

MPI-4 Fault Tolerance

Chapter text available for a limited time:

http://www.mcs.anl.gov/~wbland/ft.pdf

Permanent Home:

https://svn.mpi-forum.org/trac/mpi-forum-web/wiki/

<u>User_Level_Failure_Mitigation</u>



Failure Notification

- Local failure notification only
 - Global notification can be built on top of these semantics
- Return error class to indicate process failure
 - MPI_ERR_PROC_FAILED
- Errors are only returned if the result of the operation would be impacted by the error
 - i.e. Point-to-point with non-failed processes should work unless routing is broken
- Some processes in an operation will receive MPI_SUCCESS while others will receive MPI_ERR_PROC_FAILED
 - i.e. Collective communication will sometimes work after a failure depending on the communication topology
 - Broadcast might succeed for the top of the tree, but fail for some children
 - MPI_ALLREDUCE would always fail if the error occurred before the start of the operation
- Wildcard operations must return an error because the failed process might have been sending the message that would have matched the MPI_ANY_SOURCE.
 - Return MPI ERR PENDING for MPI IRECV.
 - If application determines that it's ok, the request can be continued after re-enabling wildcards



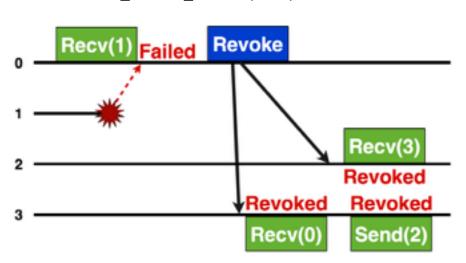
Failure Notification (cont.)

- To find out which processes have failed, use the two-phase functions:
 - MPI_COMM_FAILURE_ACK(MPI_Comm comm)
 - Internally "marks" the group of processes which are currently locally know to have failed
 - Useful for MPI_COMM_AGREE later
 - Re-enables wildcard operations on a communicator now that the user knows about the failures
 - Could be continuing old wildcard requests or new ones
 - MPI_COMM_FAILURE_GET_ACKED(MPI_Comm comm, MPI_Group *failed_grp)
 - Returns an MPI_GROUP with the processes which were marked by the previous call to MPI_COMM_FAILURE_ACK
 - Will always return the same set of processes until FAILURE ACK is called again
- Must be careful to check that wildcards should continue before starting/restarting a wildcard operation
 - Don't enter a deadlock because the failed process was supposed to send a message
- Future wildcard operations will not return errors unless a new failure occurs.



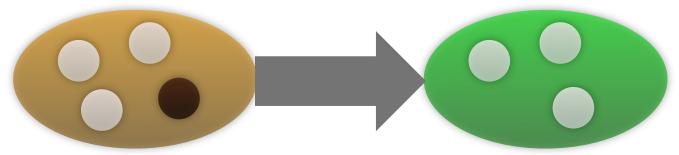
Failure Propagation

- Often unnecessary
 - Let the application discover the error as it impacts correct completion of an operation.
- When necessary, manual propagation is available.
 - MPI_COMM_REVOKE(MPI_Comm comm)
 - Interrupts all non-local MPI calls on all processes in comm.
 - Once revoked, all non-local MPI calls on all processes in comm will return MPI_ERR_REVOKED.
 - Exceptions are MPI_COMM_SHRINK and MPI_COMM_AGREE (later)
 - Necessary for deadlock prevention
 - Example on right



Failure Recovery

- Some applications will not need recovery.
 - Point-to-point applications can keep working and ignore the failed processes.
- If collective communications are required, a new communicator must be created.
 - MPI_Comm_shrink(MPI_Comm *comm, MPI_Comm *newcomm)
 - Creates a new communicator from the old communicator excluding failed processes
 - If a failure occurs during the shrink, it is also excluded.
 - No requirement that comm has a failure. In this case, it will act identically to MPI Comm dup.
- Can also be used to validate knowledge of all failures in a communicator.
 - Shrink the communicator, compare the new group to the old one, free the new communicator (if not needed).
 - Same cost as querying all processes to learn about all failures



Additional Field for Division / Organization/Sponsor/Meeting name

Fault Tolerant Consensus

- Sometimes it is necessary to decide if an algorithm is done.
 - MPI_COMM_AGREE(MPI_comm comm, int *flag);
 - Performs fault tolerant agreement over boolean flag
 - Non-acknowledged, failed processes cause MPI_ERR_PROC_FAILED.
 - Will work correctly over a revoked communicator.
 - Expensive operation. Should be used sparingly.
 - Can also pair with collectives to provide global return codes if necessary.
- Can also be used as a global failure detector
 - Very expensive way of doing this, but possible.
- Also includes a non-blocking version



MPI RMA

- MPI_WIN_REVOKE
 - Provides same functionality as MPI_COMM_REVOKE
- The state of memory targeted by any process in an epoch in which operations raised an error related to process failure is undefined.
 - Local memory targeted by remote read operations is still valid.
 - It's possible that an implementation can provide stronger semantics.
 - If so, it should do so and provide a description.
 - We may revisit this in the future if a portable solution emerges.
- MPI_WIN_FREE has the same semantics as MPI_COMM_FREE



Passive Target Locks

- Without restricting lock implementations, it's difficult to define the status of a lock after a failure
 - With some lock implementations, the library doesn't know who is holding the lock at any given time.
 - If the process holding the lock fails, the implementation might not be able to recover that lock portably.
 - Some other process should be notified of the failure and recovery can continue from there (probably with MPI_WIN_REVOKE).
 - If the implementation can get around the failure, it should try to do so and mask the failure.



MPI I/O

- When an error is returned, the file pointer associated with the call is undefined.
 - Local file pointers can be set manually
 - Application can use MPI_COMM_AGREE to determine the position of the pointer
 - Shared file pointers are broken
- MPI_FILE_REVOKE
 - Provides same functionality as MPI_COMM_REVOKE
- MPI FILE CLOSE has similar to semantics to MPI COMM FREE

