Lecture 15: Game Theory CS486/686 Intro to Artificial Intelligence

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Outline

- Game Theory
- Normal form games
 - Strictly dominated strategies
 - Pure strategy Nash equilibria
 - Mixed Nash equilibria



Multi-agent Decision Making

- Sequential Decision Making
 - Markov Decision Processes
 - Reinforcement Learning
 - Multi-Armed Bandits

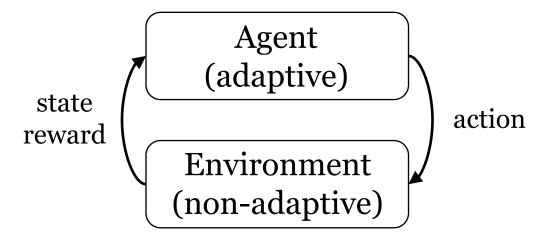
All in single agent environments

- Real world environments: usually more than one agent?
 - Each agent needs to account for other agents' actions/behaviours



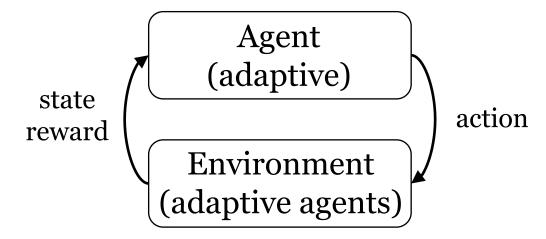
Reinforcement Learning

Single agent



Assumption: stationary transition function $Pr(s_{t+1}|s_t, a_t, s_{t-1}, a_{t-1}, s_0, a_0) = Pr(s_{t+1}|s_t, a_t)$

Multiple agents



Non-stationary transition function $\Pr(s_{t+1}|s_t, a_t, s_{t-1}, a_{t-1}, s_0, a_0) \neq \Pr(s_{t+1}|s_t, a_t)$



Game

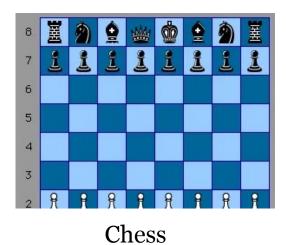
- **Game**: Any scenario where outcomes depend on actions of two or more rational and self-interested players
 - Players (Decision Makers)
 - Agents within the game (observe states and take actions)
 - Rational
 - Agents choose their best actions (unless exploring)
 - Self-interested
 - Only care about their own benefits
 - May/May not harm others



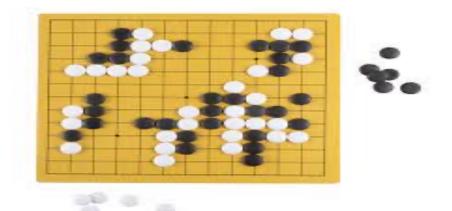
Which of these are games?



Atari



Solitaire



Go



Game Theory

• **Game Theory**: Mathematical model of strategic interactions among multiple rational agents in a game

Interaction:

- One agent directly affects other agent(s)
- Reward for one agent depends on other agent(s)

Strategic:

Agents maximize their reward by taking into account their influence (through actions) on the game

Multiple:

At-least two agents



Game Theory Applications

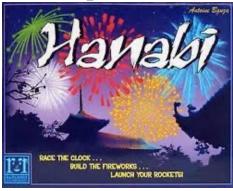
- Auctions
- Diplomacy
- Negotiations
- Sports analytics
- Autonomous Driving
- Conversational agents



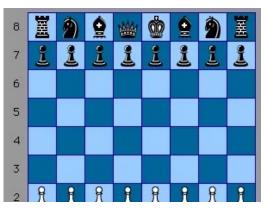
Categorization of Games

- Games can be
 - Cooperative: agents have a common goal
 - **Competitive:** agents have conflicting goals
 - Mixed: in between cooperative and competitive (agents have different goals, but they are not conflicting)

Cooperative



Competitive







Normal Form Games

- Set of **agents**: I = 1, 2, ..., N, where $N \ge 2$
- Set of **actions** for each agent: $A_i = \{a_i^1, ..., a_i^m\}$
 - Game outcome is a **strategy profile (joint action)**: $a = (a_1, ..., a_n)$
 - Total space of joint actions: $a \in \{A_1 \times A_2 \times \cdots \times A_N\}$
- **Reward function** for each agent: R_i : $A \to \Re$, where $A = \{A_1 \times A_2 \times \cdots \times A_N\}$
- No state
- Horizon: h = 1



Example: Even or Odd

Agent 2

One Two

Zero-sum game:

$$\sum_{i=1}^{n} R_i(a_1, ..., a_n) = 0$$

$$I = \{1,2\}$$

 $A_i = \{One, Two\}$
An outcome is (One, Two)
 $R_1(One, Two) = -3$ and $R_2(One, Two) = 3$



Examples of strategic games

Baseball or Soccer

Chicken

 \mathbf{B}

S

Cross Turn

B 2,1 0,0 S 0,0 1,2

Cross Turn

-1,-1	10,0
0,10	5,5









Coordination Game

Anti-Coordination Game



Example: Prisoner's Dilemma







Confess

Don't Confess

Confess

Don't Confess

-5,-5	0,-10
-10,0	-1,-1



Playing a game

- We now know how to describe a game
- Next step Playing the game!
- Recall, agents are rational
 - Let p_i be agent i's beliefs about what its opponents will do
 - Agent i is rational if it chooses to play a_i^* where

$$a_i^* = argmax_{a_i} \sum_{a_i} R(a_i, a_{-i}) p_i(a_{-i})$$

Notation:
$$a_{-i} = (a_1, ..., a_{i-1}, a_{i+1}, ..., a_n)$$



Dominated Strategies

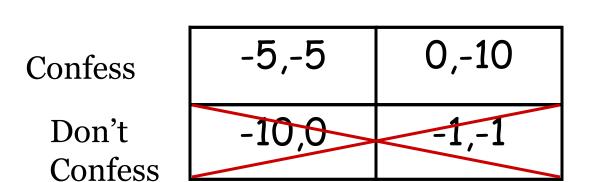
• **Definition**: A strategy s_i is *strictly dominated* if

$$\exists a_i', \forall a_{-i}, R_i(a_i, a_{-i}) < R(a_i', a_{-i})$$

- A rational agent will never play a strictly dominated strategy!
 - This allows us to solve some games!



Example: Prisoner's Dilemma

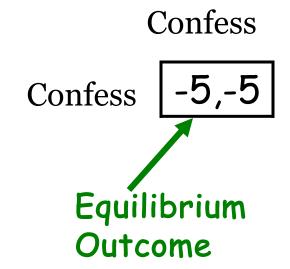


Confess

Confess Don't Confess

Confess





Don't Confess

Strict Dominance does not capture the whole picture

	Α	В	С
Α	0,4	4,0	5,3
В	4,0	0,4	5,3
C	3,5	3,5	6,6

What strict dominance eliminations can we do?



Nash Equilibrium

- Sometimes an agent's best-response depends on the strategies other agents are playing
- A strategy profile, a^* , is a **Nash equilibrium** if no agent has incentive to deviate from its strategy *given that others do not deviate*:

$$\forall i, a_i, R_i(a_i^*, a_{-i}^*) \ge R_i(a_i, a_{-i}^*)$$



Nash Equilibrium

• Equivalently, a^* is a Nash equilibrium iff $\forall i \ a_i^* = argmax_{a_i}R_i(a_i, a_{-i}^*)$

_	Α	В	С
A	0,4	4,0	5,3
В	4,0	0,4	5,3
C	3,5	3,5	6,6

(C,C) is a Nash equilibrium because:

$$R_1(C,C) = \max\{R_1(A,C), R_1(B,C), R_1(C,C)\}$$

AND

$$R_2(C,C) = \max\{R_2(C,A), R_2(C,B), R_2(C,C)\}$$

Exercise 1

What are the Nash Equilibria?

B

S



B

 2,1
 0,0

 0,0
 1,2



S



Exercise 2

What are the Nash Equilibria?

Agent 2

One Two
One 2,-2 -3,3
Agent 1
Two -3,3 4,-4

(Mixed) Nash Equilibria

- Mixed strategy σ_i : σ_i defines a probability distribution over A_i
- Strategy profile: $\sigma = (\sigma_1, ..., \sigma_n)$
- Expected utility: $R_i(\sigma) = \sum_a (\prod_j \sigma(a_j)) R_i(a)$
- Nash Equilibrium: σ^* is a (mixed) Nash equilibrium if

$$\forall i \ R_i(\sigma_i^*, \sigma_{-i}^*) \ge R_i(\sigma_i', \sigma_{-i}^*) \ \forall \sigma_i'$$



Finding Mixed Nash Equilibria

- Two players: $\boldsymbol{\sigma} = (\sigma_1, \sigma_2)$
 - Let $p = \sigma_1$ and $q = \sigma_2$
- At the equilibrium:
 - p^* should be best strategy given q^* : $p^* = argmax_p R_1(p, q^*)$
 - q^* should be best strategy given p^* : $q^* = argmax_qR_2(p^*, q)$
- Solve system of equations:



Exercise 2 Revisited

One A
$$q = Pr(one)$$
 $q = Pr(one)$ $q = Pr(one)$

$$p = Pr(one)$$
 $q = Pr(one)$

How do we determine p and q?

$$R_{A}(p,q) = 2pq - 3p(1-q) - 3(1-p)q + 4(1-p)(1-q)$$

$$R_{B}(p,q) = -2pq + 3p(1-q) + 3(1-p)q - 4(1-p)(1-q)$$

$$\frac{\partial}{\partial p} R_{A}(p,q) = 12q - 7 \rightarrow q = \frac{7}{12}$$

$$\frac{\partial}{\partial q} R_{B}(p,q) = -12p + 7 \rightarrow p = \frac{7}{12}$$



Exercise 3

B S

B 2,1 0,0 **S** 0,0 1,2

This game has 3 Nash equilibria (2 pure strategy Nash equilibria and 1 mixed strategy Nash equilibrium). Find them.

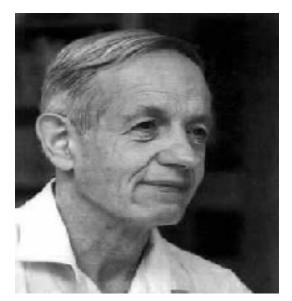


Mixed Nash Equilibrium

• **Theorem** (Nash 1950):

Every game in which the strategy sets $A_1,...,A_n$ have a finite number of elements has a mixed strategy equilibrium.

John Nash Nobel Prize in Economics (1994)





Other Useful Theorems

Theorem: In an n-player pure strategy game, if iterated elimination of strictly dominated strategies eliminates all but the strategies $(a_1^*,...,a_n^*)$ then these strategies are the unique Nash equilibria of the game

• **Theorem:** Any Nash equilibrium will survive iterated elimination of strictly dominated strategies.



Nash Equilibrium

- Interpretations:
 - Focal points, self-enforcing agreements, stable social convention, consequence of rational inference..
- Criticisms
 - They may not be unique
 - Ways of overcoming this: Refinements of equilibrium concept, Mediation, Learning
 - They may be hard to find
 - People don't always behave based on what equilibria would predict

