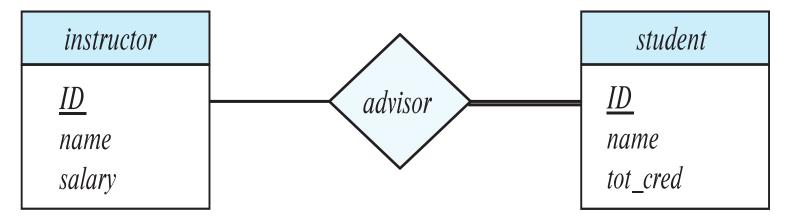
Lecture 12 - Database Modelling & Database Design (contd...)

CS211 - Introduction to Database

Participation Constraints

■ **Total participation** (**indicated by double line**): every entity in the entity set participates in at least one relationship in the relationship set. It represents **NOT NULL** constraint.



participation of student in advisor relation is total

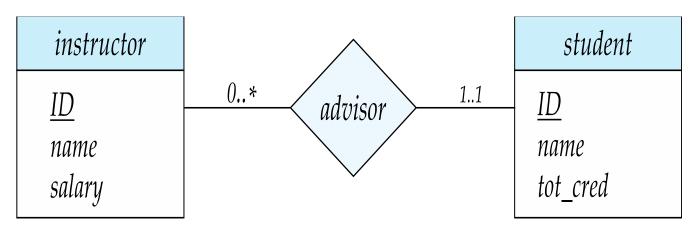
Example: every *student* must have an associated instructor

■ Partial participation: some entities may not participate in any relationship in the relationship set. It represents the attribute values can be NULL.

Example: participation of *instructor* in *advisor* is partial

Notation for Expressing More Complex Constraints

- A line may have an associated minimum and maximum cardinality, shown in the form I..h, where I is the minimum and h the maximum cardinality
 - A minimum value of 0 indicates partial participation.
 - A minimum value of 1 indicates total participation.
 - A maximum value of 1 indicates that the entity participates in at most one relationship.
 - A maximum value of * indicates no limit.
- Example



• Instructor can advise ZERO or more students. A student must have 1 advisor; cannot have multiple advisors

Primary Key Constraints

- Primary keys provide a way to specify how entities and relations are distinguished. We will consider:
 - Entity sets
 - Relationship sets.
 - Weak entity sets

Primary Key for Entity Sets

- By definition, individual entities are distinct.
- From database perspective, the differences among entities must be expressed in terms of their attributes.

■ A **key** for an entity is a set of attributes that suffice to distinguish entities from each other.

Primary Key for Relationship Sets

- To distinguish among the various relationships of a relationship set we use the individual primary keys of the entities in the relationship set.
 - \circ Let R be a relationship set involving entity sets $E_1, E_2, ..., E_n$.
 - Let $primary-key(E_i)$ denote the set of attributes that forms the primary key for entity set E_i .
- **Case-1:** If the relationship set R has no attributes associated with it, then the set of attributes primary-key(E_1) \cup primary-key(E_2) \cup · · · \cup primary-key(E_n) makes the primary key for R.
- Case-2: If the relationship set R has attributes a_1, a_2, \ldots, a_m associated with it, then the set of attributes $primary-key(E_1) \cup primary-key(E_2) \cup \cdots \cup primary-key(E_n) \cup \{a_1, a_2, \ldots, a_m\}$ makes the primary key for R.

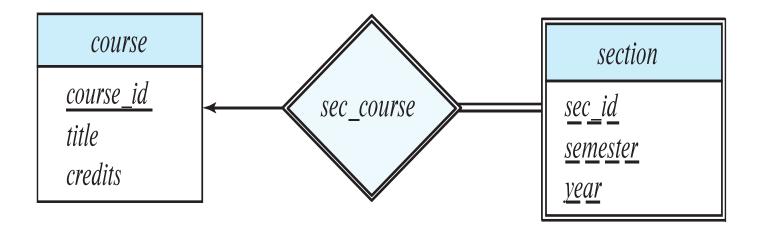
Weak Entity Sets

- An entity set that has a primary key is termed a **strong entity set**.
- An entity set that does not have sufficient attributes to form a primary key is termed a weak entity set.
- For a weak entity set to be meaningful, it must be associated with another entity set, called the identifying or owner entity set.
- The weak entity set is said to be **existence dependent** on the identifying entity set.
- The relationship associating the weak entity set with the identifying entity set is called the **identifying relationship**.

Weak Entity Sets (Discriminator)

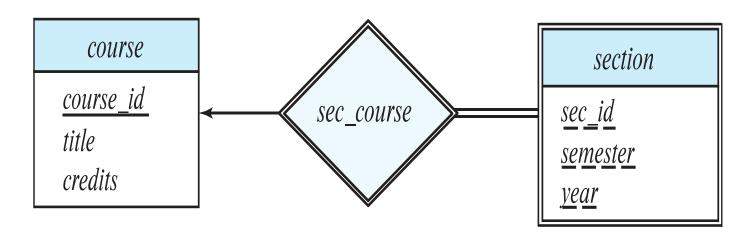
■ The discriminator of a weak entity set is a set of its attributes which combined with the primary key of the identifying entity set makes the primary key of weak entity set.

Primary-key of Weak entity set = Discriminator + Primary-key of owner



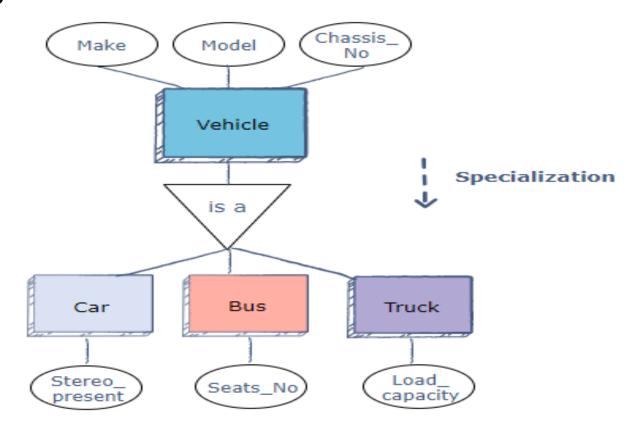
Representing Weak Entity Sets in E-R Diagram

- In E-R diagrams, a weak entity set is depicted via a double rectangle.
- We underline the discriminator of a weak entity set with a dashed line.
- The **identifying relationship set** connecting the weak entity set to the identifying strong entity set is depicted by a **double diamond**.
- Primary key for the weak entity set "section" is: (course_id, sec_id, semester, year)



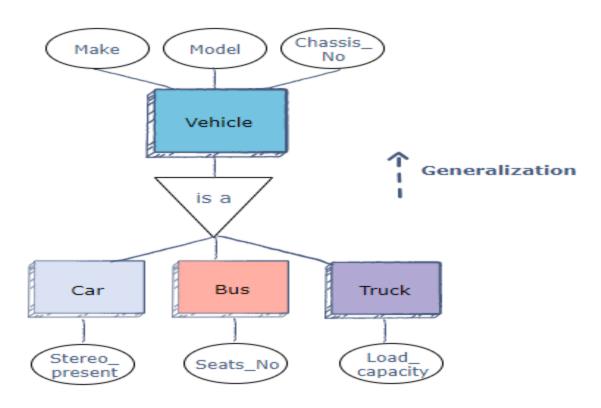
Specialization

- **Specialization** is a **top-down** approach in which a higher-level entity is **divided** into multiple **specialized** lower-level entities.
- In addition to **sharing the attributes** of the higher-level entity, these lower-level entities have **specific attributes** of their own.



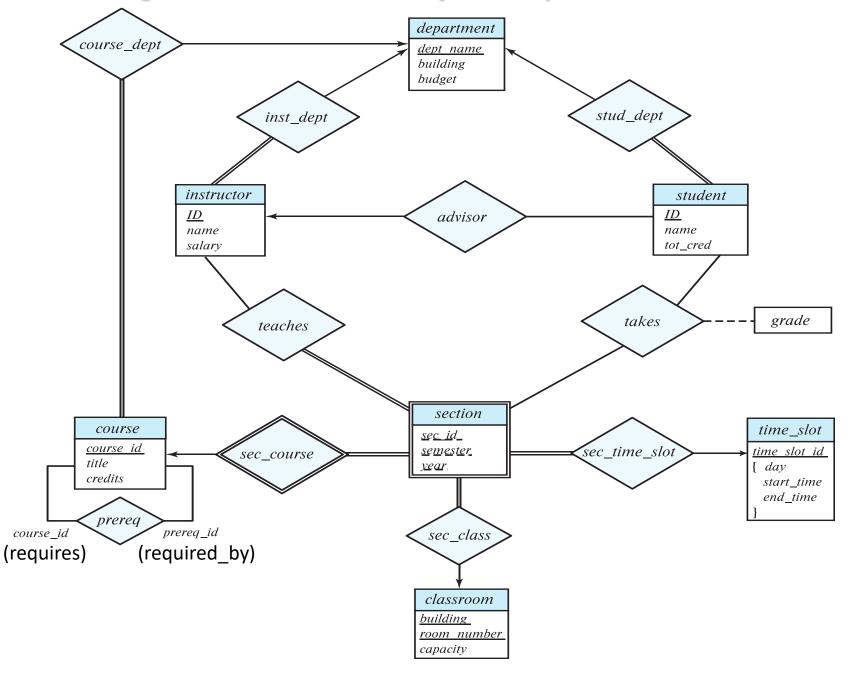
Generalization

• **Generalization** is a **bottom-up** approach in which multiple specialized lower-level entities are combined to form a single **higher-level entity**.



 Generalization is usually used to find common attributes among entities to form a generalized entity.

E-R Diagram for a University Enterprise



- Entity set
- Attribute type
- Entity set primary-key
- Relationship set
- Relationship set primary-key
- Relationship attribute
- Entity role
- Degree of relationship
- Cardinality constraint
- Participation constrain
- Weak entity
- Identifying relationship
- Owner entity set
- Discriminator

Database Design

Reduction to Relation Schemas

• Entity sets and relationship sets are expressed as relation schemas that represent the contents of the database.

 For each entity set and relationship set there is a unique schema that is assigned the name of the corresponding entity set or relationship set.

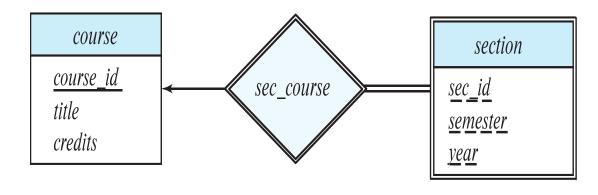
 Each schema has a number of columns (generally corresponding to attributes), which have unique names.

Reduction of Entity Sets

- A strong entity set reduces to a schema with the same attributes.
 - the primary key of the entity set serves as the primary key of the resulting schema



For schemas derived from a weak entity set, the combination of the primary key of the strong entity set and the discriminator of the weak entity set serves as the primary key of the schema.



section (course_id, sec_id, sem, year)

Reduction of Entity Sets

■ **Composite attributes** are flattened out by creating a separate attribute for each component attribute.

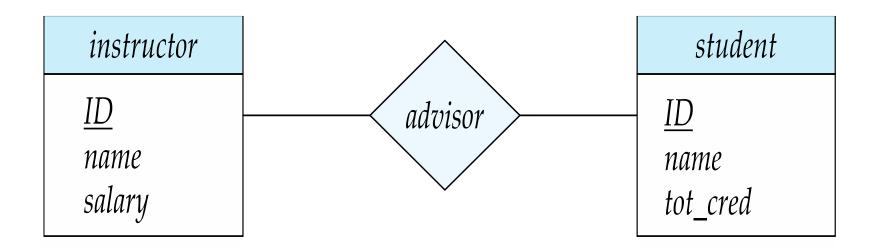
instructor IDname first_name middle_initial last_name address street street_number street name apt_number city state zip { phone_number } date_of_birth age()

Reduction of Entity Sets

- A multivalued attribute M of an entity E is represented by a separate schema EM.
- Schema EM has attributes corresponding to the primary key of E and an attribute corresponding to multivalued attribute M.
- Example: Multivalued attribute phone_number of instructor is represented by a schema:
 inst_phone= (ID, phone_number)
- Each value of the multivalued attribute maps to a separate tuple of the relation on schema EM.
 - For example, an *instructor* entity with primary key 22222 and phone numbers 456-7890 and 123-4567 maps to two tuples:

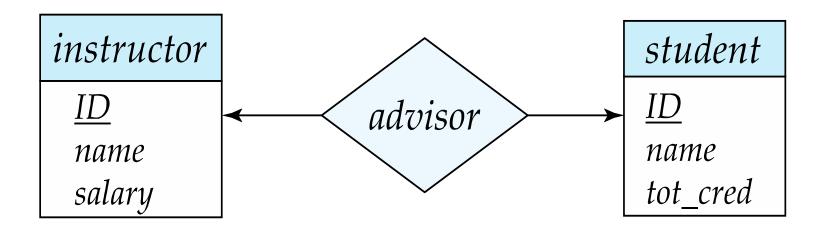
(22222, 456-7890) and (22222, 123-4567)

For a binary many-to-many relationship, the union of the primary-key attributes of all the participating entity sets becomes the primary key.



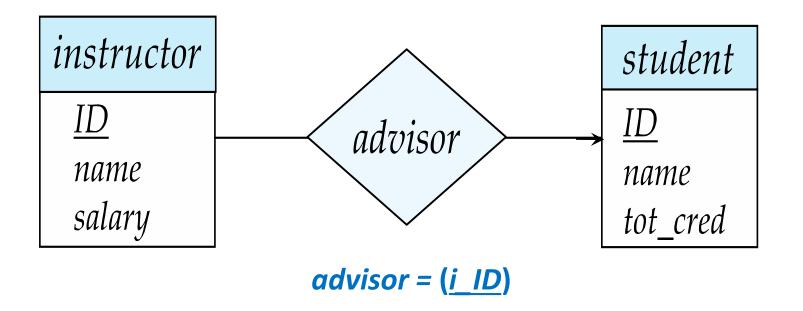
advisor = (i ID, s ID)

■ For a binary one-to-one relationship set, the primary key of either entity set can be chosen as the primary key. The choice can be made arbitrarily.



 $advisor = (\underline{s} \underline{ID})$

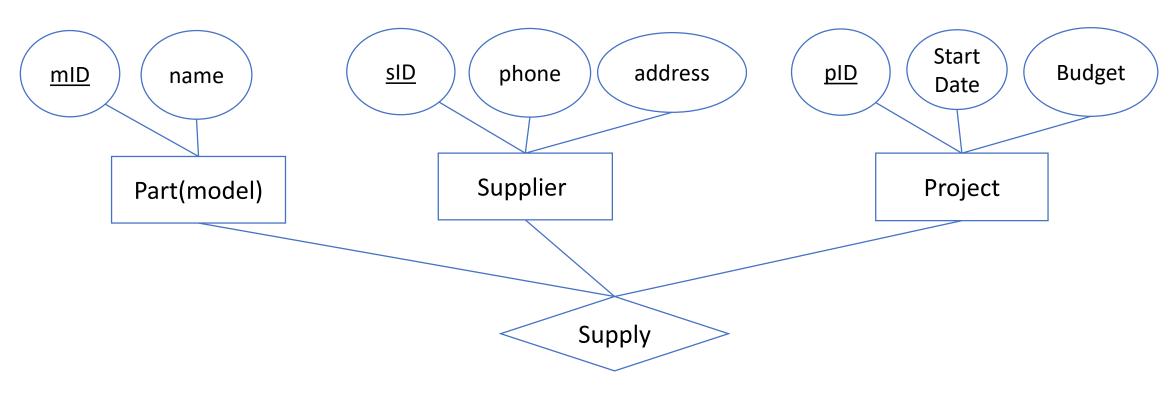
■ For a binary many-to-one or one-to-many relationship set, the primary key of the entity set on the "many" side of the relationship set serves as the primary key.



■ For an *n*-ary relationship set without any arrows on its edges, the union of the primary key-attributes of all the participating entity sets becomes the primary key.

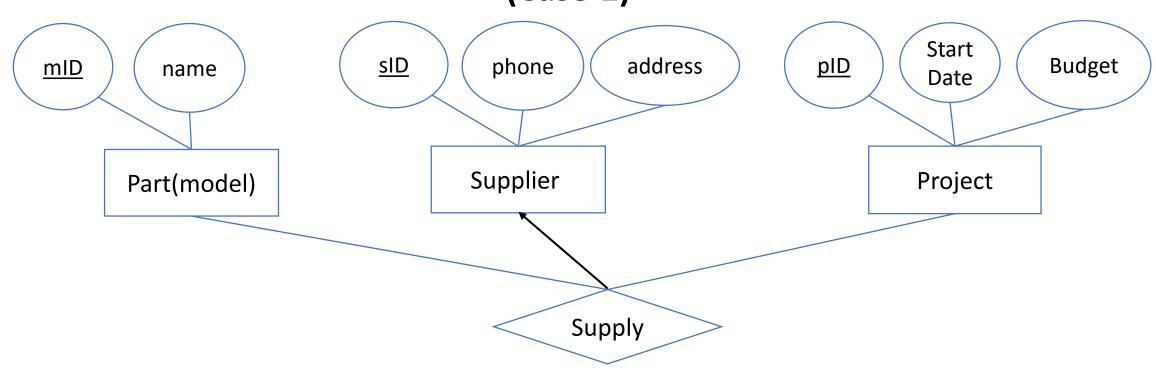
For an n-ary relationship set with an arrow on one of its edges, the primary keys of the entity sets not on the "arrow" side of the relationship set serve as the primary key for the schema.

Multi-degree relationship → schema (Case-1)



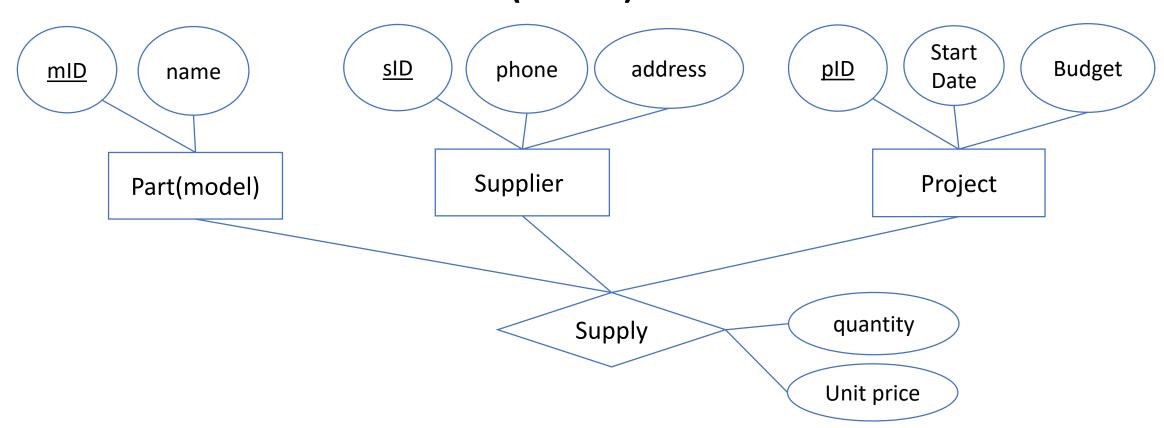
supply(mID, sID, pID)

Multi-degree relationship → schema (Case-2)



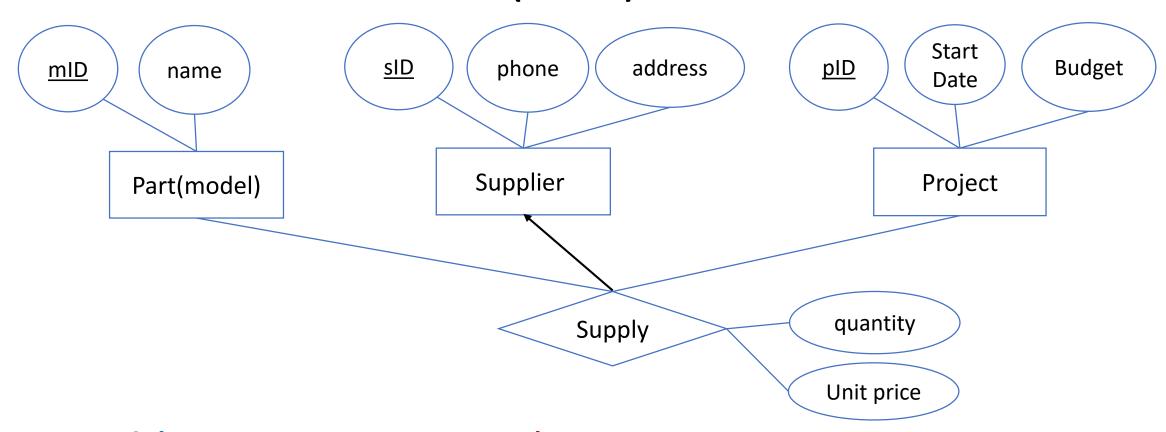
supply(mID, pID)

Multi-degree relationship → schema (Case-3)



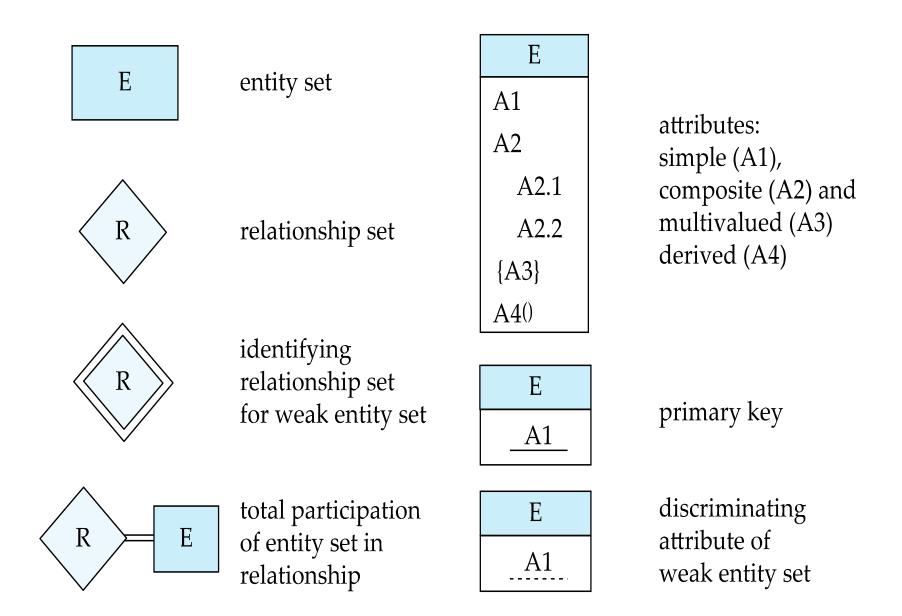
supply(mID, sID, pID, quantity, unitprice)

Multi-degree relationship → schema (Case-4)

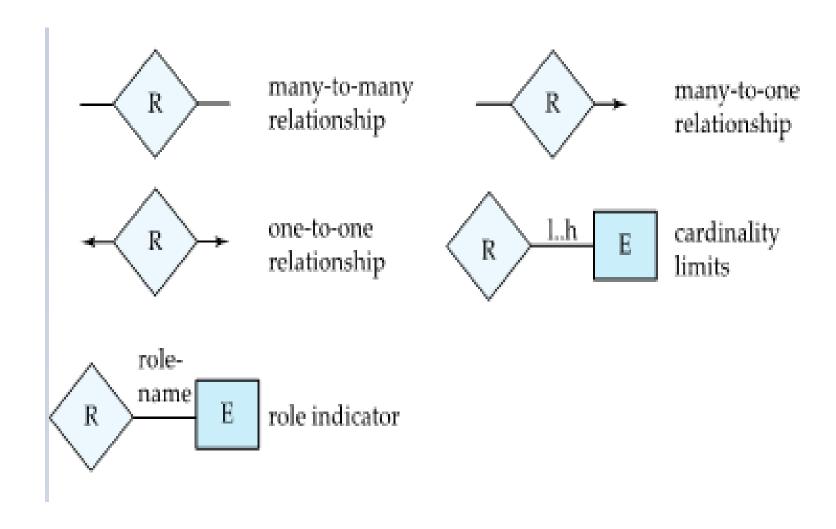


supply(mID, pID, quantity, unitprice)

Summary of Symbols Used in E-R Notation

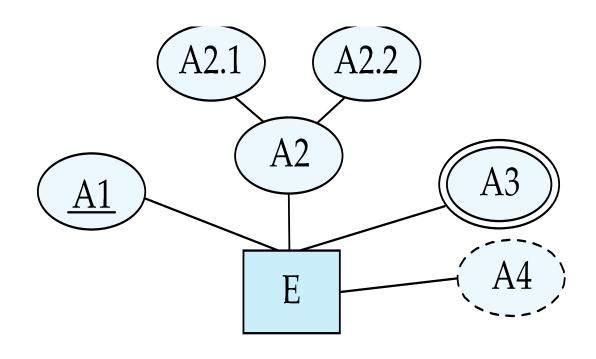


Summary of Symbols Used in E-R Notation

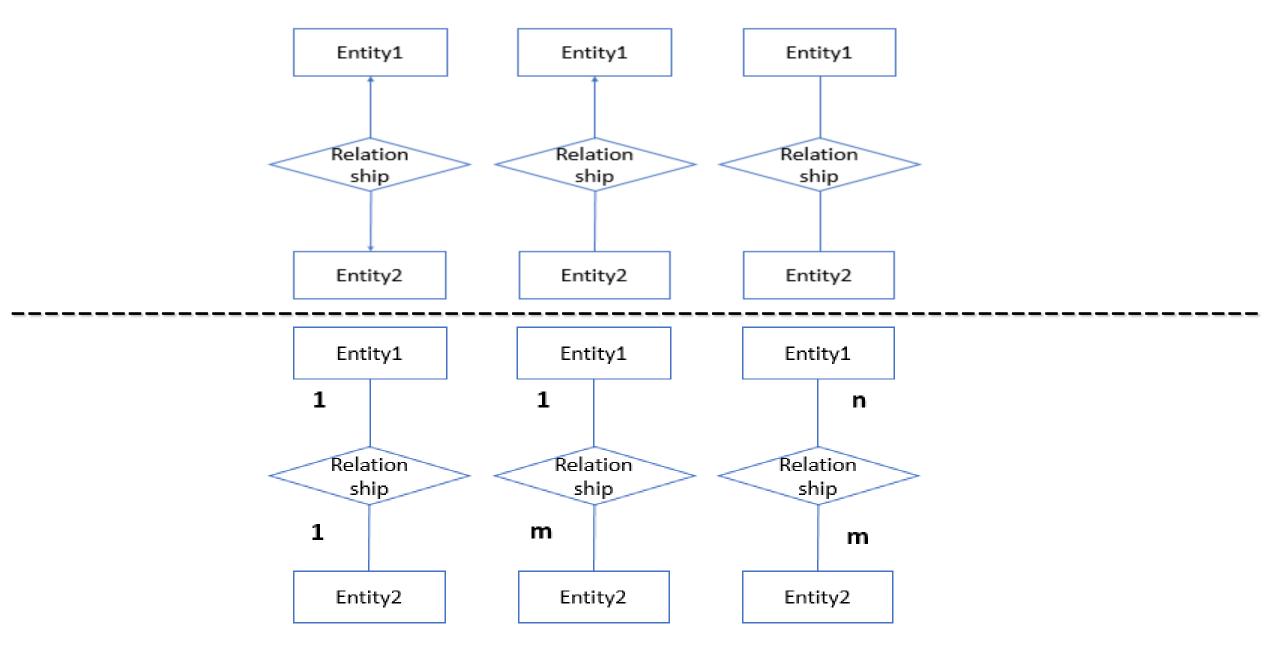


Alternative ER Notations

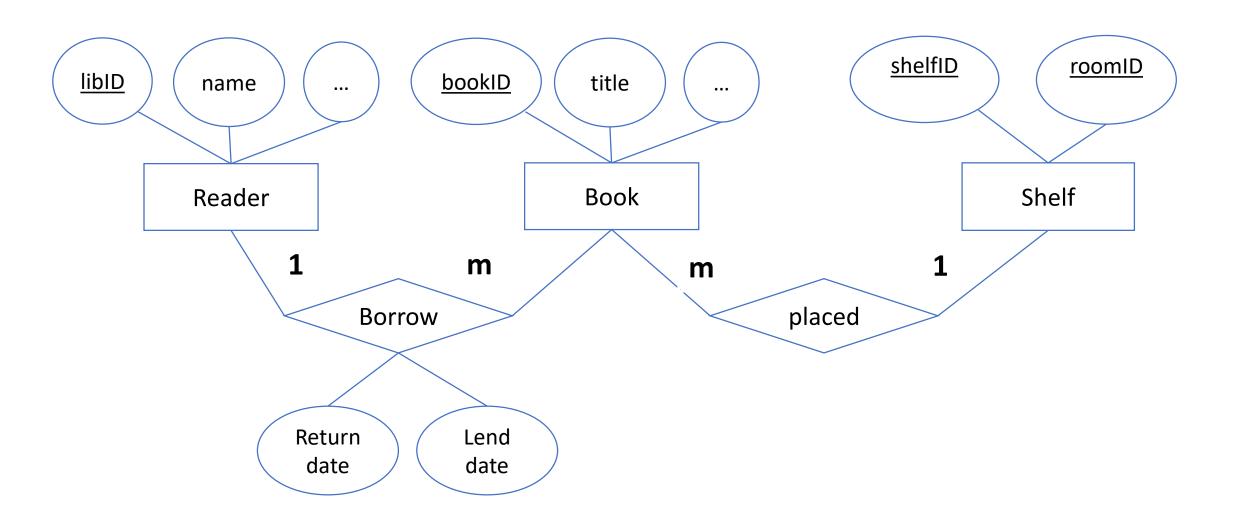
entity set E with simple attribute A1, composite attribute A2, multivalued attribute A3, derived attribute A4, and primary key A1



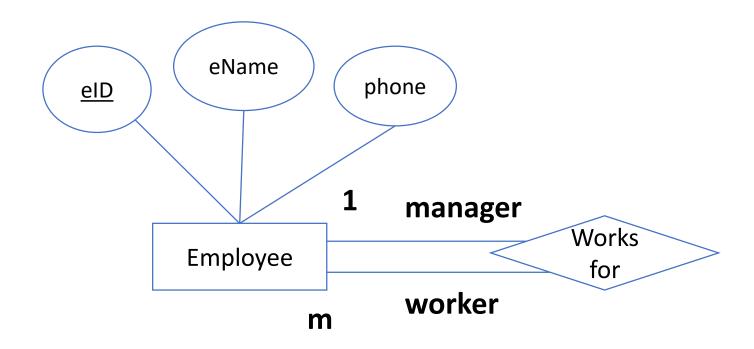
Alternative ER Notations



Example – cardinality



Example – role



Alternative ER Notations

Chen

many-to-many R R E2 E1 E2 E1 relationship one-to-one R R E1 E2 E1 E2 relationship many-to-one \mathbf{R} * R E1 E2 E2 E1 relationship participation R E2 R in R: total (E1) E1 **E**1 E2 and partial (E2)

IDE1FX (Crows feet notation)

Case study – Inventory management

- Parts
- Supplier
 - Supply parts
- Project
 - Use parts
- Warehouse
 - Store parts
- Warehouse keeper
 - Manage warehouse
 - One keeper manages multiple warehouses

Step 1 – determine Entity sets

Parts

Supplier

Project

Warehouse

• Warehouse keeper

Parts)

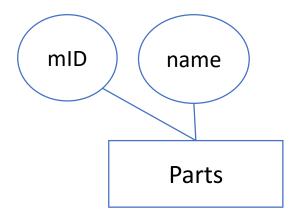
Supplier

Project

Warehouse

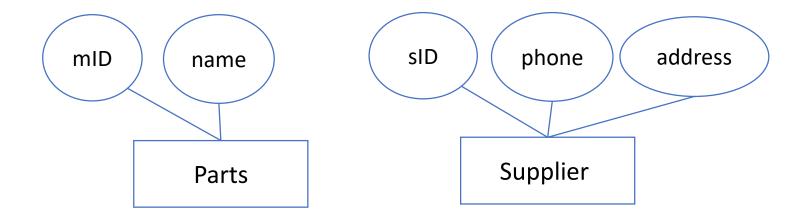
Keeper

Step 2 – determine Attributes



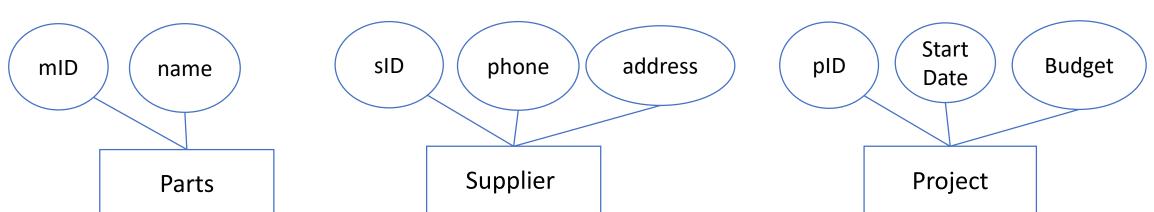
- Parts
- Supplier
- Project
- Warehouse
- Warehouse keeper

Step 2 – determine Attributes



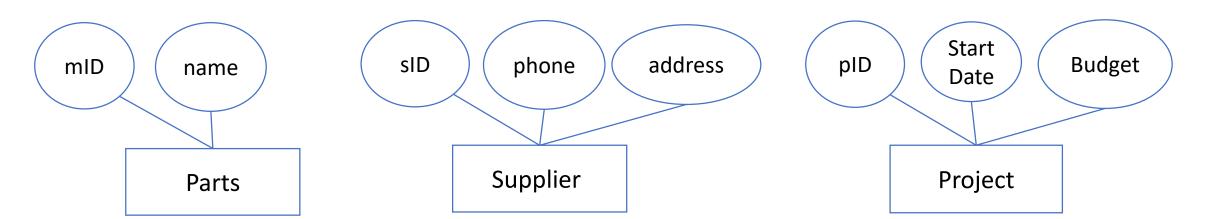
- Parts
- Supplier
- Project
- Warehouse
- Warehouse keeper

Step 2 – determine Attributes



- Parts
- Supplier
- Project
- Warehouse
- Warehouse keeper

Step 2 – determine Attributes



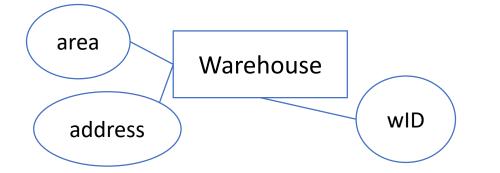
Parts

• Supplier

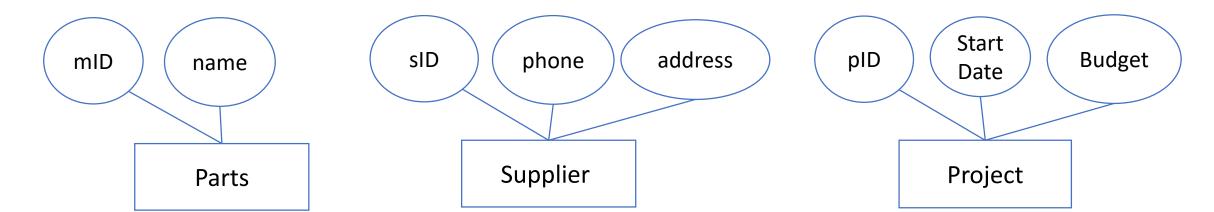
Project

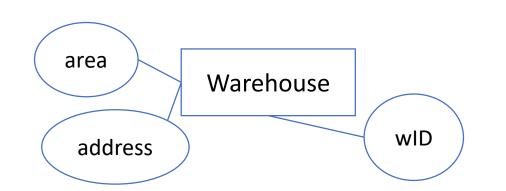
Warehouse

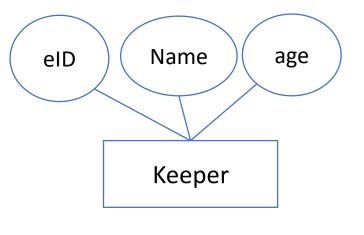
• Warehouse keeper



Step 2 – determine Attributes







Parts

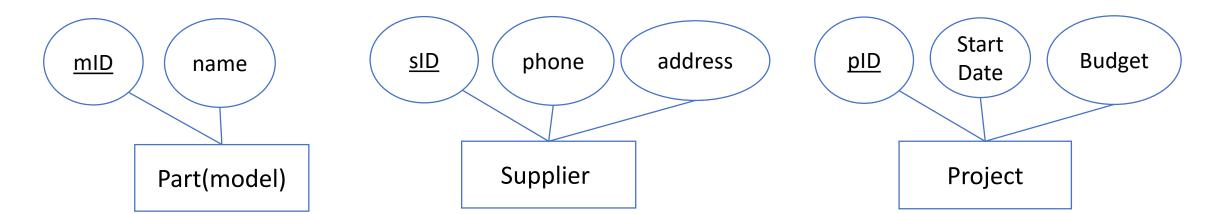
Supplier

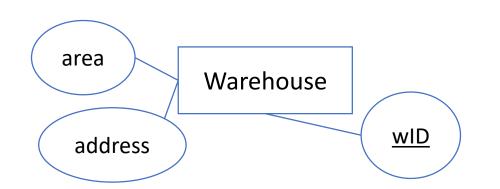
Project

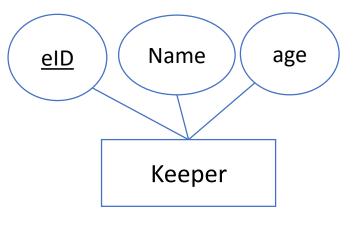
Warehouse

Warehouse keeper

Step 3 – determine keys







Parts

Supplier

Project

Warehouse

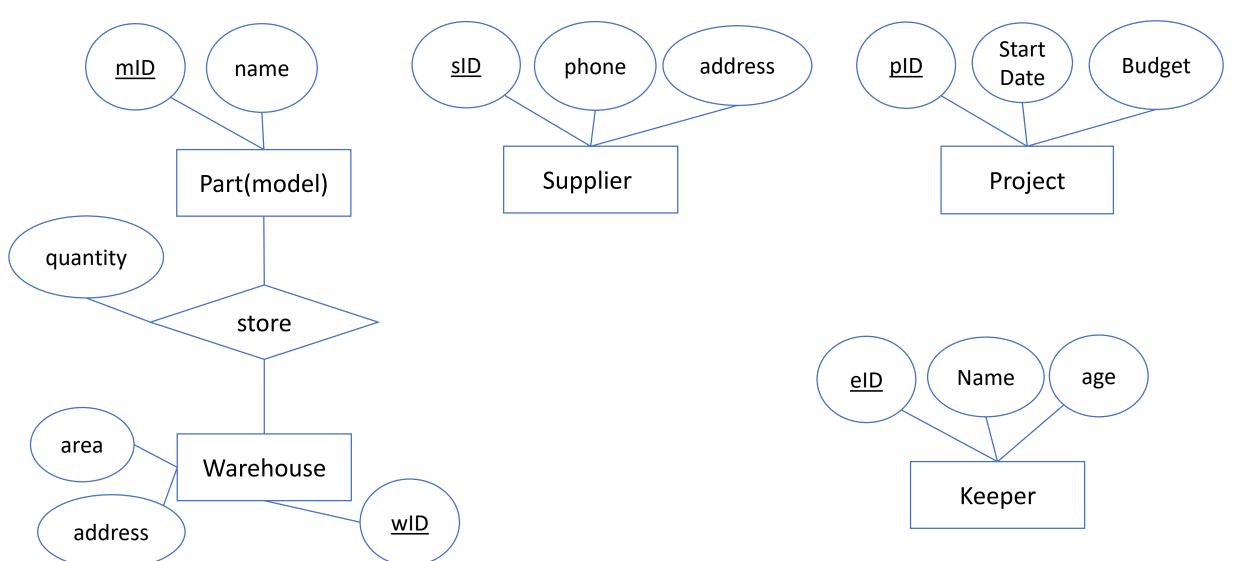
Warehouse keeper

Step 4 – establish relationships

Supplier

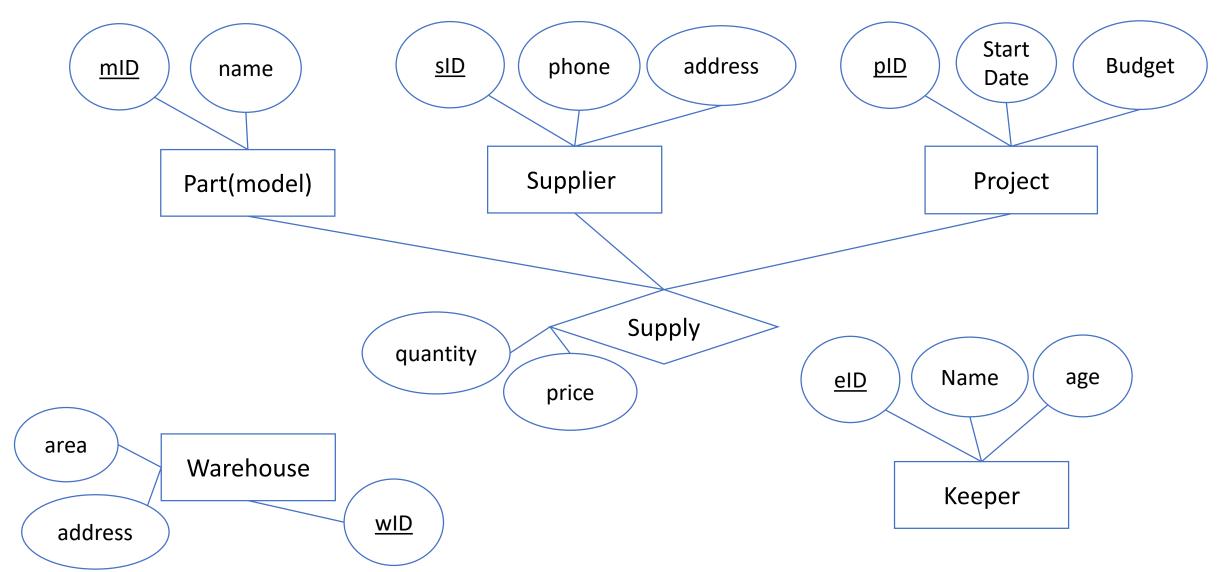
Parts

- SuppliedProject
- Warehouse
- Warehouse keeper

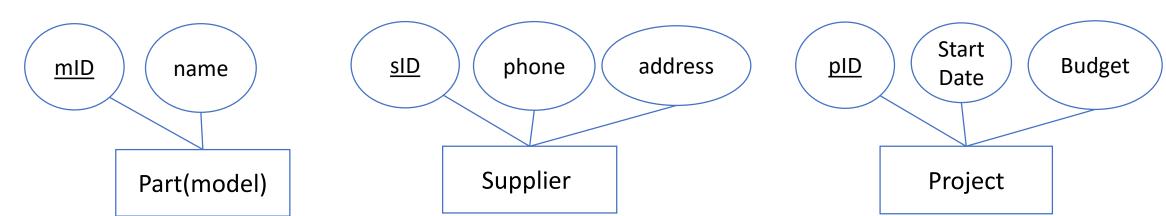




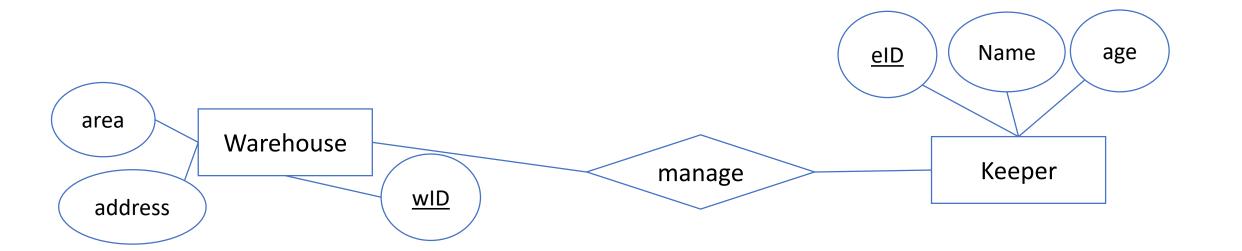
- Parts
- Supplier
- Project
- Warehouse
- Warehouse keeper



- Parts
- Supplier
- Project
- Warehouse
- Warehouse keeper

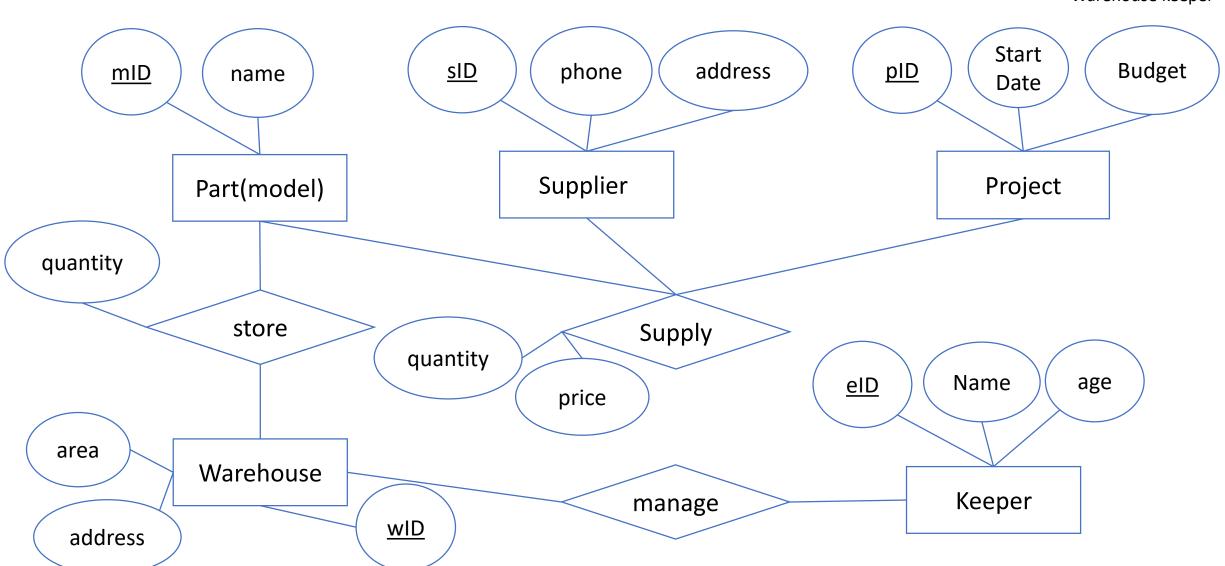


Step 4 – establish relationships





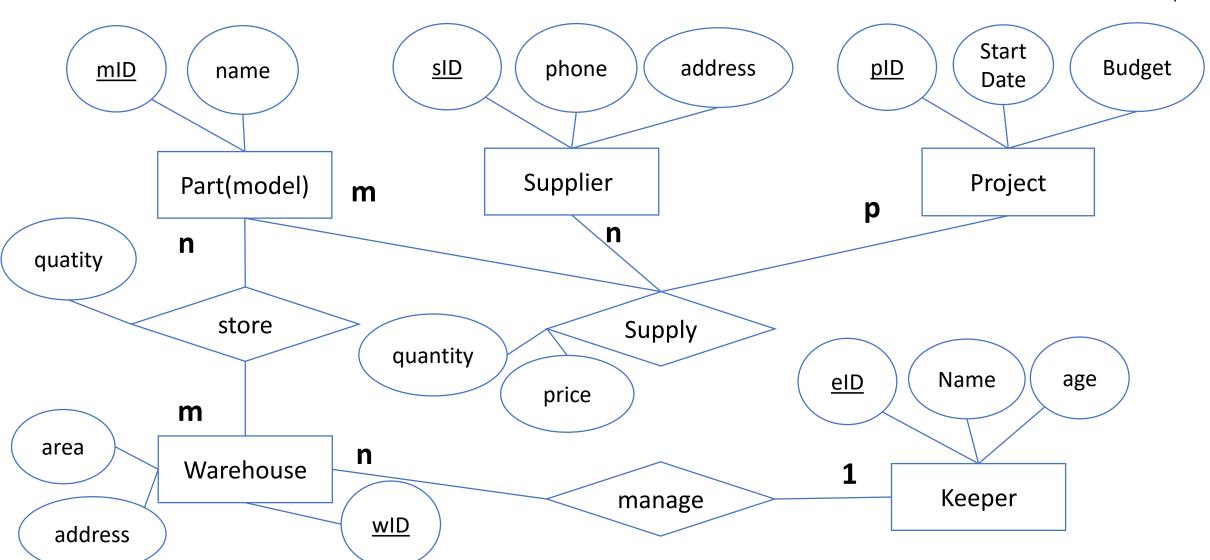
- Supplier
- Project
- Warehouse
- Warehouse keeper



Step 4 – establish relationships



- Supplier
- Project
- Warehouse
- Warehouse keeper



Step 5 – determine cardinality

Database for TV Companies (SELF STUDY)

- Store data about the TV series that the company produces
- Data include information about actors who play in the series, and directors who direct the episodes of the series
 - ☐ Actors and directors are employed by the company
 - ☐ A TV series are divided into episodes
 - ☐ Each episode may be transmitted at several occasions
 - ☐An actor is hired to participate in a series, but may participate in many series
 - ☐ Each episode of a series is directed by one of the directors, but different episodes may be directed by different directors.

Examples of database queries

- Which actors play in the series The Night Of?
- In which series does the actor Carey Mulligan participate?
- Which actors participate in more than one series?
- How many times has the first episode of the series Wild Lies been transmitted? At what times?
- How many directors are employed by the company?
- Which directors has directed the greatest number of episodes?

Entity sets

TV company

- TV company
- TV series
- Episode
- Employee
- Actor
- Director

Employee

TV series

Director

Episode

Actor

Attributes & Key

