

# CS 449: Lab 2

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adapted from CMU 15-213 recitation 2

# Announcements



My email: wel104@pitt.edu



Office hour: Mon/Wed 2:30–4:30



Data lab's deadline is extended

# Understanding bit-wise operator

- Output 1 if exactly one 1?
- Preserve 1/0? How to choose the mask?
- Extract MSD?  $\sim$

<b>&amp;</b>	0	1
0	<b>0</b>	<b>0</b>
1	<b>0</b>	<b>1</b>

← **AND (&)** outputs a 1 only when both input bits are 1.

**OR (|)** outputs a 1 when either input bit is 1. →

<b> </b>	0	1
0	<b>0</b>	<b>1</b>
1	<b>1</b>	<b>1</b>

<b>^</b>	0	1
0	<b>0</b>	<b>1</b>
1	<b>1</b>	<b>0</b>

← **XOR (^)** outputs a 1 when either input is *exclusively* 1.

**NOT (~)** outputs the opposite of its input. →

<b>~</b>	
0	<b>1</b>
1	<b>0</b>

# Great Realities

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1

Ints are not integers,  
floats are not reals!

2

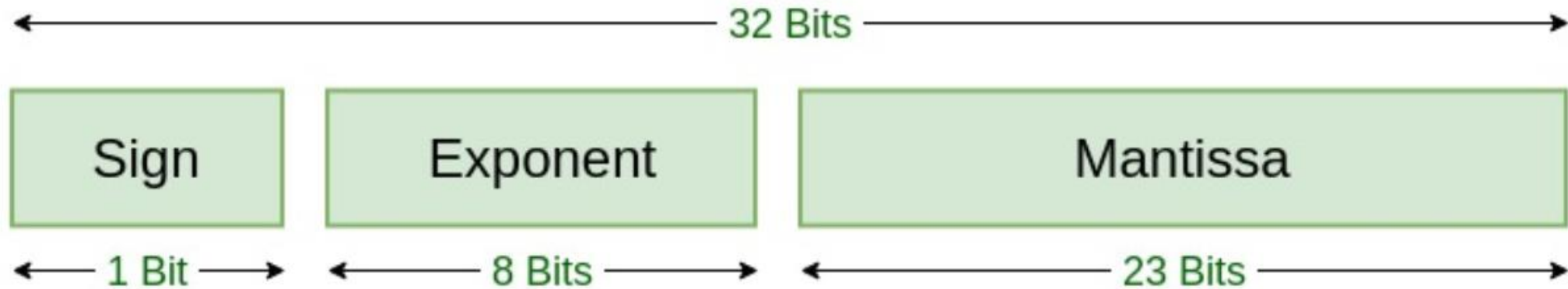
Although you're not likely to  
write assembly (cs 447), it's  
good for understanding  
machine-level executions

3

Don't feel bad if you  
think data lab is hard,  
this is still a  
programming class

# Single-precision

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Single Precision  
IEEE 754 Floating-Point Standard

**32.2 = ?**

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$$32 = 2^5 + 0$$

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$$.2^2 = 0.4 \quad \dots^2 = 0.8 \quad \dots^2 = (1).6 \quad \dots^2 = (1).2 \quad \dots^2 = 0.4$$

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$$100000.\underline{001100110011}\dots \quad 100000 = 1.00000 \cdot 2^5$$

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$$5 + 2^8 - 1 = 132 = 10000100 \leftarrow \text{we get exponent}$$

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$$32.2 > 0 \rightarrow \text{we get sign} = 0$$

---

$$\text{Mantissa} = 0000000\underline{11}\dots$$

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$$0 \ 10000100 \ 00000001100110011001100$$

# 24.0 = ?

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- "Both the argument and result are passed as unsigned int's, but they are to be interpreted as the bit-level representations of single-precision floating point values"
- $11000 = 1.1000 * 2 \text{ to the ?}$  What operator does this remind you of?
- Recall: use `fshow()` to help you...or go through the basics

# Power-of-2 with SHIFT: an optimization of runtime

## Multiplication

- $U \ll k$  gives  $u * 2^k$

## Division

- Unsigned (logical shift)  $U \gg k$  gives  $\text{floor}(u/2^k)$
- ~~Signed (arithmetic shift)  $X \gg k$  gives  $\text{RoundDown}(x/2^k)$~~

What could be wrong here?

What about addition/subtraction?



A whole area of modular arithmetic  
algorithms here...

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# Modular arithmetic (machine implementation), overflow

Visualizing  $\text{Add}_w(u, v) = u + ((2^w - v) \bmod 2^w)$

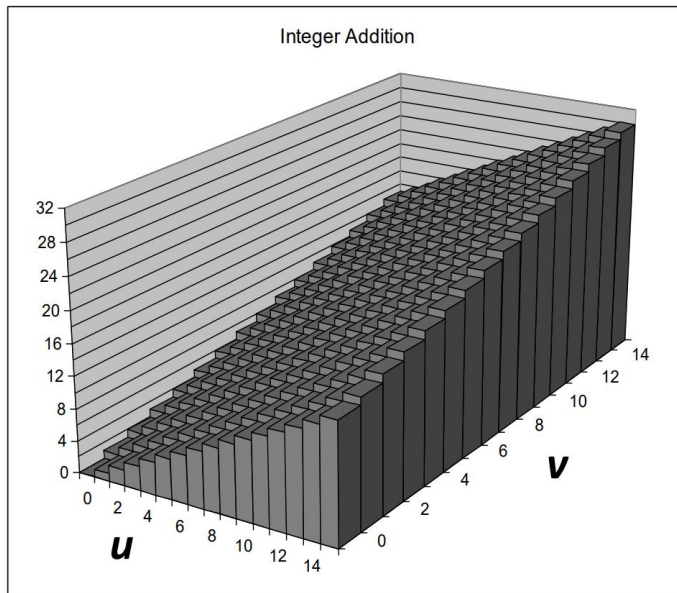
$$= u - v \bmod 2^w$$



$2^{-(w+1)}$  ---overflow-->  $2^{-w}$

$$\text{Uadd}_w(u, v) = (u + v) \bmod 2^w$$

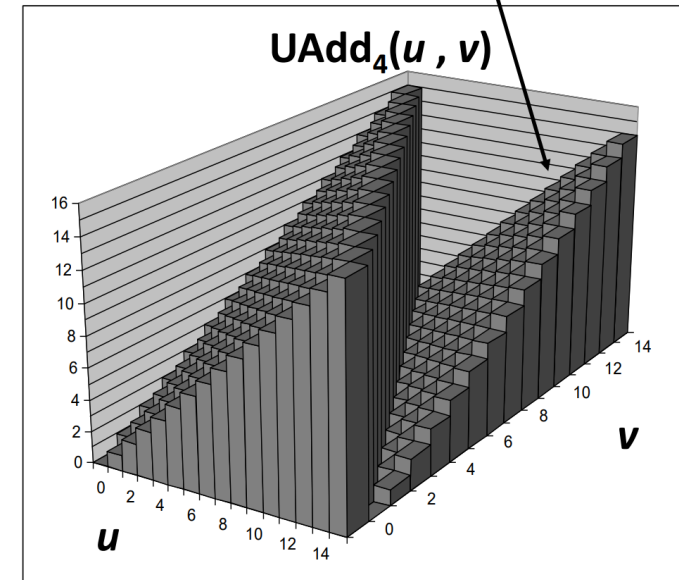
**Add<sub>4</sub>(u, v)**



$$2^5 = 32$$

**Overflow**

**UAdd<sub>4</sub>(u, v)**



$$2^4 = 16$$

# Memory

- A 1-dimensional array of BYTES
- Array index = address of the byte
- More than one byte? What does it look like?

# Example

Computer		Programmers		
Address	Content	Name	Type	Value
90000000	00	sum	int (4 bytes)	000000FF (255 <sub>10</sub> )
90000001	00			
90000002	00			
90000003	FF			
90000004	FF	age	short (2 bytes)	FFFF (-1 <sub>10</sub> )
90000005	FF			
90000006	1F	average	double (8 bytes)	1FFFFFFFFFFFFFFFFF (4.45015E-308 <sub>10</sub> )
90000007	FF			
90000008	FF			
90000009	FF			
9000000A	FF			
9000000B	FF			
9000000C	FF	ptrSum	int* (4 bytes)	90000000
9000000D	FF			
9000000E	90			
9000000F	00			
90000010	00			
90000011	00			

Note: All numbers in hexadecimal

# Try it yourself! -- sizeof()

Char/unsigned  
char

Int/unsigned int

Short/unsigned  
short

Long/unsigned  
long

Double...Is there  
unsigned double?

Float...Is there  
unsigned float?

Pointer\*

Even struct,  
Union, enum!

# Memory (cont)

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- Remember 447? (Big Endian, little endian)

**Big Endian**

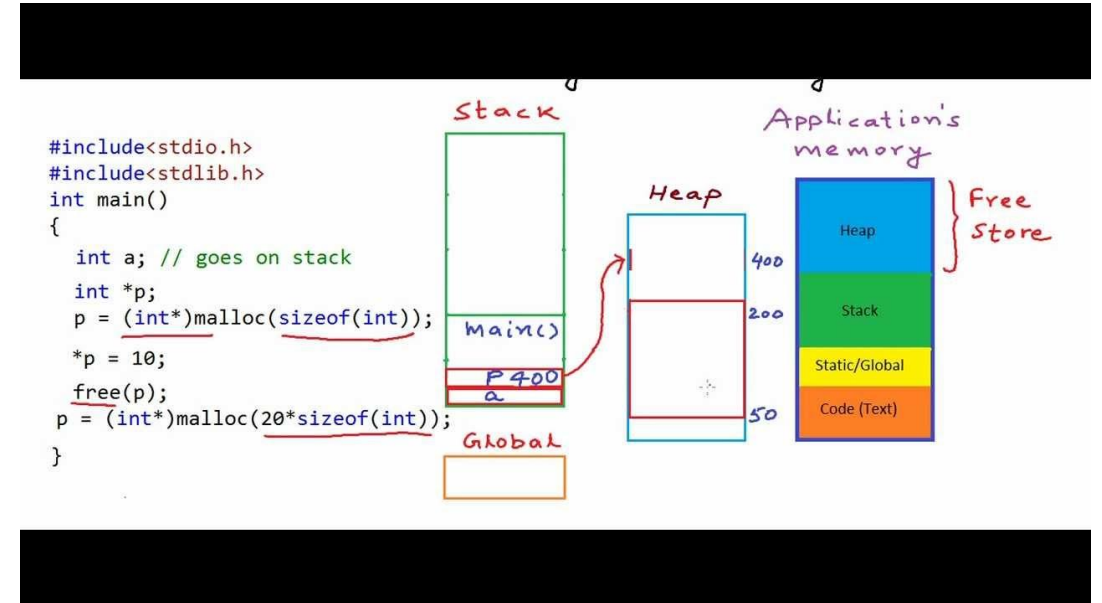
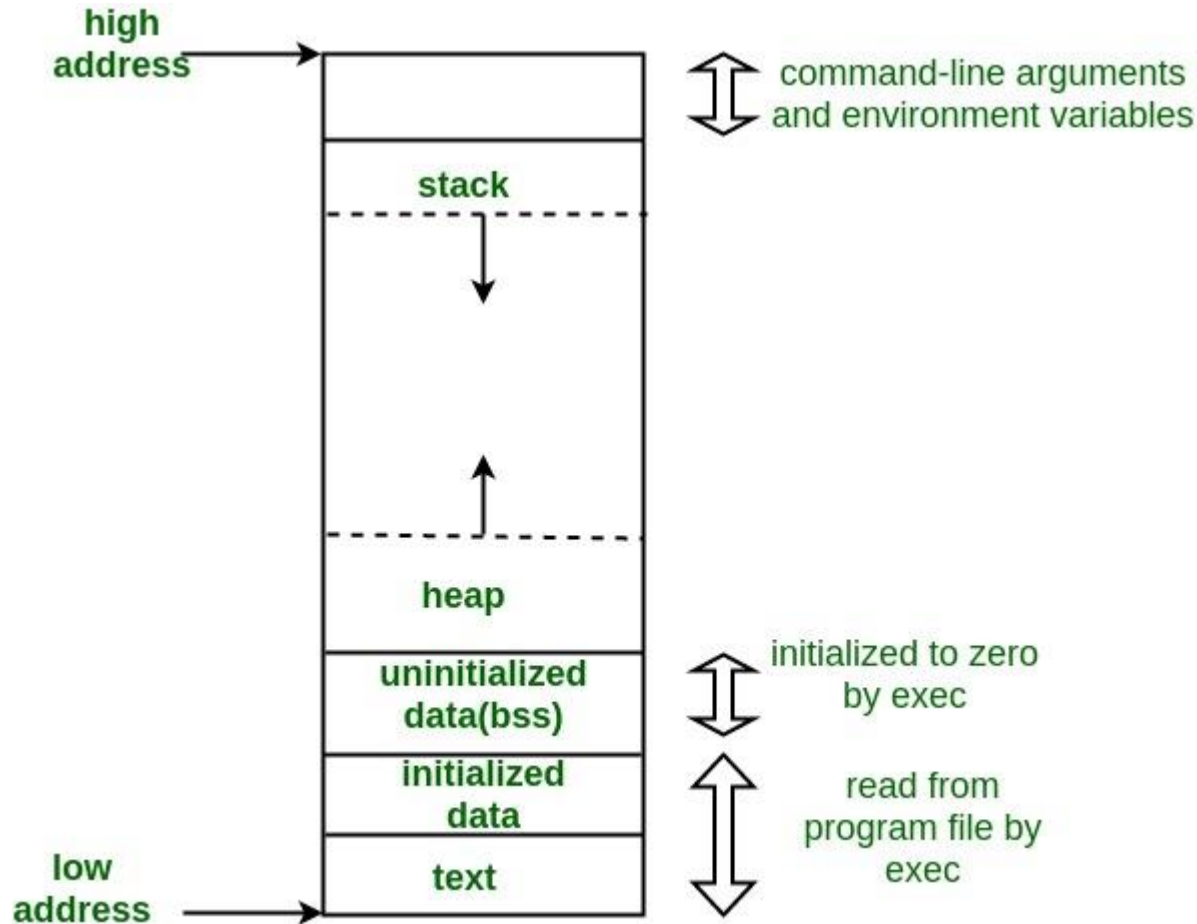
		0x100	0x101	0x102	0x103		
		01	23	45	67		

**Little Endian**

		0x100	0x101	0x102	0x103		
		67	45	23	01		

- Programs refer to data by address
- System provide "private" address space to processes\*

# Memory diagram (more on this later in the semester)





# Pointers: & and \*

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- Remember...memory has addresses
- Variable to hold address
- &: reference operator
- \*: dereference operator
- Show me some code
  - <https://colab.research.google.com/drive/1Bi0yKj4ueKj7frQ9ewCKhGehhhhQihv2>
- Wait, what about strings?



### Exercise:

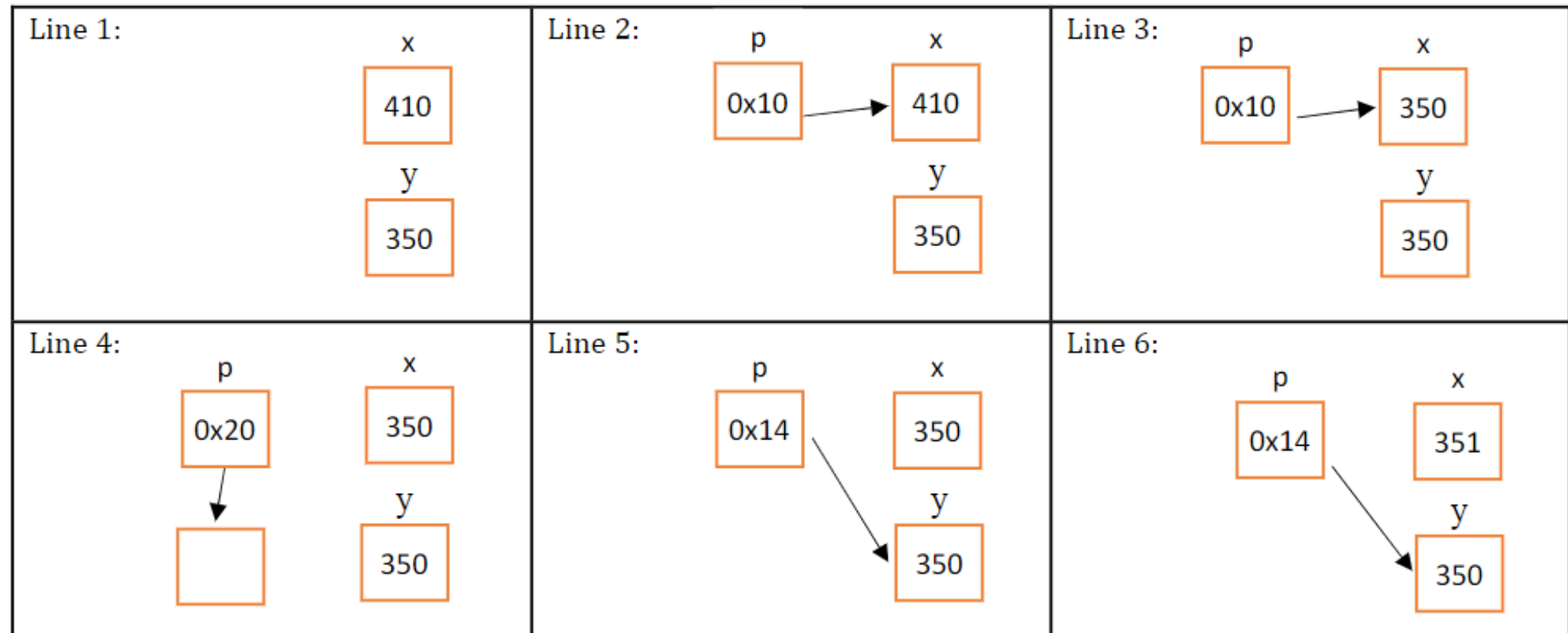
Draw out the memory diagram after sequential execution of each of the lines below:

```
int main(int argc, char **argv) {  
    int x = 410, y = 350;    // assume &x = 0x10, &y = 0x14  
    int *p = &x;            // p is a pointer to an integer  
    *p = y;  
    p = p + 4;  
    p = &y;  
    x = *p + 1;  
}
```

Line 1:	Line 2:	Line 3:
Line 4:	Line 5:	Line 6:

Draw out the memory diagram after sequential execution of each of the lines below:

```
int main(int argc, char **argv) {  
    int x = 410, y = 350;    // assume &x = 0x10, &y = 0x14  
    int *p = &x;            // p is a pointer to an integer  
    *p = y;  
    p = p + 4;  
    p = &y;  
    x = *p + 1;  
}
```



# Arrays (More examples on this...maybe next recitation)

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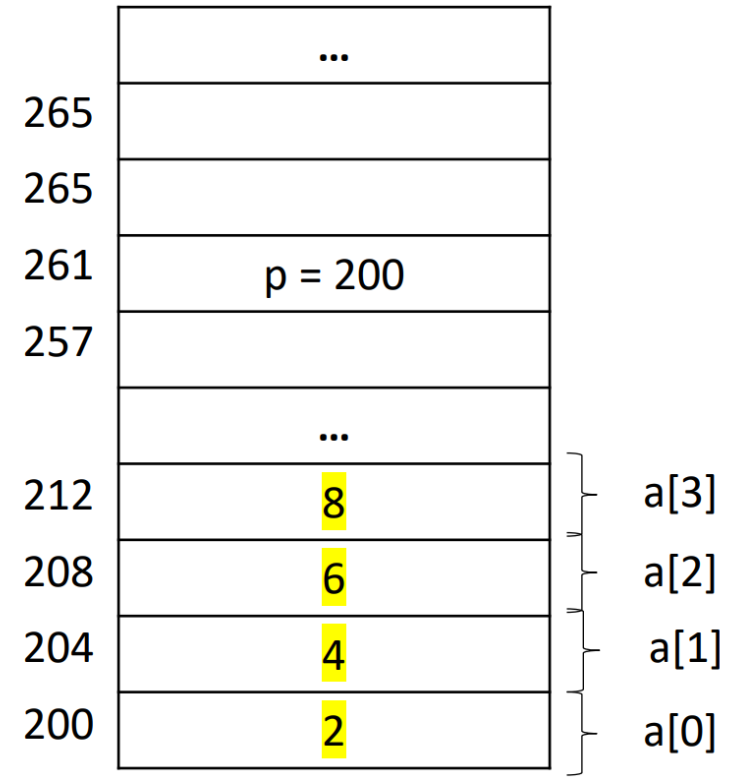
```
int a[] = {2, 4, 6, 8};
```

```
int *p = a; //Equivalent to *p = &a[0]
```

```
printf("%d , %d", (p+1), *(p+1)); //204, 4
```

```
printf("%d , %d", (a+1), *(a+1)); //204, 4
```

```
printf("%d , %d", &a[1], a[1]); //204, 4
```



# "String" as "Array"

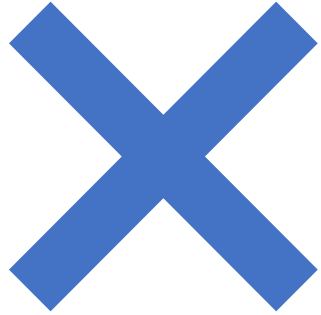
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- Represented by "array" of chars
- Each char encoded in ASCII format
  - Need to end with `\0`: C doesn't know when a string ends
    - If not...warning: you don't know what will happen!
- But don't get too comfortable saying "array" like you are still coding in Java...

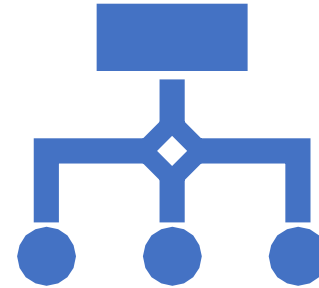
	0	1	2	3	4	5
str	G	e	e	k	s	\0
Address	0x23452	0x23453	0x23454	0x23455	0x23456	0x23457

# Arrays ain't real

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DO NOT PASS ARRAY AS  
PARAMETERS



Use POINTERS (in fact, array gets converted to  
pointer if you are determined to break the rules...)



*Classic examples*

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Next hw (I don't have  
it yet): Start early!  
Ask questions!