Assumption: there are no zero column in A and B.
game is non-degenerated.
$\overline{P}_{1} = \left\{ (x, v) \in \mathbb{R}^{m} \times \mathbb{R} \mid x \geq 0, \underline{B}^{T} \times x \leq v \cdot 1 \right\}.$
$\overline{P}_{2} = \{ (y, u) \in \mathbb{R}^{n} \times \mathbb{R} \mid y > 0, Ay \leq u \parallel y^{T} \parallel = 1 \}.$
$X^TBy = (B^Tx)^Ty \leq \max_{\text{equals when } y = [0,, 0]^T} \leq V$
Hence $\max \{x^T B y\} = \max \{(B^T x)\} \leq V$; $x^T B y \leq V \neq y$
Hence $\max \{x^T B y\} = \max \{(B^T x)\} \leq V$; $x^T B y \leq V \neq y$ $\langle x, B y \rangle = \langle B^T x, y \rangle \leq V$.
$X = [X_1, X_2,, X_m]^T$ $X_k = 0$ if $k \in S_1$ is not played.
$(B^T \times)_L = V \Rightarrow \underline{l} \in S_2$ $\underline{y} = [0, 0, \dots, 1, \dots, 0]^T$ is best response to \times .
The Polyhedron P. lives in RMH
it has one equality constraint: \(\frac{1}{2}(\times,\nu)\), \(\tilde{1},\frac{1}{2}(\times)\);=V.
起码m维,由于(x,v). => it is m-dimensional Polyhedron. RM*R.
=) in extreme point of P1, m of the inequality constraints
have to be satisfied with equality.
An extreme point (x, v) Corresponds to a Situation

where $Supp(x) = k$, $k \le m$.
these k inequality constraints are satisfied with equality.
player 2's pure strategies best response to x. > k.
player 2 's pure strategies best response to X. > k. lie. strategy gets V. Because we assume non-degenerated game.
=> no mixed strategy has more than k pure best responses
=> # player 2 's Dure Strategies best response to X. \le \kappa. i.e. strategy gets V.
j.e. = k
Normalized, so that $u = v = 1$.