Math 482: Linear Programming, Fall 2020

Due Friday, September 11, 6PM CST

Homework 2

1. Set up a linear program for the problem below. Do not solve. (Because there will be many similar constraints, it's fine if you just provide an example of each type of constraint, and say "do the same thing for every club" or "do the same thing for every pair of clubs".)

Every student at a certain school is a member of at least one (maybe more) of its five clubs: Athletics Club, Book Club, Chess Club, Drama Club, and Ethics Club. Moreover, each club is very large, and includes at least $\frac{1}{2}$ of the students at the school.

The school gives a "friendship award" to the two clubs that have the largest overlap in students. Every student that's a member of *both* clubs (not just one) will receive the award.

What is the smallest possible fraction of students that can receive the friendship award?

(Hint: to minimize the largest overlap between two clubs, minimize an auxiliary variable z that's greater than or equal to the size of every overlap. You will need to think carefully about what your other variables are in this linear program. Think about a Venn diagram, splitting up the students into 32 possible categories, depending on the set of clubs in which they are members.)

2. The linear program

maximize
$$2x - 2y$$

subject to $x - 3y \le 3$
 $-4x + y \le 4$
 $x - 2y \le 6$
 $x, y \ge 0$

is unbounded. Use the simplex method to find a ray (a starting point (x_0, y_0) and a direction (u, v)) along which every point is a feasible solution, and the objective value increases arbitrarily far.

3. Use the two-phase simplex method to solve

$$\begin{aligned} & \underset{x_1, x_2, x_3 \in \mathbb{R}}{\text{minimize}} & & 3x_1 - x_2 \\ & \text{subject to} & & x_1 + x_2 + & x_3 = 5 \\ & & 2x_1 + x_2 - 2x_3 \geq 6 \\ & & x_1 + x_2 - & x_3 \leq 1 \\ & & x_1, x_2, x_3 \geq 0. \end{aligned}$$

4. Consider the following linear program:

- (a) Perform two iterations of the simplex method using the following pivoting rule: choose the entering variable with the highest reduced cost. When both rows are valid leaving variables (in which case they'll always be tied for the smallest ratio: both ratios will always be 0) choose the basic variable for the first row as the leaving variable.
- (b) Comparing the resulting tableau to the original tableau, argue that the simplex method with this pivoting rule will cycle forever, returning to the same tableau every six steps.
- 5. (Only 4-credit students need to do this problem.)

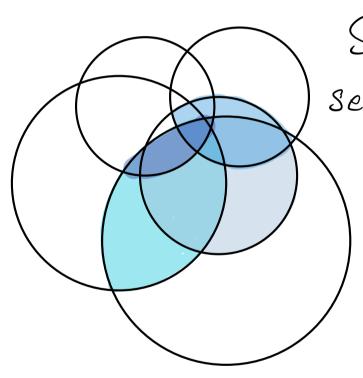
Consider a linear program of the form

$$\begin{array}{ll}
\text{maximize} & \mathbf{c}^{\mathsf{T}} \mathbf{x} \\
\mathbf{x} \in \mathbb{R}^{n} & \text{subject to} & A\mathbf{x} \leq \mathbf{b}.
\end{array}$$

Suppose that points $\mathbf{x} \in \mathbb{R}^n$ and $\mathbf{y} \in \mathbb{R}^n$ are both feasible solutions of this linear program.

- (a) Show that if both **x** and **y** are optimal solutions, then $\frac{1}{2}\mathbf{x} + \frac{1}{2}\mathbf{y}$ is an optimal solution.
- (b) Show that if $\frac{1}{2}\mathbf{x} + \frac{1}{2}\mathbf{y}$ is an optimal solution, then both \mathbf{x} and \mathbf{y} are optimal solutions.

2



Set XE as the number of E

$$\mathcal{X}$$
, $\mathcal{Z} \in \mathbb{R}^{+}$

Sit.
$$x_{inj} \leq z, \forall i, j \in S, i \neq j$$
.
 $x_i \geq \frac{1}{2}, \forall i \in S$

$$\sum_{i \in S} \chi_i - \sum_{i \in S} \sum_{j \in S, j \neq i} \chi_{i \land j} + \sum_{i \in S} \sum_{j \in S} \chi_{i \land j \land k}$$

$$i \in S \qquad i \in S \quad j \in S, j \neq i \qquad i \notin S \quad j \in S \quad k \in S \quad k \neq i \quad k \neq i \quad k \neq j$$

$$+ X_{ANBACADAE} = 1.$$
 $X, 2 > 0$

2.
$$\max_{x,y \in \mathbb{R}} 2x - 2y$$
subject to $x - 3y \le 3$
 $-4x + y \le 4$
 $x - 2y \le 6$
 $x, y \ge 0$
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Pivot on X, remove S,

	χ	y	S.	S.	53	
X	/	- 3	/	0	0	3
Sz	0	-//	4	1	0	16
 S}	0	/	-/	U	/	3
-8	0	4	-2	0	0	-6

 $(3+8,\frac{1}{3}8.0.16.3)$ (5,51,51,53)=(0.16.3)

Di	put	on Y	rev	no VQ	S3	
	χ	γ '	S.	Sz	S ₃	esign
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Sz	0	0	-7	(//	49 is unbounded.
<u> </u>	0	/ /	1-//	U	/	3
- 2	0	O	2	0	-4	-18

 $(x, y) = (12, 3) \quad (S_1, S_1, S_2, S_3) = (0, 49, 0)$ $(12 + 8, 3 + \frac{1}{2}S, 0, 49, 0) \quad S_3 = 0$ $2 = 24 + 28 - 6 - 8 = 18 + 8 \quad x = 12 + 28,$ = (2, 1)

$ray: (X_0, Y_0) = (/2,3) (u,v) = (/,\frac{1}{2})$

 $3. \qquad \chi_{3} = 3 - \chi_{1} - \chi_{2}$ $\underset{x_{1}, x_{2}, x_{3} \in \mathbb{R}}{\text{minimize}} \quad 3x_{1} - x_{2}$ $\underset{x_{1}, x_{2}, x_{3} \in \mathbb{R}}{\text{minimize}} \quad 3x_{1} + x_{2} + x_{3} = 5$ $\underset{x_{1}, x_{1}, \chi_{1}, \chi_{1$

subject to $x_1 + x_2 + x_3 = 5$ $2x_1 + x_2 - 2x_3 \ge 6$ $x_1 + x_2 - x_3 \le 1$ $x_1, x_2, x_3 \ge 0.$ $x_1 + x_2 + x_3 = 5$ $2x_1 + x_2 - x_3 \le 1$ $x_1, x_2, x_3 \ge 0.$ $x_1 + x_2 + x_3 = 5$ $x_1 + x_2 - x_3 \le 1$ $x_1, x_2, x_3 \ge 0.$ $x_1 + x_2 + x_3 = 5$ $x_1 + x_2 + x_3 = 5$ $x_1 + x_2 + x_3 = 5$ $x_1 + x_2 + x_3 = 6$ $x_2 + x_3 = 6$ $x_1 + x_2 + x_3 = 6$ $x_2 + x_3 = 6$ $x_1 + x_2 + x_3 = 6$ $x_1 + x_2 + x_3 = 6$ $x_2 + x_3 = 6$ $x_1 + x_2 + x_3 = 6$ $x_2 + x_3 = 6$ $x_1 + x_2 + x_3 = 6$ $x_2 + x_3 = 6$ $x_1 + x_2 + x_3 = 6$ $x_2 + x_3 = 6$ $x_1 + x_2 + x_3 = 6$ $x_2 + x_3 = 6$ $x_1 + x_2 + x_3 = 6$ $x_2 + x_3 = 6$ $x_3 + x_4 = 6$ $x_4 + x_4 = 6$ $x_3 + x_4 = 6$ $x_4 + x_4 = 6$

X,,X,, X,, X,, X,, 20

no feasible solution

determine whether it has

feasible solution.

	χ,	χ,	χ,	S,	2,	χ,	a Xa	X3	
χa	/	/	/	0	0	/	0	0	5
7/2	2				0		/	0	6
73ª	/	/	-/	0	/	0	0	/	/
- 2	3	-/	0	O	0	U	0	0	0
- Za	0	0	0	0	0	/	/	/	U

	χ,	χ,	χ,	S,	2,	χ,	χ,	X ₃		
χa	/	/	/	0	0	/	0	0	5	Pivot on Ti
X2ª	2	/	-2	-/	0	0	/	0	6	7a
χ, α	/	/	-/	0	/	0	0	/	/	remove 13
- 2	3	-/	0	O	0	U	0	0	0	
- Zª	1-4	-3	2	/	-/	0	0	0	-12	

	χ,	χ,	Χ³	Sı	2,	X	X	X3		
χa	0	0	2	0	-/	/	0	-/	4	~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~
7/2°	0	-/	0	-/	-2	0	/	-2	4	Pivol on X3
γ,	/	/	-/	0	/	0	0	/	/	remove X,ª
- 2	0	-4	3	O	-3	U	0	-3	-3	
- Zª	0	/ .	-2	/	3	0	0	4	8-1	

					2,				
Χ3	0	0	/	0	- <u>/</u> -2	/_	0	$-\frac{1}{2}$	2
7/2°	0	-/	0	_/	-2	0	/	-2	4
γ,	/	/	0	0	4	1/2	0	1/2	3
- 2	0	-4	3	0	-3	0	0	-3	-3
- Za	0	/	0	/	2	/	0	3	-4

 $min 2^a = 4 = 3$ no feasible solution.

4. (1) max
$$X_1 - 3X_2 - 2X_4$$
 $X_1, X_1, X_2, X_3, X_4, S_1, S_2 \in \mathbb{R}$
 $\frac{1}{2}X_1 - \frac{3}{2}X_2 - \frac{3}{2}X_3 + \frac{7}{2}X_4 + S_1 = 0$
 $\frac{1}{2}X_1 - \frac{3}{2}X_2 - \frac{1}{2}X_3 + \frac{1}{2}X_4 + S_2 = 0$
 $X_1, X_2, X_3, X_4, S_1, S_2 \geq 0$
 $X_2, \frac{1}{2} - \frac{3}{2} - \frac{1}{2} = \frac{1}{2} = 0 = 1 = 0$
 $X_1, X_2, X_3, X_4, S_1, S_2 = 0$
 $X_2, X_2, X_3, X_4, S_1, S_2 = 0$
 X_1, X_2, X_3, X_4, S

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-2	0	-2	0	0	/	-3	0	
	Pivo	H S	, , ſ	emo	ve	Χş		

Pivut Sz remove X4

we are back!!!

=) cycle forever and back to orginal tableau every six steps.