

Given bipartite graph (X, Y, E)

$$\max_{x \in \mathbb{R}^{|E|}} \sum_{i,j: (i,j) \in E} x_{ij}$$

$$\text{s.t. } \sum_{j: (i,j) \in E} x_{ij} \leq 1 \quad \forall i \in X.$$

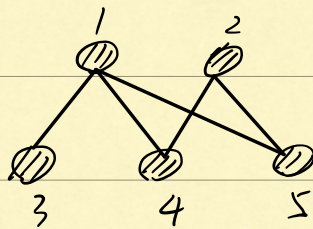
$$\sum_{i: (i,j) \in E} x_{ij} \leq 1 \quad \forall j \in Y.$$

$Ax \leq 1$ incidence matrix $(|X|+|Y|) \times |E|$

$A_{ve} = 1$ if v is an endpoint of e .

$A_{ve} = 0$ others.

Example



$$A = \begin{matrix} & \begin{matrix} (1,3) & (1,4) & (1,5) & (2,4) & (2,5) \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \\ 5 \end{matrix} & \begin{bmatrix} 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \end{bmatrix} \end{matrix}$$

if A is totally unimodular ^{by theorem previous} \Rightarrow solution is integral.

Theorem: incidence matrix A is totally unimodular.

det of submatrix = 0/1/-1.

Proof: $k=1$: $\det = 0/1$.

$k > 1$: $\det \begin{bmatrix} 0 & \square \\ 0 & \square \\ \vdots & \square \\ 0 & \square \end{bmatrix} = 0$
one col with all 0

one col with one 1: $\det \begin{bmatrix} 0 & \square \\ 0 & \square \\ \vdots & \square \\ 1 & \square \end{bmatrix} = (-1)^i \det \begin{bmatrix} \square \\ \square \\ \vdots \\ \square \end{bmatrix}_{(k-1) \times (k-1)}$

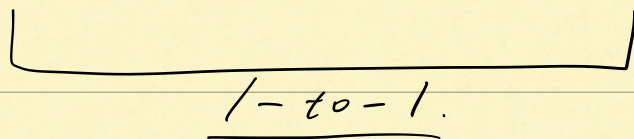
$\det \begin{bmatrix} 1 & & \\ \vdots & \square & \\ 0 & & \end{bmatrix}$
 one is in X , one is in Y
 all col with two 1. $\begin{bmatrix} 1 & & \\ \vdots & \square & \\ 0 & & \end{bmatrix}$ add all X 's rows into first row
 add all Y 's rows into last row.
 more than two is impossible. $\Rightarrow \det = 0$.

$$\begin{aligned}
 \text{Dual } \min \sum_{i \in X \cup Y} u_i \\
 \text{s.t. } u_i + u_j \geq 1, (i, j) \in E \\
 u \geq 0
 \end{aligned}$$

u describing a set of S : $i \in S$ when $u_i = 1$, $i \notin S$ when $u_i = 0$.

Vertex cover: Set of vertices that cover every edge (include at least one endpoint of each edge).
 Looking for smallest S that covers every edge.

S is a vertex cover, M is a matching.



$\# \text{ edge in } M \leq \# \text{ vertices in } S$. (Weak duality).

Theorem: $\# \text{ edge in Maximum matching}$

$= \# \text{ vertices in min vertex cover.}$

Proof: Strong duality, total unimodularity.

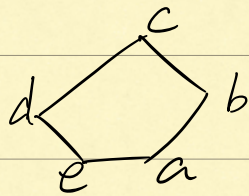
General Form: (V, E) .

Matching: share no endpoints.

Vertex cover: set of vertices covering all edges.

there is no bipartiteness \Rightarrow may not totally unimodular.

Example:



Max matching is of size 2.

$(a, b), (c, d)$

Min vertex cover is 3 vertices.

$\{a, c, d\}$.

Max matching LP \Rightarrow opt value = 2.5

Min vertex cover \Rightarrow opt value = 2.5