

Two-phase dual simplex method.

nature tableau is not primal feasible.

we want to find a basic feasible solution to begin with.

1. First optimize an artificial obj function using
(such as $x_1 + x_2 + \dots + x_n$)
dual simplex

the final tableau will be primal feasible for any obj function.

2. Use primal simplex method starting from final tableau. Add in original obj. row-reduce.

The method above depends on slack vars.

what if we are given constraints in $Ax=b, x \geq 0$.

1. First row reduce $Ax=b$, ignoring $x \geq 0$

\Rightarrow Arrive at a basic solution.

2. Now pick an artificial obj function that gives you a dual feasible tableau (ex: sum of nonbasic vars). Use dual simplex method.

Ex: Find an x . s.t.
$$\begin{cases} x_1 + 2x_2 - 3x_3 + x_4 = 5 \\ 2x_1 + 4x_2 + x_3 - 2x_4 = 2 \\ x_1, x_2, x_3, x_4 \geq 0 \end{cases}$$

row reduce
$$\begin{cases} x_1 + 2x_2 - \frac{5}{7}x_4 = \frac{11}{7} \\ x_3 - \frac{4}{7}x_4 = -\frac{8}{7} \\ x_1, x_2, x_3, x_4 \geq 0. \end{cases}$$

Basic vars $(x_1, x_3) = (\frac{11}{7}, -\frac{8}{7})$

Use Obj func: $x_2 + x_4$

	x_1	x_2	x_3	x_4	
x_1	1	2	0	$-\frac{5}{7}$	$\frac{11}{7}$
x_3	0	0	1	$-\frac{4}{7}$	$-\frac{8}{7}$
$-z$	0	1	0	1	0

then apply dual simplex method to find a primal feasible tableau.

Adding constraints.

min $x + y$

s.t. $-2x - y + S_1 = -6$

$-3x - y + S_2 = -7$ final tableau \rightarrow

$-x - 2y + S_3 = -9$

$x, y, S_1, S_2, S_3 \geq 0.$

	X	Y	S_1	S_2	S_3	
X	1	0	$-\frac{2}{3}$	0	$\frac{1}{3}$	1
S_2	0	0	$-\frac{5}{3}$	1	$\frac{1}{3}$	0
Y	0	1	$\frac{1}{3}$	0	$-\frac{2}{3}$	4
-Z	0	0	$\frac{1}{3}$	0	$\frac{1}{3}$	-5

we add an additional constraint $x \geq 2.$

$\therefore -x + S_4 = -2.$

Just add an row into the final tableau

	X	Y	S_1	S_2	S_3	S_4	
X	1	0	$-\frac{2}{3}$	0	$\frac{1}{3}$	0	1
S_2	0	0	$-\frac{5}{3}$	1	$\frac{1}{3}$	0	0
Y	0	1	$\frac{1}{3}$	0	$-\frac{2}{3}$	0	4
S_4	-1	0	0	0	0	1	-2
-Z	0	0	$\frac{1}{3}$	0	$\frac{1}{3}$	0	-5

row reduce \rightarrow

	X	Y	S_1	S_2	S_3	S_4	
X	1	0	$-\frac{2}{3}$	0	$\frac{1}{3}$	0	1
S_2	0	0	$-\frac{5}{3}$	1	$\frac{1}{3}$	0	0
Y	0	1	$\frac{1}{3}$	0	$-\frac{2}{3}$	0	4
S_4	0	0	$-\frac{2}{3}$	0	$\frac{1}{3}$	1	-1
-Z	0	0	$\frac{1}{3}$	0	$\frac{1}{3}$	0	-5

then use dual simplex method.

Shadow costs

LP: $\max C^T x$

$$\text{s.t. } Ax \leq b$$

$$x \geq 0$$

we want a small change to either c or b

Can we predict how opt sol changes.

Example: Max $2x + 3y$

$$\text{s.t. } -x + y \leq 3$$

$$x - 2y \leq 2$$

$$x + y \leq 7$$

$$x, y \geq 0$$

Opt sol is $(x, y) = (2, 5)$ opt value is 19.

Let's change $2x + 3y \rightarrow 2.1x + 3y$.

opt sol is still $(2, 5)$ opt value is 19.2

$$2x + 3y \rightarrow (2 + \delta)x + 3y$$

obj value at $(2, 5)$ is $(19 + 2\delta)$

is the lower bound on new opt value

If we change $x + y \leq 7 \rightarrow x + y \leq 7 + \delta$.

Prediction: obj val change to $19 \pm \left(\frac{5}{2}\right)\delta$.

$w = \frac{5}{2}$ represents the sensitivity of obj val to changes in corresponding constraint.

Dual vars are often called "shadow cost"

$\frac{5}{2}$ is the cost of a change in $x+y \leq 7$.