

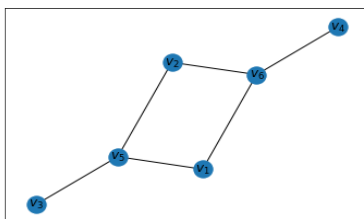
# Homework 2

Wei Ye\*

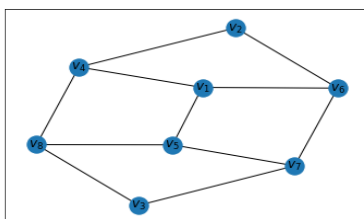
CISC5825 - Computer Algorithm

Due on Feb 13, 2023

- (a) The first graph needs 3 colors.  $\{color1 : v1, v2, v3; color2 : v4, v5; color3 : v6\}$ . The second graph needs 5 colors.  $\{color1 : v1, v2; color2 : v3; color3 : v4, v5; color4 : v6, v7; color5 : v8\}$
- (b) For the first graph vertices, see the figure below: Only two colors needed.



For the second graph vertices, see the figure below: 3 colors needed.



- (c) The two figures in the question are fit for this sub-question. For the first graph see 1, for the second, see 2
- (d) The chromatic number for the first graph is 2.  $v_1, v_2, v_3, v_6$  share the same color,  $v_4, v_5$  share the same color.
- The chromatic number for the second graph is 3.  $v_1, v_2$  share the same color,  $v_4, v_5, v_8$  share the same color, and  $v_3, v_6, v_7$  share the same color.

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\*2nd year PhD student in Economics Department at Fordham University. Email: wye22@fordham.edu

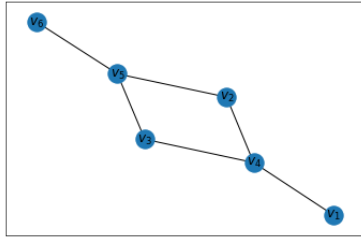


Figure 1: First Graph of question c

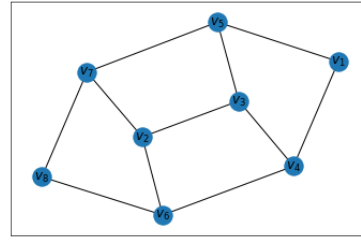


Figure 2: second graph of question c

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**Algorithm 1** Algorithm for coloring vertices

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(e) **procedure** GREEDY( $G$ , newcolor)  
      $newcolor \leftarrow \emptyset$   
     **for** each uncolored vertex  $v$  of  $G$  **do**  
         **if**  $v$  is not adjacent to any vertex in newcolor **then**  
             mark  $v$  colored  
             add  $v$  to newcolor  
         **end if**  
     **end for**  
**end procedure**

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Algorithm for Incompatible Turns

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**procedure** GREEDY(list: $T$ , newcolor)  
      $newcolor \leftarrow \emptyset$   
     **for** each uncolored  $t$  of  $T$  **do**  
         **if**  $t$  does not have common edge to any vertex in newcolors **then**  
             mark  $t$  colored  
             add  $t$  to newcolor  
         **end if**  
     **end for**  
**end procedure**

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