

CZ4079 Final Year Project

A Machine Learning-Based Approach to Time-Dependent Shortest Path Queries

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Agenda

- 1 Introduction
- 2 Preliminary Processing
- 3 Landmark Graph
- 4 Travel Time Estimation
- 5 Shortest Path Calculation
- 6 Conclusion and Q&A



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Introduction: Problem



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- A **query** $Q(u, v, t)$ that asks for a shortest path from u to v departing at time moment t

Introduction: General Approach



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- The new **machine learning-based approach** draws on collective wisdom of thousands of taxi drivers
- **Unsupervised learning** is employed to figure out the time-dependent edge costs
- A modified Dijkstra's algorithm calculates a shortest path on the fly

Introduction: Challenges

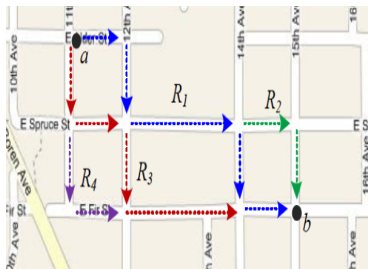


Introduction: Challenges

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- Sparse sample points



Introduction: Challenges

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- Sparse sample points
- Limited GPS accuracy

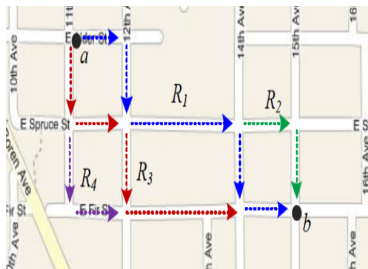


Figure 1: Examples of challenges

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Preliminary Processing: Data Description

- Is collected from Computational Sensing Lab at Tsinghua University
- Contains 83 million GPS records from 8,602 taxis in Beijing during May of 2009

Field	Explanation
CUID	ID for each taxi
UNIX_EPOCH	Unix timestamp
GPS_LONG	Longitude in WGS-84
GPS_LAT	Latitude in WGS-84
HEAD	Heading direction
SPEED	Instantaneous speed (m/s)
OCCUPIED	Hired (1) or not (0)

Table 1: A summary of the seven original fields

Preliminary Processing: Reverse Geocoding

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- GPS coordinate translation: 1.34°N , 103.68°E \rightarrow SCSE, NTU



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- China GPS shift problem: WGS84 v.s. BD09



Figure 2: An example of China GPS shift problem

Preliminary Processing: Reverse Geocoding

- GPS coordinate translation: 1.34°N , 103.68°E \rightarrow SCSE, NTU
- China GPS shift problem: WGS84 v.s. BD09
- Solution: WGS84 $\xrightarrow{\text{Baidu API}}$ BD09 $\xrightarrow{\text{Baidu API}}$ Street



Figure 2: An example of China GPS shift problem

Preliminary Processing: Outlier Detection

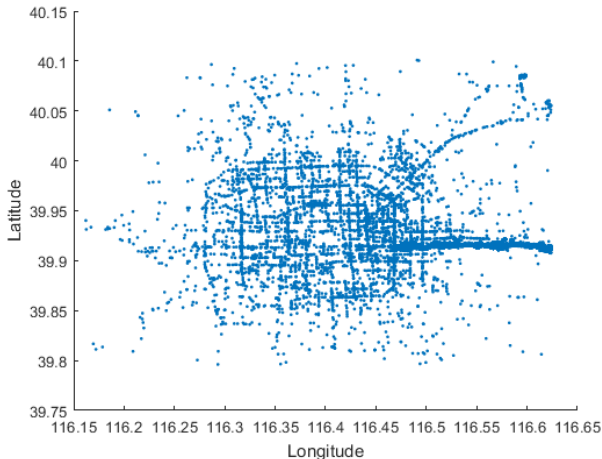


Figure 3: An example of outliers



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Theorem (*Majority Clustering Theorem*)

If a **reasonable reverse geocoder** is used to reverse-geocode a set of GPS data points which are mapped to a particular street *in reality*, then, when plotted on a 2-D plane, majority (more than 50%) of the points must be clustered together to form a rough shape that is similar to the shape of the street that they are supposed to be mapped to.

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Two-step procedure:

Outlier Detection = Outlier Identification + Outlier Removal

Outlier Identification: Clustering

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- Use real physical distance on the Earth
- Set $d_{max} = 30\text{m}$ or 50m

Preliminary Processing: Outlier Detection

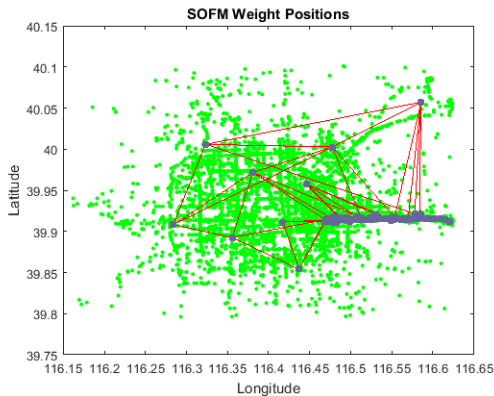


Figure 4: A plot of neuron positions after training



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Preliminary Processing: Outlier Detection

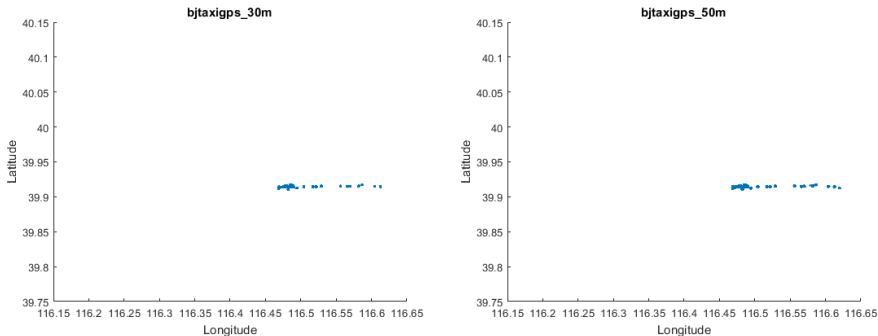


Figure 5: A plot of sample points after outlier removal

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Definition (*Landmark*)

A landmark is a road segment that is frequently traversed by taxi drivers according to the taxi GPS trajectory database.

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Step to build landmark graph

- Separate sample points into **trips**

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- Separate sample points into **trips**
- Count occurrences of each street
- Find connections between two landmarks



Landmark Graph: Trip Identification

CUID	UTC	GPS_LONG	GPS_LAT	OCCUPIED	TRIP_ID
1	1/5/2009 0:02:00	116.39616	39.81294	0	4552265
1	1/5/2009 0:04:00	116.39575	39.82296	0	4552265
1	1/5/2009 0:07:00	116.39567	39.82774	0	4552265
1	1/5/2009 17:08:00	116.30142	39.98105	1	1
1	1/5/2009 17:10:00	116.29514	39.98419	1	1
1	1/5/2009 17:11:00	116.28959	39.98289	1	1
1	1/5/2009 17:12:00	116.28087	39.97552	1	1
1	1/5/2009 17:16:00	116.26813	39.93537	1	1
1	1/5/2009 18:11:00	116.36537	39.95019	0	4552271
1	1/5/2009 18:12:00	116.36546	39.94886	0	4552271
1	1/5/2009 18:13:00	116.35927	39.94528	0	4552271

Table 2: An example of trip identification



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Landmark Graph: Frequency Counting

CUID	UTC	GPS_LONG	GPS_LAT	Street	TRIP_ID
1	1/5/2009 0:02:00	116.39616	39.81294	A	4552265
1	1/5/2009 0:04:00	116.39575	39.82296	A	4552265
1	1/5/2009 0:07:00	116.39567	39.82774	B	4552265
1	1/5/2009 17:08:00	116.30142	39.98105	C	1
1	1/5/2009 17:10:00	116.29514	39.98419	C	1
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Table 3: An illustration of frequency counting



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Landmark Graph: Construction

For each trip

- Select a landmark j
- Record intermediate streets while searching for the next landmark k
- Repeat the process starting from k until all streets are examined

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- Separate weekday's travel time from weekend's

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- Build a **predictive model** for travel time of each significant edge
- Separate weekday's travel time from weekend's
- Evaluate results against Baidu's estimates

Travel Time Estimation: Underlying Distribution

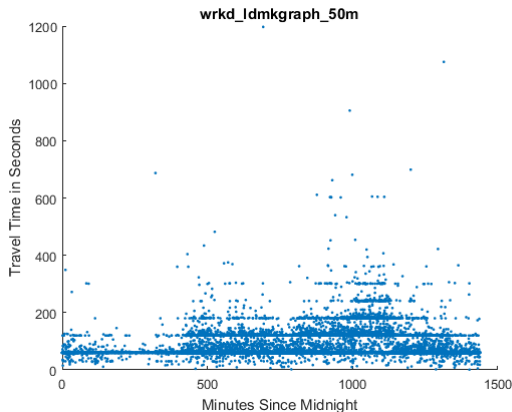


Figure 6: An example of travel time patterns

Travel Time Estimation: Underlying Distribution

Possible Explanations

- Drivers choose different routes to travel between the two landmarks

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Travel Time Estimation: Underlying Distribution

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- Drivers choose different routes to travel between the two landmarks
- Drivers have different driving skills, preferences and behaviours
- The GPS devices report locations **periodically**, therefore, durations like 60 seconds or 120 seconds are very common

Travel Time Estimation: Clustering

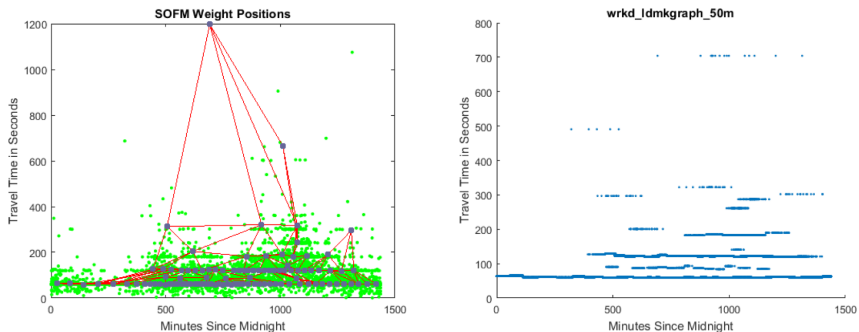


Figure 7: An illustration of travel time clustering

Travel Time Estimation: Distribution Fit

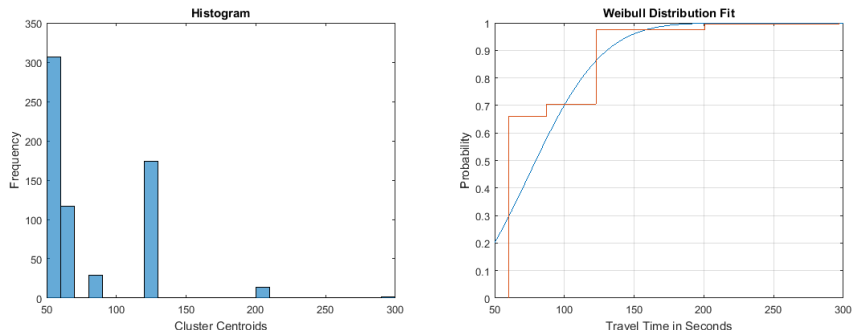


Figure 8: An illustration of fitting distribution

Definition (*Optimism Index*)

An optimism index indicates how optimistic a driver feels about his or her driving skills. A driver with an optimism index of $p\%$ usually drives faster than $(1 - p)\%$ drivers under similar road conditions.

- Calculate and store α and β for each 30-minute window

Travel Time Estimation: Implementation

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- Calculate and store α and β for each 30-minute window
- Use optimism index p to find out travel time

Travel Time Estimation: Evaluation

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Landmark Graph	RMSE	Mean Error Ratio	Mean No. of Samples Per Edge
wrkd_ldmkgraph_50m	78.84	-0.009	1824.60
wrkd_ldmkgraph_30m	87.96	-0.065	1507.56
holi_ldmkgraph_50m	87.39	-0.16	832.96
holi_ldmkgraph_30m	76.41	-0.14	681.89

Table 4: A summary of evaluation results



Travel Time Estimation: Evaluation

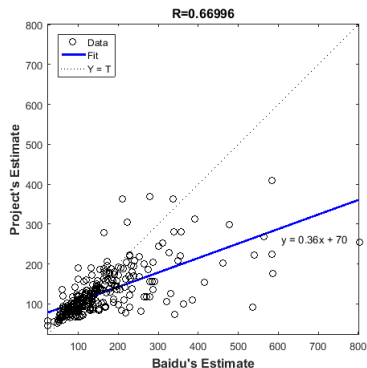
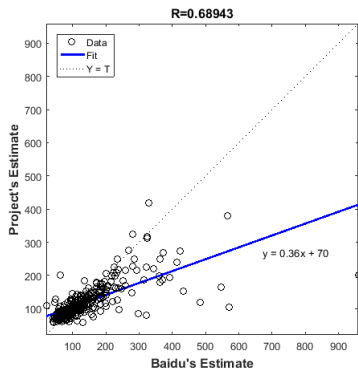


Figure 9: A plot of linear regression for weekday landmark graph

Travel Time Estimation: Evaluation

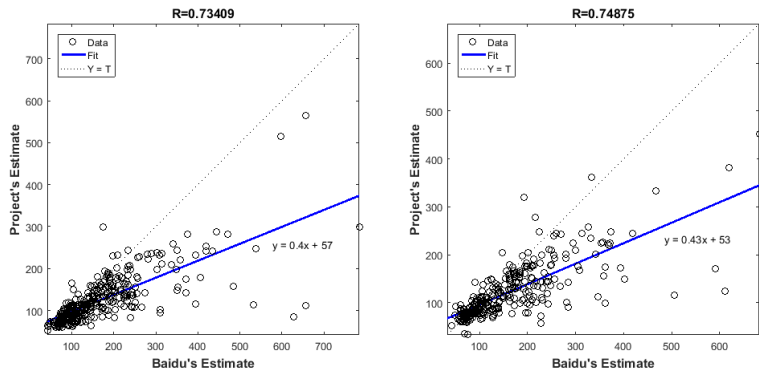


Figure 10: A plot of linear regression for weekend landmark graph

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Shortest Path Calculation: Basic Ideas

Definition (*FIFO Graph*)

A time-dependent graph $G = (V, E)$ with a dynamic weight function $w : E, t \rightarrow \mathbb{R}$ is a FIFO graph iff for every edge $(u, v) \in E$

$$\forall \Delta t \geq 0, \quad w(u, v, t_0) \leq \Delta t + w(u, v, t_0 + \Delta t) \quad (1)$$

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- Consider a landmark graph as a FIFO graph (usually true)

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- Consider a landmark graph as a FIFO graph (usually true)
- Modify Dijkstra's algorithm to calculate edge costs in an online fashion

Shortest Path Calculation: Modified Dijkstra's Algorithm

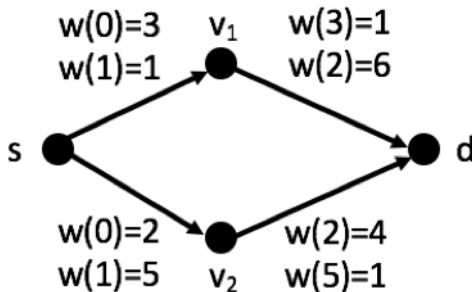


Figure 11: An example of modified Dijkstra's algorithm

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Conclusion and Q&A: Conclusion

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- Identified challenges in solving the problem



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- Identified challenges in solving the problem
- Proposed an innovative outlier detection approach
- Introduced algorithms for building landmark graphs
- Described procedures for estimating travel time
- Illustrated the modified Dijkstra's algorithm

Any questions?