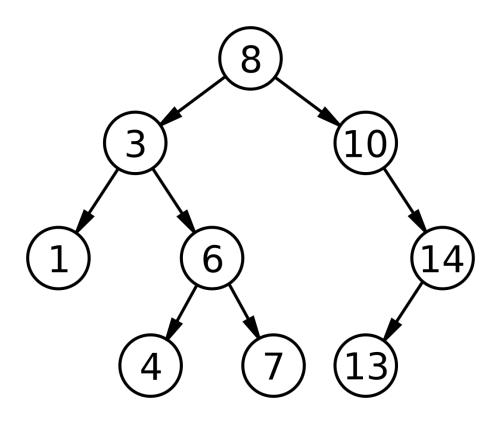
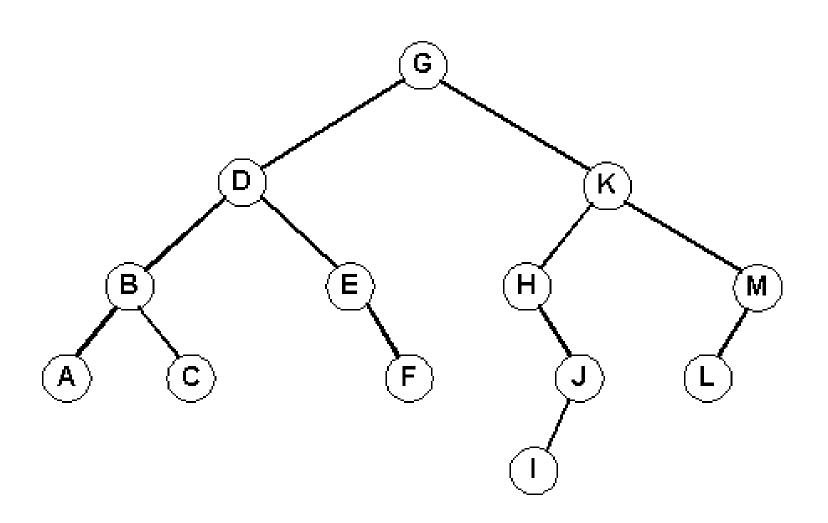
- A binary search tree (BST) is a binary tree in which
  - The values in the left subtree of a node are all less than the value in the node
  - The values in the right subtree of a node are all greater than the value of the node.
  - The subtrees of a binary search tree must themselves be binary search trees.



Left <= Node <= Right



#### **Properties and Operations**

- Note that under this definition, <u>a BST never</u> contains duplicate nodes.
- Some operations for BSTs:
  - InsertItem
  - Deleteltem
  - ItemExists
  - Traverse(Pre, In, Post)
  - Count
  - Height

### Binary Tree Algorithms

```
struct Node{
   Node *left;
   Node *right;
   int data;
};
Node *MakeNode(int Data)
   Node *node = new Node;
   node->data = Data;
   node->left = ∅;
   node->right = ∅;
   return node;
}
```

```
void FreeNode(Node *node){
   delete node;
typedef Node* Tree;
```

### Finding an Item in a BST

 State the recursive algorithm in English for finding an item.

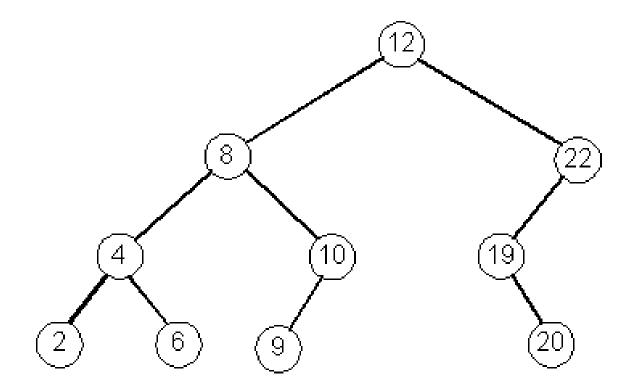
```
bool ItemExists(Tree tree, int Data){
   if (tree == 0)
        return false;
   else if (Data == tree->data)
        return true;
   else if (Data < tree->data)
        return ItemExists(tree->left, Data);
   else
        return ItemExists(tree->right, Data);
}
```

#### Insert an Item in a BST

State the recursive algorithm in English for inserting an item.

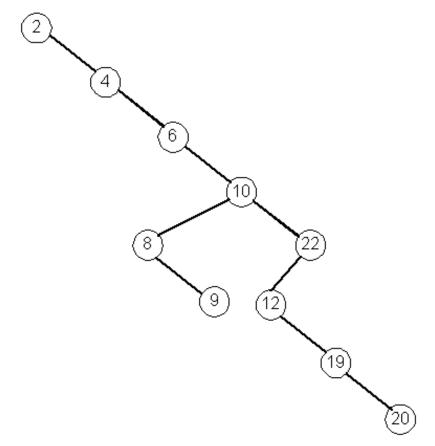
#### Insert an Item in a BST

- Create a tree using these values (in this order):
  - 12, 22, 8, 19, 10, 9, 20, 4, 2, 6



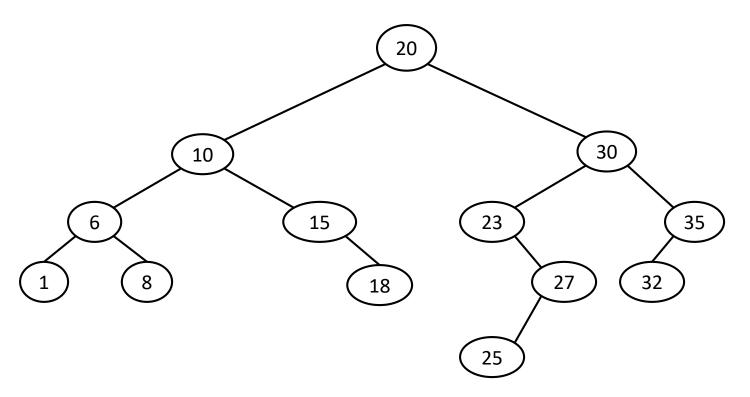
#### Insert an Item in a BST

- Create a tree using these values (in this order):
  - 2, 4, 6, 10, 8, 22, 12, 9, 19, 20



#### Deleting a Node

 The caveat of deleting a node is that, after deletion, the tree must still be a BST.



### Deleting a Node

Case 1: The node to be deleted is a leaf node.

- Case 2: The node to be deleted has an empty left child but non-empty right child.
- Case 3: The node to be deleted has an empty right child but non-empty left child.

Case 4: The node has both children non-empty.

#### Case 1: Leaf Node

- Set the parent's pointer to this node to NULL.
- Release the memory of the leaf node.

## Case 2: Empty Left Child

Replace the deleted node with its right child.

## Case 3: Empty Right Child

Replace the deleted node with its left child.

#### Case 4: Non-empty Left and Right Child

- Replace the data in the deleted node with its predecessor under in-order traversal.
- Delete the node that holds the predecessor.

#### Deleting a Node

```
void DeleteItem(Tree &tree, int Data){
    if (tree == ∅) return;
    else if (Data < tree->data)
         DeleteItem(tree->left, Data);
    else if (Data > tree->data)
         DeleteItem(tree->right, Data);
    else { // (Data == tree->data)
         if (tree->left == 0){
              Tree temp = tree;
              tree = tree->right;
              FreeNode(temp);
         else if (tree->right == 0){
              Tree temp = tree;
              tree = tree->left;
              FreeNode(temp);
         else{
              Tree pred = 0;
              FindPredecessor(tree, pred);
              tree->data = pred->data;
              DeleteItem(tree->left, tree-
              >data);
```

```
void FindPredecessor(Tree tree, Tree
&predecessor){
    predecessor = tree->left;
    while (predecessor->right != 0)
        predecessor = predecessor->right;
}
```

- We are replacing the data in the node, not the node itself
- Because the predecessor comes from the left subtree:
  - It must be <u>less</u> than everything in the right subtree.
  - It must be <u>greater</u> than everything in the left subtree.
- The recursive deletion of the predecessor's node will lead to one of the simpler cases.

## Summary

- Binary tree
- Binary search tree
  - Finding an item
  - Insertion
  - Deletion

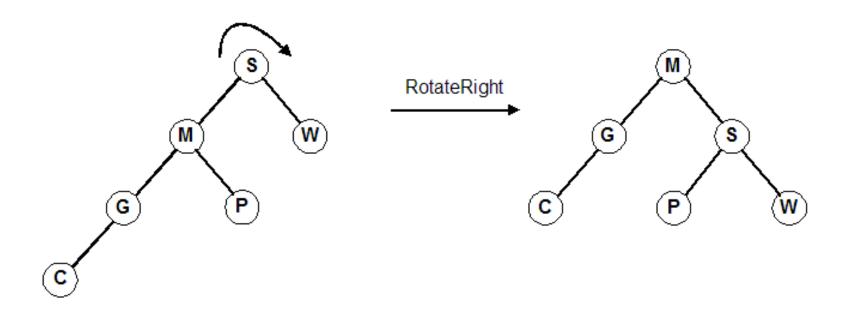
## **Rotating Nodes**

#### **Rotating Nodes**

- Rotation is a fundamental technique performed on BSTs.
- Two types of rotations:
  - Left
  - Right
- Promoting a node is the same as rotating around the node's parent.
- There is no direction in promotion.
- You can rotate about any node that has children.
- After the rotation, the sort order is preserved.
  - The resulting is STILL a BST

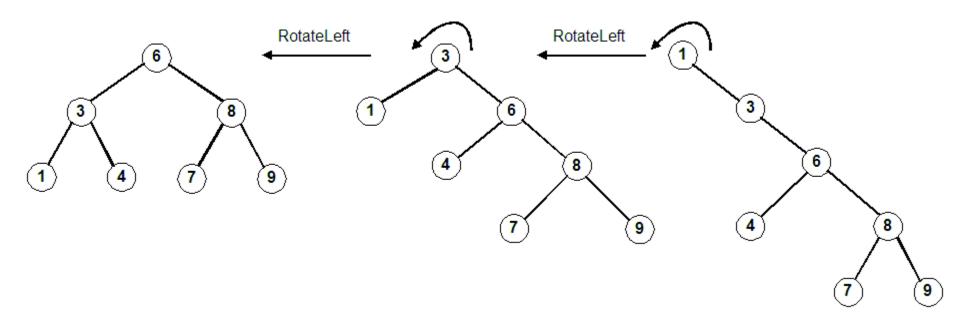
#### Right Rotation

Rotate right around the root, S (Same as promoting M)



#### Left Rotation

Rotate left twice around the root. First around
1, then around 3. (Same as promoting 3 then
6)



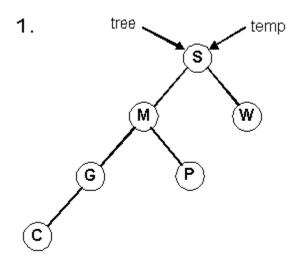
## Rotating Left and Right

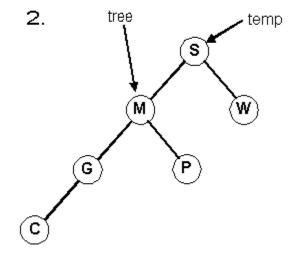
```
void RotateRight(Tree &tree){
    Tree temp = tree;
    tree = tree->left;
    temp->left = tree->right;
    tree->right = temp;
void RotateLeft(Tree &tree){
   Tree temp = tree;
   tree = tree->right;
   temp->right = tree->left;
   tree->left = temp;
}
```

## Step by Step

1.Tree temp = tree;

2.tree = tree->left;

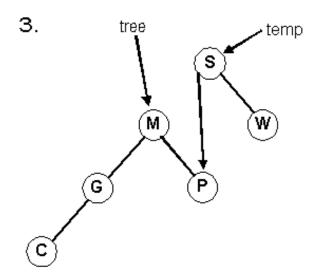


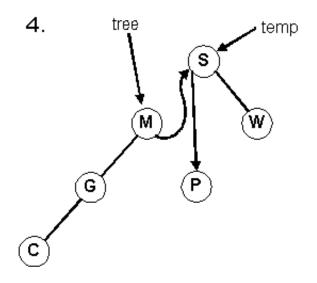


## Step by Step

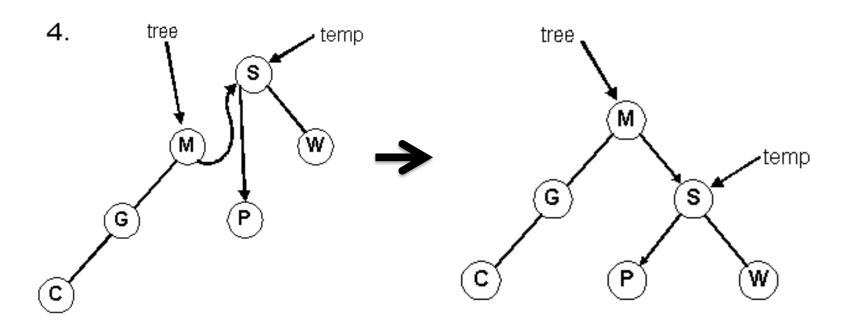
temp->left = tree->right;

tree->right = temp;





## Adjusting the Diagram



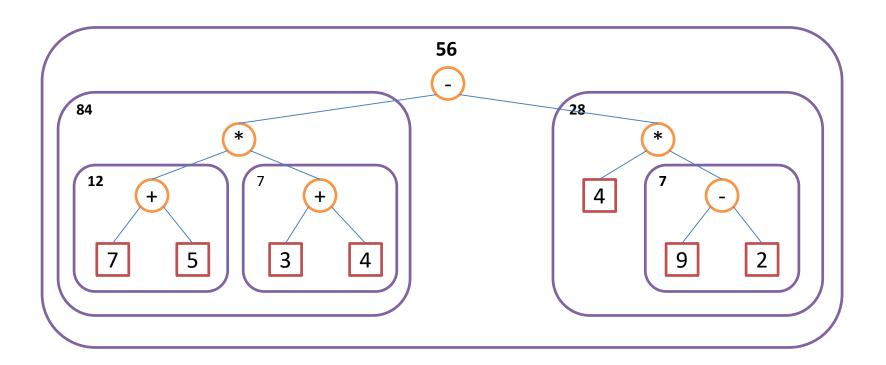
# **Expression Trees**

#### **Expression Trees**

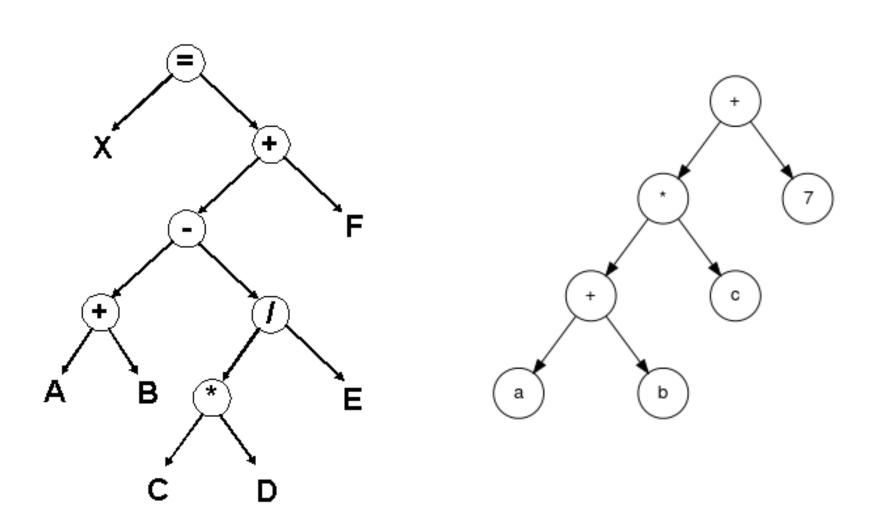
 Expression trees are like binary trees, but they are not "sorted" in the usual way.

- Expression trees are a way to
  - Solve arithmetic expressions.
  - Compile a language

## **Expression Tree**



# **Expression Trees**



## Evaluate an expression tree

```
int evaluate(Tree node){
1.
        if(node->type == NODE OPERAND)
2.
3.
              return node->opd;
        else {
4.
5.
              int left = evaluate(node->left);
6.
              int right = evaluate(node->right);
7.
8.
              switch(node->opt){
9.
                        case '+': return (left+right);
                         case '-': return (left-right);
10.
                        case '*': return (left*right);
11.
                         case '/': return (left/right);
12.
                         default: return -1;
13.
14.
15.
16.
```

• We only consider binary operators here and the list of operators you need to consider includes "+", "-", "\*" and "/".