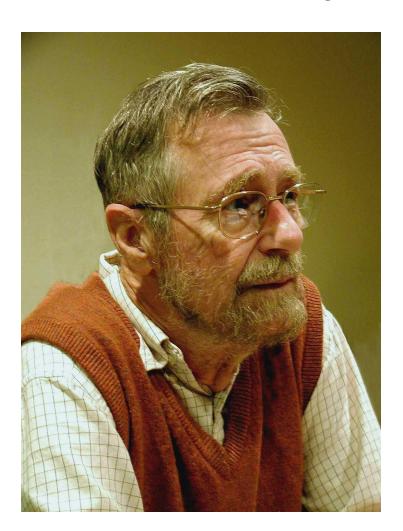
CS380 Artificial Intelligence for Games

Dijkstra's Search In Explicit Graph



Edsger Wybe Dijkstra 1930 – 2002

Main contributions:

- Dijkstra's algorithm
- Semaphore

Recipient of Turing Award (1972)

- Finds the shortest path from the source node to all other nodes in the graph with non-negative edge costs.
- Is a Best-first Search Algorithm makes the optimal choice at each step as it attempts to find the overall optimal way to solve the entire problem
- Similar to Uniform-Cost Search
 - Both use openlist implemented as a priority queue
 - DS finds shortest paths to all other nodes while UCS usually has one target node to reach
 - DS has slightly higher computational cost than UCS

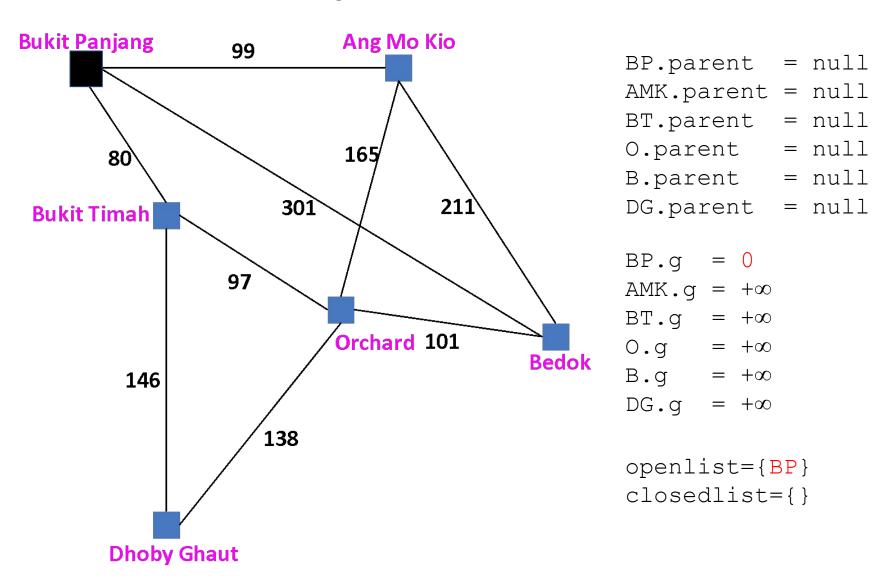
Dijkstra's Search. Initialization

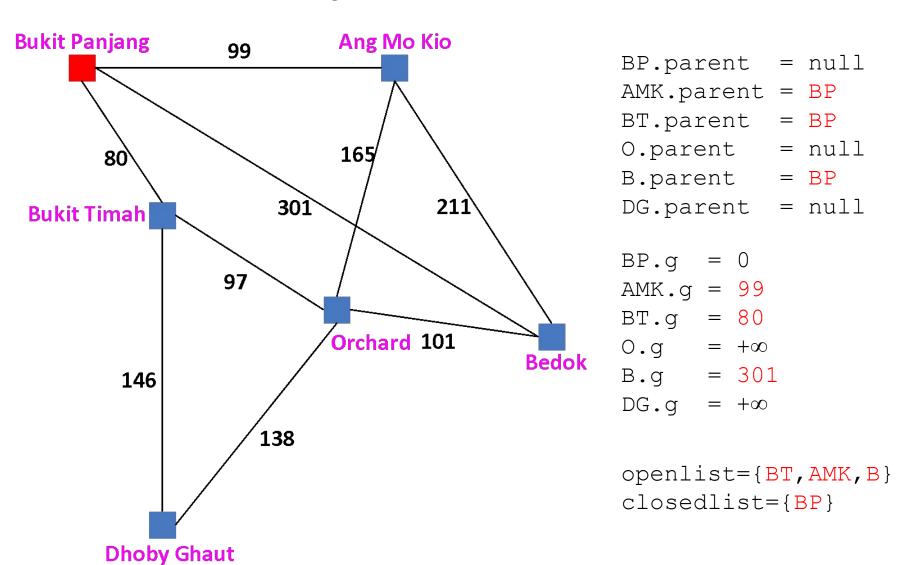
- Given:
 - graph,
 - starting and target nodes
 - empty openelist as priority queue
 - closelist as a hash table
- Initialization step:

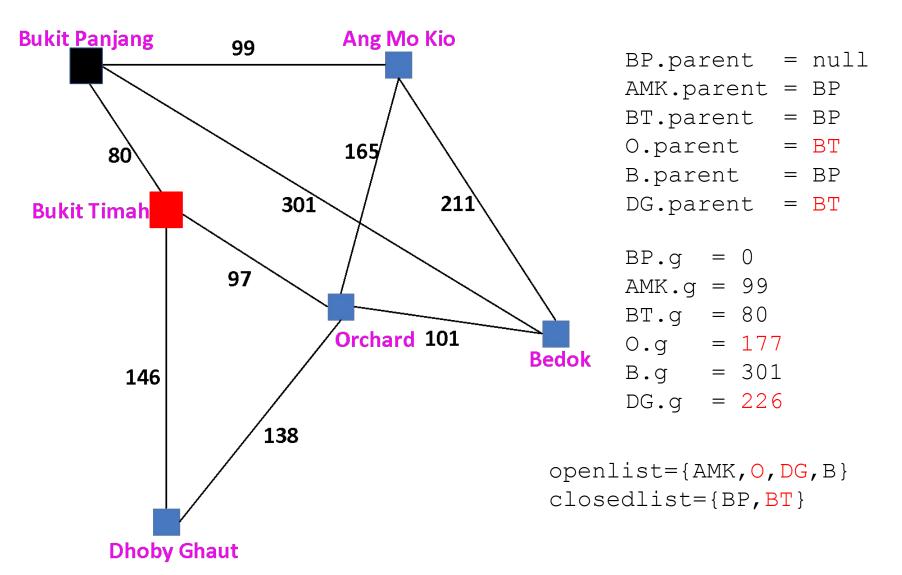
```
for_each (node of graph) {
   node.g = +∞;  // Calculated distance is set as infinite
   node.parent = null; // Previous node is undefined
}
starting.g = 0;
openlist.push(starting);
```

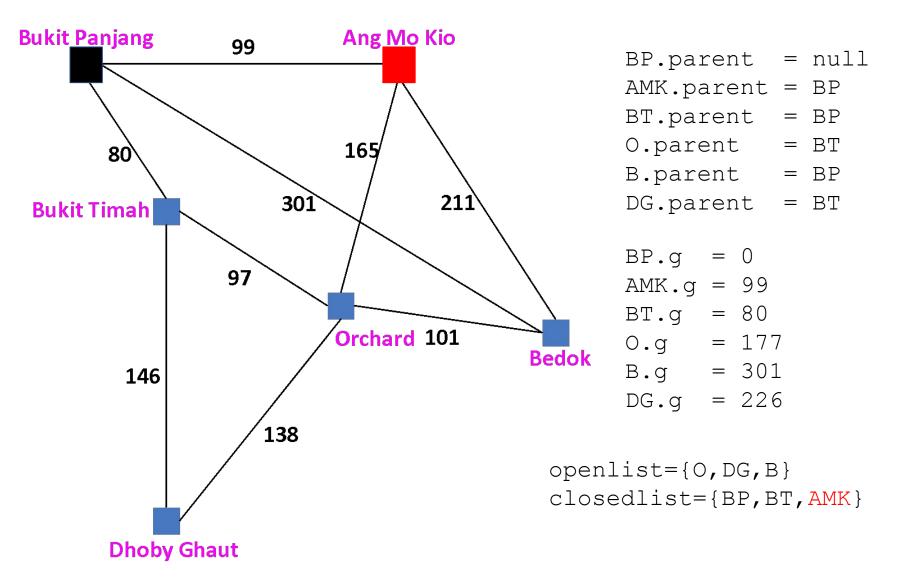
Dijkstra's Search. Main loop

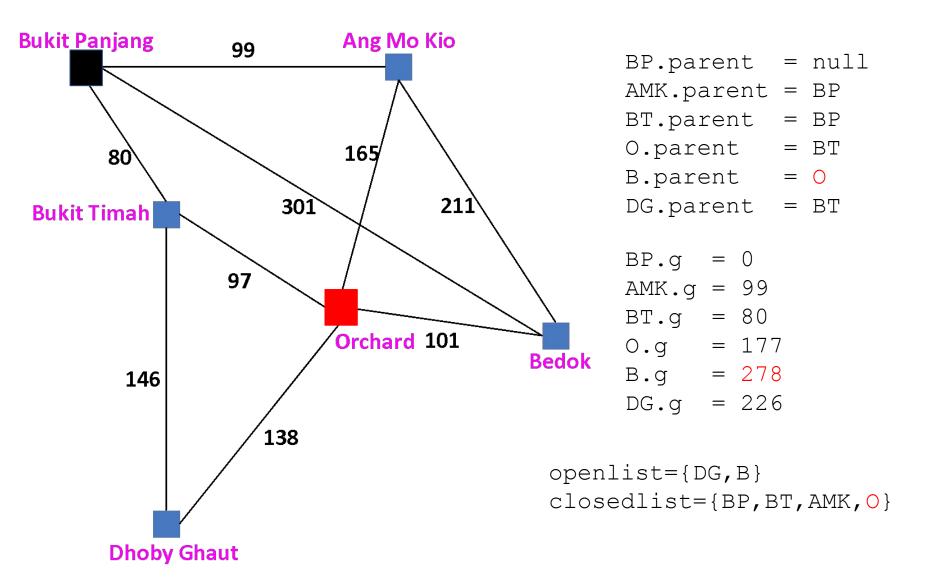
```
while (!openlist.empty()) {
   current = openlist.pop();
   closedlist.add(current); // Set as "visited"
   for each(adjacent of current)
       if (!closedlist.find(adjacent)) {
           if (!openlist.find(adjacent)) {
              adjacent.g = current.g + cost(current, adjacent);
              adjacent.parent = current;
              openlist.push (adjacent);
           } else {
              tentative g = current.g + cost(current, adjacent);
               if (tentative g < adjacent.g) {</pre>
                  adjacent.g = tentative g;
                  adjacent.parent = current;
```

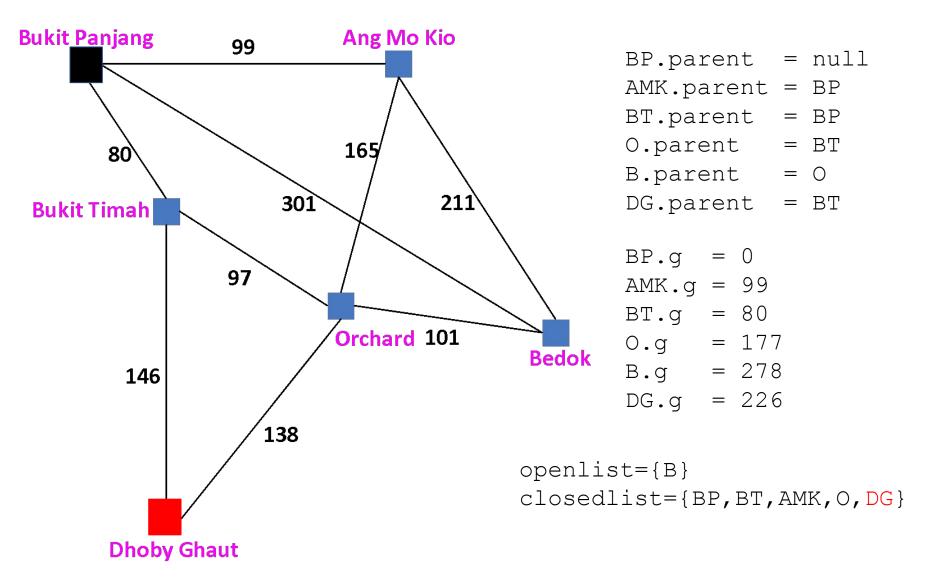


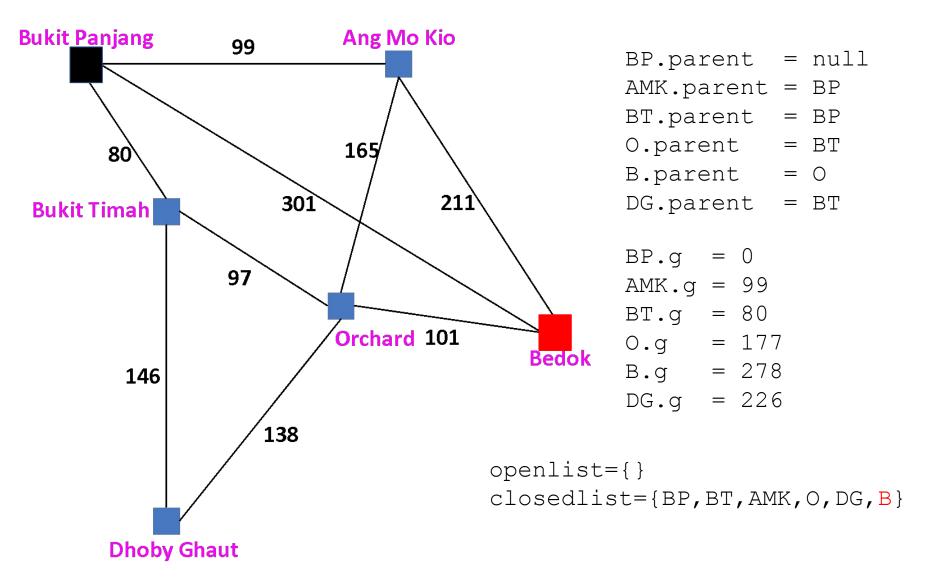












 If we are only interested in a shortest path between nodes starting and target, we can terminate the search after line

```
closedlist.add(current); // Set as "visited"
```

by testing current and target:

```
if (current == target)
  break;
```

 Now we can read the shortest path from source to target by reverse iteration:

```
list path = {};
node = target;
while (node)
{
  path.push(node);
  node = node.parent;
}
path.pop(); // Optional, remove starting point
path.reverse();
```

Complete?

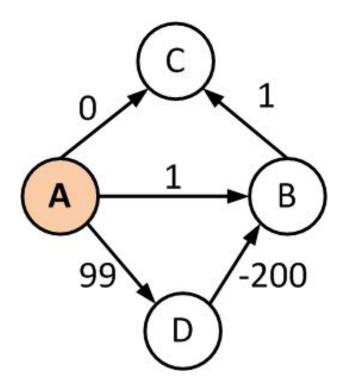
– Yes.

Optimal?

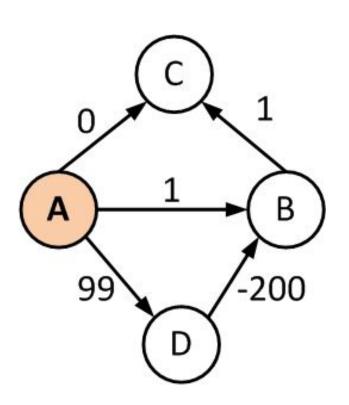
- Yes, if all states' actions have non-negative costs.
- It expands nodes in order of their optimal path cost.

 Time & Space complexity in terms of number of vertices n:
 O(n)

 Note: it may not work on graph with negative cost values



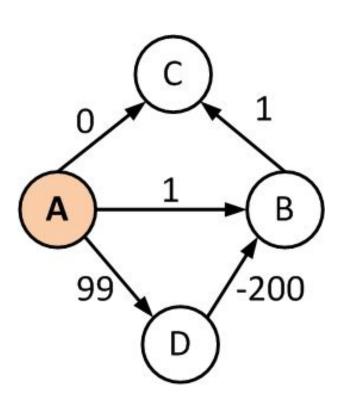
```
A.parent = null
B.parent = null
C.parent = null
D.parent = null
A.g = 0
B.q = +\infty
C.q = +\infty
D.q = +\infty
openlist={A}
closedlist={}
```



```
A.parent = null
B.parent = A
C.parent = A
D.parent = A

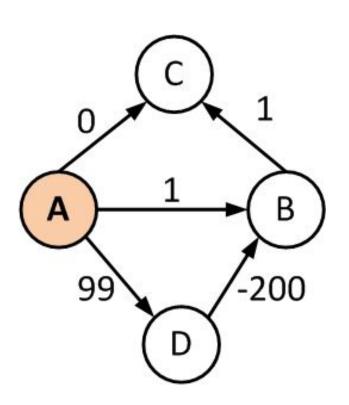
A.g = 0
B.g = 1
C.g = 0
D.g = 99

openlist={C,B,D}
closedlist={A}
```



```
A.parent = null
B.parent = A
C.parent = A
D.parent = A
D.parent = A
A.g = 0
B.g = 1
C.g = 0
D.g = 99

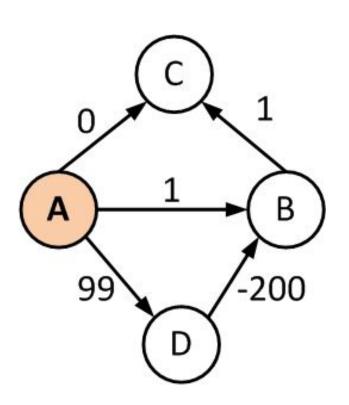
openlist={B,D}
closedlist={A,C}
```



```
A.parent = null
B.parent = A
C.parent = A
D.parent = A

A.g = 0
B.g = 1
C.g = 0
D.g = 99

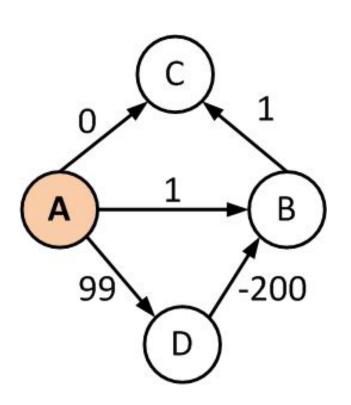
openlist={D}
closedlist={A,C,B}
```



```
A.parent = null
B.parent = D
C.parent = A
D.parent = A

A.g = 0
B.g = -101
C.g = 0
D.g = 99

openlist={}
closedlist={A,C,B,D}
```



```
A.parent = null
B.parent = D
C.parent = A
D.parent = A

A.g = 0
B.g = -101
C.g = 0 (should be -100)
D.g = 99

openlist={}
closedlist={A,C,B,D}
```

 Question: why not reduce computing shortest paths with negative edge lengths to the same problem with non-negative lengths? (by adding large constant to edge lengths)

