

**PROGRAMS
FOR THE
COSMAC ELF
GRAPHICS**

PAUL C. MOEWS

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Introduction

The programs in this booklet were written for the basic COSMAC 1802 "Elf" with Pixie graphics as described in the July, 1977 issue of *Popular Electronics*. All programs can be easily modified to work in expanded memory systems by initializing a few registers with high order memory addresses (see Registers and Expanded Memory). The programs are documented to aid understanding, but knowledge of the 1802 instruction set is presumed, and assembly language mnemonics are not provided because of space limitations.

The first four programs introduce new graphic interrupt routines which use 128 bytes or less for display purposes. This is necessary because the basic "Elf" has only 256 bytes of memory and some has to be reserved for programming. While these routines either produce a display with fewer spots or a reduced area with the same resolution as that given in the Pixie examples, they avoid showing active code. A discussion of graphic interrupt routines is provided as an aid to understanding these new routines (see Graphic Interrupt Routines).

As the basic "Elf" has so little memory a programming "trick" is used in programs 5, 7 and 8. This trick allows one to have an instruction that is varied each time a set of code is obeyed by changing the instruction after each pass through the code. See the section entitled "Variable Code" for a discussion of the method used.

Registers and Expanded Memory

The 1802 requires that register 0 be used for the direct memory access (DMA) address, and register 1 for the interrupt address. Register 2 is used as an X register in all the programs and in some programs additional X registers are employed. Register 3 is used as the main program counter in all programs and subroutines are called by the "SEP register technique" (see RCA's User Manual for the 1802 Cosmac Microprocessor, pp. 54-58). All programs make the high order DMA address the same as the high order interrupt address. Basic "Elf" owners may run the programs just as they are written. In order to work in an expanded memory system R(1), R(2), and R(3) together with any other memory addressing registers must have their high order bytes preset. For example if the first program was to be run in page 4 (block 4) of an expanded system it would be loaded just as written but to page 4 instead of page 0. The following "patch" would be inserted on page 0:

| Add. | Code | Notes |
|-------|----------|---|
| 00 00 | F8 04 | load page number to D |
| 02 | B1 B2 B3 | initialize high order bytes |
| 05 | B4 B5 | of registers used by program 1 |
| 07 | C0 04 00 | transfer control to page 4 to run program |

This patch when executed would initialize the necessary high order bytes and enter the program on page 4. Each program includes a list of registers which must be set to work in expanded systems.

Graphic Interrupt Routines

In the "Elf" the generation of the interrupt and direct memory access (DMA) signals are all done under the control of the 1861 chip and its operation must be understood when writing or modifying interrupt routines. When enabled by an I/O signal from the microprocessor the 1861 chip generates interrupt requests at intervals of 1/61 seconds (3.579545/2 MHz clock). Twenty-nine machine cycles after each interrupt the first direct memory access (DMA) request occurs. Therefore when the 1861 generates an interrupt request we must use exactly 29 machine cycles to do the housekeeping required to get ready for the direct memory access requests. Things that must be done include saving the current designations of the X and P registers and saving the con-

tents of the D register so that we can return from the interrupt routine. We can also use these instructions to set register 0 to the address where we wish the DMA to start. If we put a counter in these preliminary instructions it will be incremented 61 times a second and can serve as a convenient timer or clock.

After these 29 preliminary machine cycles we get the first "burst" of eight DMA's. Memory bytes with the addresses designated by the contents of register 0,R(0) incremented once, R(0) incremented twice, . . . to R(0) incremented seven times, are mapped to the top line of the TV screen; R(0) is left pointing + 8 memory locations from where it started (i.e. R(0) has been incremented 8 times). We are now given 6 machine cycles, which can be used as we wish, these are normally used to reset register 0. After 6 machine cycles another burst of 8 DMA's occur, followed by 6 machine cycles, and so on until a total of 128 DMA bursts have filled the TV screen. During the last four DMA bursts the EF1 flag is held low as an aid to returning from the interrupt routine.

In the one page interrupt used by the Pixie graphics the starting address of register 0 is held constant for 4 DMA bursts giving a display in which each bit represents a square and we have the display divided into a field of 64 x 32 squares.

If we were to hold the starting address of register 0 constant for 8 DMA bursts one set bit would form a rectangle twice as high as it is wide, and we would form an array of such rectangles 64 across by 16 down. An interrupt routine which does just this is given in program 1. If we chose to set 2 adjacent bits at a time we would then have an array of squares 32 across and 16 down.

In the same manner holding the starting address of register 0 constant for 16 DMA bursts would form rectangles 4 times as high as they are wide and if we set 4 bits at a time we would produce an array of squares 16 across and 8 down. Program 2 contains such an interrupt routine. Holding register 0 constant for 32 DMA bursts and setting one word at a time would give an array of squares 8 across and 4 down as in Program 3. Note that each doubling of the number of DMA bursts halves the number of store locations required for display; 256 locations for 4 DMA bursts, 128 for 8 DMA bursts, 64 for 16 DMA bursts, and 32 for 32 DMA bursts.

It isn't necessary to keep the number of DMA bursts constant at each starting address. There is nothing to prevent changing the resolutions across

the TV screen if we so desire and we can save display locations in this way. Program 4 introduces such an interrupt routine, the top part of the screen is blanked by doing 108 DMA bursts on the same 8 locations, while the last 20 DMA bursts are used to produce an array of 64 x 5 squares with the resolution of a single page interrupt routine.

Variable Code

A simple way to set up code which can be varied is to save an instruction in a register and then modify the instruction by changing the contents of the register. The register contents must, of course, be transferred to memory before the instruction can be executed. The following example is taken from program 8 where it is used to load the low order part of seven registers with seven consecutive switch bytes.

| Add. | Code | Notes |
|------|-------------|--|
| 0B | F8 1A A4 | load address of variable code to R(4) |
| 0E | F8 AF A5 | AF is the instruction to be executed in location 1A, 1st time it will be AF, 2nd time AE, then AD, etc. and finally A9 |
| 11 | 85 44 | get the instruction and store it in location 1A |
| 13 | 3F 13 37 15 | wait for in on, off |
| 17 | 6C 64 22 | read and display switch byte |
| 1A | XX | the variable instruction, first AF, etc. |
| 1B | 85 FB A9 | check, are we done? |
| 1E | 32 23 | yes—go on |
| 20 | 25 30 11 | no-change instruction and go back through code |

A similar method is used to increment one of 4 registers in program 5 as follows:

| Add. | Code | Notes |
|------|------|---|
| 48 | 8E | R(E) contains the instruction to be executed, in program 5 it is one of 1A, 1B, 1C or 1D. The instruction is here loaded to the D register. |
| 49 | 53 | R(3) is the program counter and a 53 instruction writes 1A, 1C, or 1D to the next location (4A) |

4A XX

This location becomes 1A, 1B, 1C, or 1D depending upon the contents of R(E).0 and is executed.

Variable code has the drawback that it cannot be used in read only memory. However random access memory is available in the Elf and variable code can be used to get more work out of the limited memory available. Care must be taken to maintain control over the locations changed and the instructions allowed when using this method.

Program 1—Display Routine (128 Byte)

This program introduces an interrupt routine which displays 128 memory locations on the screen. Each bit is displayed as a rectangle twice as tall as it is wide; in order to get a square, two adjacent bits must be set. Such an interrupt routine could be written by simply extending RCA's single page interrupt, however a method is used in which a counter is incremented between DMA bursts. This method is more versatile and the number of similar DMA bursts can be made equal to $5 + 3N$ where N is any integer, in this case $n = 1$ to give 8 identical sets of DMA bursts. Program 3 employs a similar routine with $N = 9$ to produce 32 identical DMA bursts. Note that the routine must not go across a page boundary and that the locations displayed are actually 7F through FE. The memory locations are, however, addressed individually as 00 through 7F in hexadecimal.

The program has three modes depending on what is set in the switches on turning to run. If 01 is set, memory is examined, enter the address desired, (remember 00 to 7F) push the in button and the contents of that location will be displayed; successive pushes of the in button step through memory. If 02 is set memory locations are changed, enter the address desired, push the in button, then enter the new contents and push in again, successive pushes of the in button change successive memory locations. If CC is set the display area is cleared and the program automatically enters the input mode.

A picture which illustrates the use of this interrupt routine is listed at the end of the program. Each line of the listing contains 64 bits (8 words) and forms one of the lines of the TV display, there are 16 lines in all. If this picture is entered by means of the program it will be seen to form as memory is filled.

Registers 1, 2, 3, 4 and 5 must be set if this pro-

gram is to be run in an expanded memory system.

| Use of Registers | | | FF we've run over displayed locations | |
|------------------|-------------------------------|-------------------|---------------------------------------|--|
| | High | Low | 3A | 3A 3E |
| R(0) | | DMA address | 3C | 85 A4 |
| R(1) | | interrupt address | 3E | 04 |
| R(2) | | X register | 3F | 31 47 |
| R(3) | | program counter | 41 | 52 64 22 |
| R(4) | | memory pointer | 44 | 14 |
| R(5) | | memory pointer | 45 | 30 33 |
| R(C) | counter for interrupt routine | | 47 | 6C 64 22 |
| R(F) | work in interrupt | work in interrupt | 4A | 54 30 44 |
| | | | 4D | 72 70 |
| | | | 4F | C4 |
| | | | 50 | 22 78 22 52 |
| | | | 54 | 9C AF |
| | | | 56 | F8 7F BF |
| | | | 59 | 91 BO |
| | | | 5B | E2 9F AO |
| | | | -- | -- |
| | | | 5E | E2 9F AO |
| | | | -- | -- |
| | | | 61 | 2F AO 8F |
| | | | -- | -- |
| | | | 64 | 32 6A 9F AO to 6A if counted out, else restore DMA address |
| | | | -- | DMA 4 |
| | | | 68 | 30 5F |
| | | | 6A | 9F AO |
| | | | -- | -- |
| | | | 6C | A0 9C 34 4D restore DMA address, and reload counter, leave routine if done |
| | | | -- | DMA 8 (except last one) |
| | | | 70 | AF 80 BF |
| | | | -- | restore counter and set new DMA address |
| | | | 73 | 30 5F |
| | | | 74-7B | new DMA 1 |
| | | | 7C-7E | back to 5F for DMA 2, etc. |
| | | | 7F-FE | unused |
| | | | | stack |
| | | | | displayed locations |

| Picture | | | | | | | | Use of Registers | |
|---------|-------------------------|------|--|--|--|------|-----|------------------|--|
| Add. | Code | | | | | High | Low | | |
| 00 | 98 00 00 60 00 60 00 00 | R(0) | DMA address | | | | | | |
| 08 | 88 00 00 20 00 10 01 80 | R(1) | interrupt address | | | | | | |
| 10 | F8 00 07 27 F8 18 01 80 | R(2) | X register, stack | | | | | | |
| 18 | 00 00 1F FE DF 98 06 01 | R(3) | program counter | | | | | | |
| 20 | F8 00 07 FF FF FO 18 06 | R(4) | memory pointer | | | | | | |
| 28 | 08 00 00 74 FF F7 F9 98 | R(5) | memory pointer | | | | | | |
| 30 | F8 00 1F 80 3F 9F FF EO | R(9) | counter for timer in interrupt routine | | | | | | |
| 38 | 00 00 F5 FF CO 7E 1E 60 | R(C) | counter for interrupt routine | | | | | | |
| 40 | F8 00 1F 80 3F 9F FF EO | R(F) | work in interrupt | | | | | | |
| 48 | A8 00 00 7F FF F7 F9 98 | | work in interrupt | | | | | | |
| 50 | F8 00 07 FF FF FO 18 06 | | | | | | | | |
| 58 | A8 00 1F FE DF 98 06 01 | | | | | | | | |
| 60 | 00 00 07 27 F8 18 01 80 | | | | | | | | |
| 68 | 88 00 00 20 00 10 01 80 | | | | | | | | |
| 70 | 88 00 00 60 00 60 00 00 | | | | | | | | |
| 78 | F8 00 00 00 00 00 00 00 | | | | | | | | |

Program 2 – Display Routine (64 Byte)

This program contains an interrupt routine which displays 64 memory locations on the screen. Each location is displayed as two adjacent squares, the screen is divided into 128 squares (16 wide by 8 high). Two sets of 64 locations are displayed so there are two displays which appear alternately, each for about 0.5 seconds.

The interrupt routine employs a counter similar to that used in program 1 except that the number of identical DMA bursts is $7 + 3N$, where N is any integer. N is set to 3 to give 16 identical DMA cycles. Again the routine must not go across a page boundary and the locations displayed are actually 7F through FE. However in this program the individual squares are the things that are addressed and they are numbered 00 through FF as shown in the diagram. Set one is numbered 00 to 7F and set 2, 80 to FF, again in hexadecimal.

The program has only two modes of operation. If CC is set in the switch register when run is turned on the displayed locations are erased, otherwise they are unaltered. Once the program is running the numbers entered are square identifiers. Enter a square identifier and push in, that square will be changed to its opposite value, that is if it is off it will be turned on, if on it will be turned off. As a trial clear the display area and enter the 74 square identifiers listed after the program.

Registers 1, 2, 3, 4 and 5 must be set if this program is to be run in an expanded memory system.

| Program Listing | | |
|-----------------|-------------|--|
| Add. | Code | Notes |
| 00 | F8 54 A1 | interrupt address |
| 03 | F8 7E A2 | stack address |
| 05 | F8 09 A3 | program counter |
| 09 | D3 E2 | establish program counter and X register |
| OB | F8 04 BC | Counter for DMA bursts, see interrupt routine |
| OE | F8 7F A5 | starting address of displayed memory locations |
| 11 | 6C 64 22 | read and display switch byte |
| 14 | FD CC | is the switch byte CC? |
| 16 | 3A 22 | unless it is go to 22 |
| 18 | 85 A4 | if it is CC clear memory locations |
| 1A | F8 00 54 | |
| 1D | 14 84 | |
| 1F | 3A 1A 24 | |
| 22 | 69 | turn on TV |
| 23 | 3F 3E | wait for in on (part of waiting loop) |
| 25 | 37 25 | wait for in off if pushed |
| 27 | 6C 64 22 | read and display switch byte |
| 2A | F6 FC 7F | shift (i.e. divide by 2) and add 7F the starting address of displayed memory |
| 2D | A4 | save result in R(4),0 |
| 2E | F0 F6 | bring back switch byte, shift to 38 if switch byte was even, else go on |
| 30 | 3B 38 | |
| 32 | 04 FB OF 54 | here if switch byte odd, x'or memory against OF and replace |

| | | | | |
|----|-------------|--|-------|--|
| 36 | 30 23 | return to waiting loop | — | DMA 2, 5, 8, 11 and 14 here |
| 38 | 04 FB FO 54 | here if switch byte even, x'or memory against FO and re- place | 69 | 32 6F 9F AO to 6F if counted out, else restore DMA address |
| 3C | 30 23 | return to waiting loop | — | DMA 3, 6, 9 and 12 here |
| 3E | 89 FD 40 | waiting loop, examine timer | 6D | 30 64 and go back to 64 |
| 41 | 33 23 | continue wait unless time up | 6F | 9F AO here if counted out, restore DMA address |
| 43 | F8 00 A9 | here if time up, reset timer | — | DMA 15 |
| 46 | 85 | get starting address of cur- rently displayed locations | 71 | AO 9C 34 52 restore DMA address, and re- load counter, leave routine if done |
| 47 | FB BF | is it set 2? | — | DMA 16 (except last one) |
| 49 | 32 4D | go to set 1 if set 2 is now being displayed | 75 | AF 80 BF restore counter and set new DMA address |
| 4B | FF 80 | else set to display set 2 | — | new DMA 1 |
| 4D | FC 7F A5 | store correct starting address | — | back to 67 for new DMA 2, etc. |
| 50 | 30 23 | back to waiting loop | 78 | 30 67 unused |
| 52 | 72 70 | return from interrupt | 7A-7B | stack |
| 54 | C4 | entry to interrupt | 7C-7E | displayed locations |
| 55 | 22 78 22 52 | save X, P and D | 7F-FE | |
| 59 | 9C AF | load counter for DMA bursts to R(F).0 | | |
| 5B | 85 BF | load starting address of dis- played locations to R(F).1 | | |
| 5D | 91 BO | load high order DMA address | | |
| 5F | 9F AO 19 | load low order DMA address, increment timer | | |
| — | | DMA 1 occurs here (1st burst only) | | |
| 62 | 30 67 9F AO | skip to 67, reenters at 64 for 2nd and subsequent bursts | | |
| — | | DMA 4, 7, 10 and 13 here | | |
| 66 | 2F AO 8F | decrement counter, restore DMA address, load counter | | |

Picture

(Note that bytes given are square identifiers.)

| | | | | | | | |
|----|----|----|----|----|----|----|----|
| 00 | 03 | 04 | 08 | 09 | OD | OE | 10 |
| 12 | 15 | 17 | 1A | 1C | 1F | 20 | 22 |
| 25 | 27 | 2A | 2F | 30 | 33 | 34 | 37 |
| 3A | 3D | 3E | 40 | 42 | 45 | 47 | 4A |
| 4C | 50 | 52 | 55 | 57 | 5A | 5C | 60 |
| 63 | 64 | 68 | 69 | 6C | 6D | 6E | 6F |
| 80 | 81 | 82 | 83 | 84 | 86 | 8B | 8C |
| 8D | 8E | 8F | 90 | 96 | 9B | AO | A6 |
| AB | B0 | B1 | B2 | B3 | B6 | BB | BC |
| BD | BE | CO | C6 | CB | DO | D6 | DB |
| E0 | E1 | E2 | E3 | E4 | E6 | E7 | E8 |
| E9 | EB | | | | | | |

Diagram — Program 2

Set No. 1

| | | | | | | | | | | | | | | | |
|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|
| 00 | 01 | 02 | 03 | 04 | 05 | 06 | 07 | 08 | 09 | 0A | 0B | 0C | 0D | 0E | 0F |
| 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 1A | 1B | 1C | 1D | 1E | 1F |
| 20 | 21 | 22 | 23 | 24 | 25 | 26 | 27 | 28 | 29 | 2A | 2B | 2C | 2D | 2E | 2F |
| 30 | 31 | 32 | 33 | 34 | 35 | 36 | 37 | 38 | 39 | 3A | 3B | 3C | 3D | 3E | 3F |
| 40 | 41 | 42 | 43 | 44 | 45 | 46 | 47 | 48 | 49 | 4A | 4B | 4C | 4D | 4E | 4F |
| 50 | 51 | 52 | 53 | 54 | 55 | 56 | 57 | 58 | 59 | 5A | 5B | 5C | 5D | 5E | 5F |
| 60 | 61 | 62 | 63 | 64 | 65 | 66 | 67 | 68 | 69 | 6A | 6B | 6C | 6D | 6E | 6F |
| 70 | 71 | 72 | 73 | 74 | 75 | 76 | 77 | 78 | 79 | 7A | 7B | 7C | 7D | 7E | 7F |

| | | | | | | | | | | | | | | | |
|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|
| 80 | 81 | 82 | 83 | 84 | 85 | 86 | 87 | 88 | 89 | 8A | 8B | 8C | 8D | 8E | 8F |
| 90 | 91 | 92 | 93 | 94 | 95 | 96 | 97 | 98 | 99 | 9A | 9B | 9C | 9D | 9E | 9F |
| A0 | A1 | A2 | A3 | A4 | A5 | A6 | A7 | A8 | A9 | AA | AB | AC | AD | AE | AF |
| B0 | B1 | B2 | B3 | B4 | B5 | B6 | B7 | B8 | B9 | BA | BB | BC | BD | BE | BF |
| C0 | C1 | C2 | C3 | C4 | C5 | C6 | C7 | C8 | C9 | CA | CB | CC | CD | CE | CF |
| D0 | D1 | D2 | D3 | D4 | D5 | D6 | D7 | D8 | D9 | DA | DB | DC | DD | DE | DF |
| E0 | E1 | E2 | E3 | E4 | E5 | E6 | E7 | E8 | E9 | EA | EB | EC | ED | EE | EF |
| F0 | F1 | F2 | F3 | F4 | F5 | F6 | F7 | F8 | F9 | FA | FB | FC | FD | FE | FF |

Set No. 2

The points in the above diagram represent squares arranged in arrays 16 wide by 8 high. An easy way to create and encode pictures is to use graph paper.

Program 3 – Display Routine (32 Byte)

The third interrupt routine displays 32 store locations on the screen. Each location is displayed as a single square and the screen is divided into 32 squares (8 wide by 4 high). Four sets of 32 locations are displayed so there are four displays which appear one after the other for about 0.5 seconds each.

This program uses an interrupt routine like that of program 1 except that 32 identical DMA bursts are produced. Again the routine must not go across a page boundary and the locations displayed are actually 7F through FE. The individual squares are the things that are addressed and they are numbered 00 through 7F as shown in the diagram.

The program has the same modes of operation as program 2. If CC is set in the switch register when run is turned on the displayed locations are erased, otherwise they are unaltered. Once the program is running the numbers entered are square identifiers. Enter a square identifier and push in, that square will be changed to its opposite value, that is if it is off it will be turned on, if on it will be turned off. As a trial the program should be entered so as to clear the display and the 78 square identifiers given at the end of the program entered.

Registers 1, 2, 3, 4 and 5 must be set if this program is to be run in an expanded memory system.

Use of Registers

| | High | Low |
|------|-------------------------------|--|
| R(0) | | DMA address |
| R(1) | | interrupt address |
| R(2) | | X register, stack |
| R(3) | | program counter |
| R(4) | | memory pointer |
| R(5) | | memory pointer |
| R(9) | | counter for timer in interrupt routine |
| R(C) | counter for interrupt routine | |
| R(F) | work in interrupt | work in interrupt |

Program Listing

| Add. | Code | Notes |
|------|----------|--|
| 00 | F8 4F A1 | interrupt address |
| 03 | F8 7E A2 | stack address |
| 06 | F8 09 A3 | program counter |
| 09 | D3 E2 | establish program counter and X register |

| | | |
|----|-------------|---|
| OB | F8 OA BC | counter for DMA bursts, see interrupt routine |
| OE | F8 7F A5 | starting address of displayed memory locations |
| 11 | 6C 64 22 | read and display switch byte is it CC |
| 14 | FD CC | no-go to 22 |
| 16 | 3A 22 | if switch byte is CC clear display locations |
| 18 | 85 A4 | |
| 1A | F8 00 54 | |
| 1D | 14 84 | |
| 1F | 3A 1A 24 | |
| 22 | 69 | turn on TV |
| 23 | 3F 36 | wait for in on (part of waiting loop) |
| 25 | 37 25 | wait for in off if pushed |
| 27 | 6C FA 7F | read switch byte, and against 7F to get legal square identifier |
| 2A | 52 64 22 | display square to be changed |
| 2D | FC 7F A4 | add starting address of locations |
| 30 | 04 FB FF 54 | pick up location, x'or against FF, put it back |
| 34 | 30 23 | return to waiting loop |
| 36 | 89 FD 20 | waiting loop, examine timer |
| 39 | 33 23 | continue wait unless time up |
| 3B | F8 00 A9 | here if time up, reset timer |
| 3E | 85 | get starting address of currently displayed locations |
| 3F | FB DF | are we on the last set? |
| 41 | 32 48 | if we are go to 48 to start over |
| 43 | 85 FC 20 | if not add 20 to starting address of displayed bytes |
| 46 | 30 4A | and return to in loop |
| 48 | F8 7F | here if starting over |
| 4A | A5 30 23 | restore R(5) and back to waiting loop |
| 4D | 72 70 | return from interrupt |
| 4F | C4 | entry to interrupt |
| 50 | 22 78 22 52 | save X, P and D |
| 54 | 9C AF | load counter for DMA bursts to R(F).0 |
| 56 | 19 | increment timer |
| 57 | 85 BF | load starting address of displayed locations |
| 59 | 91 BO | load high order DMA address |
| 5B | 9F AO | load low order DMA address |
| -- | -- | DMA 1 here (1st burst only) |

| | | | | | | | | | | |
|----|-------------|--|-------|----------|---|----|----|----|----|----|
| 5D | E2 9F AO | restore DMA address | 6F | AF 80 BF | restore counter and set new DMA address | | | | | |
| -- | | DMA 2, 5, 8, 11, 14, 17, 20, 23, 26 and 29 here | -- | | new DMA 1 here | | | | | |
| 60 | 2F AO 8F | decrement counter, restore DMA address, load counter | 72 | 30 5E | back to 5E for DMA 2, etc. | | | | | |
| -- | | DMA 3, 6, 9, 12, 15, 18, 21, 24, 27 and 30 here | 74-7B | | unused | | | | | |
| 63 | 32 69 9F AO | to 69 if counted out, else restore DMA address | 7C-7E | | stack | | | | | |
| -- | | DMA 4, 7, 10, 13, 16, 19, 22, 25 and 28 here | 7F-FE | | displayed locations | | | | | |
| 67 | 30 5E | and go back to 5E | | | Picture | | | | | |
| 69 | 9F AO | here if counted out, restore DMA address | | | (Note that bytes given are square identifiers.) | | | | | |
| -- | | DMA 31 here | 01 | 02 | 03 | 04 | 05 | 08 | 09 | OB |
| 6B | AO 9C 34 4D | restore DMA address, and reload counter, leave routine if done | OD | OE | 11 | 12 | 14 | 15 | 19 | 1A |
| -- | | DMA 32 here | 1B | 1C | 1D | 21 | 22 | 23 | 24 | 25 |
| | | | 28 | 29 | 2A | 2B | 2D | 2E | 31 | 32 |
| | | | 34 | 35 | 39 | 3A | 3B | 3C | 3D | 41 |
| | | | 42 | 43 | 44 | 45 | 48 | 49 | 4B | 4D |
| | | | 4E | 51 | 52 | 54 | 55 | 59 | 5A | 5B |
| | | | 5C | 5D | 61 | 62 | 63 | 64 | 65 | 68 |
| | | | 69 | 6B | 6C | 6D | 6E | 71 | 72 | 74 |
| | | | 75 | 79 | 7A | 7B | 7C | 7D | | |

Diagram – Program 3

Set No. 1

| | | | | | | | |
|----|----|----|----|----|----|----|----|
| 00 | 01 | 02 | 03 | 04 | 05 | 06 | 07 |
| 08 | 09 | 0A | 0B | 0C | 0D | 0E | 0F |
| 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 |
| 18 | 19 | 1A | 1B | 1C | 1D | 1E | 1F |

| | | | | | | | |
|----|----|----|----|----|----|----|----|
| 40 | 41 | 42 | 43 | 44 | 45 | 46 | 47 |
| 48 | 49 | 4A | 4B | 4C | 4D | 4E | 4F |
| 50 | 51 | 52 | 53 | 54 | 55 | 56 | 57 |
| 58 | 59 | 5A | 5B | 5C | 5D | 5E | 5F |

Set No. 3

Set No. 2

| | | | | | | | |
|----|----|----|----|----|----|----|----|
| 20 | 21 | 22 | 23 | 24 | 25 | 26 | 27 |
| 28 | 29 | 2A | 2B | 2C | 2D | 2E | 2F |
| 30 | 31 | 32 | 33 | 34 | 35 | 36 | 37 |
| 38 | 39 | 3A | 3B | 3C | 3D | 3E | 3F |

| | | | | | | | |
|----|----|----|----|----|----|----|----|
| 60 | 61 | 62 | 63 | 64 | 65 | 66 | 67 |
| 68 | 69 | 6A | 6B | 6C | 6D | 6E | 6F |
| 70 | 71 | 72 | 73 | 74 | 75 | 76 | 77 |
| 78 | 79 | 7A | 7B | 7C | 7D | 7E | 7F |

Set No. 4

The points in the above diagram represent squares arranged in arrays 8 wide by 4 high. An easy way to create and encode pictures is to use graph paper.

Program 4 — Variable Resolution Display Routine (48 Bytes)

Program 4 introduces a more complex interrupt routine. This interrupt routine displays 48 memory locations with a variable number of DMA bursts. The first 8 locations are used in producing 108 identical DMA bursts which fill most of the television screen. The last 40 store locations are used in 5 groups of 8, each group of which is used to produce 4 DMA bursts. In these last 40 locations, therefore, each bit produces a square of the same size that it would produce if a one page interrupt routine were used. That is the last 40 locations produces a display 64 squares wide and 5 squares high. The first 8 locations produce 64 bars 27 times as high as they are wide. If the first 8 locations are set to 00 the top 27/32 of the screen will be blank while the bottom 5/32 becomes the display area.

The program that uses this interrupt routine produces a "tickertape" display. The material to be displayed is coded in 11 sets of 5 bytes as shown in the diagram. After loading the program switch to run and enter bytes 5 at a time through the switches. When 5 bytes have been entered they are shifted into the tickertape display area and when all 11 sets have been entered the display area moves continuously.

Registers 1, 2, 3, 4, 5, and 6 must be set if this program is to be run in an expanded memory system.

Use of Registers

| High | Low |
|------|--|
| R(0) | DMA address |
| R(1) | interrupt address |
| R(2) | X register, stack |
| R(3) | program counter |
| R(4) | points to locations where data to be moved to display area is stored |
| R(5) | memory pointer to display area |
| R(6) | holds address of display subroutine |
| R(7) | counter, 8 bits/word |
| R(8) | counter, 5 words vertically |
| R(9) | counter for timer in interrupt routine |
| R(A) | counter, 8 words/line |
| R(B) | another 5 counter |

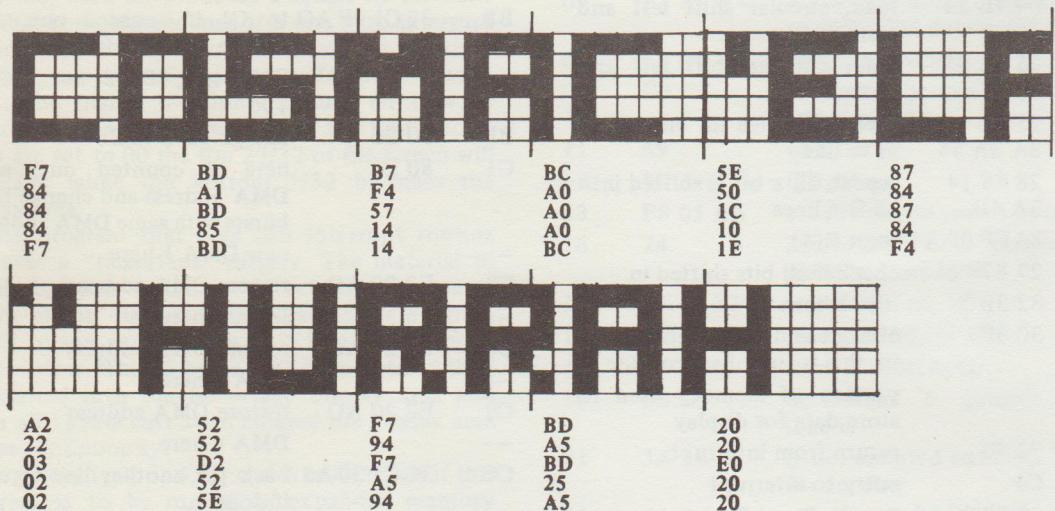
R(F) work in interrupt
(stores counter)

Program Listing

| Add. | Code | Notes |
|------|-------------|---|
| 00 | F8 A5 A1 | interrupt address |
| 03 | F8 3A A2 | X register, stack |
| 06 | F8 3C A6 | address of subroutine that shifts 5 bytes into display area |
| 09 | F8 A3 A4 | A3 is one more than address to write first data point to program counter |
| OC | F8 0F A3 | establish program counter |
| OF | D3 E2 | and X register |
| 11 | 69 | turn on TV |
| 12 | 7B | Q light on |
| 13 | F8 05 AB | set counter to read 5 bytes |
| 16 | 24 | decrement R(4)—points to place to store byte |
| 17 | 3F 17 37 19 | wait for in on, off |
| 1B | 6C 64 22 | read and display byte |
| 1E | 54 | store it in M(R(4)) |
| 1F | 2B 8B | decrement 5 counter and load it |
| 21 | 3A 16 | if not counted out back to 16 for another |
| 23 | 7A D6 | else turn Q off and call routine to shift these 5 bytes to display |
| 25 | 84 FB 6C | load R(4).0 if it equals 6C all data has been entered |
| 28 | 3A 12 | else go back for 5 more bytes |
| 2A | F8 9E A4 | here if all data loaded, restore R(4) address for display |
| 2D | D6 | this section of code repeatedly calls display subroutine to shift the message into the display area |
| 2E | 84 FF 05 | |
| 31 | A4 | |
| 32 | FB 67 3A 2D | |
| 36 | 30 2A | |
| 38 | xx xx xx | stack |
| 3B | D3 | return from subroutine, this subroutine shifts the bytes in M(R(4)), M(R(4)+1), M(R(4)+2), M(R(4)+3) and M(R(4)+4) into the display area - R(4) and the memory locations are restored on return entry to routine, establish counter 8 bits/word |
| 3C | F8 08 A7 | |

| | | | | | |
|-------|-------------|---|-------|-------------|---|
| 3F | F8 FF A5 | display location where entry starts at | B1 | 91 BO | load high order DMA address |
| 42 | F8 05 A8 | 5 words to shift in—another counter | B3 | 9F AO | load low order DMA address |
| 45 | F8 00 A9 | routine to give a brief delay | -- | | DMA 1 here (only once) |
| 48 | 89 FB 06 | | B5 | E2 9F AO | restore DMA address |
| 4B | 3A 48 | | -- | | DMA 2, 5, 8, 11, 14, . . . 104, and 107 |
| 4D | F8 08 AA | counter—8 words/line of display | B8 | 2F AO 8F | decrement counter, restore DMA address, load counter |
| 50 | 04 FE | load D, shift left to fill DF | -- | | DMA 3, 6, 9, 12, 15, . . . 105, and 108 |
| 52 | 04 7E 54 | load, circular shift left and restore | BB | 32 C1 9F AO | to C1 if counted out, else restore DMA address |
| 55 | 05 7E 55 | load, circular shift left and restore | -- | | DMA 4, 7, 10, 13, 16, . . . and 106 |
| 58 | 25 2A | repeat 8 times till bit shifted in to line | BF | 30 B6 | and go back to B6 |
| 5A | 8A 3A 55 | | C1 | 80 E2 | here if counted out, new DMA address and change to 4 bursts with same DMA address |
| 5D | 28 88 14 | repeat till a bit is shifted into all five lines | -- | | new DMA 1 here |
| 60 | 3A 4D | | C3 | E2 20 AO | restore DMA address |
| 62 | 84 FF 05 A4 | reset R(4) | -- | | DMA 2 here |
| 66 | 27 87 | check if all bits shifted in if so return | C6 | E2 20 AO | restore DMA address |
| 68 | 32 3B | | -- | | DMA 3 here |
| 6A | 30 3F | else go and shift another set in, this is end of subroutine | C9 | E2 20 AO | restore DMA address |
| 6C-A2 | -- | portion of memory used to store data for display | -- | | DMA 4 here |
| A3 | 72 70 | return from interrupt | CC | 3CC1 30 A3 | back for another line or exit if done |
| A5 | C4 | entry to interrupt | DO-FF | | the displayed locations, DO-D7 are the locations displayed with 108 identical DMA bursts and should be cleared to blank upper part of display |
| A6 | 22 78 22 52 | save X, P and D | | | |
| AA | F8 24 AF | counter for 108 DMA bursts | | | |
| AD | 19 | increment R(9), timer | | | |
| AE | F8 DO BF | load starting address for DMA | | | |

Diagram — Program 4



The above message was written and encoded on a piece of graph paper 88 squares wide and 5 squares high. Your own messages can be done in a similar way.

Program 5 — Horse Race

Four horses, represented by bars, race across the screen. The first to reach the finish line blinks to indicate the winner. The idea is to bet on the outcome.

This program uses the interrupt routine which displays 64 store locations divided into 128 squares (see program 2). The squares are addressed individually in hex (00 through 7F) as they are in set 1 of program 2. Horse 1 lights squares 10 through 1F, horse 2 lights 30 through 3F, horse 3 lights 50 - 5F, and horse 4 lights squares 70 - 7F. The current position of each horse is held in registers A through D and each horse's position is incremented randomly to produce a fair race. To run the program simply switch to run; the horses appear at the starting gate and then race across the screen.

Registers 1, 2, 3, 4, 5, 6, and 7 must be set if this program is to be run in an expanded memory system.

Use of Registers

| | High | Low |
|------|-------------------------------------|--|
| R(0) | | DMA address |
| R(1) | | interrupt address |
| R(2) | | stack |
| R(3) | | program counter |
| R(4) | | memory pointer and work |
| R(5) | | address of delay subroutine |
| R(6) | | address of display subroutine |
| R(7) | | points to table of winning positions |
| R(8) | used to generate random number | used to generate random number |
| R(9) | | counter for timer in interrupt routine |
| R(A) | work in display routine | position of horse 1 |
| R(B) | | position of horse 2 |
| R(C) | holds counter for interrupt routine | position of horse 3 |
| R(D) | | position of horse 4 |
| R(E) | | used to hold variable instruction |
| R(F) | work in interrupt | work in interrupt |

Program Listing

| Add. | Code | Notes |
|------|-------------|---|
| 00 | F8 93 A1 | interrupt address |
| 03 | F8 BE A2 | stack address |
| 06 | F8 09 A3 | assign program counter |
| 09 | D3 | program counter now R(3) |
| OA | F8 6C A5 | address of delay subroutine |
| OD | F8 77 A6 | address of display subroutine |
| 10 | F8 04 BC | counter for interrupt routine |
| 13 | F8 BF A4 | clear the display locations |
| 16 | F8 00 54 | |
| 19 | 14 84 | |
| 1B | 3A 16 24 | |
| 1E | E2 69 | R(2) is X, turn on TV |
| 20 | F8 10 AA D6 | set horses at the starting gate |
| 24 | F8 30 AB D6 | and call the display routine to |
| 28 | F8 50 AC D6 | show them |
| 2C | F8 70 AD D6 | |
| 30 | D5 | call delay routine |
| 31 | F8 62 A7 | point R(7) is address of 1st entry of table of winning positions |
| 34 | F8 1A AE | load R(E).0 with instruction 1A |
| 37 | 88 F6 A4 | this code together with instructions 8C through 90 generates a random byte in D |
| 3A | E4 98 A4 | |
| 3D | B8 F6 A4 | |
| 40 | 98 F4 | |
| 42 | E7 | |
| 43 | FA 01 | |
| 45 | 32 59 | make R(7) the X register is random no. odd or even? |
| 47 | D5 | if even go to 59 do not increment horses position, if odd go on |
| 48 | 8E 53 XX | call delay routine |
| 4B | 8E FC 70 | location 4A becomes 1A, 1B, 1C, or 1D and increments a horses position |
| 4E | 53 XX | change instruction to 8A, 8B, 8C, or 8D |
| 50 | F3 | location 4F becomes 8A, 8B, 8C, or 8D and loads a horses position |
| 51 | 32 66 | check to see if horse has won by x'or against table of winning positions |
| 53 | 8E FC 70 | if it has won go to 66 and blink |
| 56 | 53 XX D6 | else load horses position to D and call display routine |
| 59 | 8E FB 1D | is instruction 1D? |

| | | |
|----|-------------|--|
| 5C | 32 31 | yes—go to 31 to start over |
| 5E | 1E 17 30 37 | no-increment R(E) and R(7) and go back to generate another random number |
| 62 | 1F 3F 5F 7F | table of winning positions |
| 66 | 07 D6 D5 | comes here when horse has won and blinks by calling de- lay and display |
| 69 | 30 66 | |
| 6B | D3 | return from delay subroutine, |
| 6C | F8 00 A9 | R(9) is incremented in the |
| 6F | 89 FB 08 | interrupt routine |
| 72 | 3A 6F | |
| 74 | 30 6B | |
| 76 | D3 | return from display routine, this display routine works like program 2, the square identifier is passed to the sub- routine in D |
| 77 | BA F6 | enter subroutine, save D in R(A).1 and shift right |
| 79 | FC BF A4 | add starting address of dis- played locations, save in R(4) |
| 7C | 9A F6 | bring back passed byte, shift |
| 7E | 3B 86 | to 86 if passed byte even |
| 80 | 04 FB OF 54 | byte odd, load x'or against OF and restore |
| 84 | 30 76 | return from subroutine |
| 86 | 04 FB FO 54 | byte even, load, x'or against FO and restore |
| 8A | 30 76 | return from subroutine, end of subroutine |
| 8C | 18 98 | part of interrupt routine, here |
| 8E | FF 01 B8 | just before returning, add 1 to R(8).0 and subtract 1 from R(8).1 as part of random number generator |
| 91 | 72 70 | return from interrupt |
| 93 | C4 | entry to interrupt—this inter- rupt routine is very similar to |
| 94 | 22 78 22 52 | that given in program 2 but is |
| 98 | 9C AF | located in a different position |
| 9A | F8 BF BF | in memory |
| 9D | 91 BO | |
| 9F | 9F AO 19 | |
| A2 | 30 A7 9F AO | |
| A6 | 2F AO 8F | |
| A9 | 32 AF 9F AO | |
| AD | 30 A4 | |
| AF | 9F AO | |
| B1 | AO 9C 34 8C | |
| B5 | AF 80 BF | |
| B8 | 30 A7 | end of interrupt routine |

| | |
|-------|---------------------|
| BA-BB | unused |
| BC-BD | stack |
| BF-FE | displayed locations |

Program 6 — Pattern Generator

Program 6, like program 5, uses the 64 byte interrupt routine to display 8 rows of 16 squares. The area is divided into 4 quarters and used to generate patterns (see diagram).

An additive shift byte is used together with the identification numbers of up to 15 points in the upper left quarter of the display area. The logical exclusive or operation is carried out with each point and the display field and also with 3 related points in the 3 other quarters. The relationship between the points is such that the dotted lines in the diagram are "mirrors". After all of the entered points are exclusively ored the additive shift is applied to each point and the process repeats indefinitely.

To run the program put the additive shift in the switch byte and turn to run, the shift byte will be stored and displayed. Then enter the points to be displayed. Terminate entry with EE (entry is otherwise automatically terminated when 15 points have been entered) and pattern generation begins. It is suggested that a trial run be made with 11 as the shift byte and 00, 01, 12, 23, and EE as the data points.

Registers 1, 2, 3, 4, 5, 6, and 7 must be set if this program is to be run in an expanded memory system.

Note: The following changes will generate a pattern in which the points in the upper left quarter are exactly repeated in the other 3 quarters. Change memory location 65 to F4 and the table entries at locations 80 through 83 to 00, 08, 40 and 48.

Use of Registers

| High | Low |
|------|--|
| R(0) | DMA address |
| R(1) | interrupt address |
| R(2) | stack |
| R(3) | program counter |
| R(4) | memory pointer to display locations |
| R(5) | address of display subroutine |
| R(6) | points to locations of entered points |
| R(7) | points to a table of 4 constants |
| R(8) | counter |

| | | | | |
|------|--|-------------------------------------|-------------|--|
| R(9) | counter for timer in interrupt routine | | | through 83 is display subroutine |
| R(B) | work for display sub-routine | counter—4 bytes displayed at a time | 3B F8 B0 A6 | reload starting address of table of entered points |
| R(C) | holds counter for interrupt routine | | 3E 8E A8 | load no. of points to R(8).0 |
| R(D) | | holds shift byte | 40 06 FA 37 | load point, and against 37 to get into upper left quarter of display |
| R(E) | | counter, number of bytes entered | 43 52 64 22 | save it on stack, display it |
| R(F) | work in interrupt | work in interrupt | 46 F8 00 A9 | brief delay between calls to subroutine |
| | | | 49 89 FB 20 | |
| | | | 4C 3A 49 | |
| | | | 4E D5 | |

Program Listing

| Add. | Code | Notes | | |
|------|-------------|---|----------------|---|
| 00 | F8 86 A1 | interrupt address | 4F 8D | call display subroutine, display point is passed on stack |
| 03 | F8 AF A2 | stack | 50 E6 | get shift byte |
| 06 | F8 09 A3 D3 | program counter | 51 F4 56 | make pointer to table of entered bytes R(X) |
| OA | F8 04 BC | constant for interrupt routine | 53 E2 | add shift byte to table entry and write new entry |
| OD | F8 5D A5 | address of display routine | 54 16 28 | restore R(X) as 2 |
| 10 | F8 B0 A6 | starting address of table of entered points | 56 88 | shift pointer to next location and decrement counter |
| 13 | F8 BF A4 | starting address of displayed locations | 57 3A 40 | check, done with all points? |
| 16 | F8 00 AE | clear display locations and leave R(E).0 equal to 00, | 59 30 3B | no—go and do next point |
| 19 | 54 14 84 | R(E).0 will later contain number of points entered | -- | yes—start over |
| 1C | 3A 16 | | | start of display subroutine, this display routine works like program 2 except that that 4 squares are changed, the one passed on the stack and 3 related ones |
| 1E | 24 | | | restore R(X) as 2 and return to main |
| 1F | E2 69 | R(X) is 2, turn on TV | | entry point, load counter to display 4 points |
| 21 | 6C 64 22 | read and display shift byte | | load starting address of table of four constants |
| 24 | AD | save shift byte in R(D).0 | 5B E2 D3 | make table pointer R(X) |
| 25 | 3F 25 37 27 | wait for in on, off | 5D F8 04 AB | load saved point and x'or against one of the table entries |
| 29 | 6C 64 22 | read and display entered point | 60 F8 80 A7 | shift result right |
| 2C | FB EE | check, is it EE? | | add starting address of displayed locations and save in R(4) |
| 2E | 32 3B | yes—all points entered, go to 3B | 63 E7 | get result back, shift |
| 30 | F0 56 | no—save in table of entered points | 64 02 F3 BB | result even—go to 75 |
| 32 | 16 1E | and increment table pointer and counter for number of points | 67 F6 | result odd—x'or location against 0F and replace |
| 34 | 8E FB 0F | check counter, 15 points entered? | 68 FC BF A4 | go to 79 |
| 37 | 32 3B | yes—table full, go to 3B | 6B 9B F6 | |
| 39 | 30 25 | no—go back for more | 6D 3B 75 | |
| -- | | all points entered, code from 3B through 5A controls display while code from 5B | 6F 04 FB 0F 54 | |
| | | | 73 30 79 | |

| | | | | | |
|----|-------------|---|-------|-------------|--------------------------|
| 75 | 04 FB F0 54 | result even-x'or location against F0 and replace | 8D | F8 BF BF | and program 5 |
| 79 | 2B 17 8B | decrement counter, increment table pointer, examine counter | 90 | 91 B0 | |
| 7C | 3A 64 | not done—go to 64 | 92 | 9F A0 19 | |
| 7E | 30 5B | done—return from subroutine | 95 | 30 9A 9F A0 | |
| 80 | 00 0F 70 7F | table of constants, end of dis- play subroutine | 99 | 2F A0 8F | |
| 84 | 72 70 | return from interrupt | 9C | 32 A2 9F A0 | |
| 86 | C4 | entry to interrupt—this inter- rupt routine is very similar to | A0 | 30 97 | |
| 87 | 22 78 22 52 | the one used in program 2 | A2 | 9F A0 | |
| 8B | 9C AF | | A4 | A0 9C 34 84 | |
| | | | A8 | AF 80 BF | |
| | | | AB | 30 9A | end of interrupt routine |
| | | | AD-AF | | stack |
| | | | BO-BE | | entered bytes |
| | | | BF-BE | | displayed locations |

Diagram – Program 6

| | | | | | | | |
|----|----|----|----|----|----|----|----|
| 00 | 01 | 02 | 03 | 04 | 05 | 06 | 07 |
| 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 |
| 20 | 21 | 22 | 23 | 24 | 25 | 26 | 27 |
| 30 | 31 | 32 | 33 | 34 | 35 | 36 | 37 |

Program 7 – Stop Watch

This stop watch program displays hours, minutes, and seconds and uses the interrupt routine of program 4. A character table is used to generate the numbers. The operation of the program is simple; R(9) is incremented 61 times a second in the interrupt routine, and this count is propagated through registers A through E to generate seconds, minutes, and hours. Whenever the contents of registers A through E are changed a display routine is called to indicate the new time. Both the timing routine and the display routine make use of variable code and one should read the section on variable code if one has difficulty following the program. Note that with the interrupt routine of program 4 the time is displayed in the last five lines of the display while the top 27/32 of the display is blanked.

On entry to the program the time is reset to zero (however the display is not cleared) and on pushing the in button the timer is started, a second push of the in button stops it, and a third push starts it again from where it left off. To reset the stop watch turn the run switch off and then on again. The maximum time that can be shown is 9 hours, 59 minutes, and 59 seconds.

Registers 1, 2, 3, 4, 5, and 6 must be set if this program is to be run in an expanded memory system.

Use of Registers

| | High | Low |
|------|-------------------|--|
| R(0) | | DMA address |
| R(1) | | interrupt address |
| R(2) | | stack |
| R(3) | | program counter |
| R(4) | | memory pointer |
| R(5) | | memory pointer |
| R(6) | | memory pointer, also used to store variable instruction |
| R(7) | | holds variable instruction |
| R(8) | | counter in display routine |
| R(9) | | counts 61 times a second, incremented in interrupt routine |
| R(A) | | seconds counter, units |
| R(B) | | seconds counter, tens |
| R(C) | | minute counter, units |
| R(D) | | minute counter, tens |
| R(E) | | hours counter, units |
| R(F) | work in interrupt | work in interrupt |

Program Listing

| Add. | Code | Notes |
|------|-------------|---|
| 00 | F8 7E A1 | interrupt address |
| 03 | F8 49 A2 | stack |
| 06 | F8 09 A3 | program counter |
| 09 | D3 E2 | establish program counter, X register |
| 0B | F8 00 AE AD | clear registers for hours, minutes and seconds |
| 0F | AC AB AA | |
| 12 | 69 | turn on TV |
| 13 | 37 13 | wait for in off—comes here from 21 |
| 15 | 3F 15 37 17 | wait for in on, off |
| 19 | F8 00 A9 | set interrupt counter to 0 |
| 1C | 30 4F | go to display routine—note display routine always returns to 1E |
| 1E | 89 FB 3D | here we look for 61 interrupt cycles |
| 21 | 37 13 | or did someone push the in button |
| 23 | 3A 1E A9 | to 1E unless 1 second has passed—else goes on first clear R(9).0 |
| 26 | F8 3D A5 | load address of variable code to R(5) |
| 29 | F8 4A A4 | load starting address of table of 5 entries to R(4) |
| 2C | E4 | make R(4) the X register |
| 2D | F8 1A A6 | load instruction (1A) to R(6).0, first obeyed as 1A, then 1B, . . . and finally as 1E, increments registers |
| 30 | 86 53 XX | carry out above instruction |
| 33 | FC 90 55 | write AA to M(3D), first pass, second time write AB, . . . and finally write AE |
| 36 | FF 20 53 | create and obey 8A first time, then 8B, . . . and finally 8E |
| 39 | XX | contents of one of the counting registers loaded here |
| 3A | F3 | x'or against table entry, maximum count allow in register |
| 3B | 3A 4F | if register contents are not equal to maximum go to display routine |
| 3D | XX 86 | 3D becomes one of AA - AE, register which has reached maximum value is here zeroed |

| | | | | | |
|----|----------|---|----|-------------|---|
| | | out, load instruction in R(6) to see where we are | 72 | 30 68 | go and write it |
| 3F | FB 1E | if instruction is 1E we're done and should go to display routine | 74 | D9 DB DC | table of 5 locations to write characters to |
| 41 | 32 4F | | 78 | DE DF | unused |
| 43 | 16 60 | else change instruction, increment X to point to new table entry | 79 | XX XX XX | |
| 45 | 30 30 | and go back to 30 for another pass | 7C | 72 70 | return from interrupt |
| 47 | XX XX XX | stack | 7E | C4 | entry to interrupt—this interrupt routine is identical to |
| 4A | 0A 06 0A | table of maximum values allowed in the counting registers | 7F | 22 78 22 52 | that used in program 4 except |
| 4D | 06 0A | | 83 | F8 24 AF | that it is located in a different place in memory |
| -- | | the display routine starts here, it uses a character table to show the contents of registers A through E in the display locations | 86 | 19 | |
| 4F | F8 74 A6 | load starting address of table of memory locations to write to | 87 | F8 D0 BF | |
| 52 | F8 8F A7 | load variable instruction (8F) to R(7), it will be obeyed as one of 8E, 8D, 8C, 8B, or 8A | 8A | 91 B0 9F A0 | |
| 55 | 87 FB 8A | check variable instruction, if it's 8A we're done and should return | 8E | E2 9F A0 | |
| 58 | 32 1E | all done—return | 91 | 2F A0 8F | |
| 5A | 27 | more to do—decrement variable instruction | 94 | 32 9A 9F A0 | |
| 5B | 46 A4 | load display location to R(4) | 98 | 30 8F | |
| 5D | 87 53 XX | obey variable instruction, loads contents of one of the counting registers to D | 9A | 80 E2 | |
| 60 | FC A9 | add the starting address of the character table | 9C | E2 20 A0 | |
| 62 | A5 05 A5 | and pick up the address of the character to be displayed | 9F | E2 20 A0 | |
| 65 | F8 05 A8 | establish counter, 5 words per character | A2 | E2 20 A0 | |
| 68 | 45 54 | load and write word of character | A5 | 3C 9A | |
| 6A | 28 88 | decrement no. words/character | A7 | 30 7C | end of interrupt routine |
| 6C | 32 55 | to 55 for next register contents if done | -- | | locations A9 through CF are the character table |
| 6E | 84 | else load R(4).0 | A9 | B3 | starting address of 0 |
| 6F | FC 08 A4 | and set address for next word to write | AA | C6 | starting address of 1 |
| | | | AB | BB | starting address of 2 |
| | | | AC | B9 | starting address of 3 |
| | | | AD | CB | starting address of 4 |
| | | | AE | B7 | starting address of 5 |
| | | | AF | BD | starting address of 6 |
| | | | BO | C5 | starting address of 7 |
| | | | B1 | BF | starting address of 8 |
| | | | B2 | C1 | starting address of 9 |
| | | | B3 | F0 | character 0, (5 words/character |
| | | | B4 | 90 | |
| | | | B5 | 90 | |
| | | | B6 | 90 | |
| | | | B7 | F0 | character 5 |
| | | | B8 | 80 | |
| | | | B9 | F0 | character 3 |
| | | | BA | 10 | |
| | | | BB | F0 | character 2 |
| | | | BC | 10 | |
| | | | BD | F0 | character 6 |
| | | | BE | 80 | |
| | | | BF | F0 | character 8 |
| | | | C0 | 90 | |
| | | | C1 | F0 | character 9 |
| | | | C2 | 90 | |
| | | | C3 | F0 | |
| | | | C4 | 10 | |
| | | | C5 | F0 | character 7 |
| | | | C6 | 10 | character 1 |

| | |
|----|----|
| C7 | 10 |
| C8 | 10 |
| C9 | 10 |
| CA | 10 |
| CB | A0 |
| CC | A0 |
| CD | F0 |
| CE | 20 |
| CF | 20 |

character 4

display locations are DO through FF, XX means doesn't matter, other locations must be set to give

legible display

Add. Code

| | |
|----|-------------------------|
| D0 | 00 00 00 00 00 00 00 00 |
| D8 | 00 XX 00 XX XX 00 XX XX |
| E0 | 00 XX 20 XX XX 20 XX XX |
| E8 | 00 XX 00 XX XX 00 XX XX |
| F0 | 00 XX 20 XX XX 20 XX XX |
| F8 | 00 XX 00 XX XX 00 XX XX |

Note: It is probably easiest to clear memory before loading this program; then set locations E2, E5, F2 and F5 to 20.

Program 8—Twelve Hour Clock

This clock program is similar to the stop watch program but displays a settable twelve hour clock instead of a stop watch. It was not possible to use the interrupt routine of program 4 and a one page interrupt is used instead. Register 9 is incremented within the interrupt routine to act as the timer. The count in register 9 is propagated through registers A to F to generate seconds, minutes and hours.

The clock is a 12 hour one and changes from 12:59:59 to 01:00:00 at 1:00 AM and 1:00 PM. To set the clock each digit is set individually. For example say we wish to set the time 12:03:48. First turn to run, then:

set tens digit of hours (01) to switch byte, push in,
 set units digit of hours (02) to switch byte, push in,
 set tens digit of minutes (00) to switch byte, push in,
 set units digit of minutes (03) to switch byte, push in,
 set tens digit of seconds (04) to switch byte, push in,
 set units digit of seconds (08) to switch byte, push in,
 and finally to start clock,
 set 00 on switch byte, push in.

Registers 1, 2, 3, 4, 5, and 6 must be set if this program is to be run in an expanded memory system.

Use of Registers

Register use is the same as in program 7 except that register F is used as the tens of hours counter and is not employed in the interrupt routine.

Program Listing

| Add. | Code | Notes |
|------|-------------|--|
| 00 | F8 8D A1 | interrupt |
| 03 | F8 4D A2 | stack |
| 06 | F8 09 A3 | program counter |
| 09 | D3 E2 | establish program counter, X register |
| 0B | F8 1A A4 | address of location whose contents will be variable code |
| 0E | F8 AF A5 | variable instruction to R(5).0 first time will be AF |
| 11 | 85 54 | write instruction to 1A, first time AF |
| 13 | 3F 13 37 15 | wait for in on, off |
| 17 | 6C 64 22 | read and display switch byte |
| 1A | XX | variable instruction, saves switch byte in a register to set clock, first AF, then AE, etc. and finally A9 |

| | | |
|----|----------|--|
| 1B | 85 FB A9 | check is instruction A9 |
| 1E | 32 23 | yes—go to 23 and start clock |
| 20 | 25 30 11 | no—change variable instruction and go back to read another switch byte. |
| 23 | 69 | ready to start clock—turn on TV |
| 24 | 89 FB 3D | get timer from interrupt—is it 61 |
| 27 | 3A 24 | no-go and check again |
| 29 | A9 | yes—reset interrupt timer and go on |
| -- | | This portion of code—2A through 53 is almost identical to that of program 7 (26 through 4E) it propagates the second count through the registers |
| 2A | F8 41 A5 | load address of variable code to R(5) |
| 2D | F8 4E A4 | load starting address of table of 6 entries to R(4) |
| 30 | E4 | make R(4) the X register |
| 31 | F8 1A A6 | load instruction (1A) to R(6).0 first obeyed as 1A, then 1B, and finally 1F, increments counting registers |
| 34 | 86 53 XX | carry out above instruction |
| 37 | FC 90 55 | write AA to location 41, first pass, then AB, . . . and finally write AF |
| 3A | FF 20 53 | create and obey 8A first time then 8B, . . . and finally 8F |
| 3D | XX | 8A through 8F obeyed here, loads one of counting registers to D |
| 3E | F3 | x'or against table entry, maximum count allowed in register |
| 3F | 3A 54 | if register contents are not equal to maximum go to check hours and then to display routine |
| 41 | XX 86 | 41 becomes one of AA—AF, register which has reached maximum value is here zeroed out, load instruction in R(6) to see where we are |
| 43 | FB 1F | if instruction is 1F we're done |
| 45 | 32 54 | and go to check hours and then to display routine |

| | | | | | |
|----|----------|---|-------|-------------|--|
| 47 | 16 60 | else change instruction, increment X to point to new table entry | 73 | A5 05 A5 | and pick up the address of the character to be displayed |
| 49 | 30 34 | and go back to 34 for another pass | 76 | F8 05 A8 | establish counter, 5 words per character |
| 4B | XX XX XX | stack | 79 | 45 54 | load and write word of character |
| 4E | 0A 06 0A | table of maximum values allowed in the counting registers | 7B | 28 88 | decrement no. words/character |
| 51 | 06 0A 0A | | 7D | 32 66 | to 66 for next register contents if done |
| -- | | comes here to check hours, following code resets 13:00:00 to 01:00:00, else does nothing | 7F | 84 | else load R(4).0 |
| 54 | 8F F6 | load hours counter, tens, shift right | 80 | FC 08 A4 | and set address for next word to write |
| 56 | 3B 60 | to display if DF is zero, else go on | 83 | 30 79 | go and write it |
| 58 | 8E FB 03 | loads hours counter, units, is it 3? | 85 | D8 D9 DB | table of 6 locations to write characters to |
| 5B | 3A 60 | if not go to display | 88 | DC DE DF | end of display routine—begin one page interrupt |
| 5D | AF AE 1E | if R(E).0 is 3 reset hours to 01 | 8B | 72 70 | return from interrupt |
| -- | | this display routine is almost identical to the one in program 7, it returns to location 24 | 8D | C4 | entry to interrupt |
| 60 | F8 85 A6 | entry point, load starting address of table of memory locations to write to | 8E | 22 78 22 52 | save X, P and D |
| 63 | F8 90 A7 | load variable instruction (90) to R(7), it be obeyed as one of 8F, 8E, 8D, 8C, 8B, or 8A | 92 | E2 E2 19 | increment R(9) |
| 66 | 87 FB 8A | check variable instruction, if it's 8A we're done and should return | 95 | 91 B0 | load high order DMA address |
| 69 | 32 24 | all done—return | 97 | F8 00 A0 | load low order DMA address |
| 6B | 27 | more to do—decrement variable instruction | 9A | 80 E2 | DMA address to D |
| 6C | 46 A4 | load display location to R(4) | -- | | DMA 1, here |
| 6E | 87 53 XX | obey variable instruction, loads contents of one of the counting registers to D | 9C | E2 20 A0 | restore DMA address |
| 71 | FC A9 | add the starting address of the character table | -- | | DMA 2, here |
| | | | 9F | E2 20 A0 | restore DMA address |
| | | | -- | | DMA 3, here |
| | | | A2 | E2 20 A0 | restore DMA address |
| | | | -- | | DMA 4, here |
| | | | A5 | 3C 9A | not done, back to set new DMA address |
| | | | A7 | 30 8B | done, return |
| | | | A9-FF | | The character table and display locations are the same as for program 7. |

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Dear Elf owner:

Thank you for your order for my graphics booklet; hope you like it.

I would like to start work on a third booklet but have yet to chose a topic. Two things seem reasonable to me:

1.) more ambitious programs which would require expanded memory - TV games or possibly mathematical packages.

2.) programs which would require a connection to the outside world via the Q line and flags, or via input and output ports.

If you have suggestions about programs that you would like to see written and that would have a reasonably wide appeal I would like to hear them.

When you read the graphics booklet you will see that xx is used fairly often. This symbol means that the contents of the memory location involved don't matter, but can be anything. Of course something has to be there, and when loading memory these positions can be filled with 00 as an aid to debugging. If the program fails to work the memory locations which remain set at 00 can help to find the place where an error in loading has occurred.

Paul Moews