

OS 之道

操作系统简介

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Think OS

A Brief Introduction to Operating Systems

Version 0.7.4

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The L^AT_EX source for this book is available from <http://greenteapress.com/thinkos>.

序

在很多计算机科学课程中, 操作系统始终是一个高级话题. 学生在学习它之前, 可能已经学习了如何用 C 编程, 甚至很可能已经学过计算机体系架构的课程. 通常本课程的目标是让学生了解操作系统的设计和实现, 同时期待其中一些学生在这个领域持续研究, 或者编写部分操作系统.

这本书希望可以满足不同人群的不同目标, 我编写此书是为了讲授欧林学院的软件系统的课程.

多数来上课的学生已经学过了如何用 Python 编程. 所以一个目标是能够帮助他们学习 C 语言. 对于这部分内容, 我使用 O'Reilly Media 出版社出版, Griffiths 和 Griffiths 编写的 *Head First C* 一书, 来充实本书的内容.

我的学生中很少会从事编写操作系统, 但是很多人将用 C 来编写底层应用, 以及进入嵌入式领域工作. 本课程包括操作系统, 网络, 数据库, 以及嵌入式系统, 但同时着重强调了程序员需要了解的主题.

本书不会假设你学过计算机架构. 随着课程学习, 我将会解释我们所需要的内容.

如果这本书是成功的, 那么它应该能够让你更好地理解程序运行时发生了什么, 以及做什么可以令程序运行更加高效.

第一章介绍了编译语言和解释语言的差异, 以及编译器工作原理. 推荐阅读: *Head First C* 第 1 章.

第二章解释了操作系统如何使用进程隔离程序运行.

第三章解释了虚拟内存和地址转换. 推荐阅读: *Head First C* 第 2 章.

第四章介绍了文件系统和数据流. 推荐阅读: *Head First C* 第 3 章.

第五章讲述了数字、字母和其他值如何编码, 并介绍了位运算符.

第六章介绍了动态内存管理的使用及其工作原理. 推荐阅读: *Head First C* 第 6 章.

第七章介绍了缓存和内存层次结构.

第八章介绍了多任务和调度.

第九章介绍了 POSIX 线程和互斥锁. 推荐阅读: *Head First C* 第 12 章以及 *Little Book of Semaphores* 第 1 和 2 章.

第十章介绍了 POSIX 条件变量和生产者/消费者问题.

第十一章介绍了 POSIX 信号量的使用以及在 C 语言中的实现.

补充

本书目前还是早期草稿版本. 我还在努力撰写文本, 但尚缺乏一些配图. 所以某些地方, 如果配图完备了, 我确信相关解释会得到改善.

0.1 代码使用

本书的代码示例可以从 <https://github.com/AllenDowney/ThinkOS> 下载. Git 是个版本管理系统, 可以让你跟踪项目文件. Git 管理的文件集合叫做仓库. GitHub 是一个提供 Git 仓库存储和方便的 web 界面的托管服务.

我的仓库的 GitHub 主页提供了多种使用代码的方法:

- 你可以点击 **Fork** 按钮, 在 GitHub 基于我的仓库创建一个副本. 如果你还没有 GitHub 账户, 你需要创建一个. 创建副本后, 你便拥有了自己的仓库, 便可以跟踪学习本书时所编写的代码了. 然后, 你可以克隆这个仓库, 这意味着文件会被复制到你的计算机上.
- 或者你也可以克隆我的仓库. 这样便无需 GitHub 账号, 但你便不能将你的变更写回 GitHub 了.
- 如果你一点都不想用 Git, 你可以使用 GitHub 页面的按钮下载 Zip 文件.

贡献者清单

如果你有建议或更正, 请发送邮件至 downey@allendowney.com. 如果我根据你的反馈进行更正, 我将把你添加到贡献者清单 (除非你要求省略).

如果您在反馈时包含了出现错误的部分语句, 那么我便可以快速定位. 提供页码或者章节编号也可以, 只是不太容易处理. 谢谢!

- 感谢参加 Olin 学院 2014 年春季软件系统课程的学生们, 他们在本书的早期草稿阶段, 进行了大量测试. 同时纠正了很多错误, 并提出相当多的建议. 我非常感激他们的开创精神!
- James P Giannoules 指出了一处复制粘贴错误.
- Andy Engle 指出了 GB 和 GiB 之间的区别.
- Aashish Karki 指出了一些语法问题.

其他发现错别字和其他错误的人包括 Jim Tyson, Donald Robertson, Jeremy Vermast, Yuzhong Huang 和 Ian Hill.

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Chapter 1

编译

1.1 编译语言和解释语言

人们通常将编程语言分为编译型或解释型。“编译”意味着程序被翻译为机器语言,然后硬件执行;“解释”则意味着程序由软件解释器读取并执行。通常认为 C 是编译型语言,而 Python 是解释型语言。但是两者差异并非泾渭分明。

首先,很多语言是既可以被编译也可以被解释的。例如,也有 C 解释器和 Python 编译器。其次,像 Java 一样,有些语言会使用混合方法,将程序编译为一种中间语言,然后在解释器中运行翻译后的程序。Java 使用了一种叫做 Java 字节码的中间语言,和机器语言类似,但它需要由软件解释器执行,也就是 Java 虚拟机 (JVM)。

所以被编译还是被解释,并不是一个语言的内在特征;然而无论如何,编译语言和解释语言之间确实会存在一些普遍的差异。

1.2 静态类型

多数解释语言支持的是动态类型,但是编译语言通常受限于静态类型。在静态类型语言中,你可以通过查看程序便能区分变量引用的类型。而在动态类型语言中,直到程序运行,你才能确切知道变量类型。通常,静态是指编译时发生的事情,而动态是指运行时发生的事情。

例如,可以用 Python 编写如下函数:

```
def add(x, y):  
    return x + y
```

看这个代码, 你很难在运行时确定 `x` 和 `y` 所引用的类型. 这个函数可能被调用多次, 每次都可能是不同类型的值. 任何支持加法运算的值都可以运行; 其他类型则会引发异常或运行时错误.

而在 C 中, 同样函数编写如下::

```
int add(int x, int y) {  
    return x + y;  
}
```

函数第一行包括对参数和返回值的**类型声明**: `x` 和 `y` 被声明为了整数, 这就意味着在编译时, 我们可以检查这种类型是否支持加法运算 (此处支持). 返回值也同样被声明为了整数.

因为这些声明的存在, 在程序其他位置调用这个函数时, 编译器便可以检查提供的参数是否具有正确的类型, 以及返回值是否使用正确.

这些检查发生在程序运行之前, 所以错误可以被提前发现. 更重要的, 可以在未执行的程序部分发现错误. 此外, 这些检查在运行时便无需发生, 这也是编译语言比解释语言运行更快的一个原因.

编译时声明类型还能节省空间. 在动态语言中, 程序运行时, 变量名称需要保存在内存中, 同时它们会被程序频繁访问. 例如, 在 Python 中, 内置函数 `locals` 会返回一个包含变量名称的字典. 下面一个用 Python 解释器执行的示例:

```
>>> x = 5  
>>> print locals()  
{'x': 5, '__builtins__': <module '__builtin__' (built-in)>,  
  '__name__': '__main__', '__doc__': None, '__package__': None}
```

这说明在程序运行时, 变量名是存储在内存中的 (以及一些作为默认运行环境一部分的其他值).

编译语言中, 变量名存在于编译时, 而不是运行时. 编译器会为每个变量选择一个位置, 同时将这些位置记录为编译后的程序的一部分.¹ 变量的位置被称为**地址**. 运行时, 每个变量的值被存储在它的地址上, 变量的名称则不会被存储 (除非编译器出于调试的目的而添加它们).

1.3 编译流程

作为程序员, 你应该对编译期间发生了什么有个大致脉络. 如果你熟悉这个过程, 便可以帮你了解错误信息, 调试代码, 以及避免常见的陷阱.

编译步骤如下:

¹这只是一个简单描述, 稍后我们会详细说明.

1. 预处理: C 是一种包括**预处理指令**的语言, 这些指令是在编译之前起效的. 比如, `#include` 指令会将另一个文件的源代码插入到指令的位置.
2. 解析: 解析期间, 编译器会读取源码并构建程序的内在结构, 也就是**抽象语法树**. 在此期间检查出来的错误, 通常是语法错误.
3. 静态检查: 编译器会检查变量和值的类型是否正确, 调用函数是否使用了正确数量和类型的参数, 等等. 此阶段侦测到的错误通常是**静态语义**错误.
4. 代码生成: 编译器会读取程序的内在结构, 并生成机器码或字节码.
5. 链接: 如果程序使用了库中定义的值和函数, 编译器需要找到合适的库, 并囊括所需的代码.
6. 优化: 在整个流程的各个阶段, 编译器可以转换程序, 生成运行更快, 或更加节约空间的代码. 大部分的优化是一些简单的修改, 可以消除明显的耗费, 但有些编译器可以执行更复杂的分析和代码转换.

通常运行 `gcc` 时, 它会执行所有这些步骤, 生成一个可执行文件. 比如, 下面是一个极简的 C 程序代码:

```
#include <stdio.h>
int main()
{
    printf("Hello World\n");
}
```

如果你将上述代码保存为一个名为 `hello.c` 的文件, 你便能像这样对其进行编译和执行:

```
$ gcc hello.c
$ ./a.out
```

默认情况下, `gcc` 将可执行代码存储在名为 `a.out` (起初表示 “assembler output”) 的文件中. 第二行会运行可执行文件. 前缀 `./` 是告诉 `shell` 在当前目录中查找.

通常最好使用 `-o` 标识为可执行文件提供一个合适的名称:

```
$ gcc hello.c -o hello
$ ./hello
```

1.4 目标码

使用 `-c` 标识, 可以令 `gcc` 编译程序并生成机器码, 但并不会链接库中函数, 也就不会生成可执行文件:

```
$ gcc hello.c -c
```

结果是一个名为 `hello.o` 的文件, 其中 `o` 表示**目标码**, 也就是编译后的程序. 目标码不可执行, 但可以链接到可执行文件中.

UNIX 命令 `nm` 可以读取目标文件, 并生成涉及其所定义和应用的名称的信息. 例如:

```
$ nm hello.o
0000000000000000 T main
                 U puts
```

输出结果表明 `hello.o` 定义了名称 `main`, 以及名为 `puts` 的函数, 这个名字表示 “put string”. 在这个例子中, `gcc` 通过将庞大复杂的 `printf` 函数, 替换为相对简单的 `puts`, 以进行优化.

你可以使用 `-O` 标识来控制 `gcc` 进行多大程度的优化. 默认情况下, 它优化得很少, 这会令调试相对简单. 而选项 `-O1` 会开启常用且安全得优化. 更高得数字会开启更多得优化, 而这也就会需要更长得编译时间.

理论上, 优化除了加速程序之外, 不应该改变程序的行为. 但如果你的程序存在微妙的错误, 那你可能会发现, 优化会令错误暴露或消失. 通常在开发新代码时, 最好关闭优化. 一旦程序正常运行, 并能通过适当的测试, 你便可以开启优化, 并确认测试仍然通过.

1.5 汇编码

类似于 `-c` 标识, `-S` 标识会令 `gcc` 编译程序并生成汇编码, 而这基本上是一种人可读的机器码形式.

```
$ gcc hello.c -S
```

结果是一个名为 `hello.s` 的文件, 看起来像这样:

```
        .file      "hello.c"
        .section   .rodata
.LC0:
        .string    "Hello World"
        .text
        .globl     main
        .type      main, @function

main:
.LFB0:
        .cfi_startproc
        pushq %rbp
        .cfi_def_cfa_offset 16
```



```

        .cfi_offset 6, -16
        movq %rsp, %rbp
        .cfi_def_cfa_register 6
        movl $.LC0, %edi
        call puts
        movl $0, %eax
        popq %rbp
        .cfi_def_cfa 7, 8
        ret
        .cfi_endproc
.LFE0:
        .size      main, .-main
        .ident      "GCC: (Ubuntu/Linaro 4.7.3-1ubuntu1) 4.7.3"
        .section    .note.GNU-stack,"",@progbits

```

gcc 通常会被设置, 为你运行的机器生成相应的代码, 所以对我来说, 它会生成 x86 汇编语言, 这可以在 Intel, AMD 和其他厂商的多种处理器上运行. 如果你在不同架构上运行, 你可能会得到不同的代码.

1.6 预处理

现在我们在编译过程中向后再走一步. 你可以用 `-E` 标识, 仅仅只运行预处理器:

```
$ gcc hello.c -E
```

结果是预处理器的输出. 这个例子中, 它包含了 `stdio.h` 中包含的代码, 以及 `stdio.h` 中包含的所有文件, 依此类推. 在我的机器上, 结果有 800 多行代码. 由于几乎每个 C 程序都包括 `stdio.h`, 所以这 800 行代码会被编译多次. 如果像很多 C 程序一样, 你也包含了 `stdlib.h`, 那么结果要超过 1800 行代码了.

1.7 理解错误

现在我们知道了编译过程的各个步骤, 便更容易理解错误信息了. 比如, 如果 `#include` 指令中存在错误, 你会从预处理器得到下面这条信息:

```
hello.c:1:20: fatal error: stdio.h: No such file or directory
compilation terminated.
```

如果存在语法错误, 你会从编译器得到下面信息:

```
hello.c: In function 'main':
hello.c:6:1: error: expected ';' before '}' token
```

如果使用了任何标准库都没有定义的函数, 你会从链接器得到下面错误信息:

```
/tmp/cc7iAUbN.o: In function `main':  
hello.c:(.text+0xf): undefined reference to `printf'  
collect2: error: ld returned 1 exit status
```

ld 是 UNIX 链接器的名称, 如此命名是因为 “loading” 是编译过程中和链接关系密切的另一步骤.

一旦程序运行, C 便几乎不会进行运行时检查, 所以你可能只会看到极少的运行时错误. 如果你除以零, 或者执行了其他非法的浮点操作, 你会得到一个 “Floating point exception(浮点异常).” 同时, 如果你尝试读取或写入一个错误的内存位置, 你会遇到 “Segmentation fault(段错误).”

Chapter 2

进程

2.1 抽象和虚拟化

在我们谈论进程之前,我想先定义几个词汇:

- 抽象: 抽象是某个复杂事物的简化表达. 例如, 当你开车时, 你知道当你左转方向盘时, 车会左转, 反之亦然. 当然, 方向盘连接着一系列机械和(通常是) 液压系统, 这些会使车轮转动, 同时车轮会以复杂的方式和路面互动. 作为一个司机, 你通常无需考虑这些细节. 你只需要一个简单的逻辑模型来进行驾驶. 你的这个逻辑模型便是一种抽象.

同样地, 当你使用网络浏览器时, 你知道当你点击链接时, 浏览器会显示该链接所指的页面. 使这一切成为可能的软件和网络通信是复杂的. 但作为一个用户, 你不必了解其细节.

软件工程中一个重要部分就是设计这样的抽象, 这些抽象可以使用户和其他程序员能够使用强大而复杂的系统, 而无需了解其实现的细节.

- 虚拟化: 一种重要的抽象便是虚拟化, 这是创建一个令人满意的幻觉的过程.

例如, 很多公共图书馆会参与图书馆之间的协作, 允许其相互借书. 当我请求一本图书时, 有时这本书是在我当地的图书馆的书架上, 但其他时候, 它需要从另外的地方转移过来. 无论哪种情况, 当书籍可以领取时, 我都会接到通知. 我不必知道它来自哪里, 我也不需要知道我的图书馆里有哪些书. 整个系统创建了一个假象, 让我相信我的图书馆拥有世界上的任意一本图书.

我当地的图书馆里实际可用的书可能很少, 但对我来说, 可用的虚拟书籍包括合作的图书馆中的每一本书.

再举个例子，大多数计算机只连接到一个网络，但该网络会连接到其他网络，依此类推。我们所称的互联网是一组网络和协议，它们可以将数据包从一个网络转发到下一个网络。从用户或程序员的角度来看，该系统表现得好像互联网上的每台计算机都与其他计算机相连。所以物理连接的数量虽然很少，但虚拟连接的数量是非常大的。

“虚拟”一词通常在虚拟机的上下文中使用，虚拟机是一种软件，会给你一种假象，仿佛你运行在一个运行着特定操作系统的专有机上。而实际上，虚拟机可能和很多其他虚拟机一起运行在同一台计算机上，这些虚拟机可能运行着不同的操作系统。

在虚拟化的上下文中，我们有时称真实发生的事情为“物理事件”，而虚拟发生的事情为“逻辑事件”或者“抽象事件”。

2.2 隔离

工程学最重要的一个原则便是隔离：当你设计一个包括多个部件的系统时，最好将其相互之间进行隔离，如此一个部件的改变就不会给其他组件带来不良影响。

操作系统的一个重要目标便是将每个运行的程序和其他程序进行隔离，如此程序员便无需关注任何交互的可能。提供这种隔离性的软件对象叫做**进程**。

进程是一个表示当前运行程序的软件对象。我说的“软件对象”是指面向对象编程中的对象；通常，对象包含数据以及操作数据的方法。一个进程便是一个包含下面数据的对象：

- 程序的文本，通常是一系列机器语言指令。
- 程序相关的数据，包括静态数据（编译时分配）和动态数据（运行时分配）。
- 任何未完成的输入/输出操作的状态。比如，如果进程正在等待从磁盘读取数据，或者等待网络数据到达，这些操作的状态便是进程的一部分。
- 程序的硬件状态，包括寄存器中的数据，状态信息以及程序计数器。程序计数器用来标识当前正在执行哪个指令。

通常一个进程运行一个程序，但也可能会加载并运行新的程序。

还有可能，也是很常见的情况是，多个进程中运行同一个程序。在这种情况下，这些进程会共享相同的程序代码，但通常数据和硬件状态不同。

多数操作系统会提供一组基础功能，以使进程之间相互隔离：

- 多任务: 多数操作系统都具有一种能力, 可以在任何时刻将运行中的进程中断, 并保存其硬件状态, 然后稍后恢复进程. 通常, 程序员不必考虑这些中断. 程序的行为就像在专用处理器上连续运行一样, 只是指令之间的时间是不可预测的.
- 虚拟内存: 多数操作系统都会创建这样的假象, 即每个进程都有自己的一段空间, 与其他进程都隔离开来. 同样地, 程序员一般无需关注虚拟内存如何工作; 他们可以按照每个程序都有一段专门的内存来进行操作.
- 设备抽象: 运行在同一个计算机上的进程, 共享磁盘驱动器, 网络接口, 图形卡, 以及其他硬件. 如果进程不经协商便直接和硬件交互, 便会引发混乱. 例如, 本来要传递给某个进程的网络数据, 可能被其他进程读取, 或多个进程向硬盘上同一个位置存储数据. 这便有赖于操作系统提供适当的抽象, 以维护秩序.

作为程序员, 你并不需要知道这些功能实现的详情. 但如果你有兴趣, 你会发现在各种隐喻之下, 有很多有趣的事情正在发生. 如果你知晓发生的细节, 那你将成为更好的程序员.

2.3 UNIX 进程

当写这本书时, 我最关注的进程是我的文本编辑器, emacs. 每隔一段时间, 我会切换到终端窗口. 这是一个运行 UNIX shell 的窗口, 它提供了命令行界面.

当我移动鼠标, 窗口管理器便会被唤醒, 查看鼠标是否在终端窗口上, 然后又会唤醒终端. 终端唤醒 shell. 如果我在 shell 中输入 `make`, 它会创建一个新的进程来执行 Make, 创建另一个进程运行 LaTeX, 然后再创建一个进程来显示结果.

如果我需要查找一些东西, 我可能会切换到另一个桌面, 这也会再次唤醒窗口管理器. 如果我点击网络浏览器的图标, 窗口管理器会创建一个进程, 运行网络浏览器. 有些浏览器, 比如 Chrome, 会为每个窗口和每个标签页都创建新的进程.

这些还只是我知道的进程. 同时, 还有很多其他进程在**后台**运行. 其中很多进程正在执行与操作系统相关的操作.

UNIX 命令 `ps` 可以打印正在运行的进程的信息. 如果你在终端执行它, 你可能会看到像这样的东西:

PID	TTY	TIME	CMD
2687	pts/1	00:00:00	bash
2801	pts/1	00:01:24	emacs
24762	pts/1	00:00:00	ps

第一列是唯一的数字格式进程 ID. 第二列是创建进程的终端; “TTY” 代表电传打印机, 这是最早的机械终端.

第三列是进程使用的处理器总时间, 以小时、分钟和秒为单位. 在这个例子中, `bash` 是解释输入命令的 shell 的名称, `emacs` 是我的文本编辑器, `ps` 是生成此输出的程序.

默认情况下, `ps` 仅列出和当前终端相关的进程. 如果你使用 `-e` 标志, 你便能获取每个进程 (包括属于其他用户的进程, 就我看来, 这是一个安全漏洞).

在我的系统上, 目前有 233 个进程. 下面是其中一些:

PID	TTY	TIME	CMD
1	?	00:00:17	init
2	?	00:00:00	kthreadd
3	?	00:00:02	ksoftirqd/0
4	?	00:00:00	kworker/0:0
8	?	00:00:00	migration/0
9	?	00:00:00	rcu_bh
10	?	00:00:16	rcu_sched
47	?	00:00:00	cpuset
48	?	00:00:00	khelper
49	?	00:00:00	kdevtmpfs
50	?	00:00:00	netns
51	?	00:00:00	bdi-default
52	?	00:00:00	kintegrityd
53	?	00:00:00	kblockd
54	?	00:00:00	ata_sff
55	?	00:00:00	khubd
56	?	00:00:00	md
57	?	00:00:00	devfreq_wq

`init` 是操作系统启动时创建的第一个进程. 它会创建许多其他进程, 然后闲置, 直到它创建的进程都结束.

`kthreadd` 是操作系统用来创建**线程**的进程. 后续我们将更详细地讨论线程, 当前你可以将线程看作是一种进程. 开头的 `k` 代表**内核 (kernel)**, 这是操作系统符合核心功能 (比如创建线程) 的部分. 结尾的 `d` 代表**守护程序 (daemon)**, 这是运行在后台, 并提供操作系统服务的进程的另一个名称. 在此, “daemon” 是指一种乐于助人的精神, 没有任何邪恶的意思.

根据名称, 你可以推断出 `ksoftirqd` 也是一个内核守护程序; 具体来说, 它处理软件中断请求, 即 “软中断 (soft IRQ)”

`kworker` 是一个内核创建的工作进程, 用来为内核处理一些任务.

通常有多个进程运行着内核服务. 在我的系统上, 现在有 8 个 `ksoftirqd` 进程和 35 个 `kworker` 进程.

我不会详细讨论其他进程, 但如果你感兴趣, 你可以搜索更多关于它们的信息. 你应该在你的系统上运行 `ps`, 并将结果和我的进行对比.

Chapter 3

虚拟内存

3.1 一点信息论知识

比特 (bit) 是一个二进制的数位, 也是一个信息单元. 如果你有一比特, 便可以表示两种可能性中的一种, 通常写作 0 和 1. 如果你有两比特, 那边有 4 种组合, 分别是 00, 01, 10, 和 11. 通常, 如果你有 b 个比特, 你便可以表示 2^b 个值. 一个**字节 (byte)** 是 8 个比特, 所以可以表示 256 个值.

反过来想, 假设你想储存一个字母. 字母表有 26 个字母, 那么你需要几个比特呢? 如果用 4 个比特, 你可以表示 16 个值, 数量不够. 如果用 5 个比特, 你可以表示多达 32 个值, 这足够表示所有字母了, 而且还有剩余.

通常, 如果你想表示 N 个值, 你应该选择满足 $2^b \geq N$ 的最小 b 值. 两侧以 2 为底取对数, 便可以得到 $b \geq \log_2 N$.

假设我掷一枚硬币, 并告诉你结果, 那么我便是给了你一比特的信息. 如果我掷一个六面骰子, 并告诉你结果, 那我便是给了你 $\log_2 6$ 个比特的信息. 通常, 如果发生的概率是 N 中的 1 个, 那么结果便包含 $\log_2 N$ 个比特信息.

同样的, 如果结果出现的概率是 p , 那么信息量便是 $-\log_2 p$. 这个数量也被称为结果的**自信息**, 它是用来衡量结果的惊奇程度, 所以也被称为**惊奇度**. 如果你的马仅有 $1/16$ 的概率获胜, 但是赢了, 你可以 4 比特的信息 (以及赔付). 但如果获胜的机率是 75%, 那么胜利的消息只包含 0.42 比特.

直观来说, 出乎意料的消息反而包含了大量的信息. 相反, 如果你已经确定某件事情, 那么证实它仅会提供少量信息.

在本书的诸多主题中, 我们需要在比特数, b , 和其能编码的值的数量 $N = 2^b$ 之间熟练转换.

3.2 内存和存储

进程运行时，其大部分数据存储在**主内存**，这通常是一种随机访问存储 (RAM)。当前多数计算机上，主存储都是**易失的**，也就意味着一旦计算机关机，主存储的内容便会丢失。一台典型的桌面计算机有 2–8 GiB 内存。GiB 表示 “gibibyte,” 也就是 2^{30} 字节。

进程需要读写文件，而文件往往存储在硬盘驱动器 (HDD) 或者固态硬盘中。这些存储设备都是**非易失的**，因此可以用作长期存储。目前一台典型桌面计算机会有 500GB 到 2TB 容量的硬盘。其中 GB 表示 “gigabyte,” 也就是 10^9 字节。TB 表示 “terabyte,” 也就是 10^{12} 字节。

你可能已经注意到，我描述主内存大小时，用二进制单位 GiB，描述硬盘大小时，用十进制单位 GB 和 TB。由于一些历史和科技的原因，内存通常用二进制单位衡量，而硬盘则用十进制单位。在本书中，我将谨慎区分二进制和十进制单位，但你也应该注意到，“gigabyte” 和缩写 GB 通常是混用的。

在非正式的使用中，“内存” 一词有时候会指 HDD 和 SSD，以及 RAM。但这些设备的属性是不同的，所以我们需要仔细区分它们。我将使用**存储**来指代 HDD 和 SSD。

3.3 地址空间

主内存中的每个字节都由一个整数**物理地址**指定。有效的物理地址集被称为**物理地址空间**。通常从 0 到 $N - 1$ ，其中 N 是内存的大小。在 1 GiB 物理内存的系统上，最高的有效地址是 $2^{30} - 1$ ，即十进制下的 1,073,741,823，或者十六进制下的 0x3fff ffff (前缀的 0x 表示十六进制数)。

然而，大部分的操作系统会提供**虚拟内存**，这意味着程序永远不必处理物理地址，也不必知道由多少物理地址可用。

相反，程序会使用从 0 到 $M - 1$ 编号的**虚拟地址**，其中 M 是有效虚拟地址的数量。虚拟地址空间的大小是由操作系统和其运行的硬件共同决定的。

你可能听别人谈论 32 位系统和 64 位系统。这些术语表示的是寄存器的大小，通常也是虚拟地址的大小。在 32 位系统上，虚拟地址为 32 位，也就是虚拟地址空间从 0 到 0xffff ffff。所以地址空间的大小为 2^{32} 字节，或 4 GiB。

而在 64 位系统上，虚拟地址空间是 2^{64} 字节，或 $2^4 \cdot 1024^6$ 字节。这是 16 exbibytes，约为当前物理内存的十亿倍。虚拟地址空间比物理内存大这么多，这看起来很奇怪，但是我们稍后就会看到它是如何工作的。

当程序从内存读取或者向内存写入值时，它会生成虚拟地址。在操作系统的帮助下，访问内存之前，硬件会被转换成物理地址。这种转换是基于每个进程

进行的, 所以即使两个进程生成同样的虚拟地址, 它们也会映射到物理内存的不同位置.

最后, 虚拟内存是操作系统隔离进程的一种重要手段. 通常, 一个进程是无法访问另一个进程拥有的数据的, 因为如果某块物理内存已经被分配给了其他进程, 那么该进程便无法生成映射到这块物理内存的虚拟地址.

3.4 内存段

The data of a running process is organized into five segments:

- The **code segment** contains the program text; that is, the machine language instructions that make up the program.
- The **static segment** contains immutable values, like string literals. For example, if your program contains the string "Hello, World", those characters will be stored in the static segment.
- The **global segment** contains global variables and local variables that are declared **static**.
- The **heap segment** contains chunks of memory allocated at run time, most often by calling the C library function `malloc`.
- The **stack segment** contains the call stack, which is a sequence of stack frames. Each time a function is called, a stack frame is allocated to contain the parameters and local variables of the function. When the function completes, its stack frame is removed from the stack.

The arrangement of these segments is determined partly by the compiler and partly by the operating system. The details vary from one system to another, but in the most common arrangement:

- The text segment is near the “bottom” of memory, that is, at addresses near 0.
- The static segment is often just above the text segment, that is, at higher addresses.
- The global segment is often just above the static segment.
- The heap is often above the global segment. As it expands, it grows up toward larger addresses.

- The stack is near the top of memory; that is, near the highest addresses in the virtual address space. As the stack expands, it grows down toward smaller addresses.

To determine the layout of these segments on your system, try running this program, which is in `aspace.c` in the repository for this book (see Section 0.1).

```
#include <stdio.h>
#include <stdlib.h>

int global;

int main ()
{
    int local = 5;
    void *p = malloc(128);
    char *s = "Hello, World";

    printf ("Address of main is %p\n", main);
    printf ("Address of global is %p\n", &global);
    printf ("Address of local is %p\n", &local);
    printf ("p points to %p\n", p);
    printf ("s points to %p\n", s);
}
```

`main` is the name of a function; when it is used as a variable, it refers to the address of the first machine language instruction in `main`, which we expect to be in the text segment.

`global` is a global variable, so we expect it to be in the global segment. `local` is a local variable, so we expect it to be on the stack.

`s` refers to a “string literal”, which is a string that appears as part of the program (as opposed to a string that is read from a file, input by a user, etc.). We expect the location of the string to be in the static segment (as opposed to the pointer, `s`, which is a local variable).

`p` contains an address returned by `malloc`, which allocates space in the heap. “malloc” stands for “memory allocate.”

The format sequence `%p` tells `printf` to format each address as a “pointer”, so it displays the results in hexadecimal.

When I run this program, the output looks like this (I added spaces to make it easier to read):

```
Address of main is 0x      40057d
Address of global is 0x     60104c
Address of local is 0x7ffe6085443c
p points to      0x      16c3010
s points to      0x      4006a4
```

As expected, the address of `main` is the lowest, followed by the location of the string literal. The location of `global` is next, then the address `p` points to. The address of `local` is much bigger.

The largest address has 12 hexadecimal digits. Each hex digit corresponds to 4 bits, so it is a 48-bit address. That suggests that the usable part of the virtual address space is 2^{48} bytes.

As an exercise, run this program on your computer and compare your results to mine. Add a second call to `malloc` and check whether the heap on your system grows up (toward larger addresses). Add a function that prints the address of a local variable, and check whether the stack grows down.

3.5 Static local variables

Local variables on the stack are sometimes called **automatic**, because they are allocated automatically when a function is called, and freed automatically when the function returns.

In C there is another kind of local variable, called **static**, which is allocated in the global segment. It is initialized when the program starts and keeps its value from one function call to the next.

For example, the following function keeps track of how many times it has been called.

```
int times_called()
{
    static int counter = 0;
    counter++;
    return counter;
}
```

The keyword `static` indicates that `counter` is a static local variable. The initialization happens only once, when the program starts.

If you add this function to `aspace.c` you can confirm that `counter` is allocated in the global segment along with global variables, not in the stack.



图 3.1: Diagram of the address translation process.

3.6 Address translation

How does a virtual address (VA) get translated to a physical address (PA)? The basic mechanism is simple, but a simple implementation would be too slow and take too much space. So actual implementations are a bit more complicated.

Most processors provide a memory management unit (MMU) that sits between the CPU and main memory. The MMU performs fast translation between VAs and PAs.

1. When a program reads or writes a variable, the CPU generates a VA.
2. The MMU splits the VA into two parts, called the page number and the offset. A “page” is a chunk of memory; the size of a page depends on the operating system and the hardware, but common sizes are 1–4 KiB.
3. The MMU looks up the page number in the translation lookaside buffer (TLB) and gets the corresponding physical page number. Then it combines the physical page number with the offset to produce a PA.
4. The PA is passed to main memory, which reads or writes the given location.

The TLB contains cached copies of data from the page table (which is stored in kernel memory). The page table contains the mapping from virtual page numbers to physical page numbers. Since each process has its own page table, the TLB has to make sure it only uses entries from the page table of the process that’s running.

Figure 3.1 shows a diagram of this process. To see how it all works, suppose that the VA is 32 bits and the physical memory is 1 GiB, divided into 1 KiB pages.

- Since 1 GiB is 2^{30} bytes and 1 KiB is 2^{10} bytes, there are 2^{20} physical pages, sometimes called “frames.”
- The size of the virtual address space is 2^{32} B and the size of a page is 2^{10} B, so there are 2^{22} virtual pages.
- The size of the offset is determined by the page size. In this example the page size is 2^{10} B, so it takes 10 bits to specify a byte on a page.
- If a VA is 32 bits and the offset is 10 bits, the remaining 22 bits make up the virtual page number.
- Since there are 2^{20} physical pages, each physical page number is 20 bits. Adding in the 10 bit offset, the resulting PAs are 30 bits.

So far this all seems feasible. But let’s think about how big a page table might have to be. The simplest implementation of a page table is an array with one entry for each virtual page. Each entry would contain a physical page number, which is 20 bits in this example, plus some additional information about each frame. So we expect 3–4 bytes per entry. But with 2^{22} virtual pages, the page table would require 2^{24} bytes, or 16 MiB.

And since we need a page table for each process, a system running 256 processes would need 2^{32} bytes, or 4 GiB, just for page tables! And that’s just with 32-bit virtual addresses. With 48- or 64-bit VAs, the numbers are ridiculous.

Fortunately, we don’t actually need that much space, because most processes don’t use even a small fraction of their virtual address space. And if a process doesn’t use a virtual page, we don’t need an entry in the page table for it.

Another way to say the same thing is that page tables are “sparse”, which implies that the simple implementation, an array of page table entries, is a bad idea. Fortunately, there are several good implementations for sparse arrays.

One option is a multilevel page table, which is what many operating systems, including Linux, use. Another option is an associative table, where each entry includes both the virtual page number and the physical page number. Searching an associative table can be slow in software, but in hardware we can search the entire table in parallel, so associative arrays are often used to represent the page table entries in the TLB.

You can read more about these implementations at http://en.wikipedia.org/wiki/Page_table; you might find the details interesting. But the fundamental idea is that page tables are sparse, so we have to choose a good implementation for sparse arrays.

I mentioned earlier that the operating system can interrupt a running process, save its state, and then run another process. This mechanism is called a **context switch**. Since each process has its own page table, the operating system has to work with the MMU to make sure each process gets the right page table. In older machines, the page table information in the MMU had to be replaced during every context switch, which was expensive. In newer systems, each page table entry in the MMU includes the process ID, so page tables from multiple processes can be in the MMU at the same time.

Chapter 4

Files and file systems

When a process completes (or crashes), any data stored in main memory is lost. But data stored on a hard disk drive (HDD) or solid state drive (SSD) is “persistent;” that is, it survives after the process completes, even if the computer shuts down.

Hard disk drives are complicated. Data is stored in blocks, which are laid out in sectors, which make up tracks, which are arranged in concentric circles on platters.

Solid state drives are simpler in one sense, because blocks are numbered sequentially, but they raise a different complication: each block can be written a limited number of times before it becomes unreliable.

As a programmer, you don’t want to deal with these complications. What you want is an appropriate abstraction of persistent storage hardware. The most common abstraction is called a “file system.”

Abstractly:

- A “file system” is a mapping from each file’s name to its contents. If you think of the names as keys, and the contents as values, a file system is a kind of key-value database (see https://en.wikipedia.org/wiki/Key-value_database).
- A “file” is a sequence of bytes.

File names are usually strings, and they are usually “hierarchical”; that is, the string specifies a path from a top-level directory (or folder), through a series of subdirectories, to a specific file.

The primary difference between the abstraction and the underlying mechanism is that files are byte-based and persistent storage is block-based. The operating system translates byte-based file operations in the C library into block-based operations on storage devices. Typical block sizes are 1–8 KiB.

For example, the following code opens a file and reads the first byte:

```
FILE *fp = fopen("/home/downey/file.txt", "r");
char c = fgetc(fp);
fclose(fp);
```

When this code runs:

1. `fopen` uses the filename to find the top-level directory, called `/`, the subdirectory `home`, and the sub-subdirectory `downey`.
2. It finds the file named `file.txt` and “opens” it for reading, which means it creates a data structure that represents the file being read. Among other things, this data structure keeps track of how much of the file has been read, called the “file position”.

In DOS, this data structure is called a File Control Block, but I want to avoid that term because in UNIX it means something else. In UNIX, there seems to be no good name for it. It is an entry in the open file table, so I will call it an `OpenFileTableEntry`.

3. When we call `fgetc`, the operating system checks whether the next character of the file is already in memory. If so, it reads the next character, advances the file position, and returns the result.
4. If the next character is not in memory, the operating system issues an I/O request to get the next block. Disk drives are slow, so a process waiting for a block from disk is usually interrupted so another process can run until the data arrives.
5. When the I/O operation is complete, the new block of data is stored in memory, and the process resumes. It reads the first character and stores it as a local variable.
6. When the process closes the file, the operating system completes or cancels any pending operations, removes data stored in memory, and frees the `OpenFileTableEntry`.

The process for writing a file is similar, but there are some additional steps. Here is an example that opens a file for writing and changes the first character.

```
FILE *fp = fopen("/home/downey/file.txt", "w");  
fputc('b', fp);  
fclose(fp);
```

When this code runs:

1. Again, `fopen` uses the filename to find the file. If it does not already exist, it creates a new file and adds an entry in the parent directory, `/home/downey`.
2. The operating system creates an `OpenFileTableEntry` that indicates that the file is open for writing, and sets the file position to 0.
3. `fputc` attempts to write (or re-write) the first byte of the file. If the file already exists, the operating system has to load the first block into memory. Otherwise it allocates a new block in memory and requests a new block on disk.
4. After the block in memory is modified, it might not be copied back to the disk right away. In general, data written to a file is “buffered”, which means it is stored in memory and only written to disk when there is at least one block to write.
5. When the file is closed, any buffered data is written to disk and the `OpenFileTableEntry` is freed.

To summarize, the C library provides the abstraction of a file system that maps from file names to streams of bytes. This abstraction is built on top of storage devices that are actually organized in blocks.

4.1 Disk performance

I mentioned earlier that disk drives are slow. On current HDDs, the average time to read a block from disk to memory might be 5–25 ms (see https://en.wikipedia.org/wiki/Hard_disk_drive_performance_characteristics). SSDs are faster, taking 25 μ s to read a 4 KiB block and 250 μ s to write one (see <http://en.wikipedia.org/wiki/Ssd#Controller>).

To put these numbers in perspective, let’s compare them to the clock cycle of the CPU. A processor with clock rate 2 GHz completes one clock cycle every 0.5 ns. The time to get a byte from memory to the CPU is typically around

100 ns. If the processor completes one instruction per clock cycle, it would complete 200 instructions while waiting for a byte from memory.

In one microsecond, it would complete 2000 instructions, so while waiting 25 μ s for a byte from an SSD, it would complete 50,000.

In one millisecond, it would complete 2,000,000 instructions, so while waiting 20 ms for a byte from a HDD, it might complete 40 million. If there's nothing for the CPU to do while it waits, it would be idle. That's why the operating system generally switches to another process while it is waiting for data from disk.

The gap in performance between main memory and persistent storage is one of the major challenges of computer system design. Operating systems and hardware provide several features intended to “fill in” this gap:

- **Block transfers:** The time it takes to load a single byte from disk is 5–25 ms. By comparison, the additional time to load an 8 KiB block is negligible. So systems generally try to read large blocks each time they access the disk.
- **Prefetching:** Sometimes the operating system can predict that a process will read a block and start loading it before it is requested. For example, if you open a file and read the first block, there is a good chance you will go on to read the second block. The operating system might start loading additional blocks before they are requested.
- **Buffering:** As I mentioned, when you write a file, the operating system stores the data in memory and only writes it to disk later. If you modify the block several times while it is in memory, the system only has to write it to disk once.
- **Caching:** If a process has used a block recently, it is likely to use it again soon. If the operating system keeps a copy of the block in memory, it can handle future requests at memory speed.

Some of these features are also implemented in hardware. For example, some disk drives provide a cache that stores recently-used blocks, and many disk drives read more than one block at a time, even if only one is requested.

These mechanisms generally improve the performance of programs, but they don't change the behavior. Usually programmers don't have to think about them, with two exceptions: (1) if the performance of a program is unexpectedly bad, you might have to know something about these mechanisms to diagnose

the problem, and (2) when data is buffered, it can be harder to debug a program. For example, if a program prints a value and then crashes, the value might not appear, because it might be in a buffer. Similarly, if a program writes data to disk and then the computer loses power, the data might be lost if it is in a cache and not yet on disk.

4.2 Disk metadata

The blocks that make up a file might be arranged contiguously on disk, and file system performance is generally better if they are, but most operating systems don't require contiguous allocation. They are free to place a block anywhere on disk, and they use various data structures to keep track of them.

In many UNIX file systems, that data structure is called an “inode,” which stands for “index node”. More generally, information about files, including the location of their blocks, is called “metadata”. (The content of the file is data, so information about the file is data about data, hence “meta”.)

Since inodes reside on disk along with the rest of the data, they are designed to fit neatly into disk blocks. A UNIX inode contains information about a file, including the user ID of the file owner; permission flags indicating who is allowed to read, write, or execute it; and timestamps that indicate when it was last modified and accessed. In addition, it contains block numbers for the first 12 blocks that make up the file.

If the block size is 8 KiB, the first 12 blocks make up 96 KiB. On most systems, that's big enough for a large majority of files, but it's definitely not big enough for all of them. That's why the inode also contains a pointer to an “indirection block”, which contains nothing but pointers to other blocks.

The number of pointers in an indirection block depends on the sizes of the blocks and the block numbers, but it is often 1024. With 1024 block numbers and 8 KiB blocks, an indirection block can address 8 MiB. That's big enough for all but the largest files, but still not big enough for all.

That's why the inode also contains a pointer to a “double indirection block”, which contains pointers to indirection blocks. With 1024 indirection blocks, we can address 8 GiB.

And if that's not big enough, there is (finally) a triple indirection block, which contains pointers to double indirection blocks, yielding a maximum file size of 8 TiB. When UNIX inodes were designed, that seemed big enough to serve for a long time. But that was a long time ago.

As an alternative to indirection blocks, some file systems, like FAT, use a File Allocation Table that contains one entry for each block, called a “cluster” in this context. A root directory contains a pointer to the first cluster in each file. The FAT entry for each cluster points to the next cluster in the file, similar to a linked list. For more details, see http://en.wikipedia.org/wiki/File_Allocation_Table.

4.3 Block allocation

File systems have to keep track of which blocks belong to each file; they also have to keep track of which blocks are available for use. When a new file is created, the file system finds an available block and allocates it. When a file is deleted, the file system makes its blocks available for re-allocation.

The goals of the block allocation system are:

- Speed: Allocating and freeing blocks should be fast.
- Minimal space overhead: The data structures used by the allocator should be small, leaving as much space as possible for data.
- Minimal fragmentation: If some blocks are left unused, or some are only partially used, the unused space is called “fragmentation”.
- Maximum contiguity: Data that is likely to be used at the same time should be physically contiguous, if possible, to improve performance.

It is hard to design a file system that achieves all of these goals, especially since file system performance depends on “workload characteristics” like file sizes, access patterns, etc. A file system that is well tuned for one workload might not perform as well for another.

For this reason, most operating systems support several kinds of file systems, and file system design is an active area of research and development. In the last decade, Linux systems have migrated from ext2, which was a conventional UNIX file system, to ext3, a “journaling” file system intended to improve speed and contiguity, and more recently to ext4, which can handle larger files and file systems. Within the next few years, there might be another migration to the B-tree file system, Btrfs.

4.4 Everything is a file?

The file abstraction is really a “stream of bytes” abstraction, which turns out to be useful for many things, not just file systems.

One example is the UNIX pipe, which is a simple form of inter-process communication. Processes can be set up so that output from one process is taken as input into another process. For the first process, the pipe behaves like a file open for writing, so it can use C library functions like `fputs` and `fprintf`. For the second process, the pipe behaves like a file open for reading, so it uses `fgets` and `fscanf`.

Network communication also uses the stream of bytes abstraction. A UNIX socket is a data structure that represents a communication channel between processes on different computers (usually). Again, processes can read data from and write data to a socket using “file” handling functions.

Reusing the file abstraction makes life easier for programmers, since they only have to learn one API (application program interface). It also makes programs more versatile, since a program intended to work with files can also work with data coming from pipes and other sources.

Chapter 5

More bits and bytes

5.1 Representing integers

You probably know that computers represent numbers in base 2, also known as binary. For positive numbers, the binary representation is straightforward; for example, the representation for 5_{10} is $b101$.

For negative numbers, the most obvious representation uses a sign bit to indicate whether a number is positive or negative. But there is another representation, called “two’s complement” that is much more common because it is easier to work with in hardware.

To find the two’s complement of a negative number, $-x$, find the binary representation of x , flip all the bits, and add 1. For example, to represent -5_{10} , start with the representation of 5_{10} , which is $b00000101$ if we write the 8-bit version. Flipping all the bits and adding 1 yields $b11111011$.

In two’s complement, the leftmost bit acts like a sign bit; it is 0 for positive numbers and 1 for negative numbers.

To convert from an 8-bit number to 16-bits, we have to add more 0’s for a positive number and add 1’s for a negative number. In effect, we have to copy the sign bit into the new bits. This process is called “sign extension”.

In C all integer types are signed (able to represent positive and negative numbers) unless you declare them **unsigned**. The difference, and the reason this declaration is important, is that operations on unsigned integers don’t use sign extension.

5.2 Bitwise operators

People learning C are sometimes confused about the bitwise operators `&` and `|`. These operators treat integers as bit vectors and compute logical operations on corresponding bits.

For example, `&` computes the AND operation, which yields 1 if both operands are 1, and 0 otherwise. Here is an example of `&` applied to two 4-bit numbers:

```

  1100
& 1010
----
  1000

```

In C, this means that the expression `12 & 10` has the value 8.

Similarly, `|` computes the OR operation, which yields 1 if either operand is 1, and 0 otherwise.

```

  1100
| 1010
----
  1110

```

So the expression `12 | 10` has the value 14.

Finally, `^` computes the XOR operation, which yields 1 if either operand is 1, but not both.

```

  1100
^ 1010
----
  0110

```

So the expression `12 ^ 10` has the value 6.

Most commonly, `&` is used to clear a set of bits from a bit vector, `|` is used to set bits, and `^` is used to flip, or “toggle” bits. Here are the details:

Clearing bits: For any value x , $x \& 0$ is 0, and $x \& 1$ is x . So if you AND a vector with 3, it selects only the two rightmost bits, and sets the rest to 0.

```

  xxxx
& 0011
----
  00xx

```

In this context, the value 3 is called a “mask” because it selects some bits and masks the rest.

Setting bits: Similarly, for any x , $x|0$ is x , and $x|1$ is 1. So if you OR a vector with 3, it sets the rightmost bits, and leaves the rest alone:

```

xxxx
| 0011
----
xx11

```

Toggling bits: Finally, if you XOR a vector with 3, it flips the rightmost bits and leaves the rest alone. As an exercise, see if you can compute the two’s complement of 12 using \wedge . Hint: what’s the two’s complement representation of -1?

C also provides shift operators, \ll and \gg , which shift bits left and right. Each left shift doubles a number, so $5 \ll 1$ is 10, and $5 \ll 2$ is 20. Each right shift divides by two (rounding down), so $5 \gg 1$ is 2 and $2 \gg 1$ is 1.

5.3 Representing floating-point numbers

Floating-point numbers are represented using the binary version of scientific notation. In decimal notation, large numbers are written as the product of a coefficient and 10 raised to an exponent. For example, the speed of light in m/s is approximately $2.998 \cdot 10^8$.

Most computers use the IEEE standard for floating-point arithmetic. The C type `float` usually corresponds to the 32-bit IEEE standard; `double` usually corresponds to the 64-bit standard.

In the 32-bit standard, the leftmost bit is the sign bit, s . The next 8 bits are the exponent, q , and the last 23 bits are the coefficient, c . The value of a floating-point number is

$$(-1)^s c \cdot 2^q$$

Well, that’s almost correct, but there’s one more wrinkle. Floating-point numbers are usually normalized so that there is one digit before the point. For example, in base 10, we prefer $2.998 \cdot 10^8$ rather than $2998 \cdot 10^5$ or any other equivalent expression. In base 2, a normalized number always has the digit 1 before the binary point. Since the digit in this location is always 1, we can save space by leaving it out of the representation.

For example, the integer representation of 13_{10} is $b1101$. In floating point, that's $1.101 \cdot 2^3$, so the exponent is 3 and the part of the coefficient that would be stored is 101 (followed by 20 zeros).

Well, that's almost correct, but there's one more wrinkle. The exponent is stored with a "bias". In the 32-bit standard, the bias is 127, so the exponent 3 would be stored as 130.

To pack and unpack floating-point numbers in C, we can use a union and bitwise operations. Here's an example:

```
union {
    float f;
    unsigned int u;
} p;

p.f = -13.0;
unsigned int sign = (p.u >> 31) & 1;
unsigned int exp = (p.u >> 23) & 0xff;

unsigned int coef_mask = (1 << 23) - 1;
unsigned int coef = p.u & coef_mask;

printf("%d\n", sign);
printf("%d\n", exp);
printf("0x%x\n", coef);
```

This code is in `float.c` in the repository for this book (see Section 0.1).

The union allows us to store a floating-point value using `p.f` and then read it as an unsigned integer using `p.u`.

To get the sign bit, we shift the bits to the right 31 places and then use a 1-bit mask to select only the rightmost bit.

To get the exponent, we shift the bits 23 places, then select the rightmost 8 bits (the hexadecimal value `0xff` has eight 1's).

To get the coefficient, we need to extract the 23 rightmost bits and ignore the rest. We do that by making a mask with 1s in the 23 rightmost places and 0s on the left. The easiest way to do that is by shifting 1 to the left by 23 places and then subtracting 1.

The output of this program is:

1

```
130
0x500000
```

As expected, the sign bit for a negative number is 1. The exponent is 130, including the bias. And the coefficient, which I printed in hexadecimal, is 101 followed by 20 zeros.

As an exercise, try assembling or disassembling a `double`, which uses the 64-bit standard. See http://en.wikipedia.org/wiki/IEEE_floating_point.

5.4 Unions and memory errors

There are two common uses of C unions. One, which we saw in the previous section, is to access the binary representation of data. Another is to store heterogeneous data. For example, you could use a union to represent a number that might be an integer, float, complex, or rational number.

However, unions are error-prone. It is up to you, as the programmer, to keep track of what type of data is in the union; if you write a floating-point value and then interpret it as an integer, the result is usually nonsense.

Actually, the same thing can happen if you read a location in memory incorrectly. One way that can happen is if you read past the end of an array.

To see what happens, I'll start with a function that allocates an array on the stack and fills it with the numbers from 0 to 99.

```
void f1() {
    int i;
    int array[100];

    for (i=0; i<100; i++) {
        array[i] = i;
    }
}
```

Next I'll define a function that creates a smaller array and deliberately accesses elements before the beginning and after the end:

```
void f2() {
    int x = 17;
    int array[10];
    int y = 123;

    printf("%d\n", array[-2]);
}
```

```
    printf("%d\n", array[-1]);  
    printf("%d\n", array[10]);  
    printf("%d\n", array[11]);  
}
```

If I call `f1` and then `f2`, I get these results:

```
17  
123  
98  
99
```

The details here depend on the compiler, which arranges variables on the stack. From these results, we can infer that the compiler put `x` and `y` next to each other, “below” the array (at a lower address). And when we read past the array, it looks like we are getting values that were left on the stack by the previous function call.

In this example, all of the variables are integers, so it is relatively easy to figure out what is going on. But in general when you read beyond the bounds of an array, the values you read might have any type. For example, if I change `f1` to make an array of floats, the results are:

```
17  
123  
1120141312  
1120272384
```

The latter two values are what you get if you interpret a floating-point value as an integer. If you encountered this output while debugging, you would have a hard time figuring out what’s going on.

5.5 Representing strings

Related issues sometimes come up with strings. First, remember that C strings are null-terminated. When you allocate space for a string, don’t forget the extra byte at the end.

Also, the letters *and numbers* in C strings are encoded in ASCII. The ASCII codes for the digits “0” through “9” are 48 through 57, *not* 0 through 9. The ASCII code 0 is the NUL character that marks the end of a string. And the ASCII codes 1 through 9 are special characters used in some communication protocols. ASCII code 7 is a bell; on some terminals, printing it makes a sound.

The ASCII code for the letter “A” is 65; the code for “a” is 97. Here are those codes in binary:

65 = b0100 0001

97 = b0110 0001

A careful observer will notice that they differ by a single bit. And this pattern holds for the rest of the letters; the sixth bit (counting from the right) acts as a “case bit”, 0 for upper-case letters and 1 for lower case letters.

As an exercise, write a function that takes a string and converts from lower-case to upper-case by flipping the sixth bit. As a challenge, you can make a faster version by reading the string 32 or 64 bits at a time, rather than one character at a time. This optimization is made easier if the length of the string is a multiple of 4 or 8 bytes.

If you read past the end of a string, you are likely to see strange characters. Conversely, if you write a string and then accidentally read it as an int or float, the results will be hard to interpret.

For example, if you run:

```
char array[] = "allen";  
float *p = array;  
printf("%f\n", *p);
```

You will find that the ASCII representation of the first 8 characters of my name, interpreted as a double-precision floating point number, is 69779713878800585457664.

Chapter 6

Memory management

C provides 4 functions for dynamic memory allocation:

- `malloc`, which takes an integer size, in bytes, and returns a pointer to a newly-allocated chunk of memory with (at least) the given size. If it can't satisfy the request, it returns the special pointer value `NULL`.
- `calloc`, which is the same as `malloc` except that it also clears the newly allocated chunk; that is, it sets all bytes in the chunk to 0.
- `free`, which takes a pointer to a previously allocated chunk and deallocates it; that is, it makes the space available for future allocation.
- `realloc`, which takes a pointer to a previously allocated chunk and a new size. It allocates a chunk of memory with the new size, copies data from the old chunk to the new, frees the old chunk, and returns a pointer to the new chunk.

This API is notoriously error-prone and unforgiving. Memory management is one of the most challenging parts of designing large software systems, which is why most modern languages provide higher-level memory management features like garbage collection.

6.1 Memory errors

The C memory management API is a bit like Jasper Beardly, a minor character on the animated television program *The Simpsons*; in a few episodes, he appears as a strict substitute teacher who imposes corporal punishment — a “paddlin” — for all infractions.

Here are some of things a program can do that deserve a paddling:

- If you access (read or write) any chunk that has not been allocated, that's a paddling.
- If you free an allocated chunk and then access it, that's a paddling.
- If you try to free a chunk that has not been allocated, that's a paddling.
- If you free the same chunk more than once, that's a paddling.
- If you call `realloc` with a chunk that was not allocated, or was allocated and then freed, that's a paddling.

It might not sound difficult to follow these rules, but in a large program a chunk of memory might be allocated in one part of the program, used in several other parts, and freed in yet another part. So changes in one part of the program can require changes in many other parts.

Also, there might be many aliases, or references to the same allocated chunk, in different parts of the program. The chunk should not be freed until all references to the chunk are no longer in use. Getting this right often requires careful analysis across all parts of the program, which is difficult and contrary to fundamental principles of good software engineering.

Ideally, every function that allocates memory should include, as part of the documented interface, information about how that memory is supposed to be freed. Mature libraries often do this well, but in the real world, software engineering practice often falls short of this ideal.

To make matters worse, memory errors can be difficult to find because the symptoms are unpredictable. For example:

- If you read a value from an unallocated chunk, the system *might* detect the error, trigger a runtime error called a “segmentation fault”, and stop the program. Or, the program might read unallocated memory without detecting the error; in that case, the value it gets is whatever happened to be stored at the accessed location, which is unpredictable, and might be different each time the program runs.
- If you write a value to an unallocated chunk, and don't get a segmentation fault, things are even worse. After you write a value to an invalid location, a long time might pass before it is read and causes problems. At that point it will be very difficult to find the source of the problem.

And things can be even worse than that! One of the most common problems with C-style memory management is that the data structures used to implement `malloc` and `free` (which we will see soon) are often stored along with the allocated chunks. So if you accidentally write past the end of a dynamically-allocated chunk, you are likely to mangle these data structures. The system usually won't detect the problem until later, when you call `malloc` or `free`, and those functions fail in some inscrutable way.

One conclusion you should draw from this is that safe memory management requires design and discipline. If you write a library or module that allocates memory, you should also provide an interface to free it, and memory management should be part of the API design from the beginning.

If you use a library that allocates memory, you should be disciplined in your use of the API. For example, if the library provides functions to allocate and deallocate storage, you should use those functions and not, for example, call `free` on a chunk you did not `malloc`. And you should avoid keeping multiple references to the same chunk in different parts of your program.

Often there is a trade-off between safe memory management and performance. For example, the most common source of memory errors is writing beyond the bounds of an array. The obvious remedy for this problem is bounds checking; that is, every access to the array should check whether the index is out of bounds. High-level libraries that provide array-like structures usually perform bounds checking. But C arrays and most low-level libraries do not.

6.2 Memory leaks

There is one more memory error that may or may not deserve a paddling. If you allocate a chunk of memory and never free it, that's a "memory leak".

For some programs, memory leaks are ok. For example, if your program allocates memory, performs computations on it, and then exits, it is probably not necessary to free the allocated memory. When the program exits, all of its memory is deallocated by the operating system. Freeing memory immediately before exiting might feel more responsible, but it is mostly a waste of time.

But if a program runs for a long time and leaks memory, its total memory use will increase indefinitely. At that point, a few things might happen:

- At some point, the system runs out of physical memory. On systems without virtual memory, the next call to `malloc` will fail, returning `NULL`.

- On systems with virtual memory, the operating system can move another process's pages from memory to disk and then allocate more space to the leaking process. I explain this mechanism in Section 7.8.
- There might be a limit on the amount of space a single process can allocate; beyond that, `malloc` returns `NULL`.
- Eventually, a process might fill its virtual address space (or the usable part). After that, there are no more addresses to allocate, so `malloc` returns `NULL`.

If `malloc` returns `NULL`, but you persist and access the chunk you think you allocated, you get a segmentation fault. For this reason, it is considered good style to check the result from `malloc` before using it. One option is to add a condition like this after every `malloc` call:

```
void *p = malloc(size);
if (p == NULL) {
    perror("malloc failed");
    exit(-1);
}
```

`perror` is declared in `stdio.h`; it prints an error message and additional information about the last error that occurred.

`exit`, which is declared in `stdlib.h`, causes the process to terminate. The argument is a status code that indicates how the process terminated. By convention, status code 0 indicates normal termination and -1 indicates an error condition. Sometimes other codes are used to indicate different error conditions.

Error-checking code can be a nuisance, and it makes programs harder to read. You can mitigate these problems by wrapping library function calls and their error-checking code in your own functions. For example, here is a `malloc` wrapper that checks the return value.

```
void *check_malloc(int size)
{
    void *p = malloc (size);
    if (p == NULL) {
        perror("malloc failed");
        exit(-1);
    }
    return p;
}
```

Because memory management is so difficult, most large programs, like web browsers, leak memory. To see which programs on your system are using the most memory, you can use the UNIX utilities `ps` and `top`.

6.3 Implementation

When a process starts, the system allocates space for the text segment and statically allocated data, space for the stack, and space for the heap, which contains dynamically allocated data.

Not all programs allocate data dynamically, so the initial size of the heap might be small or zero. Initially the heap contains only one free chunk.

When `malloc` is called, it checks whether it can find a free chunk that's big enough. If not, it has to request more memory from the system. The function that does that is `sbrk`, which sets the “program break”, which you can think of as a pointer to the end of the heap.

When `sbrk` is called, the OS allocates new pages of physical memory, updates the process's page table, and sets the program break.

In theory, a program could call `sbrk` directly (without using `malloc`) and manage the heap itself. But `malloc` is easier to use and, for most memory-use patterns, it runs fast and uses memory efficiently.

To implement the memory management API (that is, the functions `malloc`, `free`, `calloc`, and `realloc`), most Linux systems use `ptmalloc`, which is based on `dlmalloc`, written by Doug Lea. A short paper that describes key elements of the implementation is available at <http://gee.cs.oswego.edu/dl/html/malloc.html>.

For programmers, the most important elements to be aware of are:

- The run time of `malloc` does not usually depend on the size of the chunk, but might depend on how many free chunks there are. `free` is usually fast, regardless of the number of free chunks. Because `calloc` clears every byte in the chunk, the run time depends on chunk size (as well as the number of free chunks).

`realloc` is sometimes fast, if the new size is smaller than the current size, or if space is available to expand the existing chunk. If not, it has to copy data from the old chunk to the new; in that case, the run time depends on the size of the old chunk.

- **Boundary tags:** When `malloc` allocates a chunk, it adds space at the beginning and end to store information about the chunk, including its size and the state (allocated or free). These bits of data are called “boundary tags”. Using these tags, `malloc` can get from any chunk to the previous chunk and the next chunk in memory. In addition, free chunks are chained into a doubly-linked list; each free chunk contains pointers to the next and previous chunks in the “free list”.

The boundary tags and free list pointers make up `malloc`’s internal data structures. These data structures are interspersed with program data, so it is easy for a program error to damage them.

- **Space overhead:** Boundary tags and free list pointers take up space. The minimum chunk size on most systems is 16 bytes. So for very small chunks, `malloc` is not space efficient. If your program requires large numbers of small structures, it might be more efficient to allocate them in arrays.
- **Fragmentation:** If you allocate and free chunks with varied sizes, the heap will tend to become fragmented. That is, the free space might be broken into many small pieces. Fragmentation wastes space; it also slows the program down by making memory caches less effective.
- **Binning and caching:** The free list is sorted by size into bins, so when `malloc` searches for a chunk with a particular size, it knows what bin to search in. If you free a chunk and then immediately allocate a chunk with the same size, `malloc` will usually be fast.

Chapter 7

Caching

7.1 How programs run

In order to understand caching, you have to understand how computers execute programs. For a deep understanding of this topic, you should study computer architecture. My goal in this chapter is to provide a simple model of program execution.

When a program starts, the code (or text) is usually on a hard disk or solid state drive. The operating system creates a new process to run the program, then the “loader” copies the text from storage into main memory and starts the program by calling `main`.

While the program is running, most of its data is stored in main memory, but some of the data is in registers, which are small units of memory on the CPU. These registers include:

- The program counter, or PC, which contains the address (in memory) of the next instruction in the program.
- The instruction register, or IR, which contains the machine code instruction currently executing.
- The stack pointer, or SP, which contains the address of the stack frame for the current function, which contains its parameters and local variables.
- General-purpose registers that hold the data the program is currently working with.

- A status register, or flag register, that contains information about the current computation. For example, the flag register usually contains a bit that is set if the result of the previous operation was zero.

When a program is running, the CPU executes the following steps, called the “instruction cycle”:

- Fetch: The next instruction is fetched from memory and stored in the instruction register.
- Decode: Part of the CPU, called the “control unit”, decodes the instruction and sends signals to the other parts of the CPU.
- Execute: Signals from the control unit cause the appropriate computation to occur.

Most computers can execute a few hundred different instructions, called the “instruction set”. But most instructions fall into a few general categories:

- Load: Transfers a value from memory to a register.
- Arithmetic/logic: Loads operands from registers, performs a mathematical operation, and stores the result in a register.
- Store: Transfers a value from a register to memory.
- Jump/branch: Changes the program counter, causing the flow of execution to jump to another location in the program. Branches are usually conditional, which means that they check a flag in the flag register and jump only if it is set.

Some instructions sets, including the ubiquitous x86, provide instructions that combine a load and an arithmetic operation.

During each instruction cycle, one instruction is read from the program text. In addition, about half of the instructions in a typical program load or store data. And therein lies one of the fundamental problems of computer architecture: the “memory bottleneck”.

In current computers, a typical core is capable of executing an instruction in less than 1 ns. But the time it takes to transfer data to and from memory is about 100 ns. If the CPU has to wait 100 ns to fetch the next instruction, and another 100 ns to load data, it would complete instructions 200 times slower than what’s theoretically possible. For many computations, memory is the speed limiting factor, not the CPU.

7.2 Cache performance

The solution to this problem, or at least a partial solution, is caching. A “cache” is a small, fast memory that is physically close to the CPU, usually on the same chip.

Actually, current computers typically have several levels of cache: the Level 1 cache, which is the smallest and fastest, might be 1–2 MiB with a access times near 1 ns; the Level 2 cache might have access times near 4 ns, and the Level 3 might take 16 ns.

When the CPU loads a value from memory, it stores a copy in the cache. If the same value is loaded again, the CPU gets the cached copy and doesn’t have to wait for memory.

Eventually the cache gets full. Then, in order to bring something new in, we have to kick something out. So if the CPU loads a value and then loads it again much later, it might not be in cache any more.

The performance of many programs is limited by the effectiveness of the cache. If the instructions and data needed by the CPU are usually in cache, the program can run close to the full speed of the CPU. If the CPU frequently needs data that are not in cache, the program is limited by the speed of memory.

The cache “hit rate”, h , is the fraction of memory accesses that find data in cache; the “miss rate”, m , is the fraction of memory accesses that have to go to memory. If the time to process a cache hit is T_h and the time for a cache miss is T_m , the average time for each memory access is

$$hT_h + mT_m$$

Equivalently, we could define the “miss penalty” as the extra time to process a cache miss, $T_p = T_m - T_h$. Then the average access time is

$$T_h + mT_p$$

When the miss rate is low, the average access time can be close to T_h . That is, the program can perform as if memory ran at cache speeds.

7.3 Locality

When a program reads a byte for the first time, the cache usually loads a “block” or “line” of data that includes the requested byte and some of its

neighbors. If the program goes on to read one of the neighbors, it will already be in cache.

As an example, suppose the block size is 64 B; you read a string with length 64, and the first byte of the string happens to fall at the beginning of a block. When you load the first byte, you incur a miss penalty, but after that the rest of the string will be in cache. After reading the whole string, the hit rate will be 63/64, about 98%. If the string spans two blocks, you would incur 2 miss penalties. But even then the hit rate would be 62/64, or almost 97%. If you then read the same string again, the hit rate would be 100%.

On the other hand, if the program jumps around unpredictably, reading data from scattered locations in memory, and seldom accessing the same location twice, cache performance would be poor.

The tendency of a program to use the same data more than once is called “temporal locality”. The tendency to use data in nearby locations is called “spatial locality”. Fortunately, many programs naturally display both kinds of locality:

- Most programs contain blocks of code with no jumps or branches. Within these blocks, instructions run sequentially, so the access pattern has spatial locality.
- In a loop, programs execute the same instructions many times, so the access pattern has temporal locality.
- The result of one instruction is often used immediately as an operand of the next instruction, so the data access pattern has temporal locality.
- When a program executes a function, its parameters and local variables are stored together on the stack; accessing these values has spatial locality.
- One of the most common processing patterns is to read or write the elements of an array sequentially; this pattern also has spatial locality.

The next section explores the relationship between a program’s access pattern and cache performance.

7.4 Measuring cache performance

When I was a graduate student at U.C. Berkeley I was a teaching assistant for Computer Architecture with Brian Harvey. One of my favorite exercises

involved a program that iterates through an array and measures the average time to read and write an element. By varying the size of the array, it is possible to infer the size of the cache, the block size, and some other attributes.

My modified version of this program is in the `cache` directory of the repository for this book (see Section 0.1).

The important part of the program is this loop:

```
    iters = 0;
    do {
        sec0 = get_seconds();

        for (index = 0; index < limit; index += stride)
            array[index] = array[index] + 1;

        iters = iters + 1;
        sec = sec + (get_seconds() - sec0);

    } while (sec < 0.1);
```

The inner `for` loop traverses the array. `limit` determines how much of the array it traverses; `stride` determines how many elements it skips over. For example, if `limit` is 16 and `stride` is 4, the loop would access elements 0, 4, 8, and 12.

`sec` keeps track of the total CPU time used by the inner loop. The outer loop runs until `sec` exceeds 0.1 seconds, which is long enough that we can compute the average time with sufficient precision.

`get_seconds` uses the system call `clock_gettime`, converts to seconds, and returns the result as a `double`:

```
double get_seconds(){
    struct timespec ts;
    clock_gettime(CLOCK_PROCESS_CPUTIME_ID, &ts);
    return ts.tv_sec + ts.tv_nsec / 1e9;
}
```

To isolate the time to access the elements of the array, the program runs a second loop that is almost identical except that the inner loop doesn't touch the array; it always increments the same variable:

```
    iters2 = 0;
    do {
        sec0 = get_seconds();
```

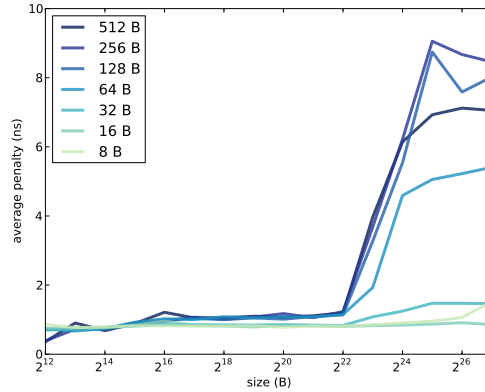


图 7.1: Average miss penalty as a function of array size and stride.

```
for (index = 0; index < limit; index += stride)
    temp = temp + index;
```

```
iters2 = iters2 + 1;
sec = sec - (get_seconds() - sec0);
```

```
} while (iters2 < iters);
```

The second loop runs the same number of iterations as the first. After each iteration, it *subtracts* the elapsed time from `sec`. When the loop completes, `sec` contains the total time for all array accesses, minus the total time it took to increment `temp`. This difference is the total miss penalty incurred by all accesses. Finally, we divide by the number of accesses to get the average miss penalty per access, in ns:

```
sec * 1e9 / iters / limit * stride
```

If you compile and run `cache.c` you should see output like this:

```
Size:    4096 Stride:      8 read+write: 0.8633 ns
Size:    4096 Stride:     16 read+write: 0.7023 ns
Size:    4096 Stride:     32 read+write: 0.7105 ns
Size:    4096 Stride:     64 read+write: 0.7058 ns
```

If you have Python and `matplotlib` installed, you can use `graph_data.py` to graph the results. Figure 7.1 shows the results when I ran it on a Dell Optiplex 7010. Notice that the array size and stride are reported in bytes, not number of array elements.

Take a minute to consider this graph, and see what you can infer about the cache. Here are some things to think about:

- The program reads through the array many times, so it has plenty of temporal locality. If the entire array fits in cache, we expect the average miss penalty to be near 0.
- When the stride is 4 bytes, we read every element of the array, so the program has plenty of spatial locality. If the block size is big enough to contain 64 elements, for example, the hit rate would be 63/64, even if the array does not fit in cache.
- If the stride is equal to the block size (or greater), the spatial locality is effectively zero, because each time we read a block, we only access one element. In that case we expect to see the maximum miss penalty.

In summary, we expect good cache performance if the array is smaller than the cache size *or* if the stride is smaller than the block size. Performance only degrades if the array is bigger than the cache *and* the stride is large.

In Figure 7.1, cache performance is good, for all strides, as long as the array is less than 2^{22} B. We can infer that the cache size is near 4 MiB; in fact, according to the specs, it is 3 MiB.

When the stride is 8, 16, or 32 B, cache performance is good. At 64 B it starts to degrade, and for larger strides the average miss penalty is about 9 ns. We can infer that the block size near 128 B.

Many processors use “multi-level caches” that include a small, fast cache and a bigger, slower cache. In this example, it looks like the miss penalty increases a little when the array size is bigger than 2^{14} B, so it’s possible that this processor also has a 16 KB cache with an access time less than 1 ns.

7.5 Programming for cache performance

Memory caching is implemented in hardware, so most of the time programmers don’t need to know much about it. But if you know how caches work, you can write programs that use them more effectively.

For example, if you are working with a large array, it might be faster to traverse the array once, performing several operations with each element, rather than traversing the array several times.

If you are working with a 2-D array, it might be stored as an array of rows. If you traverse through the elements, it would be faster to go row-wise, with stride equal to the element size, rather than column-wise, with stride equal to the row length.

Linked data structures don't always exhibit spatial locality, because the nodes aren't necessarily contiguous in memory. But if you allocate many nodes at the same time, they are usually co-located in the heap. Or, even better, if you allocate an array of nodes all at once, you know they will be contiguous.

Recursive strategies like mergesort often have good cache behavior because they break big arrays into smaller pieces and then work with the pieces. Sometimes these algorithms can be tuned to take advantage of cache behavior.

For applications where performance is critical, it is possible to design algorithms tailored to the size of the cache, the block size, and other hardware characteristics. Algorithms like that are called "cache-aware". The obvious drawback of cache-aware algorithms is that they are hardware-specific.

7.6 The memory hierarchy

At some point during this chapter, a question like the following might have occurred to you: "If caches are so much faster than main memory, why not make a really big cache and forget about memory?"

Without going too far into computer architecture, there are two reasons: electronics and economics. Caches are fast because they are small and close to the CPU, which minimizes delays due to capacitance and signal propagation. If you make a cache big, it will be slower.

Also, caches take up space on the processor chip, and bigger chips are more expensive. Main memory is usually dynamic random-access memory (DRAM), which uses only one transistor and one capacitor per bit, so it is possible to pack more memory into the same amount of space. But this way of implementing memory is slower than the way caches are implemented.

Also main memory is usually packaged in a dual in-line memory module (DIMM) that includes 16 or more chips. Several small chips are cheaper than one big one.

The trade-off between speed, size, and cost is the fundamental reason for caching. If there were one memory technology that was fast, big, and cheap, we wouldn't need anything else.

The same principle applies to storage as well as memory. Solid state drives (SSD) are fast, but they are more expensive than hard drives (HDD), so they tend to be smaller. Tape drives are even slower than hard drives, but they can store large amounts of data relatively cheaply.

The following table shows typical access times, sizes, and costs for each of these technologies.

Device	Access time	Typical size	Cost
Register	0.5 ns	256 B	?
Cache	1 ns	2 MiB	?
DRAM	100 ns	4 GiB	\$10 / GiB
SSD	10 μ s	100 GiB	\$1 / GiB
HDD	5 ms	500 GiB	\$0.25 / GiB
Tape	minutes	1–2 TiB	\$0.02 / GiB

The number and size of registers depends on details of the architecture. Current computers have about 32 general-purpose registers, each storing one “word”. On a 32-bit computer, a word is 32 bits or 4 B. On a 64-bit computer, a word is 64 bits or 8 B. So the total size of the register file is 100–300 B.

The cost of registers and caches is hard to quantify. They contribute to the cost of the chips they are on, but consumers don’t see that cost directly.

For the other numbers in the table, I looked at the specifications for typical hardware for sale from online computer hardware stores. By the time you read this, these numbers will be obsolete, but they give you an idea of what the performance and cost gaps looked like at one point in time.

These technologies make up the “memory hierarchy” (note that this use of “memory” also includes storage). Each level of the hierarchy is bigger and slower than the one above it. And in some sense, each level acts as a cache for the one below it. You can think of main memory as a cache for programs and data that are stored permanently on SSDs and HDDs. And if you are working with very large datasets stored on tape, you could use hard drives to cache one subset of the data at a time.

7.7 Caching policy

The memory hierarchy suggests a framework for thinking about caching. At every level of the hierarchy, we have to address four fundamental questions of caching:

- Who moves data up and down the hierarchy? At the top of the hierarchy, register allocation is usually done by the compiler. Hardware on the CPU

handles the memory cache. Users implicitly move data from storage to memory when they execute programs and open files. But the operating system also moves data back and forth between memory and storage. At the bottom of the hierarchy, administrators move data explicitly between disk and tape.

- What gets moved? In general, block sizes are small at the top of the hierarchy and bigger at the bottom. In a memory cache, a typical block size is 128 B. Pages in memory might be 4 KiB, but when the operating system reads a file from disk, it might read 10s or 100s of blocks at a time.
- When does data get moved? In the most basic cache, data gets moved into cache when it is used for the first time. But many caches use some kind of “prefetching”, meaning that data is loaded before it is explicitly requested. We have already seen one form of prefetching: loading an entire block when only part of it is requested.
- Where in the cache does the data go? When the cache is full, we can’t bring anything in without kicking something out. Ideally, we want to keep data that will be used again soon and replace data that won’t.

The answers to these questions make up the “cache policy”. Near the top of the hierarchy, cache policies tend to be simple because they have to be fast and they are implemented in hardware. Near the bottom of the hierarchy, there is more time to make decisions, and well-designed policies can make a big difference.

Most cache policies are based on the principle that history repeats itself; if we have information about the recent past, we can use it to predict the immediate future. For example, if a block of data has been used recently, we expect it to be used again soon. This principle suggests a replacement policy called “least recently used,” or LRU, which removes from the cache a block of data that has not been used recently. For more on this topic, see http://en.wikipedia.org/wiki/Cache_algorithms.

7.8 Paging

In systems with virtual memory, the operating system can move pages back and forth between memory and storage. As I mentioned in Section 6.2, this mechanism is called “paging” or sometimes “swapping”.

Here’s how the process works:

1. Suppose Process A calls `malloc` to allocate a chunk. If there is no free space in the heap with the requested size, `malloc` calls `sbrk` to ask the operating system for more memory.
2. If there is a free page in physical memory, the operating system adds it to the page table for Process A, creating a new range of valid virtual addresses.
3. If there are no free pages, the paging system chooses a “victim page” belonging to Process B. It copies the contents of the victim page from memory to disk, then it modifies the page table for Process B to indicate that this page is “swapped out”.
4. Once the data from Process B is written, the page can be reallocated to Process A. To prevent Process A from reading Process B’s data, the page should be cleared.
5. At this point the call to `sbrk` can return, giving `malloc` additional space in the heap. Then `malloc` allocates the requested chunk and returns. Process A can resume.
6. When Process A completes, or is interrupted, the scheduler might allow Process B to resume. When Process B accesses a page that has been swapped out, the memory management unit notices that the page is “invalid” and causes an interrupt.
7. When the operating system handles the interrupt, it sees that the page is swapped out, so it transfers the page back from disk to memory.
8. Once the page is swapped in, Process B can resume.

When paging works well, it can greatly improve the utilization of physical memory, allowing more processes to run in less space. Here’s why:

- Most processes don’t use all of their allocated memory. Many parts of the text segment are never executed, or execute once and never again. Those pages can be swapped out without causing any problems.
- If a program leaks memory, it might leave allocated space behind and never access it again. By swapping those pages out, the operating system can effectively plug the leak.
- On most systems, there are processes like daemons that sit idle most of the time and only occasionally “wake up” to respond to events. While they are idle, these processes can be swapped out.

- A user might have many windows open, but only a few are active at a time. The inactive processes can be swapped out.
- Also, there might be many processes running the same program. These processes can share the same text and static segments, avoiding the need to keep multiple copies in physical memory.

If you add up the total memory allocated to all processes, it can greatly exceed the size of physical memory, and yet the system can still behave well.

Up to a point.

When a process accesses a page that's swapped out, it has to get the data back from disk, which can take several milliseconds. The delay is often noticeable. If you leave a window idle for a long time and then switch back to it, it might start slowly, and you might hear the disk drive working while pages are swapped in.

Occasional delays like that might be acceptable, but if you have too many processes using too much space, they start to interfere with each other. When Process A runs, it evicts the pages Process B needs. Then when B runs, it evicts the pages A needs. When this happens, both processes slow to a crawl and the system can become unresponsive. This scenario is called “thrashing”.

In theory, operating systems could avoid thrashing by detecting an increase in paging and blocking or killing processes until the system is responsive again. But as far as I can tell, most systems don't do this, or don't do it well; it is often left to users to limit their use of physical memory or try to recover when thrashing occurs.

Chapter 8

Multitasking

In many current systems, the CPU contains multiple cores, which means it can run several processes at the same time. In addition, each core is capable of “multitasking”, which means it can switch from one process to another quickly, creating the illusion that many processes are running at the same time.

The part of the operating system that implements multitasking is the “kernel”. In a nut or seed, the kernel is the innermost part, surrounded by a shell. In an operating system, the kernel is the lowest level of software, surrounded by several other layers, including an interface called a “shell.” Computer scientists love extended metaphors.

At its most basic, the kernel’s job is to handle interrupts. An “interrupt” is an event that stops the normal instruction cycle and causes the flow of execution to jump to a special section of code called an “interrupt handler”.

A **hardware interrupt** is caused when a device sends a signal to the CPU. For example, a network interface might cause an interrupt when a packet of data arrives, or a disk drive might cause an interrupt when a data transfer is complete. Most systems also have timers that cause interrupts at regular intervals, or after an elapsed time.

A **software interrupt** is caused by a running program. For example, if an instruction cannot complete for some reason, it might trigger an interrupt so the condition can be handled by the operating system. Some floating-point errors, like division by zero, are handled using interrupts.

When a program needs to access a hardware device, it makes a **system call**, which is similar to a function call, except that instead of jumping to the beginning of the function, it executes a special instruction that triggers an

interrupt, causing the flow of execution to jump to the kernel. The kernel reads the parameters of the system call, performs the requested operation, and then resumes the interrupted process.

8.1 Hardware state

Handling interrupts requires cooperation between hardware and software. When an interrupt occurs, there might be several instructions running on the CPU, data stored in registers, and other **hardware state**.

Usually the hardware is responsible for bringing the CPU to a consistent state; for example, every instruction should either complete or behave as if it never started. No instruction should be left half complete. Also, the hardware is responsible for saving the program counter (PC), so the kernel knows where to resume.

Then, usually, it is the responsibility of the interrupt handler to save the rest of the hardware state before it does anything that might modify it, and then restore the saved state before the interrupted process resumes.

Here is an outline of this sequence of events:

1. When the interrupt occurs, the hardware saves the program counter in a special register and jumps to the appropriate interrupt handler.
2. The interrupt handler stores the program counter and the status register in memory, along with the contents of any data registers it plans to use.
3. The interrupt handler runs whatever code is needed to handle the interrupt.
4. Then it restores the contents of the saved registers. Finally, it restores the program counter of the interrupted process, which has the effect of jumping back to the interrupted instruction.

If this mechanism works correctly, there is generally no way for the interrupted process to know there was an interrupt, unless it detects the change in time between instructions.

8.2 Context switching

Interrupt handlers can be fast because they don't have to save the entire hardware state; they only have to save registers they are planning to use.

But when an interrupt occurs, the kernel does not always resume the interrupted process. It has the option of switching to another process. This mechanism is called a “context switch”.

In general, the kernel doesn’t know which registers a process will use, so it has to save all of them. Also, when it switches to a new process, it might have to clear data stored in the memory management unit (see Section 3.6). And after the context switch, it might take some time for the new process to load data into the cache. For these reasons, context switches are relatively slow, on the order of thousands of cycles, or a few microseconds.

In a multi-tasking system, each process is allowed to run for a short period of time called a “time slice” or “quantum”. During a context switch, the kernel sets a hardware timer that causes an interrupt at the end of the time slice. When the interrupt occurs, the kernel can switch to another process or allow the interrupted process to resume. The part of the operating system that makes this decision is the “scheduler”.

8.3 The process life cycle

When a process is created, the operating system allocates a data structure that contains information about the process, called a “process control block” or PCB. Among other things, the PCB keeps track of the process state, which is one of:

- Running, if the process is currently running on a core.
- Ready, if the process could be running, but isn’t, usually because there are more runnable processes than cores.
- Blocked, if the process cannot run because it is waiting for a future event like network communication or a disk read.
- Done, if the process has completed, but has exit status information that has not been read yet.

Here are the events that cause a process to transition from one state to another:

- A process is created when the running program executes a system call like `fork`. At the end of the system call, the new process is usually ready. Then the scheduler might resume the original process (the “parent”) or start the new process (the “child”).

- When a process is started or resumed by the scheduler, its state changes from ready to running.
- When a process is interrupted and the scheduler chooses not to let it resume, its state changes from running to ready.
- If a process executes a system call that cannot complete immediately, like a disk request, it becomes blocked and the scheduler usually chooses another process.
- When an operation like a disk request completes, it causes an interrupt. The interrupt handler figures out which process was waiting for the request and switches its state from blocked to ready. Then the scheduler may or may not choose to resume the unblocked process.
- When a process calls `exit`, the interrupt handler stores the exit code in the PCB and changes the process's state to done.

8.4 Scheduling

As we saw in Section 2.3 there might be hundreds of processes on a computer, but usually most of them are blocked. Most of the time, there are only a few processes that are ready or running. When an interrupt occurs, the scheduler decides which process to start or resume.

On a workstation or laptop, the primary goal of the scheduler is to minimize response time; that is, the computer should respond quickly to user actions. Response time is also important on a server, but in addition the scheduler might try to maximize throughput, which is the number of requests that complete per unit of time.

Usually the scheduler doesn't have much information about what processes are doing, so its decisions are based on a few heuristics:

- Processes might be limited by different resources. A process that does a lot of computation is probably CPU-bound, which means that its run time depends on how much CPU time it gets. A process that reads data from a network or disk might be I/O-bound, which means that it would run faster if data input and output went faster, but would not run faster with more CPU time. Finally, a process that interacts with the user is probably blocked, most of the time, waiting for user actions.

The operating system can sometimes classify processes based on their past behavior, and schedule them accordingly. For example, when an

interactive process is unblocked, it should probably run immediately, because a user is probably waiting for a reply. On the other hand, a CPU-bound process that has been running for a long time might be less time-sensitive.

- If a process is likely to run for a short time and then make a blocking request, it should probably run immediately, for two reasons: (1) if the request takes some time to complete, we should start it as soon as possible, and (2) it is better for a long-running process to wait for a short one, rather than the other way around.

As an analogy, suppose you are making an apple pie. The crust takes 5 minutes to prepare, but then it has to chill for half an hour. It takes 20 minutes to prepare the filling. If you prepare the crust first, you can prepare the filling while the crust is chilling, and you can finish the pie in 35 minutes. If you prepare the filling first, the process takes 55 minutes.

Most schedulers use some form of priority-based scheduling, where each process has a priority that can be adjusted up or down over time. When the scheduler runs, it chooses the runnable process with the highest priority.

Here are some of the factors that determine a process's priority:

- A process usually starts with a relatively high priority so it starts running quickly.
- If a process makes a request and blocks before its time slice is complete, it is more likely to be interactive or I/O-bound, so its priority should go up.
- If a process runs for an entire time slice, it is more likely to be long-running and CPU-bound, so its priority should go down.
- If a task blocks for a long time and then becomes ready, it should get a priority boost so it can respond to whatever it was waiting for.
- If process A is blocked waiting for process B, for example if they are connected by a pipe, the priority of process B should go up.
- The system call `nice` allows a process to decrease (but not increase) its own priority, allowing programmers to pass explicit information to the scheduler.

For most systems running normal workloads, scheduling algorithms don't have a substantial effect on performance. Simple scheduling policies are usually good enough.

8.5 Real-time scheduling

However, for programs that interact with the real world, scheduling can be very important. For example, a program that reads data from sensors and controls motors might have to complete recurring tasks at some minimum frequency and react to external events with some maximum response time. These requirements are often expressed in terms of “tasks” that must be completed before “deadlines”.

Scheduling tasks to meet deadlines is called “real-time scheduling”. For some applications, a general-purpose operating system like Linux can be modified to handle real-time scheduling. These modifications might include:

- Providing richer APIs for controlling task priorities.
- Modifying the scheduler to guarantee that the process with highest priority runs within a fixed amount of time.
- Reorganizing interrupt handlers to guarantee a maximum completion time.
- Modifying locks and other synchronization mechanisms (coming up in the next chapter) to allow a high-priority task to preempt a lower-priority task.
- Choosing an implementation of dynamic memory allocation that guarantees a maximum completion time.

For more demanding applications, especially in domains where real-time response is a matter of life and death, “real-time operating systems” provide specialized capabilities, often with much simpler designs than general purpose operating systems.

Chapter 9

Threads

When I mentioned threads in Section 2.3, I said that a thread is a kind of process. Now I will provide a more careful explanation.

When you create a process, the operating system creates a new address space, which includes the text segment, static segment, and heap; it also creates a new “thread of execution”, which includes the program counter and other hardware state, and the call stack.

The processes we have seen so far are “single-threaded”, which means that only one thread of execution runs in each address space. In this chapter, you will learn about “multi-threaded” processes that have multiple threads running in the same address space.

Within a single process, all threads share the same text segment, so they run the same code. But different threads often run different parts of the code.

And they share the same static segment, so if one thread changes a global variable, other threads see the change. They also share the heap, so threads can share dynamically-allocated chunks.

But each thread has its own stack, so threads can call functions without interfering with each other. Usually threads don’t access each other’s local variables (and sometimes they can’t).

The example code for this chapter is in the repository for this book, in a directory named `counter`. For information on downloading this code, see Section 0.1.

9.1 Creating threads

The most popular threading standard used with C is POSIX Threads, or Pthreads for short. The POSIX standard defines a thread model and an interface for creating and controlling threads. Most versions of UNIX provide an implementation of Pthreads.

Using Pthreads is like using most C libraries:

- You include headers files at the beginning of your program.
- You write code that calls functions defined by Pthreads.
- When you compile the program, you link it with the Pthread library.

For my examples, I include the following headers:

```
#include <stdio.h>
#include <stdlib.h>
#include <pthread.h>
#include <semaphore.h>
```

The first two are standard; the third is for Pthreads and the fourth is for semaphores. To compile with the Pthread library in `gcc`, you can use the `-l` option on the command line:

```
gcc -g -O2 -o array array.c -lpthread
```

This compiles a source file named `array.c` with debugging info and optimization, links with the Pthread library, and generates an executable named `array`.

9.2 Creating threads

The Pthread function that creates threads is called `pthread_create`. The following function shows how to use it:

```
pthread_t make_thread(void *(*entry)(void *), Shared *shared)
{
    int n;
    pthread_t thread;

    n = pthread_create(&thread, NULL, entry, (void *)shared);
    if (n != 0) {
        perror("pthread_create failed");
    }
}
```

```
        exit(-1);
    }
    return thread;
}
```

`make_thread` is a wrapper I wrote to make `pthread_create` easier to use, and to provide error-checking.

The return type from `pthread_create` is `pthread_t`, which you can think of as an id or “handle” for the new thread.

If `pthread_create` succeeds, it returns 0 and `make_thread` returns the handle of the new thread. If an error occurs, `pthread_create` returns an error code and `make_thread` prints an error message and exits.

The parameters of `make_thread` take some explaining. Starting with the second, `Shared` is a structure I defined to contain values shared between threads. The following `typedef` statement creates the new type:

```
typedef struct {
    int counter;
} Shared;
```

In this case, the only shared variable is `counter`. `make_shared` allocates space for a `Shared` structure and initializes the contents:

```
Shared *make_shared()
{
    Shared *shared = check_malloc(sizeof (Shared));
    shared->counter = 0;
    return shared;
}
```

Now that we have a shared data structure, let’s get back to `make_thread`. The first parameter is a pointer to a function that takes a `void` pointer and returns a `void` pointer. If the syntax for declaring this type makes your eyes bleed, you are not alone. Anyway, the purpose of this parameter is to specify the function where the execution of the new thread will begin. By convention, this function is named `entry`:

```
void *entry(void *arg)
{
    Shared *shared = (Shared *) arg;
    child_code(shared);
    pthread_exit(NULL);
}
```

The parameter of `entry` has to be declared as a `void` pointer, but in this program we know that it is really a pointer to a `Shared` structure, so we can typecast it accordingly and then pass it along to `child_code`, which does the real work.

As a simple example, `child_code` prints the value of the shared counter and increments it.

```
void child_code(Shared *shared)
{
    printf("counter = %d\n", shared->counter);
    shared->counter++;
}
```

When `child_code` returns, `entry` invokes `pthread_exit` which can be used to pass a value to the thread that joins with this thread. In this case, the child has nothing to say, so we pass `NULL`.

Finally, here is the code that creates the child threads:

```
int i;
pthread_t child[NUM_CHILDREN];

Shared *shared = make_shared(1000000);

for (i=0; i<NUM_CHILDREN; i++) {
    child[i] = make_thread(entry, shared);
}
```

`NUM_CHILDREN` is a compile-time constant that determines the number of child threads. `child` is an array of thread handles.

9.3 Joining threads

When one thread wants to wait for another thread to complete, it invokes `pthread_join`. Here is my wrapper for `pthread_join`:

```
void join_thread(pthread_t thread)
{
    int ret = pthread_join(thread, NULL);
    if (ret == -1) {
        perror("pthread_join failed");
        exit(-1);
    }
}
```

The parameter is the handle of the thread you want to wait for. All the wrapper does is call `pthread_join` and check the result.

Any thread can join any other thread, but in the most common pattern the parent thread creates and joins all child threads. Continuing the example from the previous section, here's the code that waits on the children:

```
for (i=0; i<NUM_CHILDREN; i++) {  
    join_thread(child[i]);  
}
```

This loop waits for the children one at a time in the order they were created. There is no guarantee that the child threads complete in that order, but this loop works correctly even if they don't. If one of the children is late, the loop might have to wait, and other children might complete in the meantime. But regardless, the loop exits only when all children are done.

If you have downloaded the repository for this book (see Section 0.1), you'll find this example in `counter/counter.c`. You can compile and run it like this:

```
$ make counter  
gcc -Wall counter.c -o counter -lpthread  
$ ./counter
```

When I ran it with 5 children, I got the following output:

```
counter = 0  
counter = 0  
counter = 1  
counter = 0  
counter = 3
```

When you run it, you will probably get different results. And if you run it again, you might get different results each time. What's going on?

9.4 Synchronization errors

The problem with the previous program is that the children access the shared variable, `counter`, without synchronization, so several threads can read the same value of `counter` before any threads increment it.

Here is a sequence of events that could explain the output in the previous section:

```
Child A reads 0  
Child B reads 0
```

```
Child C reads 0
Child A prints  0
Child B prints  0
Child A sets counter=1
Child D reads 1
Child D prints  1
Child C prints  0
Child A sets counter=1
Child B sets counter=2
Child C sets counter=3
Child E reads 3
Child E prints  3
Child D sets counter=4
Child E sets counter=5
```

Each time you run the program, threads might be interrupted at different points, or the scheduler might choose different threads to run, so the sequence of events, and the results, will be different.

Suppose we want to impose some order. For example, we might want each thread to read a different value of `counter` and increment it, so that the value of `counter` reflects the number of threads that have executed `child_code`.

To enforce that requirement, we can use a “mutex”, which is an object that guarantees “mutual exclusion” for a block of code; that is, only one thread can execute the block at a time.

I have written a small module called `mutex.c` that provides mutex objects. I’ll show you how to use it first; then I’ll explain how it works.

Here’s a version of `child_code` that uses a mutex to synchronize threads:

```
void child_code(Shared *shared)
{
    mutex_lock(shared->mutex);
    printf("counter = %d\n", shared->counter);
    shared->counter++;
    mutex_unlock(shared->mutex);
}
```

Before any thread can access `counter`, it has to “lock” the mutex, which has the effect of barring all other threads. Suppose Thread A has locked the mutex and is in the middle of `child_code`. If Thread B arrives and executes `mutex_lock`, it blocks.

When Thread A is done, it executes `mutex_unlock`, which allows Thread B to proceed. In effect, the threads line up to execute `child_code` one at a time,

so they can't interfere with each other. When I run this code with 5 children, I get:

```
counter = 0
counter = 1
counter = 2
counter = 3
counter = 4
```

And that satisfies the requirements. In order for this solution to work, I have to add the Mutex to the Shared struct:

```
typedef struct {
    int counter;
    Mutex *mutex;
} Shared;
```

And initialize it in `make_shared`

```
Shared *make_shared(int end)
{
    Shared *shared = check_malloc(sizeof(Shared));
    shared->counter = 0;
    shared->mutex = make_mutex();    //-- this line is new
    return shared;
}
```

The code in this section is in `counter_mutex.c`. The definition of `Mutex` is in `mutex.c`, which I explain in the next section.

9.5 Mutex

My definition of `Mutex` is a wrapper for a type called `pthread_mutex_t`, which is defined in the POSIX threads API.

To create a POSIX mutex, you have to allocate space for a `pthread_mutex_t` type and then call `pthread_mutex_init`.

One of the problems with this API is that `pthread_mutex_t` behaves like a structure, so if you pass it as an argument, it makes a copy, which makes the mutex behave incorrectly. To avoid that, you have to pass `pthread_mutex_t` by address.

My code makes it easier to get that right. It defines a type, `Mutex`, which is just a more readable name for `pthread_mutex_t`:

```
#include <pthread.h>
```

```
typedef pthread_mutex_t Mutex;
```

Then it defines `make_mutex`, which allocates space and initializes the mutex:

```
Mutex *make_mutex()
{
    Mutex *mutex = check_malloc(sizeof(Mutex));
    int n = pthread_mutex_init(mutex, NULL);
    if (n != 0) perror_exit("make_lock failed");
    return mutex;
}
```

The return value is a pointer, which you can pass around as an argument without causing unwanted copying.

The functions to lock and unlock the mutex are simple wrappers for POSIX functions:

```
void mutex_lock(Mutex *mutex)
{
    int n = pthread_mutex_lock(mutex);
    if (n != 0) perror_exit("lock failed");
}

void mutex_unlock(Mutex *mutex)
{
    int n = pthread_mutex_unlock(mutex);
    if (n != 0) perror_exit("unlock failed");
}
```

This code is in `mutex.c` and the header file `mutex.h`.

Chapter 10

Condition variables

Many simple synchronization problems can be solved using mutexes as shown in the previous chapter. In this chapter I introduce a bigger challenge, the well-known “Producer-Consumer problem”, and a new tool to solve it, the condition variable.

10.1 The work queue

In some multi-threaded programs, threads are organized to perform different tasks. Often they communicate with each other using a queue, where some threads, called “producers”, put data into the queue and other threads, called “consumers”, take data out.

For example, in applications with a graphical user interface, there might be one thread that runs the GUI, responding to user events, and another thread that processes user requests. In that case, the GUI thread might put requests into a queue and the “back end” thread might take requests out and process them.

To support this organization, we need a queue implementation that is “thread safe”, which means that both threads (or more than two) can access the queue at the same time. And we need to handle the special cases when the queue is empty and, if the size of the queue is bounded, when the queue is full.

I’ll start with a simple queue that is not thread safe, then we’ll see what goes wrong and fix it. The code for this example is in the repository for this book, in a folder called `queue`. The file `queue.c` contains a basic implementation of a circular buffer, which you can read about at https://en.wikipedia.org/wiki/Circular_buffer.

Here's the structure definition:

```
typedef struct {
    int *array;
    int length;
    int next_in;
    int next_out;
} Queue;
```

`array` is the array that contains the elements of the queue. For this example the elements are ints, but more generally they would be structures that contain user events, items of work, etc.

`length` is the length of the array. `next_in` is an index into the array that indicates where the next element should be added; similarly, `next_out` is the index of the next element that should be removed.

`make_queue` allocates space for this structure and initializes the fields:

```
Queue *make_queue(int length)
{
    Queue *queue = (Queue *) malloc(sizeof(Queue));
    queue->length = length + 1;
    queue->array = (int *) malloc(length * sizeof(int));
    queue->next_in = 0;
    queue->next_out = 0;
    return queue;
}
```

The initial value for `next_out` needs some explaining. Since the queue is initially empty, there is no next element to remove, so `next_out` is invalid. Setting `next_out == next_in` is a special case that indicates that the queue is empty, so we can write:

```
int queue_empty(Queue *queue)
{
    return (queue->next_in == queue->next_out);
}
```

Now we can add elements to the queue using `queue_push`:

```
void queue_push(Queue *queue, int item) {
    if (queue_full(queue)) {
        perror_exit("queue is full");
    }

    queue->array[queue->next_in] = item;
```

```
    queue->next_in = queue_incr(queue, queue->next_in);  
}
```

If the queue is full, `queue_push` prints an error message and exits. I will explain `queue_full` soon.

If the queue is not full, `queue_push` inserts the new element and then increments `next_in` using `queue_incr`:

```
int queue_incr(Queue *queue, int i)  
{  
    return (i+1) % queue->length;  
}
```

When the index, `i`, gets to the end of the array, it wraps around to 0. And that's where we run into a tricky part. If we keep adding elements to the queue, eventually `next_in` wraps around and catches up with `next_out`. But if `next_in == next_out`, we would incorrectly conclude that the queue was empty.

To avoid that, we define another special case to indicate that the queue is full:

```
int queue_full(Queue *queue)  
{  
    return (queue_incr(queue, queue->next_in) == queue->next_out);  
}
```

If incrementing `next_in` lands on `next_out`, that means we can't add another element without making the queue seem empty. So we stop one element before the "end" (keeping in mind that the end of the queue can be anywhere, not necessarily the end of the array).

Now we can write `queue_pop`, which removes and returns the next element from the queue:

```
int queue_pop(Queue *queue) {  
    if (queue_empty(queue)) {  
        perror_exit("queue is empty");  
    }  
  
    int item = queue->array[queue->next_out];  
    queue->next_out = queue_incr(queue, queue->next_out);  
    return item;  
}
```

If you try to pop from an empty queue, `queue_pop` prints an error message and exits.

10.2 Producers and consumers

Now let's make some threads to access this queue. Here's the producer code:

```
void *producer_entry(void *arg) {
    Shared *shared = (Shared *) arg;

    for (int i=0; i<QUEUE_LENGTH-1; i++) {
        printf("adding item %d\n", i);
        queue_push(shared->queue, i);
    }
    pthread_exit(NULL);
}
```

Here's the consumer code:

```
void *consumer_entry(void *arg) {
    int item;
    Shared *shared = (Shared *) arg;

    for (int i=0; i<QUEUE_LENGTH-1; i++) {
        item = queue_pop(shared->queue);
        printf("consuming item %d\n", item);
    }
    pthread_exit(NULL);
}
```

Here's the parent code that starts the threads and waits for them

```
pthread_t child[NUM_CHILDREN];

Shared *shared = make_shared();

child[0] = make_thread(producer_entry, shared);
child[1] = make_thread(consumer_entry, shared);

for (int i=0; i<NUM_CHILDREN; i++) {
    join_thread(child[i]);
}
```

And finally here's the shared structure that contains the queue:

```
typedef struct {
    Queue *queue;
} Shared;
```

```
Shared *make_shared()
```

```

{
    Shared *shared = check_malloc(sizeof(Shared));
    shared->queue = make_queue(Queue_LENGTH);
    return shared;
}

```

The code we have so far is a good starting place, but it has several problems:

- Access to the queue is not thread safe. Different threads could access `array`, `next_in`, and `next_out` at the same time and leave the queue in a broken, “inconsistent” state.
- If the consumer is scheduled first, it finds the queue empty, print an error message, and exits. We would rather have the consumer block until the queue is not empty. Similarly, we would like the producer to block if the queue is full.

In the next section, we solve the first problem with a **Mutex**. In the following section, we solve the second problem with condition variables.

10.3 Mutual exclusion

We can make the queue thread safe with a mutex. This version of the code is in `queue_mutex.c`.

First we add a **Mutex** pointer to the queue structure:

```

typedef struct {
    int *array;
    int length;
    int next_in;
    int next_out;
    Mutex *mutex;           //-- this line is new
} Queue;

```

And initialize the **Mutex** in `make_queue`:

```

Queue *make_queue(int length) {
    Queue *queue = (Queue *) malloc(sizeof(Queue));
    queue->length = length;
    queue->array = (int *) malloc(length * sizeof(int));
    queue->next_in = 0;
    queue->next_out = 0;
    queue->mutex = make_mutex();    //-- new
    return queue;
}

```

Next we add synchronization code to `queue_push`:

```
void queue_push(Queue *queue, int item) {
    mutex_lock(queue->mutex);    //-- new
    if (queue_full(queue)) {
        mutex_unlock(queue->mutex);    //-- new
        perror_exit("queue is full");
    }

    queue->array[queue->next_in] = item;
    queue->next_in = queue_incr(queue, queue->next_in);
    mutex_unlock(queue->mutex);    //-- new
}
```

Before checking whether the queue is full, we have to lock the `Mutex`. If the queue is full, we have to unlock the `Mutex` before exiting; otherwise the thread would leave it locked and no other threads could proceed.

The synchronization code for `queue_pop` is similar:

```
int queue_pop(Queue *queue) {
    mutex_lock(queue->mutex);
    if (queue_empty(queue)) {
        mutex_unlock(queue->mutex);
        perror_exit("queue is empty");
    }

    int item = queue->array[queue->next_out];
    queue->next_out = queue_incr(queue, queue->next_out);
    mutex_unlock(queue->mutex);
    return item;
}
```

Note that the other `Queue` functions, `queue_full`, `queue_empty`, and `queue_incr` do not try to lock the mutex. Any thread that calls these functions is required to lock the mutex first; this requirement is part of the documented interface for these functions.

With this additional code, the queue is thread safe; if you run it, you should not see any synchronization errors. But it is likely that the consumer will exit at some point because the queue is empty, or the producer will exit because the queue is full, or both.

The next step is to add condition variables.

10.4 Condition variables

A condition variable is a data structure associated with a condition; it allows threads to block until the condition becomes true. For example, `thread_pop` might want check whether the queue is empty and, if so, wait for a condition like “queue not empty”.

Similarly, `thread_push` might want to check whether the queue is full and, if so, block until it is not full.

I’ll handle the first condition here, and you will have a chance to handle the second condition as an exercise.

First we add a condition variable to the `Queue` structure:

```
typedef struct {
    int *array;
    int length;
    int next_in;
    int next_out;
    Mutex *mutex;
    Cond *nonempty;    //-- new
} Queue;
```

And initialize it in `make_queue`:

```
Queue *make_queue(int length)
{
    Queue *queue = (Queue *) malloc(sizeof(Queue));
    queue->length = length;
    queue->array = (int *) malloc(length * sizeof(int));
    queue->next_in = 0;
    queue->next_out = 0;
    queue->mutex = make_mutex();
    queue->nonempty = make_cond();    //-- new
    return queue;
}
```

Now in `queue_pop`, if we find the queue empty, we don’t exit; instead we use the condition variable to block:

```
int queue_pop(Queue *queue) {
    mutex_lock(queue->mutex);
    while (queue_empty(queue)) {
        cond_wait(queue->nonempty, queue->mutex);    //-- new
    }
}
```

```

    int item = queue->array[queue->next_out];
    queue->next_out = queue_incr(queue, queue->next_out);
    mutex_unlock(queue->mutex);
    cond_signal(queue->nonfull);    //-- new
    return item;
}

```

`cond_wait` is complicated, so let's take it slow. The first argument is the condition variable; in this case, the condition we are waiting for is "queue not empty". The second argument is the mutex that protects the queue.

When the thread that locked the mutex calls `cond_wait`, it unlocks the mutex and then blocks. This is important. If `cond_wait` did not unlock the mutex before blocking, no other thread would be able to access the queue, no more items could be added, and the queue would always be empty.

So while the consumer is blocked on `nonempty`, the producer can run. Let's see what happens when the producer runs `queue_push`:

```

void queue_push(Queue *queue, int item) {
    mutex_lock(queue->mutex);
    if (queue_full(queue)) {
        mutex_unlock(queue->mutex);
        perror_exit("queue is full");
    }
    queue->array[queue->next_in] = item;
    queue->next_in = queue_incr(queue, queue->next_in);
    mutex_unlock(queue->mutex);
    cond_signal(queue->nonempty);    //-- new
}

```

Just as before, `queue_push` locks the `Mutex` and checks whether the queue is full. Assuming it is not, `queue_push` adds a new element to the queue and then unlocks the `Mutex`.

But before returning, it does one more thing: it "signals" the condition variable `nonempty`.

Signalling a condition variable usually indicates that the condition is true. If there are no threads waiting on the condition variable, the signal has no effect.

If there are threads waiting on the condition variable, one of them gets unblocked and resumes execution of `cond_wait`. But before the awakened thread can return from `cond_wait`, it has to wait for and lock the `Mutex`, again.

Now go back to `queue_pop` and see what happens when the thread returns from `cond_wait`. It loops back to the top of the while loop and checks the

condition again. I'll explain why in just a second, but for now let's assume that the condition is true; that is, the queue is not empty.

When the consumer thread exits the while loop, we know two things: (1) the condition is true, so there is at least one item in the queue, and (2) the `Mutex` is locked, so it is safe to access the queue.

After removing an item, `queue_pop` unlocks the mutex and returns.

In the next section I'll show you how my `Cond` code works, but first I want to answer two frequently-asked questions:

- Why is `cond_wait` inside a while loop rather than an if statement; that is, why do we have to check the condition again after returning from `cond_wait`?

The primary reason you have to re-check the condition is the possibility of an intercepted signal. Suppose Thread A is waiting on `nonempty`. Thread B adds an item to the queue and signals `nonempty`. Thread A wakes up and tries to lock the mutex, but before it gets the chance, Evil Thread C swoops in, locks the mutex, pops the item from the queue, and unlocks the mutex. Now the queue is empty again, but Thread A is not blocked any more. Thread A could lock the mutex and returns from `cond_wait`. If Thread A does not check the condition again, it would try to pop an element from an empty queue, and probably cause an error.

- The other question that comes up when people learn about condition variables is “How does the condition variable know what condition it is associated with?”

This question is understandable because there is no explicit connection between a `Cond` structure and the condition it relates to. The connection is implicit in the way it is used.

Here's one way to think of it: the condition associated with a `Cond` is the thing that is false when you call `cond_wait` and true when you call `cond_signal`.

Because threads have to check the condition when they return from `cond_wait`, it is not strictly necessary to call `cond_signal` only when the condition is true. If you have reason to think the condition *might* be true, you could call `cond_signal` as a suggestion that now is a good time to check.

10.5 Condition variable implementation

The Cond structure I used in the previous section is a wrapper for a type called `pthread_cond_t`, which is defined in the POSIX threads API. It is very similar to Mutex, which is a wrapper for `pthread_mutex_t`. Both wrappers are defined in `utils.c` and `utils.h`.

Here's the typedef:

```
typedef pthread_cond_t Cond;
```

`make_cond` allocates space, initializes the condition variable, and returns a pointer:

```
Cond *make_cond() {
    Cond *cond = check_malloc(sizeof(Cond));
    int n = pthread_cond_init(cond, NULL);
    if (n != 0) perror_exit("make_cond failed");

    return cond;
}
```

And here are the wrappers for `cond_wait` and `cond_signal`.

```
void cond_wait(Cond *cond, Mutex *mutex) {
    int n = pthread_cond_wait(cond, mutex);
    if (n != 0) perror_exit("cond_wait failed");
}

void cond_signal(Cond *cond) {
    int n = pthread_cond_signal(cond);
    if (n != 0) perror_exit("cond_signal failed");
}
```

At this point there should be nothing too surprising there.

Chapter 11

Semaphores in C

Semaphores are a good way to learn about synchronization, but they are not as widely used, in practice, as mutexes and condition variables.

Nevertheless, there are some synchronization problems that can be solved simply with semaphores, yielding solutions that are more demonstrably correct.

This chapter presents a C API for working with semaphores and my code for making it easier to work with. And it presents a final challenge: can you write an implementation of a semaphore using mutexes and condition variables?

The code for this chapter is in directory `semaphore` in the repository for this book (see Section 0.1).

11.1 POSIX Semaphores

A semaphore is a data structure used to help threads work together without interfering with each other.

The POSIX standard specifies an interface for semaphores; it is not part of Pthreads, but most UNIXes that implement Pthreads also provide semaphores.

POSIX semaphores have type `sem_t`. As usual, I put a wrapper around `sem_t` to make it easier to use. The interface is defined in `sem.h`:

```
typedef sem_t Semaphore;
```

```
Semaphore *make_semaphore(int value);  
void semaphore_wait(Semaphore *sem);  
void semaphore_signal(Semaphore *sem);
```

`Semaphore` is a synonym for `sem_t`, but I find it more readable, and the capital letter reminds me to treat it like an object and pass it by pointer.

The implementation of these functions is in `sem.c`:

```
Semaphore *make_semaphore(int value)
{
    Semaphore *sem = check_malloc(sizeof(Semaphore));
    int n = sem_init(sem, 0, value);
    if (n != 0) perror_exit("sem_init failed");
    return sem;
}
```

`make_semaphore` takes the initial value of the semaphore as a parameter. It allocates space for a `Semaphore`, initializes it, and returns a pointer to `Semaphore`.

`sem_init` returns 0 if it succeeds and -1 if anything goes wrong. One nice thing about using wrapper functions is that you can encapsulate the error-checking code, which makes the code that uses these functions more readable.

Here is the implementation of `semaphore_wait`:

```
void semaphore_wait(Semaphore *sem)
{
    int n = sem_wait(sem);
    if (n != 0) perror_exit("sem_wait failed");
}
```

And here is `semaphore_signal`:

```
void semaphore_signal(Semaphore *sem)
{
    int n = sem_post(sem);
    if (n != 0) perror_exit("sem_post failed");
}
```

I prefer to call this operation “signal” rather than “post”, although both terms are common.

Here’s an example that shows how to use a semaphore as a mutex:

```
Semaphore *mutex = make_semaphore(1);

semaphore_wait(mutex);
    // protected code goes here
semaphore_signal(mutex);
```

When you use a semaphore as a mutex, you usually initialize it to 1 to indicate that the mutex is unlocked; that is, one thread can pass the semaphore without blocking.

Here I am using the variable name `mutex` to indicate that the semaphore is being used as a mutex. But remember that the behavior of a semaphore is not the same as a Pthread mutex.

11.2 Producers and consumers with semaphores

Using these semaphore wrapper functions, we can write a solution to the Producer-Consumer problem from Section 10.2. The code in this section is in `queue_sem.c`.

Here's the new definition of `Queue`, replacing the mutex and condition variables with semaphores:

```
typedef struct {
    int *array;
    int length;
    int next_in;
    int next_out;
    Semaphore *mutex;      //-- new
    Semaphore *items;      //-- new
    Semaphore *spaces;     //-- new
} Queue;
```

And here's the new version of `make_queue`:

```
Queue *make_queue(int length)
{
    Queue *queue = (Queue *) malloc(sizeof(Queue));
    queue->length = length;
    queue->array = (int *) malloc(length * sizeof(int));
    queue->next_in = 0;
    queue->next_out = 0;
    queue->mutex = make_semaphore(1);
    queue->items = make_semaphore(0);
    queue->spaces = make_semaphore(length-1);
    return queue;
}
```

`mutex` is used to guarantee exclusive access to the queue; the initial value is 1, so the mutex is initially unlocked.

`items` is the number of items in the queue, which is also the number of consumer threads that can execute `queue_pop` without blocking. Initially there are no items in the queue.

`spaces` is the number of empty spaces in the queue, which is the number of producer threads that can execute `queue_push` without blocking. Initially the number of spaces is the capacity of the queue, which is `length-1`, as explained in Section 10.1.

Here is the new version of `queue_push`, which is run by producer threads:

```
void queue_push(Queue *queue, int item) {
    semaphore_wait(queue->spaces);
    semaphore_wait(queue->mutex);

    queue->array[queue->next_in] = item;
    queue->next_in = queue_incr(queue, queue->next_in);

    semaphore_signal(queue->mutex);
    semaphore_signal(queue->items);
}
```

Notice that `queue_push` doesn't have to call `queue_full` any more; instead, the semaphore keeps track of how many spaces are available and blocks producers if the queue is full.

Here is the new version of `queue_pop`:

```
int queue_pop(Queue *queue) {
    semaphore_wait(queue->items);
    semaphore_wait(queue->mutex);

    int item = queue->array[queue->next_out];
    queue->next_out = queue_incr(queue, queue->next_out);

    semaphore_signal(queue->mutex);
    semaphore_signal(queue->spaces);

    return item;
}
```

This solution is explained, using pseudo-code, in Chapter 4 of *The Little Book of Semaphores*.

Using the code in the repository for this book, you should be able to compile and run this solution like this:

```
$ make queue_sem
$ ./queue_sem
```

11.3 Make your own semaphores

Any problem that can be solved with semaphores can also be solved with condition variables and mutexes. We can prove that's true by using condition variables and mutexes to implement a semaphore.

Before you go on, you might want to try this as an exercise: write functions that implement the semaphore API in `sem.h` using condition variables and mutexes. In the repository for this book, you'll find my solution in `mysem_soln.c` and `mysem_soln.h`.

If you have trouble getting started, you can use the following structure definition, from my solution, as a hint:

```
typedef struct {
    int value, wakeups;
    Mutex *mutex;
    Cond *cond;
} Semaphore;
```

`value` is the value of the semaphore. `wakeups` counts the number of pending signals; that is, the number of threads that have been woken but have not yet resumed execution. The reason for wakeups is to make sure that our semaphores have Property 3, described in *The Little Book of Semaphores*.

`mutex` provides exclusive access to `value` and `wakeups`; `cond` is the condition variable threads wait on if they wait on the semaphore.

Here is the initialization code for this structure:

```
Semaphore *make_semaphore(int value)
{
    Semaphore *semaphore = check_malloc(sizeof(Semaphore));
    semaphore->value = value;
    semaphore->wakeups = 0;
    semaphore->mutex = make_mutex();
    semaphore->cond = make_cond();
    return semaphore;
}
```

11.3.1 Semaphore implementation

Here is my implementation of semaphores using POSIX mutexes and condition variables:

```
void semaphore_wait(Semaphore *semaphore)
{
    mutex_lock(semaphore->mutex);
    semaphore->value--;

    if (semaphore->value < 0) {
        do {
            cond_wait(semaphore->cond, semaphore->mutex);
        } while (semaphore->wakeups < 1);
        semaphore->wakeups--;
    }
    mutex_unlock(semaphore->mutex);
}
```

When a thread waits on the semaphore, it has to lock the mutex before it decrements `value`. If the value of the semaphore becomes negative, the thread blocks until a “wakeup” is available. While it is blocked, the mutex is unlocked, so another thread can signal.

Here is the code for `semaphore_signal`:

```
void semaphore_signal(Semaphore *semaphore)
{
    mutex_lock(semaphore->mutex);
    semaphore->value++;

    if (semaphore->value <= 0) {
        semaphore->wakeups++;
        cond_signal(semaphore->cond);
    }
    mutex_unlock(semaphore->mutex);
}
```

Again, a thread has to lock the mutex before it increments `value`. If the semaphore was negative, that means threads are waiting, so the signalling thread increments `wakeups` and signals the condition variable.

At this point one of the waiting threads might wake up, but the mutex is still locked until the signalling thread unlocks it.

At that point, one of the waiting threads returns from `cond_wait` and checks

whether a wakeup is still available. If not, it loops and waits on the condition variable again. If so, it decrements `wakeups`, unlocks the mutex, and exits.

One thing about this solution that might not be obvious is the use of a `do...while` loop. Can you figure out why it is not a more conventional `while` loop? What would go wrong?

The problem is that with a `while` loop this implementation would not have Property 3. It would be possible for a thread to signal and then run around and catch its own signal.

With the `do...while` loop, it is guaranteed¹ that when a thread signals, one of the waiting threads will get the signal, even if the signalling thread runs around and gets the mutex before one of the waiting threads resumes.

¹Well, almost. It turns out that a well-timed spurious wakeup (see http://en.wikipedia.org/wiki/Spurious_wakeup) can violate this guarantee.