

Feasibility problem

Given A and b , is $\{x \in \mathbb{R}^n; Ax \leq b\}$ nonempty?

1. Put into equational form.
2. Find a basic solution (not necessarily feasible).
3. invent an objective function that gives a dual feasible tableau.
4. Use dual simplex method to make it primal feasible.

feasibility problem is as hard as optimization.

$$\begin{cases} \text{optimization} & \max \{C^T x : Ax \leq b, x \geq 0\} \\ \text{feasible} & \{(x, u) \in \mathbb{R}^{n+m} : C^T x = u^T b, Ax \leq b, u^T A \geq C^T, \\ & x \geq 0, u \geq 0\} \end{cases}$$

Fourier-Motzkin elimination.

Simpler but slow to prove infeasible

idea: Reduce on n -var problem to an equivalent $n-1$ var problem. Repeat left with a 1 var feasibility problem.

x_1, \dots, x_n , want to eliminate x_n

①. solve each inequality for x_n

$$x_n \leq \frac{b - a_1 x_1 - a_2 x_2 - \dots - a_{n-1} x_{n-1}}{a_n} \text{ or } x_n \geq \frac{b - a_1 x_1 - a_2 x_2 - \dots - a_{n-1} x_{n-1}}{a_n}$$

end up with a collection of lower and upper

bounds $x_n \geq L_1, x_n \geq L_2, \dots, x_n \geq L_k, x_n \leq U_1, \dots, x_n \leq U_m$.

it is possible to find x_n if and only if $\max\{L_1, \dots, L_k\} \leq \min\{U_1, \dots, U_m\}$.

$L_i \leq x_n \leq U_j$ for all i, j .

We get new system $L_1 \leq U_1, L_1 \leq U_2, \dots, L_1 \leq U_m$

$$L_2 \leq U_1, \dots$$

$$L_k \leq U_1, \dots, L_k \leq U_m$$

the system only have $n-1$ vars.

Example:

$$\begin{cases} x - y \geq 3 \\ x + 2y \geq 4 \\ x + y \leq 7 \\ x, y \geq 0 \end{cases} \xrightarrow{\text{Eliminate } y} \begin{cases} y \leq x - 3 \\ y \geq \frac{1}{2}(4 - x) \\ y \leq 7 - x \\ y \geq 0, x \geq 0 \end{cases}$$

$$\begin{cases} 0 \leq 7-x \\ 0 \leq x-3 \\ \frac{1}{2}(4-x) \leq 7-x \\ \frac{1}{2}(4-x) \leq x-3 \\ x \geq 0 \end{cases} \longrightarrow \begin{cases} x \leq 7 \\ 3 \leq x \\ x \leq 10 \\ \frac{10}{3} \leq x \\ x \geq 0 \end{cases} \longrightarrow \frac{10}{3} \leq x \leq 7$$

this system is feasible. to find a y pick $x=4$

$$\Rightarrow \begin{cases} y \leq 1 \\ y \geq 0 \\ y \leq 3 \\ y \geq 0 \end{cases} \Rightarrow 0 \leq y \leq 1. \quad (1.4) \text{ is feasible.}$$

Complexity

How many inequalities?

Worst behavior case: lower bound and upper bound split exactly in half each time.

Starting with 8 inequalities in n vars:

$2^{a_0} = 2^3$	$8 = 4 + 4$	n
$(\frac{2^3}{2})^2 = 2^4 = 2^{a_1}$	$4 \times 4 = 16 = 8 + 8$	$n-1$
$(\frac{2^4}{2})^2 = 2^6 = 2^{a_2}$	$8 \times 8 = 64 = 32 + 32$	$n-2$
$2^{10} = 2^{a_3}$		

$$2^{18} = 2^{a_4}$$

3, 4, 6, 10, 18

2^{2^k+2} inequalities $n-k$ vars.

$a_{n+1} = 2(a_n - 1)$

$a_{n+1} - 2 = 2(a_n - 2) \Rightarrow a_n = 2^k + 2$

we can change all inequalities into " \leq "
 and the coefficient of x_n is "1" or "-1"
 add the " $+x_n$ " inequality and " $-x_n$ " inequality.

$$\begin{cases} x-y \geq 3 \\ x+2y \geq 4 \\ x+y \leq 7 \\ x, y \geq 0 \end{cases} \longrightarrow \begin{cases} -x+y \leq -3 \\ -x-2y \leq -4 \\ x+y \leq 7 \\ -x \leq 0 \\ -y \leq 0 \end{cases} \longrightarrow \begin{cases} 2y \leq 4 \\ -y \leq 3 \\ y \leq 7 \end{cases} \longrightarrow -3 \leq y \leq 2$$

pick $y=1$ $\begin{cases} -x \leq -4 \\ -x \leq -2 \\ x \leq 6 \\ -x \leq 0 \end{cases} \longrightarrow 4 \leq x \leq 6$ pick $x=4$.
 $(1, 4)$ is feasible.

Farkas's Lemma.

For any system of inequalities $Ax \leq b$, either
 there exists a feasible solution x , or we can
 take a linear comb of inequalities to derive

a contradiction: there exists vector $u \geq 0$ such
that $u^T A = 0$, but $u^T b < 0$