

Decidability of bisimulation equivalence for equational graphs of finite out-degree

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Abstract

The bisimulation problem for equational graphs of finite out-degree is shown to be decidable. We reduce this problem to the η -bisimulation problem for deterministic rational (vectors of) boolean series on the alphabet of a dpda \mathcal{M} . We then exhibit a complete formal system for deducing equivalent pairs of such vectors.

Keywords: bisimulation; equational graphs; deterministic pushdown automata; rational languages; matrix semi-groups; decidability; complete formal systems.

1 Introduction

1.1 Motivations

1.1.1 Processes

In the context of concurrency theory, several notions of “behaviour of a process” and “behavioural equivalence between processes” have been proposed. Among them, the notion of *bisimulation* equivalence seems to play a prominent role (see [14]). The question of whether this equivalence is *decidable* or not for various classes of infinite processes has been the subject of many works in the last ten years (see [4] and the recent results [10, 22, 11]). The aim of this work is to show decidability of the bisimulation equivalence for the class of all processes defined by pushdown automata with decreasing ϵ -transitions (we assume here that ϵ -transitions are *not* visible, which implies that the graphs of the processes considered here, might have infinite in-degree). This problem was raised in [4] (see Problem 6.2 of this reference) and in [22, article version] (as the bisimulation-problem for processes “of type -1”).

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1.1.2 Infinite graphs

A wide class of graphs enjoying interesting decidability properties has been defined in [6, 2, 3] (see [7] for a survey). In particular it is known that the problem

“are \mathcal{G}, \mathcal{H} isomorphic?”

is decidable for pairs \mathcal{G}, \mathcal{H} of equational graphs. It seems quite natural to investigate whether the problem

“are \mathcal{G}, \mathcal{H} bisimilar?”

is decidable for pairs \mathcal{G}, \mathcal{H} of equational graphs. We show here that this problem is decidable for equational graphs of finite out-degree.

1.1.3 Formal languages

Another classical equivalence relation between processes is the notion of *language* equivalence. The decidability of language equivalence for *deterministic* pushdown automata has been recently established in [19] (see also in [21, 20] shorter expositions of this result). It was first noticed in [1] that, in the case of deterministic processes, language equivalence and bisimulation equivalence are identical. Moreover deterministic pushdown automata can always be normalized (with preservation of the language) in such a way that ϵ -transitions are all decreasing. Hence the main result of this work is a generalisation of the decidability of the equivalence problem for dpda's.

1.1.4 Mathematical generality

The present work is a common generalization of three different results: the results of [22, 11] establishing decidability of the bisimulation equivalence for two non-deterministic sub-classes of the class considered here, and the result of [19] dealing only with deterministic pda's (or processes). More precisely, the present work *extends the notions* developed in [19] so as to obtain