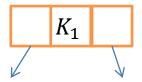
## 2-3 Search Tree

#### 2-3 Search Trees

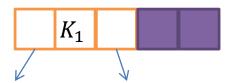
- Each node can contain 1 or 2 keys.
- Each node has 2 or 3 children, hence 2-3 trees.
- The keys in the nodes are ordered from small to large.
- All leaves are at the same (bottom most) level, meaning we always add at the bottom.

### 2-3 Search Tree Node

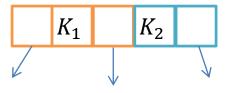
```
struct Node23{
  Node23 *left, *middle, *right;
  Key key1, key2;
};
```



2-node (not showing empty)



2-node (showing empty)



3-node

$$K_1 < K_2$$

## Properties of 2-3 Search Trees

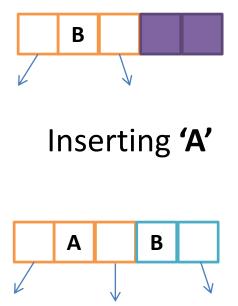
• 2-3 search trees guarantee to be balanced at all times.

Searches are O(log N) in worst case.

- Balance is maintained during insertion
  - Splitting nodes, worst case O(log N), average case
     O(1)

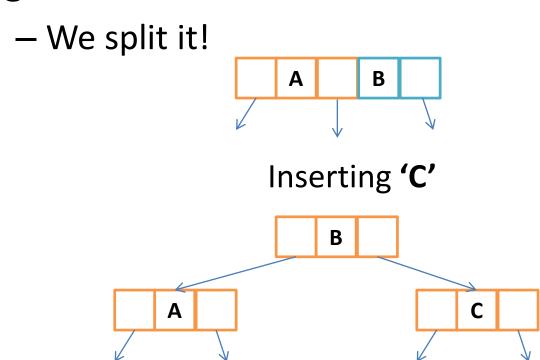
#### Insertion

 If the node you insert is a 2-node, simply grow the node to a 3-node



#### Insertion

• If the node you insert is a 3-node, we cannot grow the node more



#### Insertion

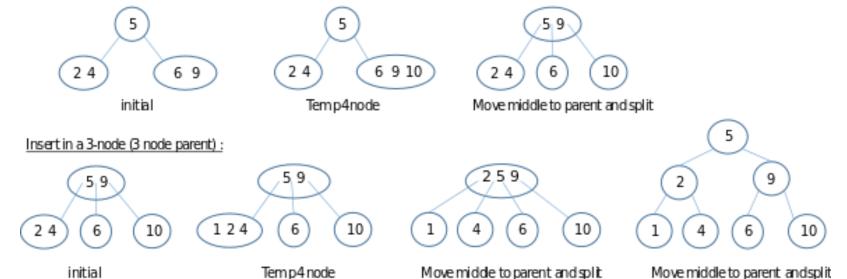
- Splitting this way is called bottom-up balancing
  - Insert the node at the bottom-most level at correct location.
  - If the node is a 3-node, split it and pass the middle key to the parent.
    - If the parent is also a 3-node, split the parent and pass the middle key up
      - Etc...
  - Eventually, the root will also be a 3-node and splitting it will grow the tree one level.

#### **Insertion - Cases**

#### Insert in a 2-node:

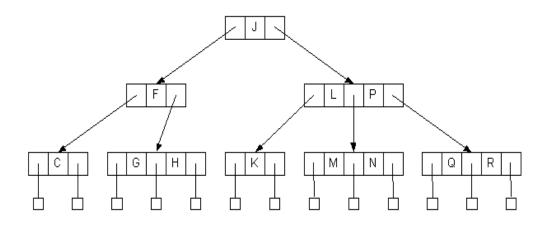


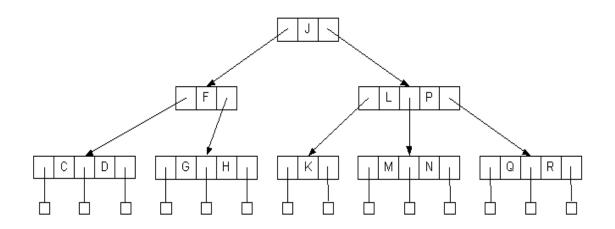
#### Insert in a 3-node (2 node parent):



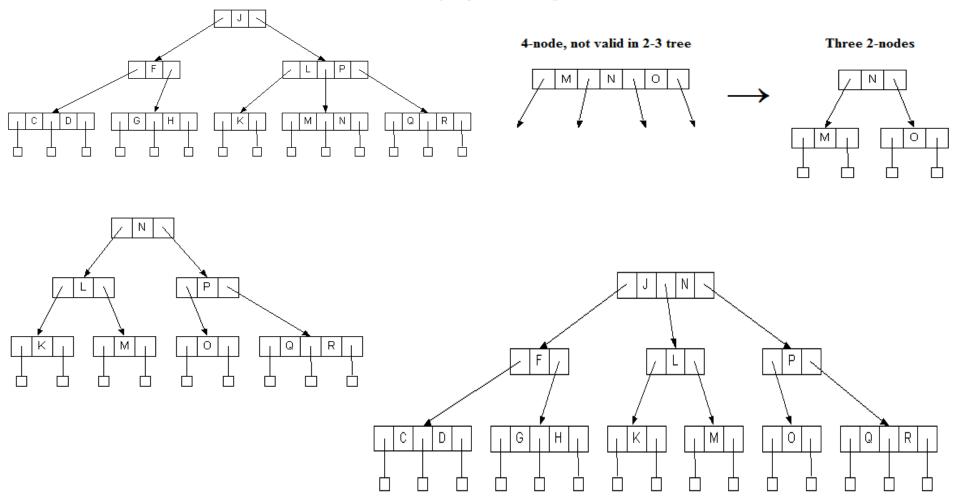
# Example

### Insert D





#### Insert O



#### **Practice**

- Build a BST with 50, 70, 20, 60, 40, 10
- Build a 2-3 tree with 50, 70, 20, 60, 40, 10

What is the height of the tree in both the trees?

 Height of the tree is proportional to the access time of a nodes

## Summary

- 2-3 Trees are always balanced.
- Nodes are ALWAYS inserted at the bottom-most level.
- Balance is maintained by splitting full nodes and passing up the middle node.
- This makes the tree's height increase by one, only when the root is split.

## 2-3-4 Search Tree

## **Basic Properties**

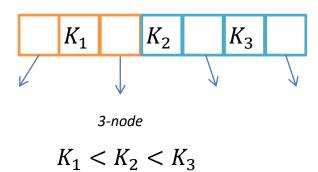
- Similar to 2-3 trees
- Nodes can contain 1, 2, or 3 keys.
- Nodes can has 2, 3, 4 children, hence 2-3-4 tree.
  - Each can have <u>at most</u> 4 children.
- Similarly to 2-3 trees, 2-3-4 trees are guaranteed to be always balanced.
- Balancing algorithm also relies on Splitting nodes
- Number of splits in the worst-case is O(log N)
  - When is the worst-case?
- Average number of splits is very few.

## **Balancing Algorithm**

- Balancing also occurs on insertion.
- Modifying the algorithms for balancing can produce better efficiency.
- Previously, with 2-3 Trees, we have seen bottom-up balancing.
- We will see top-down balancing
  - As you go down the tree to insert a node, <u>split any full node</u>.
  - A full node is a 4-node.

#### 2-3-4 Node

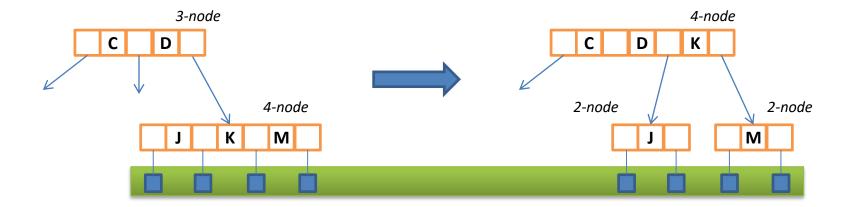
```
struct Node234
{
  Node23 *left, *midleft, *midright, *right;
  Key key1, key2, key3;
};
```



How to split a 2-3-4 node?

## Advantage of Splitting 2-3-4 Trees

- Splitting a node is cleaner.
- Splitting a 4-node into two 2-nodes preserves the number of child links.
- Changes do not have to be propagated. Change remains local to split.



## **Top-Down Balancing**

- Split nodes on the way down.
  - Guarantees that each node we pass through is not a 4-node.
  - When we reach the bottom, we will not be on a 4-node (think about it)
- This way, we only traverse the tree once, when inserting/balancing.
- After each insertion, check if the root is a 4-node
  - If it is, split it directly. This will avoid to do it at next insertion.
  - Splitting the root is the only way to grow the tree.

## Example

Build a 2-3-4 tree with the sequence: 3, 1, 5,
 19, 15, 20, 13, 10, 4, 17, 18

#### Practice

Build a 2-3-4 tree with the sequence: 6, 3, 9, 4,
5, 8, 11, 2, 1, 7, 10, 12, 14.

#### **Next Lecture**

Red-Black Tree