

OS 之道

操作系统简介

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Think OS

A Brief Introduction to Operating Systems

Version 0.7.4

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The L^AT_EX source for this book is available from <http://greenteapress.com/thinkos>.

序

在很多计算机科学课程中, 操作系统始终是一个高级话题. 学生在学习它之前, 可能已经学习了如何用 C 编程, 甚至很可能已经学过计算机体系架构的课程. 通常本课程的目标是让学生了解操作系统的设计和实现, 同时期待其中一些学生在这个领域持续研究, 或者编写部分操作系统.

这本书希望可以满足不同人群的不同目标, 我编写此书是为了讲授欧林学院的软件系统的课程.

多数来上课的学生已经学过了如何用 Python 编程. 所以一个目标是能够帮助他们学习 C 语言. 对于这部分内容, 我使用 O'Reilly Media 出版社出版, Griffiths 和 Griffiths 编写的 *Head First C* 一书, 来充实本书的内容.

我的学生中很少会从事编写操作系统, 但是很多人将用 C 来编写底层应用, 以及进入嵌入式领域工作. 本课程包括操作系统, 网络, 数据库, 以及嵌入式系统, 但同时着重强调了程序员需要了解的主题.

本书不会假设你学过计算机架构. 随着课程学习, 我将会解释我们所需要的内容.

如果这本书是成功的, 那么它应该能够让你更好地理解程序运行时发生了什么, 以及做什么可以令程序运行更加高效.

第一章介绍了编译语言和解释语言的差异, 以及编译器工作原理. 推荐阅读: *Head First C* 第 1 章.

第二章解释了操作系统如何使用进程隔离程序运行.

第三章解释了虚拟内存和地址转换. 推荐阅读: *Head First C* 第 2 章.

第四章介绍了文件系统和数据流. 推荐阅读: *Head First C* 第 3 章.

第五章讲述了数字、字母和其他值如何编码, 并介绍了位运算符.

第六章介绍了动态内存管理的使用及其工作原理. 推荐阅读: *Head First C* 第 6 章.

第七章介绍了缓存和内存层次结构.

第八章介绍了多任务和调度.

第九章介绍了 POSIX 线程和互斥锁. 推荐阅读: *Head First C* 第 12 章以及 *Little Book of Semaphores* 第 1 和 2 章.

第十章介绍了 POSIX 条件变量和生产者/消费者问题.

第十一章介绍了 POSIX 信号量的使用以及在 C 语言中的实现.

补充

本书目前还是早期草稿版本. 我还在努力撰写文本, 但尚缺乏一些配图. 所以某些地方, 如果配图完备了, 我确信相关解释会得到改善.

0.1 代码使用

本书的代码示例可以从 <https://github.com/AllenDowney/ThinkOS> 下载. Git 是个版本管理系统, 可以让你跟踪项目文件. Git 管理的文件集合叫做仓库. GitHub 是一个提供 Git 仓库存储和方便的 web 界面的托管服务.

我的仓库的 GitHub 主页提供了多种使用代码的方法:

- 你可以点击 **Fork** 按钮, 在 GitHub 基于我的仓库创建一个副本. 如果你还没有 GitHub 账户, 你需要创建一个. 创建副本后, 你便拥有了自己的仓库, 便可以跟踪学习本书时所编写的代码了. 然后, 你可以克隆这个仓库, 这意味着文件会被复制到你的计算机上.
- 或者你也可以克隆我的仓库. 这样便无需 GitHub 账号, 但你便不能将你的变更写回 GitHub 了.
- 如果你一点都不想用 Git, 你可以使用 GitHub 页面的按钮下载 Zip 文件.

贡献者清单

如果你有建议或更正, 请发送邮件至 downey@allendowney.com. 如果我根据你的反馈进行更正, 我将把你添加到贡献者清单 (除非你要求省略).

如果您在反馈时包含了出现错误的部分语句, 那么我便可以快速定位. 提供页码或者章节编号也可以, 只是不太容易处理. 谢谢!

- 感谢参加 Olin 学院 2014 年春季软件系统课程的学生们, 他们在本书的早期草稿阶段, 进行了大量测试. 同时纠正了很多错误, 并提出相当多的建议. 我非常感激他们的开创精神!
- James P Giannoules 指出了一处复制粘贴错误.
- Andy Engle 指出了 GB 和 GiB 之间的区别.
- Aashish Karki 指出了一些语法问题.

其他发现错别字和其他错误的人包括 Jim Tyson, Donald Robertson, Jeremy Vermast, Yuzhong Huang 和 Ian Hill.

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Chapter 1

编译

1.1 编译语言和解释语言

人们通常将编程语言分为编译型或解释型。“编译”意味着程序被翻译为机器语言,然后硬件执行;“解释”则意味着程序由软件解释器读取并执行。通常认为 C 是编译型语言,而 Python 是解释型语言。但是两者差异并非泾渭分明。

首先,很多语言是既可以被编译也可以被解释的。例如,也有 C 解释器和 Python 编译器。其次,像 Java 一样,有些语言会使用混合方法,将程序编译为一种中间语言,然后在解释器中运行翻译后的程序。Java 使用了一种叫做 Java 字节码的中间语言,和机器语言类似,但它需要由软件解释器执行,也就是 Java 虚拟机 (JVM)。

所以被编译还是被解释,并不是一个语言的内在特征;然而无论如何,编译语言和解释语言之间确实会存在一些普遍的差异。

1.2 静态类型

多数解释语言支持的是动态类型,但是编译语言通常受限于静态类型。在静态类型语言中,你可以通过查看程序便能区分变量引用的类型。而在动态类型语言中,直到程序运行,你才能确切知道变量类型。通常,**静态**是指在编译时发生的事情,而**动态**是指运行时发生的事情。

例如,可以用 Python 编写如下函数:

```
def add(x, y):  
    return x + y
```

看这个代码, 你很难在运行时确定 `x` 和 `y` 所引用的类型. 这个函数可能被调用多次, 每次都可能是不同类型的值. 任何支持加法运算的值都可以运行; 其他类型则会引发异常或**运行时错误**.

而在 C 中, 同样函数编写如下::

```
int add(int x, int y) {  
    return x + y;  
}
```

函数第一行包括对参数和返回值的**类型声明**: `x` 和 `y` 被声明为了整数, 这就意味着在编译时, 我们可以检查这种类型是否支持加法运算 (此处支持). 返回值也同样被声明为了整数.

因为这些声明的存在, 在程序其他位置调用这个函数时, 编译器便可以检查提供的参数是否具有正确的类型, 以及返回值是否使用正确.

这些检查发生在程序运行之前, 所以错误可以被提前发现. 更重要的, 可以在未执行的程序部分发现错误. 此外, 这些检查在运行时便无需发生, 这也是编译语言比解释语言运行更快的一个原因.

编译时声明类型还能节省空间. 在动态语言中, 程序运行时, 变量名称需要保存在内存中, 同时它们会被程序频繁访问. 例如, 在 Python 中, 内置函数 `locals` 会返回一个包含变量名称的字典. 下面一个用 Python 解释器执行的示例:

```
>>> x = 5  
>>> print locals()  
{'x': 5, '__builtins__': <module '__builtin__' (built-in)>,  
  '__name__': '__main__', '__doc__': None, '__package__': None}
```

这说明在程序运行时, 变量名是存储在内存中的 (以及一些作为默认运行环境一部分的其他值).

编译语言中, 变量名存在于编译时, 而不是运行时. 编译器会为每个变量选择一个位置, 同时将这些位置记录为编译后的程序的一部分.¹ 变量的位置被称为**地址**. 运行时, 每个变量的值被存储在它的地址上, 变量的名称则不会被存储 (除非编译器出于调试的目的而添加它们).

1.3 编译流程

作为程序员, 你应该对编译期间发生了什么有个大致脉络. 如果你熟悉这个过程, 便可以帮你了解错误信息, 调试代码, 以及避免常见的陷阱.

编译步骤如下:

¹这只是一个简单描述, 稍后我们会详细说明.

1. 预处理: C 是一种包括**预处理指令**的语言, 这些指令是在编译之前起效的. 比如, `#include` 指令会将另一个文件的源代码插入到指令的位置.
2. 解析: 解析期间, 编译器会读取源码并构建程序的内在结构, 也就是**抽象语法树**. 在此期间检查出来的错误, 通常是语法错误.
3. 静态检查: 编译器会检查变量和值的类型是否正确, 调用函数是否使用了正确数量和类型的参数, 等等. 此阶段侦测到的错误通常是**静态语义**错误.
4. 代码生成: 编译器会读取程序的内在结构, 并生成机器码或字节码.
5. 链接: 如果程序使用了库中定义的值和函数, 编译器需要找到合适的库, 并囊括所需的代码.
6. 优化: 在整个流程的各个阶段, 编译器可以转换程序, 生成运行更快, 或更加节约空间的代码. 大部分的优化是一些简单的修改, 可以消除明显的耗费, 但有些编译器可以执行更复杂的分析和代码转换.

通常运行 `gcc` 时, 它会执行所有这些步骤, 生成一个可执行文件. 比如, 下面是一个极简的 C 程序代码:

```
#include <stdio.h>
int main()
{
    printf("Hello World\n");
}
```

如果你将上述代码保存为一个名为 `hello.c` 的文件, 你便能像这样对其进行编译和执行:

```
$ gcc hello.c
$ ./a.out
```

默认情况下, `gcc` 将可执行代码存储在名为 `a.out` (起初表示 “assembler output”) 的文件中. 第二行会运行可执行文件. 前缀 `./` 是告诉 `shell` 在当前目录中查找.

通常最好使用 `-o` 标识为可执行文件提供一个合适的名称:

```
$ gcc hello.c -o hello
$ ./hello
```

1.4 目标码

使用 `-c` 标识, 可以令 `gcc` 编译程序并生成机器码, 但并不会链接库中函数, 也就不会生成可执行文件:

```
$ gcc hello.c -c
```

结果是一个名为 `hello.o` 的文件, 其中 `o` 表示**目标码**, 也就是编译后的程序. 目标码不可执行, 但可以链接到可执行文件中.

UNIX 命令 `nm` 可以读取目标文件, 并生成涉及其所定义和应用的名称的信息. 例如:

```
$ nm hello.o
0000000000000000 T main
                 U puts
```

输出结果表明 `hello.o` 定义了名称 `main`, 以及名为 `puts` 的函数, 这个名字表示 “put string”. 在这个例子中, `gcc` 通过将庞大复杂的 `printf` 函数, 替换为相对简单的 `puts`, 以进行优化.

你可以使用 `-O` 标识来控制 `gcc` 进行多大程度的优化. 默认情况下, 它优化得很少, 这会令调试相对简单. 而选项 `-O1` 会开启常用且安全得优化. 更高得数字会开启更多得优化, 而这也就会需要更长得编译时间.

理论上, 优化除了加速程序之外, 不应该改变程序的行为. 但如果你的程序存在微妙的错误, 那你可能会发现, 优化会令错误暴露或消失. 通常在开发新代码时, 最好关闭优化. 一旦程序正常运行, 并能通过适当的测试, 你便可以开启优化, 并确认测试仍然通过.

1.5 汇编码

类似于 `-c` 标识, `-S` 标识会令 `gcc` 编译程序并生成汇编码, 而这基本上是一种人可读的机器码形式.

```
$ gcc hello.c -S
```

结果是一个名为 `hello.s` 的文件, 看起来像这样:

```
.file          "hello.c"
.section       .rodata
.LC0:
.string       "Hello World"
.text
.globl        main
.type         main, @function
main:
.LFB0:
.cfi_startproc
pushq %rbp
.cfi_def_cfa_offset 16
```



```

        .cfi_offset 6, -16
        movq %rsp, %rbp
        .cfi_def_cfa_register 6
        movl $.LC0, %edi
        call puts
        movl $0, %eax
        popq %rbp
        .cfi_def_cfa 7, 8
        ret
        .cfi_endproc
.LFE0:
        .size      main, .-main
        .ident      "GCC: (Ubuntu/Linaro 4.7.3-1ubuntu1) 4.7.3"
        .section    .note.GNU-stack,"",@progbits

```

gcc 通常会被设置, 为你运行的机器生成相应的代码, 所以对我来说, 它会生成 x86 汇编语言, 这可以在 Intel, AMD 和其他厂商的多种处理器上运行. 如果你在不同架构上运行, 你可能会得到不同的代码.

1.6 预处理

现在我们在编译过程中向后再走一步. 你可以用 `-E` 标识, 仅仅只运行预处理器:

```
$ gcc hello.c -E
```

结果是预处理器的输出. 这个例子中, 它包含了 `stdio.h` 中包含的代码, 以及 `stdio.h` 中包含的所有文件, 依此类推. 在我的机器上, 结果有 800 多行代码. 由于几乎每个 C 程序都包括 `stdio.h`, 所以这 800 行代码会被编译多次. 如果像很多 C 程序一样, 你也包含了 `stdlib.h`, 那么结果要超过 1800 行代码了.

1.7 理解错误

现在我们知道了编译过程的各个步骤, 便更容易理解错误信息了. 比如, 如果 `#include` 指令中存在错误, 你会从预处理器得到下面这条信息:

```
hello.c:1:20: fatal error: stdio.h: No such file or directory
compilation terminated.
```

如果存在语法错误, 你会从编译器得到下面信息:

```
hello.c: In function 'main':
hello.c:6:1: error: expected ';' before '}' token
```

如果使用了任何标准库都没有定义的函数, 你会从链接器得到下面错误信息:

```
/tmp/cc7iAUbN.o: In function `main':  
hello.c:(.text+0xf): undefined reference to `printf'  
collect2: error: ld returned 1 exit status
```

ld 是 UNIX 链接器的名称, 如此命名是因为 “loading” 是编译过程中和链接关系密切的另一步骤.

一旦程序运行, C 便几乎不会进行运行时检查, 所以你可能只会看到极少的运行时错误. 如果你除以零, 或者执行了其他非法的浮点操作, 你会得到一个 “Floating point exception(浮点异常).” 同时, 如果你尝试读取或写入一个错误的内存位置, 你会遇到 “Segmentation fault(段错误).”

Chapter 2

进程

2.1 抽象和虚拟化

在我们谈论进程之前,我想先定义几个词汇:

- 抽象: 抽象是某个复杂事物的简化表达. 例如, 当你开车时, 你知道当你左转方向盘时, 车会左转, 反之亦然. 当然, 方向盘连接着一系列机械和(通常是) 液压系统, 这些会使车轮转动, 同时车轮会以复杂的方式和路面互动. 作为一个司机, 你通常无需考虑这些细节. 你只需要一个简单的逻辑模型来进行驾驶. 你的这个逻辑模型便是一种抽象.

同样地, 当你使用网络浏览器时, 你知道当你点击链接时, 浏览器会显示该链接所指的页面. 使这一切成为可能的软件和网络通信是复杂的. 但作为一个用户, 你不必了解其细节.

软件工程中一个重要部分就是设计这样的抽象, 这些抽象可以使用户和其他程序员能够使用强大而复杂的系统, 而无需了解其实现的细节.

- 虚拟化: 一种重要的抽象便是虚拟化, 这是创建一个令人满意的幻觉的过程.

例如, 很多公共图书馆会参与图书馆之间的协作, 允许其相互借书. 当我请求一本图书时, 有时这本书是在我当地的图书馆的书架上, 但其他时候, 它需要从另外的地方转移过来. 无论哪种情况, 当书籍可以领取时, 我都会接到通知. 我不必知道它来自哪里, 我也不需要知道我的图书馆里有哪些书. 整个系统创建了一个假象, 让我相信我的图书馆拥有世界上的任意一本图书.

我当地的图书馆里实际可用的书可能很少, 但对我来说, 可用的虚拟书籍包括合作的图书馆中的每一本书.

再举个例子,大多数计算机只连接到一个网络,但该网络会连接到其他网络,依此类推.我们所称的互联网是一组网络和协议,它们可以将数据包从一个网络转发到下一个网络.从用户或程序员的角度来看,该系统表现得好像互联网上的每台计算机都与其他计算机相连.所以物理连接的数量虽然很少,但虚拟连接的数量是非常大的.

“虚拟”一词通常在虚拟机的上下文中使用,虚拟机是一种软件,会给你一种假象,仿佛你运行在一个运行着特定操作系统的专有机上.而实际上,虚拟机可能和很多其他虚拟机一起运行在同一台计算机上,这些虚拟机可能运行着不同的操作系统.

在虚拟化的上下文中,我们有时称真实发生的事情为“物理事件”,而虚拟发生的事情为“逻辑事件”或者“抽象事件”.

2.2 隔离

工程学最重要的一个原则便是隔离:当你设计一个包括多个部件的系统时,最好将其相互之间进行隔离,如此一个部件的改变就不会给其他组件带来不良影响.

操作系统的一个重要目标便是将每个运行的程序和其他程序进行隔离,如此程序员便无需关注任何交互的可能.提供这种隔离性的软件对象叫做**进程**.

进程是一个表示当前运行程序的软件对象.我说的“软件对象”是指面向对象编程中的对象;通常,对象包含数据以及操作数据的方法.一个进程便是一个包含下面数据的对象:

- 程序的文本,通常是一系列机器语言指令.
- 程序相关的数据,包括静态数据(编译时分配)和动态数据(运行时分配).
- 任何未完成的输入/输出操作的状态.比如,如果进程正在等待从磁盘读取数据,或者等待网络数据到达,这些操作的状态便是进程的一部分.
- 程序的硬件状态,包括寄存器中的数据,状态信息以及程序计数器.程序计数器用来标识当前正在执行哪个指令.

通常一个进程运行一个程序,但也可能会加载并运行新的程序.

还有可能,也是很常见的情况是,多个进程中运行同一个程序.在这种情况下,这些进程会共享相同的程序代码,但通常数据和硬件状态不同.

多数操作系统会提供一组基础功能,以使进程之间相互隔离:

- 多任务: 多数操作系统都具有一种能力, 可以在任何时刻将运行中的进程中断, 并保存其硬件状态, 然后稍后恢复进程. 通常, 程序员不必考虑这些中断. 程序的行为就像在专用处理器上连续运行一样, 只是指令之间的时间是不可预测的.
- 虚拟内存: 多数操作系统都会创建这样的假象, 即每个进程都有自己的一段空间, 与其他进程都隔离开来. 同样地, 程序员一般无需关注虚拟内存如何工作; 他们可以按照每个程序都有一段专门的内存来进行操作.
- 设备抽象: 运行在同一个计算机上的进程, 共享磁盘驱动器, 网络接口, 图形卡, 以及其他硬件. 如果进程不经协商便直接和硬件交互, 便会引发混乱. 例如, 本来要传递给某个进程的网络数据, 可能被其他进程读取, 或多个进程向硬盘上同一个位置存储数据. 这便有赖于操作系统提供适当的抽象, 以维护秩序.

作为程序员, 你并不需要知道这些功能实现的详情. 但如果你有兴趣, 你会发现在各种隐喻之下, 有很多有趣的事情正在发生. 如果你知晓发生的细节, 那你将成为更好的程序员.

2.3 UNIX 进程

当写这本书时, 我最关注的进程是我的文本编辑器, emacs. 每隔一段时间, 我会切换到终端窗口. 这是一个运行 UNIX shell 的窗口, 它提供了命令行界面.

当我移动鼠标, 窗口管理器便会被唤醒, 查看鼠标是否在终端窗口上, 然后又会唤醒终端. 终端唤醒 shell. 如果我在 shell 中输入 `make`, 它会创建一个新的进程来执行 Make, 创建另一个进程运行 LaTeX, 然后再创建一个进程来显示结果.

如果我需要查找一些东西, 我可能会切换到另一个桌面, 这也会再次唤醒窗口管理器. 如果我点击网络浏览器的图标, 窗口管理器会创建一个进程, 运行网络浏览器. 有些浏览器, 比如 Chrome, 会为每个窗口和每个标签页都创建新的进程.

这些还只是我知道的进程. 同时, 还有很多其他进程在后台运行. 其中很多进程正在执行与操作系统相关的操作.

UNIX 命令 `ps` 可以打印正在运行的进程的信息. 如果你在终端执行它, 你可能会看到像这样的东西:

PID	TTY	TIME	CMD
2687	pts/1	00:00:00	bash
2801	pts/1	00:01:24	emacs
24762	pts/1	00:00:00	ps

第一列是唯一的数字格式进程 ID。第二列是创建进程的终端;“TTY”代表电传打印机,这是最早的机械终端。

第三列是进程使用的处理器总时间,以小时、分钟和秒为单位。在这个例子中, `bash` 是解释输入命令的 shell 的名称, `emacs` 是我的文本编辑器, `ps` 是生成此输出的程序。

默认情况下, `ps` 仅列出和当前终端相关的进程。如果你使用 `-e` 标志,你便能获取每个进程 (包括属于其他用户的进程,就我看来,这是一个安全漏洞)。

在我的系统上,目前有 233 个进程。下面是其中一些:

PID	TTY	TIME	CMD
1	?	00:00:17	init
2	?	00:00:00	kthreadd
3	?	00:00:02	ksoftirqd/0
4	?	00:00:00	kworker/0:0
8	?	00:00:00	migration/0
9	?	00:00:00	rcu_bh
10	?	00:00:16	rcu_sched
47	?	00:00:00	cpuset
48	?	00:00:00	khelper
49	?	00:00:00	kdevtmpfs
50	?	00:00:00	netns
51	?	00:00:00	bdi-default
52	?	00:00:00	kintegrityd
53	?	00:00:00	kblockd
54	?	00:00:00	ata_sff
55	?	00:00:00	khubd
56	?	00:00:00	md
57	?	00:00:00	devfreq_wq

`init` 是操作系统启动时创建的第一个进程。它会创建许多其他进程,然后闲置,直到它创建的进程都结束。

`kthreadd` 是操作系统用来创建线程的进程。后续我们将更详细地讨论线程,当前你可以将线程看作是一种进程。开头的 `k` 代表内核 (**k**ernel),这是操作系统符合核心功能 (比如创建线程) 的部分。结尾的 `d` 代表守护程序 (**d**aemon),这是运行在后台,并提供操作系统服务的进程的另一个名称。在此,“daemon”是指一种乐于助人的精神,没有任何邪恶的意思。

根据名称,你可以推断出 `ksoftirqd` 也是一个内核守护程序;具体来说,它处理软件中断请求,即“软中断 (soft IRQ)”

`kworker` 是一个内核创建的工作进程,用来为内核处理一些任务。

通常有多个进程运行着内核服务. 在我的系统上, 现在有 8 个 `ksoftirqd` 进程和 35 个 `kworker` 进程.

我不会详细讨论其他进程, 但如果你感兴趣, 你可以搜索更多关于它们的信息. 你应该在你的系统上运行 `ps`, 并将结果和我的进行对比.

Chapter 3

虚拟内存

3.1 一点信息论知识

比特 (bit) 是一个二进制的数位, 也是一个信息单元. 如果你有一比特, 便可以表示两种可能性中的一种, 通常写作 0 和 1. 如果你有两比特, 那边有 4 种组合, 分别是 00, 01, 10, 和 11. 通常, 如果你有 b 个比特, 你便可以表示 2^b 个值. 一个**字节 (byte)** 是 8 个比特, 所以可以表示 256 个值.

反过来想, 假设你想储存一个字母. 字母表有 26 个字母, 那么你需要几个比特呢? 如果用 4 个比特, 你可以表示 16 个值, 数量不够. 如果用 5 个比特, 你可以表示多达 32 个值, 这足够表示所有字母了, 而且还有剩余.

通常, 如果你想表示 N 个值, 你应该选择满足 $2^b \geq N$ 的最小 b 值. 两侧以 2 为底取对数, 便可以得到 $b \geq \log_2 N$.

假设我掷一枚硬币, 并告诉你结果, 那么我便是给了你一比特的信息. 如果我掷一个六面骰子, 并告诉你结果, 那我便是给了你 $\log_2 6$ 个比特的信息. 通常, 如果发生的概率是 N 中的 1 个, 那么结果便包含 $\log_2 N$ 个比特信息.

同样的, 如果结果出现的概率是 p , 那么信息量便是 $-\log_2 p$. 这个数量也被称为结果的**自信息**, 它是用来衡量结果的惊奇程度, 所以也被称为**惊奇度**. 如果你的马仅有 $1/16$ 的概率获胜, 但是赢了, 你可以 4 比特的信息 (以及赔付). 但如果获胜的机率是 75%, 那么胜利的消息只包含 0.42 比特.

直观来说, 出乎意料的消息反而包含了大量的信息. 相反, 如果你已经确定某件事情, 那么证实它仅会提供少量信息.

在本书的诸多主题中, 我们需要在比特数, b , 和其能编码的值的数量 $N = 2^b$ 之间熟练转换.

3.2 内存和存储

进程运行时，其大部分数据存储在**主内存**，这通常是一种随机访问存储 (RAM)。当前多数计算机上，主存储都是**易失的**，也就意味着一旦计算机关机，主存储的内容便会丢失。一台典型的桌面计算机有 2–8 GiB 内存。GiB 表示 “gibibyte,” 也就是 2^{30} 字节。

进程需要读写文件，而文件往往存储在硬盘驱动器 (HDD) 或者固态硬盘中。这些存储设备都是**非易失的**，因此可以用作长期存储。目前一台典型桌面计算机会有 500GB 到 2TB 容量的硬盘。其中 GB 表示 “gigabyte,” 也就是 10^9 字节。TB 表示 “terabyte,” 也就是 10^{12} 字节。

你可能已经注意到，我描述主内存大小时，用二进制单位 GiB，描述硬盘大小时，用十进制单位 GB 和 TB。由于一些历史和科技的原因，内存通常用二进制单位衡量，而硬盘则用十进制单位。在本书中，我将谨慎区分二进制和十进制单位，但你也应该注意到，“gigabyte” 和缩写 GB 通常是混用的。

在非正式的使用中，“内存” 一词有时候会指 HDD 和 SSD，以及 RAM。但这些设备的属性是不同的，所以我们需要仔细区分它们。我将使用**存储**来指代 HDD 和 SSD。

3.3 地址空间

主内存中的每个字节都由一个整数**物理地址**指定。有效的物理地址集被称为**物理地址空间**。通常从 0 到 $N - 1$ ，其中 N 是内存的大小。在 1 GiB 物理内存的系统上，最高的有效地址是 $2^{30} - 1$ ，即十进制下的 1,073,741,823，或者十六进制下的 0x3fff ffff (前缀的 0x 表示十六进制数)。

然而，大部分的操作系统会提供**虚拟内存**，这意味着程序永远不必处理物理地址，也不必知道由多少物理地址可用。

相反，程序会使用从 0 到 $M - 1$ 编号的**虚拟地址**，其中 M 是有效虚拟地址的数量。虚拟地址空间的大小是由操作系统和其运行的硬件共同决定的。

你可能听别人谈论 32 位系统和 64 位系统。这些术语表示的是寄存器的大小，通常也是虚拟地址的大小。在 32 位系统上，虚拟地址为 32 位，也就是虚拟地址空间从 0 到 0xffff ffff。所以地址空间的大小为 2^{32} 字节，或 4 GiB。

而在 64 位系统上，虚拟地址空间是 2^{64} 字节，或 $2^4 \cdot 1024^6$ 字节。这是 16 exbibytes，约为当前物理内存的十亿倍。虚拟地址空间比物理内存大这么多，这看起来很奇怪，但是我们稍后就会看到它是如何工作的。

当程序从内存读取或者向内存写入值时，它会生成虚拟地址。在操作系统的帮助下，访问内存之前，硬件会被转换成物理地址。这种转换是基于每个进程

进行的, 所以即使两个进程生成同样的虚拟地址, 它们也会映射到物理内存的不同位置.

最后, 虚拟内存是操作系统隔离进程的一种重要手段. 通常, 一个进程是无法访问另一个进程拥有的数据的, 因为如果某块物理内存已经被分配给了其他进程, 那么该进程便无法生成映射到这块物理内存的虚拟地址.

3.4 内存段

运行中得进程的数据会被组织成五个段:

- **代码段** 包含程序文本; 也就是构成程序的机器语言指令.
- **静态段** 包含不可变的值, 像字符串常量. 比如, 如果你的程序包含字符串 "Hello, World", 那这些字符会被保存在静态段.
- **全局段** 包含全局变量和被声明为 **static(静态类型)** 的局部变量.
- **堆段** 包含运行时被分配的内存块, 通常是通过调用 C 库函数 **malloc** 所分配.
- **栈段** 包含调用栈, 也就是一系列的堆栈帧. 每次函数被调用, 都会为其分配一个包含函数的参数和局部变量的堆栈帧. 当函数运行结束, 其堆栈帧会从栈中移除.

段的排序部分由编译器决定, 部分由操作系统决定. 具体细节取决于操作系统, 但常见的排列方式包括::

- 文本段位于内存接近“底部”的位置, 也就是靠近地址 0 的地方.
- 静态段通常在文本段上方, 也就是高地址位置.
- 全局段通常位于静态段上方.
- 堆通常在全局段上方. 随着堆的扩张, 它会向更大地址增长.
- 栈靠近内存的顶部; 也就是虚拟地址空间中最高地址附近. 栈扩展时, 会朝着更小地址方向增长.

若要确定你系统上这些段的布局, 尝试运行这个程序, 在本书仓库中的 `aspace.c` 文件中 (详见第 0.1 节).

```
#include <stdio.h>
#include <stdlib.h>

int global;

int main ()
{
    int local = 5;
    void *p = malloc(128);
    char *s = "Hello, World";

    printf ("Address of main is %p\n", main);
    printf ("Address of global is %p\n", &global);
    printf ("Address of local is %p\n", &local);
    printf ("p points to %p\n", p);
    printf ("s points to %p\n", s);
}
```

`main` 是这个函数的名称; 当作为变量使用时, 它指的是 `main` 中的第一条机器语言指令的地址, 同时我们希望它在文本段中.

`global` 是全局变量, 所以我们希望它在全局段中. `local` 是局部变量, 我们期望它存在于栈中.

`s` 指向“字符串文本”, 它是程序的一部分 (与从文件读取, 或用户输入的字符串等不同). 我们希望字符串的位置在静态段中 (与指针 `s` 不同, 它是局部变量).

`p` 包含 `malloc` 返回的地址, 这个函数会在堆中分配空间. “`malloc`”表示“memory allocate(内存分配)”.

格式序列`%p` 会告诉 `printf` 将每个地址格式化为“指针”, 因此它会以十六进制显示结果.

当我运行这个程序, 输出如下 (我增加了些空格, 以便于阅读):

```
Address of main is   0x      40057d
Address of global is 0x      60104c
Address of local is  0x7ffe6085443c
p points to          0x      16c3010
s points to          0x      4006a4
```

和预计的一样, `main` 的地址是最低的, 然后是字符串文本的位置. 接着是 `global` 的位置, 再然后是 `p` 指向的地址. `local` 的地址要大得多.

最大的地址有 12 位十六进制数. 每个十六进制数对应 4 个比特, 所以一共是 48 比特地址. 这便表示可用虚拟地址空间的大小为 2^{48} 字节.

做个练习, 在你电脑上运行这个程序, 并和我的结果进行比较. 再添加一个 `malloc`, 检查你系统上的堆是否向上增长了 (向更大地址方向). 再添加一个输出局部变量地址的函数, 并检查栈是否向下变化了.

3.5 静态局部变量

栈上的局部变量有时被称为**自动变量**, 因为它们会在函数调用时, 自动分配, 在函数返回时, 自动释放.

在 C 语言中, 存在另一种局部变量, 会在全局段中进行分配, 称为**静态变量**, 它会在程序启动时被初始化, 同时在函数调用之间保持其值不变.

例如, 下面的函数会跟踪其被调用的次数.

```
int times_called()
{
    static int counter = 0;
    counter++;
    return counter;
}
```

关键字 `static` 表示 `counter` 是一个静态局部变量. 初始化仅会在程序启动时, 发生一次.

如果你将此函数添加到 `aspace.c` 中, 你便可以确认 `counter` 是在全局段中与全局变量一起分配的, 而不是在栈中分配的.

3.6 地址转换

虚拟地址 (VA) 如何能转换成物理地址 (PA)? 基本机制很简单, 但简单的实现会太慢, 并消耗太多空间. 所以实际的实现有点复杂.

大部分处理器会提供位于 CPU 和主存之间的内存管理单元 (MMU). MMU 会执行 VA 和 PA 之间的快速转换.

1. 当程序读取或写入一个变量时, CPU 会生成一个 VA.
2. MMU 将 VA 分成两部分, 分别是页号和偏移量. “页”是内存的一个块, 页的大小取决于操作系统和硬件, 常见大小一般是 1-4 KiB.



图 3.1: 地址转换过程示意图.

3. MMU 会从页表缓存 (TLB) 查找页号, 并获得相应的物理地址页号. 然后, 它将物理页号和偏移量结合起来, 构成 PA.
4. PA 会传递给主存, 主存基于给定位置进行读写.

TLB 包含来自页表 (存储在内核内存中) 数据的缓存副本. 页表包含虚拟页号到物理页号的映射. 由于每个进程都有自己的页表, 因此 TLB 必需确保它只使用正在运行的进程的页表中的条目.

图 3.1 是这个过程的示意图. 若要明了其工作原理, 假设 VA 是 32 位, 物理内存是 1 GiB, 并划分为了 1 KiB 的页.

- 因为 1 GiB 是 2^{30} 字节, 1 KiB 是 2^{10} 字节, 所以有 2^{20} 个物理页, 有时也称为“帧”.
- 虚拟地址空间大小是 2^{32} B, 单个页的大小是 2^{10} B, 所以有 2^{22} 个虚拟页.
- 偏移量的大小由页面大小决定. 这个例子中, 页面大小是 2^{10} B, 所以需要 10 位来定位页面上的字节.
- 如果 VA 是 32 位, 同时偏移量是 10 位, 剩下的 22 位便构成了虚拟页号.
- 由于有 2^{20} 个物理页, 每个物理页号是 20 位. 加上 10 位的偏移量, 得到的 PA 便是 30 位.

至此, 一切似乎看起来可行. 但让我们想一下, 页表可能有多大. 页表的最简单实现是一个数组, 其中每个条目对应一个虚拟页. 每个条目包含物理页号, 本例中是 20 位, 以及关于每帧的附加信息. 所以我们预计每个条目需要 3-4 字节的空间. 由于存在 2^{22} 个虚拟页, 页表将需要 2^{24} 字节, 即 16 MiB.

由于每个进程都需要一个页表, 那么运行 256 个进程的系统, 便需要 2^{32} 字节, 即 4GiB, 还只是页表空间! 而这仅仅是针对 32 位虚拟地址的情况. 对于 48-或 64-位的虚拟地址来说, 这个数字就太吓人了.

幸运的是, 我们实际上并不需要这么多空间, 因为多数进程甚至用不到虚拟地址空间的一小部分. 而且, 如果一个进程不使用虚拟页, 我们便不需要在页面中为其分配条目.

另一种表达同样观点的方式是页表是“稀疏的”, 而这最简单的实现方式, 即一个页表条目数组, 但这是一个糟糕的注意. 幸运的是, 对于稀疏数组, 有其他高效的实现方案.

一种选择是多级页表, 这是许多操作系统 (包括 Linux) 使用的方案. 另一种选择是关联表, 其中每个条目包括虚拟页号和物理页号. 在软件中检索关联表很慢, 但在硬件中, 我们可以并行检索整个表, 所以其通常用于表示 TLB 中的页表条目.

你也可以从http://en.wikipedia.org/wiki/Page_table了解更多关于这些实现的信息; 你可能对这些细节感兴趣. 但核心思想为页表是稀疏的, 所以我们需要为稀疏数组选择一个好的实现方案.

前面我提到操作系统可以中断运行中的进程, 保存其状态, 然后运行另一个进程. 这个机制叫做**上下文切换**. 因为每个进程都有自己的页表, 操作系统需要和 MMU 协作, 从而确保每个进程都能获得正确的页表. 在老的机器上, MMU 中的页表信息需要在每次上下文切换时进行替换, 而这是很昂贵的. 在新的系统中, MMU 中的每个页表条目都包含进程 ID, 因此多个进程的页表可以同时存在于 MMU 中.

Chapter 4

文件和文件系统

当进程结束 (或崩溃), 存储在主内存中的任何数据都会丢失. 但存储在硬盘驱动器 (HDD) 或者固态硬盘驱动器上的数据是“持久的”; 也就是说进程结束, 甚至计算机关闭, 也会存在.

硬盘驱动器是复杂的, 数据以块的形式存储, 这些块存在于扇区上, 扇区构成磁道, 磁道则排列在盘片的同心圆上.

固态硬盘在某种意义上相对简单, 因为块会按照顺序编号, 但它们也因此引入了一个不一样的复杂性: 每个块只能写入有限的次数, 然后就会变得不可靠.

作为个程序员, 你不会想要处理这些复杂性. 你想要的是对持久存储硬件的合理抽象. 而最常见的抽象就是“文件系统”.

抽象地说:

- “文件系统”是将每个文件名称映射到其内容的一种方式. 如果你将名称视为键, 内容当作值, 文件系统便是一种键-值数据库. (参见 https://en.wikipedia.org/wiki/Key-value_database).
- 一个“文件”便是一个字节序列.

文件名通常为字符串, 而且一般是“分层次的”; 也就是说, 该字符串指定了一个从顶级目录 (或文件夹) 到特定文件的一系列子目录的路径.

抽象层以及底层机制之间最主要的区别是, 文件是基于字节的, 而持久层存储是基于块的. 操作系统将 C 库中基于字节的文件操作转换为对存储设备上基于块的操作. 典型的块大小一般为 1-8 KiB.

例如, 下面代码会打开文件并读取第一个字节:

```
FILE *fp = fopen("/home/downey/file.txt", "r");  
char c = fgetc(fp);  
fclose(fp);
```

上述代码运行时:

1. `fopen` 会使用文件名中的/ 确定顶级目录, 以及子目录 `home`, 以及下一级子目录 `downey`.
2. 它会找到名为 `file.txt` 的文件, 然后“打开 (opens)”文件以进行读取, 也就表示, 这会创建一个表示正在被读取的文件的数据结构. 这个数据结构除了其他事情, 还会跟踪文件已读取了多少, 也就是“文件位置”.

在 DOS 中, 这个数据结构叫做文件控制块 (FCB), 但我希望尽量避免使用这个术语, 因为在 UNIX 中, 它有其他含义. 在 UNIX 中, 似乎没有一个相应的好名称. 它是打开文件表中的一个条目, 所以我称其为打开文件表条目.

3. 当我调用 `fgetc` 时, 操作系统会检查文件的下一个字符是否已经存在于内存中. 如果已存在, 则会读取下一个字符, 更新文件位置, 并返回结果.
4. 如果下一个字符不在内存中, 操作系统会发出一个 I/O 请求, 获取下一个块. 磁盘驱动器速度较慢, 所以一个等待从磁盘获取块的进程通常会被中断, 以便另一个进程运行, 直到数据返回, 再回到等待的进程.
5. 当 I/O 操作结束, 新的数据块会被存储到内存, 同时进程继续运行. 读取第一个字符, 并将其存储为局部变量.
6. 当进程关闭文件, 操作系统会完成或者取消所有挂起的操作, 移除存储在内存中的数据, 并释放打开文件表条目 (`OpenFileTableEntry`).

写文件的过程类似, 只是会多一些额外步骤. 下面是一个打开文件进行写入并修改第一个字符的示例.

```
FILE *fp = fopen("/home/downey/file.txt", "w");  
fputc('b', fp);  
fclose(fp);
```

当代码运行时:

1. 再次, `fopen` 使用文件名定位文件. 如果文件不存在, 则会创建一个新文件, 并在父目录 `/home/downey` 中添加一个条目.
2. 操作系统会创建一个打开文件表条目, 以标明此文件以写入方式打开, 同时设置文件位置为 0.

3. `fputc` 会尝试写入 (或重写) 文件的第一个字节. 如果文件已经存在, 操作系统需要将第一个块加载入内存. 否则, 它会在内存中分配一个新块, 同时在磁盘上申请一个新块.
4. 内存中的块被修改后不会立刻被复制回磁盘. 通常, 写入文件的数据是“缓冲的”, 意味着它会保存在内存, 直到至少有一个块需要写入, 才会写入磁盘.
5. 当文件关闭, 任何缓冲区的数据都会被写入磁盘, 同时释放打开文件表条目.

综上所述, C 库提供了一个从文件名映射到字节流的文件系统的抽象. 这个抽象建立在实际以块进行组织的存储设备之上.

4.1 磁盘性能

我之前提过磁盘驱动器很慢. 对于当前的 HDD 来说, 从磁盘读取一个块到内存的平均时间大约 5–25 毫秒 (参见 https://en.wikipedia.org/wiki/Hard_disk_drive_performance_characteristics). 而 SSD 则快很多, 读一个 4 KiB 的块, 只需要 25 μ s, 写一个块需要 250 μ s (参见 <http://en.wikipedia.org/wiki/Ssd#Controller>).

为了更好地理解这些数字, 我们将其和 CPU 的时钟周期进行比较. 时钟频率为 2GHz 的处理器每 0.5 ns 完成一个时钟周期. 从内存获取一个字节到 CPU 的时间通常约为 100 ns. 如果处理器每个时钟周期完成一条指令, 那么在等待从内存获取一个字节的的时间可以完成 200 条指令.

在 1 微秒内, 处理器可以完成 2000 条指令, 所以等待从 SSD 获取一字节的 25 μ s 内, 可以完成 50,000 条指令.

在 1 毫秒内, 处理器可以完成 2,000,000 条指令, 所以等待从 HDD 获取一字节的 20 ms 内, 处理器可以完成 4 千万条指令. 如果在等待期间, CPU 没有其他任务执行, 则它会处于空闲状态. 这也是为什么操作系统通常在等待磁盘数据时切换到另一个进程的原因.

主存和持久化存储之间的性能差距是计算机系统设计的主要挑战之一. 操作系统和硬件提供了几种特性, 以“填补”这个差距:

- 块传输: 从磁盘加载单个字节的时间是 5–25 ms. 相比之下, 加载 8 KiB 块的额外时间可以忽略不急. 所以系统通常每次访问磁盘时, 尽量读取大块数据.

- 预取：有时候操作系统可以预测一个进程将会读取的块，从而在请求之前便进行加载。比如，如果你打开一个文件并读取第一个块，那么你很可能会继续读取第二个块。操作系统可能会在被请求之前，先加载其他块。
- 缓冲区：正如我之前提到的，当你写入一个文件，操作系统会将数据存储在内存中，稍后再将其写入磁盘。如果当其在内存中时，你多次修改了块，系统只需要往磁盘写入一次。
- 缓存：如果进程最近使用了某个块，很可能再次使用它。如果操作系统在内存中保留了该块的副本，那便可以以内存速度处理将来的请求。

这些功能的有些实现是在硬件中。比如，有些磁盘驱动器会提供一个缓存，来存储最近使用过的块，同时很多磁盘驱动器即使只被请求了一个块，也会一次夺取多个块。

这些机制通常会提高程序的性能，但不会改变其行为。通常程序员不需要考虑它们，除非出现以下两种情况：(1) 如果程序性能很糟糕，你可能需要了解这些机制以便诊断问题，以及 (2) 当数据被缓冲存储了，程序调试会变得困难。例如，如果一个程序打印一个值，然后崩溃。该值可能不会出现，因为它可能在缓冲区。同样，如果一个程序往磁盘写数据，然后计算机断电了，如果数据还在缓存中，尚未写入磁盘，那么数据很可能会丢失。

4.2 磁盘元数据

构成文件的块在磁盘上可能是连续排列，如此的话，文件系统的性能一般会更好，但多数操作系统不强求连续分配。它们可以在磁盘的任意位置自由放置数据块，并使用各种数据结构对其进行跟踪。

在很多 UNIX 文件系统中，该数据结构被称为“inode”，即“index node(索引节点)”。通常，关于文件的信息，包括其数据块的位置，被称为“元数据”。(文件的内容是数据，所以关于文件的信息便是关于数据的数据，因此称为“元数据”。)

由于 inode 一般与其他数据一起存储在磁盘上，它们被设计成了完美适应磁盘块。UNIX 的 inode 包含文件的信息，包括文件所有者的用户 ID；表示谁可以读写或执行该文件的权限标注；以及指示文件上次修改和访问的时间戳。此外，它还包含组成文件的前 12 个数据块编号。

如果块大小是 8 KiB，那么前 12 个块总共占用 96 KiB 空间。多数系统中，这对于绝大多数文件来说已经足够大，但仍不够容纳所有文件。这也是为何 inode 会包含指向“间接块”的指针的原因，该间接块仅包含指向其他块的指针。

间接块中的指针数量取决于块大小和块编号, 但通常是 1024. 对于具有 1024 个块编号, 并且块大小为 8 KiB 时, 那么一个间接块可以寻址 8 MiB 的空间. 这对于大多数大文件来说, 已经足够大, 但仍然不够容纳所有文件.

这也是为何 inode 也包含指向“双重间接块”的指针的缘故. 双重间接块包含指向间接块的指针. 通过 1024 个间接块, 我们便能寻址 8GiB 空间.

同时如果还是不够大, (最后) 还有三重间接块, 也就是包含指向双重间接块指针的块, 从而使文件大小达到 8 TiB. 当 UNIX 的 inode 设计时, 这似乎在很长时间内都足够大了. 但那是很久之前了.

有些文件系统, 比如 FAT, 会使用一个被称为“簇”的文件分配表作为间接块的替代方案, 这个表包含每个块的一个条目. 根目录包含每个文件中第一个簇的指针. 每个簇的 FAT 条目都会指向文件中的下一个簇, 和链表类似. 更多详细信息, 请参考 http://en.wikipedia.org/wiki/File_Allocation_Table.

4.3 块分配

文件系统需要跟踪哪些块属于某个文件; 也需要跟踪哪些块可以被使用. 当创建新文件时, 文件系统会找到一个可用的块并进行分配. 当删除了文件, 文件系统会将块标记为可用, 以重新分配.

块分配系统的目标是:

- 速度: 分配和释放块要快速.
- 最小化空间开销: 分配器使用的数据结构要尽量小, 为数据保留尽量多的空间.
- 最小化碎片: 如果有些块不再使用, 或者有些仅部分使用, 那么未使用的空间便是“碎片”.
- 最大化连续性: 如果可以, 一起使用的数据最好在物理上是连续的, 以提高性能.

很难设计出一个满足上述所有目标的文件系统, 特别是文件系统的性能取决于诸如文件大小, 访问模式等的“负载特征”. 对于个工作负载优化过的文件系统, 并不一定适用于另一个.

因此, 大部分操作系统支持多种文件系统, 同时文件系统设计也是一个研究和开发的活跃领域. 过去十年, Linux 系统已经从传统的 UNIX 文件系统 ext2, 迁移到了 ext3, ext3 是一个“日志”文件系统, 旨在提高速度和连续性, 最近又迁移到了 ext4, 这是一个可以处理更大的文件的文件系统. 在可预计的将来, 可能又会迁移到 B-树文件系统, Btrfs.

4.4 一切皆文件？

文件抽象实际是一个“字节流”抽象，这种抽象对于很多事情都有用，而不仅仅是文件系统。

一个例子便是 UNIX 管道，这是一个简单的进程间通信方式。进程可以设置为将一个进程的输出作为另一个进程的输入。对于第一个进程来说，这个管道就像一个以写入方式打开的文件，所以它可以使用 C 库函数，如 `fputs` 和 `fprintf`。对于第二个进程来说，管道就像一个以读取方式打开的文件，所以它可以使用 `fgets` 和 `fscanf`。

网络通信也使用字节流抽象。UNIX 套接字便是一个表示不同计算机的进程间通信通道的数据结构。同样，进程可以使用“文件”处理函数，从套接字读取以及写入数据。

复用文件抽象令程序员的工作变得简单，因为他们只需要学习一个 API(应用程序编程接口)。同时也使程序更加灵活，因为原本设计用于文件的程序，也可以处理来自管道和其他来源的数据。

Chapter 5

位和字节

5.1 整数表示

你可能知道计算机是用二进制表示数字。对于正数来说，二进制表示很简单；例如，十进制的 5_{10} 可以用二进制表示为 $b101$ 。

对于负数来说，最直观表示是采用一个符号位来标识数字的正负。但也有其他表示方法，称为“补码”，因为其在硬件处理中更容易，所以更加常用。

若要找到一个负数 $-x$ 的补码，需要先找到 x 的二进制表示，然后翻转所有未，再加 1。比如，若要表示 -5_{10} ，可以从表示 5_{10} 开始，如果写成 8 位版本，即 $b00000101$ 。将所有位翻转并加 1，得到 $b11111011$ 。

在二进制补码中，最左侧的位充当符号位；正数的符号位是 0，负数的符号位为 1。

若要将 8 位数字转为 16 位，我们需要为正数添加更多的 0，为负数添加更多的 1。实际上，我们需要将符号位复制到新的位上。这个过程叫做“符号扩展”。

在 C 语言中，所有的整数类型都是有符号位的（能够表示正数和负数），除非声明其为无符号位。这个声明的区别以及重要性在于，无符号位的整数不需要符号扩展。

5.2 位运算符

人们学习 C 时，有时会对位运算符 `&` 和 `|` 满腹疑惑。这些运算符将整数视为位向量，并对相应的位进行逻辑运算。

比如，`&` 进行 AND(且) 运算，这会在两个操作数都是 1 时，结果为 1，否则为 0。下面是对两个 4-位数字应用 `&` 运算的例子：

```

  1100
& 1010
----
  1000

```

在 C 中, 这表示表达式 `12 & 10` 的结果为 8.

同样的, `|` 会进行 OR(或) 运算, 在两个操作数的其中一个为 1 时, 结果为 1, 否则为 0.

```

  1100
| 1010
----
  1110

```

所以表达式 `12 | 10` 的结果为 14.

最后, 运算符 `^` 表示 XOR(异或) 操作, 在两个操作数有且只有一个为 1 时, 结果为 1.

```

  1100
^ 1010
----
  0110

```

所以表达式 `12 ^ 10` 的结果为 6.

最常见的用法, `&` 用来从位向量中清除一组位, `|` 用来设置位, 而 `^` 用来翻转或“切换”位. 下面是详细信息:

清除位: 对于任意 x , $x \& 0$ 结果是 0, 而 $x \& 1$ 的结果是 x . 所以如果你用 3 对一个向量执行 AND 操作, 则只会选择最右边的两位, 并将其他位置设置为 0.

```

  xxxx
& 0011
----
  00xx

```

在这里, 3 叫做“掩码”, 因为它选择一些位, 并掩盖其余的位.

设置位: 同样, 对于任意 x , $x | 0$ 结果是 x , 而 $x | 1$ 结果是 1. 所以如果你用 3 对一个向量执行 OR 运算, 则会设置最右边的位, 并保持其余位不变:

```

  xxxx
| 0011
----
  xx11

```

翻转位: 最后, 如果你用 3 对一个向量执行 XOR 操作, 它会翻转最右边的位, 同时保持其余位不变. 作为练习, 看看你是否可以使用 `^` 计算出 12 的二进制补码. 提示: -1 的二进制补码表现形式是什么?

C 语言也提供了移位操作 `<<` 和 `>>`, 用于左移位和右移位. 每次左移会将数字翻倍, 所以 `5 << 1` 结果是 10, `5 << 2` 等于 20. 每次右移将会令数字除 2(向下取整), 所以 `5 >> 1` 的结果为 2, 同时 `2 >> 1` 结果为 1.

5.3 浮点数表示

浮点数使用二进制科学计数法表示. 在十进制表示法中, 大数字会被写成系数与 10 的指数幂相乘的形式. 例如, 光速用 m/s 表示, 大约为 $2.998 \cdot 10^8$.

多数计算机使用 IEEE 标准进行浮点数运算. C 语言中的 `float` 类型通常对应 32 位的 IEEE 标准; `double` 通常对应 64 位标准.

在 32 位标准中, 最左侧的位是符号位 s . 其后 8 位是指数 q , 最后的 23 位是系数 c . 那么一个浮点数的值便可以表示为:

$$(-1)^s c \cdot 2^q$$

基本正确, 但是小有差异. 浮点数通常是规范化的, 以确保小数点前有一个数字. 例如, 在十进制中, 我们一般写成 $2.998 \cdot 10^8$, 而不是 $2998 \cdot 10^5$, 或者其他等价表达式. 在二进制中, 规范的数字总是在二进制点前有一个数字. 由于该位置的数字始终是 1, 我们便可以在表示中省略它, 从而节省空间.

比如, 13_{10} 的整数表示是 `b1101`. 浮点数表示便是 $1.101 \cdot 2^3$, 指数是 3, 系数部分存储为 101(后跟 20 个零).

这基本正确, 但存在一个小细节. 指数是用“偏移”存储的. 在 32-位标准中, 偏移是 127, 所以指数 3 会被存储为 130.

在 C 中打包和解包浮点数, 我们可以用联合位运算操作. 这是一个例子:

```
union {
    float f;
    unsigned int u;
} p;

p.f = -13.0;
unsigned int sign = (p.u >> 31) & 1;
unsigned int exp = (p.u >> 23) & 0xff;

unsigned int coef_mask = (1 << 23) - 1;
unsigned int coef = p.u & coef_mask;

printf("%d\n", sign);
```

```
printf("%d\n", exp);  
printf("0x%x\n", coef);
```

此代码存在于本书配套仓库中的 `float.c` 文件中 (见章节 0.1).

联合运算符允许我们使用 `p.f` 存储浮点数, 然后用 `p.u` 将其读取为无符号整数.

要获取符号位, 我们需要向右移 31 位, 然后用 1 位掩码来选择最右边的位.

若要获取指数, 我们右移 23 位, 然后选择最右侧的 8 位 (十六进制 `0xff` 有 8 个 1).

要获取系数, 我们需要提取最右侧 23 位并忽略其余位. 我们通过创建一个右侧 23 位均为 1, 左侧为 0 的掩码来实现. 最容易的方法是将 1 左移 23 位, 然后减去 1.

程序输出为:

```
1  
130  
0x500000
```

如预期一样, 负数的符号位是 1. 指数是 130, 包括偏差. 系数用十六进制表示, 是 101, 后跟 20 个零.

做个练习, 尝试聚合或者拆散用 64 位表示的 `double`, 请参阅http://en.wikipedia.org/wiki/IEEE_floating_point.

5.4 联合体和内存错误

C 联合体有两种常见用途. 一种是前面章节看到的, 用于访问数据的二进制表示. 另一种是用于存储异构数据. 比如, 你可以使用一个联合体来表示一个可能是整数, 浮点数, 复数或有理数的数字.

然而, 联合体是存在错误风险的. 这往往取决于你, 作为程序员, 往往需要跟踪联合体中数据类型; 如果你写了一个浮点数, 然后将其解释为了整数, 得到的结果便毫无意义了.

实际上, 如果你错误读取了内存中的某个位置, 也会发生同样的事情. 导致这种问题发生的一种方式是你读取了超出数组末尾的位置.

若要看到会发生什么, 我先创建一个在堆栈上分配数组, 并用 0 到 99 的数字进行填充的函数.

```
void f1() {
    int i;
    int array[100];

    for (i=0; i<100; i++) {
        array[i] = i;
    }
}
```

下一步我会定义一个函数, 以创建一个更小的数组, 并有意访问开始之前和结束之后的元素.

```
void f2() {
    int x = 17;
    int array[10];
    int y = 123;

    printf("%d\n", array[-2]);
    printf("%d\n", array[-1]);
    printf("%d\n", array[10]);
    printf("%d\n", array[11]);
}
```

如果我先调用 `f1`, 然后调用 `f2`, 我会得到下面结果:

```
17
123
98
99
```

这里细节取决于编译器, 它会在堆栈安排变量的位置. 从这些结果, 我们可以推断出编译器将 `x` 和 `y` 放在了一起, 位于数组的“下方” (一个低地址的位置). 当我们读取超出数组的元素时, 似乎我们得到了前一个函数调用而留在堆栈上的值.

在这个例子中, 所有的变量都是整数, 所以相对容易弄清楚发生了什么. 但是通常当你读取超出数组边界的位置时, 你读到的值可能是任何类型. 例如, 如果我修改 `f1` 函数, 创建一个浮点数数组, 结果如下:

```
17
123
1120141312
1120272384
```

后两个值是你将浮点数作为整数解释得到的值. 如果你调试时遇到这种输出, 你很难确定发生了什么.

5.5 字符串表示

有时和字符串相关的问题也会出现. 首先, 记住 C 语言中的字符串是以 null 结尾的. 当你为字符串分配空间时, 不要忘记在末尾多留一个字节.

另外, C 语言中字符串内的字母以及数字是用 ASCII 编码的. 数字 “0” 到 “9” 的 ASCII 编码是 48 到 57, 而不是 0 到 9. ASCII 码中的 0 是 NULL 字符, 用来标识字符串的结束. ASCII 码中的 1 到 9 是用于某些通讯协议的特殊字符. ASCII 码中的 7 是一个响铃符; 在某些终端上, 打印它会发出声音.

字母 “’A” 的 ASCII 码是 65; “a” 的编码是 97. 下面是这些编码的二进制表示:

65 = b0100 0001

97 = b0110 0001

仔细观察者会注意到他们只有一个位存在差异. 其他字母也是这种模式; 第六位 (从右边数) 充当的是 “大小写位”, 大写字母是 0, 小写字母是 1.

做个练习, 写个函数, 使其接收一个字符串, 通过反转第六位, 将其从小写转成大写. 作为挑战, 你可以通过一次性读取 32 位或 64 位的字符串, 而不是逐个字符读取, 来创建更快的版本. 如果字符串长度是 4 或 8 字节的倍数, 优化过程会更容易实现.

如果你读取超出了字符串的末尾, 你很可能会看到奇怪的字符. 相反, 如果你写入一个字符串, 然后意外将其当作整数或者浮点数读取, 结果会很难解释.

例如, 如果你运行:

```
char array[] = "allen";  
float *p = array;  
printf("%f\n", *p);
```

你会发现将我名字的前 8 个字符的 ASCII 表示, 解释为双精度的浮点数字, 结果是 69779713878800585457664.

Chapter 6

内存管理

C 语言为动态内存分配提供了 4 个函数:

- `malloc` 接收一个以字节为单位的整数, 返回一个 (至少) 指定大小的新分配内存块. 如果无法满足此条件, 则返回特殊指针 `NULL`.
- `calloc` 和 `malloc` 一样, 不同之处在于它会清除新分配的块, 也就是说, 它会将块中的所有字节设置为 0.
- `free` 会接受先前分配块的指针, 并释放它; 也就是说, 它可以将空间用于未来分配.
- `realloc` 会接收一个之前分配块的指针和一个新的大小. 它会根据新的大小分配内存块, 并将数据从旧的块拷贝到新的块, 释放旧的块, 返回新的块的指针.

这个 API 是出了名的容易出错且无法容忍. 内存管理一直是设计大型软件系统所面临的最大挑战之一, 而这也是为什么大多数现代语言提供类似垃圾回收的高级内存管理的特性的原因.

6.1 内存错误

C 语言的内存管理 API 和动画节目 *The Simpsons* 中的一个小人物 Jasper Beardly 有点像; 在少数的几集中, 他作为一个严厉的代课老师而出现, 对所有的违规行为都施加体罚——“打屁股”.

下面是程序会做的一些需要被打屁股的事情:

- 访问 (读写) 任意未分配的块.

- 访问已释放的分配内存.
- 尝试释放未分配的内存.
- 重复释放同一块内存.
- 对未分配的内存, 或者分配后已释放的内存调用 `realloc` 函数.

遵循这些规则看起来不难, 但在大型程序中, 内存块会被某段程序分配, 又被不同部分代码使用, 再被其他部分释放. 所以, 程序中一部分的修改, 往往涉及其他很多地方.

同时, 程序的不同部分可能有很多对同一块内存的别称, 或者引用. 直到所有对某个内存块的引用都不再使用, 内存才能释放. 正确处理这点, 往往需要在程序所有地方都仔细分析, 这很困难, 而这也违反了良好软件工程的基本原则.

理想情况是, 每个分配内存的函数, 都应该在文档化的接口中包含如何释放该内存的信息. 成熟的库通常在这点做的很好, 但现实世界中, 软件工程实践往往和理想状态差之甚远.

更糟糕的是, 内存错误因为各种状况难以预测, 往往很难定位. 比如:

- 如果你从一个未分配的内存块中读取值, 系统可能会检测到错误, 触发一个叫“分段错误”的运行时错误, 并终止程序. 或者, 程序也可能并不会报错, 直接读取未分配的内存; 这种情况下, 获取的值是访问位置存储的值, 而这是无法预测的, 程序每次执行结果可能都不同.
- 如果你向一个未分配内存块写入值, 也没有遇到分段错误, 事情可能更糟糕. 你向一个无效的位置写入一个值, 可能经过很长时间才会读取, 并引发问题. 而这时很难发现问题的源头.

事情可能更糟! C 风格的内存管理中一个最常见的问题是用于实现 `malloc` 和 `free` (我们很快会看到) 的数据结构通常和分配的内存块存储在一起. 因此, 如果你意外地在动态分配内存的末尾之后写入, 你很可能会破坏这些数据结构. 系统通常会在你调用 `malloc` 或 `free`, 同时这些函数莫名其妙失败时, 才会检测到问题.

从中, 你可以得到的结论是, 安全内存管理需要设计和纪律. 如果你编写了一个库或者模块来分配内存, 你也应该提供释放它的接口, 同时, 内存管理应该从开始便作为 API 设计的一部分.

如果你使用一个库分配内存, 你应该在使用 API 时遵循纪律. 比如, 如果库提供了分配和释放存储的函数, 你应该先分配而不是先释放, 避免在一个没有执行 `malloc` 的内存块上调用 `free` 函数.

通常需要在安全内存管理和性能之间进行权衡。比如，内存错误最常见的一个来源便是超出数组边界进行写入。解决这个问题最明显的方法是边界检查；即每次数组访问都应该检查索引是否超出边界。提供类似数组结构的高级库通常会执行边界检查。但 C 数组和大多数底层库不会执行此类检查。

6.2 内存泄漏

还有一种内存错误，可能需要注意，也可能不需要在意。如果你分配了一块内存，但从未释放，这便是“内存泄漏”。

对有些程序来说，内存泄漏可以接受。比如，你的程序分配内存，执行计算，然后退出，那么可能没有必要释放分配的内存。当程序退出时，全部内存都会被操作系统释放掉。在退出之前立刻释放内存可能感觉更加负责，但这基本上是浪费时间。

但如果程序运行很长时间且存在内存泄漏，它的总内存使用将无限增长。这时候，会发生以下情况：

- 某些时候，系统会用光物理内存。在没有虚拟内存的系统上，再调用 `malloc` 会失败，返回 `NULL`。
- 在具有虚拟内存的系统上，操作系统可以将另一个进程的页从内存移动到磁盘，然后为泄漏的进程分配其他空间。我在 7.8 节解释了这种机制。
- 单个进程可能存在分配空间的限额；超出限额，`malloc` 会返回 `NULL`。
- 最终，一个进程可能会填满其虚拟地址空间（或可用部分）。然后，便没有了跟多地址可供分配，`malloc` 则返回 `NULL`。

如果 `malloc` 返回 `NULL`，但是你依然访问你认为已经分配的内存块，那么会导致分段错误。因此，良好的编程风格应该是在使用 `malloc` 的结果之前先检查它。一个方法是在每个 `malloc` 调用后添加一个类似下面的条件：

```
void *p = malloc(size);
if (p == NULL) {
    perror("malloc failed");
    exit(-1);
}
```

`perror` 声明在 `stdio.h` 中；它会打印一个错误消息以及关于最后发生错误的附加信息。

`exit` 也在 `stdlib.h` 中声明，它会导致进程终止。参数是一个状态码，表示进程终止方式。按照惯例，状态码 0 表示正常终止，-1 表示错误条件。有时用不同代码表示不同错误条件。

错误代码检查可能令人厌烦, 而且令程序难以阅读. 你可以通过将库函数调用以及错误代码检查封装进自己的函数中, 来减轻这些问题. 比如, 下面是一个检查返回值的 `malloc` 封装函数.

```
void *check_malloc(int size)
{
    void *p = malloc (size);
    if (p == NULL) {
        perror("malloc failed");
        exit(-1);
    }
    return p;
}
```

因为内存管理如此之难, 大多数的大型程序, 比如 web 浏览器, 会泄漏内存. 想要查看系统上哪些程序使用了最多的内存, 你可以用 UNIX 程序 `ps` 和 `top`.

6.3 实现

当一个进程启动时, 系统会为文本段和静态分配的数据分配空间, 为堆栈分配空间, 以及包含动态分配数据的堆空间.

并非所有程序都是动态分配数据, 因此堆的初始大小可能很小, 甚至是零. 初始时, 堆仅包含一个空闲块.

当调用 `malloc` 时, 函数会检查是否可以找到足够大的空闲块. 如果无法找到, 便需要从系统申请更多内存. 执行此操作的函数是 `sbrk`, 它会设置“程序断点”, 你可以将其视为一个指向堆末尾的指针.

当 `sbrk` 被调用, 操作系统会分配新的物理内存的页, 更新进程的页表, 并设置程序断点.

理论上, 一个程序可以直接调用 `sbrk`(无需使用 `malloc`), 并自行管理堆. 但 `malloc` 更易于使用, 同时对于大多数内存使用模式来说, 它速度更快, 而且内存使用率高.

若要实现内存管理 API (即, 函数 `malloc`, `free`, `calloc`, 和 `realloc`), 大部分 Linux 系统使用 `ptmalloc`, 其源于 Doug Lea 写的 `dlmalloc`. 至于描述实现的关键要点的简短论文可以从<http://gee.cs.oswego.edu/dl/html/malloc.html>获取.

对于程序员来说, 需要关注的最重要的要素是:

- `malloc` 的运行时间通常并不取决于块的大小, 但可能取决于存在多少空闲块. 同时, 无论有多少空闲块, `free` 通常都很快. 而因为 `calloc` 需要清除块中的每个字节, 所以其运行时间取决于块大小 (也就是空闲块的数量).

`realloc` 有时很快, 比如要分配的空间小于当前块大小, 或者有可用空间来扩展当前块. 如果不是这样, 便需要将数据从旧块复制到新块; 这种情况下, 运行时间便取决于旧块的大小.

- 边界标签 (Boundary tags): 当 `malloc` 分配块时, 会在开始和结尾添加空间, 以存储关于块的信息, 包括其大小和状态 (占用或空闲). 这些数据位被称为 “边界标签”. `malloc` 通过使用这些标签, 可以轻易从任何一个块, 获取内存中其前后相邻的块. 此外, 空闲块被链接为一个双向链表; 每个空闲块包含在 “空闲列表” 中指向下一个和前一个块的指针. 边界标签和空闲列表指针便组成了 `malloc` 的内部数据结构. 这些数据结构和程序数据交错在一起, 所以程序错误很容易便会破坏他们.
- 空间开销 (Space overhead): 边界标签和空闲列表指针会占用空间. 大部分系统的最小块大小是 16 字节. 因此对于非常小的块, `malloc` 无法高效利用空间. 如果你的程序需要大量小结构, 可能在数组中分配它们比较高效.
- 碎片化 (Fragmentation): 如果你以各种大小分配和释放块, 则堆很可能会变得碎片化. 也就是, 空闲块很可能会被拆分成很多小碎片. 碎片化会浪费空间; 同时也会通过降低内存缓存的效率, 而减慢程序运行速度.
- 分箱和缓存 (Binning and caching): 空闲列表会按照大小分箱, 所以当 `malloc` 检索特定大小的块时, 知道在哪个箱中进行检索. 如果你释放一个块, 然后立刻分配一个同样大小的块, `malloc` 通常会很快.

Chapter 7

缓存

7.1 程序如何运行

为了理解缓存, 你需要了解计算机是如何执行程序的. 要想对这个主题有深入了解, 你需要学习计算机体系架构. 本章的目标是提供一个程序执行的简单模型.

当程序启动时, 代码 (或文本) 通常是位于硬盘或固态硬盘上. 操作系统会创建一个进程来运行程序, 然后“加载器”会从存储中将文本拷贝到主内存中, 并通过调用 `main` 以启动程序.

当程序运行时, 大部分的数据存储在主内存中, 但有些数据存储在寄存器中, 寄存器是 CPU 上的小的内存单元. 这些寄存器包括:

- 程序计数器 (PC), 其包含程序中下一条指令的内存地址.
- 指令寄存器 (IR), 其包含当前运行的机器码指令.
- 栈指针 (SP), 其包含当前函数 (包括其参数和局部变量) 的栈帧地址.
- 通用寄存器, 用于保存程序当前处理的数据.
- 状态寄存器, 或标志寄存器, 其包含当前计算的有关信息. 比如, 标志寄存器通常包含一个位, 如果前一操作的结果为零, 则该位便被设置.

当程序运行时, CPU 执行以下步骤, 被称为“指令周期”:

- 取指 (Fetch): 从内存获取下一指令, 并存储到指令寄存器中.
- 解码 (Decode): CPU 的一部分, 叫做“控制单元”, 解码指令并向 CPU 其他部分发送信号.

- 执行 (Execute): 来自控制单元的信号会引发适当计算的执行.

多数计算机可以执行几百个不同的指令, 叫做“指令集”. 但是大部分的指令可以归纳为以下几种通用类别:

- 加载 (Load): 将值从内存传输到寄存器.
- 算术/逻辑 (Arithmetic/logic): 从寄存器加载操作数, 执行数学运算, 将结果存储到寄存器.
- 存储 (Store): 将值从寄存器传输到内存.
- 跳转/分支 (Jump/branch): 更改程序计数器的值, 令执行流程跳转到程序的另一个位置. 分支通常是有条件的, 也就意味着它们会检查标志寄存器中的标志, 且仅在标志被设置时才跳转.

一些指令集, 包括无处不在的 x86, 提供了一些将加载和算术运算组合在一起的指令.

在每个指令周期中, 一条指令都是从程序文本中读取. 此外, 典型的程序中大约一半的指令是加载或存储数据. 因此正是在这里, 存在计算机架构的一个基本问题: “内存瓶颈”.

在现在计算机中, 典型的核心处理器可以在不到 1 纳秒内执行一条指令. 而它与内存交换数据则需要 100 纳秒. 如果 CPU 需要等待 100 纳秒来获取下一条指令, 又等待 100 纳秒以加载数据, 那么它完成指令的速度相比于其理论可行速度要慢 200 倍. 对于许多计算任务来说, 限制速度的因素是内存, 而不是 CPU.

7.2 缓存性能

这个问题的解决方案, 或至少一个部分解决方案, 便是缓存. “缓存”是一种小型, 高速的内存, 在物理上靠近 CPU, 通常位于同一个芯片上.

实际上, 现代计算机通常有多级缓存: 一级缓存, 是最小最快的, 可能有 1-2 MiB, 访问时间接近 1 纳秒; 二级缓存访问时间可能接近 4 纳秒; 而三级缓存可能需要 16 纳秒.

当 CPU 从内存加载数据, 会将拷贝保存在缓存中. 如果同样的值需要再次加载, CPU 会取用缓存中的拷贝, 无需等待内存.

最终缓存会被填满. 然后, 为了引入新数据, 我们需要先淘汰一些旧数据. 所以如果 CPU 加载了一个值, 很久之后再次加载, 这个值可能已经不在缓存中了.

许多程序的性能受限于缓存效率. 如果 CPU 需要的指令和数据通常都在缓存中, 程序便能以接近 CPU 满速的性能执行. 如果 CPU 频繁需要不在缓存中的数据, 程序便受限于内存速度了.

缓存的“命中率 (hit rate)”, h , 是在内存访问时找到数据在缓存中的比例; “未命中率 (miss rate)”, m , 是在内存访问时必须访问内存的比例. 如果处理缓存命中的时间是 T_h , 同时缓存未命中的时间是 T_m , 每次内存访问的平均时间为

$$hT_h + mT_m$$

同样的, 我们可以将“未命中惩罚”定义为处理未命中缓存的额外时间, $T_p = T_m - T_h$. 那么平均访问时间便是

$$T_h + mT_p$$

当未命中率较低时, 平均访问时间会接近于 T_h . 也就是说, 程序运行中内存访问的速度接近于缓存速度.

7.3 局部性

当程序首次读取一个字节, 缓存通常会加载包含请求字节及其相邻字节的“块”或“行”数据. 如果程序继续读取一个相邻字节, 那么已经在缓存中了.

举个例子, 假设块大小是 64 B; 你读取一个长度 64 的字符串, 字符串的首个字节恰好位于块的开头. 当你加载第一个字节时, 会遇到缺失惩罚, 但之后的整个字符串都会存在于缓存中. 读取整个字符串, 命中率将会是 63/64, 大约 98%. 如果字符串横跨两个块, 你会遇到 2 次缺失惩罚. 但即便如此, 命中率仍将是 62/64, 接近 97%. 如果你之后再次读取相同字符串, 命中率将是 100%.

另一方面, 如果程序以不可预料的方式跳转, 从内存的分散位置读取数据, 并几乎不会再次访问同样位置, 那缓存性能将会变差.

程序多次使用相同数据的趋势叫做“时间局部性”. 而使用临近位置数据的趋势叫做“空间局部性”. 幸运的是, 很多程序天然地表现出了两种局部性:

- 多数程序包含没有跳转和分支的代码块. 在这些块中, 指令顺序运行, 所以访问模式具有空间局部性.
- 在循环中, 程序多次执行相同的指令, 因此访问模式具有空间局部性.
- 一个指令的结果通常会立刻用作下一个指令的操作数, 所以数据访问模式具有时间局部性.

- 当程序执行函数时，其参数和局部变量都被存储在堆栈上；访问这些值便具有空间局部性。
- 读取或写入数组元素的一个最常见的模式便是顺序操作；这种模式也具有空间局部性。

下一节会探讨程序的访问模式和缓存性能之间的关系。

7.4 测量缓存性能

When I was a graduate student at U.C. Berkeley I was a teaching assistant for Computer Architecture with Brian Harvey. One of my favorite exercises involved a program that iterates through an array and measures the average time to read and write an element. By varying the size of the array, it is possible to infer the size of the cache, the block size, and some other attributes.

My modified version of this program is in the `cache` directory of the repository for this book (see Section 0.1).

The important part of the program is this loop:

```
iters = 0;
do {
    sec0 = get_seconds();

    for (index = 0; index < limit; index += stride)
        array[index] = array[index] + 1;

    iters = iters + 1;
    sec = sec + (get_seconds() - sec0);

} while (sec < 0.1);
```

The inner `for` loop traverses the array. `limit` determines how much of the array it traverses; `stride` determines how many elements it skips over. For example, if `limit` is 16 and `stride` is 4, the loop would access elements 0, 4, 8, and 12.

`sec` keeps track of the total CPU time used by the inner loop. The outer loop runs until `sec` exceeds 0.1 seconds, which is long enough that we can compute the average time with sufficient precision.

`get_seconds` uses the system call `clock_gettime`, converts to seconds, and returns the result as a `double`:

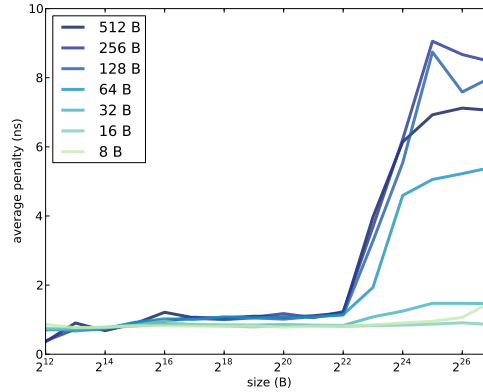


图 7.1: Average miss penalty as a function of array size and stride.

```
double get_seconds(){
    struct timespec ts;
    clock_gettime(CLOCK_PROCESS_CPUTIME_ID, &ts);
    return ts.tv_sec + ts.tv_nsec / 1e9;
}
```

To isolate the time to access the elements of the array, the program runs a second loop that is almost identical except that the inner loop doesn't touch the array; it always increments the same variable:

```
    iters2 = 0;
    do {
        sec0 = get_seconds();

        for (index = 0; index < limit; index += stride)
            temp = temp + index;

        iters2 = iters2 + 1;
        sec = sec - (get_seconds() - sec0);

    } while (iters2 < iters);
```

The second loop runs the same number of iterations as the first. After each iteration, it *subtracts* the elapsed time from `sec`. When the loop completes, `sec` contains the total time for all array accesses, minus the total time it took to increment `temp`. This difference is the total miss penalty incurred by all accesses. Finally, we divide by the number of accesses to get the average miss penalty per access, in ns:

```
sec * 1e9 / iters / limit * stride
```

If you compile and run `cache.c` you should see output like this:

```
Size:    4096 Stride:      8 read+write: 0.8633 ns
Size:    4096 Stride:     16 read+write: 0.7023 ns
Size:    4096 Stride:     32 read+write: 0.7105 ns
Size:    4096 Stride:     64 read+write: 0.7058 ns
```

If you have Python and `matplotlib` installed, you can use `graph_data.py` to graph the results. Figure 7.1 shows the results when I ran it on a Dell Optiplex 7010. Notice that the array size and stride are reported in bytes, not number of array elements.

Take a minute to consider this graph, and see what you can infer about the cache. Here are some things to think about:

- The program reads through the array many times, so it has plenty of temporal locality. If the entire array fits in cache, we expect the average miss penalty to be near 0.
- When the stride is 4 bytes, we read every element of the array, so the program has plenty of spatial locality. If the block size is big enough to contain 64 elements, for example, the hit rate would be 63/64, even if the array does not fit in cache.
- If the stride is equal to the block size (or greater), the spatial locality is effectively zero, because each time we read a block, we only access one element. In that case we expect to see the maximum miss penalty.

In summary, we expect good cache performance if the array is smaller than the cache size *or* if the stride is smaller than the block size. Performance only degrades if the array is bigger than the cache *and* the stride is large.

In Figure 7.1, cache performance is good, for all strides, as long as the array is less than 2^{22} B. We can infer that the cache size is near 4 MiB; in fact, according to the specs, it is 3 MiB.

When the stride is 8, 16, or 32 B, cache performance is good. At 64 B it starts to degrade, and for larger strides the average miss penalty is about 9 ns. We can infer that the block size near 128 B.

Many processors use “multi-level caches” that include a small, fast cache and a bigger, slower cache. In this example, it looks like the miss penalty increases a little when the array size is bigger than 2^{14} B, so it’s possible that this processor also has a 16 KB cache with an access time less than 1 ns.

7.5 Programming for cache performance

Memory caching is implemented in hardware, so most of the time programmers don't need to know much about it. But if you know how caches work, you can write programs that use them more effectively.

For example, if you are working with a large array, it might be faster to traverse the array once, performing several operations with each element, rather than traversing the array several times.

If you are working with a 2-D array, it might be stored as an array of rows. If you traverse through the elements, it would be faster to go row-wise, with stride equal to the element size, rather than column-wise, with stride equal to the row length.

Linked data structures don't always exhibit spatial locality, because the nodes aren't necessarily contiguous in memory. But if you allocate many nodes at the same time, they are usually co-located in the heap. Or, even better, if you allocate an array of nodes all at once, you know they will be contiguous.

Recursive strategies like mergesort often have good cache behavior because they break big arrays into smaller pieces and then work with the pieces. Sometimes these algorithms can be tuned to take advantage of cache behavior.

For applications where performance is critical, it is possible to design algorithms tailored to the size of the cache, the block size, and other hardware characteristics. Algorithms like that are called "cache-aware". The obvious drawback of cache-aware algorithms is that they are hardware-specific.

7.6 The memory hierarchy

At some point during this chapter, a question like the following might have occurred to you: "If caches are so much faster than main memory, why not make a really big cache and forget about memory?"

Without going too far into computer architecture, there are two reasons: electronics and economics. Caches are fast because they are small and close to the CPU, which minimizes delays due to capacitance and signal propagation. If you make a cache big, it will be slower.

Also, caches take up space on the processor chip, and bigger chips are more expensive. Main memory is usually dynamic random-access memory (DRAM), which uses only one transistor and one capacitor per bit, so it is possible to pack

more memory into the same amount of space. But this way of implementing memory is slower than the way caches are implemented.

Also main memory is usually packaged in a dual in-line memory module (DIMM) that includes 16 or more chips. Several small chips are cheaper than one big one.

The trade-off between speed, size, and cost is the fundamental reason for caching. If there were one memory technology that was fast, big, and cheap, we wouldn't need anything else.

The same principle applies to storage as well as memory. Solid state drives (SSD) are fast, but they are more expensive than hard drives (HDD), so they tend to be smaller. Tape drives are even slower than hard drives, but they can store large amounts of data relatively cheaply.

The following table shows typical access times, sizes, and costs for each of these technologies.

Device	Access time	Typical size	Cost
Register	0.5 ns	256 B	?
Cache	1 ns	2 MiB	?
DRAM	100 ns	4 GiB	\$10 / GiB
SSD	10 μ s	100 GiB	\$1 / GiB
HDD	5 ms	500 GiB	\$0.25 / GiB
Tape	minutes	1–2 TiB	\$0.02 / GiB

The number and size of registers depends on details of the architecture. Current computers have about 32 general-purpose registers, each storing one “word”. On a 32-bit computer, a word is 32 bits or 4 B. On a 64-bit computer, a word is 64 bits or 8 B. So the total size of the register file is 100–300 B.

The cost of registers and caches is hard to quantify. They contribute to the cost of the chips they are on, but consumers don't see that cost directly.

For the other numbers in the table, I looked at the specifications for typical hardware for sale from online computer hardware stores. By the time you read this, these numbers will be obsolete, but they give you an idea of what the performance and cost gaps looked like at one point in time.

These technologies make up the “memory hierarchy” (note that this use of “memory” also includes storage). Each level of the hierarchy is bigger and slower than the one above it. And in some sense, each level acts as a cache for

the one below it. You can think of main memory as a cache for programs and data that are stored permanently on SSDs and HDDs. And if you are working with very large datasets stored on tape, you could use hard drives to cache one subset of the data at a time.

7.7 Caching policy

The memory hierarchy suggests a framework for thinking about caching. At every level of the hierarchy, we have to address four fundamental questions of caching:

- Who moves data up and down the hierarchy? At the top of the hierarchy, register allocation is usually done by the compiler. Hardware on the CPU handles the memory cache. Users implicitly move data from storage to memory when they execute programs and open files. But the operating system also moves data back and forth between memory and storage. At the bottom of the hierarchy, administrators move data explicitly between disk and tape.
- What gets moved? In general, block sizes are small at the top of the hierarchy and bigger at the bottom. In a memory cache, a typical block size is 128 B. Pages in memory might be 4 KiB, but when the operating system reads a file from disk, it might read 10s or 100s of blocks at a time.
- When does data get moved? In the most basic cache, data gets moved into cache when it is used for the first time. But many caches use some kind of “prefetching”, meaning that data is loaded before it is explicitly requested. We have already seen one form of prefetching: loading an entire block when only part of it is requested.
- Where in the cache does the data go? When the cache is full, we can’t bring anything in without kicking something out. Ideally, we want to keep data that will be used again soon and replace data that won’t.

The answers to these questions make up the “cache policy”. Near the top of the hierarchy, cache policies tend to be simple because they have to be fast and they are implemented in hardware. Near the bottom of the hierarchy, there is more time to make decisions, and well-designed policies can make a big difference.

Most cache policies are based on the principle that history repeats itself; if we have information about the recent past, we can use it to predict the immediate

future. For example, if a block of data has been used recently, we expect it to be used again soon. This principle suggests a replacement policy called “least recently used,” or LRU, which removes from the cache a block of data that has not been used recently. For more on this topic, see http://en.wikipedia.org/wiki/Cache_algorithms.

7.8 Paging

In systems with virtual memory, the operating system can move pages back and forth between memory and storage. As I mentioned in Section 6.2, this mechanism is called “paging” or sometimes “swapping”.

Here’s how the process works:

1. Suppose Process A calls `malloc` to allocate a chunk. If there is no free space in the heap with the requested size, `malloc` calls `sbrk` to ask the operating system for more memory.
2. If there is a free page in physical memory, the operating system adds it to the page table for Process A, creating a new range of valid virtual addresses.
3. If there are no free pages, the paging system chooses a “victim page” belonging to Process B. It copies the contents of the victim page from memory to disk, then it modifies the page table for Process B to indicate that this page is “swapped out”.
4. Once the data from Process B is written, the page can be reallocated to Process A. To prevent Process A from reading Process B’s data, the page should be cleared.
5. At this point the call to `sbrk` can return, giving `malloc` additional space in the heap. Then `malloc` allocates the requested chunk and returns. Process A can resume.
6. When Process A completes, or is interrupted, the scheduler might allow Process B to resume. When Process B accesses a page that has been swapped out, the memory management unit notices that the page is “invalid” and causes an interrupt.
7. When the operating system handles the interrupt, it sees that the page is swapped out, so it transfers the page back from disk to memory.

8. Once the page is swapped in, Process B can resume.

When paging works well, it can greatly improve the utilization of physical memory, allowing more processes to run in less space. Here's why:

- Most processes don't use all of their allocated memory. Many parts of the text segment are never executed, or execute once and never again. Those pages can be swapped out without causing any problems.
- If a program leaks memory, it might leave allocated space behind and never access it again. By swapping those pages out, the operating system can effectively plug the leak.
- On most systems, there are processes like daemons that sit idle most of the time and only occasionally "wake up" to respond to events. While they are idle, these processes can be swapped out.
- A user might have many windows open, but only a few are active at a time. The inactive processes can be swapped out.
- Also, there might be many processes running the same program. These processes can share the same text and static segments, avoiding the need to keep multiple copies in physical memory.

If you add up the total memory allocated to all processes, it can greatly exceed the size of physical memory, and yet the system can still behave well.

Up to a point.

When a process accesses a page that's swapped out, it has to get the data back from disk, which can take several milliseconds. The delay is often noticeable. If you leave a window idle for a long time and then switch back to it, it might start slowly, and you might hear the disk drive working while pages are swapped in.

Occasional delays like that might be acceptable, but if you have too many processes using too much space, they start to interfere with each other. When Process A runs, it evicts the pages Process B needs. Then when B runs, it evicts the pages A needs. When this happens, both processes slow to a crawl and the system can become unresponsive. This scenario is called "thrashing".

In theory, operating systems could avoid thrashing by detecting an increase in paging and blocking or killing processes until the system is responsive again. But as far as I can tell, most systems don't do this, or don't do it well; it is often left to users to limit their use of physical memory or try to recover when thrashing occurs.

Chapter 8

Multitasking

In many current systems, the CPU contains multiple cores, which means it can run several processes at the same time. In addition, each core is capable of “multitasking”, which means it can switch from one process to another quickly, creating the illusion that many processes are running at the same time.

The part of the operating system that implements multitasking is the “kernel”. In a nut or seed, the kernel is the innermost part, surrounded by a shell. In an operating system, the kernel is the lowest level of software, surrounded by several other layers, including an interface called a “shell.” Computer scientists love extended metaphors.

At its most basic, the kernel’s job is to handle interrupts. An “interrupt” is an event that stops the normal instruction cycle and causes the flow of execution to jump to a special section of code called an “interrupt handler”.

A **hardware interrupt** is caused when a device sends a signal to the CPU. For example, a network interface might cause an interrupt when a packet of data arrives, or a disk drive might cause an interrupt when a data transfer is complete. Most systems also have timers that cause interrupts at regular intervals, or after an elapsed time.

A **software interrupt** is caused by a running program. For example, if an instruction cannot complete for some reason, it might trigger an interrupt so the condition can be handled by the operating system. Some floating-point errors, like division by zero, are handled using interrupts.

When a program needs to access a hardware device, it makes a **system call**, which is similar to a function call, except that instead of jumping to the beginning of the function, it executes a special instruction that triggers an

interrupt, causing the flow of execution to jump to the kernel. The kernel reads the parameters of the system call, performs the requested operation, and then resumes the interrupted process.

8.1 Hardware state

Handling interrupts requires cooperation between hardware and software. When an interrupt occurs, there might be several instructions running on the CPU, data stored in registers, and other **hardware state**.

Usually the hardware is responsible for bringing the CPU to a consistent state; for example, every instruction should either complete or behave as if it never started. No instruction should be left half complete. Also, the hardware is responsible for saving the program counter (PC), so the kernel knows where to resume.

Then, usually, it is the responsibility of the interrupt handler to save the rest of the hardware state before it does anything that might modify it, and then restore the saved state before the interrupted process resumes.

Here is an outline of this sequence of events:

1. When the interrupt occurs, the hardware saves the program counter in a special register and jumps to the appropriate interrupt handler.
2. The interrupt handler stores the program counter and the status register in memory, along with the contents of any data registers it plans to use.
3. The interrupt handler runs whatever code is needed to handle the interrupt.
4. Then it restores the contents of the saved registers. Finally, it restores the program counter of the interrupted process, which has the effect of jumping back to the interrupted instruction.

If this mechanism works correctly, there is generally no way for the interrupted process to know there was an interrupt, unless it detects the change in time between instructions.

8.2 Context switching

Interrupt handlers can be fast because they don't have to save the entire hardware state; they only have to save registers they are planning to use.

But when an interrupt occurs, the kernel does not always resume the interrupted process. It has the option of switching to another process. This mechanism is called a “context switch”.

In general, the kernel doesn’t know which registers a process will use, so it has to save all of them. Also, when it switches to a new process, it might have to clear data stored in the memory management unit (see Section 3.6). And after the context switch, it might take some time for the new process to load data into the cache. For these reasons, context switches are relatively slow, on the order of thousands of cycles, or a few microseconds.

In a multi-tasking system, each process is allowed to run for a short period of time called a “time slice” or “quantum”. During a context switch, the kernel sets a hardware timer that causes an interrupt at the end of the time slice. When the interrupt occurs, the kernel can switch to another process or allow the interrupted process to resume. The part of the operating system that makes this decision is the “scheduler”.

8.3 The process life cycle

When a process is created, the operating system allocates a data structure that contains information about the process, called a “process control block” or PCB. Among other things, the PCB keeps track of the process state, which is one of:

- Running, if the process is currently running on a core.
- Ready, if the process could be running, but isn’t, usually because there are more runnable processes than cores.
- Blocked, if the process cannot run because it is waiting for a future event like network communication or a disk read.
- Done, if the process has completed, but has exit status information that has not been read yet.

Here are the events that cause a process to transition from one state to another:

- A process is created when the running program executes a system call like `fork`. At the end of the system call, the new process is usually ready. Then the scheduler might resume the original process (the “parent”) or start the new process (the “child”).

- When a process is started or resumed by the scheduler, its state changes from ready to running.
- When a process is interrupted and the scheduler chooses not to let it resume, its state changes from running to ready.
- If a process executes a system call that cannot complete immediately, like a disk request, it becomes blocked and the scheduler usually chooses another process.
- When an operation like a disk request completes, it causes an interrupt. The interrupt handler figures out which process was waiting for the request and switches its state from blocked to ready. Then the scheduler may or may not choose to resume the unblocked process.
- When a process calls `exit`, the interrupt handler stores the exit code in the PCB and changes the process's state to done.

8.4 Scheduling

As we saw in Section 2.3 there might be hundreds of processes on a computer, but usually most of them are blocked. Most of the time, there are only a few processes that are ready or running. When an interrupt occurs, the scheduler decides which process to start or resume.

On a workstation or laptop, the primary goal of the scheduler is to minimize response time; that is, the computer should respond quickly to user actions. Response time is also important on a server, but in addition the scheduler might try to maximize throughput, which is the number of requests that complete per unit of time.

Usually the scheduler doesn't have much information about what processes are doing, so its decisions are based on a few heuristics:

- Processes might be limited by different resources. A process that does a lot of computation is probably CPU-bound, which means that its run time depends on how much CPU time it gets. A process that reads data from a network or disk might be I/O-bound, which means that it would run faster if data input and output went faster, but would not run faster with more CPU time. Finally, a process that interacts with the user is probably blocked, most of the time, waiting for user actions.

The operating system can sometimes classify processes based on their past behavior, and schedule them accordingly. For example, when an

interactive process is unblocked, it should probably run immediately, because a user is probably waiting for a reply. On the other hand, a CPU-bound process that has been running for a long time might be less time-sensitive.

- If a process is likely to run for a short time and then make a blocking request, it should probably run immediately, for two reasons: (1) if the request takes some time to complete, we should start it as soon as possible, and (2) it is better for a long-running process to wait for a short one, rather than the other way around.

As an analogy, suppose you are making an apple pie. The crust takes 5 minutes to prepare, but then it has to chill for half an hour. It takes 20 minutes to prepare the filling. If you prepare the crust first, you can prepare the filling while the crust is chilling, and you can finish the pie in 35 minutes. If you prepare the filling first, the process takes 55 minutes.

Most schedulers use some form of priority-based scheduling, where each process has a priority that can be adjusted up or down over time. When the scheduler runs, it chooses the runnable process with the highest priority.

Here are some of the factors that determine a process's priority:

- A process usually starts with a relatively high priority so it starts running quickly.
- If a process makes a request and blocks before its time slice is complete, it is more likely to be interactive or I/O-bound, so its priority should go up.
- If a process runs for an entire time slice, it is more likely to be long-running and CPU-bound, so its priority should go down.
- If a task blocks for a long time and then becomes ready, it should get a priority boost so it can respond to whatever it was waiting for.
- If process A is blocked waiting for process B, for example if they are connected by a pipe, the priority of process B should go up.
- The system call `nice` allows a process to decrease (but not increase) its own priority, allowing programmers to pass explicit information to the scheduler.

For most systems running normal workloads, scheduling algorithms don't have a substantial effect on performance. Simple scheduling policies are usually good enough.

8.5 Real-time scheduling

However, for programs that interact with the real world, scheduling can be very important. For example, a program that reads data from sensors and controls motors might have to complete recurring tasks at some minimum frequency and react to external events with some maximum response time. These requirements are often expressed in terms of “tasks” that must be completed before “deadlines”.

Scheduling tasks to meet deadlines is called “real-time scheduling”. For some applications, a general-purpose operating system like Linux can be modified to handle real-time scheduling. These modifications might include:

- Providing richer APIs for controlling task priorities.
- Modifying the scheduler to guarantee that the process with highest priority runs within a fixed amount of time.
- Reorganizing interrupt handlers to guarantee a maximum completion time.
- Modifying locks and other synchronization mechanisms (coming up in the next chapter) to allow a high-priority task to preempt a lower-priority task.
- Choosing an implementation of dynamic memory allocation that guarantees a maximum completion time.

For more demanding applications, especially in domains where real-time response is a matter of life and death, “real-time operating systems” provide specialized capabilities, often with much simpler designs than general purpose operating systems.

Chapter 9

Threads

When I mentioned threads in Section 2.3, I said that a thread is a kind of process. Now I will provide a more careful explanation.

When you create a process, the operating system creates a new address space, which includes the text segment, static segment, and heap; it also creates a new “thread of execution”, which includes the program counter and other hardware state, and the call stack.

The processes we have seen so far are “single-threaded”, which means that only one thread of execution runs in each address space. In this chapter, you will learn about “multi-threaded” processes that have multiple threads running in the same address space.

Within a single process, all threads share the same text segment, so they run the same code. But different threads often run different parts of the code.

And they share the same static segment, so if one thread changes a global variable, other threads see the change. They also share the heap, so threads can share dynamically-allocated chunks.

But each thread has its own stack, so threads can call functions without interfering with each other. Usually threads don’t access each other’s local variables (and sometimes they can’t).

The example code for this chapter is in the repository for this book, in a directory named `counter`. For information on downloading this code, see Section 0.1.

9.1 Creating threads

The most popular threading standard used with C is POSIX Threads, or Pthreads for short. The POSIX standard defines a thread model and an interface for creating and controlling threads. Most versions of UNIX provide an implementation of Pthreads.

Using Pthreads is like using most C libraries:

- You include headers files at the beginning of your program.
- You write code that calls functions defined by Pthreads.
- When you compile the program, you link it with the Pthread library.

For my examples, I include the following headers:

```
#include <stdio.h>
#include <stdlib.h>
#include <pthread.h>
#include <semaphore.h>
```

The first two are standard; the third is for Pthreads and the fourth is for semaphores. To compile with the Pthread library in `gcc`, you can use the `-lpthread` option on the command line:

```
gcc -g -O2 -o array array.c -lpthread
```

This compiles a source file named `array.c` with debugging info and optimization, links with the Pthread library, and generates an executable named `array`.

9.2 Creating threads

The Pthread function that creates threads is called `pthread_create`. The following function shows how to use it:

```
pthread_t make_thread(void *(*entry)(void *), Shared *shared)
{
    int n;
    pthread_t thread;

    n = pthread_create(&thread, NULL, entry, (void *)shared);
    if (n != 0) {
        perror("pthread_create failed");
    }
}
```

```
        exit(-1);
    }
    return thread;
}
```

`make_thread` is a wrapper I wrote to make `pthread_create` easier to use, and to provide error-checking.

The return type from `pthread_create` is `pthread_t`, which you can think of as an id or “handle” for the new thread.

If `pthread_create` succeeds, it returns 0 and `make_thread` returns the handle of the new thread. If an error occurs, `pthread_create` returns an error code and `make_thread` prints an error message and exits.

The parameters of `make_thread` take some explaining. Starting with the second, `Shared` is a structure I defined to contain values shared between threads. The following `typedef` statement creates the new type:

```
typedef struct {
    int counter;
} Shared;
```

In this case, the only shared variable is `counter`. `make_shared` allocates space for a `Shared` structure and initializes the contents:

```
Shared *make_shared()
{
    Shared *shared = check_malloc(sizeof (Shared));
    shared->counter = 0;
    return shared;
}
```

Now that we have a shared data structure, let’s get back to `make_thread`. The first parameter is a pointer to a function that takes a `void` pointer and returns a `void` pointer. If the syntax for declaring this type makes your eyes bleed, you are not alone. Anyway, the purpose of this parameter is to specify the function where the execution of the new thread will begin. By convention, this function is named `entry`:

```
void *entry(void *arg)
{
    Shared *shared = (Shared *) arg;
    child_code(shared);
    pthread_exit(NULL);
}
```

The parameter of `entry` has to be declared as a `void` pointer, but in this program we know that it is really a pointer to a `Shared` structure, so we can typecast it accordingly and then pass it along to `child_code`, which does the real work.

As a simple example, `child_code` prints the value of the shared counter and increments it.

```
void child_code(Shared *shared)
{
    printf("counter = %d\n", shared->counter);
    shared->counter++;
}
```

When `child_code` returns, `entry` invokes `pthread_exit` which can be used to pass a value to the thread that joins with this thread. In this case, the child has nothing to say, so we pass `NULL`.

Finally, here is the code that creates the child threads:

```
int i;
pthread_t child[NUM_CHILDREN];

Shared *shared = make_shared(1000000);

for (i=0; i<NUM_CHILDREN; i++) {
    child[i] = make_thread(entry, shared);
}
```

`NUM_CHILDREN` is a compile-time constant that determines the number of child threads. `child` is an array of thread handles.

9.3 Joining threads

When one thread wants to wait for another thread to complete, it invokes `pthread_join`. Here is my wrapper for `pthread_join`:

```
void join_thread(pthread_t thread)
{
    int ret = pthread_join(thread, NULL);
    if (ret == -1) {
        perror("pthread_join failed");
        exit(-1);
    }
}
```


The parameter is the handle of the thread you want to wait for. All the wrapper does is call `pthread_join` and check the result.

Any thread can join any other thread, but in the most common pattern the parent thread creates and joins all child threads. Continuing the example from the previous section, here's the code that waits on the children:

```
for (i=0; i<NUM_CHILDREN; i++) {  
    join_thread(child[i]);  
}
```

This loop waits for the children one at a time in the order they were created. There is no guarantee that the child threads complete in that order, but this loop works correctly even if they don't. If one of the children is late, the loop might have to wait, and other children might complete in the meantime. But regardless, the loop exits only when all children are done.

If you have downloaded the repository for this book (see Section 0.1), you'll find this example in `counter/counter.c`. You can compile and run it like this:

```
$ make counter  
gcc -Wall counter.c -o counter -lpthread  
$ ./counter
```

When I ran it with 5 children, I got the following output:

```
counter = 0  
counter = 0  
counter = 1  
counter = 0  
counter = 3
```

When you run it, you will probably get different results. And if you run it again, you might get different results each time. What's going on?

9.4 Synchronization errors

The problem with the previous program is that the children access the shared variable, `counter`, without synchronization, so several threads can read the same value of `counter` before any threads increment it.

Here is a sequence of events that could explain the output in the previous section:

```
Child A reads 0  
Child B reads 0
```

```
Child C reads 0
Child A prints 0
Child B prints 0
Child A sets counter=1
Child D reads 1
Child D prints 1
Child C prints 0
Child A sets counter=1
Child B sets counter=2
Child C sets counter=3
Child E reads 3
Child E prints 3
Child D sets counter=4
Child E sets counter=5
```

Each time you run the program, threads might be interrupted at different points, or the scheduler might choose different threads to run, so the sequence of events, and the results, will be different.

Suppose we want to impose some order. For example, we might want each thread to read a different value of `counter` and increment it, so that the value of `counter` reflects the number of threads that have executed `child_code`.

To enforce that requirement, we can use a “mutex”, which is an object that guarantees “mutual exclusion” for a block of code; that is, only one thread can execute the block at a time.

I have written a small module called `mutex.c` that provides mutex objects. I’ll show you how to use it first; then I’ll explain how it works.

Here’s a version of `child_code` that uses a mutex to synchronize threads:

```
void child_code(Shared *shared)
{
    mutex_lock(shared->mutex);
    printf("counter = %d\n", shared->counter);
    shared->counter++;
    mutex_unlock(shared->mutex);
}
```

Before any thread can access `counter`, it has to “lock” the mutex, which has the effect of barring all other threads. Suppose Thread A has locked the mutex and is in the middle of `child_code`. If Thread B arrives and executes `mutex_lock`, it blocks.

When Thread A is done, it executes `mutex_unlock`, which allows Thread B to proceed. In effect, the threads line up to execute `child_code` one at a time,

so they can't interfere with each other. When I run this code with 5 children, I get:

```
counter = 0
counter = 1
counter = 2
counter = 3
counter = 4
```

And that satisfies the requirements. In order for this solution to work, I have to add the Mutex to the Shared struct:

```
typedef struct {
    int counter;
    Mutex *mutex;
} Shared;
```

And initialize it in `make_shared`

```
Shared *make_shared(int end)
{
    Shared *shared = check_malloc(sizeof(Shared));
    shared->counter = 0;
    shared->mutex = make_mutex();    //-- this line is new
    return shared;
}
```

The code in this section is in `counter_mutex.c`. The definition of `Mutex` is in `mutex.c`, which I explain in the next section.

9.5 Mutex

My definition of `Mutex` is a wrapper for a type called `pthread_mutex_t`, which is defined in the POSIX threads API.

To create a POSIX mutex, you have to allocate space for a `pthread_mutex_t` type and then call `pthread_mutex_init`.

One of the problems with this API is that `pthread_mutex_t` behaves like a structure, so if you pass it as an argument, it makes a copy, which makes the mutex behave incorrectly. To avoid that, you have to pass `pthread_mutex_t` by address.

My code makes it easier to get that right. It defines a type, `Mutex`, which is just a more readable name for `pthread_mutex_t`:

```
#include <pthread.h>
```

```
typedef pthread_mutex_t Mutex;
```

Then it defines `make_mutex`, which allocates space and initializes the mutex:

```
Mutex *make_mutex()
{
    Mutex *mutex = check_malloc(sizeof(Mutex));
    int n = pthread_mutex_init(mutex, NULL);
    if (n != 0) perror_exit("make_lock failed");
    return mutex;
}
```

The return value is a pointer, which you can pass around as an argument without causing unwanted copying.

The functions to lock and unlock the mutex are simple wrappers for POSIX functions:

```
void mutex_lock(Mutex *mutex)
{
    int n = pthread_mutex_lock(mutex);
    if (n != 0) perror_exit("lock failed");
}

void mutex_unlock(Mutex *mutex)
{
    int n = pthread_mutex_unlock(mutex);
    if (n != 0) perror_exit("unlock failed");
}
```

This code is in `mutex.c` and the header file `mutex.h`.

Chapter 10

Condition variables

Many simple synchronization problems can be solved using mutexes as shown in the previous chapter. In this chapter I introduce a bigger challenge, the well-known “Producer-Consumer problem”, and a new tool to solve it, the condition variable.

10.1 The work queue

In some multi-threaded programs, threads are organized to perform different tasks. Often they communicate with each other using a queue, where some threads, called “producers”, put data into the queue and other threads, called “consumers”, take data out.

For example, in applications with a graphical user interface, there might be one thread that runs the GUI, responding to user events, and another thread that processes user requests. In that case, the GUI thread might put requests into a queue and the “back end” thread might take requests out and process them.

To support this organization, we need a queue implementation that is “thread safe”, which means that both threads (or more than two) can access the queue at the same time. And we need to handle the special cases when the queue is empty and, if the size of the queue is bounded, when the queue is full.

I’ll start with a simple queue that is not thread safe, then we’ll see what goes wrong and fix it. The code for this example is in the repository for this book, in a folder called `queue`. The file `queue.c` contains a basic implementation of a circular buffer, which you can read about at https://en.wikipedia.org/wiki/Circular_buffer.

Here's the structure definition:

```
typedef struct {
    int *array;
    int length;
    int next_in;
    int next_out;
} Queue;
```

`array` is the array that contains the elements of the queue. For this example the elements are ints, but more generally they would be structures that contain user events, items of work, etc.

`length` is the length of the array. `next_in` is an index into the array that indicates where the next element should be added; similarly, `next_out` is the index of the next element that should be removed.

`make_queue` allocates space for this structure and initializes the fields:

```
Queue *make_queue(int length)
{
    Queue *queue = (Queue *) malloc(sizeof(Queue));
    queue->length = length + 1;
    queue->array = (int *) malloc(length * sizeof(int));
    queue->next_in = 0;
    queue->next_out = 0;
    return queue;
}
```

The initial value for `next_out` needs some explaining. Since the queue is initially empty, there is no next element to remove, so `next_out` is invalid. Setting `next_out == next_in` is a special case that indicates that the queue is empty, so we can write:

```
int queue_empty(Queue *queue)
{
    return (queue->next_in == queue->next_out);
}
```

Now we can add elements to the queue using `queue_push`:

```
void queue_push(Queue *queue, int item) {
    if (queue_full(queue)) {
        perror_exit("queue is full");
    }

    queue->array[queue->next_in] = item;
```

```
    queue->next_in = queue_incr(queue, queue->next_in);  
}
```

If the queue is full, `queue_push` prints an error message and exits. I will explain `queue_full` soon.

If the queue is not full, `queue_push` inserts the new element and then increments `next_in` using `queue_incr`:

```
int queue_incr(Queue *queue, int i)  
{  
    return (i+1) % queue->length;  
}
```

When the index, `i`, gets to the end of the array, it wraps around to 0. And that's where we run into a tricky part. If we keep adding elements to the queue, eventually `next_in` wraps around and catches up with `next_out`. But if `next_in == next_out`, we would incorrectly conclude that the queue was empty.

To avoid that, we define another special case to indicate that the queue is full:

```
int queue_full(Queue *queue)  
{  
    return (queue_incr(queue, queue->next_in) == queue->next_out);  
}
```

If incrementing `next_in` lands on `next_out`, that means we can't add another element without making the queue seem empty. So we stop one element before the "end" (keeping in mind that the end of the queue can be anywhere, not necessarily the end of the array).

Now we can write `queue_pop`, which removes and returns the next element from the queue:

```
int queue_pop(Queue *queue) {  
    if (queue_empty(queue)) {  
        perror_exit("queue is empty");  
    }  
  
    int item = queue->array[queue->next_out];  
    queue->next_out = queue_incr(queue, queue->next_out);  
    return item;  
}
```

If you try to pop from an empty queue, `queue_pop` prints an error message and exits.

10.2 Producers and consumers

Now let's make some threads to access this queue. Here's the producer code:

```
void *producer_entry(void *arg) {
    Shared *shared = (Shared *) arg;

    for (int i=0; i<QUEUE_LENGTH-1; i++) {
        printf("adding item %d\n", i);
        queue_push(shared->queue, i);
    }
    pthread_exit(NULL);
}
```

Here's the consumer code:

```
void *consumer_entry(void *arg) {
    int item;
    Shared *shared = (Shared *) arg;

    for (int i=0; i<QUEUE_LENGTH-1; i++) {
        item = queue_pop(shared->queue);
        printf("consuming item %d\n", item);
    }
    pthread_exit(NULL);
}
```

Here's the parent code that starts the threads and waits for them

```
pthread_t child[NUM_CHILDREN];

Shared *shared = make_shared();

child[0] = make_thread(producer_entry, shared);
child[1] = make_thread(consumer_entry, shared);

for (int i=0; i<NUM_CHILDREN; i++) {
    join_thread(child[i]);
}
```

And finally here's the shared structure that contains the queue:

```
typedef struct {
    Queue *queue;
} Shared;
```

```
Shared *make_shared()
```



```

{
    Shared *shared = check_malloc(sizeof(Shared));
    shared->queue = make_queue(Queue_LENGTH);
    return shared;
}

```

The code we have so far is a good starting place, but it has several problems:

- Access to the queue is not thread safe. Different threads could access `array`, `next_in`, and `next_out` at the same time and leave the queue in a broken, “inconsistent” state.
- If the consumer is scheduled first, it finds the queue empty, print an error message, and exits. We would rather have the consumer block until the queue is not empty. Similarly, we would like the producer to block if the queue is full.

In the next section, we solve the first problem with a **Mutex**. In the following section, we solve the second problem with condition variables.

10.3 Mutual exclusion

We can make the queue thread safe with a mutex. This version of the code is in `queue_mutex.c`.

First we add a **Mutex** pointer to the queue structure:

```

typedef struct {
    int *array;
    int length;
    int next_in;
    int next_out;
    Mutex *mutex;           //-- this line is new
} Queue;

```

And initialize the **Mutex** in `make_queue`:

```

Queue *make_queue(int length) {
    Queue *queue = (Queue *) malloc(sizeof(Queue));
    queue->length = length;
    queue->array = (int *) malloc(length * sizeof(int));
    queue->next_in = 0;
    queue->next_out = 0;
    queue->mutex = make_mutex();    //-- new
    return queue;
}

```

Next we add synchronization code to `queue_push`:

```
void queue_push(Queue *queue, int item) {
    mutex_lock(queue->mutex);    //-- new
    if (queue_full(queue)) {
        mutex_unlock(queue->mutex);    //-- new
        perror_exit("queue is full");
    }

    queue->array[queue->next_in] = item;
    queue->next_in = queue_incr(queue, queue->next_in);
    mutex_unlock(queue->mutex);    //-- new
}
```

Before checking whether the queue is full, we have to lock the `Mutex`. If the queue is full, we have to unlock the `Mutex` before exiting; otherwise the thread would leave it locked and no other threads could proceed.

The synchronization code for `queue_pop` is similar:

```
int queue_pop(Queue *queue) {
    mutex_lock(queue->mutex);
    if (queue_empty(queue)) {
        mutex_unlock(queue->mutex);
        perror_exit("queue is empty");
    }

    int item = queue->array[queue->next_out];
    queue->next_out = queue_incr(queue, queue->next_out);
    mutex_unlock(queue->mutex);
    return item;
}
```

Note that the other `Queue` functions, `queue_full`, `queue_empty`, and `queue_incr` do not try to lock the mutex. Any thread that calls these functions is required to lock the mutex first; this requirement is part of the documented interface for these functions.

With this additional code, the queue is thread safe; if you run it, you should not see any synchronization errors. But it is likely that the consumer will exit at some point because the queue is empty, or the producer will exit because the queue is full, or both.

The next step is to add condition variables.

10.4 Condition variables

A condition variable is a data structure associated with a condition; it allows threads to block until the condition becomes true. For example, `thread_pop` might want check whether the queue is empty and, if so, wait for a condition like “queue not empty”.

Similarly, `thread_push` might want to check whether the queue is full and, if so, block until it is not full.

I’ll handle the first condition here, and you will have a chance to handle the second condition as an exercise.

First we add a condition variable to the `Queue` structure:

```
typedef struct {
    int *array;
    int length;
    int next_in;
    int next_out;
    Mutex *mutex;
    Cond *nonempty;    //-- new
} Queue;
```

And initialize it in `make_queue`:

```
Queue *make_queue(int length)
{
    Queue *queue = (Queue *) malloc(sizeof(Queue));
    queue->length = length;
    queue->array = (int *) malloc(length * sizeof(int));
    queue->next_in = 0;
    queue->next_out = 0;
    queue->mutex = make_mutex();
    queue->nonempty = make_cond();    //-- new
    return queue;
}
```

Now in `queue_pop`, if we find the queue empty, we don’t exit; instead we use the condition variable to block:

```
int queue_pop(Queue *queue) {
    mutex_lock(queue->mutex);
    while (queue_empty(queue)) {
        cond_wait(queue->nonempty, queue->mutex);    //-- new
    }
}
```

```

    int item = queue->array[queue->next_out];
    queue->next_out = queue_incr(queue, queue->next_out);
    mutex_unlock(queue->mutex);
    cond_signal(queue->nonfull);    //-- new
    return item;
}

```

`cond_wait` is complicated, so let's take it slow. The first argument is the condition variable; in this case, the condition we are waiting for is "queue not empty". The second argument is the mutex that protects the queue.

When the thread that locked the mutex calls `cond_wait`, it unlocks the mutex and then blocks. This is important. If `cond_wait` did not unlock the mutex before blocking, no other thread would be able to access the queue, no more items could be added, and the queue would always be empty.

So while the consumer is blocked on `nonempty`, the producer can run. Let's see what happens when the producer runs `queue_push`:

```

void queue_push(Queue *queue, int item) {
    mutex_lock(queue->mutex);
    if (queue_full(queue)) {
        mutex_unlock(queue->mutex);
        perror_exit("queue is full");
    }
    queue->array[queue->next_in] = item;
    queue->next_in = queue_incr(queue, queue->next_in);
    mutex_unlock(queue->mutex);
    cond_signal(queue->nonempty);    //-- new
}

```

Just as before, `queue_push` locks the `Mutex` and checks whether the queue is full. Assuming it is not, `queue_push` adds a new element to the queue and then unlocks the `Mutex`.

But before returning, it does one more thing: it "signals" the condition variable `nonempty`.

Signalling a condition variable usually indicates that the condition is true. If there are no threads waiting on the condition variable, the signal has no effect.

If there are threads waiting on the condition variable, one of them gets unblocked and resumes execution of `cond_wait`. But before the awakened thread can return from `cond_wait`, it has to wait for and lock the `Mutex`, again.

Now go back to `queue_pop` and see what happens when the thread returns from `cond_wait`. It loops back to the top of the while loop and checks the

condition again. I'll explain why in just a second, but for now let's assume that the condition is true; that is, the queue is not empty.

When the consumer thread exits the while loop, we know two things: (1) the condition is true, so there is at least one item in the queue, and (2) the `Mutex` is locked, so it is safe to access the queue.

After removing an item, `queue_pop` unlocks the mutex and returns.

In the next section I'll show you how my `Cond` code works, but first I want to answer two frequently-asked questions:

- Why is `cond_wait` inside a while loop rather than an if statement; that is, why do we have to check the condition again after returning from `cond_wait`?

The primary reason you have to re-check the condition is the possibility of an intercepted signal. Suppose Thread A is waiting on `nonempty`. Thread B adds an item to the queue and signals `nonempty`. Thread A wakes up and tries to lock the mutex, but before it gets the chance, Evil Thread C swoops in, locks the mutex, pops the item from the queue, and unlocks the mutex. Now the queue is empty again, but Thread A is not blocked any more. Thread A could lock the mutex and returns from `cond_wait`. If Thread A does not check the condition again, it would try to pop an element from an empty queue, and probably cause an error.

- The other question that comes up when people learn about condition variables is “How does the condition variable know what condition it is associated with?”

This question is understandable because there is no explicit connection between a `Cond` structure and the condition it relates to. The connection is implicit in the way it is used.

Here's one way to think of it: the condition associated with a `Cond` is the thing that is false when you call `cond_wait` and true when you call `cond_signal`.

Because threads have to check the condition when they return from `cond_wait`, it is not strictly necessary to call `cond_signal` only when the condition is true. If you have reason to think the condition *might* be true, you could call `cond_signal` as a suggestion that now is a good time to check.

10.5 Condition variable implementation

The Cond structure I used in the previous section is a wrapper for a type called `pthread_cond_t`, which is defined in the POSIX threads API. It is very similar to Mutex, which is a wrapper for `pthread_mutex_t`. Both wrappers are defined in `utils.c` and `utils.h`.

Here's the typedef:

```
typedef pthread_cond_t Cond;
```

`make_cond` allocates space, initializes the condition variable, and returns a pointer:

```
Cond *make_cond() {
    Cond *cond = check_malloc(sizeof(Cond));
    int n = pthread_cond_init(cond, NULL);
    if (n != 0) perror_exit("make_cond failed");

    return cond;
}
```

And here are the wrappers for `cond_wait` and `cond_signal`.

```
void cond_wait(Cond *cond, Mutex *mutex) {
    int n = pthread_cond_wait(cond, mutex);
    if (n != 0) perror_exit("cond_wait failed");
}

void cond_signal(Cond *cond) {
    int n = pthread_cond_signal(cond);
    if (n != 0) perror_exit("cond_signal failed");
}
```

At this point there should be nothing too surprising there.

Chapter 11

Semaphores in C

Semaphores are a good way to learn about synchronization, but they are not as widely used, in practice, as mutexes and condition variables.

Nevertheless, there are some synchronization problems that can be solved simply with semaphores, yielding solutions that are more demonstrably correct.

This chapter presents a C API for working with semaphores and my code for making it easier to work with. And it presents a final challenge: can you write an implementation of a semaphore using mutexes and condition variables?

The code for this chapter is in directory `semaphore` in the repository for this book (see Section 0.1).

11.1 POSIX Semaphores

A semaphore is a data structure used to help threads work together without interfering with each other.

The POSIX standard specifies an interface for semaphores; it is not part of Pthreads, but most UNIXes that implement Pthreads also provide semaphores.

POSIX semaphores have type `sem_t`. As usual, I put a wrapper around `sem_t` to make it easier to use. The interface is defined in `sem.h`:

```
typedef sem_t Semaphore;
```

```
Semaphore *make_semaphore(int value);  
void semaphore_wait(Semaphore *sem);  
void semaphore_signal(Semaphore *sem);
```

`Semaphore` is a synonym for `sem_t`, but I find it more readable, and the capital letter reminds me to treat it like an object and pass it by pointer.

The implementation of these functions is in `sem.c`:

```
Semaphore *make_semaphore(int value)
{
    Semaphore *sem = check_malloc(sizeof(Semaphore));
    int n = sem_init(sem, 0, value);
    if (n != 0) perror_exit("sem_init failed");
    return sem;
}
```

`make_semaphore` takes the initial value of the semaphore as a parameter. It allocates space for a `Semaphore`, initializes it, and returns a pointer to `Semaphore`.

`sem_init` returns 0 if it succeeds and -1 if anything goes wrong. One nice thing about using wrapper functions is that you can encapsulate the error-checking code, which makes the code that uses these functions more readable.

Here is the implementation of `semaphore_wait`:

```
void semaphore_wait(Semaphore *sem)
{
    int n = sem_wait(sem);
    if (n != 0) perror_exit("sem_wait failed");
}
```

And here is `semaphore_signal`:

```
void semaphore_signal(Semaphore *sem)
{
    int n = sem_post(sem);
    if (n != 0) perror_exit("sem_post failed");
}
```

I prefer to call this operation “signal” rather than “post”, although both terms are common.

Here’s an example that shows how to use a semaphore as a mutex:

```
Semaphore *mutex = make_semaphore(1);

semaphore_wait(mutex);
    // protected code goes here
semaphore_signal(mutex);
```


When you use a semaphore as a mutex, you usually initialize it to 1 to indicate that the mutex is unlocked; that is, one thread can pass the semaphore without blocking.

Here I am using the variable name `mutex` to indicate that the semaphore is being used as a mutex. But remember that the behavior of a semaphore is not the same as a Pthread mutex.

11.2 Producers and consumers with semaphores

Using these semaphore wrapper functions, we can write a solution to the Producer-Consumer problem from Section 10.2. The code in this section is in `queue_sem.c`.

Here's the new definition of `Queue`, replacing the mutex and condition variables with semaphores:

```
typedef struct {
    int *array;
    int length;
    int next_in;
    int next_out;
    Semaphore *mutex;      //-- new
    Semaphore *items;      //-- new
    Semaphore *spaces;     //-- new
} Queue;
```

And here's the new version of `make_queue`:

```
Queue *make_queue(int length)
{
    Queue *queue = (Queue *) malloc(sizeof(Queue));
    queue->length = length;
    queue->array = (int *) malloc(length * sizeof(int));
    queue->next_in = 0;
    queue->next_out = 0;
    queue->mutex = make_semaphore(1);
    queue->items = make_semaphore(0);
    queue->spaces = make_semaphore(length-1);
    return queue;
}
```

`mutex` is used to guarantee exclusive access to the queue; the initial value is 1, so the mutex is initially unlocked.

`items` is the number of items in the queue, which is also the number of consumer threads that can execute `queue_pop` without blocking. Initially there are no items in the queue.

`spaces` is the number of empty spaces in the queue, which is the number of producer threads that can execute `queue_push` without blocking. Initially the number of spaces is the capacity of the queue, which is `length-1`, as explained in Section 10.1.

Here is the new version of `queue_push`, which is run by producer threads:

```
void queue_push(Queue *queue, int item) {
    semaphore_wait(queue->spaces);
    semaphore_wait(queue->mutex);

    queue->array[queue->next_in] = item;
    queue->next_in = queue_incr(queue, queue->next_in);

    semaphore_signal(queue->mutex);
    semaphore_signal(queue->items);
}
```

Notice that `queue_push` doesn't have to call `queue_full` any more; instead, the semaphore keeps track of how many spaces are available and blocks producers if the queue is full.

Here is the new version of `queue_pop`:

```
int queue_pop(Queue *queue) {
    semaphore_wait(queue->items);
    semaphore_wait(queue->mutex);

    int item = queue->array[queue->next_out];
    queue->next_out = queue_incr(queue, queue->next_out);

    semaphore_signal(queue->mutex);
    semaphore_signal(queue->spaces);

    return item;
}
```

This solution is explained, using pseudo-code, in Chapter 4 of *The Little Book of Semaphores*.

Using the code in the repository for this book, you should be able to compile and run this solution like this:

```
$ make queue_sem
$ ./queue_sem
```

11.3 Make your own semaphores

Any problem that can be solved with semaphores can also be solved with condition variables and mutexes. We can prove that's true by using condition variables and mutexes to implement a semaphore.

Before you go on, you might want to try this as an exercise: write functions that implement the semaphore API in `sem.h` using condition variables and mutexes. In the repository for this book, you'll find my solution in `mysem_soln.c` and `mysem_soln.h`.

If you have trouble getting started, you can use the following structure definition, from my solution, as a hint:

```
typedef struct {
    int value, wakeups;
    Mutex *mutex;
    Cond *cond;
} Semaphore;
```

`value` is the value of the semaphore. `wakeups` counts the number of pending signals; that is, the number of threads that have been woken but have not yet resumed execution. The reason for wakeups is to make sure that our semaphores have Property 3, described in *The Little Book of Semaphores*.

`mutex` provides exclusive access to `value` and `wakeups`; `cond` is the condition variable threads wait on if they wait on the semaphore.

Here is the initialization code for this structure:

```
Semaphore *make_semaphore(int value)
{
    Semaphore *semaphore = check_malloc(sizeof(Semaphore));
    semaphore->value = value;
    semaphore->wakeups = 0;
    semaphore->mutex = make_mutex();
    semaphore->cond = make_cond();
    return semaphore;
}
```

11.3.1 Semaphore implementation

Here is my implementation of semaphores using POSIX mutexes and condition variables:

```
void semaphore_wait(Semaphore *semaphore)
{
    mutex_lock(semaphore->mutex);
    semaphore->value--;

    if (semaphore->value < 0) {
        do {
            cond_wait(semaphore->cond, semaphore->mutex);
        } while (semaphore->wakeups < 1);
        semaphore->wakeups--;
    }
    mutex_unlock(semaphore->mutex);
}
```

When a thread waits on the semaphore, it has to lock the mutex before it decrements `value`. If the value of the semaphore becomes negative, the thread blocks until a “wakeup” is available. While it is blocked, the mutex is unlocked, so another thread can signal.

Here is the code for `semaphore_signal`:

```
void semaphore_signal(Semaphore *semaphore)
{
    mutex_lock(semaphore->mutex);
    semaphore->value++;

    if (semaphore->value <= 0) {
        semaphore->wakeups++;
        cond_signal(semaphore->cond);
    }
    mutex_unlock(semaphore->mutex);
}
```

Again, a thread has to lock the mutex before it increments `value`. If the semaphore was negative, that means threads are waiting, so the signalling thread increments `wakeups` and signals the condition variable.

At this point one of the waiting threads might wake up, but the mutex is still locked until the signalling thread unlocks it.

At that point, one of the waiting threads returns from `cond_wait` and checks

whether a wakeup is still available. If not, it loops and waits on the condition variable again. If so, it decrements **wakeups**, unlocks the mutex, and exits.

One thing about this solution that might not be obvious is the use of a **do...while** loop. Can you figure out why it is not a more conventional **while** loop? What would go wrong?

The problem is that with a **while** loop this implementation would not have Property 3. It would be possible for a thread to signal and then run around and catch its own signal.

With the **do...while** loop, it is guaranteed¹ that when a thread signals, one of the waiting threads will get the signal, even if the signalling thread runs around and gets the mutex before one of the waiting threads resumes.

¹Well, almost. It turns out that a well-timed spurious wakeup (see http://en.wikipedia.org/wiki/Spurious_wakeup) can violate this guarantee.