\widehat{z}_{k+\frac{1}{2}}) := (\frac{\epsilon}{\eta} \cdot \langle a_{k+1}, z_{k+1} \rangle, 0, 0, \frac{\epsilon}{\eta}, \frac{\ell \epsilon}{\eta}, F(z_{k+\frac{1}{2}}))$$
(40)

$$\widehat{F}(\widehat{z}_{k+1}) := \left(\frac{\epsilon}{\eta} \cdot \langle a_{k+1}, z_{k+1} \rangle, 0, 0, \frac{\epsilon}{\eta}, \frac{\ell \epsilon}{\eta}, F(z_{k+1})\right) \tag{41}$$

$$\hat{a}_k := (\epsilon \cdot \langle a_k, z_k \rangle, \epsilon, 0, 0, \ell \epsilon, a_k) \tag{42}$$

$$\hat{a}_{k+\frac{1}{2}} := \left(\epsilon \cdot \left\langle a_{k+\frac{1}{2}}, z_{k+\frac{1}{2}} \right\rangle, 0, \epsilon, 0, \ell \epsilon, a_{k+\frac{1}{2}}\right) \tag{43}$$

$$\hat{a}_{k+1} := (\epsilon \cdot \langle a_{k+1}, z_{k+1} \rangle, 0, 0, \epsilon, \ell \epsilon, a_{k+1}) \tag{44}$$

Clearly, \hat{a}_k , $\hat{a}_{k+\frac{1}{2}}$, \hat{a}_{k+1} are linear independent, \hat{a}_{k+1} and $\hat{z}_{k+1} - \hat{z}_k + \eta \widehat{F}(\hat{z}_{k+\frac{1}{2}})$ are co-directed, and $\hat{a}_{k+\frac{1}{2}}$ and $\hat{z}_{k+\frac{1}{2}} - \hat{z}_k + \eta \widehat{F}(\hat{z}_k)$ are co-directed. Note that the following inequalites hold. By Equation (31)-(33), it is clear that $-a_k \in \widehat{N}(z_k) \subseteq N(z_k)$, $-a_{k+\frac{1}{2}} \in N(z_{k+\frac{1}{2}})$ and $-a_{k+1} \in N(z_{k+1})$, which further implies

$$\langle \hat{a}_i, \hat{z}_j \rangle = \langle a_i, z_j \rangle - \langle a_i, z_i \rangle \ge 0, \qquad \forall i, j \in \{k, k + \frac{1}{2}, k + 1\}$$

$$(45)$$

$$\langle \hat{a}_i, \hat{z}_i \rangle = 0, \qquad \forall i \in \{k, k + \frac{1}{2}, k + 1\}$$

$$(46)$$

$$\left\|\widehat{F}(\hat{z}_{k+1}) - \widehat{F}(\hat{z}_{k+\frac{1}{2}})\right\|^2 = \left\|F(z_{k+1}) - F(z_{k+\frac{1}{2}})\right\|^2 \le L^2 \left\|z_{k+1} - z_{k+\frac{1}{2}}\right\|^2 = \left\|\hat{z}_{k+1} - \hat{z}_{k+\frac{1}{2}}\right\|^2 \tag{47}$$

$$\left\langle \widehat{F}(\widehat{z}_{k+1}) - \widehat{F}(\widehat{z}_k), \widehat{z}_{k+1} - \widehat{z}_k \right\rangle = \left\langle F(z_{k+1}) - F(z_k), z_{k+1} - z_k \right\rangle \ge 0 \tag{48}$$

Moreover, $\langle \hat{a}_k, \hat{F}(\hat{z}_k) \rangle = \frac{\epsilon^2}{\eta} \cdot \left(\langle a_k, z_k \rangle \langle a_{k+\frac{1}{2}}, z_{k+\frac{1}{2}} \rangle + \ell^2 \right) + \langle a_k, F(z_k) \rangle$. We choose ℓ to be sufficiently large so that $\langle a_k, z_k \rangle \langle a_{k+\frac{1}{2}}, z_{k+\frac{1}{2}} \rangle + \ell^2 \geq 0$. Hence, for our choice of ℓ ,

$$\langle \hat{a}_k, \hat{F}(\hat{z}_k) \rangle \ge \langle a_k, F(z_k) \rangle \ge 0.$$
 (49)

We define function

$$\begin{split} \widehat{H}(\widehat{z}_k) &:= \left\| \widehat{F}(\widehat{z}_k) - \frac{\langle \widehat{F}(\widehat{z}_k), \widehat{a}_k \rangle}{\left\| \widehat{a}_k \right\|^2} \cdot \widehat{a}_k \right\|^2 \\ &= \left\| \widehat{F}(\widehat{z}_k) - \frac{\frac{\epsilon^2}{\eta} \cdot \left(\langle a_k, z_k \rangle \langle a_{k+\frac{1}{2}}, z_{k+\frac{1}{2}} \rangle + \ell^2 \right) + \langle a_k, F(z_k) \rangle}{\epsilon^2 (\langle a_k, z_k \rangle^2 + 1 + \ell^2) + \left\| a_k \right\|^2} \cdot \widehat{a}_k \right\|^2. \end{split}$$

Define
$$f(\epsilon) := \frac{\frac{\epsilon^2}{\eta} \cdot \left(\langle a_k, z_k \rangle \langle a_{k+\frac{1}{2}, z_{k+\frac{1}{2}}} \rangle + \ell^2 \right) + \langle a_k, F(z_k) \rangle}{\epsilon^2 (\langle a_k, z_k \rangle^2 + 1 + \ell^2) + \|a_k\|^2}$$
. We can simplify $\widehat{H}(\widehat{z}_k)$ to be

$$\epsilon^{2} \cdot \left\| \left(\frac{\langle a_{k+\frac{1}{2}}, z_{k+\frac{1}{2}} \rangle}{\eta} - f(\epsilon) \langle a_{k}, z_{k} \rangle, -f(\epsilon), \frac{1}{\eta}, 0, \frac{\ell}{\eta} - f(\epsilon) \ell \right) \right\|^{2} + \|F(z_{k}) - f(\epsilon) \cdot a_{k}\|^{2}.$$

When $a_k = (0, ..., 0)$, $r^{tan}(z_k)^2 = \|F(z_k)\|^2$ (Equation (31)) and $f(\epsilon) = \frac{\ell^2}{\eta(1+\ell^2)}$ for any $\epsilon > 0$. Therefore, $\lim_{\epsilon \to 0^+} \widehat{H}(\hat{z}_k) = \|F(z_k)\|^2 = r^{tan}(z_k)^2$, when $a_k = (0, ..., 0)$. When $a_k \neq (0, ..., 0)$, $r^{tan}(z_k)^2 = \|F(z_k) - \frac{\langle a_k, F(z_k) \rangle}{\|a_k\|^2} \cdot a_k\|^2$ (Equation (31)) and $\lim_{\epsilon \to 0^+} f(\epsilon) = \frac{\langle a_k, F(z_k) \rangle}{\|a_k\|^2}$, so $\lim_{\epsilon \to 0^+} \widehat{H}(\hat{z}_k) = \|F(z_k) - \frac{\langle a_k, F(z_k) \rangle}{\|a_k\|^2} \cdot a_k\|^2 = r^{tan}(z_k)^2$.

Similarly, we define function

$$\begin{split} \widehat{H}(\widehat{z}_{k+1}) &:= \left\| \widehat{F}(\widehat{z}_{k+1}) - \frac{\langle \widehat{F}(\widehat{z}_{k+1}), \widehat{a}_{k+1} \rangle \cdot \widehat{a}_{k+1}}{\|\widehat{a}_{k+1}\|^2} \mathbb{1} \left[\left\langle \widehat{F}(\widehat{z}_{k+1}), \widehat{a}_{k+1} \right\rangle \ge 0 \right] \right\|^2 \\ &= \left\| \widehat{F}(\widehat{z}_{k+1}) - \frac{\frac{\epsilon^2}{\eta} \cdot \left(\left\langle a_{k+1}, z_{k+1} \right\rangle^2 + 1 + \ell^2 \right) + \left\langle F(z_{k+1}), a_{k+1} \right\rangle}{\epsilon^2 (\left\langle a_{k+1}, z_{k+1} \right\rangle^2 + 1 + \ell^2) + \left\| a_{k+1} \right\|^2} \cdot \widehat{a}_{k+1} \mathbb{1} \left[\left\langle \widehat{F}(\widehat{z}_{k+1}), \widehat{a}_{k+1} \right\rangle \ge 0 \right] \right\|^2. \end{split}$$

Define $g(\epsilon) := \frac{\frac{\epsilon^2}{\eta} \cdot \left(\langle a_{k+1}, z_{k+1} \rangle^2 + 1 + \ell^2 \right) + \langle F(z_{k+1}), a_{k+1} \rangle}{\epsilon^2 (\langle a_{k+1}, z_{k+1} \rangle^2 + 1 + \ell^2) + \|a_{k+1}\|^2}$, and we can simplify $\widehat{H}(\widehat{z}_{k+1})$ to be

$$\epsilon^{2} \left(\frac{1}{\eta} - g(\epsilon) \mathbb{1} \left[\left\langle \widehat{F}(\widehat{z}_{k+1}), \widehat{a}_{k+1} \right\rangle \ge 0 \right] \right)^{2} \cdot \left\| \left(\left\langle a_{k+1}, z_{k+1} \right\rangle, 0, 0, 1, \ell \right) \right\|^{2} \\
+ \left\| F(z_{k+1}) - g(\epsilon) \cdot a_{k+1} \mathbb{1} \left[\left\langle \widehat{F}(\widehat{z}_{k+1}), \widehat{a}_{k+1} \right\rangle \ge 0 \right] \right\|^{2}.$$

When $a_{k+1} = (0, ..., 0)$, we have $g(\epsilon) = \frac{1}{\eta}$ and $\langle \widehat{F}(\hat{z}_{k+1}), \hat{a}_{k+1} \rangle \geq 0$, so $\widehat{H}(\hat{z}_{k+1}) = \|F(z_{k+1})\|^2 \geq r^{tan}(z_{k+1})^2$. When $a_{k+1} \neq (0, ..., 0)$, we have $\lim_{\epsilon \to 0^+} g(\epsilon) = \frac{\langle F(z_{k+1}), a_{k+1} \rangle}{\|a_{k+1}\|^2}$ and $\lim_{\epsilon \to 0^+} \mathbb{1}\left[\langle \widehat{F}(\hat{z}_{k+1}), \hat{a}_{k+1} \rangle \geq 0\right] = \mathbb{1}\left[\langle F(z_{k+1}), a_{k+1} \rangle \geq 0\right]$, hence

$$\lim_{\epsilon \to 0^+} \widehat{H}(\hat{z}_{k+1}) = \left\| F(z_{k+1}) - \frac{\langle F(z_{k+1}), a_{k+1} \rangle}{\|a_{k+1}\|^2} \cdot a_{k+1} \mathbb{1} \left[\langle F(z_{k+1}), a_{k+1} \rangle \ge 0 \right] \right\|^2 \ge r^{tan} (z_{k+1})^2.$$

The last inequality is because $\frac{-a_{k+1}}{\|a_{k+1}\|} \in \widehat{N}_{\mathcal{Z}}(z_{k+1})$ and Lemma 7.

If
$$r^{tan}(z_k)^2 - r^{tan}(z_{k+1})^2 < 0$$
, then

$$\lim_{\epsilon \to 0^+} \left(\widehat{H}(\widehat{z}_k) - \widehat{H}(\widehat{z}_{k+1}) \right) \le r^{tan} (z_k)^2 - r^{tan} (z_{k+1})^2 < 0.$$

This is because $\langle \widehat{F}(\hat{z}_{k+1}), \hat{a}_{k+1} \rangle$ is never smaller than $\langle F(z_{k+1}), a_{k+1} \rangle$, and the function $\mathbb{1}[x \geq 0]$ is right continuous.

Thus, we can choose a sufficiently small ϵ so that $\widehat{H}(\hat{z}_k) - \widehat{H}(\hat{z}_{k+1}) < 0$. Together with Inequalities (45)-(49), we can satisfy all properties excluding the form property using \hat{z}_k , $\hat{z}_{k+\frac{1}{2}}$, \hat{z}_{k+1} , \hat{a}_k , $\hat{a}_{k+\frac{1}{3}}$, \hat{a}_{k+1} , $\hat{F}(\hat{z}_k)$, $\hat{F}(\hat{z}_{k+\frac{1}{3}})$, $\hat{F}(\hat{z}_{k+1})$.

Now we show how to make the vectors also satisfy the form property. We perform a change of basis so that vectors \hat{a}_k , $\hat{a}_{k+\frac{1}{2}}$ and \hat{a}_{k+1} only depend on the first three coordinates. We use the Gram–Schmidt process to generate a basis, where vectors \hat{a}_k , $\hat{a}_{k+\frac{1}{2}}$, \hat{a}_{k+1} all lie in the span of the first three vector of the new basis. More formally, let N=n+5 and $\{b_i\}_{i\in[N]}$ be a sequence of orthonormal vectors produced by the Gram-Schmidt process on ordered input \hat{a}_{k+1} , $\hat{a}_{k+\frac{1}{2}}$, \hat{a}_k and $\{e_i\}_{i\in[N]}$. Let Q be the $N\times N$ matrix, where the i-th row of Q is vector b_i . Observe that any vector $z\in\mathbb{R}^N$ written in the basis $\{e_i\}_{i\in[N]}$ can be represented by the basis $\{b_i\}_{i\in[N]}$ with coefficients $Q\cdot z$.

Let $\overline{a}_k = Q \cdot \hat{a}_k$, $\overline{a}_{k+\frac{1}{2}} = Q \cdot \hat{a}_{k+\frac{1}{2}}$, $\overline{a}_{k+1} = Q \cdot \hat{a}_{k+1}$, $\overline{z}_k = Q \cdot \hat{z}_k$, $\overline{z}_{k+\frac{1}{2}} = Q \cdot \hat{z}_{k+\frac{1}{2}}$, $\overline{z}_{k+1} = Q \cdot \hat{z}_{k+1}$, $\overline{F}(\overline{z}_k) = Q \cdot \widehat{F}(\hat{z}_k)$, $\overline{F}(\overline{z}_{k+\frac{1}{2}}) = Q \cdot \widehat{F}(\hat{z}_{k+\frac{1}{2}})$ and $\overline{F}(\overline{z}_{k+1}) = Q \cdot \widehat{F}(\hat{z}_{k+1})$. Note that \overline{a}_k is the coefficients of \hat{a}_k written in the basis $\{b_i\}_{i \in [N]}$ and similar reasoning holds for the rest of the defined vectors. Clearly, for any $\hat{z}, \hat{z}' \in \mathbb{R}^N$, if we define $\overline{z} = Q \cdot \hat{z}$ and $\overline{z}' = Q \cdot \hat{z}'$, then we have

$$\langle \bar{z}, \bar{z}' \rangle = \langle \hat{z}, \hat{z}' \rangle. \tag{50}$$

Combining Equation (50) and the fact that all properties but the form property hold for vectors \hat{a}_k , $\hat{a}_{k+\frac{1}{2}}$, \hat{a}_{k+1} , \hat{z}_k , $\hat{z}_{k+\frac{1}{2}}$, \hat{z}_{k+1} , $\hat{F}(\hat{z}_k)$, $\hat{F}(\hat{z}_{k+\frac{1}{2}})$ and $\hat{F}(\hat{z}_{k+1})$, we conclude that the same properties also hold for vectors \bar{a}_k , $\bar{a}_{k+\frac{1}{2}}$, \bar{a}_{k+1} , \bar{z}_k , $\bar{z}_{k+\frac{1}{2}}$, \bar{z}_{k+1} , $\bar{F}(\bar{z}_k)$, $\bar{F}(\bar{z}_{k+\frac{1}{2}})$ and $\bar{F}(\bar{z}_{k+1})$.

Finally, by properties of the Gram-Schmidt process, the order of the vector in its input, and the fact that vectors \hat{a}_k , $\hat{a}_{k+\frac{1}{2}}$ and \hat{a}_{k+1} are linearly independent, we have $\hat{a}_{k+1} \in \operatorname{Span}(b_1)$, $\hat{a}_{k+\frac{1}{2}} \in \operatorname{Span}(b_1,b_2)$, $\hat{a}_k \in \operatorname{Span}(b_1,b_2,b_3)$, $\langle \hat{a}_{k+1},b_1 \rangle > 0$, $\langle \hat{a}_{k+\frac{1}{2}},b_2 \rangle > 0$ and $\langle \hat{a}_k,b_3 \rangle > 0$. Thus $\overline{a}_k = (\beta_1,\beta_2,b,0,\ldots,0)$, $\overline{a}_{k+\frac{1}{2}} = (\alpha,c,0,\ldots,0)$ and $\overline{a}_{k+1} = (d,0,\ldots,0)$, where $\beta_1,\beta_2,\alpha \in \mathbb{R}$ and b,c,d>0. By properly scaling \overline{a}_k , $\overline{a}_{k+\frac{1}{2}}$ and \overline{a}_{k+1} , we can make them satisfy the form property. This completes the proof.

Proof of Theorem 2: Assume towards contradiction that $r^{tan}(z_{k+1}) > r^{tan}(z_k)$, using Lemma 6 there exists numbers $\alpha, \beta_1, \beta_2 \in \mathbb{R}$ and vectors \overline{a}_k , $\overline{a}_{k+\frac{1}{2}}$, \overline{a}_{k+1} , \overline{z}_k , $\overline{z}_{k+\frac{1}{2}}$, \overline{z}_{k+1} , $\overline{F}(\overline{z}_k)$, $\overline{F}(\overline{z}_{k+\frac{1}{2}})$, $\overline{F}(\overline{z}_{k+\frac{1}{2}})$, where N = n+5 that satisfy the properties in the statement of Lemma 6 and

$$0 > \left\| \overline{F}(\overline{z}_{k}) - \frac{\langle F(\overline{z}_{k}), \overline{a}_{k} \rangle}{\|\overline{a}_{k}\|^{2}} \overline{a}_{k} \right\|^{2} - \left\| \overline{F}(\overline{z}_{k+1}) - \frac{\langle \overline{F}(\overline{z}_{k+1}), \overline{a}_{k+1} \rangle \overline{a}_{k+1}}{\|\overline{a}_{k+1}\|^{2}} \mathbb{1}[\langle \overline{F}(\overline{z}_{k+1}), \overline{a}_{k+1} \rangle \ge 0] \right\|^{2}$$

$$= \|\overline{F}(\overline{z}_{k})\|^{2} - \|\overline{F}(\overline{z}_{k+1})\|^{2} - \frac{(\beta_{1}\overline{F}(\overline{z}_{k})[1] + \beta_{2}\overline{F}(\overline{z}_{k})[2] + \overline{F}(\overline{z}_{k})[3])^{2}}{\beta_{1}^{2} + \beta_{2}^{2} + 1} + \overline{F}(\overline{z}_{k+1})[1]^{2} \mathbb{1}[\overline{F}(\overline{z}_{k+1})[1] \ge 0],$$

where we use the fact that $\bar{a}_k = (\beta_1, \beta_2, 1, 0, \dots, 0)$ and $\bar{a}_{k+1} = (1, 0, \dots, 0)$. We use TARGET to denote

$$\|\overline{F}(\overline{z}_{k})\|^{2} - \|\overline{F}(\overline{z}_{k+1})\|^{2} - \frac{(\beta_{1}\overline{F}(\overline{z}_{k})[1] + \beta_{2}\overline{F}(\overline{z}_{k})[2] + \overline{F}(\overline{z}_{k})[3])^{2}}{\beta_{1}^{2} + \beta_{2}^{2} + 1} + \overline{F}(\overline{z}_{k+1})[1]^{2}\mathbb{1}[\overline{F}(\overline{z}_{k+1})[1] \geq 0],$$

and our goal is to show that TARGET is non-negative, and thus reach a contradiction.

Our plan is to show that we can obtain a sum of quotients of SOS polynomials by adding non-positive terms to TARGET, which implies the non-negativity of TARGET. We add the non-positive terms in a few steps.

Combining Lemma 6 and the fact that $\eta > 0$ and $(\eta L)^2 \le 1$, we derive the following two inequalities:

$$\langle \eta \overline{F}(\overline{z}_{k+1}) - \eta \overline{F}(\overline{z}_k), \overline{z}_k - \overline{z}_{k+1} \rangle \le 0,$$
 (51)

$$\left\| \eta \overline{F}(\overline{z}_{k+1}) - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}}) \right\|^2 - \left\| \overline{z}_{k+1} - \overline{z}_{k+\frac{1}{2}} \right\|^2 \le 0.$$
 (52)

Equipped with these two inequalities, it is clear that

$$\eta^2 \cdot \text{TARGET} \ge \eta^2 \cdot \text{TARGET} + 2 \cdot \text{LHS of Inequality (51)} + \text{LHS of Inequality (52)}.$$
 (53)

Therefore, it is sufficient to show that the RHS of Inequality (53) is non-negative.

We take advantage of the sparsity of vectors \overline{a}_k , $\overline{a}_{k+\frac{1}{2}}$, \overline{a}_{k+1} by considering the following partition of [N]. We define $P_1 = \{1,2,3\}$ and $P_2 = \{4,5,\cdots,N\}$. For any vector $z \in \mathbb{R}^N$ and $j \in \{1,2\}$, we define $p_j(z) \in \mathbb{R}^n$ to be the vector such that $p_j(z)[i] = z[i]$ for $i \in P_j$ and $p_j(z)[i] = 0$ otherwise. We divide the RHS of Inequality (53) using the partition to Expression (54) and Expression (55):

$$\eta^{2} \| p_{2}(\overline{F}(\overline{z}_{k})) \|^{2} - \eta^{2} \| p_{2}(\overline{F}(\overline{z}_{k+1})) \|^{2} + 2\eta \langle p_{2}(\overline{F}(\overline{z}_{k+1})) - p_{2}(\overline{F}(\overline{z}_{k})), p_{2}(\overline{z}_{k}) - p_{2}(\overline{z}_{k+1}) \rangle
+ \eta^{2} \| p_{2}(\overline{F}(\overline{z}_{k+1})) - p_{2}(\overline{F}(\overline{z}_{k+\frac{1}{2}})) \|^{2} - \| p_{2}(\overline{z}_{k+1}) - p_{2}(\overline{z}_{k+\frac{1}{2}}) \|^{2},$$
(54)

and

$$\eta^{2} \| p_{1}(\overline{F}(\overline{z}_{k})) \|^{2} - \eta^{2} \| p_{1}(\overline{F}(\overline{z}_{k+1})) \|^{2} + 2\eta \langle p_{1}(\overline{F}(\overline{z}_{k+1})) - p_{1}(\overline{F}(\overline{z}_{k})), p_{1}(\overline{z}_{k}) - p_{1}(\overline{z}_{k+1}) \rangle
+ \eta^{2} \| p_{1}(\overline{F}(\overline{z}_{k+1})) - p_{1}(\overline{F}(\overline{z}_{k+\frac{1}{2}})) \|^{2} - \| p_{1}(\overline{z}_{k+1}) - p_{1}(\overline{z}_{k+\frac{1}{2}}) \|^{2}
- \frac{\eta^{2}}{\beta_{1}^{2} + \beta_{2}^{2} + 1} (\beta_{1}\overline{F}(\overline{z}_{k})[1] + \beta_{2}\overline{F}(\overline{z}_{k})[2] + \overline{F}(\overline{z}_{k})[3])^{2} + \eta^{2}\overline{F}(\overline{z}_{k+1})[1]^{2} \mathbb{1}[\overline{F}(\overline{z}_{k+1})[1] \ge 0].$$
(55)

First we show that the Expression (54) is non-negative. According to Lemma 6, $\bar{a}_{k+\frac{1}{2}}$ and $\bar{z}_{k+\frac{1}{2}} - \bar{z}_k + \eta \bar{F}(\bar{z}_k)$ are co-directed, and \bar{a}_{k+1} and $\bar{z}_{k+1} - \bar{z}_k + \eta \bar{F}(\bar{z}_{k+\frac{1}{2}})$ are co-directed. For any $i \in P_2$, we know $\bar{a}_{k+\frac{1}{2}}[i] = \bar{a}_{k+1}[i] = 0$ and thus we have

$$\begin{split} 0 &= \left\langle e_i, \overline{a}_{k+\frac{1}{2}} \right\rangle = \left\langle e_i, \overline{z}_{k+\frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_{k+1}) \right\rangle \Leftrightarrow \overline{z}_{k+\frac{1}{2}}[i] = \overline{z}_k[i] - \eta \overline{F}(\overline{z}_{k+1})[i], \\ 0 &= \left\langle e_i, \overline{a}_{k+1} \right\rangle = \left\langle e_i, \overline{z}_{k+1} - \overline{z}_k + \eta \overline{F}(\overline{z}_{k+1}) \right\rangle \Leftrightarrow \overline{z}_{k+1}[i] = \overline{z}_k[i] - \eta \overline{F}(\overline{z}_{k+1})[i], \end{split}$$

which implies that

$$p_2(\overline{z}_{k+\frac{1}{2}}) = p_2(\overline{z}_k) - \eta \cdot p_2(\overline{F}(\overline{z}_k)),$$

$$p_2(\overline{z}_{k+1}) = p_2(\overline{z}_k) - \eta \cdot p_2(\overline{F}(\overline{z}_{k+\frac{1}{k}})).$$

Intuitively, one can think of the coordinates in P_2 as the ones where the EG update is unconstrained. With the two new equalities, it is easy to verify that Expression (54) is always 0.

$$\begin{split} \eta^{2} & \| p_{2}(\overline{F}(\overline{z}_{k})) \|^{2} - \eta^{2} \| p_{2}(\overline{F}(\overline{z}_{k+1})) \|^{2} + 2\eta \left\langle p_{2}(\overline{F}(\overline{z}_{k+1})) - p_{2}(\overline{F}(\overline{z}_{k})), p_{2}(\overline{z}_{k}) - p_{2}(\overline{z}_{k+1}) \right\rangle \\ & + \eta^{2} \| p_{2}(\overline{F}(\overline{z}_{k+1})) - p_{2}(\overline{F}(\overline{z}_{k+\frac{1}{2}})) \|^{2} - \| p_{2}(\overline{z}_{k+1}) - p_{2}(\overline{z}_{k+\frac{1}{2}}) \|^{2} \\ = & \eta^{2} \| p_{2}(\overline{F}(\overline{z}_{k})) \|^{2} - \eta^{2} \| p_{2}(\overline{F}(\overline{z}_{k+1})) \|^{2} + 2\eta^{2} \left\langle p_{2}(\overline{F}(\overline{z}_{k+1})) - p_{2}(\overline{F}(\overline{z}_{k})), p_{2}(\overline{F}(\overline{z}_{k+\frac{1}{2}})) \right\rangle \\ & + \eta^{2} \| p_{2}(\overline{F}(\overline{z}_{k+1})) - p_{2}(\overline{F}(\overline{z}_{k+\frac{1}{2}})) \|^{2} - \eta^{2} \| p_{2}(\overline{F}(\overline{z}_{k}) - p_{2}(\overline{F}(\overline{z}_{k+\frac{1}{2}})) \|^{2} \\ = & 0. \end{split}$$

We now turn our attention to Expression (55) and show that it is non-negative. The analysis is more challenging here. We introduce the following six non-positive expressions, multiply each of them with a carefully chosen coefficient, then add them together with Expression (55). We finally verify that the sum is a sum of quotients of SOS polynomials implying the non-negativity of Expression (55). We believe it will be extremely challenging if not impossible for human beings to discover these non-positive expressions and their associated coefficients manually to complete this proof. We instead harness the power of the SOS programming to overcome the difficulty and make the discovery.

We first present the six non-positive expressions.

$$\overline{z}_{k+\frac{1}{2}}[1]\left(\left(\overline{z}_{k+\frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_k)\right)[1] - \alpha\left(\overline{z}_{k+\frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_k)\right)[2]\right) = 0, \tag{56}$$

$$\overline{z}_{k+1}[2]\left(\left(\overline{z}_{k+\frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_k)\right)[1] - \alpha\left(\overline{z}_{k+\frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_k)\right)[2]\right) = 0, \tag{57}$$

$$\left(\alpha\left(\overline{z}_{k}-\eta\overline{F}(\overline{z}_{k})\right)[1]+\left(\overline{z}_{k}-\eta\overline{F}(\overline{z}_{k})\right)[2]\right)\left(\alpha\overline{z}_{k+1}[1]+\overline{z}_{k+1}[2]\right)\leq0,\tag{58}$$

$$-\eta \left(\beta_1 \overline{F}(\overline{z}_k)[1] + \beta_2 \overline{F}(\overline{z}_k)[2] + \overline{F}(\overline{z}_k)[3]\right) \left(\beta_1 \overline{z}_{k+\frac{1}{2}}[1] + \beta_2 \overline{z}_{k+\frac{1}{2}}[2] + \overline{z}_{k+\frac{1}{2}}[3]\right) \le 0, \tag{59}$$

$$\overline{z}_k[1]\left(\overline{z}_k[1] - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[1]\right) \le 0, \tag{60}$$

$$-\eta \overline{F}(\overline{z}_{k+1})[1]\mathbb{1}\left[\overline{F}(\overline{z}_{k+1})[1] \le 0\right] \left(\overline{z}_{k}[1] - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[1]\right) \le 0.$$
 (61)

Equation (56) and Equation (57) follow from the combination of the fact that $\langle (1, -\alpha, 0, \cdots, 0), \overline{a}_{k+\frac{1}{2}} \rangle = 0$ and that $\overline{a}_{k+\frac{1}{2}}$ and $\overline{z}_{k+\frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_k)$ are co-directed:

$$\begin{split} &\left\langle (1, -\alpha, 0, \cdots, 0), \overline{z}_{k + \frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_k) \right\rangle = 0 \\ \Leftrightarrow &\left(\overline{z}_{k + \frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_k) \right) [1] - \alpha \left(\overline{z}_{k + \frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_k) \right) [2] = 0. \end{split}$$

Note that the LHS of Inequality (58) is equal to $\langle \overline{a}_{k+\frac{1}{2}}, \overline{z}_k - \eta \overline{F}(\overline{z}_k) \rangle \cdot \langle \overline{a}_{k+\frac{1}{2}}, \overline{z}_{k+1} \rangle$. Lemma 6 guarantees that $\langle \overline{a}_{k+\frac{1}{2}}, \overline{z}_{k+1} \rangle \geq 0$. Since $\overline{a}_{k+\frac{1}{2}}$ and $\overline{z}_k - \eta \overline{F}(\overline{z}_k) - \overline{z}_{k+\frac{1}{2}}$ are oppositely directed, and $\langle \overline{a}_{k+\frac{1}{2}}, \overline{z}_{k+\frac{1}{2}} \rangle = 0$, we have that

$$0 \ge \left\langle \overline{a}_{k+\frac{1}{2}}, \overline{z}_k - \eta \overline{F}(\overline{z}_k) - \overline{z}_{k+\frac{1}{2}} \right\rangle = \left\langle \overline{a}_{k+\frac{1}{2}}, \overline{z}_k - \eta \overline{F}(\overline{z}_k) \right\rangle.$$

Hence, $\langle \overline{a}_{k+\frac{1}{2}}, \overline{z}_k - \eta \overline{F}(\overline{z}_k) \rangle \cdot \langle \overline{a}_{k+\frac{1}{2}}, \overline{z}_{k+1} \rangle \leq 0$.

Observe that the LHS of Inequality (59) is equal to $-\eta \langle \overline{a}_k, \overline{F}(\overline{z}_k) \rangle \cdot \langle \overline{a}_k, \overline{z}_{k+\frac{1}{2}} \rangle$. Lemma 6 guarantees that $\langle \overline{a}_k, \overline{F}(\overline{z}_k) \rangle \geq 0$ and $\langle \overline{a}_k, \overline{z}_{k+\frac{1}{2}} \rangle \geq 0$.

Finally, we argue Inequality (60) and Inequality (61). Note that the LHS of Inequality (60) is equal to $\langle \overline{a}_{k+1}, \overline{z}_k \rangle \cdot \langle \overline{a}_{k+1}, \overline{z}_k - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}}) \rangle$. Since \overline{a}_{k+1} and $\overline{z}_k - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}}) - \overline{z}_{k+1}$ are oppositely directed, and $\langle \overline{a}_{k+1}, \overline{z}_{k+1} \rangle = 0$, we have that

$$0 \geq \left\langle \overline{a}_{k+1}, \overline{z}_k - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}}) - \overline{z}_{k+1} \right\rangle = \left\langle \overline{a}_{k+1}, \overline{z}_k - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}}) \right\rangle.$$

Clearly, $-\overline{F}(\overline{z}_{k+1})\mathbb{1}[\overline{F}(\overline{z}_{k+1})[1] \leq 0]$ is non-negative and $\overline{z}_k[1] = \langle \overline{a}_{k+1}, \overline{z}_k \rangle$ is also non-negative due to Lemma 6. Thus Inequality (60) and Inequality (61) hold.

Our next step is to show that the following is non-negative.

Expression (55) + 2 × (LHS of Equation (56) + LHS of Inequality (60) + LHS of Inequality (61))

+
$$\frac{2\alpha}{1+\alpha^2}$$
 × LHS of Equation (57) + $\frac{2}{1+\alpha^2}$ × LHS of Inequality (58) + $\frac{2}{1+\beta_1^2+\beta_2^2}$ × LHS of Inequality (59) (62)

We first simplify Expression (62), using the following relationship between the variables.

$$\bar{z}_k[3] = -\beta_1 \bar{z}_k[1] - \beta_2 \bar{z}_k[2],$$
(63)

$$\overline{z}_{k+\frac{1}{2}}[2] = -\alpha \overline{z}_{k+\frac{1}{2}}[1], \tag{64}$$

$$\overline{z}_{k+\frac{1}{2}}[3] = \overline{z}_k[3] - \eta \overline{F}(\overline{z}_k)[3], \tag{65}$$

$$\bar{z}_{k+1}[1] = 0,$$
 (66)

$$\bar{z}_{k+1}[2] = \bar{z}_k[2] - \eta \bar{F}(\bar{z}_{k+\frac{1}{2}})[2],$$
 (67)

$$\bar{z}_{k+1}[3] = \bar{z}_k[3] - \eta \bar{F}(\bar{z}_{k+\frac{1}{2}})[3].$$
 (68)

Equation (63), Equation (64), and Equation (66) follows by $\langle \overline{a}_i, \overline{z}_i \rangle = 0$ for $i \in \{k, k + \frac{1}{2}, k + 1\}$ due to Lemma 6. We know that (i) $\langle \overline{a}_{k+\frac{1}{2}}, e_3 \rangle = \langle \overline{a}_{k+1}, e_2 \rangle = \langle \overline{a}_{k+1}, e_3 \rangle = 0$ by the definition of $\overline{a}_{k+\frac{1}{2}}$ and \overline{a}_{k+1} , and (ii) $\overline{a}_{k+\frac{1}{2}}$ and $\overline{z}_{k+\frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_k)$ are co-directed, \overline{a}_{k+1} and $\overline{z}_{k+1} - \overline{z}_k + \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})$ are co-directed by Lemma 6, thus Equation (65), Equation (67), and Equation (68) follow from the combination of (i) and (ii).

We simplify Expression (62) by substituting $\overline{z}_k[3]$, $\overline{z}_{k+\frac{1}{2}}[2]$, $\overline{z}_{k+\frac{1}{2}}[3]$, $\overline{z}_{k+1}[1]$, $\overline{z}_{k+1}[2]$, and $\overline{z}_{k+1}[3]$ using Equations (63)-(68).

Expression (55) is equal to the sum of the following three parts.

The first part is

$$\eta^{2} \| p_{1}(\overline{F}(\overline{z}_{k})) \|^{2} - \eta^{2} \| p_{1}(\overline{F}(\overline{z}_{k+1})) \|^{2} - \frac{\eta^{2}}{\beta_{1}^{2} + \beta_{2}^{2} + 1} (\beta_{1} \overline{F}(\overline{z}_{k})[1] + \beta_{2} \overline{F}(\overline{z}_{k})[2] + \overline{F}(\overline{z}_{k})[3])^{2}
+ \eta^{2} \overline{F}(\overline{z}_{k+1})[1]^{2} \mathbb{I}[\overline{F}(\overline{z}_{k+1})[1] \ge 0]
= \sum_{i=1}^{3} (\eta^{2} \overline{F}(\overline{z}_{k})[i]^{2} - \eta^{2} \overline{F}(\overline{z}_{k+1})[i]^{2}) - \frac{\eta^{2}}{\beta_{1}^{2} + \beta_{2}^{2} + 1} (\beta_{1} \overline{F}(\overline{z}_{k})[1] + \beta_{2} \overline{F}(\overline{z}_{k})[2] + \overline{F}(\overline{z}_{k})[3])^{2}
+ \eta^{2} \overline{F}(\overline{z}_{k+1})[1]^{2} \mathbb{I}[\overline{F}(\overline{z}_{k+1})[1] \ge 0].$$
(69)

The second part is

$$2\eta \left\langle p_{1}(\overline{F}(\overline{z}_{k+1})) - p_{1}(\overline{F}(\overline{z}_{k})), p_{1}(\overline{z}_{k}) - p_{1}(\overline{z}_{k+1}) \right\rangle$$

$$= 2\eta \overline{z}_{k}[1] \left(\overline{F}(\overline{z}_{k+1})[1] - \overline{F}(\overline{z}_{k})[1] \right) + 2\eta^{2} \overline{F}(\overline{z}_{k+\frac{1}{2}})[2] \left(\overline{F}(\overline{z}_{k+1})[2] - \overline{F}(\overline{z}_{k})[2] \right)$$

$$+ 2\eta^{2} \overline{F}(\overline{z}_{k+\frac{1}{2}})[3] \left(\overline{F}(\overline{z}_{k+1})[3] - \overline{F}(\overline{z}_{k})[3] \right). \tag{70}$$

The third part is

$$\eta^{2} \| p_{1}(\overline{F}(\overline{z}_{k+1})) - p_{1}(\overline{F}(\overline{z}_{k+\frac{1}{2}})) \|^{2} - \| p_{1}(\overline{z}_{k+1}) - p_{1}(\overline{z}_{k+\frac{1}{2}}) \|^{2} \\
= \sum_{i=1}^{3} \eta^{2} \left(\overline{F}(\overline{z}_{k+1})[i] - \overline{F}(\overline{z}_{k+\frac{1}{2}})[i] \right)^{2} - \overline{z}_{k+\frac{1}{2}}[1]^{2} - \left(\overline{z}_{k}[2] - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[2] + \alpha \overline{z}_{k+\frac{1}{2}}[1] \right)^{2} \\
- \left(\eta \overline{F}(\overline{z}_{k})[3] - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[3] \right)^{2}.$$
(71)

 $2 \times$ LHS of Equation (56) is equal to

$$2\overline{z}_{k+\frac{1}{2}}[1]\Big((\overline{z}_{k+\frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_k))[1] - \alpha(\overline{z}_{k+\frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_k))[2]\Big)$$

$$= 2\overline{z}_{k+\frac{1}{2}}[1]\Big((\overline{z}_{k+\frac{1}{2}} - \overline{z}_k + \eta \overline{F}(\overline{z}_k))[1] - \alpha(-\alpha\overline{z}_{k+\frac{1}{2}}[1] - \overline{z}_k[2] + \eta \overline{F}(\overline{z}_k)[2])\Big). \tag{72}$$

 $\frac{2\alpha}{1+\alpha^2}$ × LHS of Equation (57) is equal to

$$\frac{2\alpha}{1+\alpha^{2}}\overline{z}_{k+1}[2]\left((\overline{z}_{k+\frac{1}{2}}-\overline{z}_{k}+\eta\overline{F}(\overline{z}_{k}))[1]-\alpha(\overline{z}_{k+\frac{1}{2}}-\overline{z}_{k}+\eta\overline{F}(\overline{z}_{k}))[2]\right) \\
=\frac{2\alpha}{1+\alpha^{2}}\left(\overline{z}_{k}[2]-\eta\overline{F}(\overline{z}_{k+\frac{1}{2}})[2]\right)\left((\overline{z}_{k+\frac{1}{2}}-\overline{z}_{k}+\eta\overline{F}(\overline{z}_{k}))[1]-\alpha\left(-\alpha\overline{z}_{k+\frac{1}{2}}[1]-\overline{z}_{k}[2]+\eta\overline{F}(\overline{z}_{k})[2]\right)\right).$$
(73)

 $\frac{2}{1+\alpha^2}$ × LHS of Inequality (58) is equal to

$$\frac{2}{1+\alpha^{2}} \left(\alpha \left(\overline{z}_{k} - \eta \overline{F}(\overline{z}_{k})\right)[1] + \left(\overline{z}_{k} - \eta \overline{F}(\overline{z}_{k})\right)[2]\right) \left(\alpha \overline{z}_{k+1}[1] + \overline{z}_{k+1}[2]\right) \\
= \frac{2}{1+\alpha^{2}} \left(\alpha \left(\overline{z}_{k} - \eta \overline{F}(\overline{z}_{k})\right)[1] + \left(\overline{z}_{k} - \eta \overline{F}(\overline{z}_{k})\right)[2]\right) \left(\overline{z}_{k}[2] - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[2]\right). \tag{74}$$

 $\frac{2}{1+\beta_1^2+\beta_2^2}$ × LHS of Inequality (59) is equal to

$$-\frac{2}{1+\beta_{1}^{2}+\beta_{2}^{2}}(\beta_{1}\eta\overline{F}(\overline{z}_{k})[1]+\beta_{2}\eta\overline{F}(\overline{z}_{k})[2]+\eta\overline{F}(\overline{z}_{k})[3])(\beta_{1}\overline{z}_{k+\frac{1}{2}}[1]+\beta_{2}\overline{z}_{k+\frac{1}{2}}[2]+\overline{z}_{k+\frac{1}{2}}[3])$$

$$=-\frac{2}{1+\beta_{1}^{2}+\beta_{2}^{2}}(\beta_{1}\eta\overline{F}(\overline{z}_{k})[1]+\beta_{2}\eta\overline{F}(\overline{z}_{k})[2]+\eta\overline{F}(\overline{z}_{k})[3])((\beta_{1}-\alpha\beta_{2})\overline{z}_{k+\frac{1}{2}}[1]+\overline{z}_{k}[3]-\eta\overline{F}(\overline{z}_{k})[3])$$

$$=-\frac{2}{1+\beta_{1}^{2}+\beta_{2}^{2}}(\eta\beta_{1}\overline{F}(\overline{z}_{k})[1]+\eta\beta_{2}\overline{F}(\overline{z}_{k})[2]+\eta\overline{F}(\overline{z}_{k})[3])$$

$$\cdot((\beta_{1}-\alpha\beta_{2})\overline{z}_{k+\frac{1}{2}}[1]-\beta_{1}\overline{z}_{k}[1]-\beta_{2}\overline{z}_{k}[2]-\eta\overline{F}(\overline{z}_{k})[3]). \tag{75}$$

 $2 \times$ LHS of Inequality (60) is equal to

$$2\overline{z}_{k}[1]\left(\overline{z}_{k}[1] - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[1]\right). \tag{76}$$

 $2 \times$ LHS of Inequality (61) is equal to

$$-2\eta \overline{F}(\overline{z}_{k+1})[1]\mathbb{1}\left[\overline{F}(\overline{z}_{k+1})[1] \le 0\right] \left(\overline{z}_{k}[1] - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[1]\right). \tag{77}$$

After the substitution, we need to argue that the sum of Expression (69) to (77) is a sum of quotients of SOS polynomials, which we prove by establishing the following identity.

Expression (69) + Expression (70) + Expression (71) + Expression (72) + Expression (73) + Expression (74) + Expression (75) + Expression (76) + Expression (77)

$$= \left(\overline{z}_{k}[1] - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[1] + \eta \overline{F}(\overline{z}_{k+1})[1] \cdot \mathbb{1}[\overline{F}(\overline{z}_{k+1})[1] \ge 0]\right)^{2}$$
(78)

$$+\frac{\left(\overline{z}_{k}[1] - \eta \overline{F}(\overline{z}_{k})[1] - \overline{z}_{k+\frac{1}{2}}[1]\right)^{2}}{1 + \beta_{1}^{2} + \beta_{2}^{2}}$$
(79)

$$+\frac{\left(\eta \overline{F}(\overline{z}_{k})[3] + \beta_{1}\overline{z}_{k}[1] + \beta_{2}\overline{z}_{k}[2] + (\alpha\beta_{2} - \beta_{1})\overline{z}_{k+\frac{1}{2}}[1]\right)^{2}}{1 + \beta_{1}^{2} + \beta_{2}^{2}}$$
(80)

$$+\frac{\left(\overline{z}_{k}[2]-\eta \overline{F}(\overline{z}_{k})[2]+\alpha \overline{z}_{k+\frac{1}{2}}[1]\right)^{2}}{1+\beta_{1}^{2}+\beta_{2}^{2}} \tag{81}$$

$$+\frac{\left(\beta_{1}(\overline{z}_{k}[2]-\eta \overline{F}(\overline{z}_{k})[2]+\alpha \overline{z}_{k+\frac{1}{2}}[1])-\beta_{2}(\overline{z}_{k}[1]-\eta \overline{F}(\overline{z}_{k})[1]-\overline{z}_{k+\frac{1}{2}}[1])\right)^{2}}{1+\beta_{1}^{2}+\beta_{2}^{2}}$$
(82)

 ≥ 0 .

We verify this identity by expanding both the LHS and RHS in Section F only for the case where $\overline{F}(\overline{z}_{k+1})[1] \ge 0$. Observe that only Expression (69), Expression (77) and Term (78) depend

on the sign of $\overline{F}(\overline{z}_{k+1})[1]$. It is sufficient for us to verify the case where $\overline{F}(\overline{z}_{k+1})[1] \geq 0$, as when $\overline{F}(\overline{z}_{k+1})[1] < 0$, we only need to subtract $\eta^2 \overline{F}(\overline{z}_{k+1})[1]^2 + 2\eta \overline{F}(\overline{z}_{k+1})[1](\overline{z}_k[1] - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[1])$ from both the LHS and the RHS, ¹⁶ and the identity still holds.

Hence, Expression (55) is non-negative. Combining with the non-negativity of Expression (54), we conclude that TARGET is non-negative. This completes the proof.

F Table Verification

In this section, we expand Expression (69)-(82) and verify the identity in the proof of Theorem 2. We expand Expression (69)-(82) in the tables below, In each table, rows that correspond to an expression contain the coefficients of that expression; rows that are marked by "Sum" contain the sum of the coefficients of the expressions above it.

	三/- \[a] - [a]	<u> </u>	〒/- \[a] - [a]	〒/- \[a] - [a]	
	$\eta \overline{F}(\overline{z}_k)[1] \overline{z}_k[1]$	$ \eta \overline{F}(\overline{z}_k)[1] \overline{z}_{k+\frac{1}{2}}[1] $	$ \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[1] \overline{z}_{k}[1] $	$ \eta \overline{F}(\overline{z}_{k+1})[1] \overline{z}_k[1] $	$ \eta \overline{F}(\overline{z}_k)[2] \overline{z}_k[1] $
Expression(69)	0	0	0	0	0
Expression(70)	-2	0	0	2	0
Expression(71)	0	0	0	0	0
Expression(72)	0	2	0	0	0
Expression(73)	0	0	0	0	0
Expression(74)	0	0	0	0	0
Expression(75)	$\frac{2{\beta_1}^2}{{\beta_1}^2 + {\beta_2}^2 + 1}$	$-\frac{\beta_1 (2 \beta_1 - 2 \alpha \beta_2)}{{\beta_1}^2 + {\beta_2}^2 + 1}$	0	0	$\frac{2\beta_1\beta_2}{{\beta_1}^2 + {\beta_2}^2 + 1}$
Expression(76)	0	0	-2	0	0
Sum	$-\frac{2{\beta_2}^2+2}{{\beta_1}^2+{\beta_2}^2+1}$	$\frac{2\beta_2^2 + 2\alpha\beta_1\beta_2 + 2}{\beta_1^2 + \beta_2^2 + 1}$	-2	2	$\frac{2\beta_1\beta_2}{{\beta_1}^2 + {\beta_2}^2 + 1}$
Expression(78)	0	0	-2	2	0
Expression(79)	$-rac{2}{eta_{1}^{2}+eta_{2}^{2}+1}$	$\frac{2}{\beta_1^2 + \beta_2^2 + 1}$	0	0	0
Expression(80)	0	0	0	0	0
Expression(81)	0	0	0	0	0
Expression(82)	$-\frac{2\beta_2^2}{\beta_1^2 + \beta_2^2 + 1}$	$\frac{2{\beta_2}^2 + 2\alpha\beta_1\beta_2}{{\beta_1}^2 + {\beta_2}^2 + 1}$	0	0	$\frac{2\beta_1\beta_2}{{\beta_1}^2 + {\beta_2}^2 + 1}$
Sum	$-\frac{2{\beta_2}^2+2}{{{\beta_1}}^2+{{\beta_2}}^2+1}$	$\frac{2\beta_2^2 + 2\alpha\beta_1\beta_2 + 2}{\beta_1^2 + \beta_2^2 + 1}$	-2	2	$\frac{2\beta_1\beta_2}{{\beta_1}^2 + {\beta_2}^2 + 1}$

$$\frac{\overline{16}\text{Notice that }(\overline{z}_{k}[1] - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[1] + \eta \overline{F}(\overline{z}_{k+1})[1])^{2}}{2\eta \overline{F}(\overline{z}_{k+1})[1](\overline{z}_{k}[1] - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[1])^{2} + \eta^{2} \overline{F}(\overline{z}_{k+1})[1]^{2} + 2\eta \overline{F}(\overline{z}_{k+1})[1](\overline{z}_{k}[1] - \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[1]).}$$

	$\eta \overline{F}(\overline{z}_k)[2] \overline{z}_{k+\frac{1}{2}}[\overline{z}_k]$	1]	$ \eta \overline{F}(\overline{z}_{k+\frac{1}{2}}) $	$[2] \overline{z}_k[1]$	$\eta \overline{F}(\overline{z}_{k+})$	$(z_{\frac{1}{2}})[2]\overline{z}_{k+\frac{1}{2}}[1]$	$\eta \overline{F}(\overline{z})$	$(\overline{z}_k)[3] \overline{z}_k[1]$
Expression(69)	0		0		0			0
Expression(70)	0		0		0			0
Expression(71)	0		0	0		2 α		0
Expression(72)	-2α		0		0			0
Expression(73)	0		$\frac{2\alpha}{\alpha^2+}$	 ī		-2α		0
Expression(74)	0		$ \begin{array}{c} \frac{2\alpha}{\alpha^2+1} \\ -\frac{2\alpha}{\alpha^2+1} \end{array} $		0			0
Expression(75)	$-\tfrac{\beta_2(2\beta_1\!-\!2\alpha\beta_2)}{{\beta_1}^2\!+\!{\beta_2}^2\!+\!1}$		0		0		${\beta_1}$	$\frac{2\beta_1}{^2+\beta_2^2+1}$
Expression(76)	0		0		0		PI	0
Sum	$-\frac{2\alpha(\beta_1^2+1)}{\beta_1^2+\beta_2^2+1}-\frac{2\beta_1}{{\beta_1}^2+\beta_2^2}$	$\frac{\beta_2}{\beta_2^2+1}$	0		0		$\overline{\beta_1}$	$\frac{2\beta_1}{^2+\beta_2^2+1}$
Expression(78)	0		0			0		0
Expression(79)	0		0	0		0		0
Expression(80)	0		0			0	$\overline{\beta_1}$	$\frac{2\beta_1}{^2+{\beta_2}^2+1}$
Expression(81)	$-\frac{2\alpha}{\beta_1^2+\beta_2^2+1}$		0		0			0
Expression(82)	$-\frac{\beta_1^2+\beta_2^2+1}{\beta_1^2+\beta_2+1} \\ -\frac{2\beta_1(\alpha\beta_1+\beta_2)}{\beta_1^2+\beta_2^2+1}$		0			0		0
Sum	$-\frac{2\alpha(\beta_1^2+1)}{\beta_1^2+\beta_2^2+1}-\frac{2\beta_1\beta_2}{\beta_1^2+\beta_2^2+1}$		0		0 -		$\overline{\beta_1}$	$\frac{2\beta_1}{^2+{\beta_2}^2+1}$
	$\eta \overline{F}(\overline{z}_k)[3] \overline{z}_{k+\frac{1}{2}}[1]$	$\eta \overline{F}(\bar{z})$	$(\overline{z}_k)[1] \overline{z}_k[2]$	$\eta \overline{F}(\overline{z}_k)$	$[2] \overline{z}_k[2]$	$ \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[2] $		
Expression(69)	0		0	(0	0		
Expression(70)	0		0	(0	0		
Expression(71)	0		0		0	2		
Expression(72)	0		0		0	0		
Expression(73)	0		$\frac{2\alpha}{\alpha^2+1}$	$-\frac{1}{\alpha}$	$\frac{2 \alpha^2}{(2+1)}$	$-\frac{2\alpha^2}{\alpha^2+1}$		
Expression(74)	0		$-\frac{2\alpha}{\alpha^2+1}$		$\frac{2}{2^2+1}$	$-\frac{2\alpha^2}{\alpha^2+1}$ $-\frac{2}{\alpha^2+1}$		
Expression(75)	$-rac{2eta_1-2lphaeta_2}{{eta_1}^2+{eta_2}^2+1}$	$\overline{\beta_1}$	$\frac{2\beta_1\beta_2}{^2+\beta_2^2+1}$	2 /		0		
Expression(76)	0	, -	0		0	0		
Sum	$-rac{2eta_1-2lphaeta_2}{{eta_1}^2+{eta_2}^2+1}$	$\overline{eta_1}$	$\frac{2\beta_1\beta_2}{(2+\beta_2)^2+1}$	$-\frac{2\mu}{\beta_1^2}$	$\frac{{\beta_1}^2+2}{+{\beta_2}^2+1}$	0		
Expression(78)	0		0		0	0		
Expression(79)	0		0		0	0		
Expression(80)	$-rac{2eta_1-2lphaeta_2}{eta_1^2+eta_2^2+1}$		0		0	0		
Expression(81)	0		0	$-{\beta_1^2}$	$\frac{2}{+\beta_2^2+1}$	0		
Expression(82)	0		$\frac{2\beta_1\beta_2}{(2+\beta_2)^2+1}$	$-rac{eta^2}{eta_1^2}$	$\frac{2\beta_1^2}{+\beta_2^2+1}$	0		
Sum	$-rac{2eta_1-2lphaeta_2}{{eta_1}^2+{eta_2}^2+1}$	$\overline{\beta_1}$	$\frac{2\beta_1\beta_2}{^2+{\beta_2}^2+1}$	$-\frac{2\mu}{\beta_1^2}$	$\frac{{\beta_1}^2 + 2}{+{\beta_2}^2 + 1}$	0		

	$\eta \overline{F}(\overline{z}_k)[3] \overline{z}_k[2]$	$\overline{z}_k[1]\overline{z}_{k+\frac{1}{2}}[1]$	$\overline{z}_k[1]\overline{z}_k[2]$	$\overline{z}_{k+\frac{1}{2}}[1]\overline{z}_k[2]$	
Expression(69)	0	0	0	0	
Expression(70)	0	0	0	0	
Expression(71)	0	0	0	-2α	
Expression(72)	0	-2	0	2 α	
Expression(73)	0	0	$-\frac{2\alpha}{\alpha^2+1}$	2 α	
Expression(74)	0	0	$ \begin{array}{r} -\frac{2 \alpha}{\alpha^2 + 1} \\ \frac{2 \alpha}{\alpha^2 + 1} \\ \frac{2 \alpha}{\alpha^2 + 1} \end{array} $	0	
Expression(75)	$\frac{2\beta_2}{{\beta_1}^2 + {\beta_2}^2 + 1}$	0	0	0	
Expression(76)	0	0	0	0	
Sum	$\frac{2\beta_2}{\beta_1^2 + \beta_2^2 + 1}$	-2	0	2 α	
Expression(78)	0	0	0	0	
Expression(79)	0	$-\frac{2}{\beta_1^2+\beta_2^2+1}$	0	0	
Expression(80)	$\frac{2\beta_2}{\beta_1^2 + \beta_2^2 + 1}$	$-\frac{2\beta_{1}(\beta_{1}-\alpha\beta_{2})}{\beta_{1}^{2}+\beta_{2}^{2}+1}$	$\frac{2\beta_1\beta_2}{{\beta_1}^2 + {\beta_2}^2 + 1}$	$-\frac{2\beta_2(\beta_1 - \alpha\beta_2)}{{\beta_1}^2 + {\beta_2}^2 + 1}$	
Expression(81)	0	0	0	$\frac{\frac{2\alpha}{\beta_1^2 + \beta_2^2 + 1}}{\frac{2\alpha}{\beta_1^2 + \beta_2^2 + 1}}$	
Expression(82)	0	$-\frac{2{\beta_2}^2 + 2\alpha\beta_1\beta_2}{{\beta_1}^2 + {\beta_2}^2 + 1}$	$-\frac{2\beta_1\beta_2}{{\beta_1}^2 + {\beta_2}^2 + 1}$	$\frac{2\beta_{1}(\alpha\beta_{1}+\beta_{2})}{\beta_{1}^{2}+\beta_{2}^{2}+1}$	
Sum	$\frac{\frac{2\beta_2}{\beta_1^2 + \beta_2^2 + 1}}{(\eta \overline{F}(\overline{z}_k)[1])^2}$	-2	0	2α	
		$\overline{(\eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[1])^2}$	$(\eta \overline{F}(\overline{z}_{k+1})[1]$	$ f ^2 f ^2 f ^2 f ^2$	$(\eta \overline{F}(\overline{z}_{k+1})[2])^2$
Expression(69)	$1 - \frac{{eta_1}^2}{{eta_1}^2 + {eta_2}^2 + 1}$	0	0	$1 - \frac{{\beta_2}^2}{{\beta_1}^2 + {\beta_2}^2 + 1}$	-1
Expression(70)	0	0	0	0	0
Expression(71)	0	1	1	0	1
Expression(72)	0	0	0	0	0
Expression(73)	0	0	0	0	0
Expression(74)	0	0	0	0	0
Expression(75)	0	0	0	0	0
Expression(76)	0	0	0	0	0
Sum	$\frac{{\beta_2}^2 + 1}{{\beta_1}^2 + {\beta_2}^2 + 1}$	1	1	$\frac{{\beta_1}^2 + 1}{{\beta_1}^2 + {\beta_2}^2 + 1}$	0
Expression(78)	0	1	1	0	0
Expression(79)	$\frac{1}{\beta_1^2 + \beta_2^2 + 1}$	0	0	0	0
Expression(80)	0	0	0	0	0
Expression(81)	0	0	0	$\frac{1}{\beta_1^2 + \beta_2^2 + 1}$	0
Expression(82)	$\frac{{\beta_2}^2}{{\beta_1}^2 + {\beta_2}^2 + 1}$	0	0	$\frac{\beta_1^2}{\beta_1^2+\beta_2^2+1}$	0
Sum	$\frac{{\beta_2}^2 + 1}{{\beta_1}^2 + {\beta_2}^2 + 1}$	1	1	$\frac{{\beta_1}^2 + 1}{{\beta_1}^2 + {\beta_2}^2 + 1}$	0

	$(\eta \overline{F}(\overline{z}_k)[3])$	$(\eta \overline{F}(\overline{z}_{k+1})[3])^2$	$\overline{z}_k[1]^2$	$\overline{z}_{k+\frac{1}{2}}[1]^2$	
Expression(69)	$1 - \frac{1}{{{eta_1}^2} + {{eta_2}^2}}$	1	0	0	
Expression(70)	0	0	0	0	
Expression(71)	-1	1	0	$-\alpha^2-1$	
Expression(72)	0	0	0	$2\alpha^2+2$	
Expression(73)	0	0	0	0	
Expression(74)	0	0	0	0	
Expression(75)	$\frac{2}{\beta_1^2 + \beta_2^2 + 1}$	0	0	0	
Expression(76)	0	0	2	0	
Sum	$\frac{1}{\beta_1^2 + \beta_2^2 + 1}$	0	2	$\alpha^2 + 1$	
Expression(78)	0	0	1	0	
Expression(79)	0	0	$\frac{1}{\beta_1^2 + \beta_2^2 + 1}$	$\frac{1}{\beta_1^2 + \beta_2^2 + 1}$	
Expression(80)	$\frac{1}{{\beta_1}^2 + {\beta_2}^2 + 1}$	0	$\frac{\beta_1 + \beta_2 + 1}{\beta_1^2 + \beta_2^2 + 1}$	$\frac{(\beta_{1} + \beta_{2} + 1)}{(\beta_{1} - \alpha \beta_{2})^{2}}$ $\frac{(\beta_{1} - \alpha \beta_{2})^{2}}{(\beta_{1}^{2} + \beta_{2}^{2} + 1)}$	
Expression(81)	0	0	0	$\frac{\alpha^2}{\beta_1^2 + \beta_2^2 + 1}$	
Expression(82)	0	0	$\frac{{\beta_2}^2}{{\beta_1}^2 + {\beta_2}^2 + 1}$	$\frac{\frac{\beta_1 + \beta_2 + 1}{(\beta_2 + \alpha\beta_1)^2}}{\frac{\beta_1^2 + \beta_2^2 + 1}{\beta_1^2 + \beta_2^2 + 1}}$	
Sum	$\frac{1}{{\beta_1}^2 + {\beta_2}^2 + 1}$	0	2	$\alpha^2 + 1$	
	$\overline{z}_k[2]^2$	$\overline{\eta}\overline{F}(\overline{z}_{k+\frac{1}{2}})[1]\overline{\eta}\overline{F}(\overline{z}_{k+\frac{1}{2}})$	$\overline{(1)[1] \mid \eta \overline{F}(\overline{z})}$	$\overline{(\overline{z}_k)[1] \eta \overline{F}(\overline{z}_k)}[2]$	$\overline{\eta}\overline{F}(\overline{z}_k)[1]\eta\overline{F}(\overline{z}_{k+\frac{1}{2}})[2]$
Expression(69)	0	0	-	$-\frac{2eta_1eta_2}{{eta_1}^2 + {eta_2}^2 + 1}$	0
Expression(70)	0	0	0		0
Expression(71)	-1	-2		0	0
Expression(72)	0	0		0	0
Expression(73)	$ \frac{\frac{2\alpha^2}{\alpha^2+1}}{\frac{2}{\alpha^2+1}} $	0		0	$-\frac{2\alpha}{\alpha^2+1}$ $-\frac{2\alpha}{2\alpha}$
Expression(74)	$\frac{2}{\alpha^2+1}$	0		0	$\frac{2\alpha}{\alpha^2+1}$
Expression(75)	0	0		0	0
Expression(76)	0	0		0	0
Sum	1	-2	$-\frac{2\beta_1\beta_2}{{\beta_1}^2\!+\!{\beta_2}^2\!+\!1}$		0
Expression(78)	0	-2		0	0
Expression(79)	0	0	0		0
Expression(80)	$\frac{{{eta_2}^2}}{{{eta_1}^2} + {{eta_2}^2} + 1}$	0	0		0
Expression(81)	$\frac{1}{\beta_1^2 + \beta_2^2 + 1}$	0	0		0
Expression(82)	$\frac{\beta_1^2}{\beta_1^2 + \beta_2^2 + 1}$	0	$-\frac{2\beta_1\beta_2}{{\beta_1}^2\!+\!{\beta_2}^2\!+\!1}$		0
Sum	1	-2	-	$-\frac{2\beta_1\beta_2}{{\beta_1}^2 + {\beta_2}^2 + 1}$	0

	$ \eta \overline{F}(\overline{z}_k)[1] \eta \overline{F}(\overline{z}_k)[3] $	$\overline{\eta}\overline{F}(\overline{z}_k)[2]\overline{\eta}\overline{F}(\overline{z}_{k+\frac{1}{2}})[2]$	$\overline{\eta}\overline{F}(\overline{z}_{k+\frac{1}{2}})[2]\overline{\eta}\overline{F}(\overline{z}_{k+1})[2]$	$\overline{\eta}\overline{F}(\overline{z}_k)[2]\overline{\eta}\overline{F}(\overline{z}_k)[3]$
Expression(69)	$-rac{2eta_1}{{eta_1}^2+{eta_2}^2+1}$	0	0	$-rac{2eta_2}{{eta_1}^2+{eta_2}^2+1}$
Expression(70)	0	-2	2	0
Expression(71)	0	0	-2	0
Expression(72)	0	0	0	0
Expression(73)	0	$\frac{\frac{2\alpha^2}{\alpha^2+1}}{\frac{2}{\alpha^2+1}}$	0	0
Expression(74)	0	$\frac{2}{\alpha^2+1}$	0	0
Expression(75)	$\frac{2\beta_1}{{\beta_1}^2 + {\beta_2}^2 + 1}$	0	0	$\frac{2\beta_2}{\beta_1^2 + \beta_2^2 + 1}$
Expression(76)	0	0	0	0
Sum	0	0	0	0
Expression(78)	0	0	0	0
Expression(79)	0	0	0	0
Expression(80)	0	0	0	0
Expression(81)	0	0	0	0
Expression(82)	0	0	0	0
Sum	0	0	0	0
	$n\overline{F}(\overline{z}_{1})[3] n\overline{F}(\overline{z}_{1})[3]$	\overline{S} $n\overline{F}(\overline{z}_{k+1})[3] n\overline{F}(\overline{z}_{k+1})$)[3]	

	$ \eta \overline{F}(\overline{z}_k)[3] \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[3] $	$ \eta \overline{F}(\overline{z}_{k+\frac{1}{2}})[3] \eta \overline{F}(\overline{z}_{k+1})[3] $
Expression(69)	0	0
Expression(70)	-2	2
Expression(71)	2	-2
Expression(72)	0	0
Expression(73)	0	0
Expression(74)	0	0
Expression(75)	0	0
Expression(76)	0	0
Sum	0	0
Expression(78)	0	0
Expression(79)	0	0
Expression(80)	0	0
Expression(81)	0	0
Expression(82)	0	0
Sum	0	0

G Non-Monotonicity of Several Standard Performance Measures

We conduct numerical experiments by trying to find saddle points in constrained bilinear games using EG, and verified that the following performance measures are not monotone: the (squared) natural residual, $\|z_k - z_{k+\frac{1}{2}}\|^2$, $\|z_k - z_{k+1}\|^2$, $\max_{z \in \mathcal{Z}} \langle F(z), z_k - z \rangle$, $\max_{z \in \mathcal{Z}} \langle F(z_k), z_k - z \rangle$.

All of our counterexamples are constructed by trying to find a saddle point in bilinear games of the following form:

$$\min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} x^{\top} A y - b^{\top} x - c^{\top} y \tag{83}$$

where $\mathcal{X}, \mathcal{Y} \subseteq \mathbb{R}^2$, A is a 2×2 matrix and b, c are 2-dimensional column vectors. All of the instances of the bilinear game considered in this section have $\mathcal{X}, \mathcal{Y} = [0, 10]^2$. We denote by $\mathcal{Z} = \mathcal{X} \times \mathcal{Y}$ and by $F(x,y) = \begin{pmatrix} Ay - b \\ -A^\top x + c \end{pmatrix} : \mathcal{Z} \to \mathbb{R}^n$. We remind readers that finding a saddle point of bilinear game (83), is equivalent to solving the monotone VI with operator F(z) on set \mathcal{Z} .

G.1 Non-Monotonicity of the Natural Residual and its Variants

Performance Measure: Natural Residual. Let $A = \begin{bmatrix} 1 & 2 \\ 1 & 1 \end{bmatrix}$, $b = c = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$. Running the EG method on the corresponding VI problem with step-size $\eta = 0.1$ starting at $z_0 = (0.3108455, 0.4825575, 0.4621875, 0.5768655)^T$ has the following trajectory:

$$z_1 = (0.24923465, 0.47967569, 0.43497808, 0.57458145)^T,$$

 $z_2 = (0.19396855, 0.48164918, 0.40193211, 0.56061753)^T.$

Thus we have

$$r^{nat}(z_0)^2 = 0.15170013184049996,$$

 $r^{nat}(z_1)^2 = 0.13617654362050116,$
 $r^{nat}(z_2)^2 = 0.16125792556139756.$

It is clear that the natural residual is not monotone.

Performance Measure: $\|z_k - z_{k+\frac{1}{2}}\|^2$. Note that the norm of the operator mapping defined in [Diakonikolas, 2020] is exactly $\frac{1}{\eta} \cdot \|z_k - z_{k+\frac{1}{2}}\|$. Let $A = \begin{bmatrix} 0.50676631 & 0.15042569 \\ 0.46897595 & 0.96748026 \end{bmatrix}$, $b = c = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$. Running the EG method on the corresponding VI problem with step-size $\eta = 0.1$ starting at $z_0 = (2.35037432, 0.00333996, 1.70547279, 0.71065999)^T$ has the following trajectory:

$$\begin{split} z_{\frac{1}{2}} = & (2.35325656, 0, 1.72473848, 0.64633879)^T, \\ z_{1} = & (2.35324779, 0, 1.72472791, 0.64605901)^T, \\ z_{1+\frac{1}{2}} = & (2.35612601, 0, 1.74398258, 0.58145791)^T \\ z_{2} = & (2.35612201, 0, 1.74412844, 0.5815012)^T, \\ z_{2+\frac{1}{2}} = & (2.35898819, 0, 1.76352876, 0.51694333)^T. \end{split}$$

Thus we have

$$\left\| z_0 - z_{\frac{1}{2}} \right\|^2 = 0.00452784581555656,$$

$$\left\| z_1 - z_{1 + \frac{1}{2}} \right\|^2 = 0.004552329544896258,$$

$$\left\| z_2 - z_{2 + \frac{1}{2}} \right\|^2 = 0.004552306444552208.$$

It is clear that the $||z_k - z_{k+\frac{1}{2}}||^2$ is not monotone.

Performance Measure: $||z_k - z_{k+1}||^2$. Let $A = \begin{bmatrix} 0.50676631 & 0.15042569 \\ 0.46897595 & 0.96748026 \end{bmatrix}$, $b = c = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$. Running the EG method on the corresponding VI problem with step-size $\eta = 0.1$ starting at $z_0 = (2.37003485, 0, 1.84327237, 0.25934775)^T$ has the following trajectory:

$$z_1 = (2.37267186, 0, 1.86351397, 0.1950396)^T,$$

 $z_2 = (2.37524308, 0, 1.88388624, 0.13077023)^T,$
 $z_3 = (2.37774149, 0.00426125, 1.90438549, 0.06653856)^T.$

Thus we have

$$||z_0 - z_1||^2 = 0.004552214685275266,$$

 $||z_1 - z_2||^2 = 0.004552191904998012,$
 $||z_2 - z_3||^2 = 0.004570327450598002.$

It is clear that the $||z_k - z_{k+1}||^2$ is not monotone.

G.2 Non-Monotonicity of the Gap Functions and its Variant

Performance Measure: Gap Function and $\max_{z\in\mathcal{Z}}\langle F(z),z_k-z\rangle$. Let $A=\begin{bmatrix} -0.21025101 & 0.22360196 \\ 0.40667685 & -0.2922158 \end{bmatrix}$, $b=c=\begin{bmatrix} 0\\ 0 \end{bmatrix}$. One can easily verify that $\langle F(z),z_k-z\rangle=\langle F(z_k),z_k-z\rangle=\langle F(z_k),z_k-z\rangle$ which further implies that $\max_{z\in\mathcal{Z}}\langle F(z),z_k-z\rangle=\max_{z\in\mathcal{Z}}\langle F(z_k),z_k-z\rangle=\langle F(z_k),z_k-z\rangle$. Running the EG method on the corresponding VI problem with step-size $\eta=0.1$ starting at $z_0=(0.53095379,0.29084076,0.62132986,0.49440498)$ has the following trajectory:

$$z_1 = (0.53290086, 0.28009156, 0.62151204, 0.4981395)^T,$$

 $z_2 = (0.5347502, 0.26947398, 0.62122195, 0.50222691)^T.$

One can easily verify that

$$GAP(z_0) = 0.6046398415472187,$$

 $GAP(z_1) = 0.58462873354003214,$
 $GAP(z_2) = 0.5914026255469654.$

It is clear that the duality gap is not monotone.