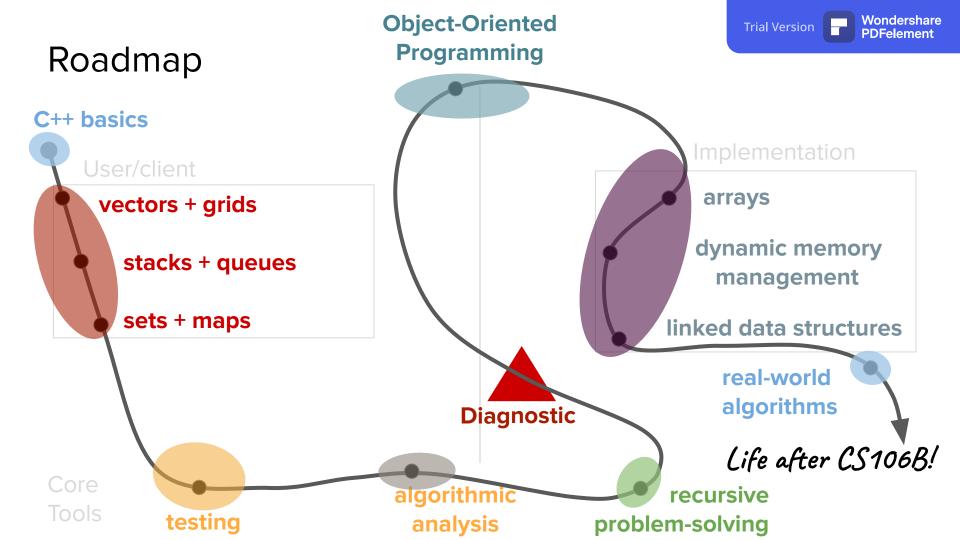
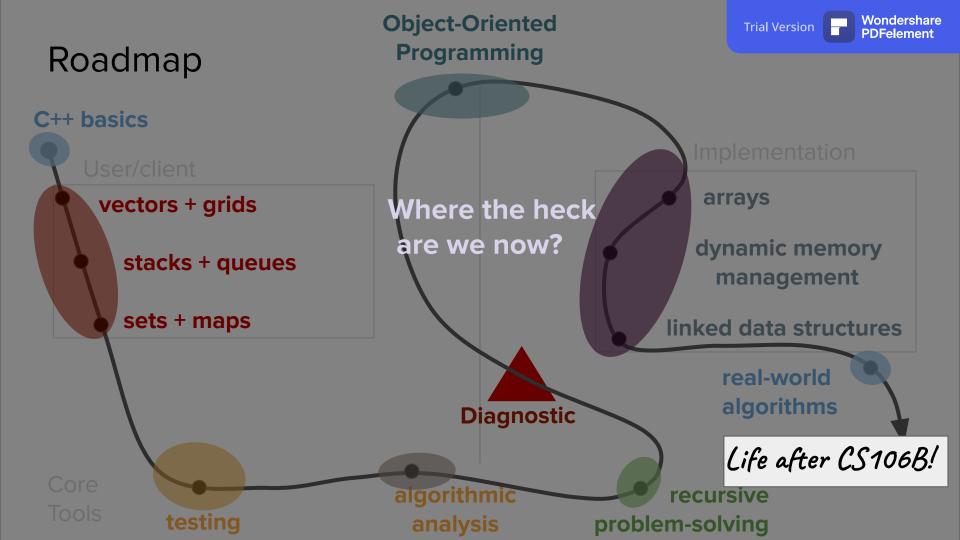
Multithreading and Parallel Computing

What do you think Trip has been up to this quarter?

(wrong answers only in the chat)



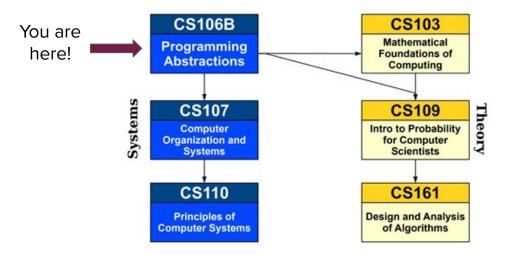






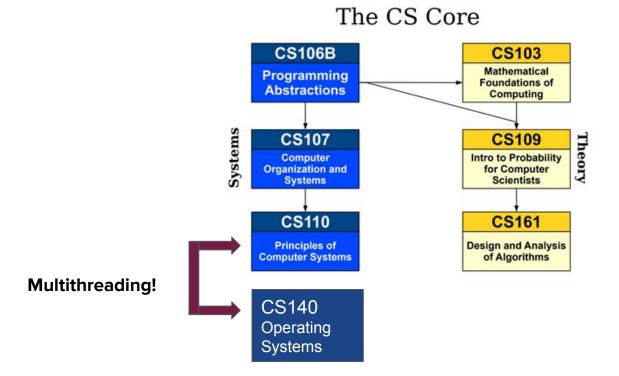


The CS Core



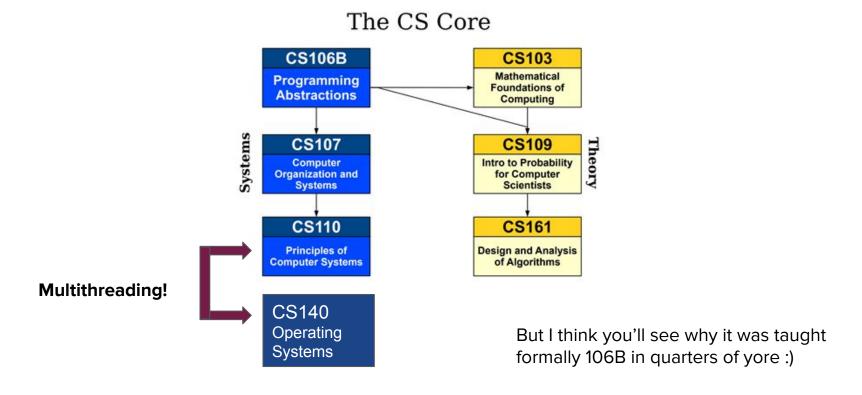


A picture you'll see again...





A picture you'll see again...



Today's question

How can we harness the cores in our computer in order to parallelize a workload safely?



Today's question

woah. How can we harness the cores in our computer in order to parallelize a workload safely?

Multiple cores?

Today's questio,

r to parallelize a kload safely?



Today's topics

1. Review (short!)

Some Computer
 Architecture (Threads & Processors)

3. Thread Safety

Review (short!)

(simple code flow)

How code is run

- How does the computer read and run your code?
 - Logically, it should read your code from top to bottom!

...but *who* is **it**? What's the thing that encapsulates and runs your code?

```
int main () {
   int yeet = 9338;
   double foo = 2.4;

   doSomeMath(yeet);

   cout << "time to go home!" << endl;
   return 0;
}</pre>
```



Line by line, top to bottom!

thread

An abstraction that represents a sequential

execution of code.

Anything that's code!

How to think about threads

- When talking about a **thread**, you'll very frequently see it referenced as a "**thread** of execution."
 - Think about the line on the right as a program's execution. You start at **main()**, which might call other functions, which might return to **main()** or call other helper functions. Although the execution flow of your program may involve many function calls, it will eventually go from the top of **main()** to the bottom.
 - The flow would almost looks like a thread, or a piece of string!



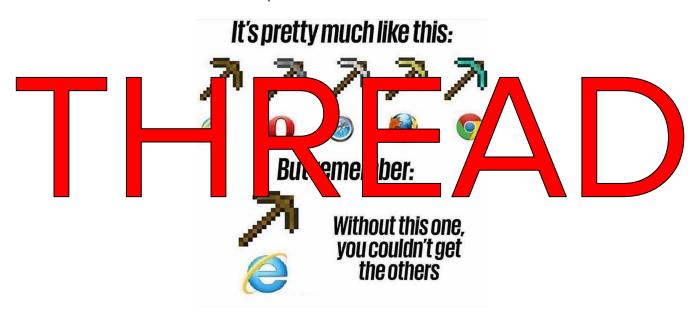
code end

- Right now, your computer probably has a few threads running right now!
 - What are some examples of threads running on your PC?

Are you on Zoom right now?



Do you have a web browser open? (Chrome, Safari?)



Are you watching TikToks during lecture?

Are you watching TikToks during lecture?

Stanford announces Charli D'Ameli as 2020 Commencement Speaker



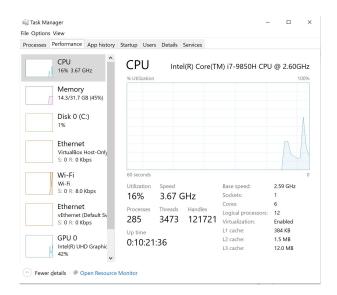
"Charli undeniably captures the same spirit of ingenuity we try to cultivate at our different schools," said President Marc Tessier-Lavigne. (Photo Edit: RICHARD COCA/The Stanford Daily)

I have been told Ms. D'Ameli is a TikTok #influencer

Question:

How many threads do you think my computer had active when I was making this slide?

- Right now, your computer is executing a bunch of threads!
 - At the time of making this slide show, my computer was handline 3473 threads!
- Many large programs (your web browsers!)
 need multiple threads to run. That's because they have so many moving parts!



Question:

When you run a program in Qt Creator, is a thread executing your code?

Answer:

Er... Yes, sort of!

Yes, when you run a program in Qt, a thread encapsulating your code is being **executed**.

However, a thread alone isn't enough to run your code!

Definitions

software

Programs and and abstractions (code). Not a physical entity.

hardware

Physical parts of a computer.

The hardware-software boundary

- A thread alone cannot run your program.
 - A thread is just software that is an abstraction for some code.
- A thread needs to work with the computer's hardware in order to run the code it encapsulates!

... but what piece of hardware does this?

Definitions

CPU (Central Processing Unit)

A piece of hardware responsible for executing instructions that make up a computer program

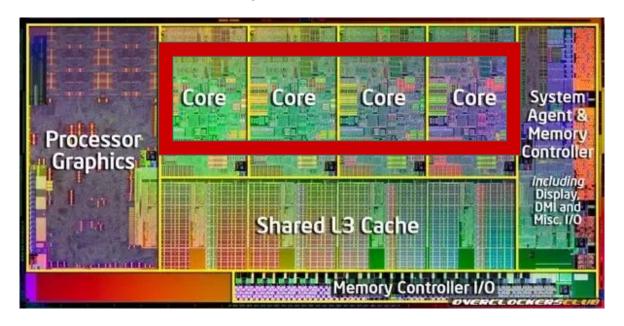
Core

An individual processor inside of a **CPU**. Each **core** is able to execute code independently of other **cores**.



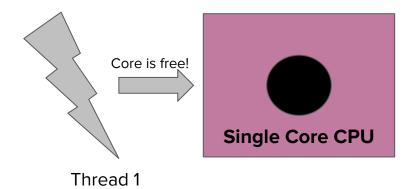
Inside a CPU...

How many concurrent programs can this CPU run?



Don't worry about the other stuff -- we just care about the cores!

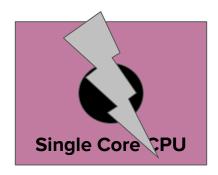
- In order for a **thread** to be able to execute some code, it must be running on a **CPU core**.
- If all cores are currently busy, a thread must wait for a core to free up before it can hop on that core and begin executing its own code!



Let's assume this computer has a CPU with only one core.

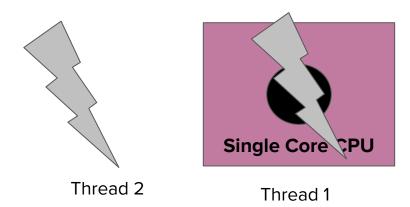
Question: if the core is free, how is anything getting done :o

- In order for a **thread** to be able to execute some code, it must be running on a **CPU core**.
- If all **cores** are currently busy, a thread must **wait** for a **core** to free up before it can hop on that **core** and begin executing its own code!

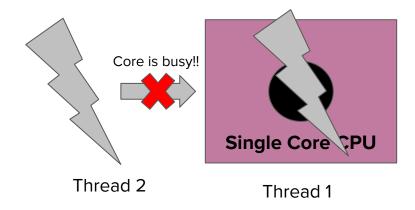


Thread 1

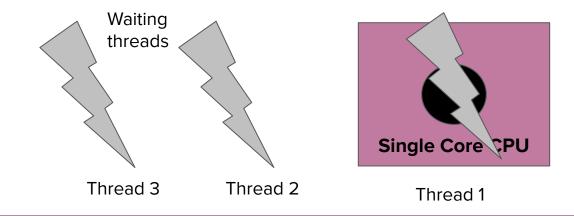
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Question:

Who decides how long a thread should be able to run on a processor? Who decides which thread should run next?

What was running when the single-core was free in the example???

Definition

Operating System

Code that manages the relationship between a computer's **hardware** and **software**.

Thread Scheduling

- The Operating System, determines both how long a thread should run on a core, AND which thread should run next.
 - Want to learn how to implement these strategies? Take CS140!
- For the purposes of this lecture, let's assume that a thread will run on a core until its program terminates or it is forced off the processor.

Code example

- Let's take a break from all of this low-level jazz and write a simple program!
- Let's say you wanted to revise your **A2 Search Engine** program by cheating and making it ping the internet with queries.
 - Such a task is called I/O Bound, because the performance bottleneck is the waiting that happens between sending your request and getting your data! (We call this, and anything involving communication with the outside world, I/O)

Code example

Let's write a program that repeatedly executes the below I/O
bound function. (Forget the search engine thing; that's just an
example of such a task).

```
static void task (int input);
```

 I've already implemented task for you; all you need to do is call it repeatedly!

Code example

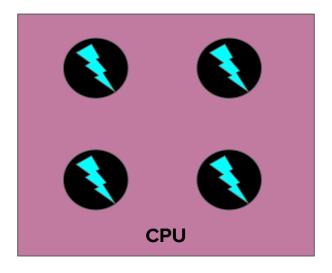
- Let's write a program that repeatedly executes an I/O bound function. (Forget the search engine thing, let's just say it's any old I/O bound function).
- I've already written the I/O bound function for you; all you need to do is call it repeatedly and store the many return values in a **Vector**.

Let's do it!

Code example

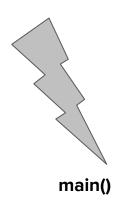
- What happened there?
 - Our code was slow as heck! This shouldn't be surprising, however. Here's what happened:

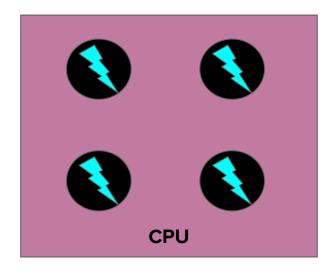


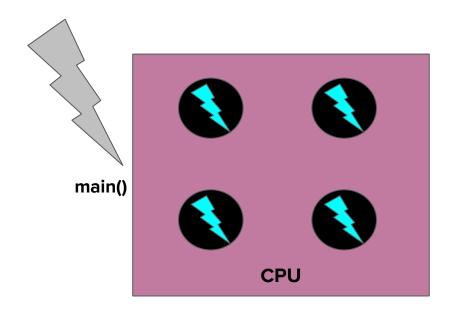


Before you run your program, your **CPU** is probably chugging away at other tasks!



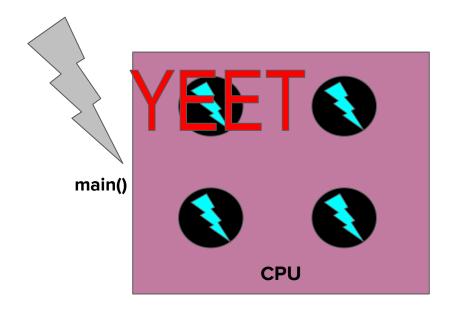




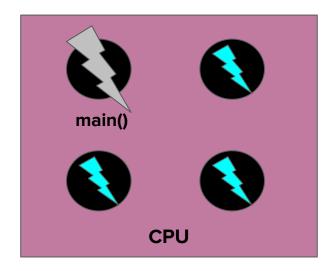


main() is a pretty important thread, so it has the power to boot another thread off a core!



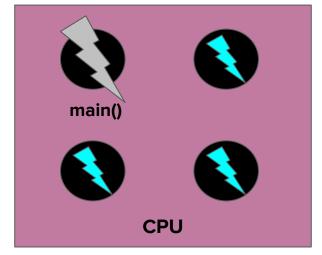






 When you call the I/O bound function task() from main(), the thread will remove itself from the processor, as it is waiting on an I/O and therefore unable to do any work. Another thread will take its place

immediately.

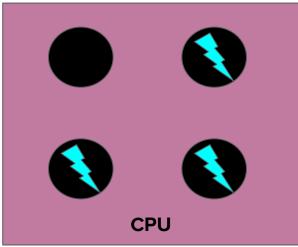


Question for yourselves: why does self-removal make sense here?

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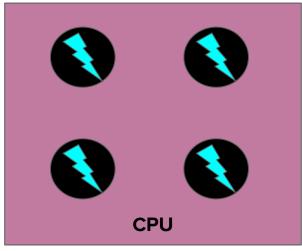




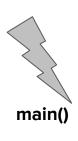
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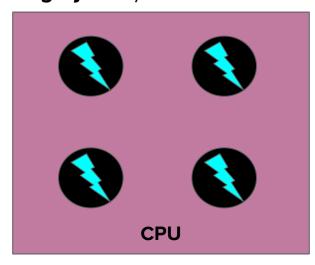
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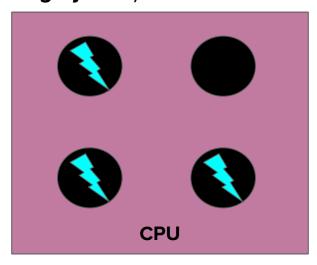
 When the I/O bound task completes, your thread will attempt to get back on a core as soon as possible in order to continue (but its order in line is up to your Operating System)





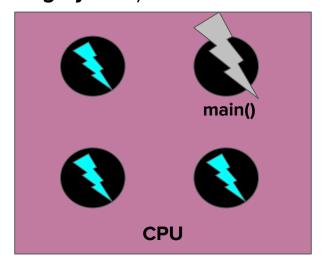
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A vacancy!

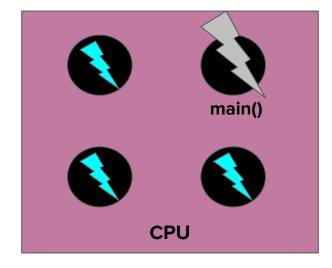
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Note how we're core agnostic.
This doesn't need to be the case in some
OS schedulers.



Questions about these events?



- This process of getting on a core, removing ourselves and waiting, and reacquiring a core happened every time we called task()
 - Can we do better?

But first...

Announcements

Announcements

- Make sure to sign up for a final presentation time slot if you haven't already!
- Assignment 6 is due tomorrow at 11:59pm PDT. Remember that this is a hard deadline and there is no grace period!
- In lecture tomorrow, we will be having an "Ask Us Anything" component for the last part of lecture. We'll be collecting questions in advance as well – if you have any burning inquiries on your mind, go ahead and fill out this Google
 form!
- Remember that there is no section this week!

Back to the action!

- This process of getting on a core, removing ourselves and waiting, and reacquiring a core happened every time we called task()
 - Can we do better?
- In the words of a sectionee last quarter...
 - "Let's parallelize this bad boy"

Multithreading

- Let's try and implement this same routine using **multithreading**.
 - That means we'll try and use multiple threads instead of one in order to parallelize the workflow!
- Before you can make threads, you'll first need to:

#include <thread>

 Bonus points: this is a standard c++ library, so no Stanford-only woes!

Multithreading

• To instantiate a thread, it's pretty simple!

```
thread newthread = thread(funcName);
```

- This should look pretty vanilla, except for the parameter!
 - funcName is the name of a the function you want to execute!
 - Let's make new threads that encapsulate task()!

Thread joining

- Woah woah, hold your horses, eager beaver. Two things to think about:
 - As soon as you instantiate a thread, it begins to run. Be sure you're ready before you dispatch them.
 - Threads are somewhat resource intensive, so when we dispatch them, we need to keep track of them so that we can clean up their memory once they've completed.
 - This is very much like the **new** and **delete** keywords you've used!

Thread joining

- After you've spawned a thread, simply call **threadName.join()** to clean it up.
 - This usually requires storing your threads in a collection! Note:
 Stanford's Vector can't store threads because it needs an update :(

Questions about creating / joining threads?

- You can call join() from your **main()** thread immediately after spawning the thread. Don't worry, **main()** will wait for your thread to finish:).
- To pass params to a thread, just include them as the subsequent parameters.

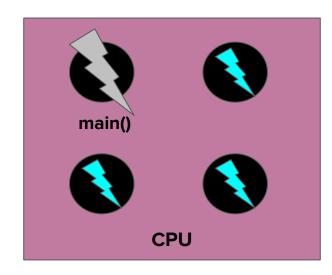
Let's Parallelize!



Wow, that was super fast!

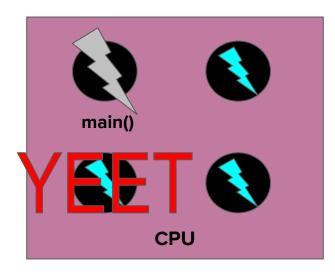
- When our main() thread spawned up a new thread, the new thread might have taken a new core on the processor!
 - o note* we don't know exactly what happened, but it could have done this!



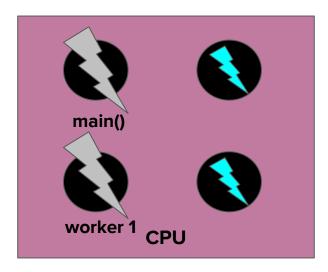


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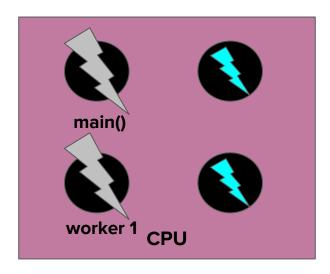




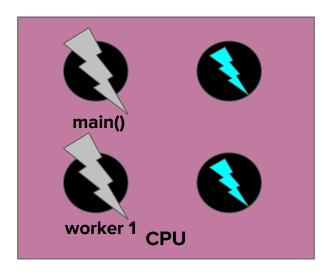
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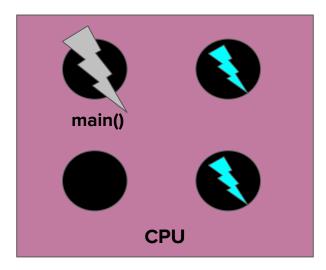
Note now that both main() and worker 1 are running concurrently!



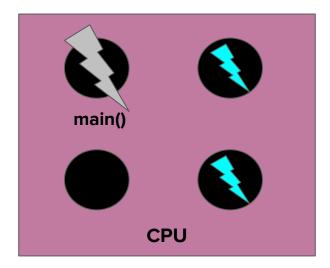
Worker 1 will start its I/O and remove itself from the core, getting replaced



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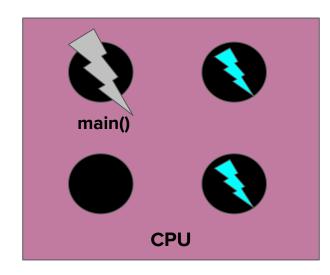


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- But lo! Who is that in the distance?



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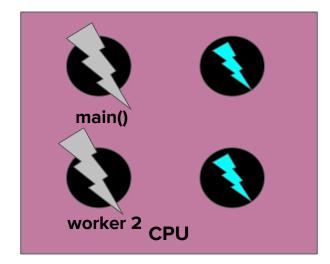




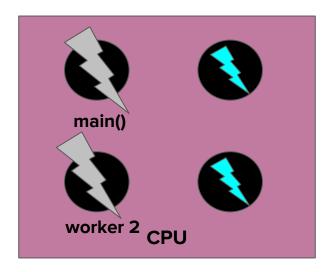
- Worker 1 will start its I/O and remove itself from the core, getting replaced
- But lo! Who is that in the distance?

While worker 1 was waiting for its I/O, main() was busy spinning up new

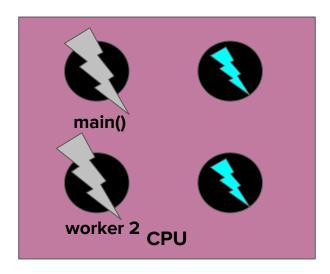
threads!



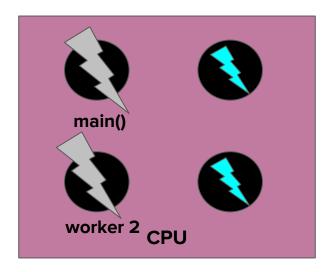
• This process will continue -- each **worker thread** will only need to be on a core for a fraction of a second, just to set up the **I/O**, and then it can leave the processor and let a new **worker thread** set up its **I/O**.



- A similar thing will happen at completion time!
 - Each thread will be able to retake a core, but the core will only be needed for a few instructions! Then the task() will finish, and a new thread will try and complete!



A fair warning -- you can't predict which worker thread will begin working first!
 It might seem like worker 1 should always start first, but the OS and CPU work in unpredictable ways!



- The example you saw was blazing fast because the **task** at hand only needed to be on the processor for a **short period of time**.
- As you can see, the process of yielding a core to another worker takes an almost imperceptible amount of time!
 - That's because your OS is doing it constantly :o
- Parallelization is less successful when you don't have long I/O waits.
 - Take CS140 to find out more :)

Questions?

Bonus! Race Conditions

- Remember when I said that we can't really determine the order that threads will run in? Let's show that!
- Let's add logging to our code to show the order that threads show up!
- It's easy! Just add a print statement inside inside task() and keep a counter variable!

Let's try it!

woah...

Definition

Race Condition

A bug that is the product of two threads "racing" against each other and operating on the same state in the incorrect order.

Bonus: Race Conditions

- Congratulations, you've experienced your first race condition!
- It turns out that **cout** is not **thread-safe**, meaning that it will not behave predictably if you have multiple threads calling it at the same time!
 - Every time you printed to the console, you had some jumbling of all 10 cout statements!

How can we fix this?

Definition

Atomic

A state that can only be observed or superseded **before** or **after** an operation occurs, **not during.**

Mutex

- To make code atomic, we can use something called a mutex.
 - Sounds like Mut(ual) Ex(clusion)!
- To make a mutex, you'll need this library:

and you'll want to declare a single mutex like this:

```
mutex m;
```

Mutex

 You'll want to make a single mutex, and pass it as a pointer to your worker threads.

```
thread t = thread (funcName, &mutexName);
```

• In order to make code **atomic**, all you need to do is wrap the code in question around these two statements:

```
mutexName->lock();
mutexName->unlock();
```

Mutex

 In order to make code atomic, all you need to do is wrap the code around these two statements:

```
mutexName->lock();
mutexName->unlock();
```

- When you lock a mutex, any other threads trying to lock that mutex will be forced to wait until you unlock it.
 - Once you **unlock**, the **Operating System** decides which thread can **lock** the **mutex** next!

Let's try it!

We're still not done!?

Why is everything 10?

```
Setting the program up...
Let's process 10 numbers!
Starting in 3...
2...
GO!!
Hello from worker 10
All done! The total time spent working was 1353 milliseconds (roughly 1 second!)
```

We're still not done!?

- Remember how we passed id by reference? (using a pointer)
- The problem is that the threads <u>share</u> the variable "i"
- This actually indicates that main() finished the for loop that created all ten threads (therefore increasing i to the max value) before a single worker could complete.
 - This should make sense because even the first worker had to wait a full second before it could print anything!
- How do we fix this?

Final thoughts

- Multithreading is an incredibly powerful tool that lets you parallelize work among your CPU's cores.
- Threads are a fundamental building block of computing that play an important role in Operating Systems!
- When using multiple threads, be wary of any data that is shared between them.
 - Using a mutex allows you to enforce atomicity in sections of code, but sometimes even that isn't enough!
 - If all of your code is atomic, there's no parallelization at all!
- If you liked this topic, CS110 and CS140 (and CS149) go into more depth :)

What's next?

