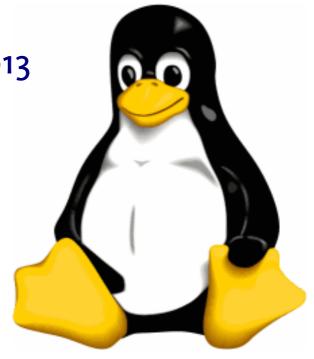


Automotive Linux Summit Fall 2013

Marc Kleine-Budde Pengutronix <mkl@pengutronix.de>

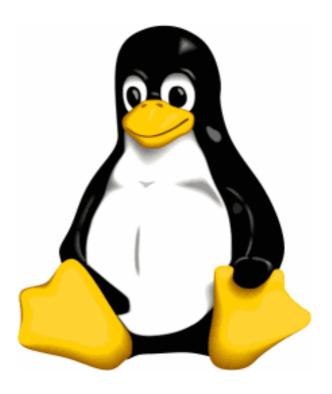






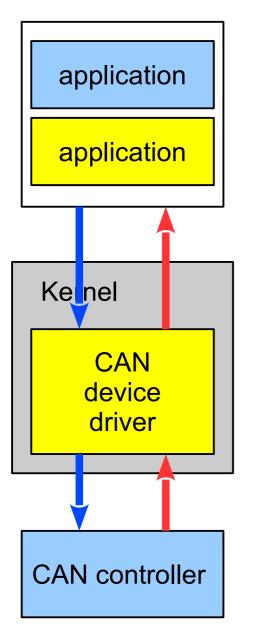
Overview

- History of CAN in Linux
- Linux networking subsystem
- CAN device driver
- Userspace example









- No standard Linux CAN driver model, all character device based
- Only single application at a time
- Higher level protocols and filtering have to be implemented in application
- CAN hardware vendor provide own driver
- Change of hardware vendor urges adaptation of CAN application





HTTP

FTP

...

----- Application layer

Linux Socket Layer

Socket layer

PF_INET

TCP

UDP



Linux Socket Layer

routing and packet scheduler



eth1



----- Packe



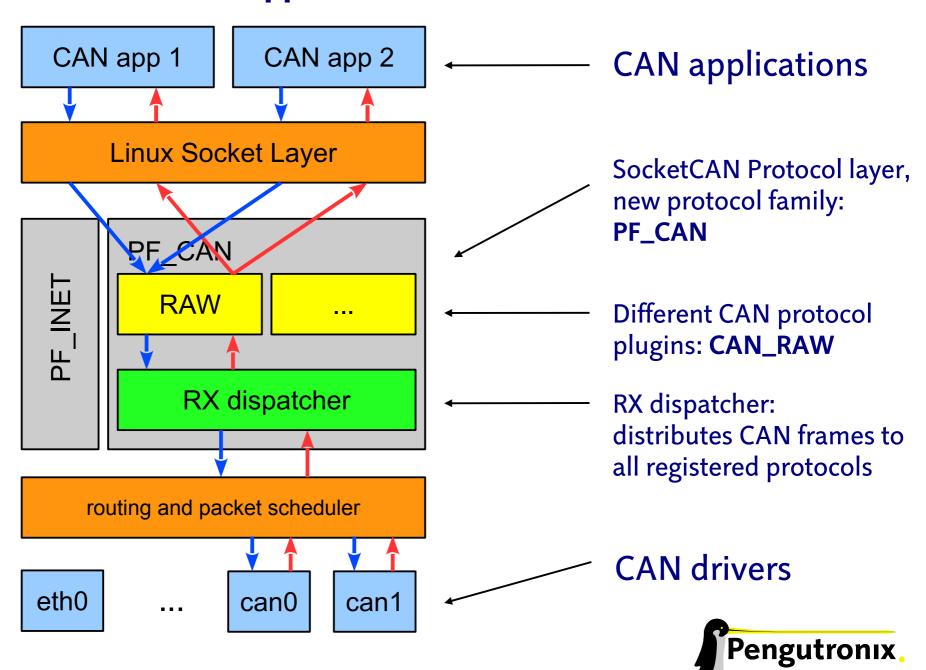
Packet layer

Protocol layer

Network driver layer



A socket-based approach





Multi-application support

- For CAN_RAW each socket can specify a filter list
- RX dispatcher implements complex filtering
- Received CAN frames are transmitted to all CAN protocols that have a matching filter
- Local originated CAN frames are looped back into the RX queue





Multi-application support - Why to loopback?

- Consider two embedded systems, each running a SocketCAN application
- They exchange messages via the CAN bus
- If these applications run on the same system, they still have to see each other's CAN messages
- Put CAN frame into RX queue after transmission has been completed
- For best results, to preserve sequence of frames, do echoing in TX complete interrupt handler





A socket-based approach - the advantages

- Use of existing and established POSIX socket API
- New protocol family PF_CAN is developed against established abstractions
- Communication protocols and filtering can be implemented inside the kernel
- CAN network device driver implements standard Ethernet driver model
- Support for multi-user and multi-application possible





Drawbacks and limitations

- More memory overhead, the Linux network stack has been designed for much larger Ethernet frames:
 64 Bytes (min. Ethernet frame) vs.
 8 byte (max. CAN frame)
- The "packet scheduler" is a shared resource among all networking devices
- Heavy Ethernet traffic can lead to delays in CAN traffic.
- No Support for hardware filtering (yet)





CAN networking device drivers

- Initialize and configure hardware
- Receive incoming CAN frames from hardware and push them into upper layer
- Obtain outgoing frame from upper layer and transmit to wire
- Almost identical to Ethernet drivers, but handle CAN instead of Ethernet frames. Make use of existing Ethernet driver model!
- Define bit rate constraints of CAN hardware in clock rate independent way



Applications and the CAN_RAW protocol

- Simplest method to access the CAN bus
- Programming interface similar to character device drivers, transfer whole CAN frames
- First create a CAN_RAW socket, then bind() to a CAN interface. Use standard systems calls to read and write CAN frames



SocketCAN - struct can_frame





Conclusion + Outlook

- Multi-application + Multi-user POSIX socket API to send and receive raw CAN frames
- Standard driver model known from Ethernet drivers
- Kernel internal infrastructure to filter, send and receive CAN frames to implement more complex protocols
- More CAN protocols
 - BCM Broadcast Manager
 - CANGW CAN Gateway
 - ISOTP ISO 15765-2
 - J1939





Questions?





Thanks!

More information:

- Inside the Kernel: Documentation/networking/can.txt
- Mailing-list: linux-can@vger.kernel.org
- Project Homepage: http://gitorious.org/linux-can
 - Upstream git
 - Userspace tools
 - Support for older Kernels

