Lecture 9 DB-Security

CS211 - Introduction to Database

DB Security

 To protect databases against compromises of their confidentiality, integrity and availability.

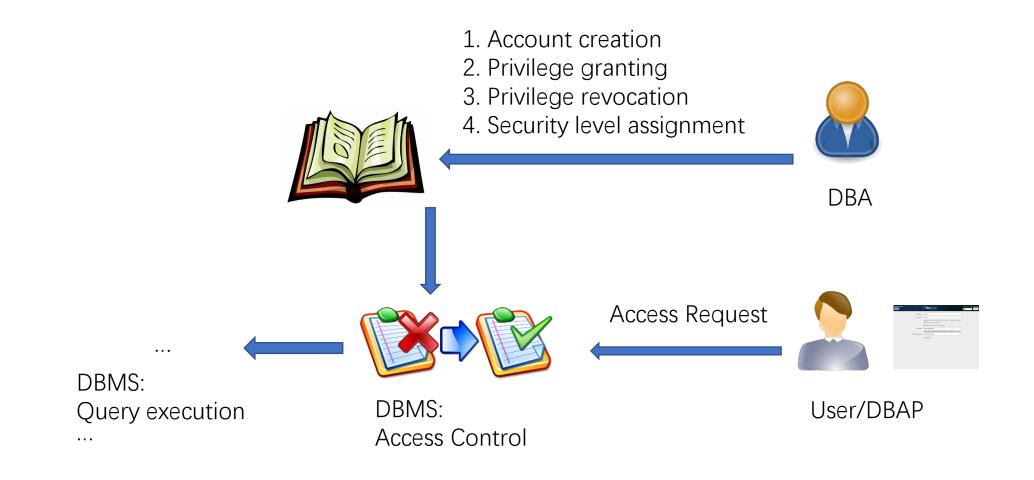
- DBMS security models
 - Authentication: a process by which you verify that someone is who they claim they are
 - Authorization: the process of establishing if the user (who is already authenticated), is permitted to have access to a resource. Authorization determines what a user is and is not allowed to do.
 - Access Control: the process of enforcing the required security for a particular resource.

Discretionary Access Control

granting and revoking privileges

- The DBA
 - granting privileges to users who need to use the system
 - classifying users and data in accordance with the policy of the organization

- Administrator account
 - 1. Account creation
 - 2. Privilege granting
 - 3. Privilege revocation
 - 4. Security level assignment
 - Action 1 is authentication, whereas 2 and 3 are discretionary authorization and 4 is used to control mandatory authorization



AccessRule

- (S O t P)
 - Subject: user/user group, account
 - Object: data (granularity: database/table/attribute/tuple/data item)
 - t: right (create/update/insert/delete/alter/query)
 - Predicate: condition
 - S own right t on O when P is satisfied
 - {AccessRule} stored in data dictionary or sys directory

Example – security requirement

employee(eID, eName, age, gender, salary, dID, manager)

User group

HR: Access all information

General office: read eName, dID

An employee: his/her own record, read only

Manager: query employee in the department

Senior Manager: query all information, read only

Example - implementation

S	О	t	P
HR	Employee	Read, del, insert, update	
General office	Employee(eName, dID)	read	
Senior manager	Employee	read	
Employee	Employee	read	eID=:userID
Department manager	Employee	read	Manager=:userID

Example – Matrix – (S O: t)

	01	O2	 On
S1	All	All	 All
S2	Read	-	 -
Sn	_	Read	 -

Given a matrix, and there is a query request: select-from-where

How do we validate the request?

List all required privileges and objects

UPDATE instructor SET workload=workload*1.1
WHERE iID in (select iID from course where cID='211')

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WHERE iID in (select iID from course where cID='211')

O	t
Instructor.workload	Update
Instructor.workload	Select
Instructor.iID	Select
Course.cID	Select
Course.iID	Select

List all privileges and objects

DELETE FROM student WHERE sID NOT IN (SELECT sID FROM rc)

List all privileges and objects

DELETE FROM student WHERE sID NOT IN (SELECT sID FROM rc)

O	t
Student	Del
Student.sid	Select/read
rc.sID	Select

Which privilege on table instructor is NOT needed

instructor(iID, iName, dID, workload)

UPDATE instructor SET workload = workload + 1 WHERE dID='03'

- 1. update(workload)
- 2. read(iName)
- 3. read(dID)
- 4. read(workload)

Example – View - (S O t P)

Data item access control

CREATE VIEW empV1 AS select * from employee

CREATE VIEW empV2 AS select eName, dID from employee (for general office)

Predicate

```
CREATE VIEW empV3 AS select * from employee where eID = :userID
```

CREATE VIEW empV4 AS select * from employee where manager = :userID

Consider table student(sID,sName,GPA),

Is it possible to grant a user modification privileges only on those students whose GPA is above average? (by creating a view)

- 1. Yes
- 2. No

Consider table student(sID,sName,GPA),

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1. Yes

2. No

```
CREATE VIEW stu_above_avg AS
select * from student where GPA >
( select avg(GPA) from student );
```

The view modification is not allowed

User levels

- User/Role levels: DBA → DBA developer → End User
- From right to left
- Level 1: select (read database/table/tuple/attribute…)
- Level 2: modify
 - Insert (tuple)
 - Update (data item)
 - Delete (tuple)
- Level 3: create
 - Create (database/table/index/view…)
 - Alter
 - Drop

- 1. Higher level has all privileges of lower level
- 2. Higher level has the right of granting

DBA – create a new user account

MariaDB/mySQL

CREATE USER 'test'@'localhost' IDENTIFIED BY 'myPassword';

CREATE USER 'test1'@'localhost' IDENTIFIED BY 'myPassword';

SQL - DCL

```
GRANT {privilege list}
ON {table/view}
TO {user/role list}
[with grant option]
```

- Privilege: select | insert | update | delete
- Only those privileges owned by the higher level user can be granted to lower level user
- With grant option: rights can be transited

GRANT all
ON proj1.*
TO test1@localhost

Example – security requirement

employee(eID, eName, age, gender, salary, dID, manager)

HR: Access all information

General office: read eName, dID

An employee: his/her own record, read only

Manager: query employee in the department

Senior Manager: query all information, read only

SQL DCL - Example

	eID
A Senior Manager	1001
A Department Manager	1021
A staff from HR	2001
A staff from General office	5001

CREATE VIEW empV1 AS select * from employee

CREATE VIEW empV2 AS select eName, dID from employee (for general office)

CREATE VIEW empV3 AS select * from employee where eID = :userID

CREATE VIEW empV4 AS select * from employee where manager = :userID

GRANT all ON employee TO 2001; GRANT select ON empV2 TO 5001; GRANT select ON empV3 TO public; GRANT select ON empV4 TO 1021;

SQL – DCL – revoke

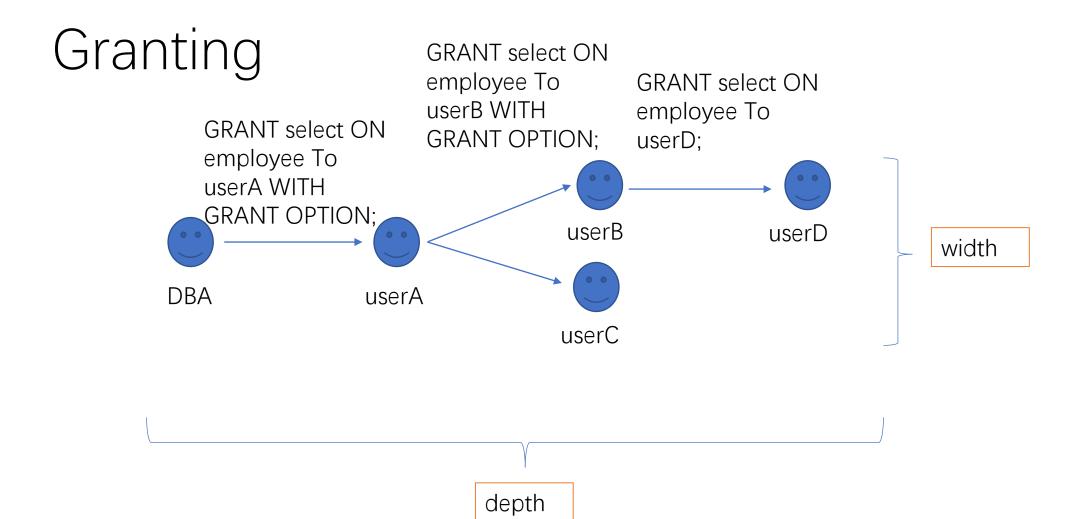
REVOKE {privilege list}
ON {table/view}
FROM {user/role list}
[cascade/restrict]

Cascade: Also revoke privileges granted from privileges being revoked (transitively), unless also granted from another source

Restrict: only revoke from directly related users(R)

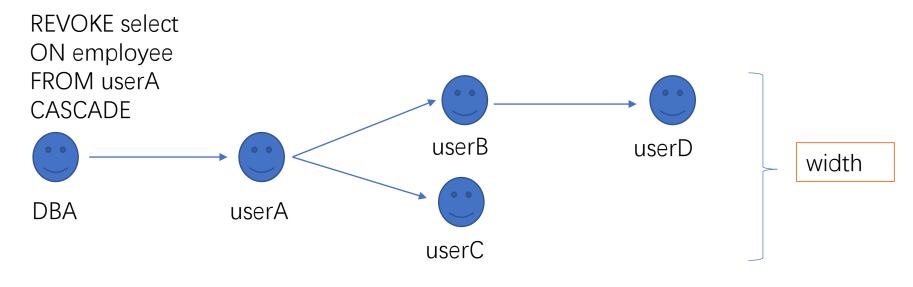
Both are not working in MariaDB/mySQL







Revoking



depth

Consider table R with owner A. Suppose we have the sequence:

- (1) A grants select(R) to B with grant option;
- (2) A grants select(R) to C with grant option;
- (3) B grants select(R) to D;
- (4) C grants select(R) to D;
- (5) A revokes select(R) from B with restrict.
- (6) D issue a select * from R; statement.

Does step (6) generate an error?

- Yes
- No

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Mandatory security

- Data object security level
 - Top secret, secret, confidential, unclassified
- User security level
 - Top secret, secret, confidential, unclassified

- User S is not permitted to read data object O unless Level(S)>=Level(O)
- User S is not permitted to write data object O when Level(S)>Level(O)
 - The level(O) is promoted as level(S) once O is written
 - If high level user write O, then low level user can not read it

Mandatory security – Implementation

Extend schema definition

R(A1:D1, A2:D2, ..., An:Dn)

R(A1:D1,C1, A2:D2,C2, ···, An:Dn,Cn, TC)

C1···Cn: security level for attributes

TC: security level for tuple