

# Microeconomic Theory II

## Preference, Utility, & Choice

Is it possible to develop a model that explains observed choice?

Primitives: unobserved "preferences".

$X$  is the choice set, or a list of alts.

(★) preference relation  $x \succeq y \Leftrightarrow x$  at least as good as  $y$

strict preference rel.  $x \succ y \Leftrightarrow x \succeq y$  but not  $y \succeq x$ .

indifference rel.  $x \sim y \Leftrightarrow x \succeq y$  and  $y \succeq x$   
( $x \neq y$  and  $y \neq x$ )

ex.

$$X = \{\{ \text{coffee, tea} \}\}$$

$$X = \{\{ \text{coffee, tea}, \{ \text{coffee, tea, wine} \} \}$$

$$x = \{ \text{coffee} \}$$

$$y = \{ \text{tea} \}$$

A  $\succeq$  is rational if it is

1. complete:  $\forall x, y \in X$  we have  $x \succeq y$ ,  $y \succeq x$ , or both.
2. transitive:  $\forall x, y, z \in X$ , if  $x \succeq y$  and  $y \succeq z$ , then  $x \succeq z$ .

If  $\succeq$  is rational, then

i.)  $>$  is irreflexive  $\Leftrightarrow x > x$  never holds  
and transitive  $\Leftrightarrow$  if  $x > y$  and  $y > z \Rightarrow x > z$

ii.)  $\sim$  is reflexive  $\Leftrightarrow x \sim x$  always holds ( $\forall x$ )  
is transitive  $\Leftrightarrow$  if  $x \sim y$  and  $y \sim z \Rightarrow x \sim z$   
is symmetric  $\Leftrightarrow x \sim y \Leftrightarrow y \sim x$

iii.) if  $x > y$  and  $y \succeq z$ , then  $x > z$

A utility fn. assigns a numerical value to each  $x \in X$  ranking these elements of  $X$  consistently w/ preferences.

Def. A fn.  $u: X \rightarrow \mathbb{R}$  is a utility fn. representing pref. rel.  $\succeq$  if  
—  $\forall x, y \in X, x \succeq y \iff u(x) \geq u(y)$

Prop. A  $\succeq$  can be represented by a utility fn. only if it is rational.

Proof. Show that if  $\exists$  a utility fn. that represents  $\succeq$ , then  $\succeq$  must be (a) complete & (b) transitive.

(a) B/c  $u(\cdot)$  is real-valued on  $X$  it must be  $\forall x, y \in X$ , either  $u(x) \geq u(y)$  or  $u(y) \geq u(x)$ . By def. above  $\Rightarrow$  either  $x \succeq y$ ,  $y \succeq x$ , or both (completeness).

(b) WLOG, suppose  $x \succeq y$  and  $y \succeq z$ . B/c  $u(\cdot)$  represents  $\succeq$  we must have  $u(x) \geq u(y)$  and  $u(y) \geq u(z) \Rightarrow u(x) \geq u(z)$ . B/c  $u(\cdot)$  represents  $\succeq$ , this  $\Rightarrow x \succeq z$  (transitivity).

## Choice

potential available  
choices  
/

A choice structure  $(\beta, C(\cdot))$  consists of

i)  $\beta$  is a set of non-empty subsets of  $X$ .

Every element of  $\beta$  is a set  $B \subset X$ . For ex.  $B \in \beta$  could be a budget set.

ii)  $C(\cdot)$  is a choice correspondence that assigns a non-empty set of chosen elements  $C(B) \subset B$  for every set  $B \in \beta$ . When single-valued, it is an individual's choice from  $B$ .

Review:  $X$  is the choice set

$\beta$  is the potential available choices

$B$  is a subset of  $\beta$

$C(B)$  yields a choice made from  $B$  when single-valued

What kind of restrictions might we need to say a choice structure is reasonable or 'exhibit basic consistency'.

- What does an unreasonable choice structure look like?

Ex. 1 Suppose  $X = \{x, y, z\}$ ,  $\beta = \{\{x, y\}, \{x, y, z\}\}$

Then  $C_1(\{x, y\}) = x$  and  $C_1(\{x, y, z\}) = x$ .

This seems fine.

$x \succ^* y$  and  
 $x \succ^* z$   
 satisfies WARP  
 b/c  $x$  &  $y$  are  
 never chosen,  
 no violation  
 is possible

a.  
 $y \succ^* x$   
 (and  $x \succ^* y$   
 $x \succ^* z$  and  
 $y \succ^* z$ )

Ex. 2. Suppose  $X = \{x, y, z\}$ ,  $\beta = \{\{x, y\}, \{x, y, z\}\}$ .

Then  $C_2(\{x, y\}) = x$  and  $C_2(\{x, y, z\}) = \{x, y\}$ .

This is problematic b/c it seems like the availability of  $z$  (not chosen) affects the choice over  $x$  and  $y$ .

b.  $x \succ^* y$   
 WARP is  
 violated

How can we rule out the type of inconsistency we're seeing in Ex. 2?

① Def.  $\beta, C(\cdot)$  satisfies the weak axiom of revealed preference (WARP) if the following holds

If, for some  $B \in \beta$  w/  $x, y \in B$ , we have  $x \in C(B)$ , then for any  $B' \in \beta$  w/  $x, y \in B'$  and  $y \in C(B')$ , then we require  $x \in C(B')$ .

In words:

If  $x$  is chosen when  $y$  is available in one case, then there cannot be another case where both are available and  $y$  is chosen, but  $x$  is not.

Does this rule out  $G_{x,2}$ ?

If we observe choices, what can we infer their consistency?

Def. Given  $\beta, C(\cdot)$ , the revealed preference rel.  $\succeq^*$  is defined by  
$$x \succeq^* y \iff \exists B \in \beta \text{ such that } x, y \in B \text{ and } x \in C(B).$$

This allows us to restate WARP as follows:

'if  $x$  is revealed at as good as  $y$ , then  $y$  cannot be revealed (strongly) preferred to  $x$ ' (if  $x \succeq^* y$ , then  $y \not\succ^* x$ .)

Lets consider our two examples in revealed preference notation.

Can we go back and forth? Does rationality  $\Rightarrow$  WARP? Does WARP  $\Rightarrow$  rationality?  
Yes. ↗ Sometimes.

This is the important question to us in Micro?

A note on notation: Suppose rational  $\succeq$  on  $X$ . If facing  $B \subset X \Rightarrow$   
 $C^*(B, \succeq) = \{x \in B : x \succeq y \ \forall y \in B\}$  is the preference-maximizing choice correspondence. (This says preference-maximizing behavior is to choose the most preferred alternatives.)



Does 'rationality'  $\Rightarrow$  WARP?

Prop. Suppose that  $\succeq$  is rational. Then  $(\beta, C^*(\cdot, \succeq))$  satisfies WARP.

Proof. Suppose for that some  $B \in \beta$  we have  $x, y \in B$  and  $x \in C^*(B, \succeq)$ .

By def. of  $C^*$  this  $\Rightarrow x \succeq y$ . Now suppose that for  $B' \in \beta$  w/  $x, y \in B'$  we have  $y \in C^*(B', \succeq)$ . This implies  $y \succeq z \forall z \in B'$  by  $C^*$ .

But we know  $x \succeq y \Rightarrow x \succeq z \forall z \in B' \Rightarrow x \in C^*(B', \succeq)$ , but this is WARP by definition!

Does WARP  $\Rightarrow$  rationality?

Def. Given  $(\beta, C(\cdot))$  we say rational  $\succeq$  rationalizes  $C(\cdot)$  relative to  $\beta$  if  $C(B) = C^*(B, \succeq) \forall B \in \beta$ ; that is if  $\succeq$  generates the choice structure  $(\beta, C(\cdot))$  when profit-maximizing choices are made.

Let's look at an example where WARP  $\nRightarrow$  rationality (and see what might be wrong).

Ex. 3 Suppose  $X = \{x, y, z\}$ ,  $\beta = \{\{x, y\}, \{y, z\}, \{x, z\}\}$ ,  $C(\{x, y\}) = x$ ,  $C(\{y, z\}) = y$ , and  $C(\{x, z\}) = z$ .

This satisfies WARP (no opp. for contradiction.)

To have rational  $\succeq$ , we would need  $x \succ y$  and  $y \succ z \Rightarrow x \succ z$  by transitivity, but this is contradicted by the above.

Prop. Arrow (1959)

If  $(B, C(\cdot))$  is a choice structure such that

i) WARP is satisfied

ii)  $\beta$  includes all subsets of  $X$  up to three elements

then  $\exists$  a rational  $\succeq$  that rationalizes  $C(\cdot)$  relative to  $\beta$ ;  
that is  $C(B) = C^*(B, \succeq) \forall B \in \beta$ . Further, this rational  $\succeq$  is  
unique!

Proof: Go through yourself.