In this section, we study the arithmetic functions.

Definition: An <u>anthretic function</u> is a function defined over integers, i.e.  $f: N \longrightarrow C$ .

Example: 11, The trivial function I: N-) C.

1(n)=1 for all  $n \in N$ .

(2) The Enler's Phi funtion:  $\phi: \mathcal{N} \rightarrow \mathcal{I}$   $\phi(m) := \# \left\{ a : 1 \leq a \leq m, g(d(a,m) = 1) \right\}$ 

(3) The divisor function:  $d: N \rightarrow \mathbb{C}$  $d(m): = \# \{a: a|m\}.$ 

(4) The Möbius function:  $\mu: N \to \mathbb{C}$ .

 $\mu(m) = \begin{cases} (-1)^r & \text{if } m = P_1 P_2 - P_r \text{ with } P_i \text{ abstitut.} \\ 0 & \text{otherwise.} \end{cases}$ 

Definition: An arithmetic function:  $f: N \to C$  is <u>multiplicative</u> if f(mn) = f(m) f(n) when gcd(m,n) = 1.

An arithmetic function  $f: N \to \mathbb{C}$  is completely multiplicative if f(mn) = f(m) f(n) for all m, n.

Remark: f completely multiplicative => multiplicative.

In fant: 11, I(m) is completely multiplicative.

(2)  $\phi(m)$  is multiplicative  $\sim$  will show later.

but not completely multiplicative.

(outer) example:  $\phi(4) = 2$   $\phi(2) = 1$  $\phi(4) = \phi(2 \cdot 2) \neq \phi(2) \cdot \phi(2)$ 

13) d(m) is multiplicative

but not completely multiplicative.

(outer) example: d(4) = 3 d(2) = 2

 $d(4) = d(2\cdot 2) + d(2) \cdot d(2)$ 

4).  $\mu(m)$  is compliative

but not completely multiplicative.

(conten) example:  $4=2\cdot 2=2^2$ 

 $\mu(4)=0$   $\mu(2)=-1.$ 

 $\mu(4) = \mu(2\cdot 2) + \mu(2) \cdot \mu(2)$ .

Notations:

sum notation; Z

product notation:

example:  $\sum_{p} p$  means: find all primes drivides n and p sum them.

$$\sum_{\text{P|10}} = 2 + 5 = 7$$

 $\lceil \lceil (1-\frac{1}{p}) \rceil$  means: miltiply all  $(1-\frac{1}{p})$  where p/n.

why the multiplicative functions are important. m= Pi pi ~ Pr with Pi diotina. Let Let f be multiplicative. Then  $f(m) = f(P_1^{\alpha_1} P_2^{\alpha_2} - P_r^{\alpha_r})$  $= f(p_1^{\alpha_1}) f(p_2^{\alpha_2}) - \cdots f(p_r^{\alpha_r})$  $= \prod_{p^{\alpha}||m|} f(p^{\alpha})$ f is totally determined by its values at prime powers More over, if f is completely multiplicative.  $f(m) = f(p_1)^{\alpha_1} f(p_2)^{\alpha_2} - f(p_r)^{\alpha_r}$  $= \prod_{p \neq || m} f(p)^{d}$ 

f is totally determined by its values at primes.

First example: Enler's Phi funtion. Theorem (11.1 Enler's Phi funtion formla) (a) If p is a prime and k=1, then  $\phi(p^k) = p^k - p^{k-1} = p^k \left(1 - \frac{1}{p}\right)$ (b) If gcd(m,n)=1, then  $\phi(mn) = \phi(m)\phi(n).$ (c) For m= P1 P2 ... Pr  $\phi(m) = P_1^{\alpha_1} P_2^{\alpha_2} - P_r^{\alpha_r} \left(1 - \frac{1}{P_1}\right) - \left(1 - \frac{1}{P_r}\right)$  $= m \cdot \prod_{p \mid m} \left( 1 - \frac{1}{p} \right)$ Proof of (c): By (a), (b)  $\phi(m) = \phi(P_1^{d_1}) \phi(P_2^{d_2}) \cdots \phi(P_r^{d_r})$  $= (p_1^{\alpha_1} - p_1^{\alpha_1-1}) (p_2^{\alpha_2} - p_2^{\alpha_2-1}) \cdots (p_r^{\alpha_r} - p_r^{\alpha_{r-1}})$ 

$$= (p_{1}^{\alpha_{1}} - p_{1}^{\alpha_{1}-1}) (p_{2}^{\alpha_{2}} - p_{2}^{\alpha_{2}}) \cdots (p_{r}-p_{r})$$

$$= p_{1}^{\alpha_{1}} (1-\frac{1}{p_{1}}) p_{2}^{\alpha_{2}} (1-\frac{1}{p_{2}}) \cdots p_{r}^{\alpha_{r}} (1-\frac{1}{p_{r}})$$

$$= m \cdot \prod_{P \mid m} \left( 1 - \frac{1}{P} \right)$$

$$\phi(p^k) = \# \{ \alpha : 1 \leq \alpha \leq p^k, \ g(d(a, p^k) = 1) \}$$

We con show:

$$\{a: 1 \le a \le p^k, p|a\} = \{p, 2p, 3p, 4p, \dots (p^{k+1}-1)p, p^k\}$$

$$\Rightarrow \#\{a: |\leq a \leq p^k, p|a\} = p^{k-1}$$

This implies:

$$\phi(p^k) = p^k - p^{k-1}$$

Let g(d(m,n) = 1.

$$A = \left\{ \alpha : |\leq \alpha \leq mn, \gcd(\alpha, mn) = 1 \right\} \qquad \phi(mn) = \# A.$$

$$B = \begin{cases} b : | \leq b \leq m, \gcd(b, m) = 1 \end{cases}$$
  $\phi(m) = \# B.$ 

(We need to show: 
$$\phi(mn) = \phi(m) \phi(n)$$
 i.e.)  
 $\# A = \# B \cdot \# C$ 

We look at the following set:

$$M = \left\{ (b,c) : | \leq b \leq m, \gcd(b,m) = 1 \right\}$$

$$| \leq c \leq n, \gcd(c,n) = 1$$

We can show:  $\#B \cdot \#C = \#M$ Therefore, it suffices to show: #A = #M.

Strategy: we construct an one-to-one map from A to M.

Définition: Let f: A > B be a map.

- f is injective if  $f(b_1) = f(b_2) \Rightarrow b_1 = b_2$
- of is surjective if for any  $b \in B$ , we can find  $a \in A$  such that f(a) = b.
- · f is a <u>bijection</u> if f is both injective and sujective.

  one—to one map

Let A, B be finite sets. If there is an one-to-one map  $f: A \rightarrow B$ , then #A = #B. We construct the following manp: f: A -> M  $\begin{cases} a: & |\leq a \leq mn \\ g:d(a,mn)=1 \end{cases} \longrightarrow \begin{cases} (b,c): & |\leq b \leq m \ g:d(b,n)=1 \\ |\leq c \leq n \ g:d(c,n)=1 \end{cases}$  $\alpha \mapsto (a \pmod{m}, a \pmod{n}).$ We need to show f is both injective and sujective. • injective: Let  $a_1, a_2 \in A$  with  $f(a_1) = f(a_2)$ (ne need to show:  $a_1=a_2$ )  $\left( \alpha_{1} \left( \text{mod } m \right), \alpha_{1} \left( \text{mod } n \right) \right) = \left( \alpha_{2} \left( \text{mod } m \right), \alpha_{3} \left( \text{mod } n \right) \right)$  $\Rightarrow$   $\alpha_1 \equiv \alpha_2 \pmod{m}$  $\alpha_1 \equiv \alpha_2 \pmod{n}$ .

 $gcd(m,n)=1 \Rightarrow \alpha_1 \equiv \alpha_2 \pmod{mn}$ 

 $| \leq \alpha_1 \leq mn \quad | \leq \alpha_2 \leq mn \quad \Rightarrow \quad \alpha_1 = \alpha_2.$ • Surjective: (let  $(b,c) \in M$ , then we can find  $a \in A$ )

such that  $(a \pmod{m}, a \pmod{n}) = (b,c)$ . We look at the linear congruent quation:  $my = (C-b) \pmod{n}$  $gcd(m,n)=1 \Rightarrow we can find y, such that$  $my_1 \equiv (c-b) \pmod{n}$ . Set  $X = my_1 + b$ .  $\chi = b \pmod{m}$  $\chi = m y_1 + b \equiv (c - b + b) \pmod{n} \equiv C \pmod{n}$ a between I and mn such that  $\alpha = x \pmod{mn}$ .  $\alpha \equiv \chi \pmod{m} \equiv b \pmod{m}$  $0 \equiv \chi \pmod{n} \equiv c \pmod{n}$ Q.

Proof of (c): We showed: the map between  $A = \begin{cases} a: |\leq a \leq mn, & \gcd(a, mn) = 1 \end{cases}$   $M = \begin{cases} (b,c): |\leq b \leq m & \gcd(b, m) = 1 \\ |\leq c \leq n & \gcd(c, n) = 1 \end{cases}$ is one to one. Therefore # A = # M  $# A = \phi(mn)$  $\# M = \#B \cdot \#C = \phi(m) \cdot \phi(n)$  $\Rightarrow$   $\phi(mn) = \phi(m) \phi(n)$ . The "surjective" pourt can be generalized to the following theorem: Theorem (11.2 Chinese Remainder Theorem) Let m, n be integers with gcd (m,n)=1. Let b,c be integers Then the simultaneous congrnences  $X \equiv b \pmod{m}$   $X \equiv C \pmod{n}$ has exactly one solution with  $0 \le x \le mn$ .